

CS 202, Spring 2021

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SECTION: 2

Homework 1 – Algorithm Efficiency and Sorting

Question 1

(a) Showing $f(n) = 5n^3 + 4n^2 + 10 = O(n^4)$

We need to find two positive constants: c and n_0 such that:

$$0 \leq 5n^3 + 4n^2 + 10 \leq cn^4 \text{ for all } n \geq n_0$$

$$5/n + 4/n^2 + 10/n^4 \leq c \text{ for all } n \geq n_0$$

Choose $c = 3$ and $n_0 = 3$

- $5n^3 + 4n^2 + 10 \leq 3n^4$ for all $n \geq 3$

(b) Tracing sorting algorithms to sort integer array [24, 8, 51, 28, 20, 29, 21, 17, 38, 27] in ascending order.

– **Insertion sort ($O(n^2)$)**

```
void insertionSort( int arr[], int arrSize) {
    for ( int notSorted = 1; notSorted < arrSize; ++notSorted) {
        int    key = arr[notSorted]; // get current item to compare
        int    index = notSorted;    // get location
        for ( ; (index > 0) && (arr[index-1] > key); --index)
            arr[index] = arr[index-1]; // shift while item is bigger than key
        arr[index] = key; // insert current key
    }
}
```

[24, 8, 51, 28, 20, 29, 21, 17, 38, 27]

- Select first element as sorted ($\text{arr}[0]$) // **notSorted = 1**

[24, 8, 51, 28, 20, 29, 21, 17, 38, 27]

- Key to compare $\Rightarrow 8$ // **key = arr[1] & index = 1**
- Shift 24 to right, insert 8 to index 0 // **as arr[0] > 8 (1 comparison)**
- Increase sorted part by 1 // **++notSorted (2)**

[8, 24, 51, 28, 20, 29, 21, 17, 38, 27]

- Key to compare $\Rightarrow 51$ // **key = arr[2] & index = 2**
- Compare once, there exists no larger element in sorted part // **arr[1] is not smaller than 51 (1 comparison)**

- Increase sorted part by 1 (continue iteration) **++notSorted (3)**

[8, 24, 51, 28, 20, 29, 21, 17, 38, 27]

- Key to compare => 28 // **key = arr[3] & index = 3**
- Shift 51 to right, insert 28 to index 2 // **as arr[2] > 28, arr[1] not greater than 28 (2 comparisons)**
- Increase sorted part by 1 // **++notSorted (4)**

[8, 24, 28, 51, 20, 29, 21, 17, 38, 27]

- Key to compare => 20 // **key = arr[4] & index = 4**
- Shift 51, 28, 24 to right respectively, insert 20 to index 1 // **up to arr[0] not greater than 20 (4 comparisons)**
- Increase sorted part by 1 // **++notSorted (5)**

[8, 20, 24, 28, 51, 29, 21, 17, 38, 27]

- Key to compare => 29 // **key = arr[5] & index = 5**
- Shift 51 to right, insert 29 to index 4 // **up to arr[3] not greater than 29 (2 comparisons)**
- Increase sorted part by 1 // **++notSorted (6)**

[8, 20, 24, 28, 29, 51, 21, 17, 38, 27]

- Key to compare => 21 // **key = arr[6] & index = 6**
- Shift 51, 29, 28, 24 to right respectively, insert 21 to index 2 // **up to arr[3] not greater than 21 (5 comparisons)**
- Increase sorted part by 1 // **++notSorted (7)**

[8, 20, 21, 24, 28, 29, 51, 17, 38, 27]

- Key to compare => 17 // **key = arr[7] & index = 7**
- Shift 51, 29, 28, 24, 21, 20 to right respectively, insert 17 to index 1 // **up to arr[0] not greater than 17 (7 comparisons)**
- Increase sorted part by 1 // **++notSorted (8)**

[8, 17, 20, 21, 24, 28, 29, 51, 38, 27]

- Key to compare => 38 // **key = arr[8] & index = 8**
- Shift 51 to right, insert 38 to index 7 // **up to arr[6] not greater than 38 (2 comparisons)**
- Increase sorted part by 1 // **++notSorted (9)**

[8, 17, 20, 21, 24, 28, 29, 38, 51, 27]

