

L3D

## Revision lecture :

1

= Postscript to lectures 1-29:

If we work in ZFC then Gödel's Completeness theorem + Compactness theorem.

hold for arbitrary  $\mathcal{L}$

(Also versions of Löwenheim-Skolem, ...)

Different

Exam: Past papers '18, '19, '20, '23, '24 on Maths Central (B13)  
- used to be called P65.

MATH60132

pass mark 40

4 ques



MATH70132

pass: 50

+ 1 qu.

(on mastery)

Examinable: Lectures, problem sheets (inc. prob. classes)

Non-examinable: Certain long/complicated proofs:

Th. 2.3.3, Lemma 2.3.8, Th. 2.3.10, Th. 2.5.3, Lemma 2.6.3

Th. 3.1.6, Th. 3.5.2, Th. 3.6.1.

(Numbering as in typed notes.)

# Propositional logic .

Formulas ; truth tables / propositional valuations

i.e. of formulas ; adequacy of sets of connectives

'disjunctive normal form' (1.1.9 / 10)

Formal system L : axioms, proof, theorem  $\vdash_L \phi$

- a few examples

Deductions, consequences

$\sum \vdash_L \phi$

Deduction theorem

$\rightarrow$  more theorems of L .

Soundness (+ generalisation)

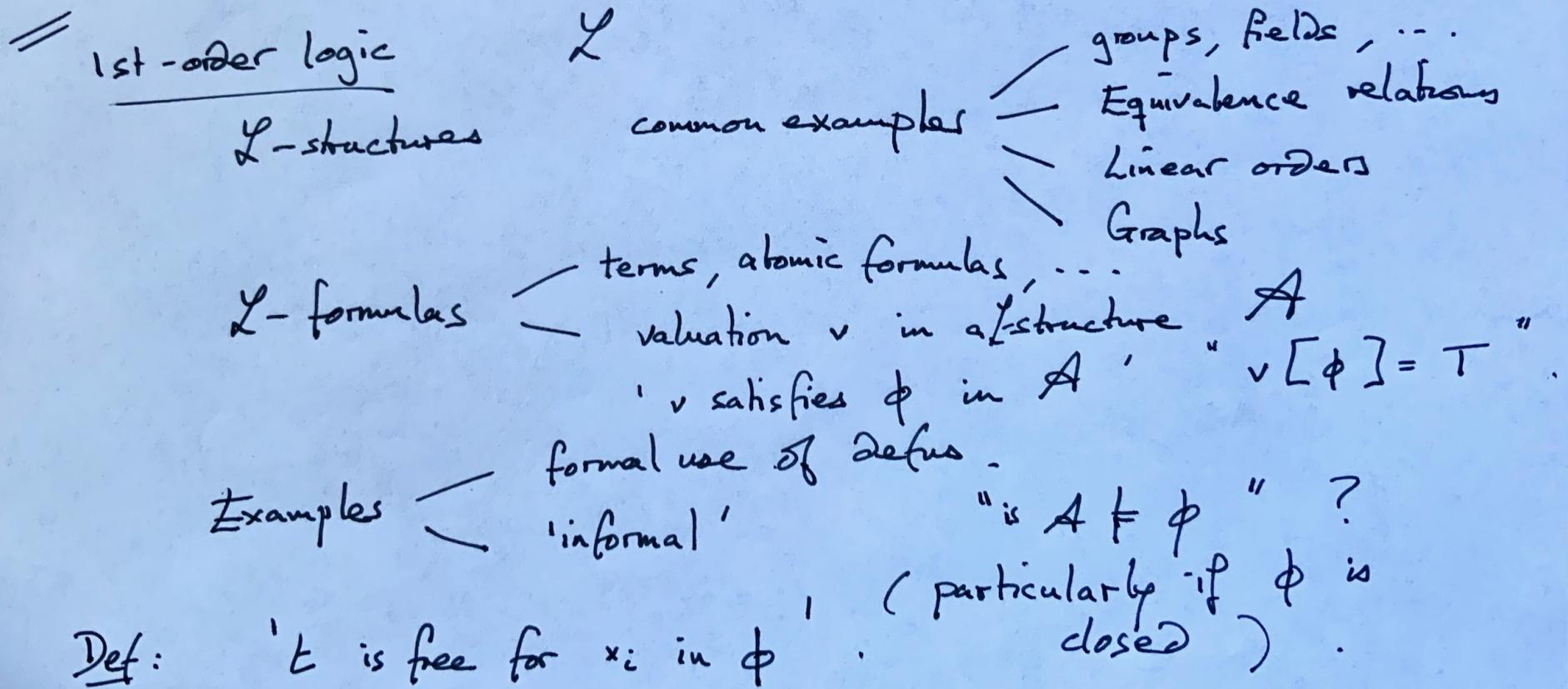
if  $\sum \vdash_L \phi$  then for every val.  $v$

with  $v(\sum) = T$ , we have  $v(\phi) = T$ .

Completeness / Adequacy thm. for L . Converse  
of this .

Key Lemma: If  $\Sigma$  is consistent &  $\Sigma \not\vdash \phi$  (3)  
then  $\Sigma \cup \{\phi\}$  is consistent.

Lindenbaum Lemma: ... there is  $\Sigma^* \supseteq \Sigma$  which is consistent  
and complete.



Formal system  $K_L$  — Axioms, Deduction rules  
MP + Gen (4)

$$\vdash_{K_L} \phi$$

$$\Sigma \vdash_{K_L} \phi$$

(restriction on  
use of Gen).

Deduction thm : as for L

Soundness + generalisation

Gödel Completeness Th + Generalisations

Hard part : 2.5-3 Gödel Existence Thm.

Compactness thm.

= Equality    Normal  $L^=$ -structures ; Axioms for equality

Compactness thm. for normal  $L^=$ -structures.

' Downward Lowenheim-Skolem thm.'

Application : Axiomatising  $\text{Th}((\mathbb{Q}; \leq))$  d.l.o.v.e.

(5)

Set theoryCardinality  $\approx$ 

Cantor - Schröder - Bernstein

Examples

ZF axioms(in  $= \in$ )

Orderings + operations :  $+$   $\times$  reverse lexicographic  
 $\approx$   $\equiv$  similar (isomorphic)

Well orderingsOrdinals - generalise natural numbers.Ordinals  $\alpha, \beta$   $\alpha < \beta \Leftrightarrow \alpha \in \beta \Leftrightarrow \alpha \subsetneq \beta$ .

Transfinite induction + recursion

↑ treat informally

$$\alpha \geq \omega \Rightarrow |\alpha \times \alpha| = |\alpha|$$

AC :  $\left\langle \begin{matrix} \text{WO} ; \text{ZL} \\ \text{Cardinality of } X ; |X|=1 \text{ cardinal} \end{matrix} \right\rangle$ ; Cardinal arithmetic.