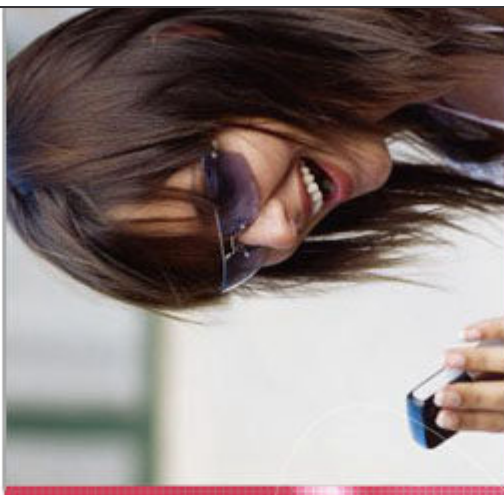


Performance comparison between a 2 phases and a 3 phases Negotiation protocol

Antoine Pichot, Alejandro Gaspar
April, 2007



Agenda

1. Problem

2. Model

3. Results

1

The Problem

The context : Co-allocation

Multiple Computing farms

A (G)MPLS network

How to reserve two (or more) resources at the same time ?

I.e. : A network connexion & a CPU



Solutions

- VIOLA like (2Phases commit protocol)
- WS-Agreement based Negotiation (3Phases Commit Protocol)
 - Cf O. Waeldrich & W. Ziegler draft @OGF18

3 Phases Commit Protocol vs 2 Phases

- Capability and availability check

Phase 1

- Pre-reservation

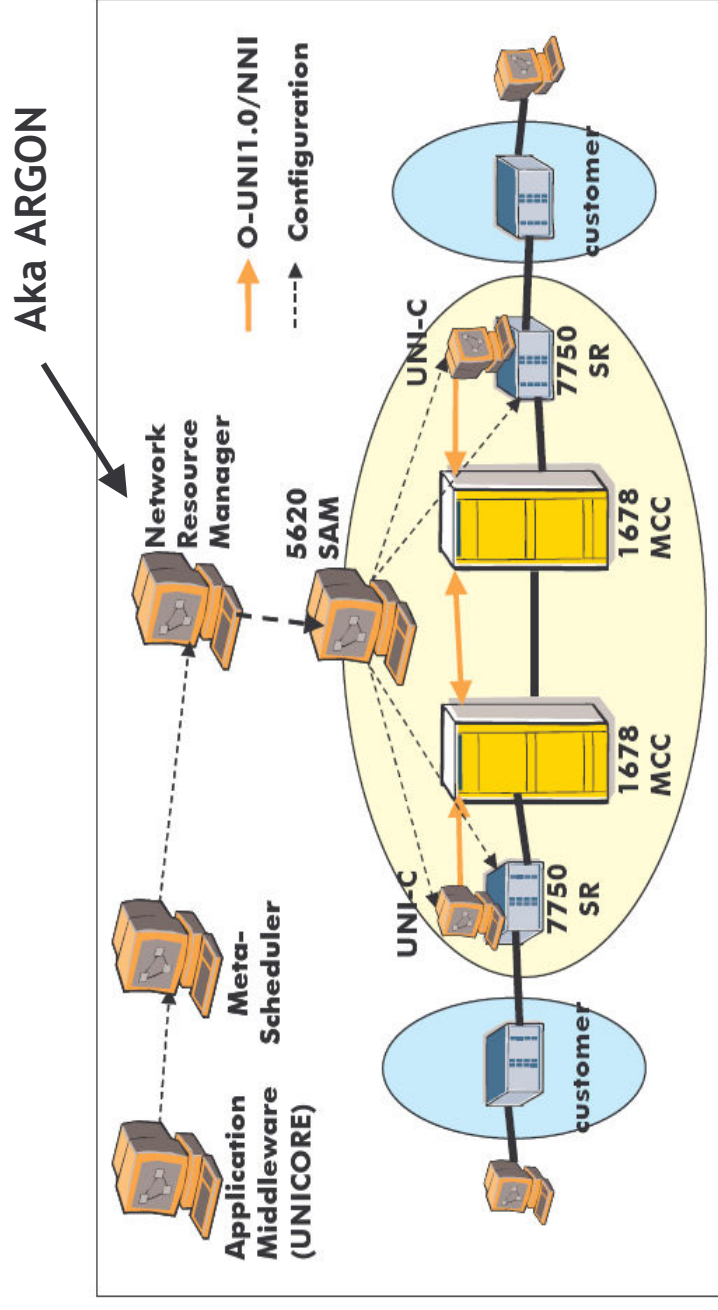
- Resource are reserved with short reservation lifetime
 - (No penalties if reservation is cancelled at this stage)

