

Non-convergence Analysis of Probabilistic Direct Search

2nd Derivative-Free Optimization Symposium

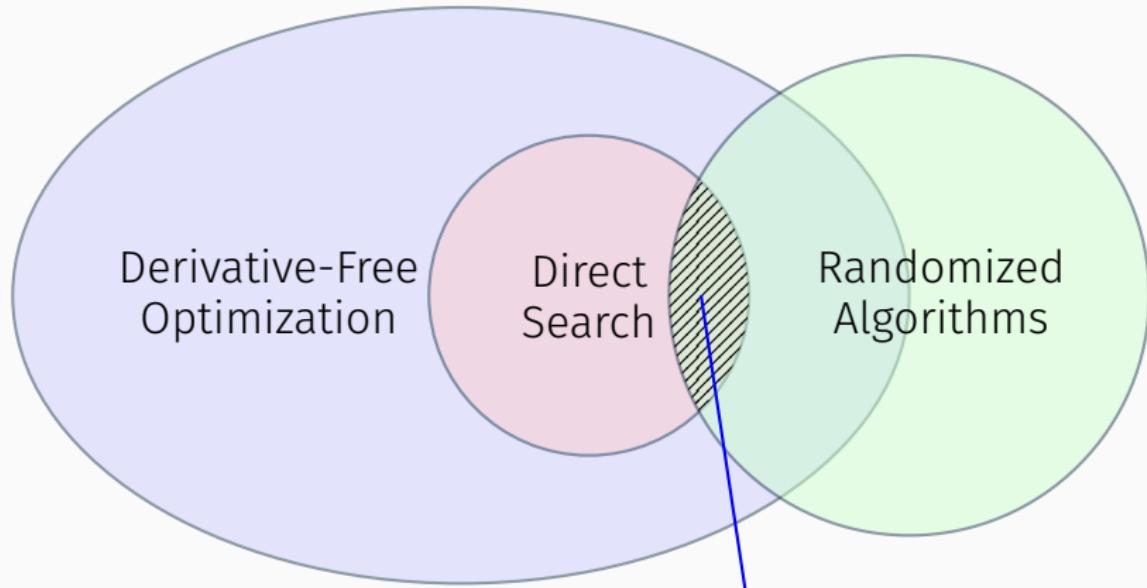
Cunxin Huang

Supervised by Prof. Xiaojun Chen and Dr. Zaikun Zhang

Padova, Italy June 28, 2024

The Hong Kong Polytechnic University

Brief introduction to Probabilistic Direct Search



The algorithm we consider in this talk:
Probabilistic Direct Search (PDS)
(Gratton, Royer, Vicente, and Zhang 2015)

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Solution: I will show some computation works at the end.

What is Derivative-Free Optimization and why

Derivative-Free Optimization (DFO)

- Do not use derivatives (first-order info.), only use function values
- Also called: zeroth-order/black-box/simulation-based optimization

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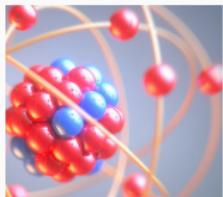
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Nuclear Physics



Machine Learning



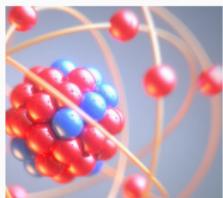
Circuit Design

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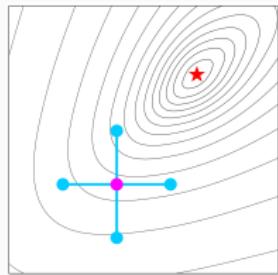
Difficulties

- Problems are often **noisy** (naive finite difference?)
- Each function evaluation is **expensive** (e.g., PDE simulation)

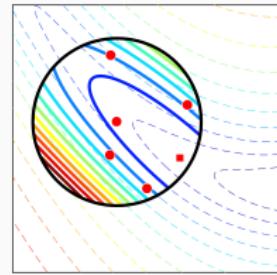
Direct-search methods and model-based methods

How to determine iterates?

