

## Computationally Hard Problems – Fall 2016 Assignment 2

**Date:** 06.09.2016, **Due date:** 12.09.2016, 21:00

The following exercises are **not** mandatory:

**Exercise 2.1:** Prove Proposition 2.14 in the lecture notes, which is also repeated below, and determine the constants  $c_p$  and  $c_e$ . You may assume that the running time for input size  $\ell$  is  $\ell^a$  for a positive constant  $a > 0$  in the case of polynomial running time, and  $b^\ell$  for a positive constant  $b > 1$  in the case of exponential running time.

**Proposition 2.14:** Suppose an algorithm has been implemented on some machine. With the current hardware, this can solve problems up to size  $n$  in a fixed time  $t_0$ . Now suppose that the speed of the hardware is doubled. If the running time of the algorithm is polynomial then there is a constant  $c_p > 1$  such that one can now solve problems of size  $c_p \cdot n$  in time  $t_0$ . If the running time of the algorithm is exponential then there is a constant  $c_e > 0$  such that one can now solve problems of size  $c_e + n$  in time  $t_0$ .

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End of Exercise 1

**Exercise 2.2:** Show how to convert a decision algorithm  $A_d$  for the following problem into a polynomial-time optimization algorithm  $A_o$ . The input to the problem consists of an undirected graph  $G = (V, E)$  and a positive integer  $k$ . The decision algorithm  $A_d$  answers YES if there is a set  $V' \subseteq V$  of cardinality  $k$  such that  $\forall v, w \in V': \{v, w\} \notin E$ . This problem is called INDEPENDENTSET since there are no edges between the vertices in  $V'$ .

We now want to solve the optimization problem, that is, we are only given the undirected graph  $G$  and want to find in it an independent set  $V'$  of maximum cardinality.

- a) Describe an algorithm  $A_o$  which solves the optimization problem as described above. The running time of  $A_o$  has to be polynomial in the input size and  $A_o$  may make calls to  $A_d$ . Recall that such a call counts as one step.
- b) Argue that your algorithm is correct.
- c) Show the running time of the algorithm.

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End of Exercise 2

Continued on next page.

**Exercise 2.3:** Consider the following problem.

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**Problem:** [MINIMUMTESTSET]

**Input:** You want to test an electronic circuit. The circuit can have  $n$  different errors,  $e_1, \dots, e_n$ . You can perform  $m$  different tests  $T_1, \dots, T_m$ . Every test can detect some, but not necessarily all, errors.

**Output for the optimizing version:** A minimum set of tests that detects all possible errors.

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We suggest the following algorithm for selecting a minimum test set: Select a test  $T_i$  which detects a maximum number of errors. Continue by selecting tests which detect a maximum number of errors, which have not already been detected by previously selected tests.

Prove or disprove that the output of this algorithm is a minimum test set.

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End of Exercise 3

The following exercise is **mandatory**:

**Exercise 2.4:** We consider the following problem.

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**Problem:** [REFUTATION] Given is a disjunctive form consisting of  $k$  monomials  $m_1, \dots, m_k$  over  $n$  boolean variables  $x_1, \dots, x_n$  (as on Exercise Sheet 1). The task is to decide if there is a truth assignment to the variables such that the truth value of the disjunctive form is *false*.

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You are given a decision algorithm  $A_d$  that solves this problem, i. e., for each instance to REFUTATION,  $A_d$  answers YES if there is an assignment that makes the truth value of the disjunctive form *false*; otherwise it answers NO.

- a) Describe an algorithm  $A_o$  which solves the optimization problem, that is, which finds a truth assignment making the disjunctive form *false* if one exists. The running time has to be polynomial in the input size and  $A_o$  may make calls to  $A_d$ . Such calls count as one basic computational step.
- b) Argue that your algorithm is correct.
- c) Prove that the running time of the algorithm is bounded from above by a polynomial. Any polynomial is sufficient; you need not look for a polynomial of minimal degree. Recall that a call to  $A_d$  counts one step.

**Note:** The input to REFUTATION is a disjunctive form of monomials over  $n$  boolean variables, nothing else. In particular, a legal input cannot specify specific settings of variables.

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End of Exercise 4