

# Admin

HW1 and proposal feedback released

- Check your grades
- Generally, not regrading unless there are calculation errors

HW2 still being finalized

- Due last day of reading week

Project due last day of reading weeks

# Lecture 10: Scaling up

## Part II

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But where's the **machine learning**?

Not the **focus** of this class,  
But we will **use ML to solve games**

# Scaling up 2p0s game solvers

## Games with perfect information

- A little bit of history
- AlphaGo, Muzero, AlphaZero

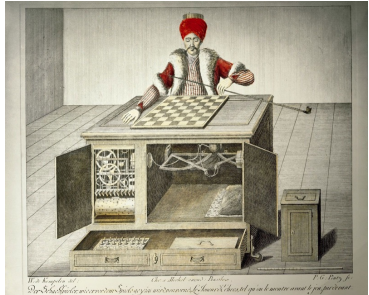
## Games with imperfect information

- **Subgame solving/search [Focus of Today's Lecture]**

# 2p0s Perfect Information Games

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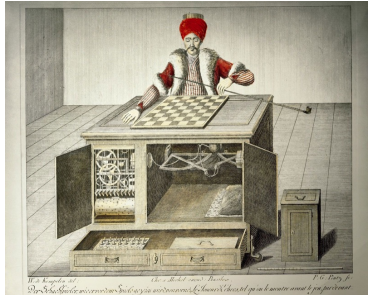
## Mech Turk [Chess]



~1770

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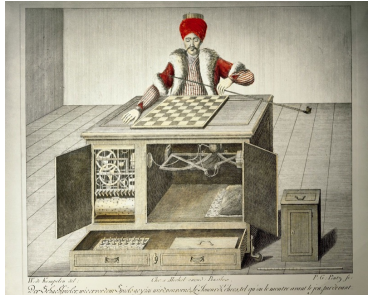


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look, magic!

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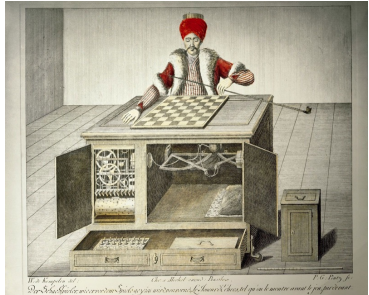
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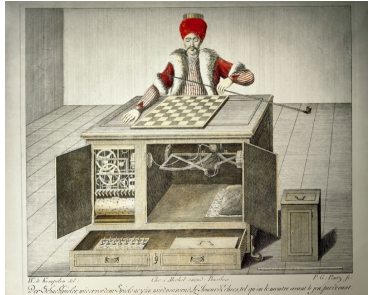
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**look, magic!**

**Harnessing the  
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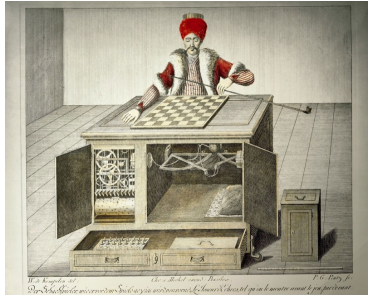
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**Pattern  
recognition**

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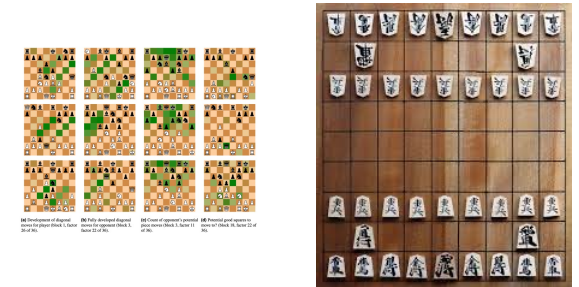
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AlphaGo  
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2016

MuZero  
[many games]



2019

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2019

General  
game-playing

# What happens with huge game trees?

Full tree search requires time linear in \_\_\_\_ ?

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- Chinese Chess, Shogi, Go

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- Chinese Chess, Shogi, Go

Rely on **evaluation function**  $\tilde{V}(s)$

- At every state, tell me how “good” a state is from P1’s perspective, assuming both players play optimally from here onwards
- $\tilde{V}(s)$  approximates  $V(s)$  using some features of  $s$ 
  - Material, Pawn structure, math

# What happens with huge game trees?

Heuristics: be selective in depth/branching or early termination

- Quiescence search
- Null-moves
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Tradeoff between “high-tech” evaluation functions and searching deeper

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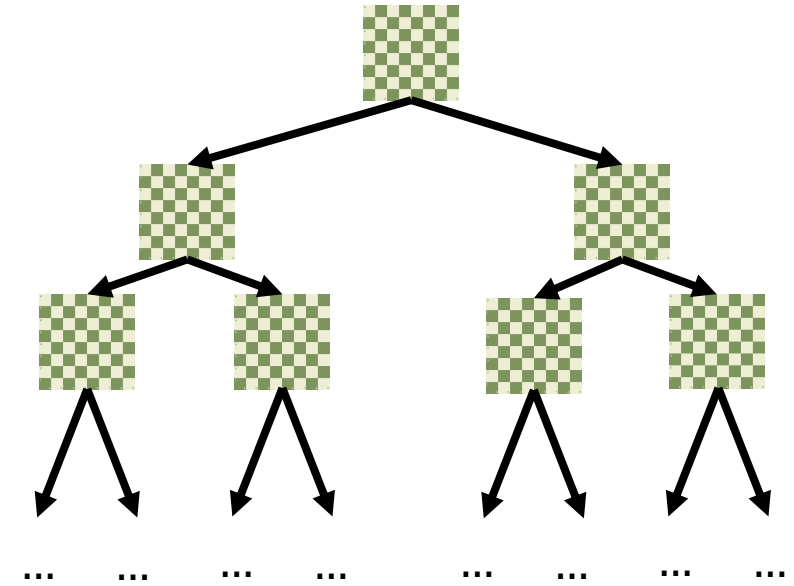
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Monte Carlo Tree Search

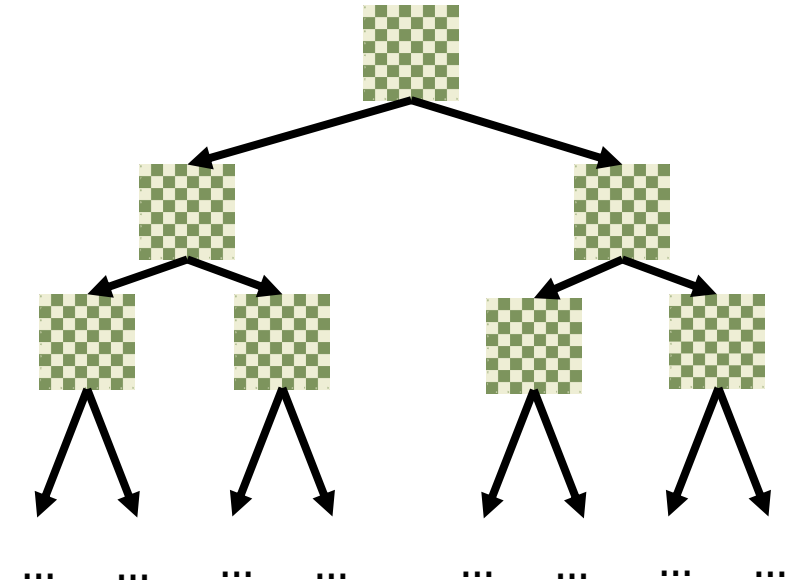
# Game Solving with Function Approximation



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Consider zero-sum games

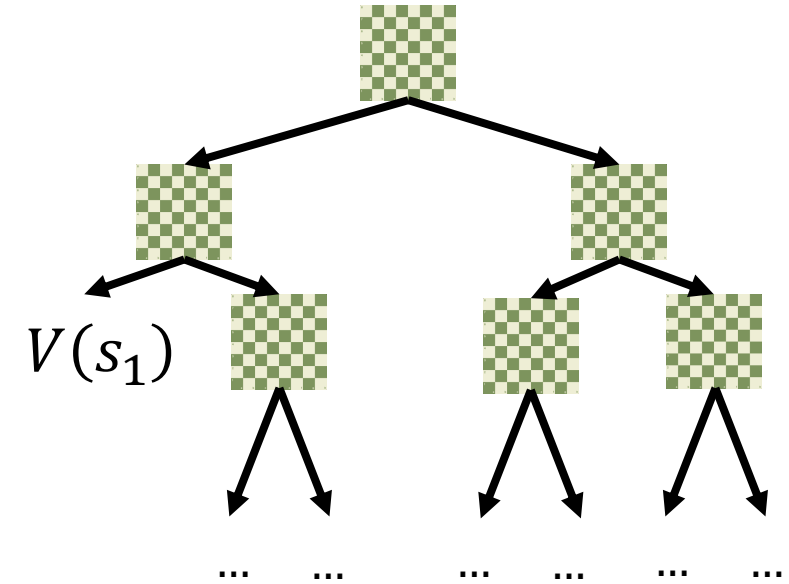
- Game tree too large to traverse explicitly
- **Value Function**  $V(s)$  approximates how “good or bad” each state is



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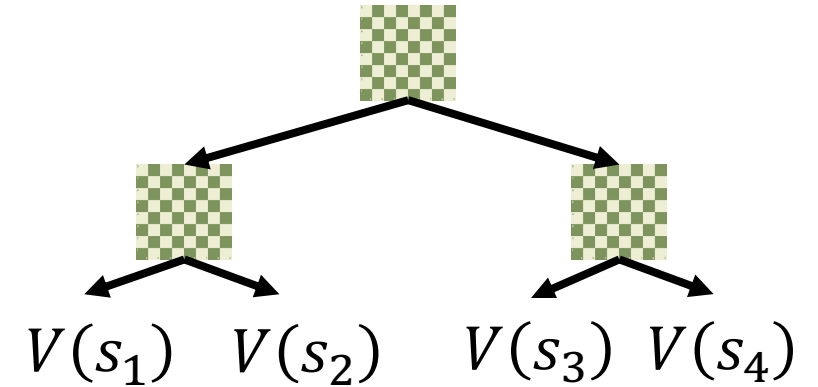




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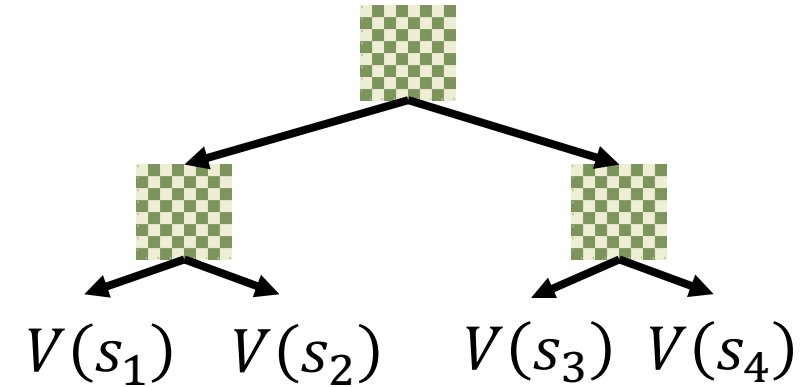
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Previously (e.g., Deep Blue)

- Handcrafted evaluation function  $V(s)$  based on heuristics from experts



# Game Solving with Function Approximation

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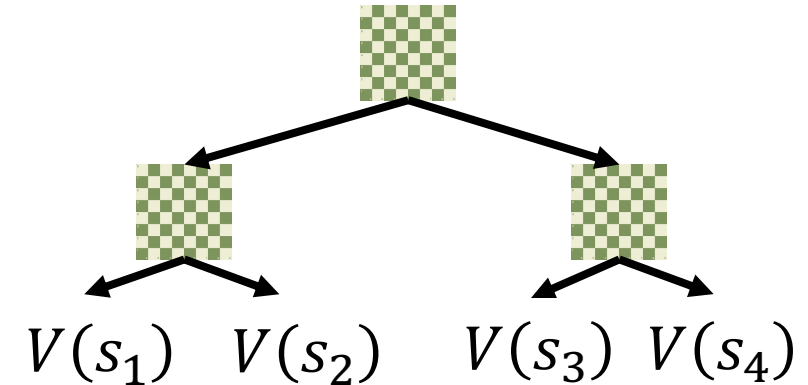
- Game tree too large to traverse explicitly
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Previously (e.g., Deep Blue)

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Today (e.g., AlphaGo/Zero, DeepStack)

- **Learn** how good each state  $s$  is, generalize to states not seen before
- $V_\phi(s)$  is a network parameterized by  $\phi$  approximating the value of  $s$



# Some observations

Relies on the idea of a value  $V(s)$  or  $Q(s, a)$  and their approximates

- Where did these come from, practically?
- Should we even be as accurate as possible?
  - Example:  $V(s)$  for chess is technically restricted to be in  $\{-1, 0, 1\}$ , yet in practice we still evaluate positions in terms of “centipawns”. Same goes for most win/loss games

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**Warning: hard to be  
true in practice**

What happens if value function is “slightly off?”, i.e.,  $|\tilde{V}(s) - V(s)| < \epsilon$ ?

- Will a deeper search always yield a better immediate action at the root?
  - No (see Nau 1982, Wilson et. al., 2018) but very rare in practice
  - Looking ahead deeper is probably the right thing to do
- Are there any other guarantees?
  - Yes, though a lot weaker, we’ll see that later

# Subgame Decomposition

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# What about imperfect information EFGs?

Let's start with the simplest task

Can we even **store** or summarize  
“values” in subgames?

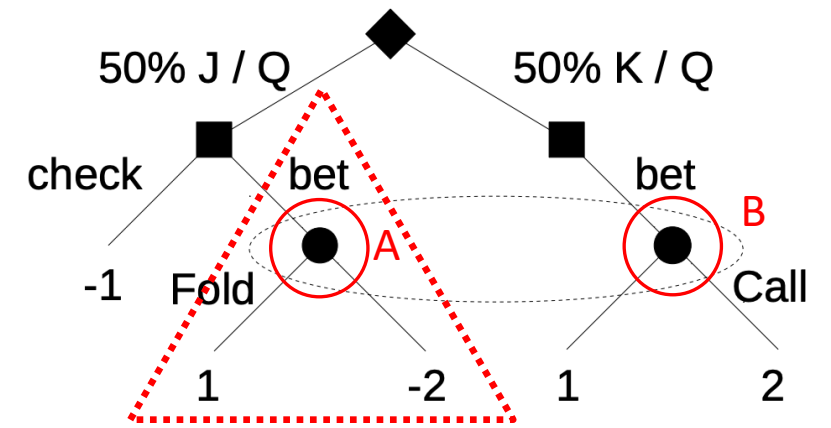
# Attempt 1: Replacing w/ subtree value

Consider subtree rooted at A (or B)

- Take EFG formed by the sub-EFG  $G'$  rooted at A (or B)
- If  $s, s'$  in same infoset in  $G$  and are both in  $G'$ , then they are the same infoset in  $G'$  (and vice versa)
- Solve  $G'$  to get value of vertex  $v$

Replace payoffs at by  $v_A$  (and  $v_B$ )

Does this work?



$v_A = -2$



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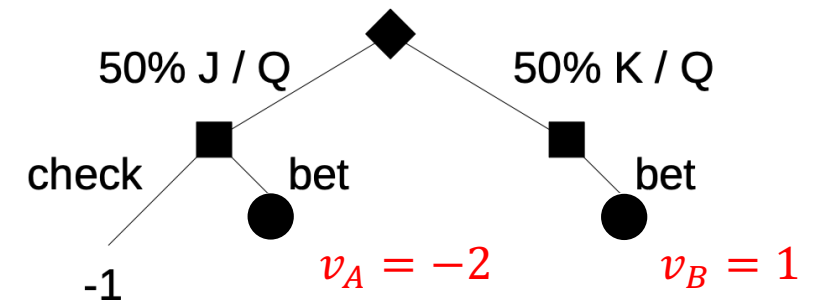
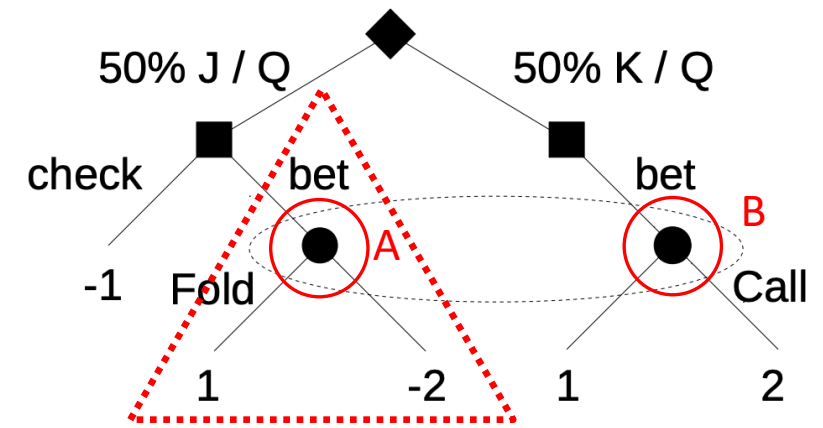
Replace payoffs at by  $v_A$  (and  $v_B$ )

Does this work?

Recall true Nash has value  $1/3$

No. New value under “truncated” game is 0

- Information set at the root is not respected
- In this instance, P1 is “respecting” P2 too much



# Attempt 2: Replace w/ “Nash node value”

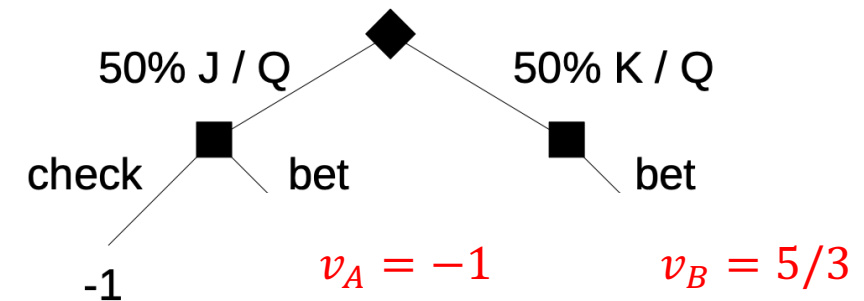
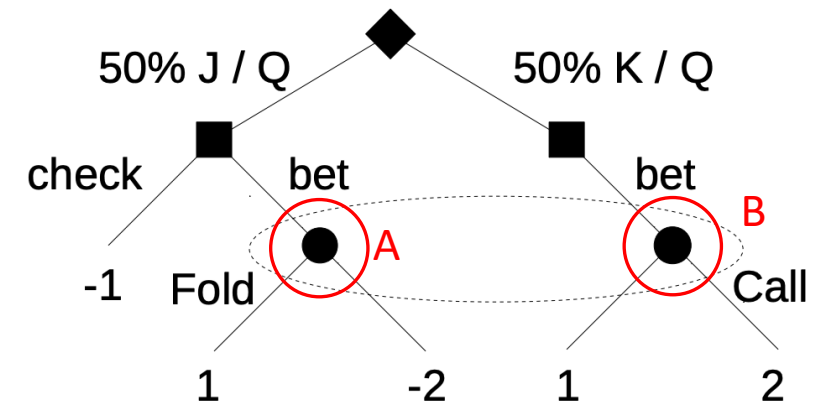
P2's strategy (at both A, B) is

- Fold w.p. 1/3, Call w.p. 2/3 (make sure you know why!)
- “Value” at A is  $1/3 - 4/3 = -1$

By this logic, P1 can do *anything* if it got J

- Value calculated here is 2/3
- This is **higher** than in the true game!

What went wrong here?



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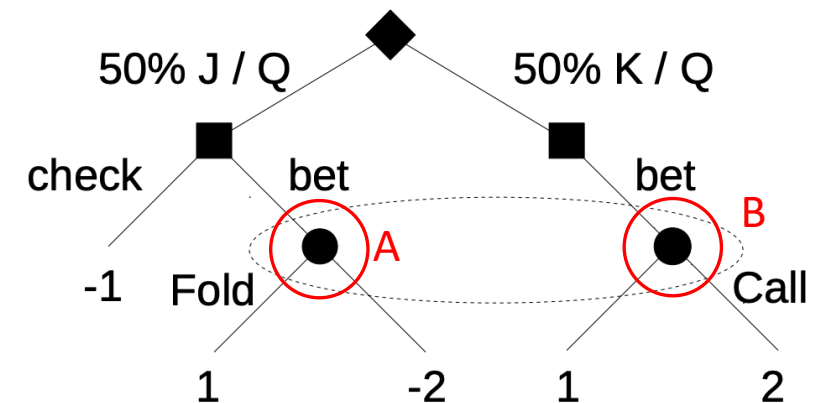
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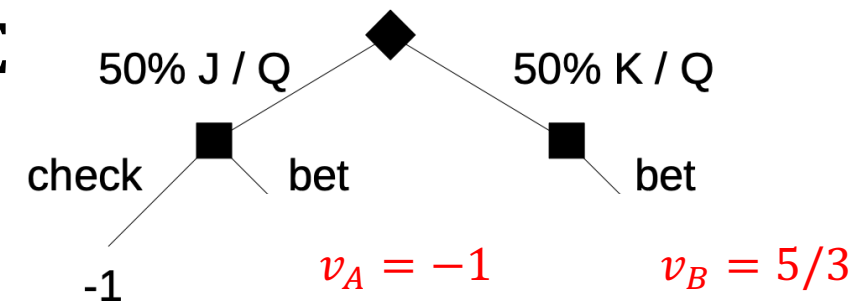
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We assumed that P2 continues playing its NE



# Important reminder!

That a strategy does well against the opponent's NE is **not** a sign that your strategy is optimal, or even good

- Holds even for normal-form games
- Need to consider **exploitability**

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Very, very, very common mistake that I see in papers all the time

- “My agent’s performance is good against an opponent that is not exploiting me”
- “My agent’s performance is good against an opponent that ... I myself trained”

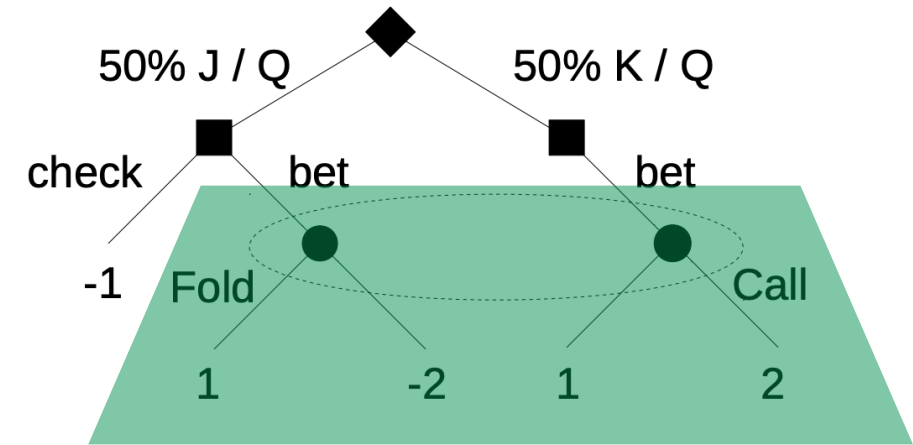
What makes subgame solving in EFG tricky

Need to deal with both imperfect information **and** the fact that there is another player is taking advantage of this!

