

Mathematical Logic (MATH70132)
Mastery Material Problem Sheet

[1] Suppose $n \in \mathbb{N}$. The first-order language with equality \mathcal{L}_n^- has constant symbols c_1, \dots, c_n and no other relation, function or constant symbols (apart from equality). Write down a set T_n of closed \mathcal{L}_n^- -formulas whose normal models are precisely infinite sets in which the constant symbols c_1, \dots, c_n are interpreted as distinct elements. Use Vaught's Test (Theorem 8.18 in Cori - Lascar) to prove that T_n is complete.

[2] The language with equality \mathcal{L}_c^- has equality, a single binary relation symbol R and a constant symbols c_1, c_2 . How many non-isomorphic countable normal models are there in which R is interpreted as a dense linear ordering without endpoints? How many countable normal models are there which are not elementarily equivalent? What happens if the language has n constant symbols (rather than two), for $n \in \mathbb{N}$?

[3] Suppose \mathcal{L}^- is a language with equality and \mathcal{M} is a normal \mathcal{L}^- -structure with domain M . An *automorphism* of \mathcal{M} is an isomorphism $\alpha : \mathcal{M} \rightarrow \mathcal{M}$. Note that in this case, if $\phi(x_1, \dots, x_n)$ is an \mathcal{L}^- -formula and $a_1, \dots, a_n \in M$, then

$$\mathcal{M} \models \phi[a_1, \dots, a_n] \Leftrightarrow \mathcal{M} \models \phi[\alpha(a_1), \dots, \alpha(a_n)].$$

(i) With the above notation, suppose \mathcal{N} is a substructure of \mathcal{M} (with domain N) having the following property. For all $a_1, \dots, a_n \in N$ and $b \in M$ there is an automorphism α of \mathcal{M} with $\alpha(a_i) = a_i$, for $i \leq n$ and $\alpha(b) \in N$. Using the Tarski-Vaught Test, prove that \mathcal{N} is an elementary substructure of \mathcal{M} .

(ii) With T_n as in [1], show that if \mathcal{M} is a normal model of T_n , then every infinite substructure of \mathcal{M} is an elementary substructure.

(iii) Suppose \mathcal{M} is a vector space (in the usual language of vector spaces over a field F) and \mathcal{N} is a subspace of infinite dimension. Prove that \mathcal{N} is an elementary substructure of \mathcal{M} .

[4] Let \mathcal{L} be the usual language of groups and let $(\mathcal{M}_i : i \in I)$ be a family of groups. Suppose \mathcal{F} is a non-principal ultrafilter on I and consider the ultraproduct $\mathcal{M} = \prod_{i \in I} \mathcal{M}_i / \mathcal{F}$.

(i) Prove that $\{(a_i : i \in I) \in \prod_{i \in I} \mathcal{M}_i : \{i \in I : a_i = 1\} \in \mathcal{F}\}$ is a normal subgroup of $\prod_{i \in I} \mathcal{M}_i$ and that the quotient group by this is isomorphic to \mathcal{M} . (Here we are denoting by 1 the identity element of a group).

(ii) Let I be the set of prime numbers and \mathcal{M}_i the cyclic group of order i . Prove that \mathcal{M} is elementarily equivalent to $\langle \mathbb{Q}; + \rangle$ (the additive group of rational numbers).

[5] Let \mathcal{R} denote the structure $\langle \mathbb{R}; \leq, +, -, \cdot, 0, 1 \rangle$ in the usual language of rings with an ordering. Let \mathcal{F} be a non-principal ultrafilter on ω and consider the ultrapower $\mathcal{R}^* = \mathcal{R}^\omega / \mathcal{F}$. Say why we can regard \mathcal{R}^* as an elementary extension of \mathcal{R} . Decide which of the following are true, giving reasons for your answers.

(i) \mathcal{R}^* is a field.

(ii) Every polynomial of odd degree with coefficients in \mathcal{R}^* has a root in \mathcal{R}^* .

(iii) For every $r \in \mathcal{R}^*$ there is $n \in \mathbb{N}$ with $r < n$.

(iv) Every non-empty subset of \mathcal{R}^* which is bounded above in \mathcal{R}^* has a least upper bound in \mathcal{R}^* .