| | Page No. |
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| | Assignment No. 3 |
| au-1 Ans- | What is the Normalization? Explain its types |
| | NORMALIZATION |
| -10 h | Normalization is a process of organizing the data in database to avoid data redundancy, insertion anomaly, update anomaly and deletion anomaly. |
| | The purpose of normalization is to produce a stable set of relations that is a faithfull model of the operations of the enterprise. By following the principles of normalization, we can achieve a design of highly flexibility, allowing the model to be extended. |
| | TYPES OF NORMALIZATION |
| 1) | First Normal Form |
| ii) | Second Normal form (2NF) |
| 11) | Third Normal form (3NF) |
| (V) | Boyce - Codd Normal form (BCNF) |
| vi) | fourth Normal form (4NF) fifth Normal form (5NF) |
| in | first Normal form: |
| | if and only if all underlying domains contain atomic values only. |

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| Fire | st Normal | form - | | |
|------|-----------|----------|-----|--|
| | SID | ACTIVITY | PEE | |
| | 100 | SKING | 200 | |
| 1 | 150 | SWIMMING | 150 | |
| | 175 | SQUASH | 50 | |

ii) Second Normal form:

A relation is in second normal form

(2NF) and every non-key attribute is fully

dependent on the key. If the key is a

single attribute, then the relation is automatically in the 2NF.

SC(S#, CITY) and CS (CITY, STATUS)

| Į | S# | CITY | 8 23 3 3 | | |
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| 1 | \$1 | | CITY | STATUS | |
| | 34 | Delhi | Dehi | 1.30 | |
| - | 52 | Karnal | Karnal | 10 | |
| | 53 | Rohtak | | 10 | |
| | 1 34 | Delhi | Rohtak | 40 | |
| ı | | 1 2000 | | | |

(3NF) if and only if it is 2NF and every nonkey attribute is nontransitively dependent on the primary pey.

3NF was designed to improve the database processing while minimizing the storage costs.

9th 18 yeary beneficial for online transactions processing applications with the need of heavy order entry.

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Boyce - Codd Normal form (BCNF):

Unfortunately, even relations in third

normal form can have anomalies. Consider

the ABVISOR (SID, Major, Fname) relation.

Since students can have several majors, SID

does not determine major. Further, since

students can have several advisors, SID

does not determine fname. Thus SID

cannot be a key.

Key (primary): (SID, Major)

Key (candidate): (SID, Frame)

Functional dependencies: Frame -> Major

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| | DIS | Major | - Fname | |
|---|-----|---------|---------|--|
| - | | | | |
| - | 100 | MATH | CAUCHY | |
| | 150 | PHYSICS | JUNG | |
| | 200 | MATH | RIEMANN | |
| - | 250 | MATH | CAUCHY | |
| | 200 | PHYSICS | CHETAN | |
| | | | | |

V) fourth Normal form (4NF):

STUDENT

| The state of the s | | | |
|--|------|------------|----------|
| | SID | Major | Activity |
| | 200 | MUSIC | SWIMMING |
| tol. | 200 | ACCOUNTING | SWIMMING |
| | -200 | MUSIC | TENNIS |
| - | 200 | ACCOUNTING | TENNIS |
| 18 | 250 | MATH | JOGGING |
| | | | |

SID -> Activity

The relation is in BCNF (2NF becomes it is all key, 3 NF because it has no non-key determinants). However, as we have

seen, it has anomalies.
We can eliminate these anomalies by decompose
STUDENT relations into two relations.

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STU-MAJOR (SID, Major) and STU-ACT (SID, Activity)

Key: (SID, Major)

| == | - | | |
|-----|-----|------------|---|
| | SID | Fname | - |
| | 200 | MUSIC | |
| | 200 | ACCOUNTING | |
| | 250 | MATH | |
| 100 | | | * |

Key: (SID, Activity)

| _ | | |
|-------------------------|-----|----------|
| | SID | fname |
| | 200 | SKING |
| | 200 | SWIMMING |
| | 200 | TENNIS |
| | 250 | JOGGING |
| THE OWNER OF THE OWNER, | | |

This observation leads to the definition of YNF:
A relation is in YNF if it is in BCNF
and has no multivalued dependencies.

vi) fifth Normal Form (5NF):

Consider relation SPI shown below this relation is "all key" and involves no FDS on MUDS and so in YNF.

It shows

a) the three projections SP, PJ & JS of SPJ and b) the effect of joining SP & PJ over P# & then joining the result & JS over (J#,S#).

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| 17 | SPJ | | | | | | | | | | |
| | | | | 54 | | P# | J# | | | | |
| | | | | 31 | | PL | J2 | | | | |
| | | | | 31 | | P2 | JL | | | | |
| | | | 52 | | | PI | 71 | | | | |
| | | | | 31 | | PI | JL | | | | |
| | | | | | | | | | | | |
| SP | | | | P5 | | | | 35 | | | |
| | 5# | 149 | | DE | P# | J# | | J# | S7# | | |
| | 31 | P1 | | | PI | 52 | | 52 | S1 | | |
| | 51 | P2 | | | P1 | 151 | | 31 | 51 | | |
| | 52 | P1 | | | P2 | 71 | | 31 | 52 | | |
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| | 7 | S2 S2 | PI | | 2) | | 51 | 1 82 | 1 71 | - | |
| (640) | 16 | 74 | PL | 1 3 | 1 | 100 | 52 | PL | J1 | - | |
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| THAT WE SEE THE WAS TO SELECT THE | | | | | | | 6 1 | 100 | _ | | |
| | if the pain (S1, p1 > appears in SP | | | | | | | | | _ | |
| | and the pain (j1, j1) appears in PJ and the pain (j1, s1) appears in JS | | | | | | | | _ | | |
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| Que-2 Au- | What is meant by functional dependency? Give example of functional dependency. |
| April - | FUNCTIONAL DEPENDENCY |
| | A functional dependency is a constraint between the two sets of attributes in a relation from a database. It describes a relationship between attributes in a relation. |
| | A dependency (Example) FD: x -> y means that the value of y are determined by the values of x. Two rows |
| | of tuples having the same value of x will necessarily have the same value of y. Types of functional dependencies are: |
| | · Trivial dependency · Non trivial dependency · Transitive dependency |
| Que-3 Ans - | What is trigger? Explain. |
| | TRIGGER: A trigger is a statement that a system executes as a side effect of a modification to the database. A trigger can also be |
| 10 TO | |

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procedure that is executed whonever an artim, such as data modification, takes place.

A trigger is always defined on a table, and is said to have fired whonever the data in the underlying table is affected by any of the Data pranipulation language (DML) statements. - INSERT, UPDATE or BEIETE.

To design a trigger, we must:

is to be executed. This is broken up into an event that causes the trigger to be checked and a condition that must be satisfied for trigger execution to proceed.

ii) Specify the actions to be taken when the trigger executes.