# Insight #1: Subtrees are not enough

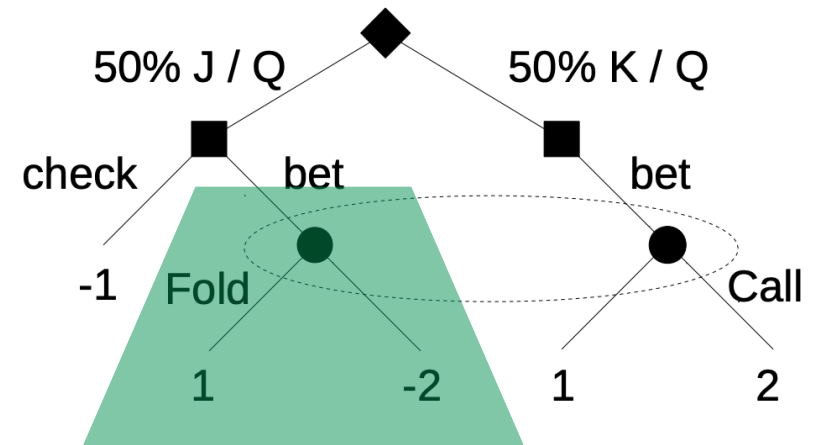
We need to look at **subgames**

- Closed under “in-same-infoset” relationship
- Collection of subtrees (forest) such that
  - all descendants are included AND
  - If  $s, s'$  are in same infoset, and  $s$  is in subgame, then  $s'$  is also in same subgame
- Intuitively: subgame cannot have infosets “poking out”



Subgame now “respects” infosets

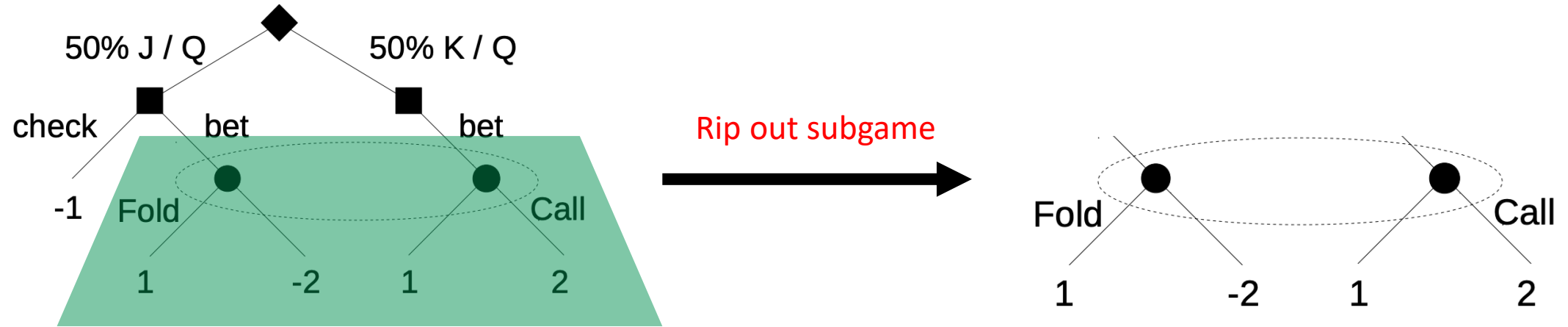
- But not yet enough --- we can’t solve a subgame “independently” from the rest of the game!



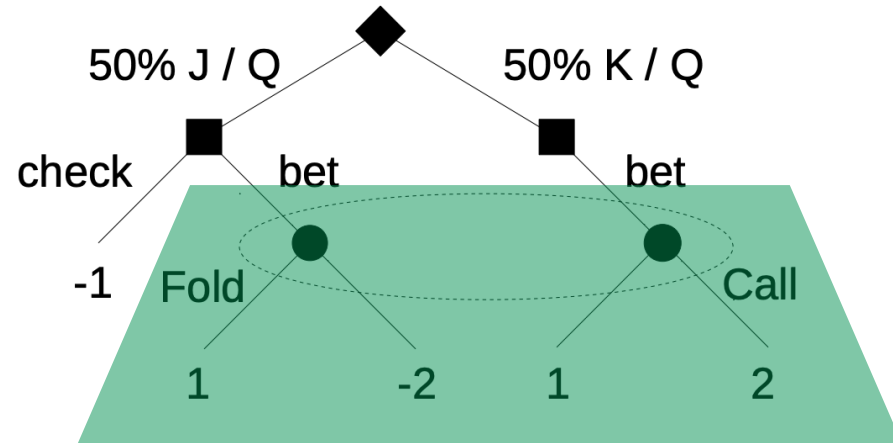
Not allowed!



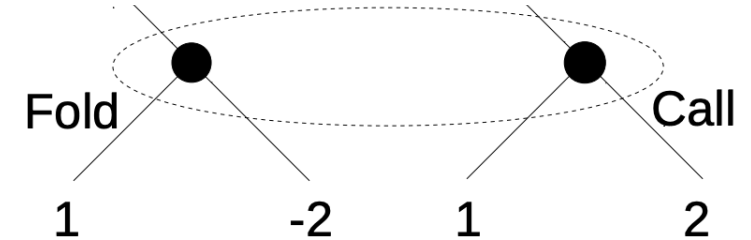
# Can we “solve” subgames alone?



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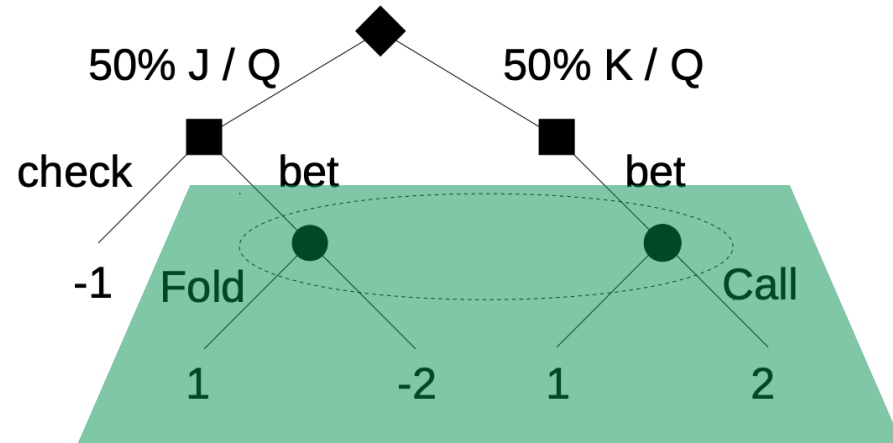


Rip out subgame

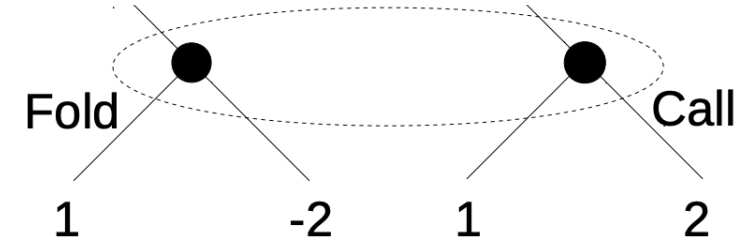


No. Why?

# Can we “solve” subgames alone?



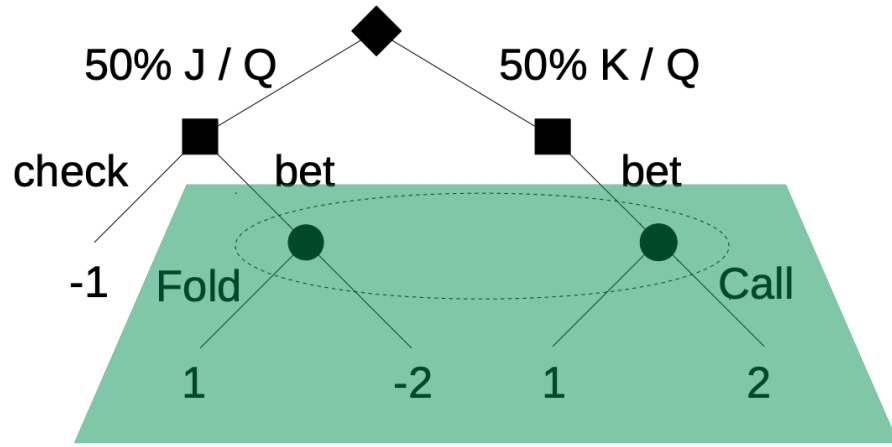
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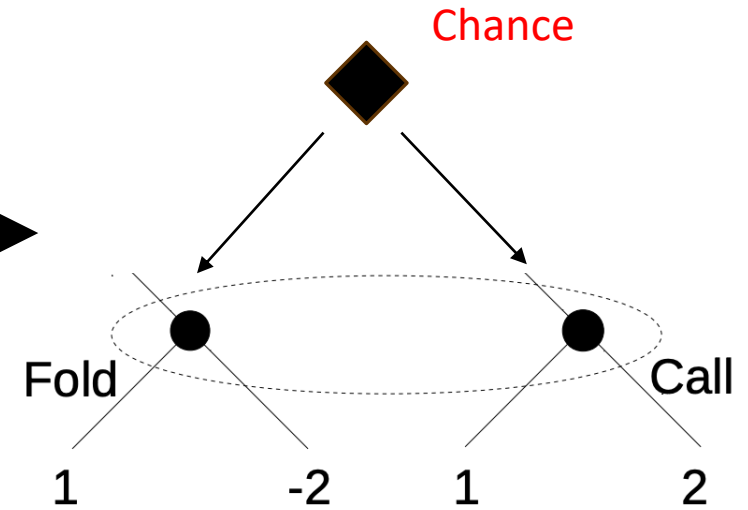
No. Why?

Where is the root?

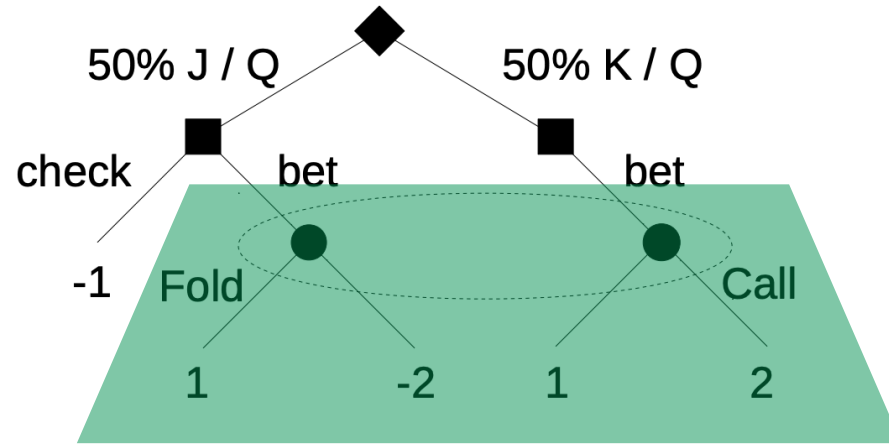
# Adding in chance nodes



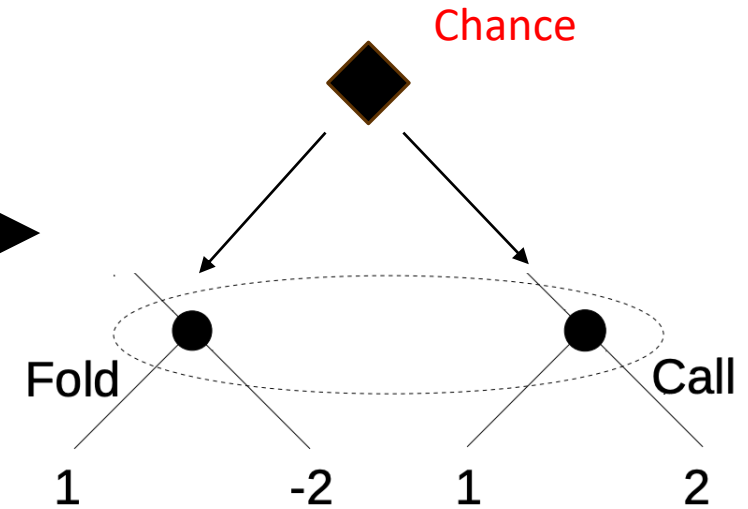
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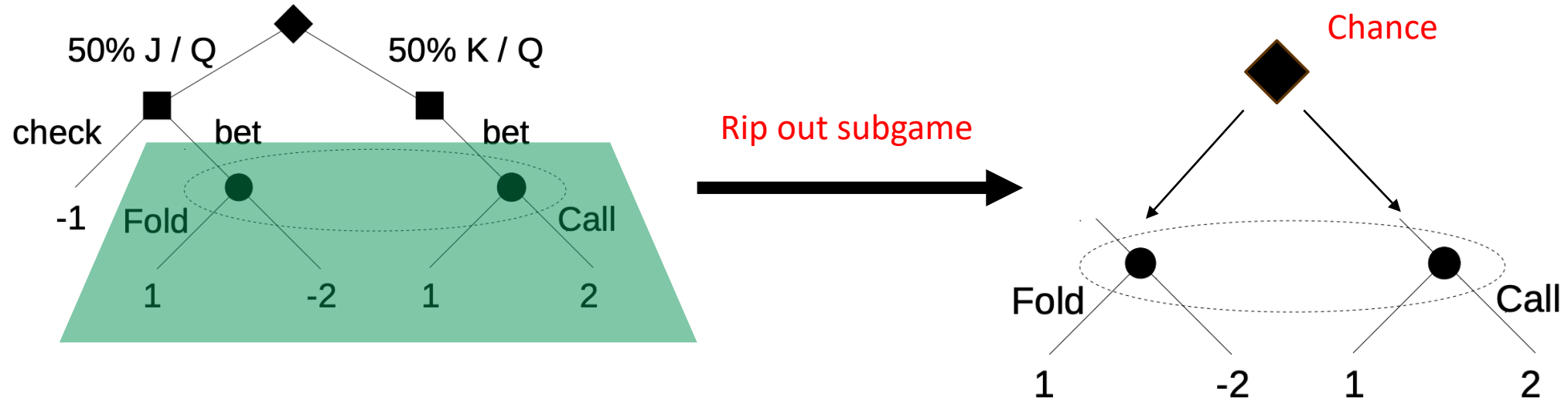


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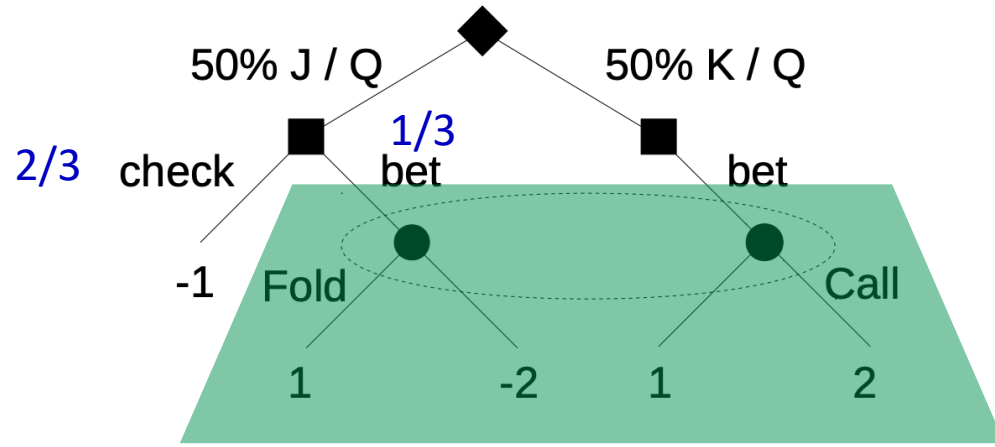


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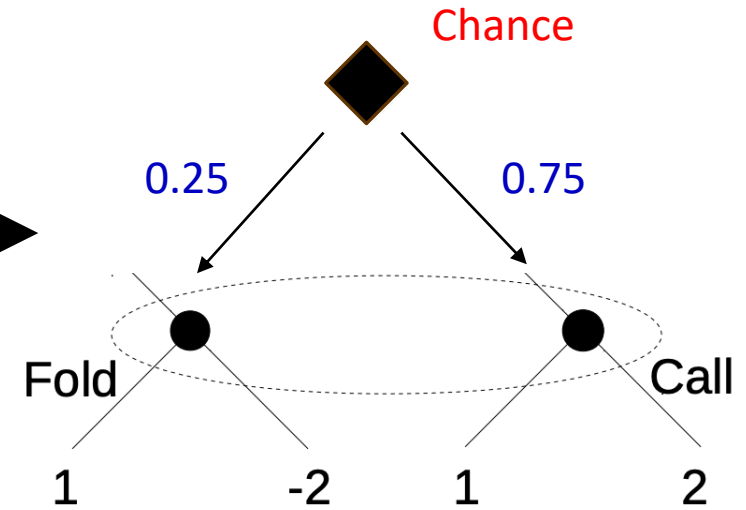
What are the chance probabilities?

- The probabilities dramatically affect solution of subgame!
- Do we use the probabilities of reaching **under Nash**?
- Or...?

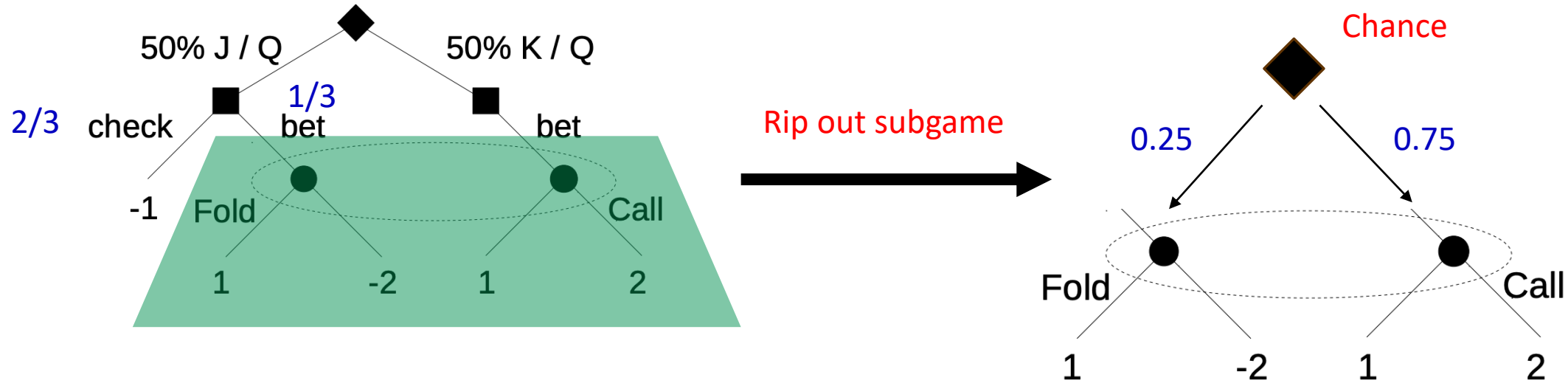
# Let's try using NE as probabilities



Rip out subgame



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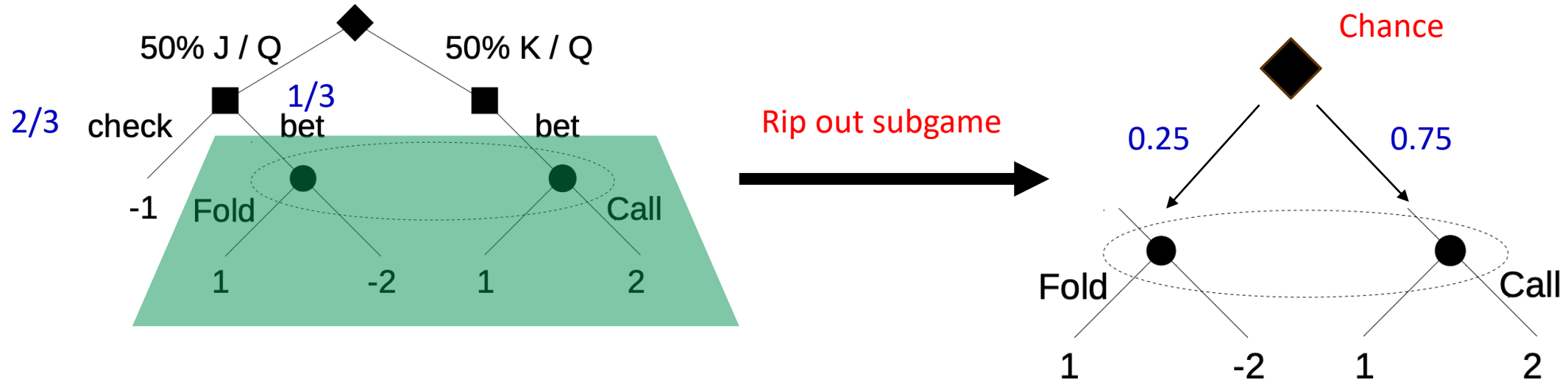


P2 is now indifferent between folding as calling (as expected)

- But... What can the solution of this game be useful for? Not much...
- **Any strategy** of P2 in subgame is a Nash!



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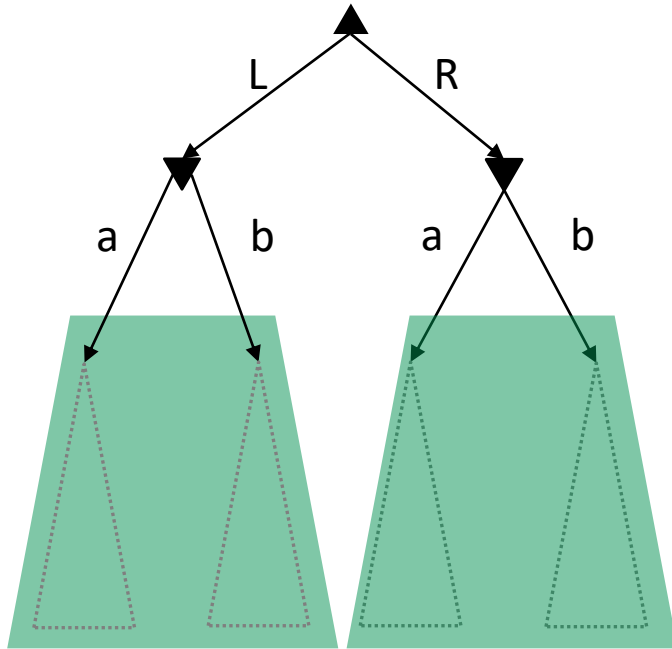
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Still a useful idea to keep in mind!

