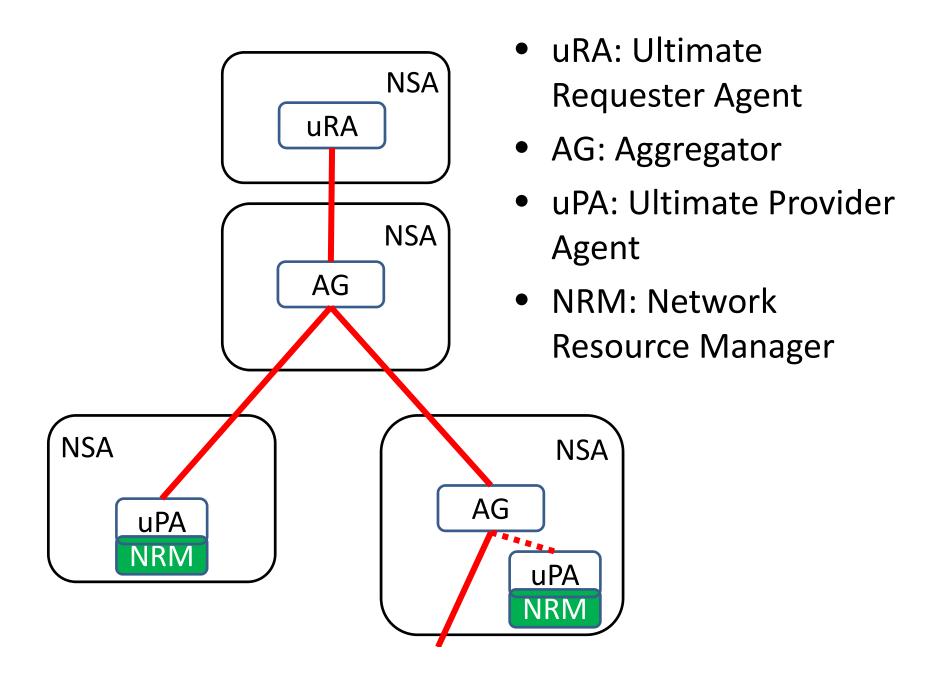


NSI CS Protocol State Machines and Message Handler

NSA: uRA, Aggregator and uPA

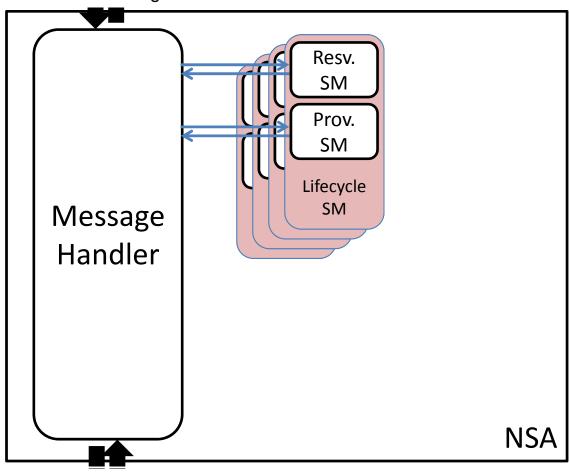


State Machines and Message Handler

- Behavior of NSI CS protocol is modeled as state machines and message handler
- State Machines:
 - RSM: Reservation State Machine
 - PSM: Provision State Machine
 - ASM: Activation State Machine
 - LSM: Lifecycle State Machine
- Aggregator:
 - can talk to upstream and downstream NSAs
 - Has RSM, PSM and LSM
- uPA
 - Can talk to upstream NSAs only
 - Has RSM, PSM, ASM and LSM