- Direct-search methods: “simple” comparison of function values
- Model-based methods: build a surrogate of the objective function



Direct-search methods¹



Model-based methods²

¹Source: Kolda, Lewis, and Torczon 2003

²Source: Larson, Menickelly, and Wild 2019

Probabilistic Direct Search (PDS): a simplified framework

Algorithm 1: Direct Search based on **sufficient decrease**

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if $f(x_k + \alpha_k d_k) < f(x_k) - c\alpha_k^2$ **then**

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 (Move and expand step size)

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Typical choice of $\{\mathcal{D}_k\}$ (Gratton, Royer, Vicente, and Zhang 2015):

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N.B.: typical choice in the deterministic case is $\{\pm e_i\}_{i=1}^n$, Coordinate Search (CS)

Illustration of how PDS works

$\mathcal{D}_k = \{d_1, d_2\}$, where $d_\ell \stackrel{\text{i.i.d.}}{\sim} \mathbf{U}(\mathcal{S}^1)$

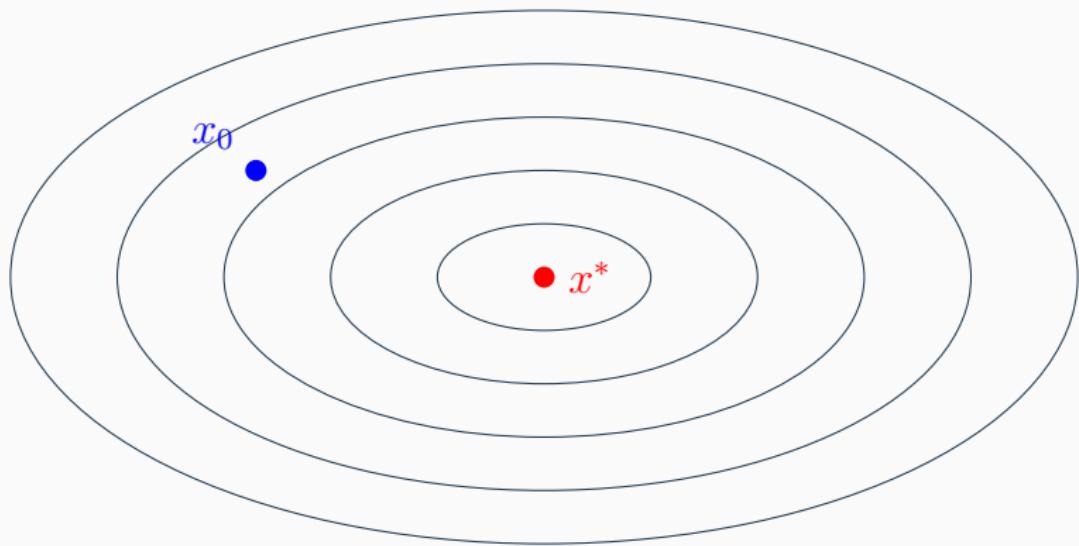


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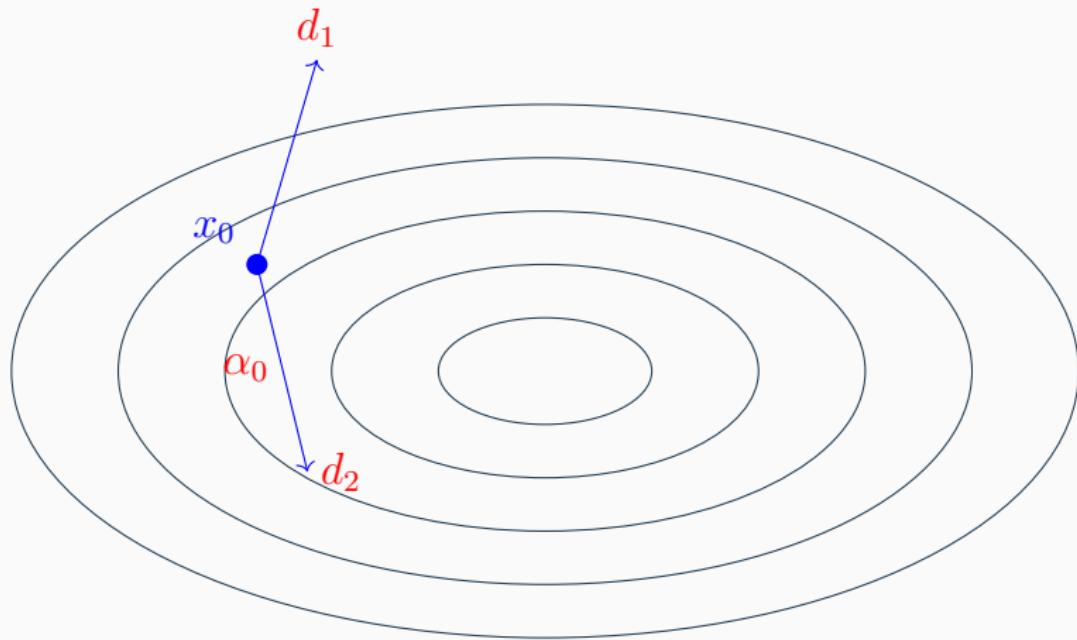