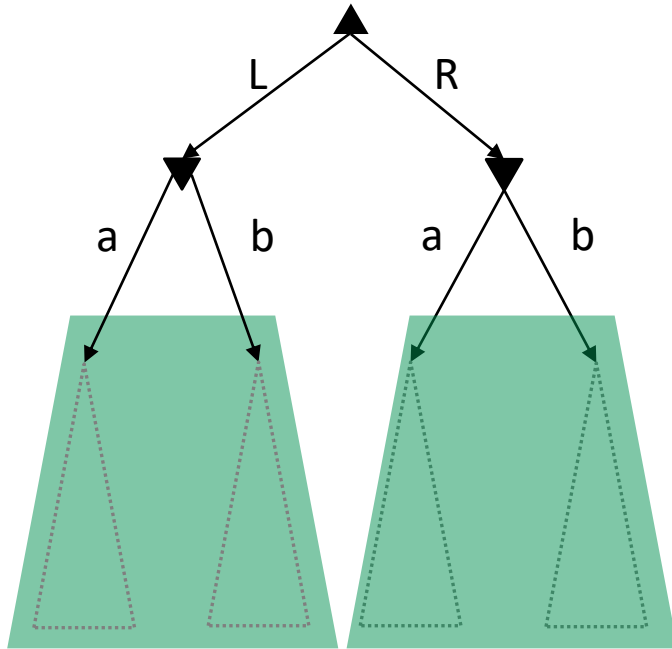
# Subgame Solving

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# Framework for Subgame solving

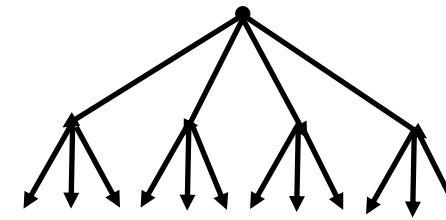


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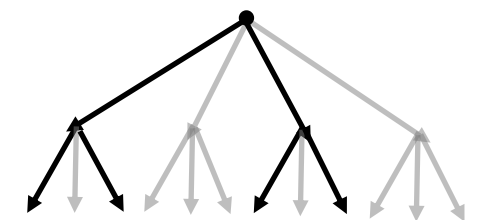


## Blueprint

- Can be any baseline strategy
- Typically, the 'best' strategy which can be computed offline based on some small, abstracted variant of game



Full Game

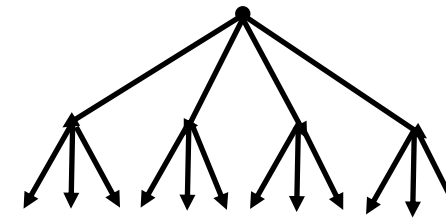
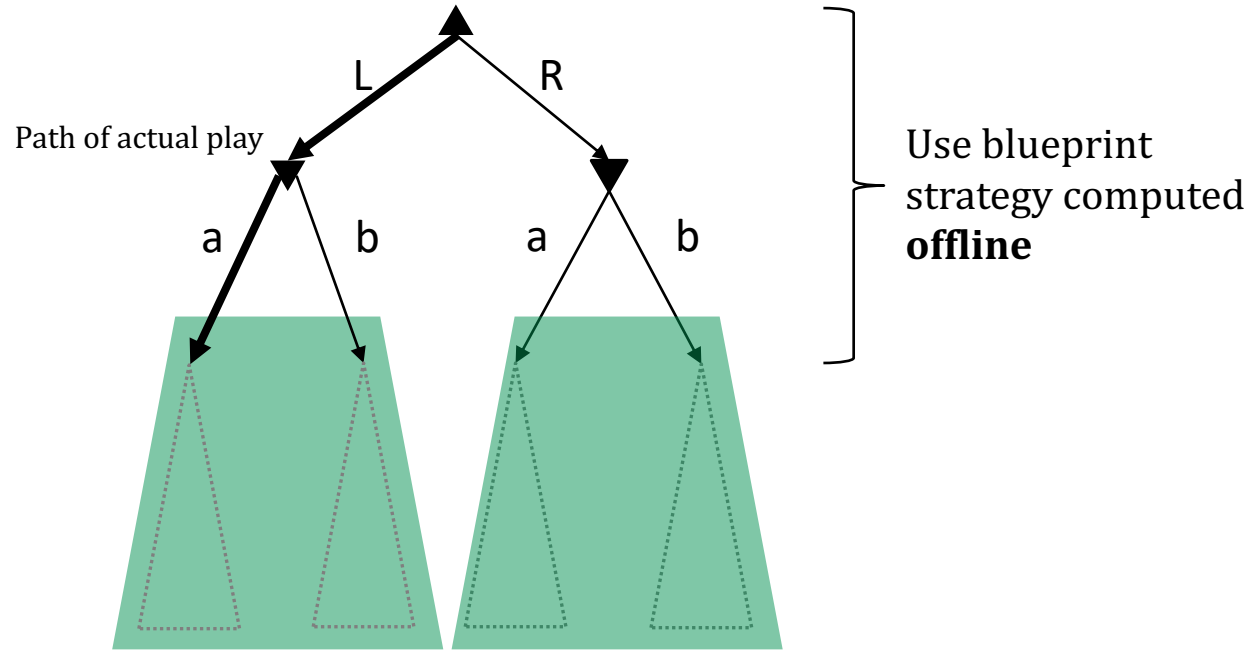


Abstraction to compute blueprint

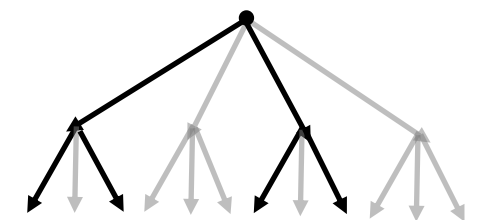
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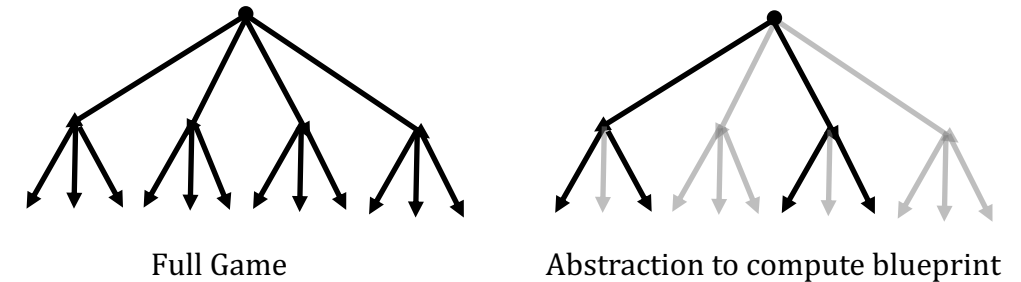
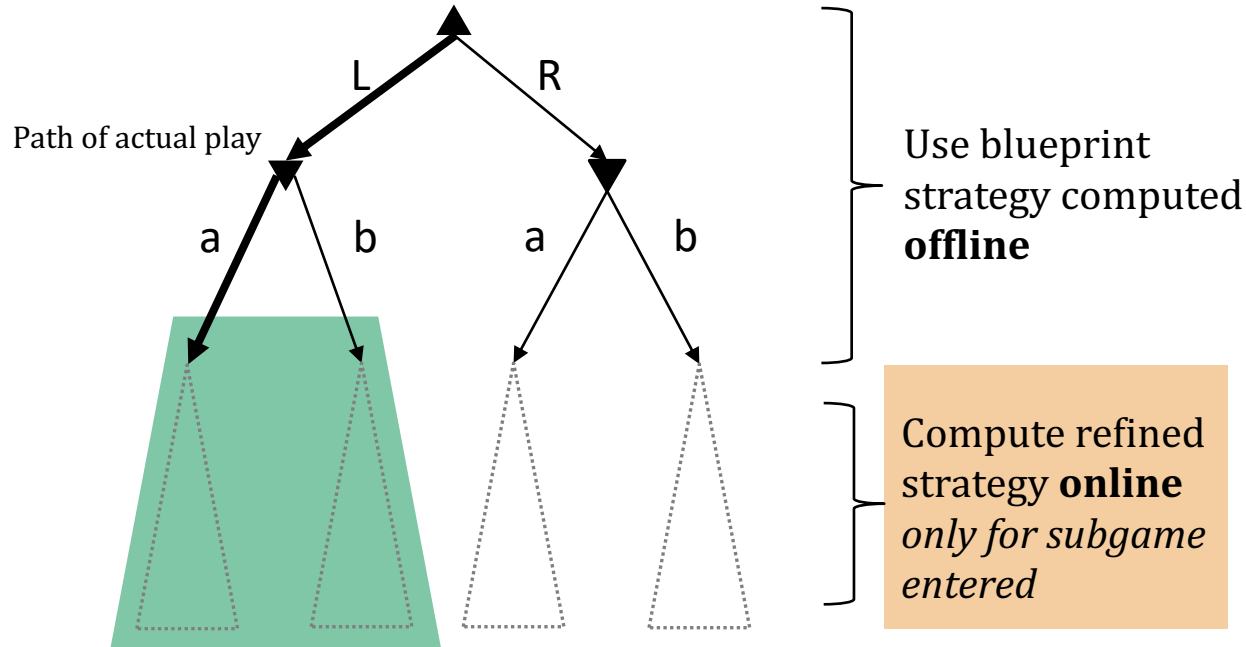


Abstraction to compute blueprint

# Framework for Subgame solving

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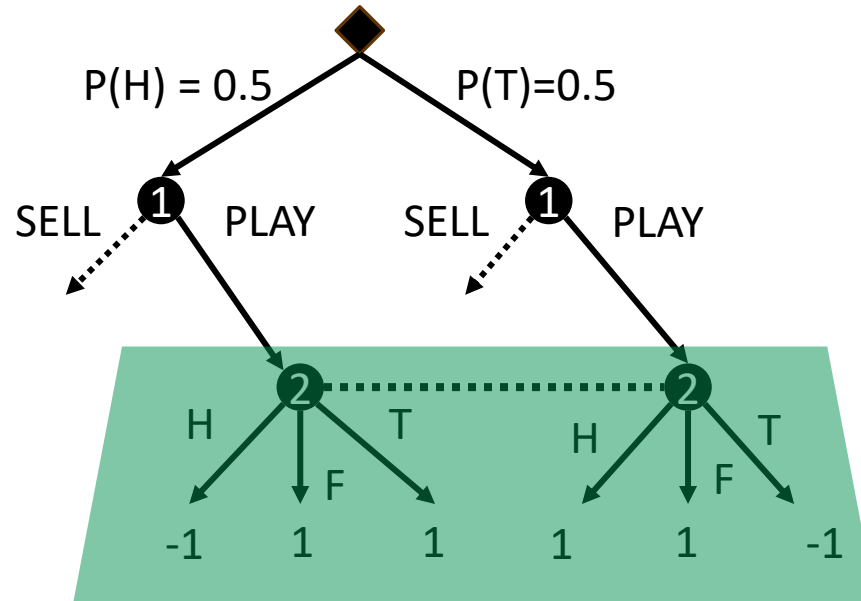
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## Refinement (of a blueprint)

- An attempt at improving P1's strategy in a subgame from the blueprint
- Followed for remainder of game

# Example: Coin Toss

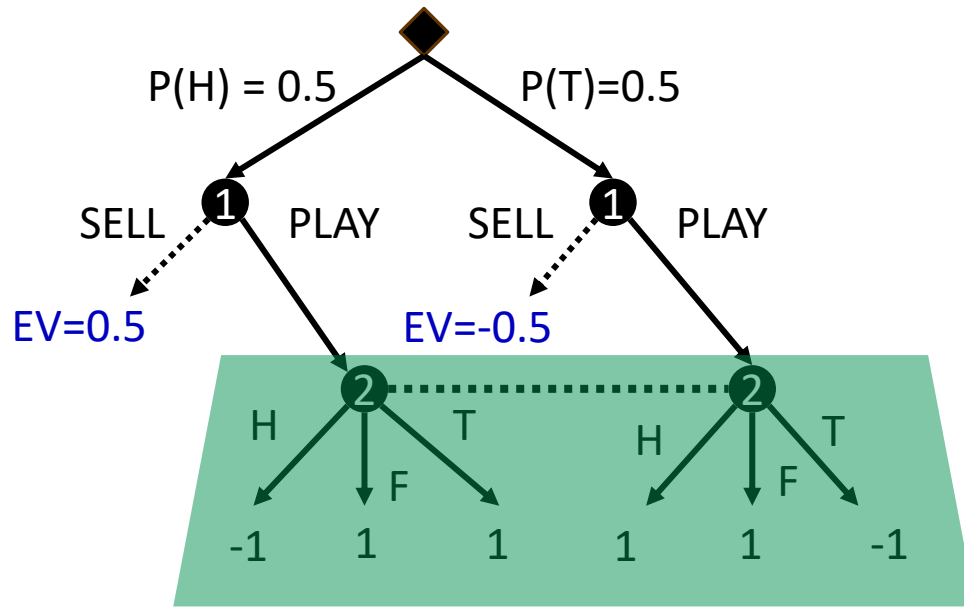


H=Heads  
T=Tails  
F=Forfeit

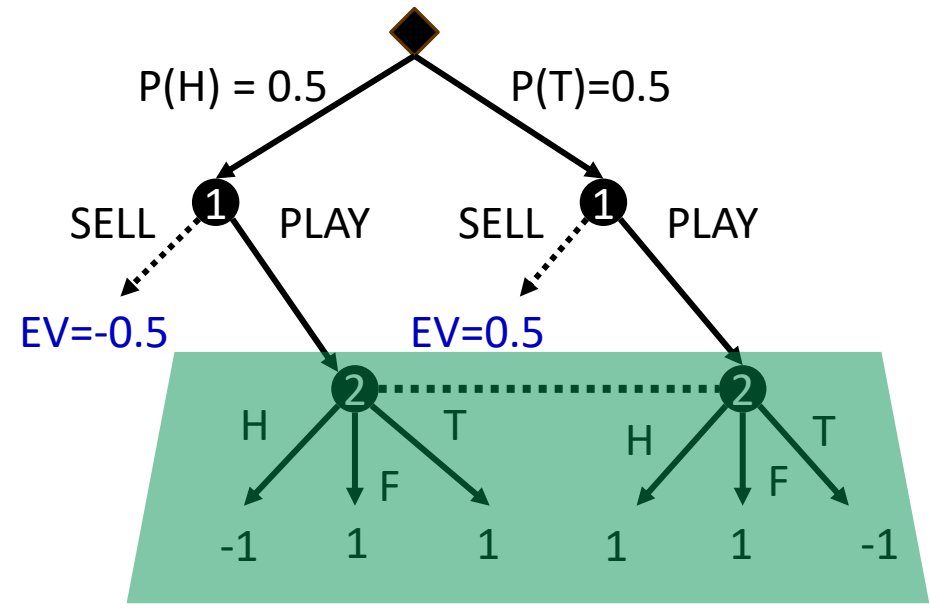
Rewards for for P1

NOTE: P2 is doing the subgame solving here (to follow the paper's convention)

# Subgame cannot be solved independently



Game 1

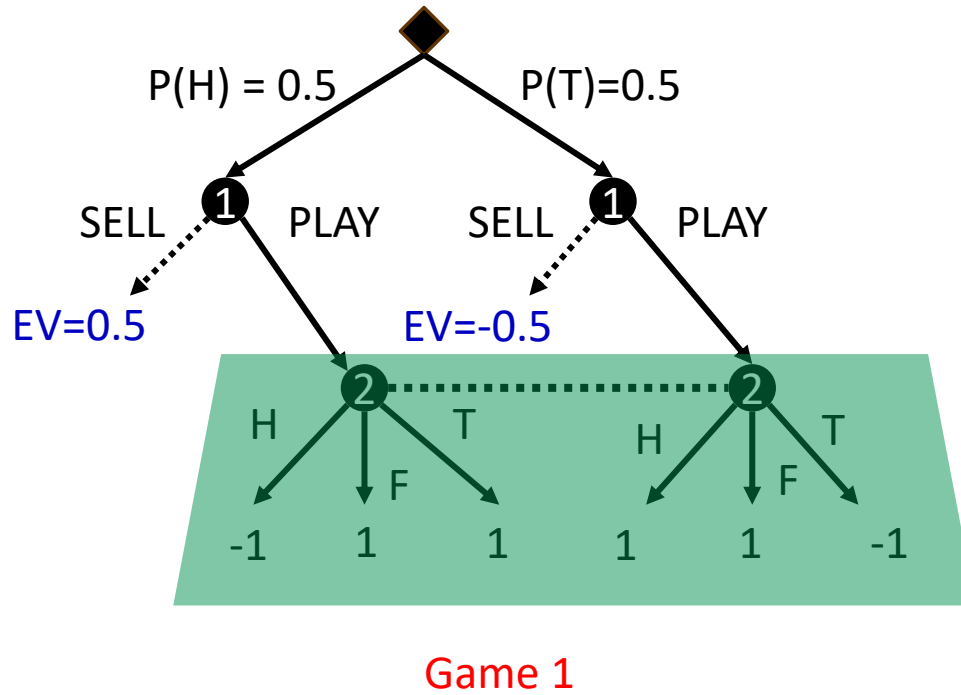


Game 2

H=Heads  
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F=Forfeit

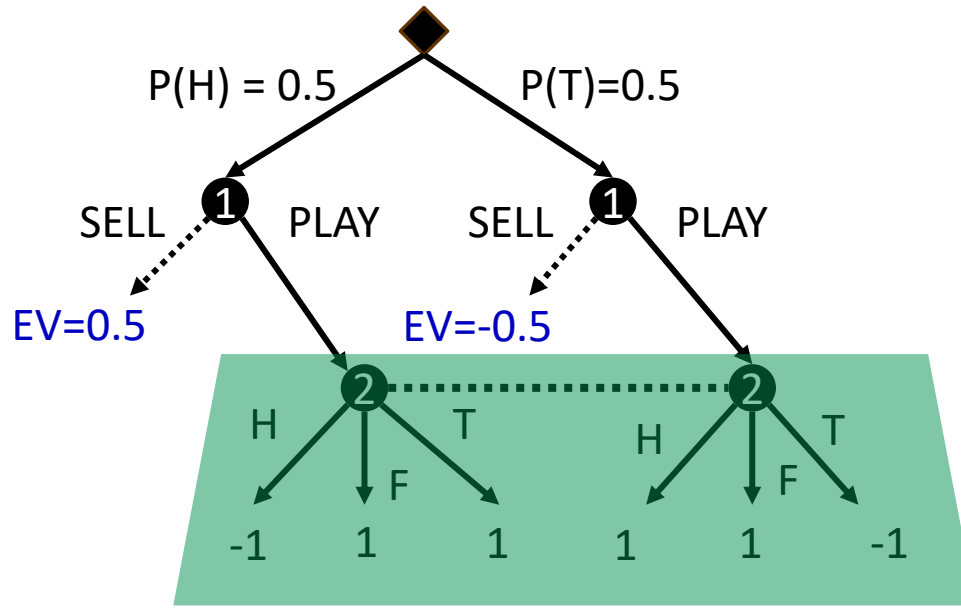


# Subgame cannot be solved independently II



# Subgame cannot be solved independently II

What is the NE of this game?

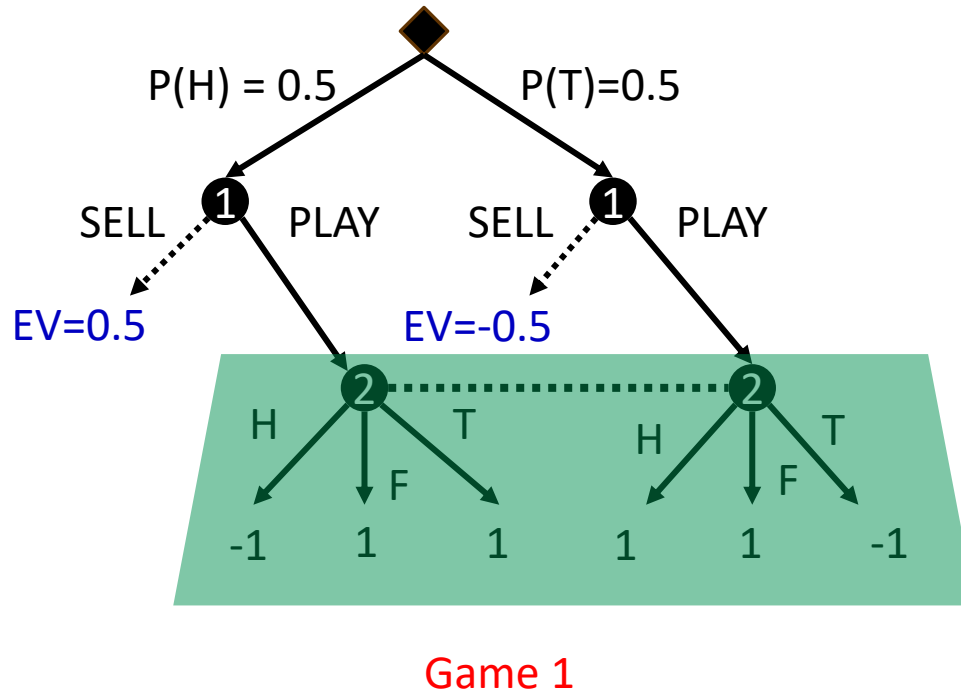


H=Heads  
T=Tails  
F=Forfeit

# Subgame cannot be solved independently II

What is the NE of this game?

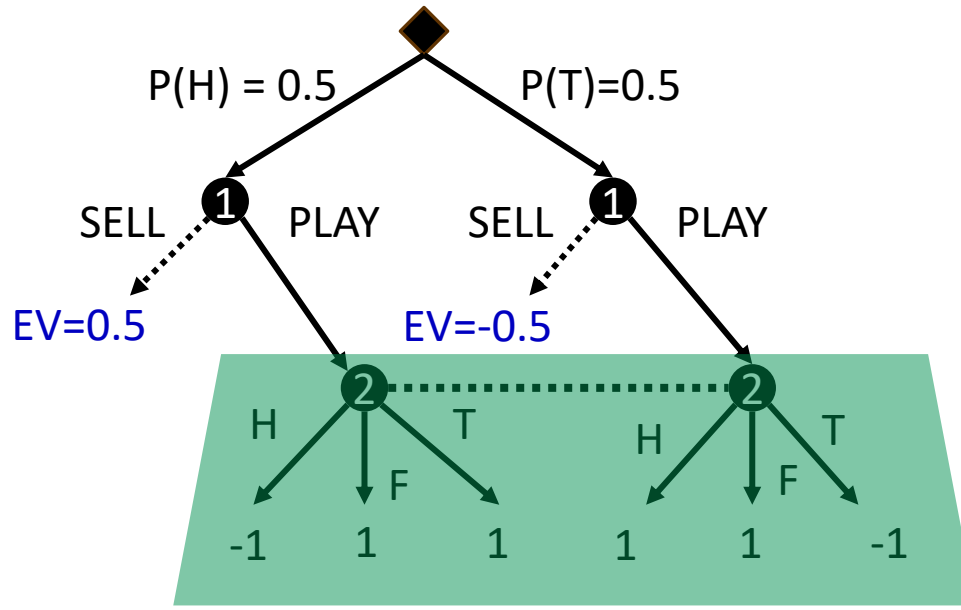
- If P2 plays heads, P1's BR will sell if coin is heads, and play if tails  $\rightarrow 0.75$  on avg



# Subgame cannot be solved independently II

What is the NE of this game?

- If P2 plays heads, P1's BR will sell if coin is heads, and play if tails  $\rightarrow 0.75$  on avg
- If P2 plays tails, P1's BR will sell if coin is tails, and play if heads  $\rightarrow 0.25$  on avg



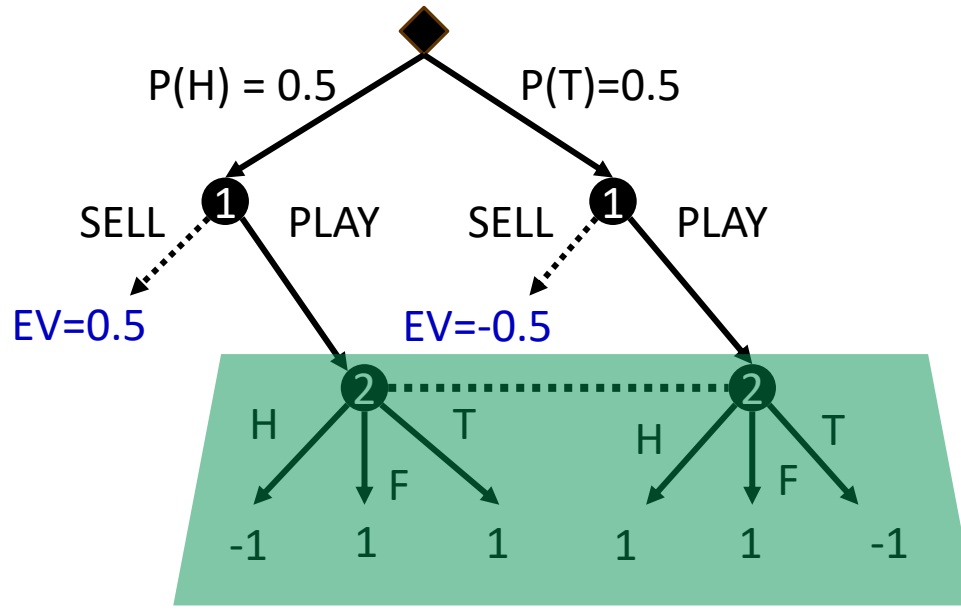
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T=Tails  
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Game 1

# Subgame cannot be solved independently II

What is the NE of this game?