- Key to compare => 27 // **key = arr[9] & index = 9**
- Shift 51, 38, 29, 28 to right, insert 27 to index 5 // **up to arr[4] not greater than 38 (5 comparisons)**
- Increase sorted part by 1 // **++notSorted (10)**

[8, 17, 20, 21, 24, 27, 28, 29, 38, 51] //THE ARRAY IS SORTED... (notSorted == arrSize)

– Bubble sort ($O(n^2)$)

// the characteristic of the algorithm is making comparisons with next element

// it is useful as it gives the opportunity to exit immediately

```
void bubbleSort( int arr[], int arrSize) {
    flagSorted = false; // helps for immediate exit

    for ( int cont = 1; (cont< arrSize) && ! flagSorted; ++cont) {
        flagSorted = true;
        for ( int index = 0; arrSize - cont; ++index) {
            int nextIndex = index + 1;
            if ( arr[index] > arr[nextIndex]) {
                swap( arr[index], arr[nextIndex]);
                flagSorted = false; // change signal
            }
        }
    }
}
```

[24, 8, 51, 28, 20, 29, 21, 17, 38, 27]

- flagSorted = false // initialize
- **cont = 1**, flagSorted = true, index = 0, inner loop: for (0->9)
- nextIndex = 1
- swap 24 & 8 (arr[0] > arr[1]) => [8, 24, 51, 28, 20, 29, 21, 17, 38, 27]
- **flagSorted = false**, index++ (1), nextIndex (2)
- no swap 24 & 51 !(arr[1] > arr[2]) => [8, 24, 51, 28, 20, 29, 21, 17, 38, 27]
- index++ (2), nextIndex (3)
- swap 51 & 28 (arr[2] > arr[3]) => [8, 24, 28, 51, 20, 29, 21, 17, 38, 27]
- flagSorted = false, index++ (3), nextIndex (4)
- swap 51 & 20 (arr[3] > arr[4]) => [8, 24, 28, 20, 51, 29, 21, 17, 38, 27]
- flagSorted = false, index++ (4), nextIndex (5)
- swap 51 & 29 (arr[4] > arr[5]) => [8, 24, 28, 20, 29, 51, 21, 17, 38, 27]
- flagSorted = false, index++ (5), nextIndex (6)
- swap 51 & 21 (arr[5] > arr[6]) => [8, 24, 28, 20, 29, 21, 51, 17, 38, 27]
- flagSorted = false, index++ (6), nextIndex (7)
- swap 51 & 17 (arr[6] > arr[7]) => [8, 24, 28, 20, 29, 21, 17, 51, 38, 27]
- flagSorted = false, index++ (7), nextIndex (8)
- swap 51 & 38 (arr[7] > arr[8]) => [8, 24, 28, 20, 29, 21, 17, 38, 51, 27]
- flagSorted = false, index++ (8), nextIndex (9)
- swap 51 & 27 (arr[8] > arr[9]) => [8, 24, 28, 20, 29, 21, 17, 38, 27, 51]
- flagSorted = false, index++ (9) (exit inner loop as !(index<9))
- **cont = 2**, flagSorted = true, index = 0, arrSize - cont = 8
- inner loop: for (0->8) [8, 24, 28, 20, 29, 21, 17, 38, 27, 51] (51 is fixed)
- nextIndex = 1
- no swap 8 & 24 !(arr[0] > arr[1]) => [8, 24, 28, 20, 29, 21, 17, 38, 27, 51]
- index++ (1), nextIndex (2)

- no swap 24 & 28 $!(arr[1] > arr[2]) \Rightarrow [8, 24, 28, 20, 29, 21, 17, 38, 27, 51]$
- index++ (2), nextIndex (3)
- swap 28 & 20 ($arr[2] > arr[3]$) $\Rightarrow [8, 24, 20, 28, 29, 21, 17, 38, 27, 51]$
- **flagSorted = false**, index++ (3), nextIndex (4)
- no swap 28 & 29 $!(arr[3] > arr[4]) \Rightarrow [8, 24, 20, 28, 29, 21, 17, 38, 27, 51]$
- index++ (4), nextIndex (5)
- swap 29 & 21 ($arr[4] > arr[5]$) $\Rightarrow [8, 24, 20, 28, 21, 29, 17, 38, 27, 51]$
- index++ (5), nextIndex (6)
- swap 29 & 17 ($arr[5] > arr[6]$) $\Rightarrow [8, 24, 20, 28, 21, 17, 29, 38, 27, 51]$
- index++ (6), nextIndex (7)
- no swap 29 & 38 $!(arr[6] > arr[7]) \Rightarrow [8, 24, 20, 28, 21, 17, 29, 38, 27, 51]$
- index++ (7), nextIndex (8)
- swap 38 & 27 ($arr[7] > arr[8]$) $\Rightarrow [8, 24, 20, 28, 21, 17, 29, 27, 38, 51]$
- index++ (8) (**exit inner loop as $!(index < 8)$**)