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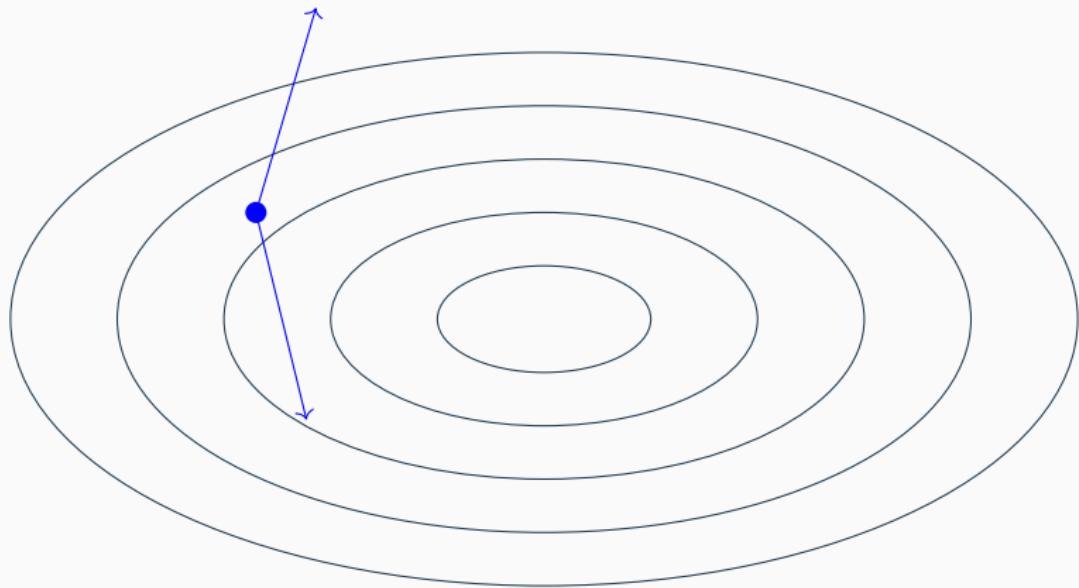


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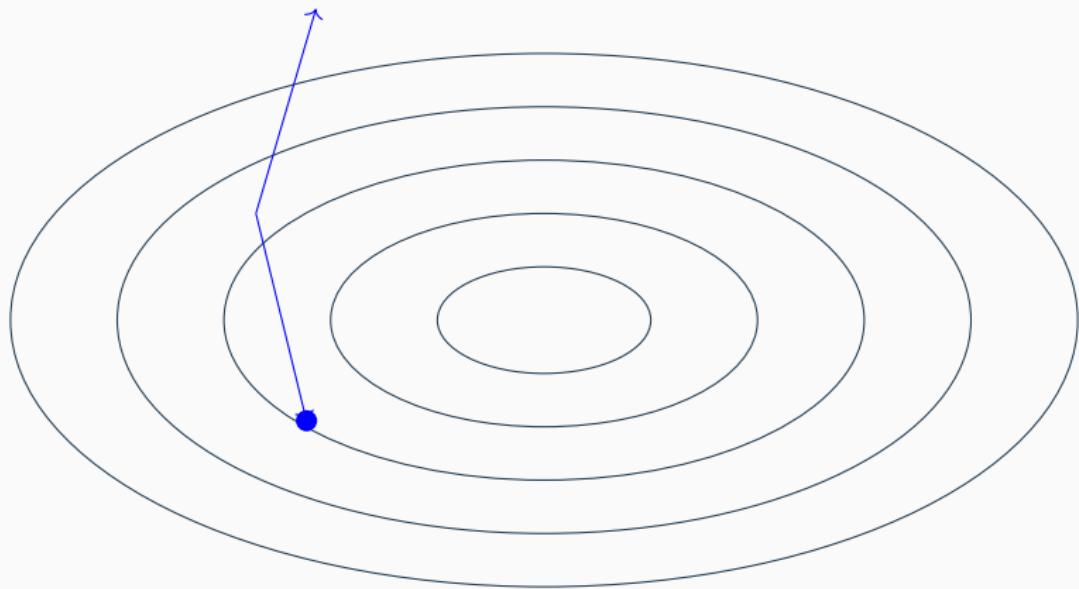