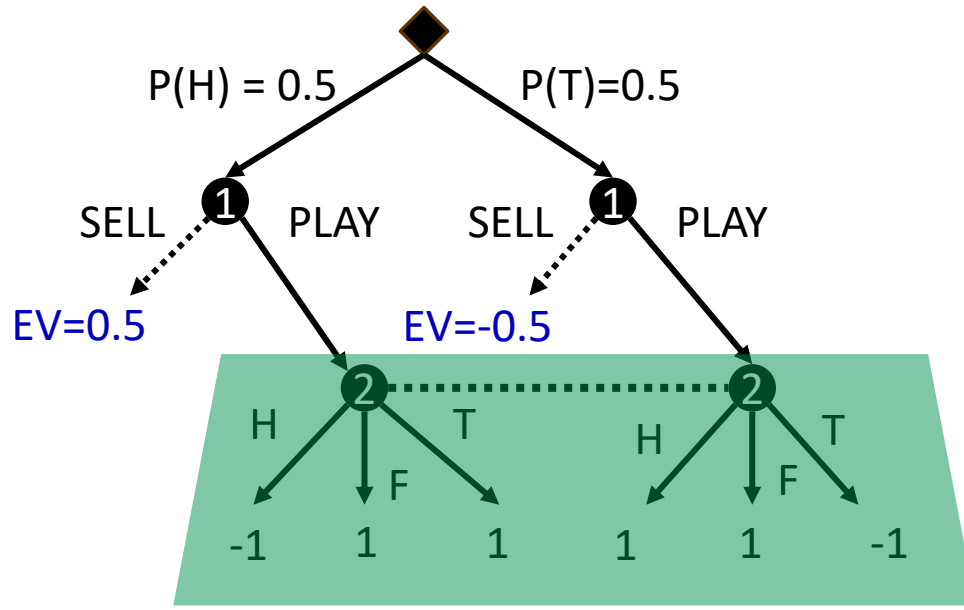
- If P2 plays heads, P1's BR will sell if coin is heads, and play if tails  $\rightarrow 0.75$  on avg
- If P2 plays tails, P1's BR will sell if coin is tails, and play if heads  $\rightarrow 0.25$  on avg
- NE is for P2 to play heads wp 0.25, tails 0.75  $\rightarrow 0$  on avg (show!), better for P2



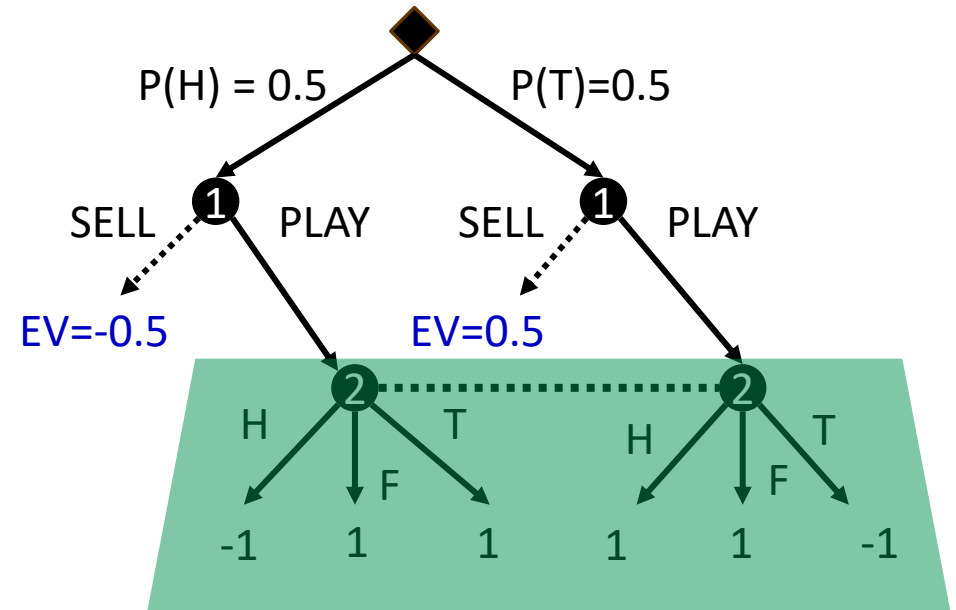
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Game 1

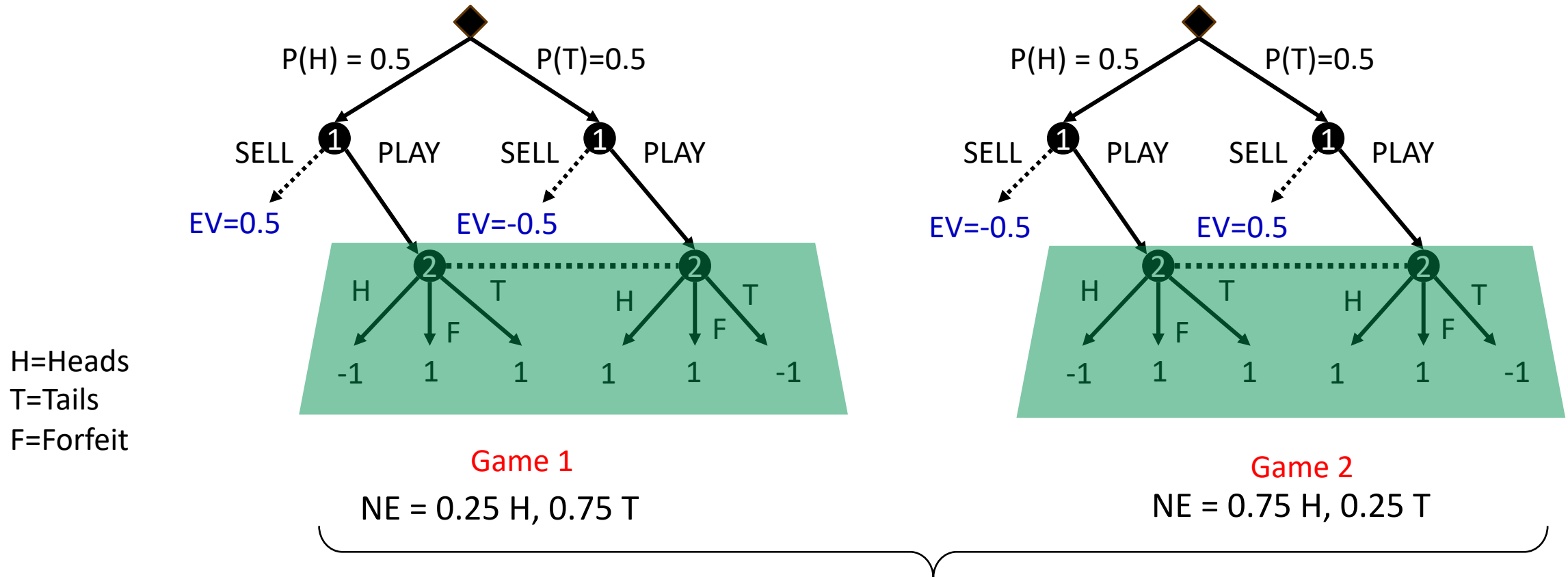
# Subgame cannot be solved independently III



H=Heads  
T=Tails  
F=Forfeit



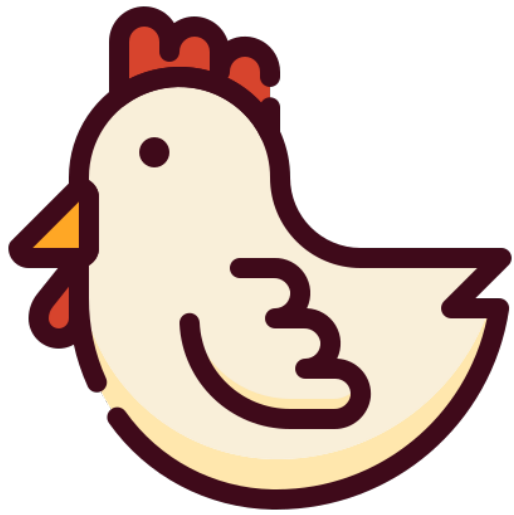
# Subgame cannot be solved independently III



Subgame is identical, but solution is different

Parts of the game **outside** subgame affects solution within subgame

# A chicken and egg problem...?



Optimal strategy in subgame

Depends on



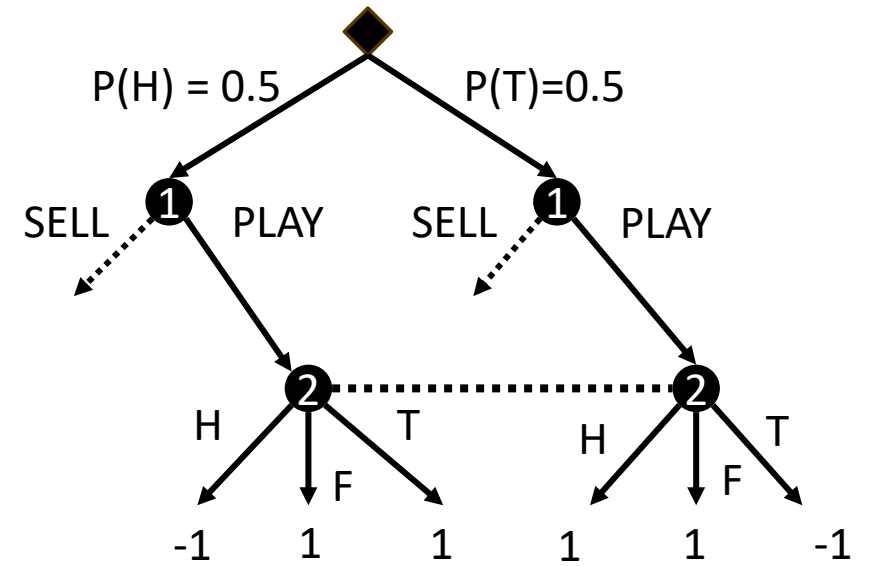
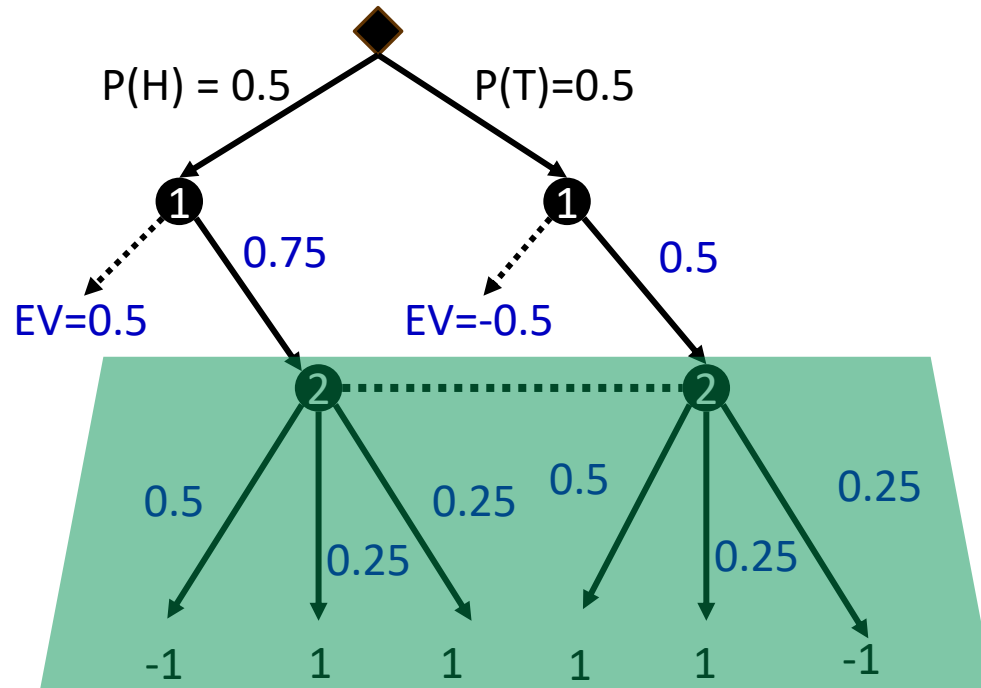
Best response of opponent  
(including in pre-subgame)



# Naïve/unsafe subgame solving

---

# Coin Toss Blueprint

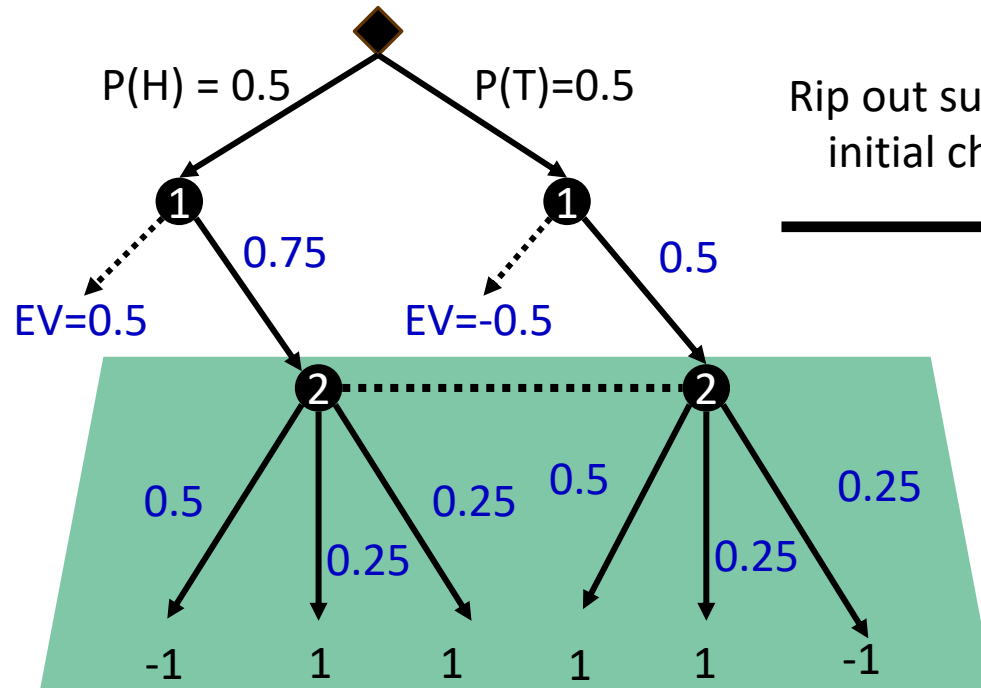


"Annotated Game"

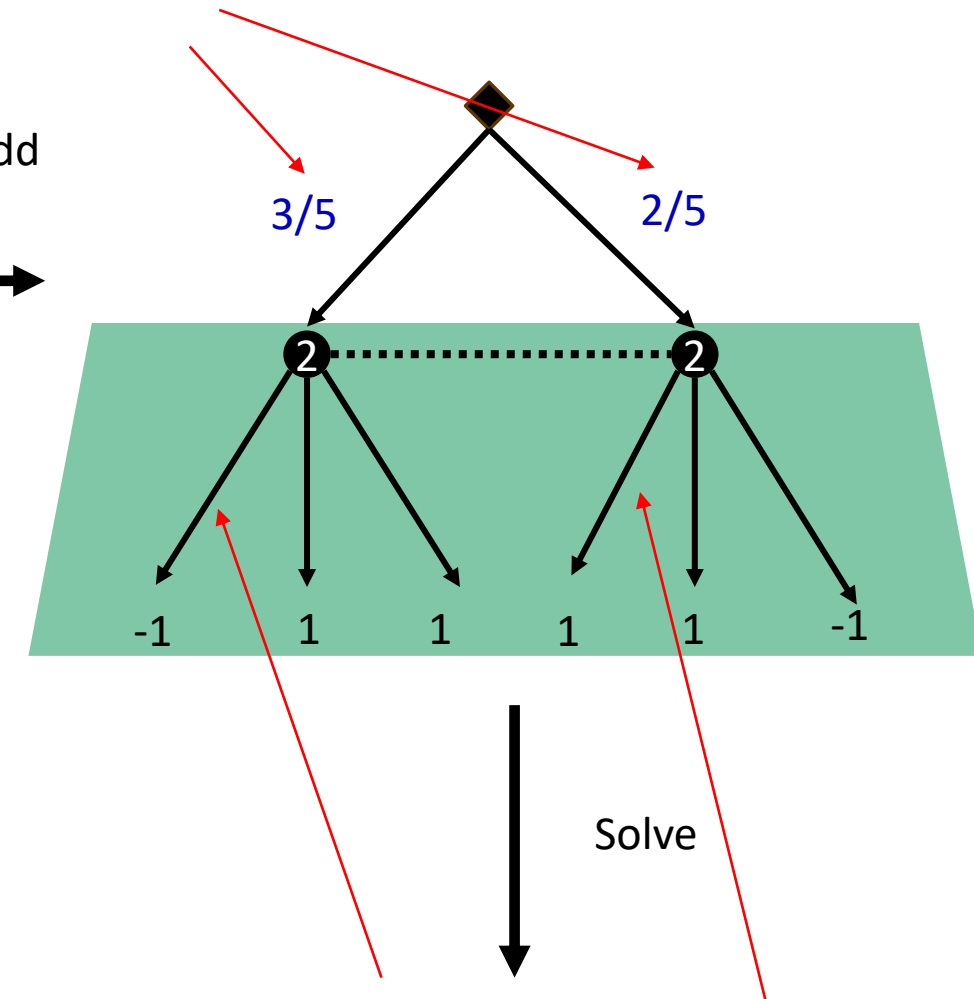
Note: Blueprint is obviously not optimal (it chooses to forfeit), but that is besides the point

# Naïve subgame solving

Normalized probabilities of reaching these nodes

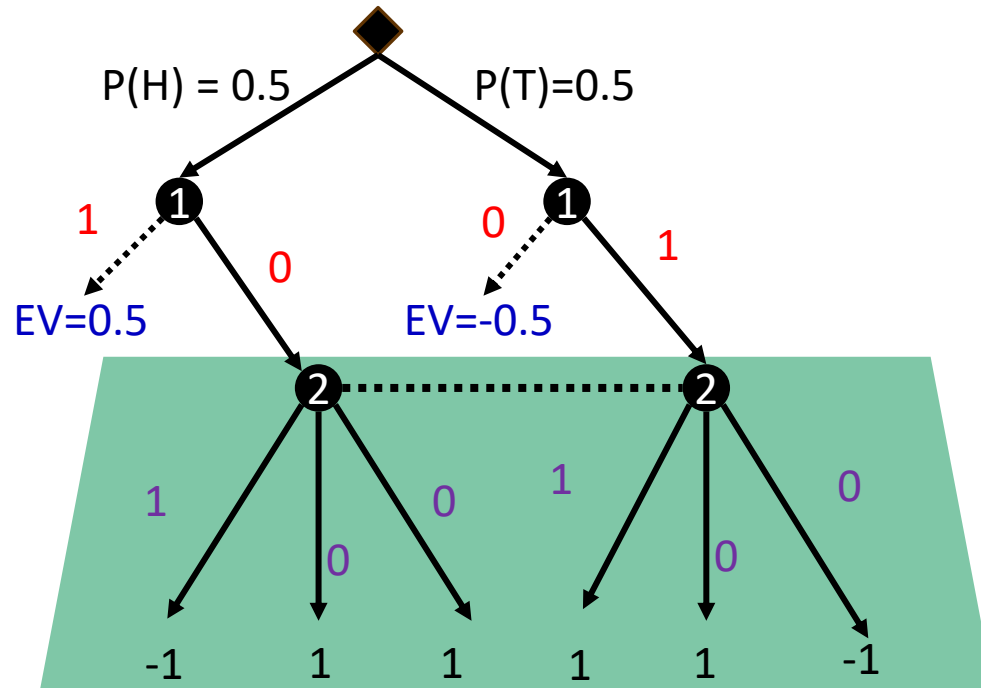


Rip out subgame + add initial chance node



NE in subgame: Play **heads** all the time

# But what does your opponent think?



What would P1 do, if it **knew** **you are going to do a refinement this way?**

- P1 best responds to the refined strategy!

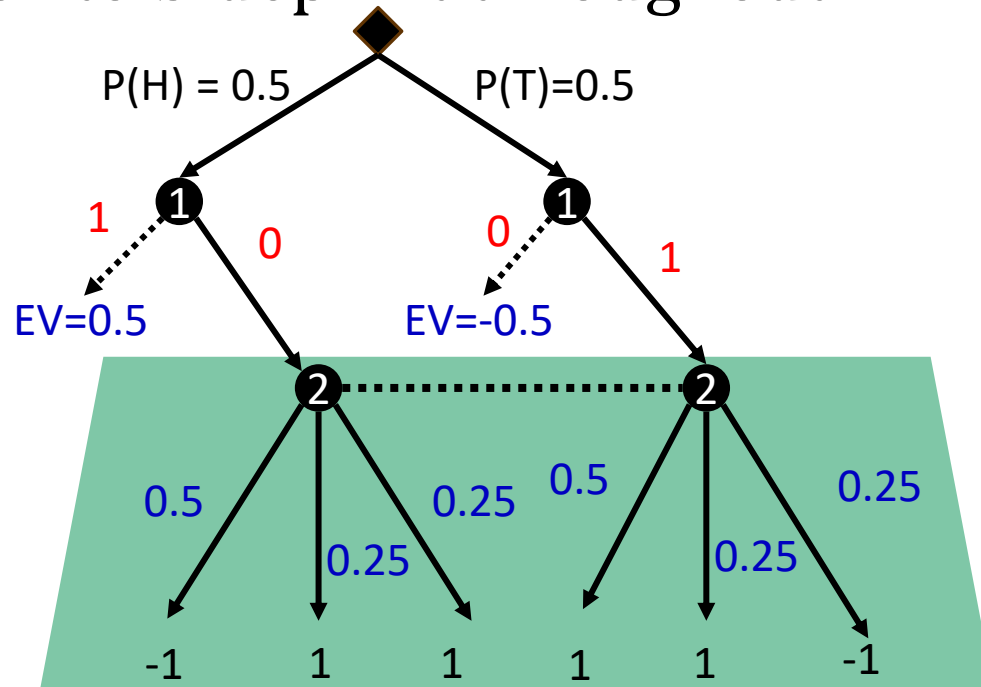
Plays only when coin toss is tails

Expected utility =  $P(H) * 0.5 + P(T) * 1 = 0.75$

- Reminder: P1 (the opponent) is maximizing

# What if you ditched subgame solving?

Stick to blueprint throughout



Note: as P2, you cannot control what P1 plays

What would P1 do, if it **knew** **you were sticking to blueprint?**

- Best responds to blueprint

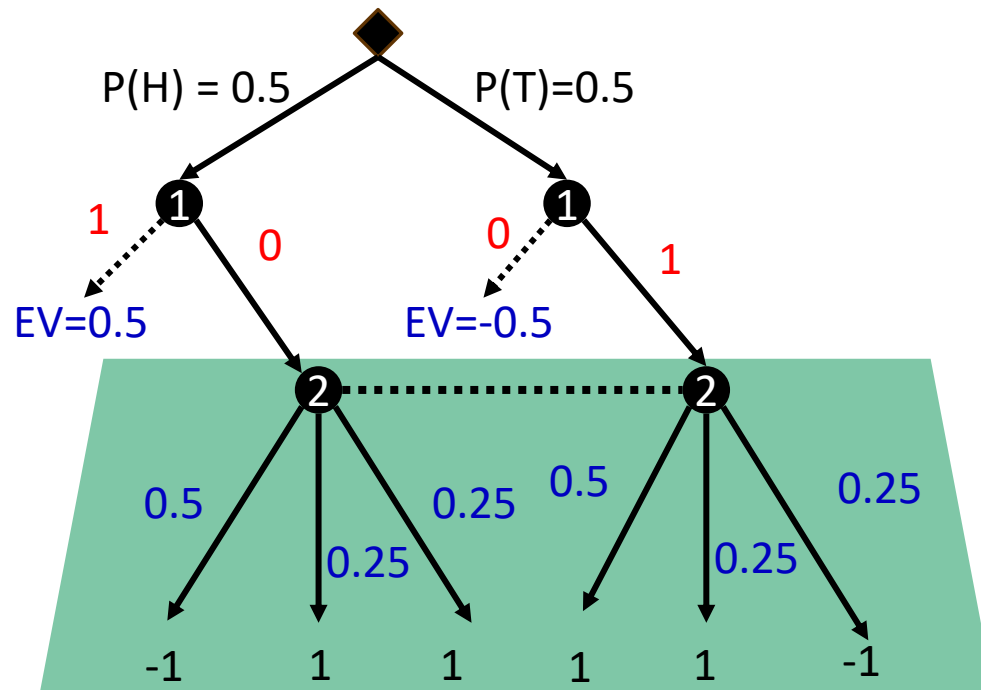
Plays only when coin toss is tails

- Same as before

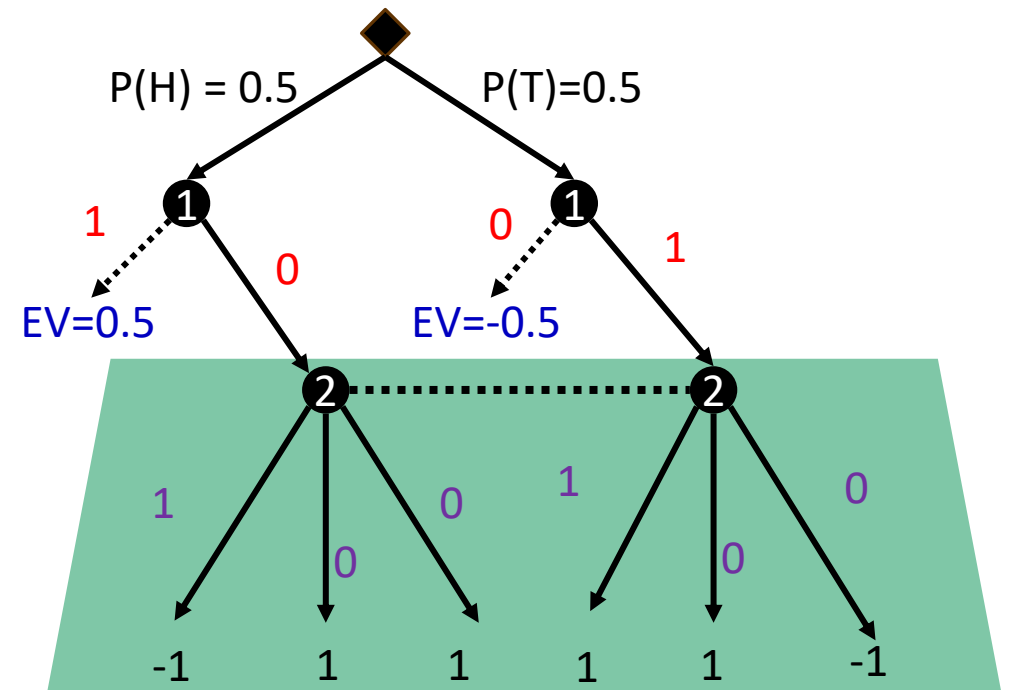
Expected utility =  $P(H) * 0.5 + P(T) * (0.5 + 0.25 - 0.25) = 0.5$