- **cont = 3, flagSorted = true**, index = 0, arrSize - cont = 7
- **inner loop: for (0->7)** **[8, 24, 20, 28, 21, 17, 29, 27, 38, 51]** (38 is fixed)
- nextIndex = 1
- no swap 8 & 24 $!(arr[0] > arr[1]) \Rightarrow [8, 24, 20, 28, 21, 17, 29, 27, 38, 51]$
- index++ (1), nextIndex (2)
- swap 24 & 20 ($arr[1] > arr[2]$) $\Rightarrow [8, 20, 24, 28, 21, 17, 29, 27, 38, 51]$
- **flagSorted = false**, index++ (2), nextIndex (3)
- no swap 24 & 28 $!(arr[2] > arr[3]) \Rightarrow [8, 20, 24, 28, 21, 17, 29, 27, 38, 51]$
- index++ (3), nextIndex (4)
- swap 28 & 21 ($arr[3] > arr[4]$) $\Rightarrow [8, 20, 24, 21, 28, 17, 29, 27, 38, 51]$
- index++ (4), nextIndex (5)
- swap 28 & 17 ($arr[4] > arr[5]$) $\Rightarrow [8, 20, 24, 21, 17, 28, 29, 27, 38, 51]$
- index++ (5), nextIndex (6)
- no swap 28 & 29 $!(arr[5] > arr[6]) \Rightarrow [8, 20, 24, 21, 17, 28, 29, 27, 38, 51]$
- index++ (6), nextIndex (7)
- swap 29 & 27 ($arr[6] > arr[7]$) $\Rightarrow [8, 20, 24, 21, 17, 28, 27, 29, 38, 51]$
- index++ (7) (**exit inner loop as $!(index < 7)$**)

- **cont = 4, flagSorted = true**, index = 0, arrSize - cont = 6
- **inner loop: for (0->6)** **[8, 20, 24, 21, 17, 28, 27, 29, 38, 51]** (29 is fixed)
- nextIndex = 1
- no swap 8 & 20 $!(arr[0] > arr[1]) \Rightarrow [8, 20, 24, 21, 17, 28, 27, 29, 38, 51]$
- index++ (1), nextIndex (2)
- no swap 20 & 24 $!(arr[0] > arr[1]) \Rightarrow [8, 20, 24, 21, 17, 28, 27, 29, 38, 51]$
- index++ (2), nextIndex (3)
- swap 24 & 21 ($arr[2] > arr[3]$) $\Rightarrow [8, 20, 21, 24, 17, 28, 27, 29, 38, 51]$
- **flagSorted = false**, index++ (3), nextIndex (4)
- swap 24 & 17 ($arr[3] > arr[4]$) $\Rightarrow [8, 20, 21, 17, 24, 28, 27, 29, 38, 51]$
- index++ (4), nextIndex (5)
- no swap 24 & 28 $!(arr[4] > arr[5]) \Rightarrow [8, 20, 21, 17, 24, 28, 27, 29, 38, 51]$
- index++ (5), nextIndex (6)
- swap 28 & 27 ($arr[5] > arr[6]$) $\Rightarrow [8, 20, 21, 17, 24, 27, 28, 29, 38, 51]$
- index++ (6) (**exit inner loop as $!(index < 6)$**)