Phase 2

- Commitment

- Resource are reserved whenever needed
 - (penalties if cancelled)

VIOLA architecture



VIOLA reservation process

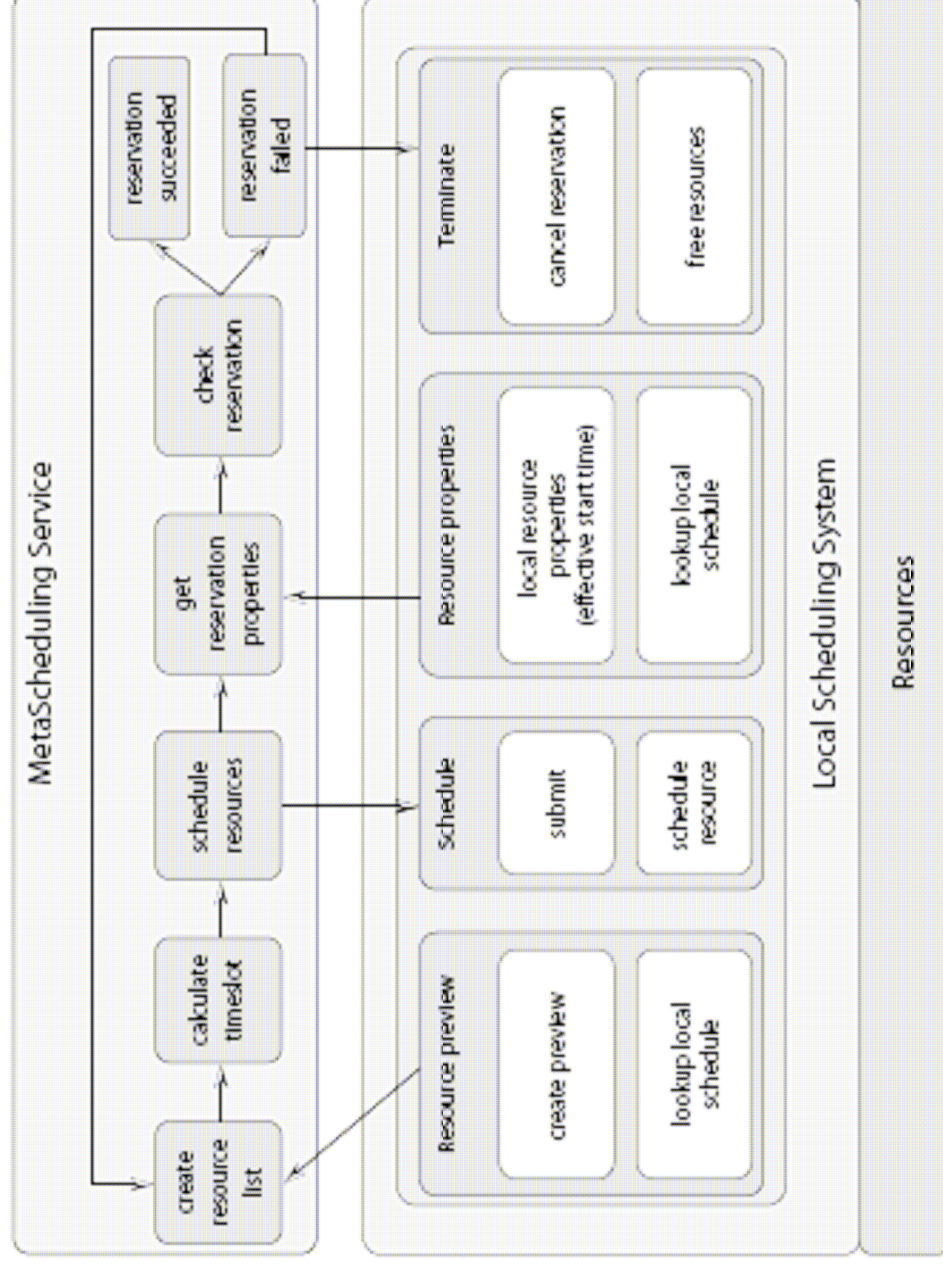
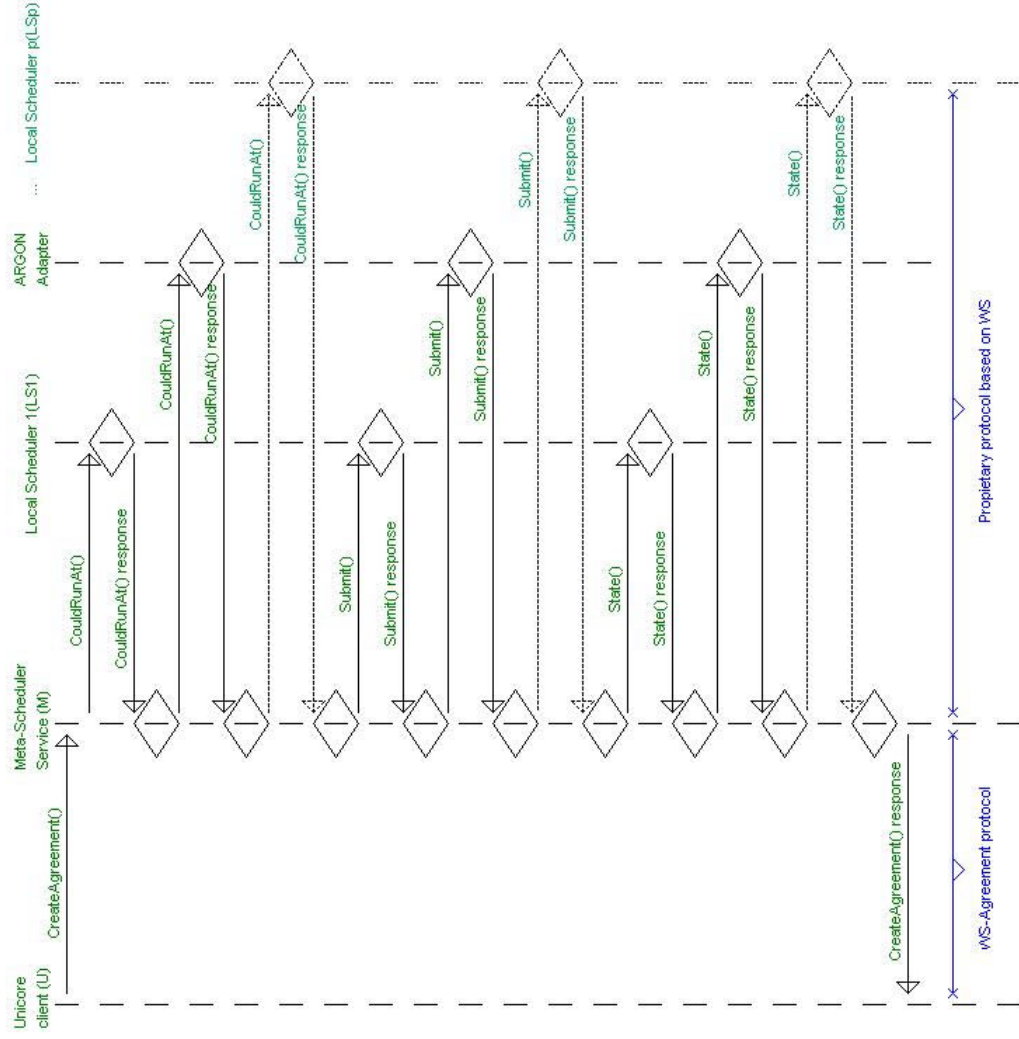


Figure 1: The negotiation process

VIOLA's Meta-Scheduler reservation



2

The Model Explained

Model origin

Gurbani V.K., Jagadeesan L., Mendiratta V.B.,
“Characterizing session initiation protocol (SIP) network performance and
reliability”,
International service availability symposium, April 2005