- Reminder: P1 (the opponent) is maximizing

# Naïve subgame solving is counterproductive!

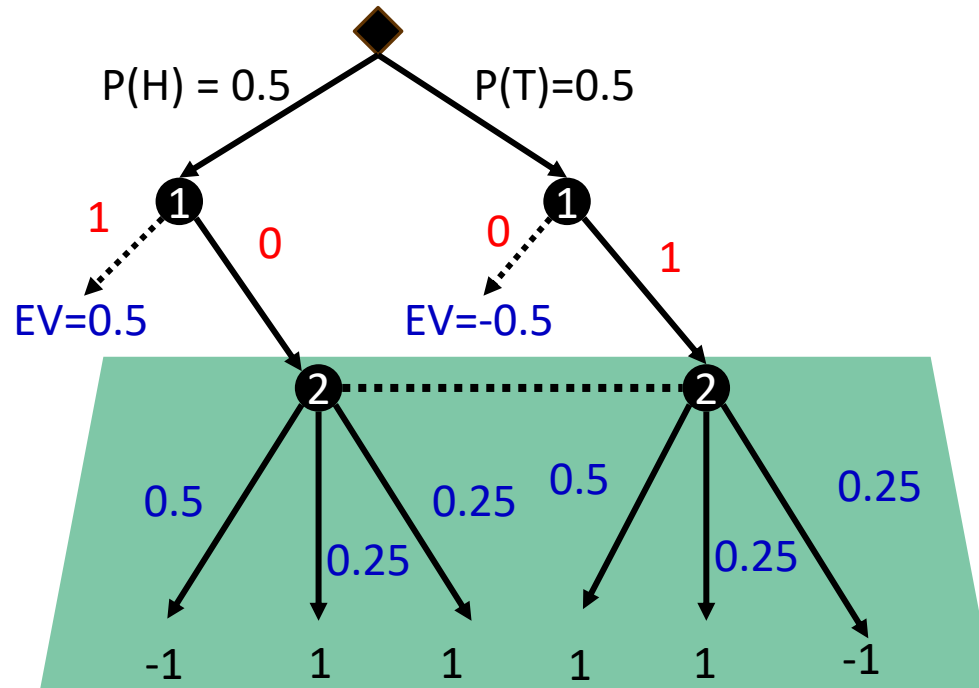


If we stuck to the blueprint



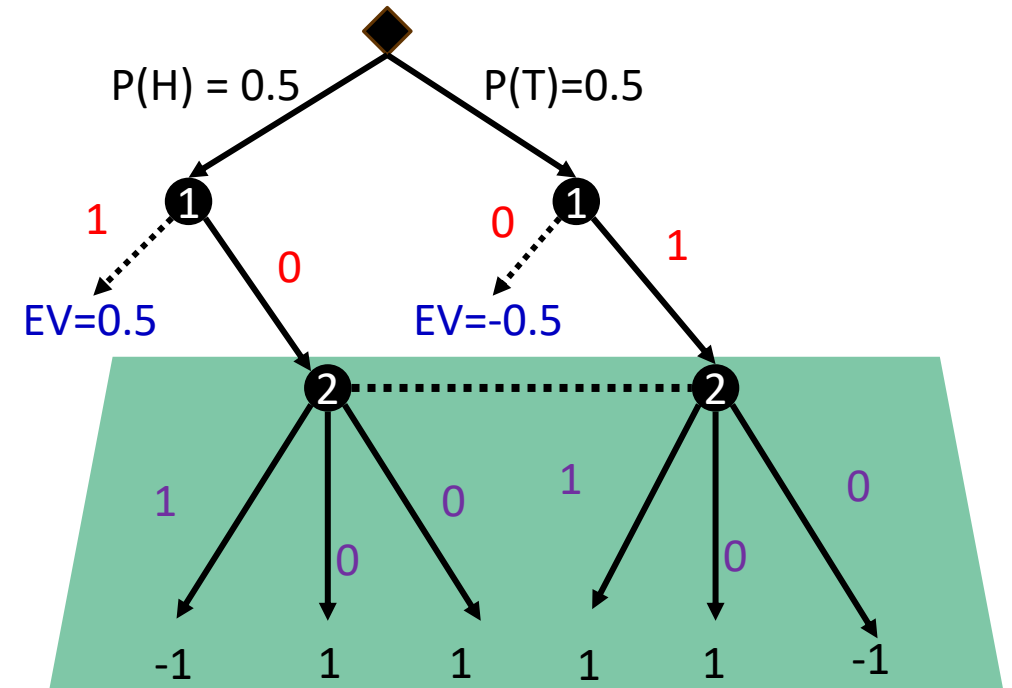
If we did refinement based on the BP

# Naïve subgame solving is counterproductive!



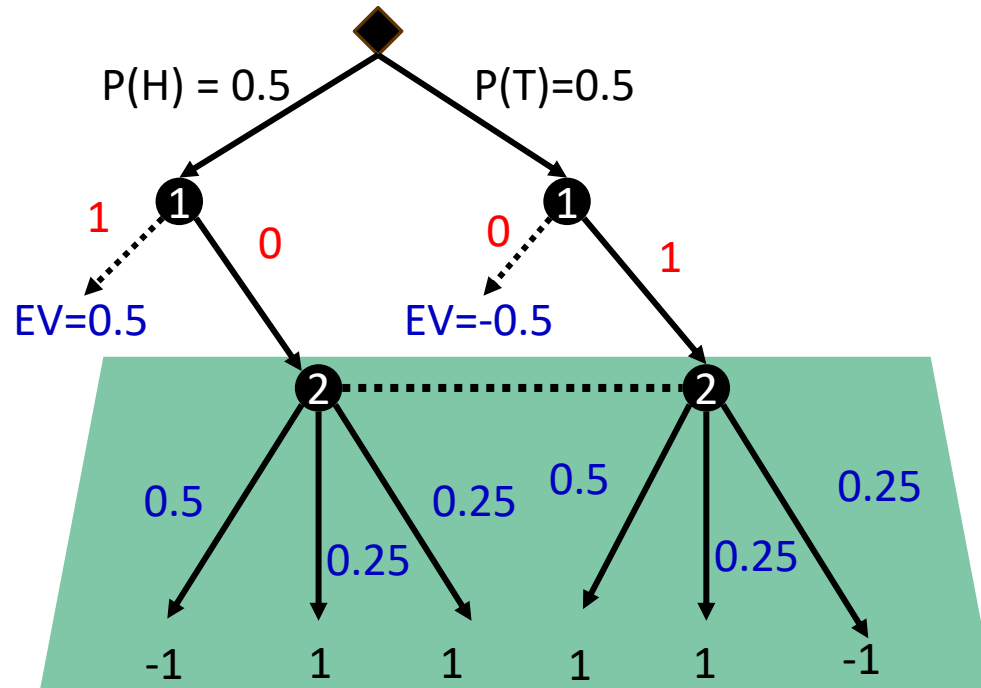
If we stuck to the blueprint

0.5



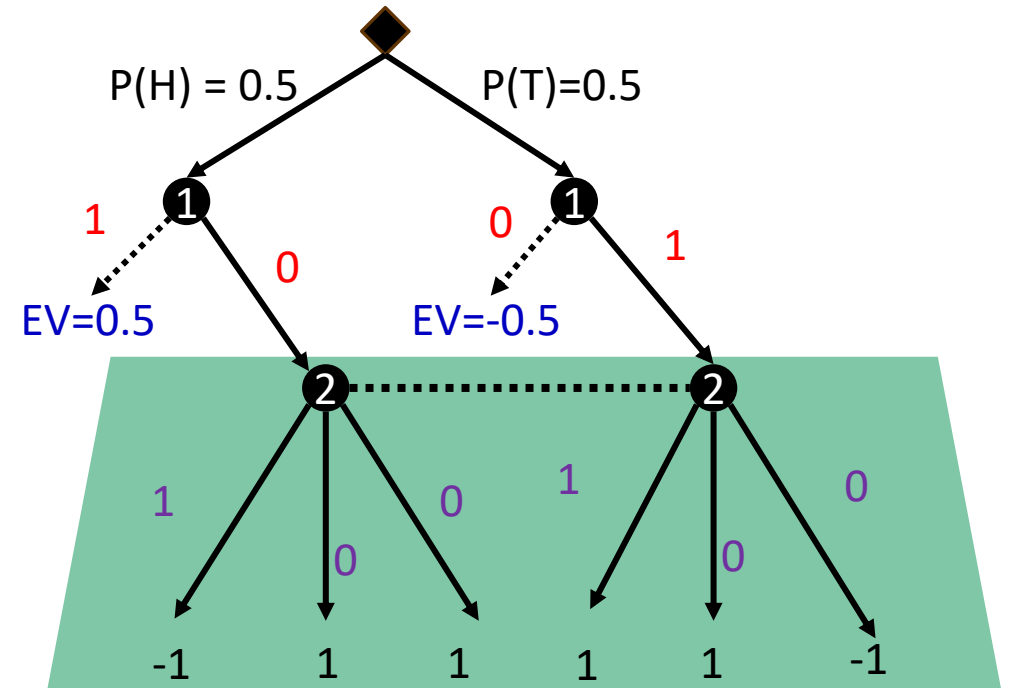
If we did refinement based on the BP

# Naïve subgame solving is counterproductive!



If we stuck to the blueprint

0.5

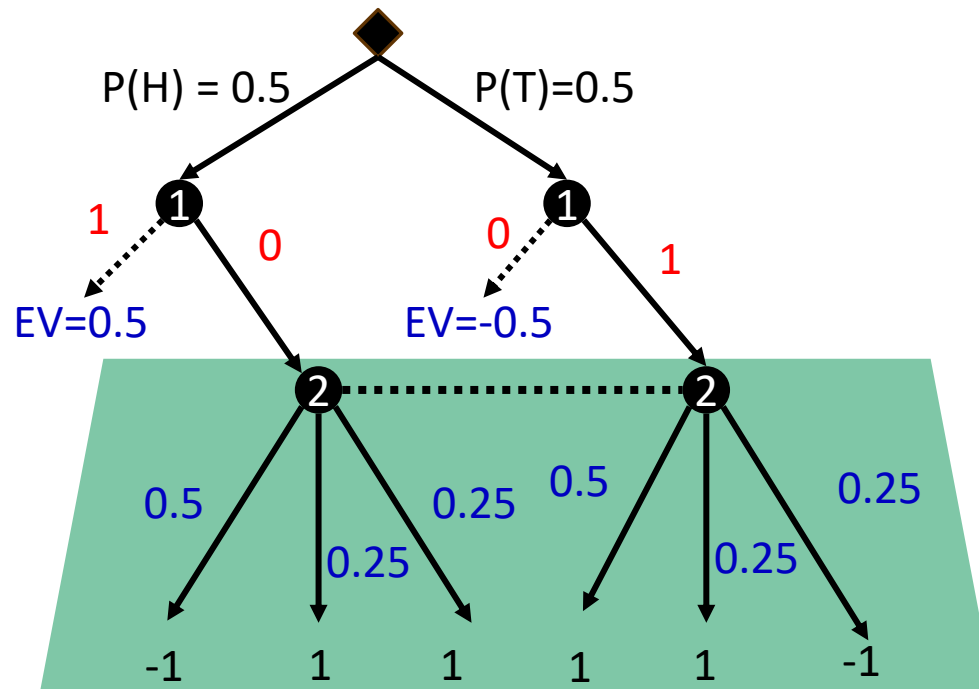


If we did refinement based on the BP

0.75



# Naïve subgame solving is counterproductive!

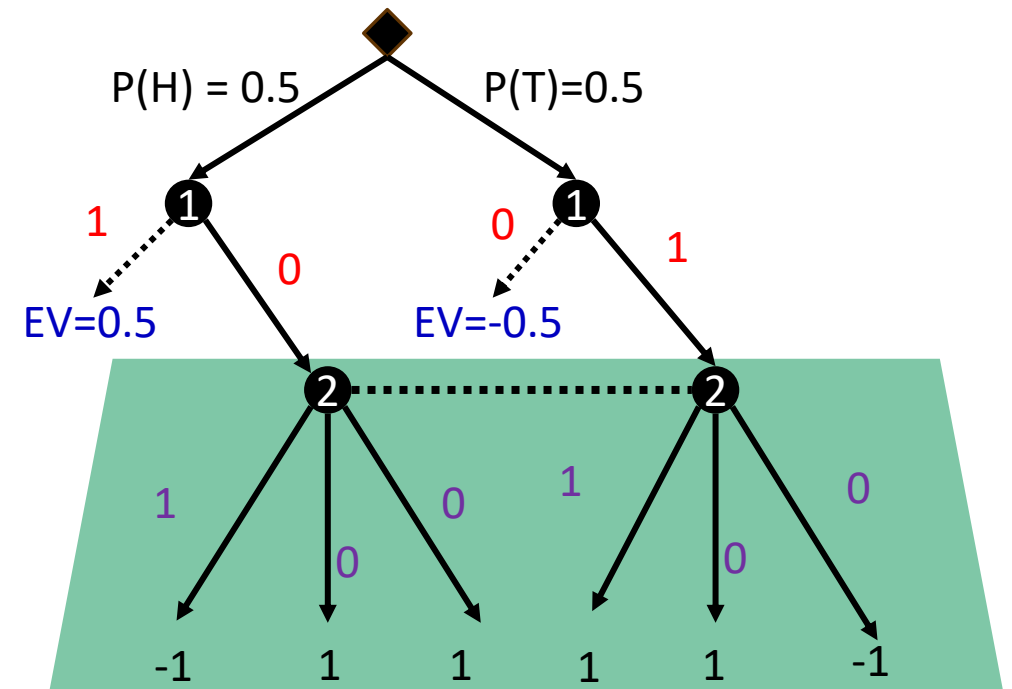


## If we stuck to the blueprint

0.5

 $\angle$ 

0.75



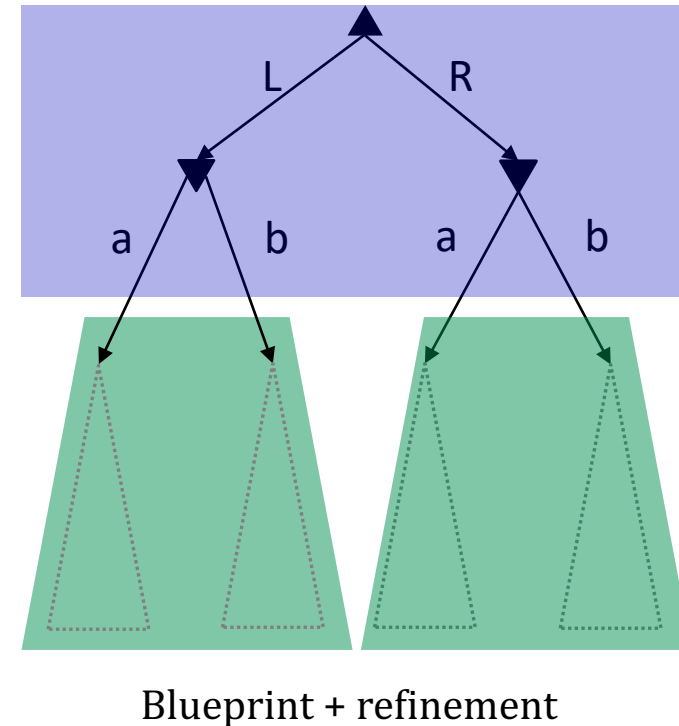
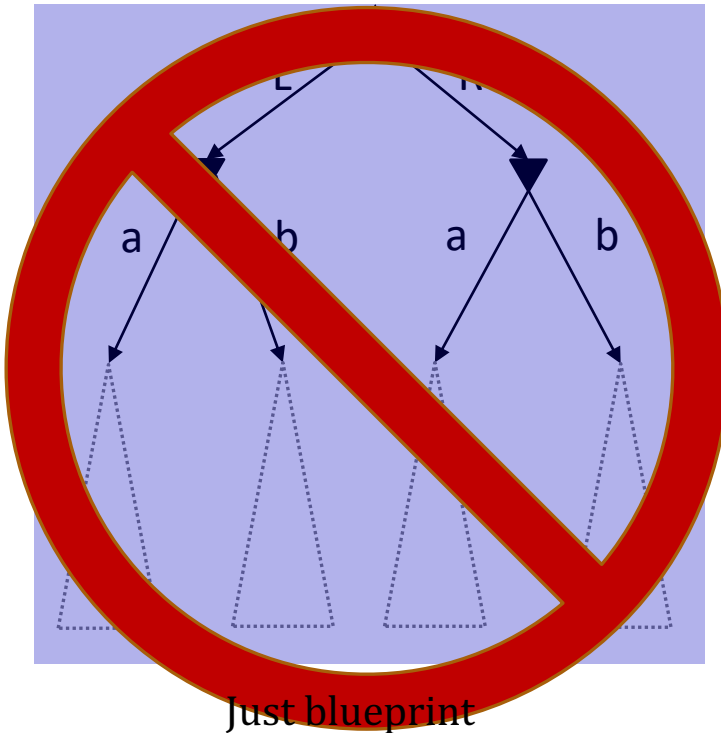
If we did refinement based on the BP

Is better than  
(from P2's perspective)

# Insight #2: Opponent BRs to refined strategy

Since we are dealing with Nash, we are thinking of the worst-case opponent. From the opponent's viewpoint, you are not playing the blueprint, but rather, **blueprint + refined strategy**

- As if refinement “source code” was uploaded online



# Safe Refinements

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Safety: Refined strategy performs **no worse than blueprint**

- Refinements are never counterproductive
- No harm in “dumping” extra compute on refinements

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Naïve subgame solving is **unsafe**

Note: in **perfect information 2p0s games** naïve subgame solving is **trivially safe**. Why?