- **cont = 5, flagSorted = true**, index = 0, arrSize - cont = 5
- **inner loop: for (0->5)** [8, 20, 21, 17, 24, 27, **28, 29, 38, 51**] (28 is fixed)
- nextIndex = 1
- no swap 8 & 20 !(arr[0] > arr[1]) => [8, 20, 21, 17, 24, 27, **28, 29, 38, 51**]
- index++ (1), nextIndex (2)
- no swap 20 & 21 !(arr[1] > arr[2]) => [8, 20, 21, 17, 24, 27, **28, 29, 38, 51**]
- index++ (2), nextIndex (3)
- **swap** 21 & 17 (arr[2] > arr[3]) => [8, 20, **17, 21**, 24, 27, **28, 29, 38, 51**]
- **flagSorted = false**, index++ (3), nextIndex (4)
- no swap 21 & 24 !(arr[3] > arr[4]) => [8, 20, 17, 21, 24, 27, **28, 29, 38, 51**]
- index++ (4), nextIndex (5)
- no swap 24 & 27 !(arr[4] > arr[5]) => [8, 20, 17, 21, 24, 27, **28, 29, 38, 51**]
- index++ (5) (**exit inner loop as !(index<5)**)

- **cont = 6, flagSorted = true**, index = 0, arrSize - cont = 4
- **inner loop: for (0->4)** [8, 20, 17, 21, 24, **27, 28, 29, 38, 51**] (27 is fixed)
- nextIndex = 1
- no swap 8 & 20 !(arr[0] > arr[1]) => [8, 20, 17, 21, 24, **27, 28, 29, 38, 51**]
- index++ (1), nextIndex (2)
- **swap** 20 & 17 (arr[1] > arr[2]) => [8, **17, 20**, 21, 24, **27, 28, 29, 38, 51**]
- **flagSorted = false**, index++ (2), nextIndex (3)
- no swap 20 & 21 !(arr[2] > arr[3]) => [8, 17, 20, 21, 24, **27, 28, 29, 38, 51**]
- index++ (3), nextIndex (4)
- no swap 21 & 24 !(arr[3] > arr[4]) => [8, 17, 20, 21, 24, **27, 28, 29, 38, 51**]
- index++ (4) (**exit inner loop as !(index<4)**)

- **cont = 7, flagSorted = true**, index = 0, arrSize - cont = 3
- **inner loop: for (0->3)** [8, 17, 20, 21, **24, 27, 28, 29, 38, 51**] (24 is fixed)
- nextIndex = 1
- no swap 8 & 17 !(arr[0] > arr[1]) => [8, 17, 20, 21, **24, 27, 28, 29, 38, 51**]
- index++ (1), nextIndex (2)
- no swap 17 & 20 !(arr[1] > arr[2]) => [8, 17, 20, 21, **24, 27, 28, 29, 38, 51**]
- index++ (2), nextIndex (3)
- no swap 20 & 21 !(arr[2] > arr[3]) => [8, 17, 20, 21, **24, 27, 28, 29, 38, 51**]
- index++ (3) (**exit inner loop as !(index<3)**)

- ! flagSorted condition is no longer functioning as flagSorted = true after inner loop thus the outer loop terminates

[8, 17, 20, 21, 24, 27, 28, 29, 38, 51] **//THE ARRAY IS SORTED**

Question 2

-Run your executable and add the screenshot of the console to the solution of Question 2 in the pdf submission.

```
[onur.vural@dijkstra SortingAlgorithms]$ ./SortMake
The array before selection sort:
12, 7, 11, 18, 19, 9, 6, 14, 21, 3, 17, 20, 5, 12, 14, 8,
comp num is: 120
move num is: 45
The array after selection sort:
3, 5, 6, 7, 8, 9, 11, 12, 12, 14, 14, 17, 18, 19, 20, 21,

The array before merge sort:
12, 7, 11, 18, 19, 9, 6, 14, 21, 3, 17, 20, 5, 12, 14, 8,
comp num is: 46
move num is: 128
The array after merge sort:
3, 5, 6, 7, 8, 9, 11, 12, 12, 14, 14, 17, 18, 19, 20, 21,

The array before quick sort:
12, 7, 11, 18, 19, 9, 6, 14, 21, 3, 17, 20, 5, 12, 14, 8,
comp num is: 45
move num is: 102
The array after quick sort:
3, 5, 6, 7, 8, 9, 11, 12, 12, 14, 14, 17, 18, 19, 20, 21,

The array before radix sort:
12, 7, 11, 18, 19, 9, 6, 14, 21, 3, 17, 20, 5, 12, 14, 8,
The array after radix sort:
3, 5, 6, 7, 8, 9, 11, 12, 12, 14, 14, 17, 18, 19, 20, 21,

THE ANALYSIS IS BEING DONE. PLEASE WAIT...
Analysis of Selection Sort
-Array Size-      -Elapsed Time-      -compCount-      -moveCount-
```