Model explained (1)

Job requests ⇔ Clients of the queuing network

Time spent in a waiting queue ⇔ Time spent in a state inside the MS

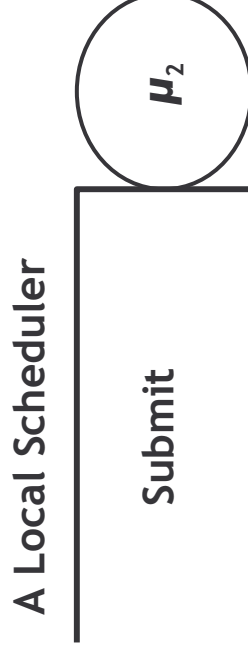
A queue ⇔ A state inside the MS

Model Explained (2), Notation convention



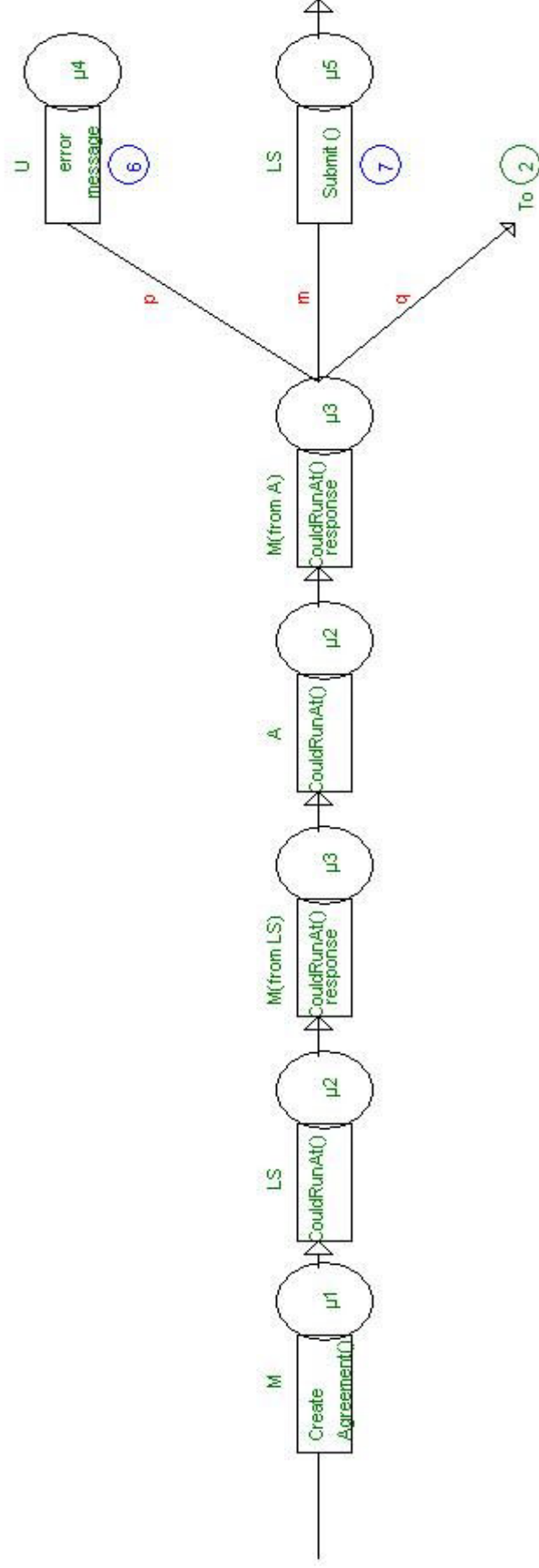
$1/\mu_1$ is the average time taken by the Meta-Scheduler to

