

Applied Complex Analysis - Lecture Eleven

Andrew Gibbs

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The trapezium rule

Trapezium rule on the unit circle

$$I = \oint_{|z|=1} \frac{f(z)}{z} dz$$

- Changing variables $z = e^{i\theta}$, for $\theta \in [0, 2\pi]$

$$I = \frac{1}{2\pi} \int_0^{2\pi} f(e^{i\theta}) d\theta.$$

- Trapezium rule approx:

$$I \approx I_N = \frac{1}{N} \sum_{j=1}^N f(z_j), \quad \text{where } z_j = e^{2\pi i j/N}, \quad \text{for } j = 1, \dots, N$$

- **Thm:** Suppose f is analytic and satisfies $|f(z)| < M$ inside the complex disk $|z| < r$ for some $r > 1$. Then

$$|I - I_N| \leq \frac{M}{r^N - 1} = O(r^{-N}) \quad \text{as } N \rightarrow \infty.$$

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$$\cos(0) = 1$$

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Application: Rounding errors

Double precision floating point numbers are stored as

$$(-1)^S \times 2^{E-1023} \times (1 + M)$$

where one bit is allocated for the sign S , 11 bits are allocated for the exponent E , and 52 bits are allocated for the Manissa M , this gives us 16 digits of precision.

- Taking differences of small numbers is OK
- Taking differences of larger numbers can be an issue
- For example:

$$1.000000012345678 - 1.000000000000000 = 0.000000012345678$$

- Sometimes, lots of digits are needed!
- Errors are exaggerated when dividing by differences of large numbers
- **Example** - removable singularity:

$$f(z) = \frac{e^z - 1 - z}{z^2}, \quad f(10^{-8}) = 0.500000001666667$$

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Approximating derivatives

- A popular (but dodgy) way to approximate derivatives is to use a 'finite difference approach', for EG:

$$f'(z) = \frac{f(z + \epsilon) - f(z - \epsilon)}{2\epsilon}, \quad \text{for small } \epsilon.$$

- Choosing the right ϵ is difficult.
- Too big \implies inaccurate
- Too small \implies rounding errors (also inaccurate!)
- for high order derivatives this can get very messy.
- Thanks to Cauchy's integral formula, we can easily approximate derivatives:

$$f^{(n)}(z) = \frac{n!}{2\pi i} \int_{\gamma} \frac{f(\xi)}{(\xi - z)^n} d\xi \approx \frac{n!}{N} \sum_{j=1}^N \frac{f(re^{2\pi i j/N})}{r^j e^{2\pi i \theta j(n-1)/N}}$$

- But, our current estimates cannot be applied here...

Trapezium rule for periodic analytic functions

Trapezium rule for periodic functions

Thm: For

$$I = \int_0^{2\pi} f(\theta) d\theta \approx I_N = \frac{1}{N} \sum_{n=1}^N f(\theta_n), \quad \theta_n = 2\pi n/N,$$

if f is 2π -periodic and analytic in the strip $S_a := \{\theta : -a < \operatorname{Im}\{\theta\} < a\}$ for $a > 0$, then

$$|I - I_N| \leq \frac{4\pi M}{e^{aN} - 1},$$

where $M := \sup_{\theta \in S_a} |f(\theta)|$.

- Convergence of real integral depends on complex behavior
- More general class: may be applied to **non-circular contour integrals** $\oint_{\gamma} f$, where f is analytic in an annulus containing γ
- Requires the parametrisation of γ to be analytic
- **Proof**

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Exactness of trapezium rule for periodic functions

Thm: If a function f has a Laurent expansion of the form

$$f(z) = \sum_{j=-N}^{N-2} a_j (z - z_0)^j,$$

for z in some annulus D , then an N -point trapezium rule I_N can exactly approximate

$$I = \oint_{\gamma} f(z) dz,$$

where γ is a closed anti-clockwise-oriented contour in D .

- Analogous to N -point Gauss quadrature integrating degree $(2N - 1)$ -degree integrals exactly
 - Not particularly useful - but a helpful way to check your code!
 - Proof
 - Example
 - Numerical examples

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The argument principle

For f meromorphic and g analytic in Ω ,

$$\frac{1}{2\pi i} \oint_{\gamma} \frac{f'(z)}{f(z)} g(z) dz = \sum_{a \in \{\text{zeros of } f\}} g(a) m_a - \sum_{b \in \{\text{poles of } f\}} g(b) m_b.$$

where m_a and m_b represent the order of the zeros and poles respectively, γ is a closed contour in Ω with no loops, such that $f(z) \neq 0$ for $z \in \gamma$.

- **Proof**
- Consequences and applications
- Approximation trick when $g = 1$

Application: Root-finding