-The performanceAnalysis function needs to produce an output similar to the one given on the next page. Include this output to the answer of Question 2 in the pdf submission.

```
Analysis of Selection Sort
-Array Size-      -Elapsed Time-      -compCount-      -moveCount-
Randomized Inputs:
6000              80ms              17997000          17997
10000             240ms              49995000          29997
14000             470ms              97993000          41997
18000             780ms              161991000         53997
22000             1160ms             241989000         65997
26000             1610ms             337987000         77997
30000             2140ms             449985000         89997
Ascending Inputs:
6000              90ms              17997000          17997
10000             250ms              49995000          29997
14000             490ms              97993000          41997
18000             820ms              161991000         53997
22000             1230ms             241989000         65997
26000             1710ms             337987000         77997
30000             2290ms             449985000         89997
Descending Inputs:
6000              90ms              17997000          17997
10000             250ms              49995000          29997
14000             490ms              97993000          41997
18000             790ms              161991000         53997
22000             1190ms             241989000         65997
26000             1670ms             337987000         77997
30000             2210ms             449985000         89997
```

```
-----
Analysis of Merge Sort
-Array Size-      -Elapsed Time-      -compCount-      -moveCount-
Randomized Inputs:
6000              0ms              67827             151616
10000             10ms             120389            267232
14000             10ms             175419            387232
18000             10ms             231986            510464
22000             10ms             290108            638464
26000             10ms             348868            766464
30000             10ms             408597            894464
Ascending Inputs:
6000              0ms              39152             151616
10000             0ms              69008             267232
14000             0ms              99360             387232
18000             0ms              130592            510464
22000             0ms              165024            638464
26000             10ms             197072            766464
30000             10ms             227728            894464
Descending Inputs:
6000              0ms              36656             151616
10000             0ms              64608             267232
14000             0ms              94256             387232
18000             10ms             124640            510464
22000             10ms             154208            638464
26000             10ms             186160            766464
30000             10ms             219504            894464
-----
```

```

-----
Analysis of Quick Sort
-Array Size-      -Elapsed Time-      -compCount-      -moveCount-
Randomized Inputs:
6000               0ms               87053             151863
10000              0ms               152142            262255
14000              0ms               220119            382840
18000              10ms              311268            493286
22000              10ms              370141            615472
26000              0ms               436221            743165
30000              10ms              525538            893423
Ascending Inputs:
6000               80ms              17997000           23996
10000              210ms              49995000           39996
14000              410ms              97993000           55996
18000              680ms              161991000          71996
22000              1020ms             241989000          87996
26000              1420ms             337987000          103996
30000              1890ms             449985000          119996
Descending Inputs:
6000              160ms              17997000           27023996
10000              440ms              49995000           75039996
14000              860ms              97993000           147055996
18000             1430ms             161991000          243071996
22000             2140ms             241989000          363087996
26000             2990ms             337987000          507103996
30000             3980ms             449985000          675119996
-----

```

```

-----
Analysis of Radix Sort
-Array Size-      -Elapsed Time-
Randomized Inputs:
6000               10ms
10000              10ms
14000              10ms
18000              20ms
22000              20ms
26000              30ms
30000              30ms
Ascending Inputs:
6000               0ms
10000              10ms
14000              10ms
18000              20ms
22000              20ms
26000              20ms
30000              30ms
Descending Inputs:
6000               10ms
10000              10ms
14000              10ms
18000              10ms
22000              20ms
26000              20ms
30000              30ms
-----