- receive the CreateAgreement message,
- process it, and
- take action



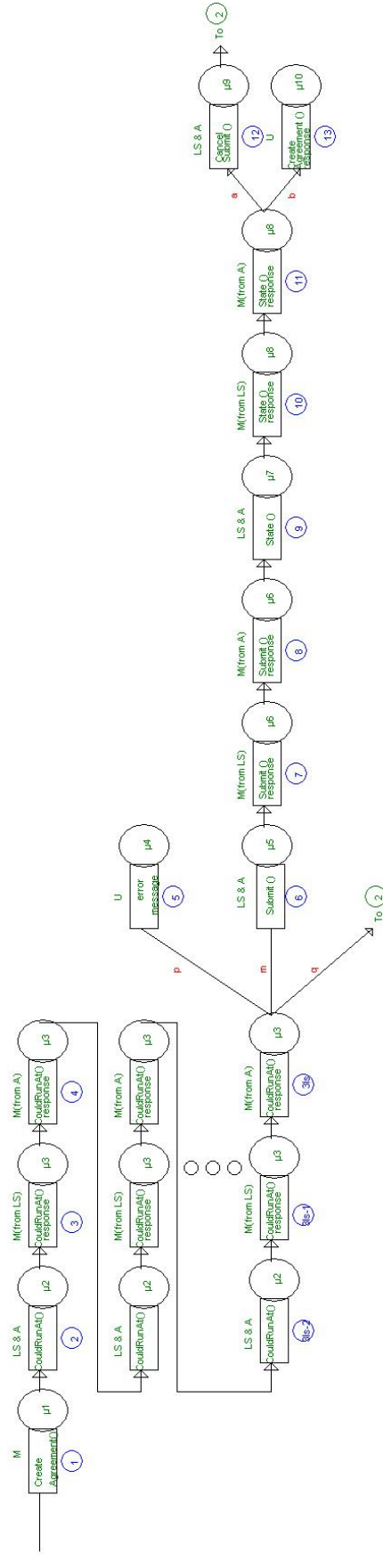
Model Explained (3), Example

Example synchronous

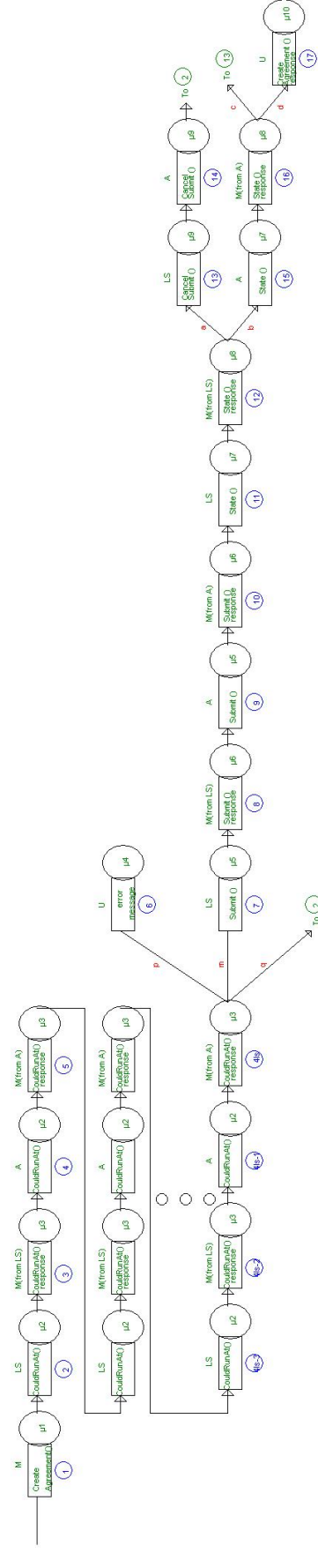


2 Phases Negotiation protocol Model

Asynchronous

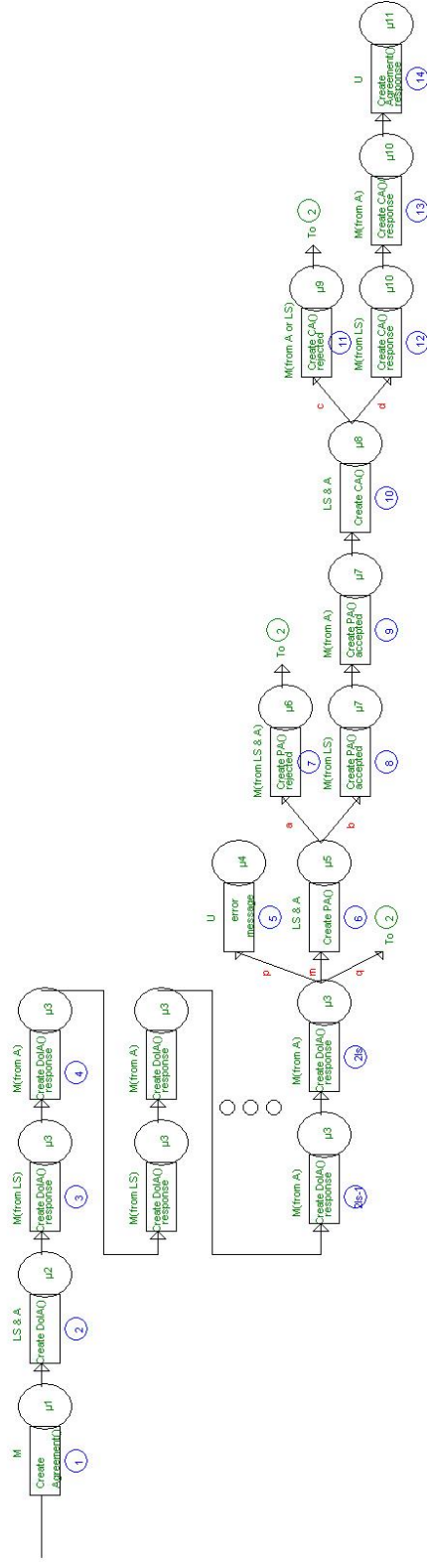


Synchronous

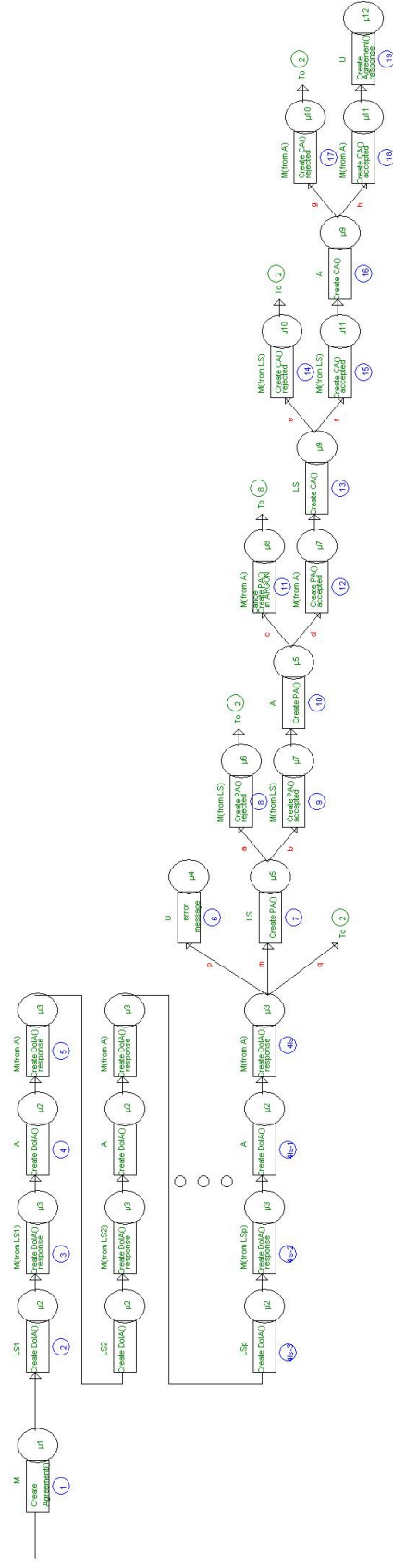


3 Phases Negotiation protocol Model

Asynchronous



Synchronous



Service Time & Error Probability example

After a few measures on a P4@2.8GHz

- Receive an XML message 1 ms
 - Check message validity..... 0.002 ms
 - Parse the message 0.715 ms
-
- +1.717ms

Probability to cancel a reservation : 10%

Probability to need to look after scheduling horizon : 10%

Those values can be modified to take more realistic values.