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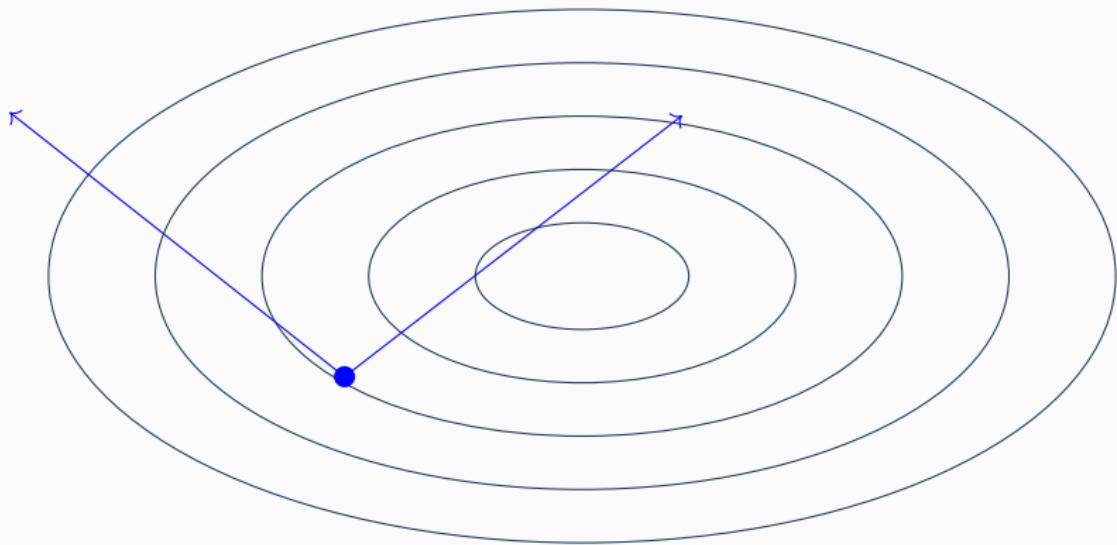


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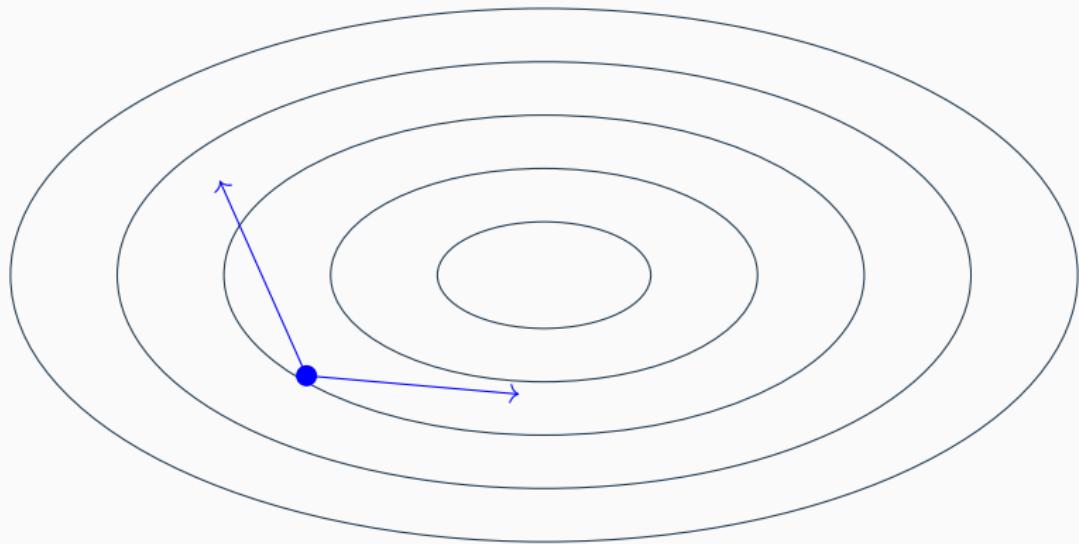