```

Question 3

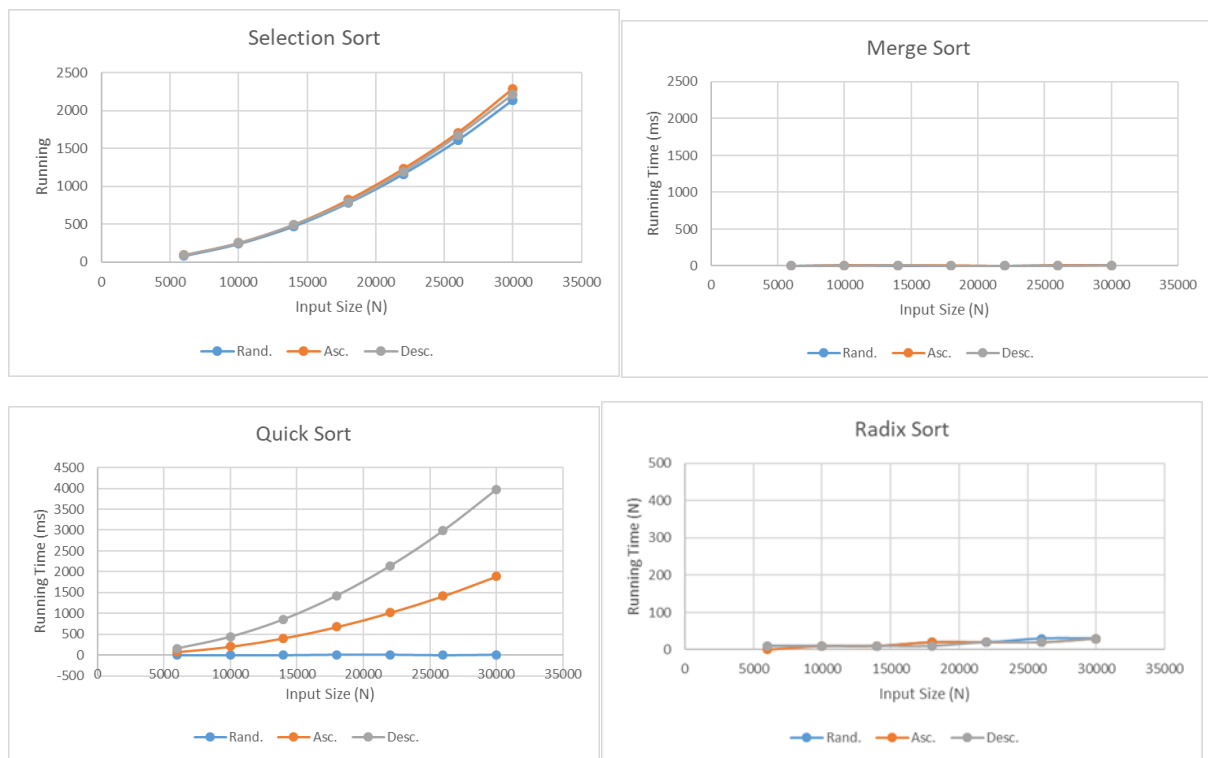


Figure 1: The relation between input size and running time for each sorting algorithm