# Safe subgame solving

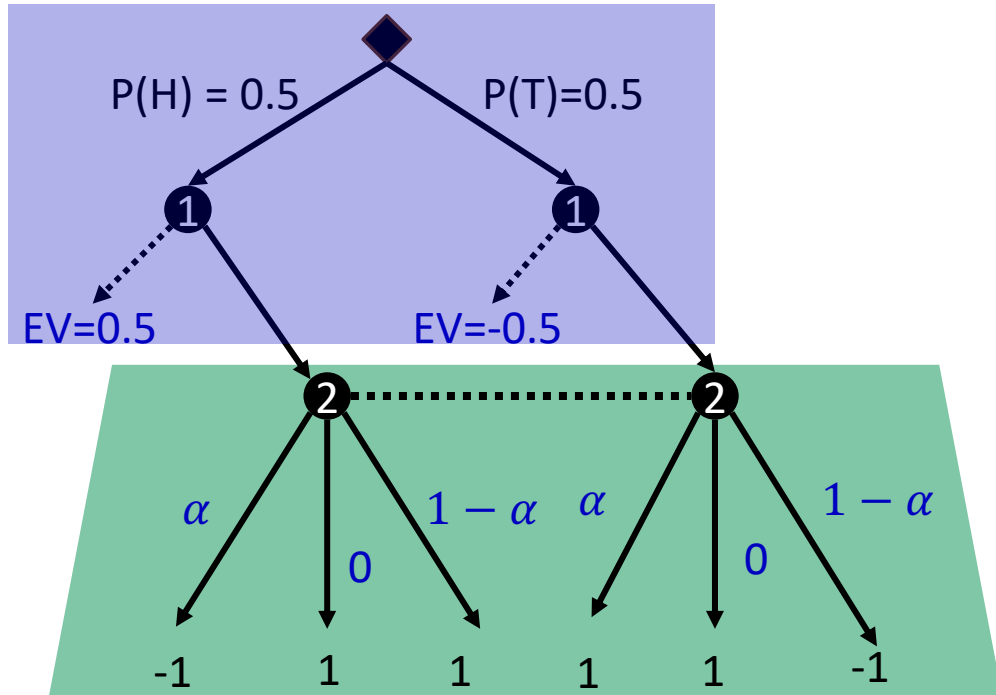
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Most material from: Safe and Nested Subgame Solving for Imperfect-Information Games (Brown and Sandholm 2017)



# Why is naïve subgame solving unsafe?

Go back to the basics: opponent is **best responding**

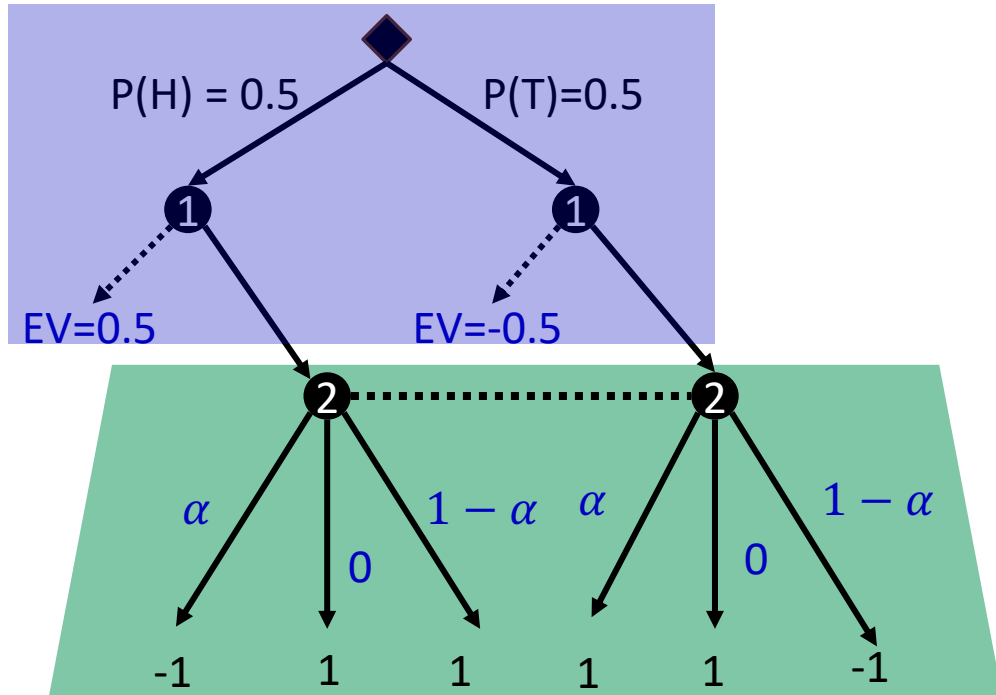


- When  $\alpha = 0$ 
  - $H \rightarrow \text{PLAY}, T \rightarrow \text{SELL}$
- When  $\alpha = 1$ 
  - $H \rightarrow \text{SELL}, T \rightarrow \text{PLAY}$
- When  $\alpha = 0.25$ ,
  - $H \rightarrow \text{INDIFFERENT}, T \rightarrow \text{INDIFFERENT}$

As  $\alpha$  increases, Heads (left branch) improves but Tails (right branch) deteriorates from P2's perspective

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As  $\alpha$  increases, Heads (left branch) improves but Tails (right branch) deteriorates from P2's perspective

**Pre-subgame** actions of opponent will change “in response” to strategy refinement!

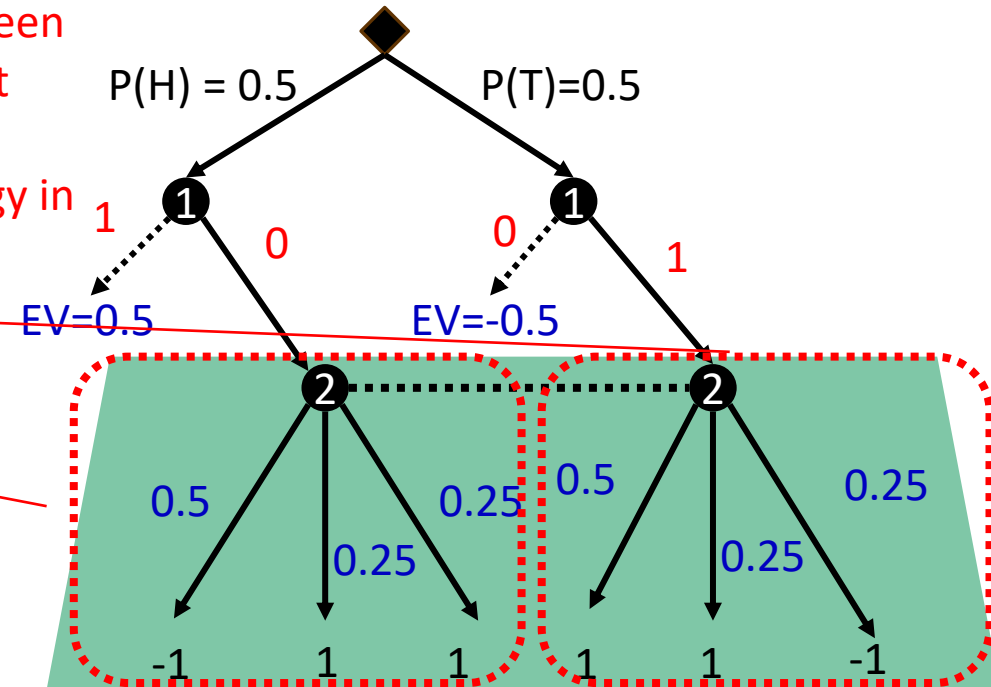
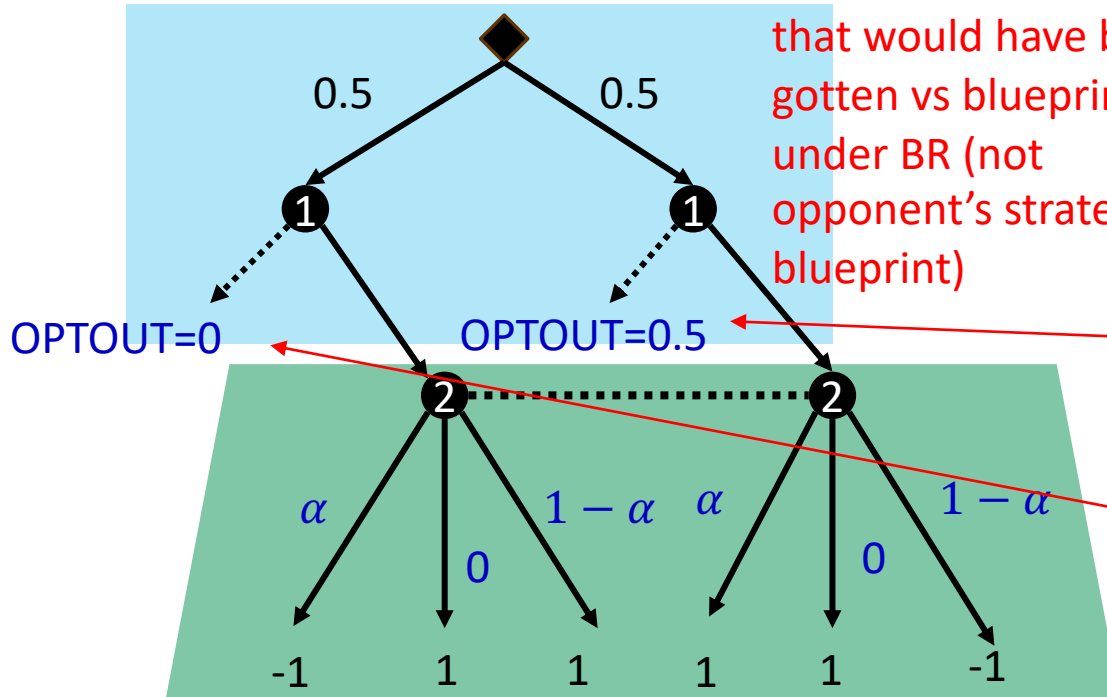
# Resolve Subgame Solving

Idea: use gadget to add in an “opt-out” option for the opponent

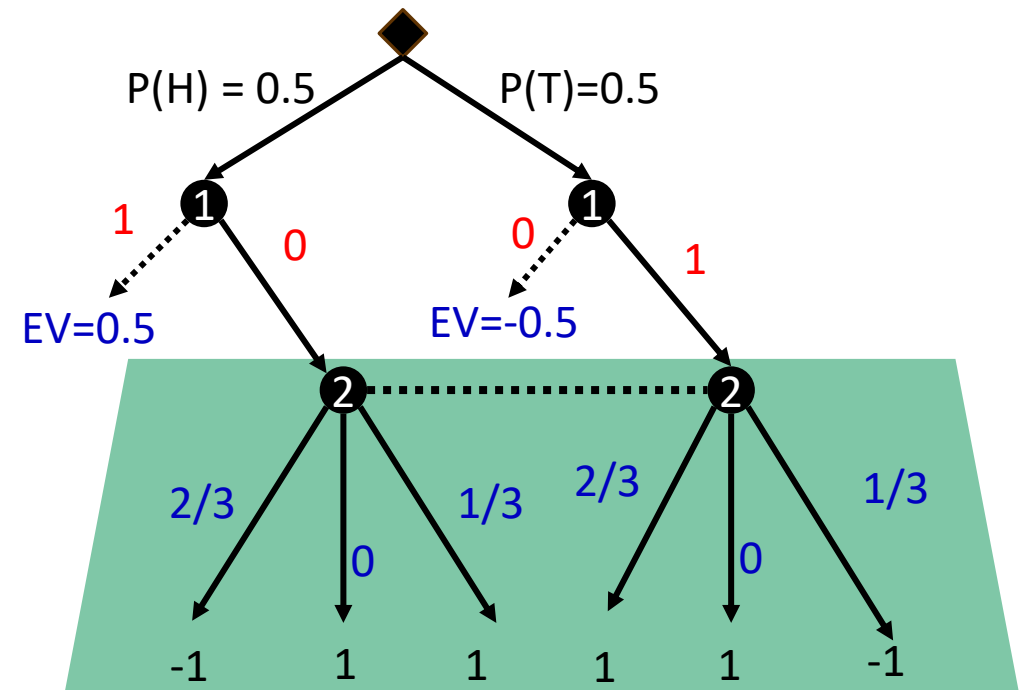
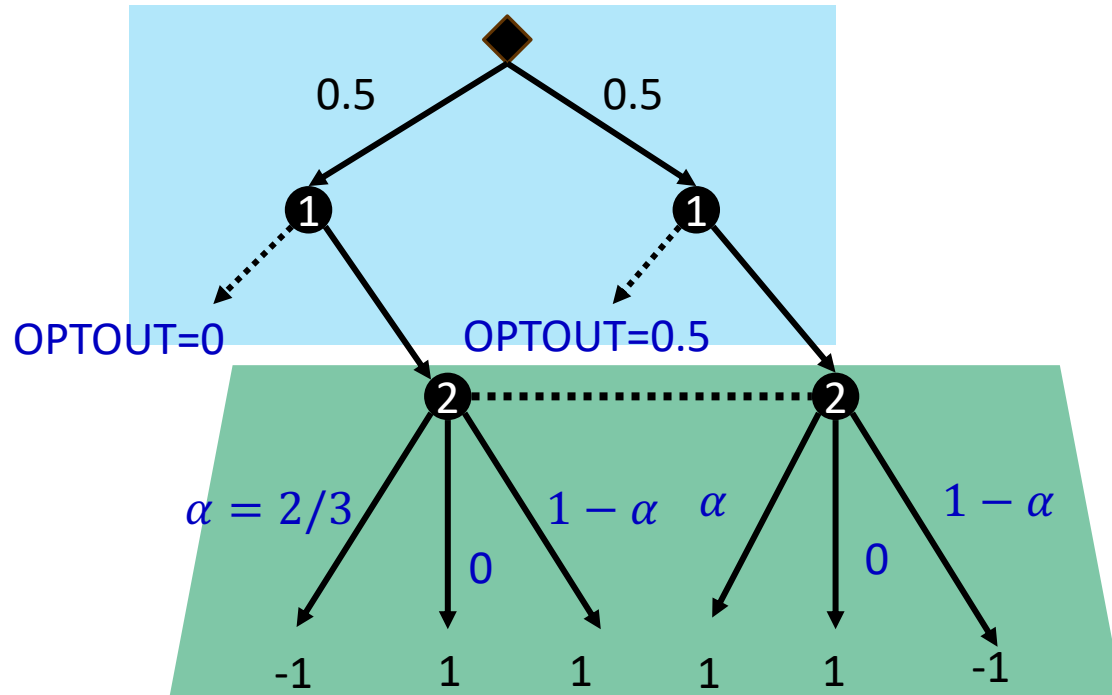
- Value of opting out is based on value under BR to blueprint
- Ensures that we cannot “hallucinate” rewards even though opponent would have changed his best response such that subgame isn’t even entered

Note: these are values that would have been gotten vs blueprint under BR (not opponent's strategy in blueprint)

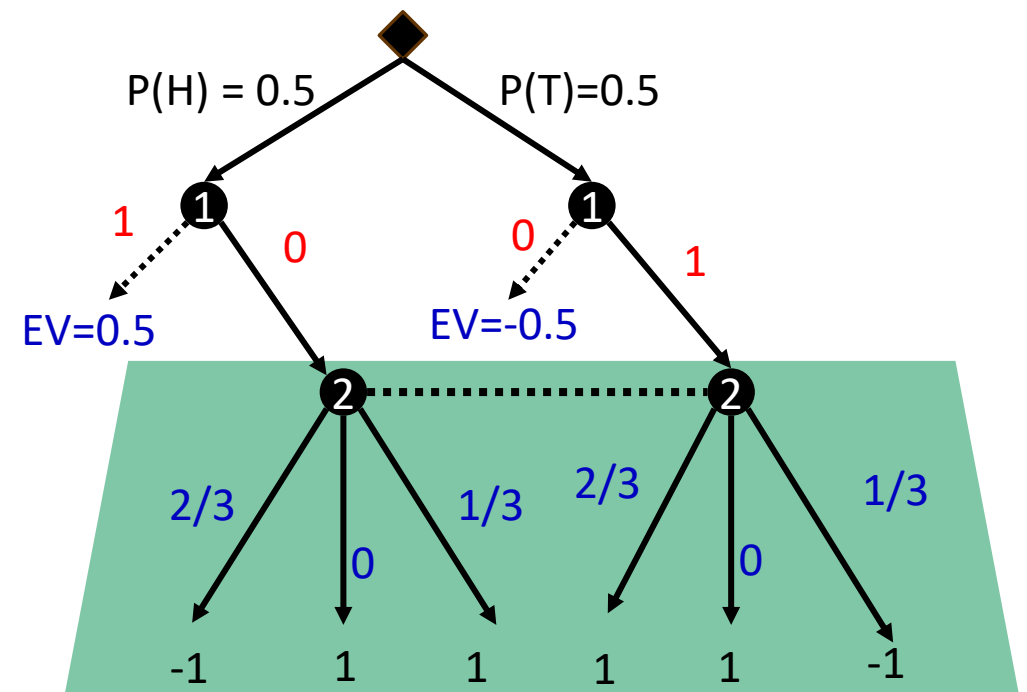
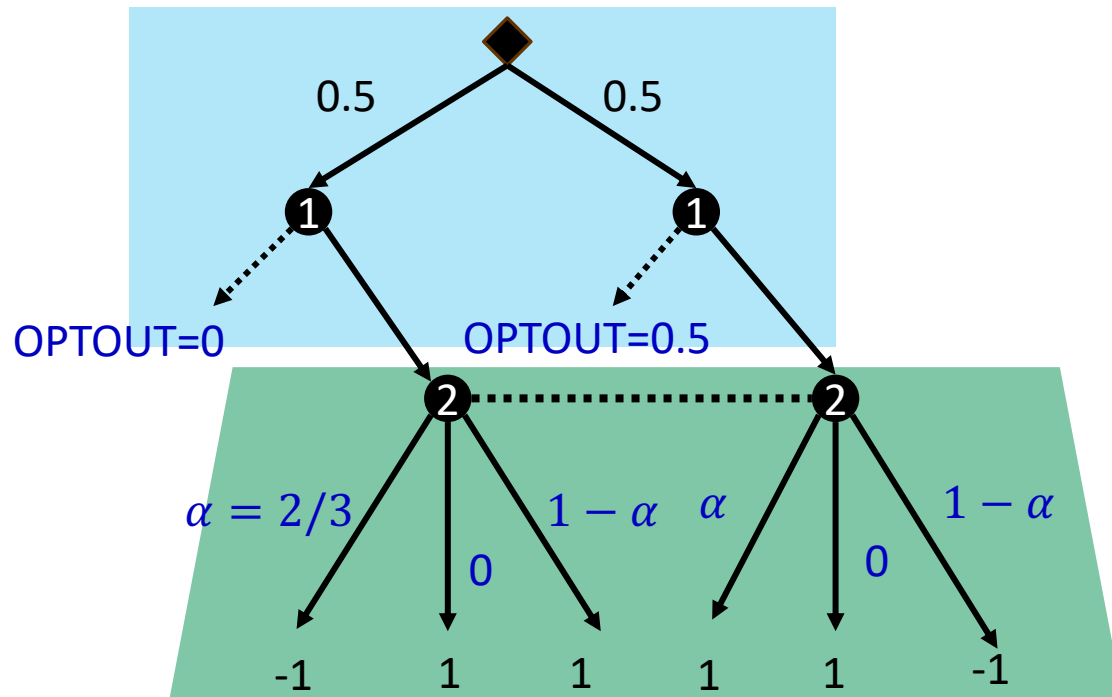
Warning: This is **not** the same game (the “pre-subgame part” in this gadget are additional vertices added and have nothing to do with the pre-subgame part in the original game)



# Examples of refined strategy I

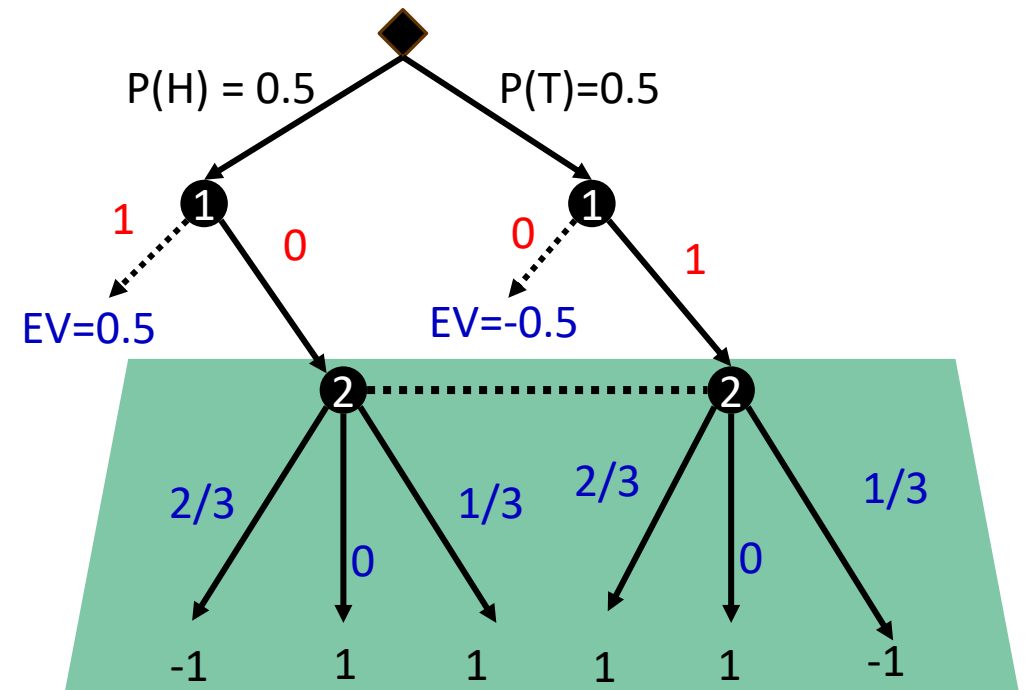
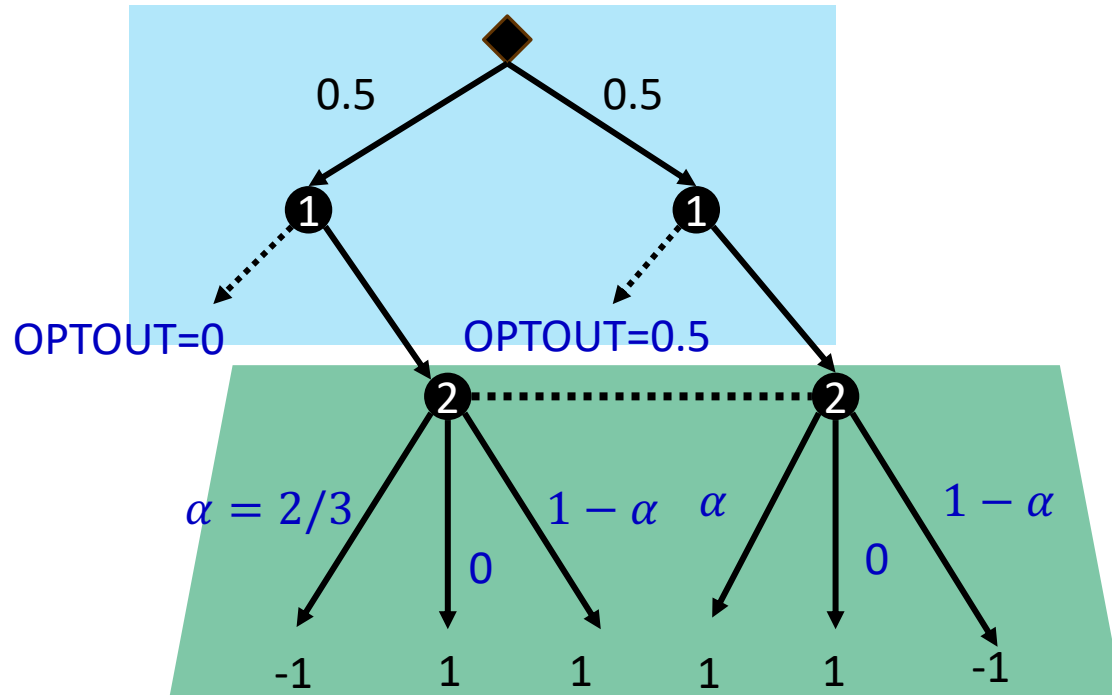


# Examples of refined strategy I



Nash of original game (taking components in subgame) is also a Nash in this game (why?)  
 What is player 1's BR here?