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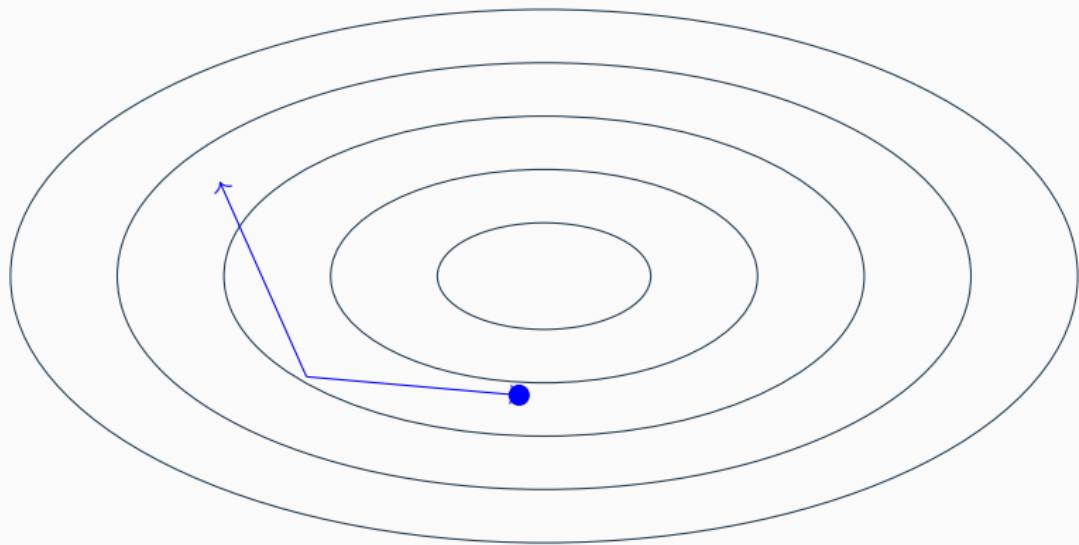


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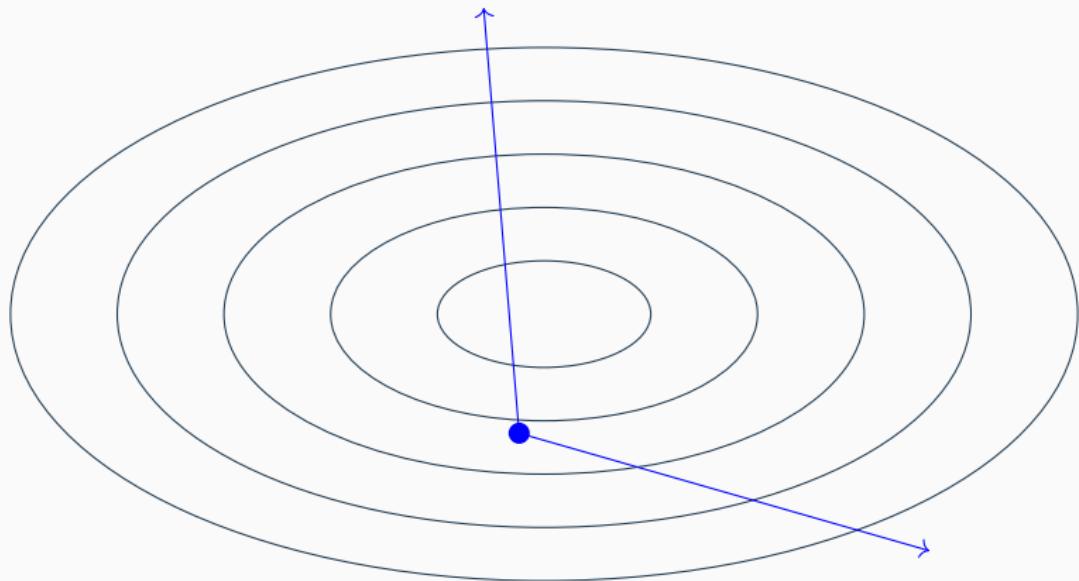