Limiting Queue

For example (3 Phases Asynchronous) :

$$\rho_2 = \frac{\lambda / \mu_2}{1 - \underbrace{(Q + (M(A + CB)))}_{\text{Service rate}}}$$

Arrival rate λ / μ_2

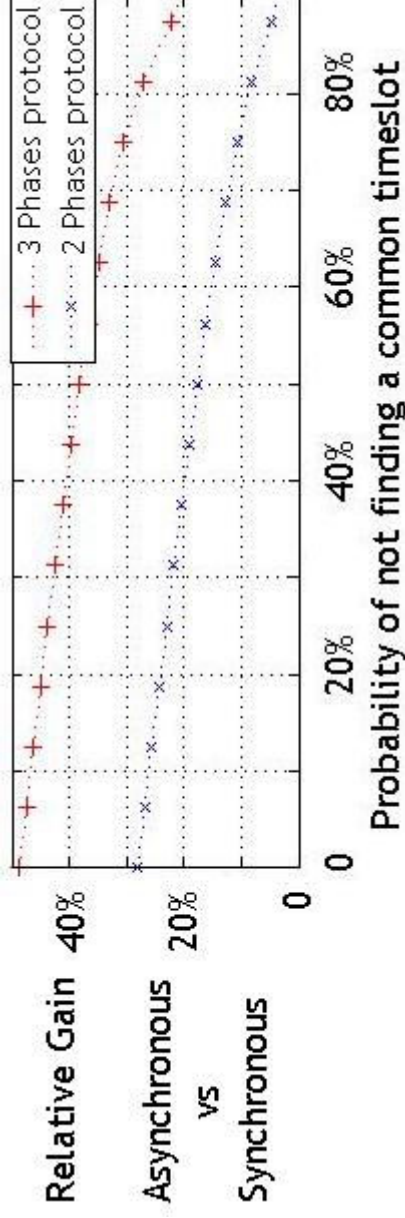
Probability to loop or to get an error

$$N = \sum_{k=1}^{NbQueues} \frac{\rho_k}{1 - \rho_k}$$

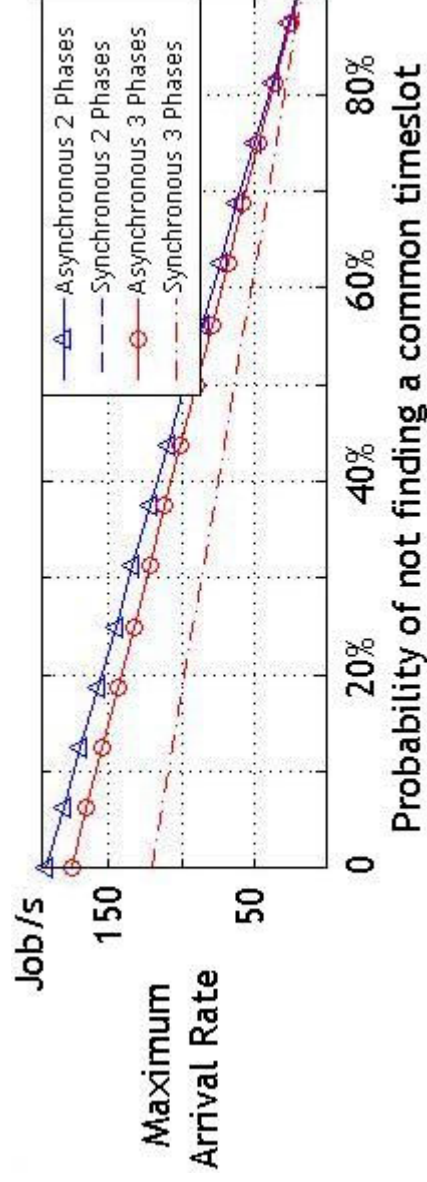
Load of Queue k ρ_k

3 Results

Maximum Job request Arrival Rate in the Meta-Scheduler

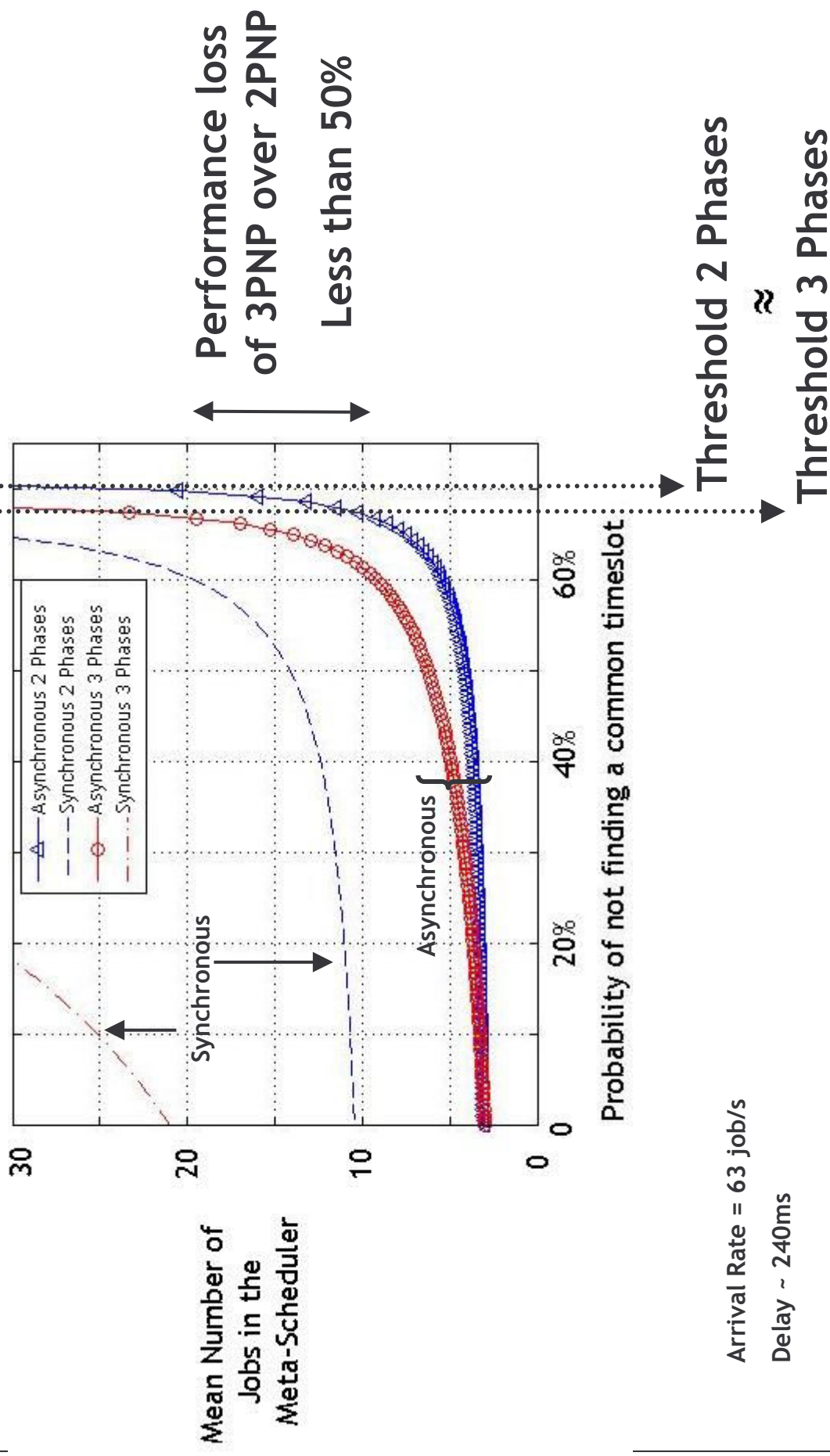


**NEED for
ASYNCHRONOUS
Implementation**



**Almost NO IMPACT on
Arrival Rate**

Mean number of Jobs in the Meta-Scheduler



Arrival Rate = 63 job/s

Delay ~ 240ms

xPNP : x Phases Negotiation Protocol

Conclusions

Need for an Asynchronous implementation (obvious)
Performance loss of a 3 Phases over a 2 Phases Negotiation protocol is

- less than 10% on the maximum job request arrival rate
- less than 50% on the total job request processing delay and memory requirement

References

Article to be published soon...
by the end of the year

For a similar model used in a different context :

Gurbani V.K., Jagadeesan L., Mendiratta V.B., “Characterizing session initiation protocol (SIP) network performance and reliability”, *International service availability symposium*, April 2005

www.alcatel-lucent.com

antoine.pichot@alcatel-lucent.fr