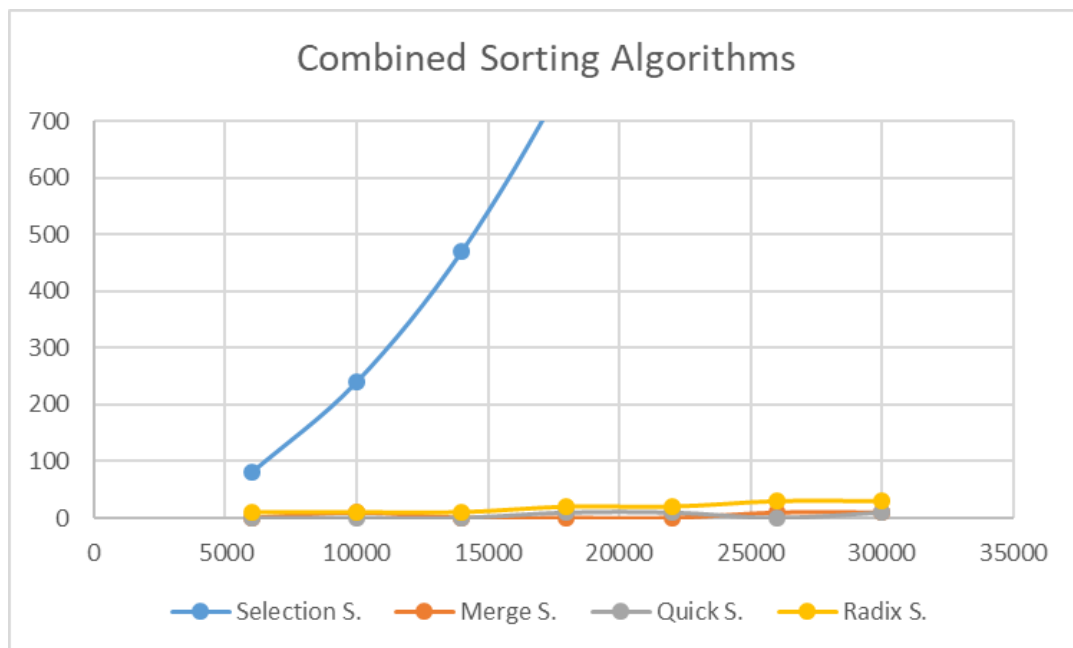


Figure 2: The sorting algorithms combined (for random inputs)

ANALYSIS REPORT:

For Selection Sort, it can be observed from the experimental results that the algorithm takes the highest running time among the four and therefore becomes very slow for very large inputs. This is in confirmatory with theoretical expectations as the algorithm proceeds by comparing all the items with each other while iterating inside two loop structure and thus becoming $O(n^2)$. This exponential relation is clearly noticable in the corresponding graph. Besides, it must be stated that the algorithm takes $O(n^2)$ for all three cases; random data, ascending data, descending data respectively. This is mostly because the algorithm does not depend on the initial organization of the data.

When it comes to Merge Sort, it turns out to be a highly efficient algorithm just as expected showing remarkably low running time overall. At this point it must be noted that although it almost looks like a constant relation $O(1)$ (the reason is that it proceeds highly efficiently under large inputs when compared to Selection Sort and to understand it's full behaviour much larger inputs than 30000 must be supplied) in the graph, it is actually $O(n \log n)$ for all three cases as the algorithm proceeds under binary recursion, splitting the array into two halves continuously. Although all cases have same time complexity, ascending and descending arrays result in a decrease on the number of key comparisons in a considerable scale as seen on the running time data on Q2.

Quick Sort Algorithm shows parallelity with theoretical expectations as it is expected to be $O(n^2)$ for worst case and $O(n \log n)$ else. The key concept in Quick Sort is not the initial order of the data but the value of the pivot. To provide more efficiency the pivot, which sepates the smaller values to one side and larger to another, must split the array in a more balanced fashion (50-50 split is idealized) while entering into recursion. This explains why ascending and descending values show $O(n^2)$ and randomized $O(n \log n)$ in the experimental findings. In a list which is already sorted, pivot value taken as the first entry means that (it is the largest or smallest value in the entire list) the split will completely be one parted (n-1 sub-list size) whereas random values show the characteristics of average case $O(n \log n)$ with the randomized pivot value.

The experimental results for Radix Sort demonstrate that it is also a highly efficient algorithm as the graph shows a linear relation $O(cn)$ (where c is digit number in this case) meaning that it grows relatively slower for high inputs when compared to algorithms with time complexity of $O(n^2)$ such as Bubble Sort or Selection Sort. But at this point it has to mentioned that although it provides remarkable efficiency, it requires to form groups that in order to hold the original data for each, resulting large memory usage (size of data * digit number). The algorithm is not concerned with the former organization of data thus for random, ascending and descending data result in same time complexity of $O(cn)$. Different from all the others, Radix Sort is not only dependant on total size but also a constant.

As a final comparison it can be observed from the graphs that the Selection Sort is the least efficient algorithm overall among others, growing exponentially with respect to input size. When it comes to random inputs, Merge Sort and Quick Sort act very similarly, and behave as $O(n \log n)$ which results in little execution time (QS is a bit faster). But for ascending and descending data, in other words sorted data, the experimental findings clearly demonstrate that Merge Sort outperforms Quick Sort as Quick Sort grows exponentially just like Selection Sort due to the selection of pivot value. In our data range (8000-30000), Radix Sort produces similar runtime results to Merge Sort and average case of Quick Sort but as it is linear ($O(cn)$), the theoretical information suggests that for more and more larger values it will grow much slower than the two and take less time if c, constant value, is not high (apart from the fact that it will consume much more space).