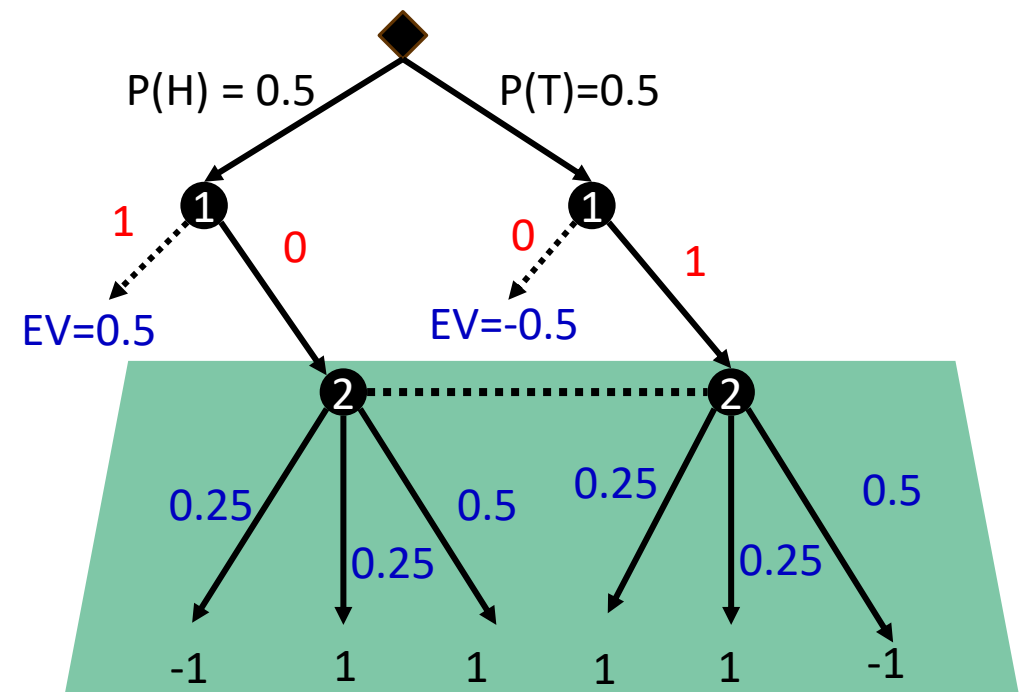
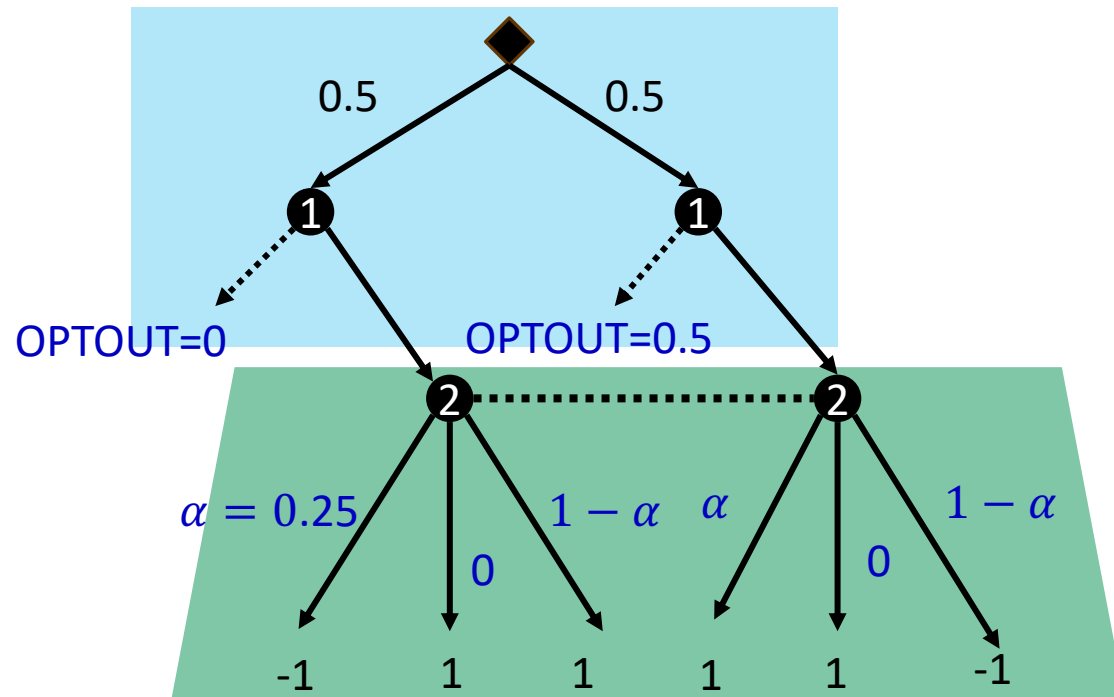
# Examples of refined strategy I



Nash of original game (taking components in subgame) is also a Nash in this game (why?)  
What is player 1's BR here?

Payoff under refined strategy is:  $0.5 * 0.5 + 0.5 * 1/3 = 0.416 < \text{Blueprint}$   
**Strict improvement from blueprint!**

# Examples of refined strategy II



The blueprint itself is also a Nash of the gadget game...  $\rightarrow$  No improvement!

# Resolve Subgame Solving is Safe

Why?

- Guarantees that probability distribution for refined strategy at “head” of subgame is the same as blueprint

But usually won't lead to meaningful improvements

- Too stringent!
- Gadget game doesn't allow for much improvement, not explicitly at least
- We solved the problem of “sacrificing one branch for the other”, but in turn made it very hard to improve



# Resolve Subgame Solving is Safe

Why?

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But usually won't lead to meaningful improvements

- Too stringent!
- Gadget game doesn't allow for much improvement, not explicitly at least
- We solved the problem of “sacrificing one branch for the other”, but in turn made it very hard to improve

How about allowing cases where **every** “head subgame” vertex's EV is improved?

# Maxmargin Subgame Solving

Caveat! We can relax this

Idea (from before): if values of all starting subtrees in subgame (under BR of opponent) do not increase, then we are **safe**

- The problem of non-safety came about because we didn't know which branch we will end up in, **since opponent BRs to us even pre-subgame**
- But if all branches are not worse than blueprint (under opponent BR), then we can't have performed worse

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What if we try to maximize the *margin*, i.e., minimum improvement from blueprint across all branches?

- Margin will never be negative because blueprint itself is feasible refinement  
→ safe!
- If margin is strictly positive, whichever branch we end up in, we will always do strictly better than before!

# Maxmargin Subgame Solving

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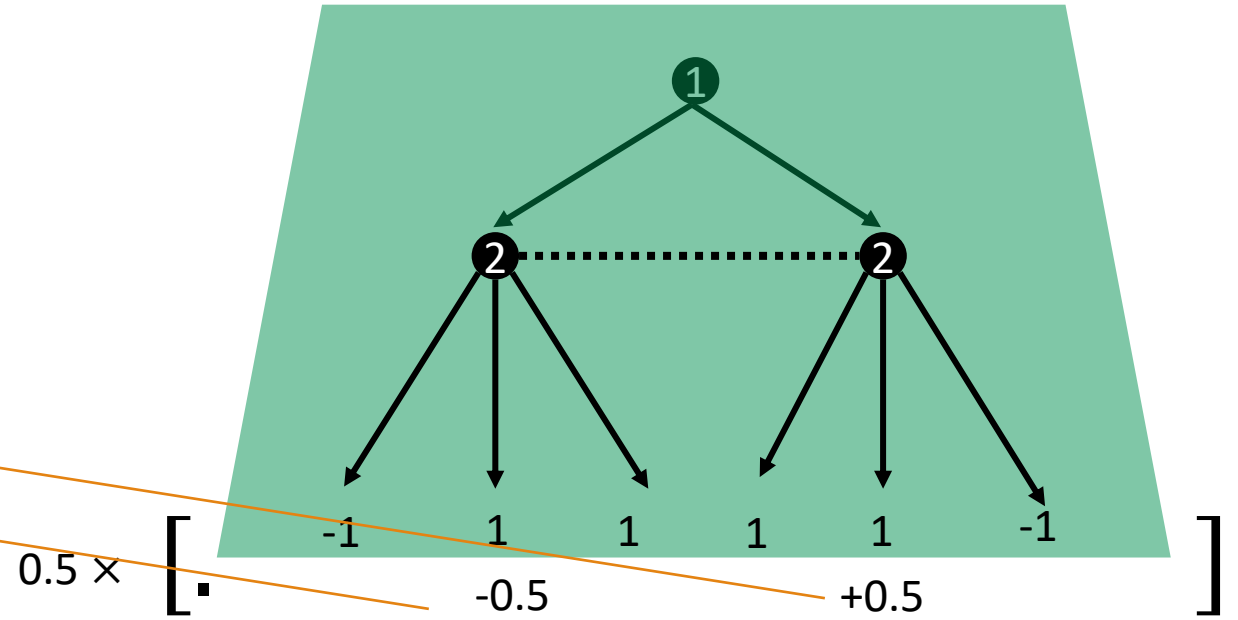
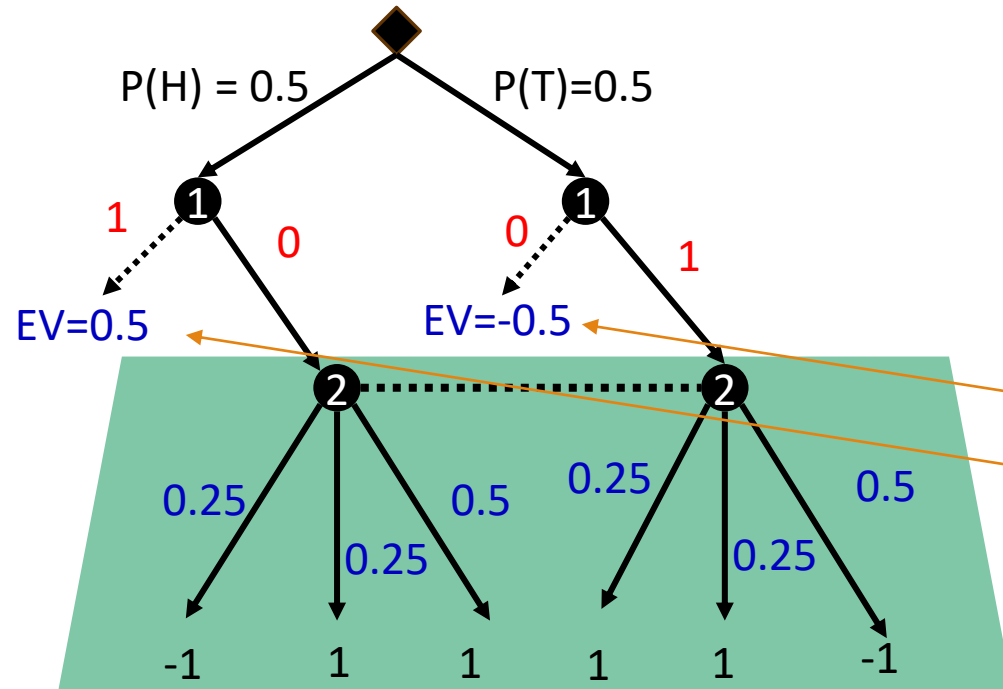
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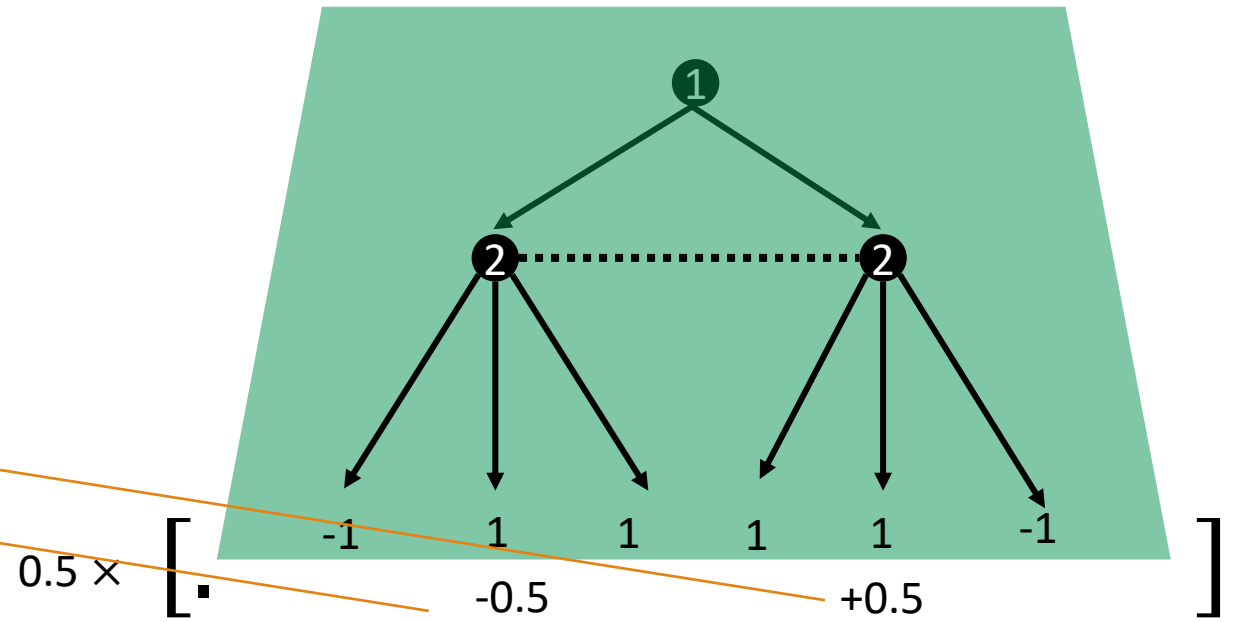
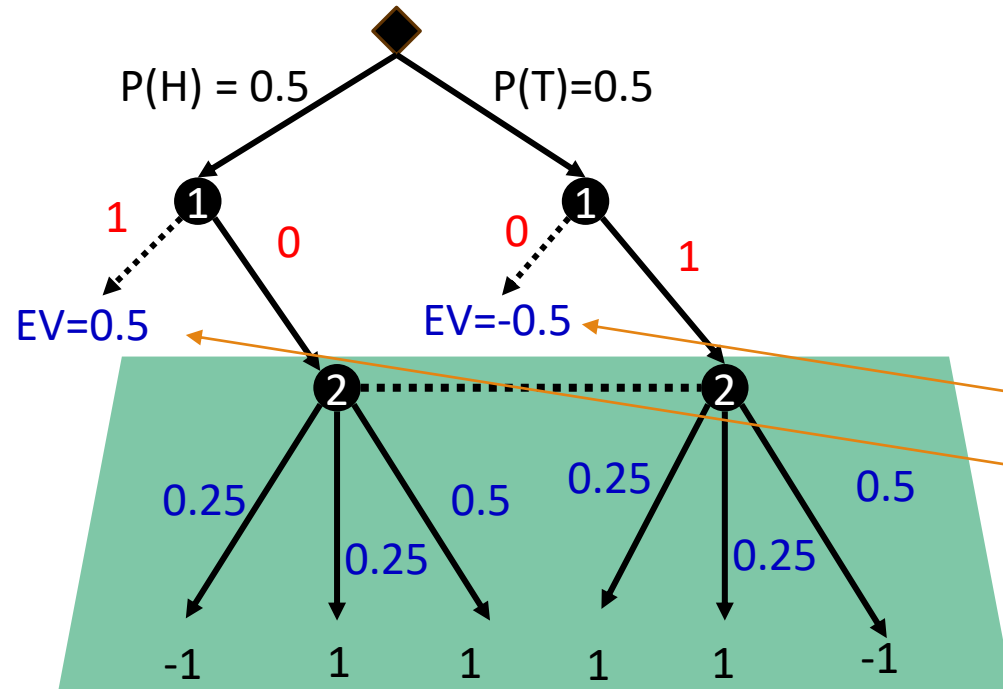
Refinement can be done via linear programming

# Maxmargin Gadget I



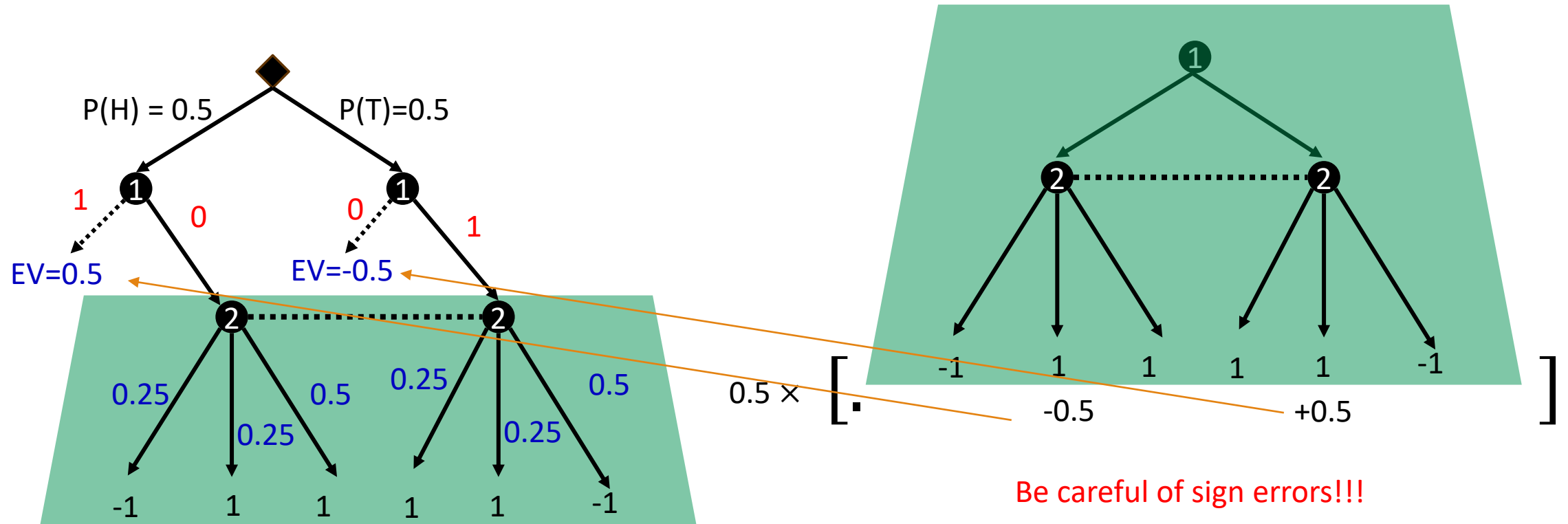
Be careful of sign errors!!!

# Maxmargin Gadget I



Why does this work?

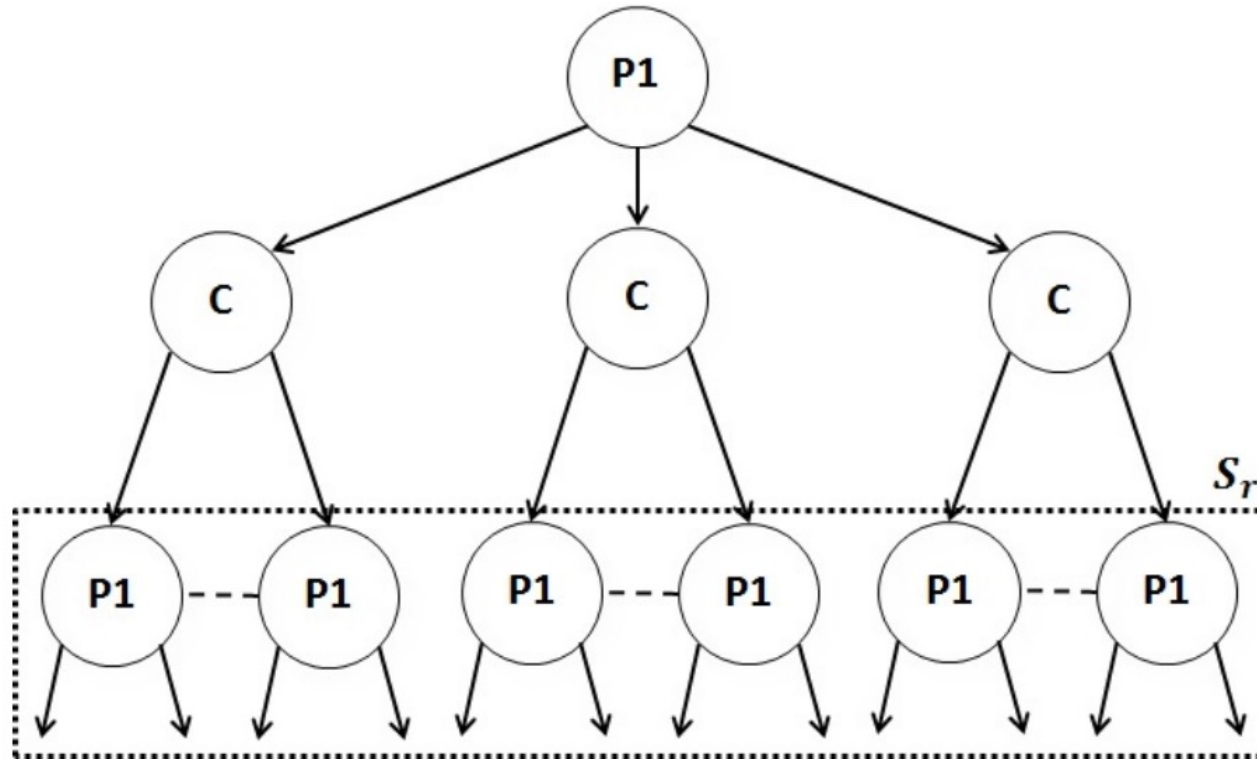
# Maxmargin Gadget I



Why does this work?

Think of maxmargin as a min-max problem, opponent now chooses the branch with the lowest margin (improvement).

# Maxmargin Gadget II



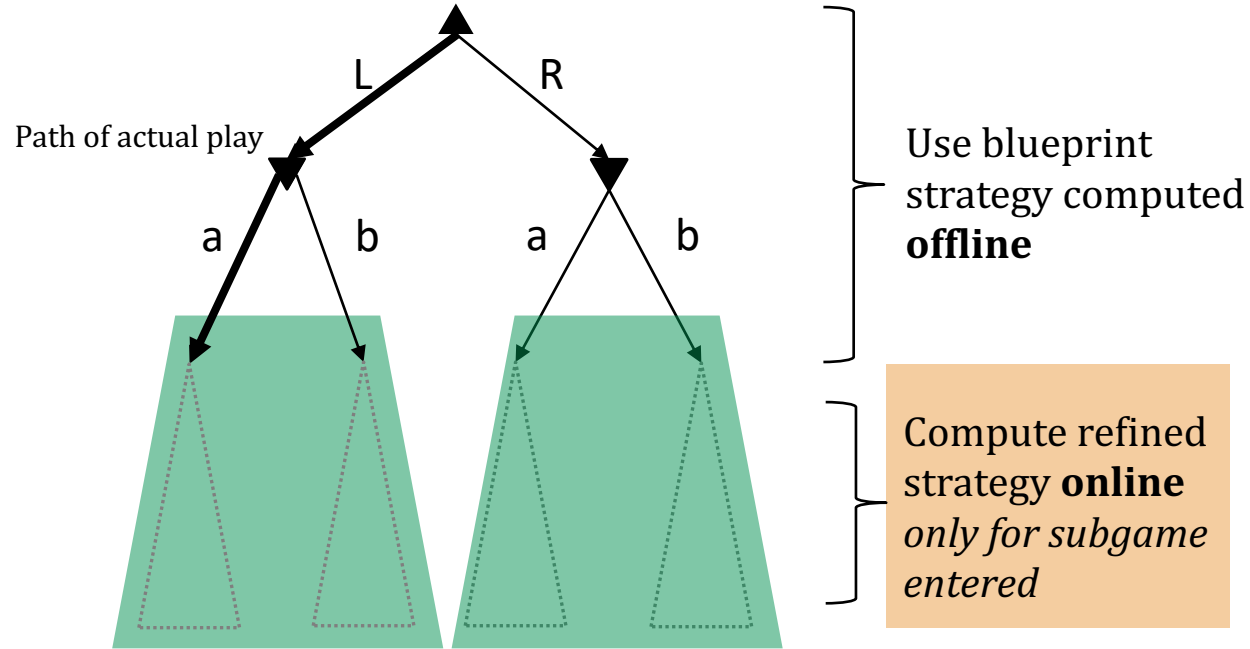
Implementation can be  
very tricky

Figure 6: An example of a gadget game in Maxmargin refinement.  $P_1$  picks the initial info set she wishes to enter  $S_r$  in. Chance then picks the particular node of the info set, and play then proceeds identically to the augmented subgame, except all  $P_1$  payoffs are shifted by the size of the alternative payoff and the alternative payoff is then removed from the augmented subgame.



# Reach subgame solving/maxmargin

Intuition: So far, only considered a single subgame at any point.