Aggregator

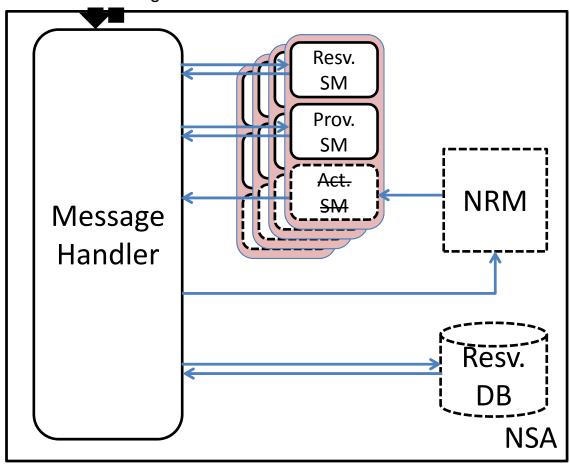




NSI message

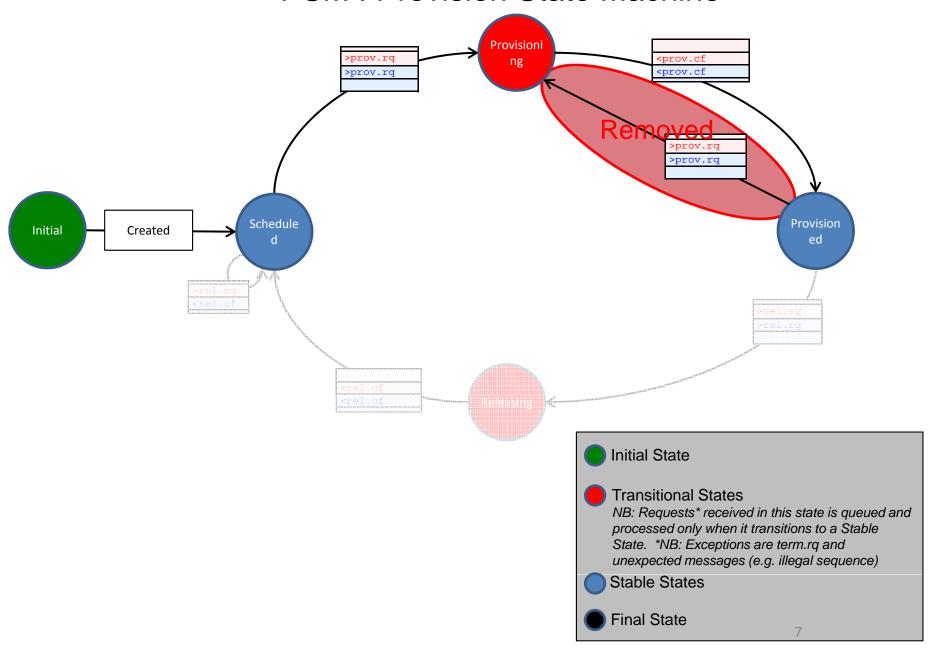
uPA





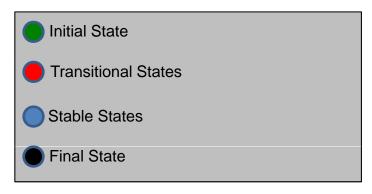
RSM: Reservation State Machine Initial rsv.rq >rsv.rq Modifyin <modify.cf Reserved <modify.cf</pre> g *1:Modify Checking, >mdfychk.rq modify.rq Modify Checked, <rsv.cf Modify Failed, Reserving **Modify Timeout** Modify Modify <mdfychk.cf Checking Checked and Reserved <mdfychk.cf modify_timeout <rsv.fl <mdfychk.fl <rsv.fl <modifyTimeout.nt</pre> <mdfychk.fl Modify timeout transitions back to Reserved. (This Modify Modify Reserve transition will only happen **Timeout** Failed Failed in uPA) **Initial State** mdfvcncl.ra Modify **Transitional States** mdfycncl.cf Reserved >mdfycncl.rq Canceling mdfycncl.cf NB: Requests* received in this state is queued and processed only when it transitions to a Stable State. *NB: Exceptions are term.rq and unexpected messages (e.g. illegal sequence) Stable States Final State

