

# CS4261/5461: Assignment for Week 11 Solutions

Due: Sunday, 9th Nov 2025, 11:59 pm SGT.

---

1. (a) Note that  $a$  receives 3 votes,  $b$  receives 4 votes,  $c$  receives 3 votes, and  $d$  receives 2 votes. Hence, the only AV committee is  $\boxed{\{a, b, c\}}$ .  
(b) The committees  $\boxed{\{a, c, d\}}$  and  $\boxed{\{b, c, d\}}$  cover all voters. They are also the only committees of size 3 to do so, since choosing  $c$  and  $d$  is necessary due to voters 4 and 6.  
(c) The utilities of the voters for  $\{a, b, c\}$  are 3, 3, 2, 1, 1, 0, respectively. Hence, the PAV score of this committee is  $2\left(1 + \frac{1}{2} + \frac{1}{3}\right) + \left(1 + \frac{1}{2}\right) + 2 \cdot 1 = \boxed{\frac{43}{6}}$ .  
(d) Each voter begins with a budget of  $k/n = 1/2$ . First,  $b$  has the highest number of approvals and is chosen. Voters 1, 2, 3, 5 pay  $1/4$  each, and has a budget of  $1/2 - 1/4 = 1/4$  left. Next,  $a$  and  $d$  are not affordable while  $c$  is affordable, so  $c$  is chosen. Voters 1, 2, 4 pay their remaining budget ( $1/4, 1/4, 1/2$ , respectively). No more candidate is affordable, so  $a$  is chosen ahead of  $d$  by the approval score tie-breaking. Hence, the only MES committee is  $\boxed{\{a, b, c\}}$ .
2. Assume for contradiction that a committee  $W$  returned by GreedyCC fails JR. This means there exists a cohesive group of voters  $S$  such that no member of  $S$  has an approved candidate in  $W$ . Let  $x$  be a candidate approved by all voters in  $S$ . Since  $|S| \geq n/k$ , adding  $x$  would have added additional coverage of at least  $n/k$  at any stage throughout the execution of GreedyCC. Since GreedyCC did not select  $x$ , in each step it must have selected a candidate with additional coverage at least  $n/k$ , so the total coverage of  $W$  is at least  $(n/k) \cdot k = n$ . Hence, every voter has an approved candidate in  $W$ , contradicting the earlier assumption that no member of  $S$  has an approved candidate in  $W$ .
3. From the condition that  $S$  is a  $t$ -cohesive group of voters, we have that  $|S| \geq t \cdot \frac{n}{k}$  and the voters in  $S$  commonly approve  $t$  candidates. Since  $W$  satisfies EJR, there exists a voter  $i_1 \in S$  such that  $u_{i_1}(W) \geq t$ . Remove  $i_1$  from  $S$ , and repeat the argument. At any stage, if  $|S| \geq \ell \cdot \frac{n}{k}$  for some nonnegative integer  $\ell \leq t$ , then  $S$  is  $\ell$ -cohesive, and EJR implies that there exists a

voter  $i \in S$  such that  $u_i(W) \geq \ell$ . Hence, we get the following utility guarantees for the voters in  $S$ :

$$\overbrace{t, t, \dots, t}^{\geq 1 \text{ time}}, \overbrace{t-1, t-1, \dots, t-1}^{n/k \text{ times}}, \overbrace{1, 1, \dots, 1}^{n/k \text{ times}}, \overbrace{0, 0, \dots, 0}^{n/k-1 \text{ times}}.$$

Observe that if we decrease one of the  $t$ 's to 0 and remove the remaining  $t$ 's (if any), the average of this sequence would be exactly  $\frac{t-1}{2}$ . Thus, the actual average of the sequence is greater than  $\frac{t-1}{2}$ .