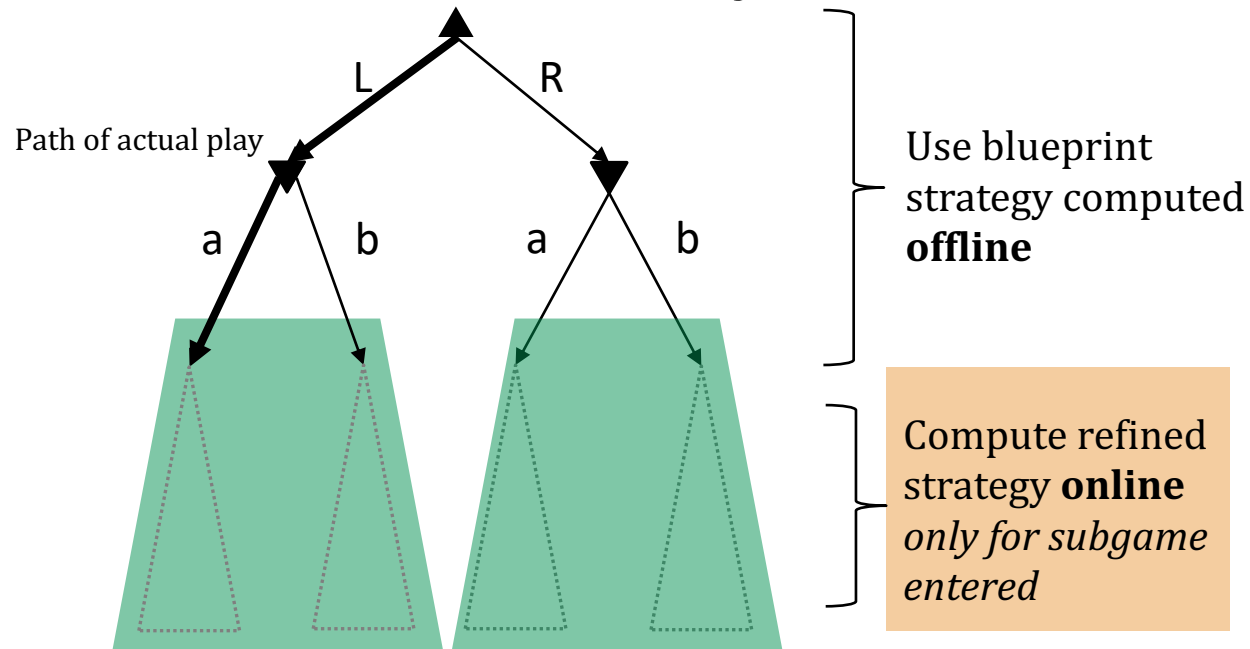


Not covered due to lack of time

In reality, many subgames will exist!

# Reach subgame solving/maxmargin

Intuition: So far, only considered a single subgame at any point.



Not covered due to lack of time

In reality, many subgames will exist!

From opponent's perspective, when finding BR, refinement is done on **all** subgames, not just one, even though from the refiner's POV, only one subgame's refinement is computed

# Subgame solving for Stackelberg eqm

---

[Optional]

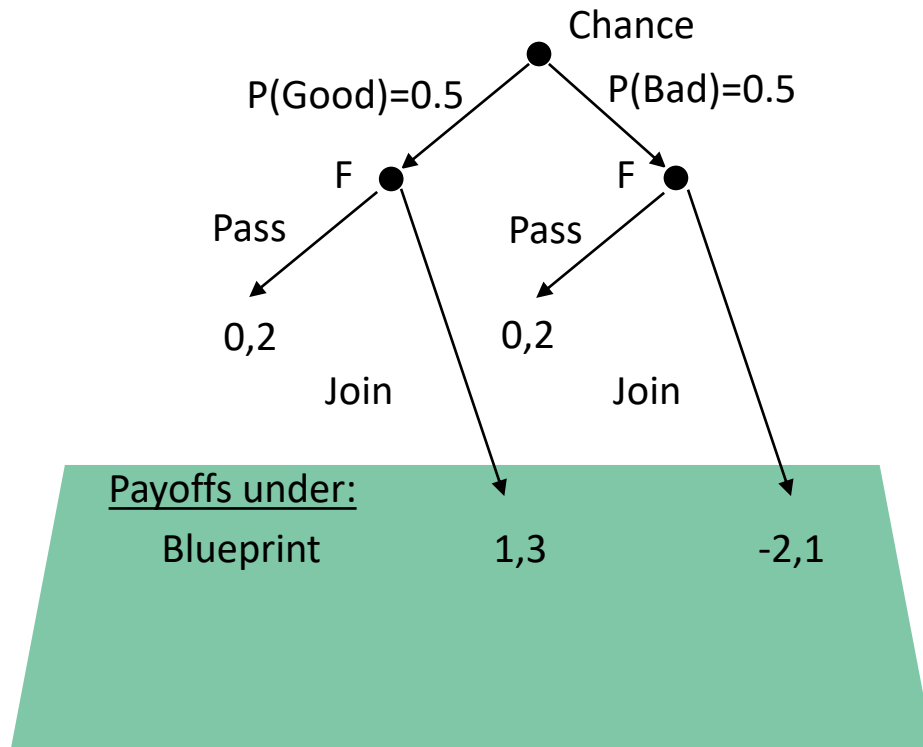
# Carrots and Sticks in Resolving

## Extended Hiring Game

Follower: Applicants applying for a job

Leader: Hiring committee

Game: How should leader conduct the hiring process?



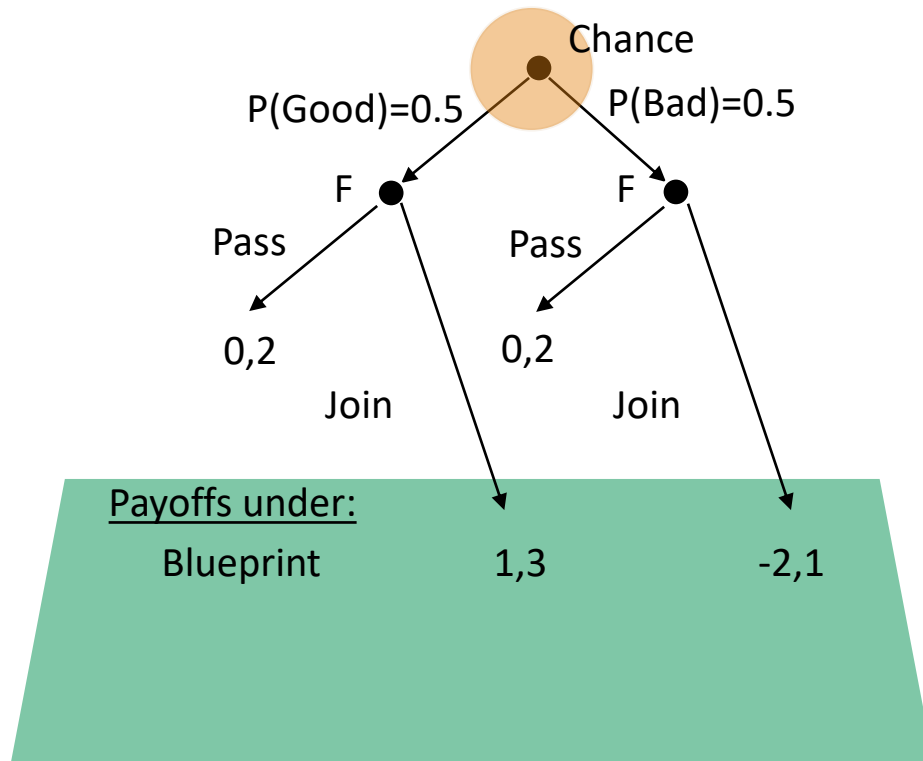
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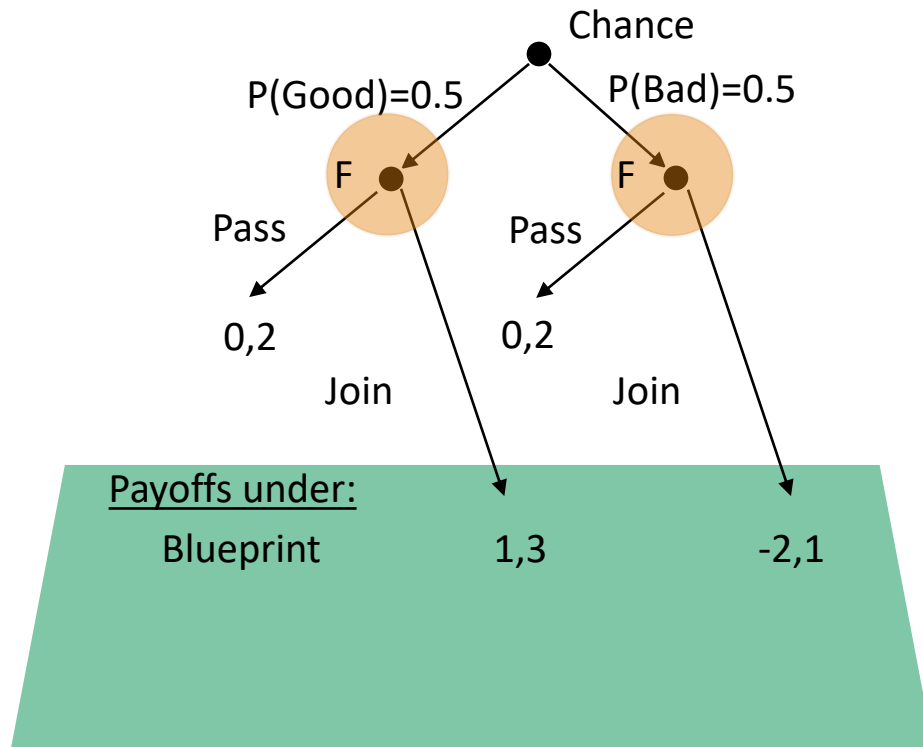
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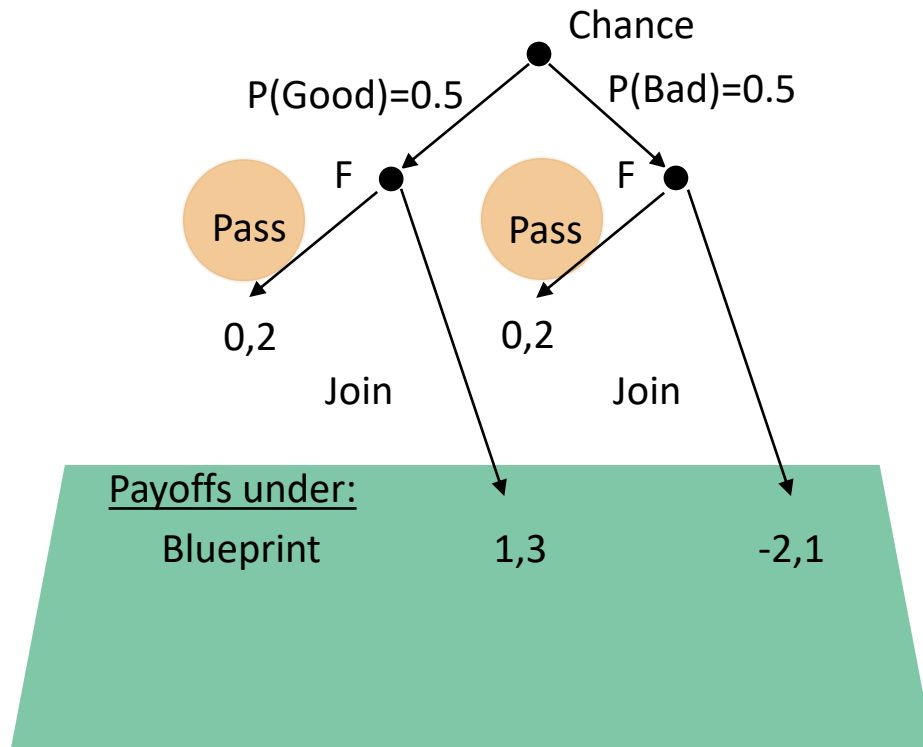
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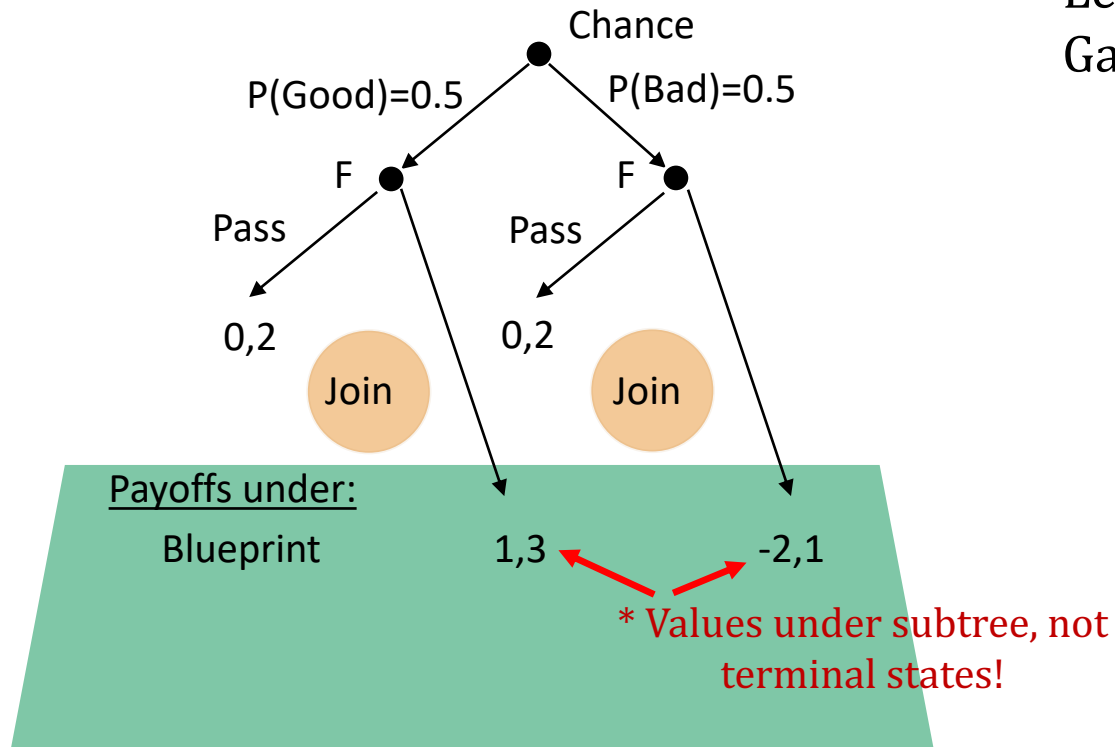
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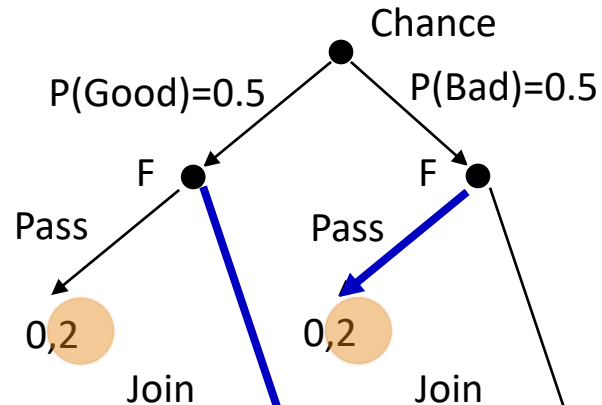
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## Extended Hiring Game

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Game: How should leader conduct the hiring process?



Payoffs under:

Blueprint

## Blueprint

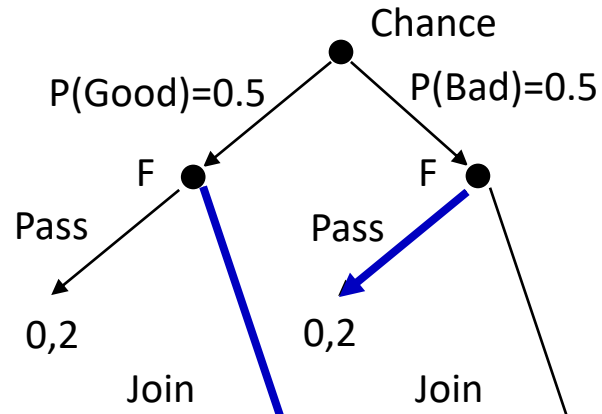
Candidates follow the **blue paths**.

Good candidates apply, bad candidates pass. 🤔

Leader expected payoff = **0.5**

# Carrots and Sticks in Resolving

## Extended Hiring Game



Payoffs under:  
Blueprint  
Unsafe Refinement

1,3  
2,1

-2,1  
-1,3

\*Better for leader regardless

Follower: Applicants applying for a job

Leader: Hiring committee

Game: How should leader conduct the hiring process?

### Blueprint

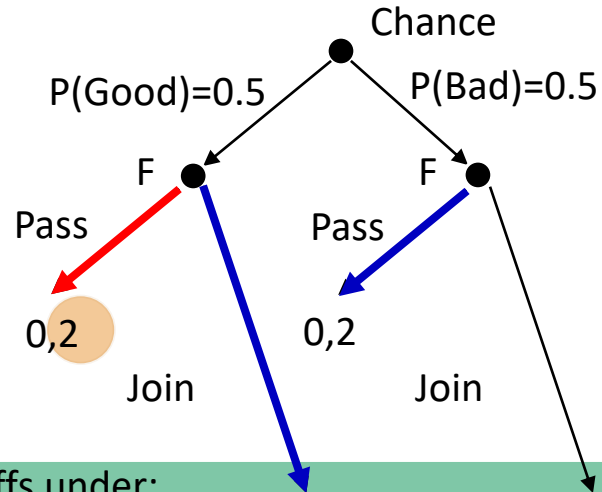
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# Carrots and Sticks in Resolving

## Extended Hiring Game



Payoffs under:

|                   |     |      |
|-------------------|-----|------|
| Blueprint         | 1,3 | -2,1 |
| Unsafe Refinement | 2,1 | -1,3 |

Follower: Applicants applying for a job

Leader: Hiring committee

Game: How should leader conduct the hiring process?

### Blueprint

Candidates follow the **blue paths**.

Good candidates apply, bad candidates pass. 🤔

Leader expected payoff = **0.5**

### Naïve Refinement

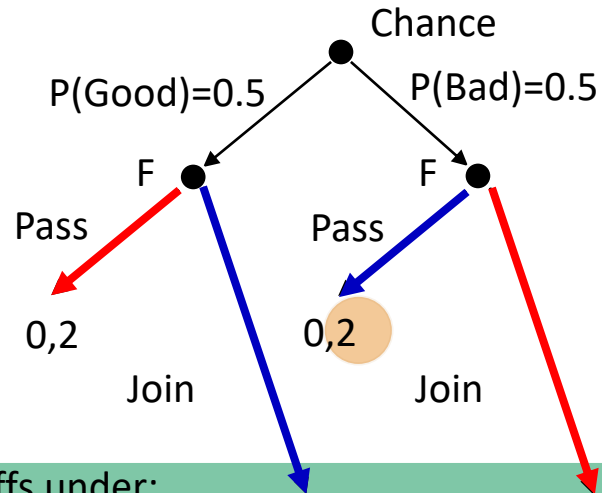
Candidates follow the **red paths**.

Good candidates pass, bad candidates apply. 😡

Leader expected payoff = **-0.5**

# Carrots and Sticks in Resolving

## Extended Hiring Game



Payoffs under:

|                   |     |      |
|-------------------|-----|------|
| Blueprint         | 1,3 | -2,1 |
| Unsafe Refinement | 2,1 | -1,3 |

Follower: Applicants applying for a job

Leader: Hiring committee

Game: How should leader conduct the hiring process?

### Blueprint

Candidates follow the **blue paths**.

Good candidates apply, bad candidates pass. 🤔

Leader expected payoff = **0.5**

### Naïve Refinement

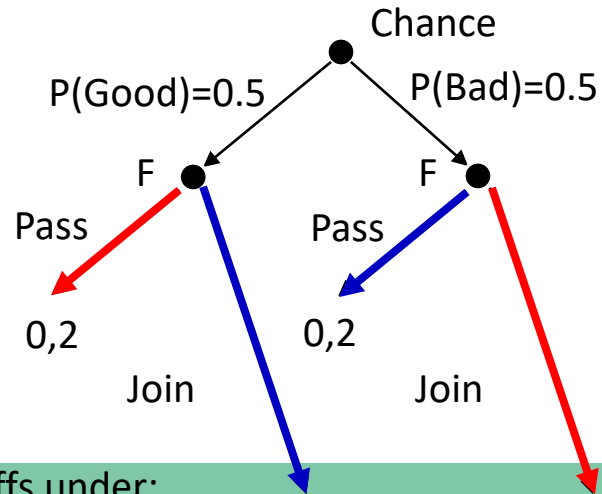
Candidates follow the **red paths**.

Good candidates pass, bad candidates apply. 😡

Leader expected payoff = **-0.5**

# Carrots and Sticks in Resolving

## Extended Hiring Game



Payoffs under:

|                   |     |      |
|-------------------|-----|------|
| Blueprint         | 1,3 | -2,1 |
| Unsafe Refinement | 2,1 | -1,3 |

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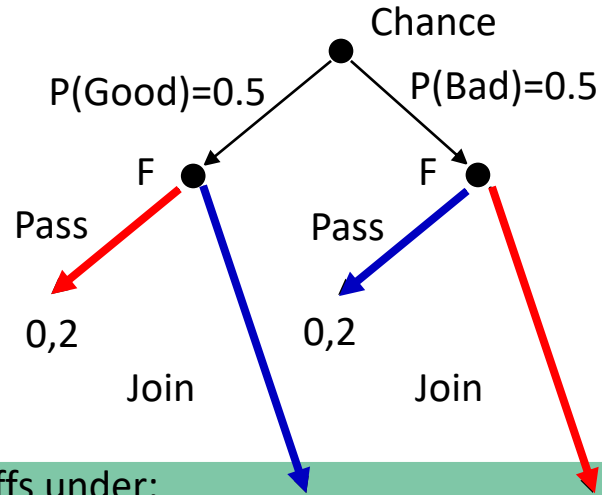
Leader expected payoff = **-0.5**

>

UNSAFE

# Carrots and Sticks in Resolving

## Extended Hiring Game



Payoffs under:

|                   |          |          |
|-------------------|----------|----------|
| Blueprint         | 1,3      | -2,1     |
| Unsafe Refinement | 2,1      | -1,3     |
| Safe Bounds       | $\geq 2$ | $\leq 2$ |



Follower: Applicants applying for a job

Leader: Hiring committee

Game: How should leader conduct the hiring process?

### Blueprint

Candidates follow the **blue paths**.  
 Good candidates apply, bad candidates pass. 🥴  
 Leader expected payoff = **0.5**

### Naïve Refinement

Candidates follow the **red paths**.  
 Good candidates pass, bad candidates apply. 😡  
 Leader expected payoff = **-0.5**

>

UNSAFE

### Safe Bounds

Candidates follow the **blue paths** 🥴  
 Blueprint obeys bounds  $\Rightarrow$  expected payoff: at least that of blueprint!

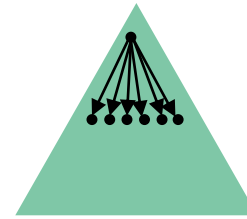
# Our algorithm

## Step 1



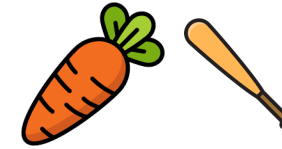
Compute safe value bounds for follower information sets

## Step 2 (only for subgame)



Solve

+



Apply naïve refinement with the enforcement of bounds

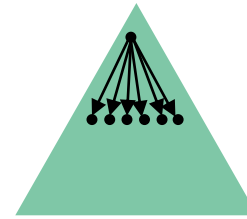
# Our algorithm

## Step 1



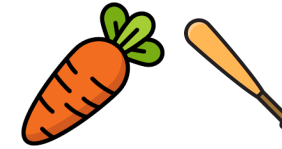
Compute safe value bounds for follower information sets

## Step 2 (only for subgame)



Solve

+



Apply naïve refinement with the enforcement of bounds

Not a standard Stackelberg problem



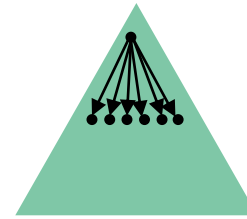
# Our algorithm

## Step 1



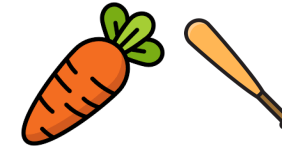
Compute safe value bounds for follower information sets

## Step 2 (only for subgame)



Solve

+



Apply naïve refinement with the enforcement of bounds

Not a standard Stackelberg problem

Step 2 can be solved by

- Mathematical programming (integer linear program)
- Transform into another Stackelberg problem using **gadget game**