PSM: Provision State Machine



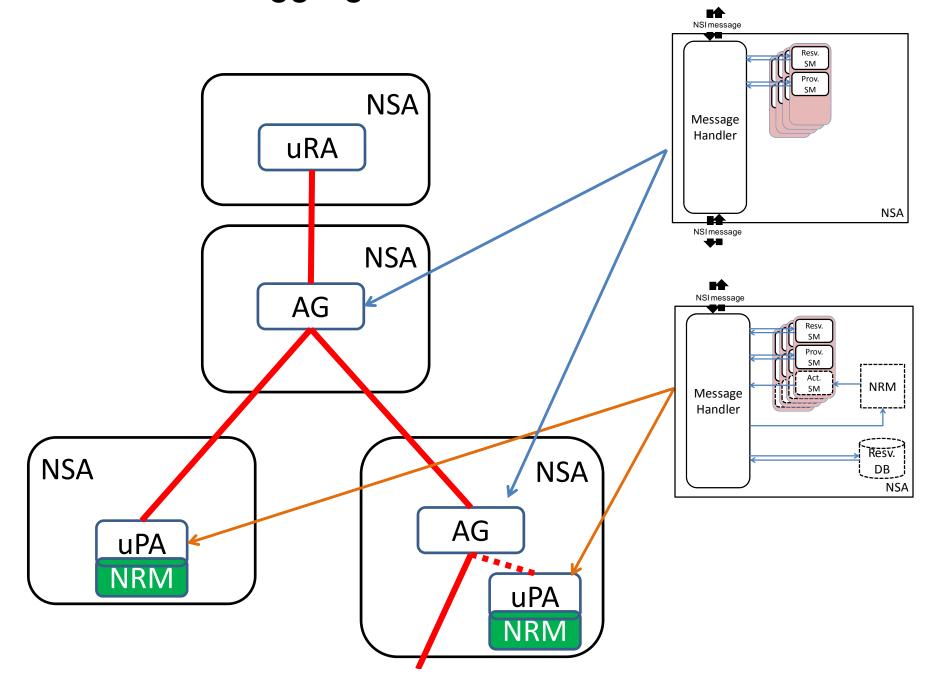
LSM : Lifecycle State Machine





07/03/2013

NSA: uRA, Aggregator and uPA



Modify

- Modify operations modify reservation
- Currently, following two changes are supported
 - Change end time of a reservation
 - Change bandwidth
- Modify is a 2-phase operation
 - 1: check availability (ModifyCheck.rq)
 - Note: resources are held
 - 2: Commit (Modify.rq) or Abort (ModifyCancel.rq)
- When committed, the reservation is updated
 - Reservation has a version number assigned by uRA, and the version number is updated when committed (uniformly increasing)

Provision and Release

- Provision state machine is independent from the reservation state machine
- Provision state:
 - Data plane should be activated if the PSM is in "Provisioned" state AND start_time < current_time < end_time
 - Data plane shall not be activated before the start_time
- A connection can be repeatedly provisioned and released

Data plane activation

- Data plane should be activated if the PSM is in "Provisioned" state AND start_time < current_time < end_time
- Activation is done at the timing of following events (if the above condition is met), using the latest reservation information
 - PSM transits to "Provisioned"
 - At the start_time
 - Reservation is updated (by commit of modify)
 - Data plane is recovered from an error
- Data plane activation/deactivation are notified by DataPlaneStateChange.nt notification messages.
- Errors are notified by a generic error message

DataPlaneStateChage.nt (1)

- PA and aggregator has DataPlaneStatus information
 - (Boolean) Active: True if data plane is active. For an aggregator, this flag is true when data plane is activated in all participating children
 - (Int) Version: For a uPA, current (latest) reservation version number.
 For an aggregator, the largest version number of the participating children. This field is valid when Active is true.
 - (Boolean) VersionConsistent: Always true for uPA. For an aggregator, If version numbers of all children are the same, This flag is true. This field is valid when Active is true.
- When a valid filed of DataPlaneStatus is changed, DataPlaneStatusChange.nt is sent up.

DataPlaneStatus

Active
Version
VersionConsistent

DataPlaneStateChage.nt(2)

- An aggregator keeps an array of statuses of its children, ChildrenDataPlaneStatus[1..n]
- Aggregator's DataPlaneStatus is determined by the following rule

ChildrenDataPlaneStatus

Active	Active		Active
Ver.	Ver.	•••	Ver.
VC	VC		VC
Child 1	Child 2	•	Child n

DataPlaneStatus

Active	
Ver.	
VC	

Notifications: Activation related

- There are no activateComplete.nt nor deactivateComplete.nt
- A general error message is used to notify following events. Those error are sent up the tree to uRA immediately
 - activateFailed: Activation failed at the time when uPA should activate its data plane
 - deactivateFailed: Deactivation failed at the time when uPA should deactivate its data plane
 - dataplaneError: Data plane is deactivate when deactivation is not expected. The error is recoverable.
 - forcedEnd: Something unrecoverable is happened in uPA/NRM

Requests which can fail

- Operations which can functionally fail are:
 - Reserve.rq
 - ModifyCheck.rq
 - Those requests fail when requested resources are not available.
- Other operation cannot fail. However, they can timeout in MTL/MH, or can be denied because they are invalid requests.
 - If a SM is at a state in which the request cannot be received, the request is denied.
 - *.na (not applicable) message is returned.

Message Handler (MH) and Message Transport Layer (MTL)

NSI stack

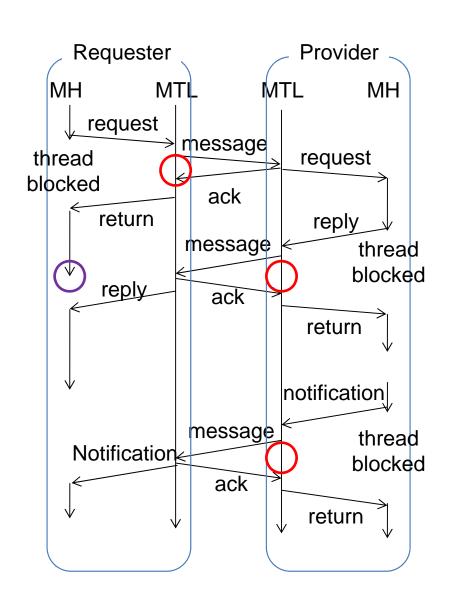
Message Handler



Message Transport Layer

- MH is a part of NSI stack, and uses MTL to send/receive messages
- MH is primarily responsible for keeping track of messaging state, e.g.
 - Who was the message sent to
 - Was the message received (i.e. ack'ed or MTL timeout)
 - Who has not replied to the message (e.g. *.cf,
 *.fl, etc)
- MTL is primarily responsible for sending and receiving messages, and notifying MH if the message was received, or if a (MTL) timeout occurs
- MTL interface (to MH) has 2 simple operations:
 - Send: blocks until ack is returned by destination MTL, or timeout happens.
 Timeout value is implementation dependent. NB: The MTL may be implemented to retry sending messages, but this is opaque to the MH
 - Receive: a thread in MH is invoked when a message is received

Message ack, reply and timeouts

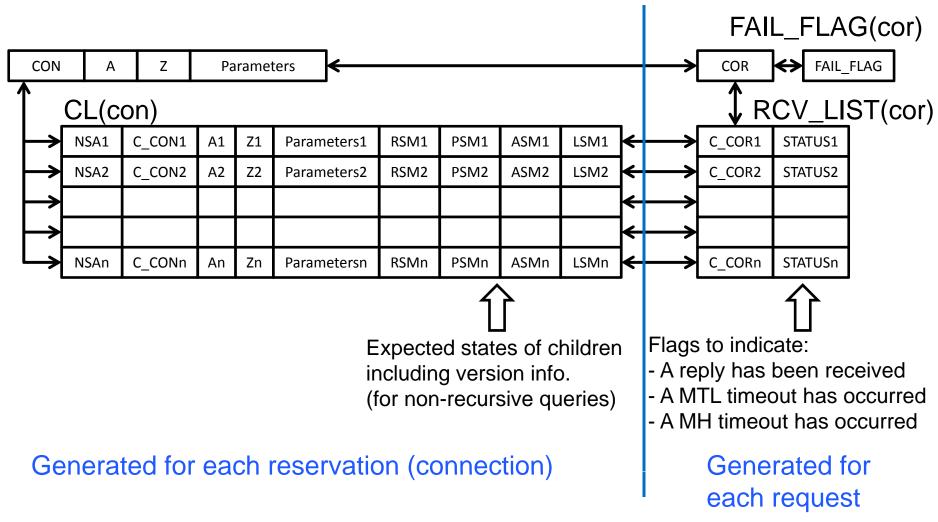


- : MTL timeout may happen
- : MH timeout may happen
- Ack is sent by MTL for each message
 - If ack is not returned in a certain period of time, MTL timeout occurs
- Reply is sent by MH (via MTL) and is either confirm, fail or not_applicable
 - MH can timeout if expected reply is not received from a child

Timeouts

- Message transport layer (MTL) timeout
 - Underlying MTL (http/tcp) initiates a MTL timeout
 - Happens when an ack is not returned for a message.
- Message Handler (MH) timeout
 - MH can timeout if a reply message is not returned in a certain period of time
- MH notifies both MTL and MH timeouts to the parent RA
- When a MTL/MH timeout is notified, uRA can either retry or terminate the connection.
 - Retry is requested by NSI_messageRetry.rq, which has the original request message's id (correlation id) as a parameter
 - MH keeps not-yet-replied requests in a table, so that it can re-send the request.

Tables an aggregator MH maintains for each reservation (connection)



modifyCancel after timeout

 After modifying operations, if a NSA is already in RESERVED state, it can receive NSI_modifyCancel.rq and reply NSI_modifyCancel.cf, but the modification is not rolled back. The system may be in an inconsistent state (different versions across the system) after those operations.

Notifications: modify timeout and MTL failure

- NSI_modifyTimeout.nt
- NSI_genericEvent.nt
 - Message delivery failure will be notified by this message (to be defined)
- When a MTL/MH timeout is notified, uRA can either retry or terminate the connection.
 - Retry is requested by NSI_messageRetry.rq, which has the original request message's id (correlation id) as a parameter

Termination and LSM

- A connection lifecycle is terminated when NSI_terminate.rq is received.
- LSM (Lifecycle State Machine) handles the terminate request.
 - Terminate request will delete the RSM, PSM and ASM, but the LSM should be there to send/receive terminate request and confirm messages
 - uPA may delete RSM, PSM and ASM when it isseues fcd_end, but LSM cannot be deleted.