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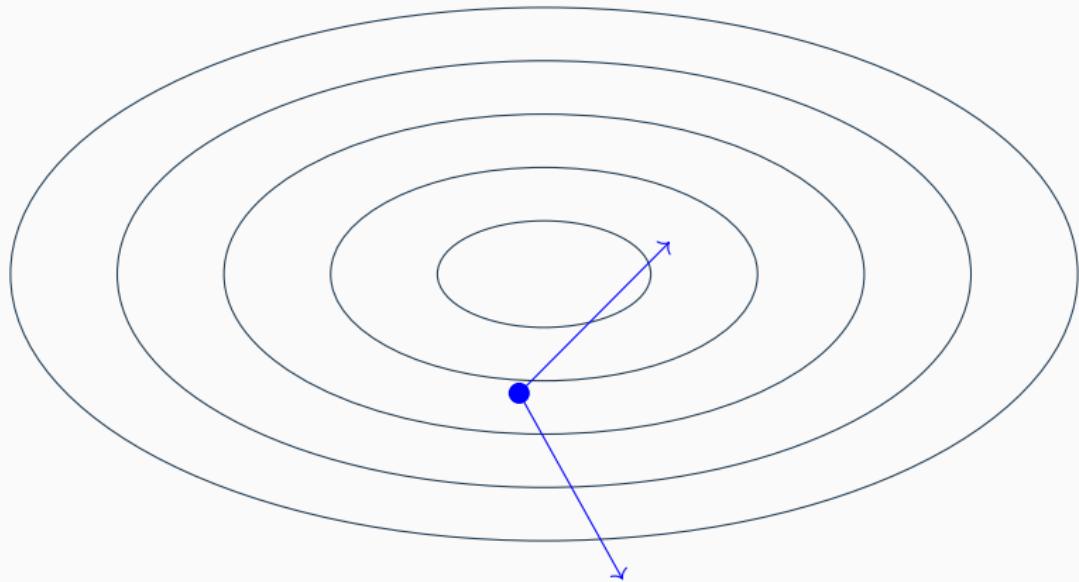


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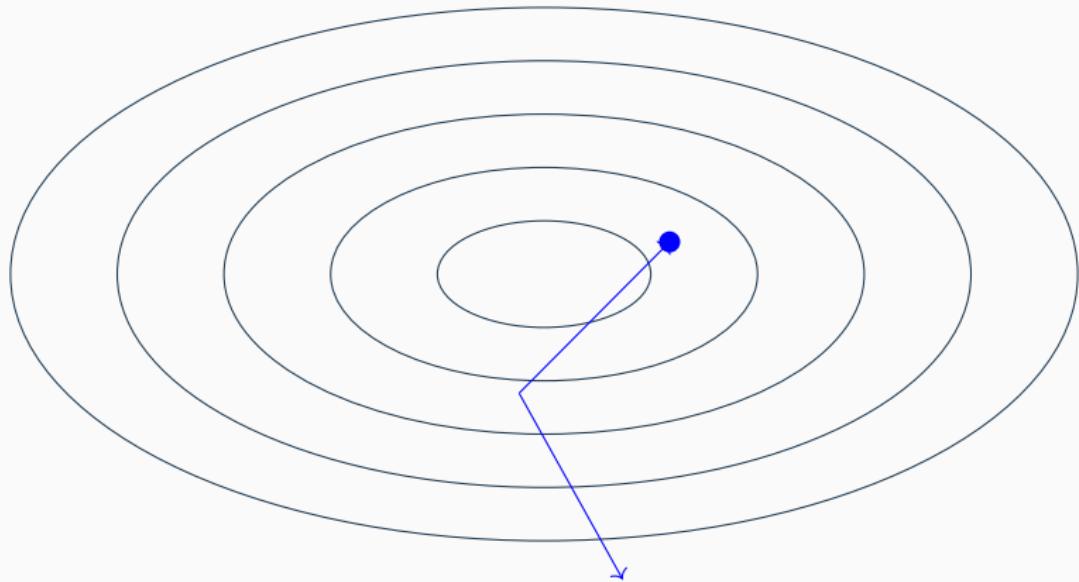


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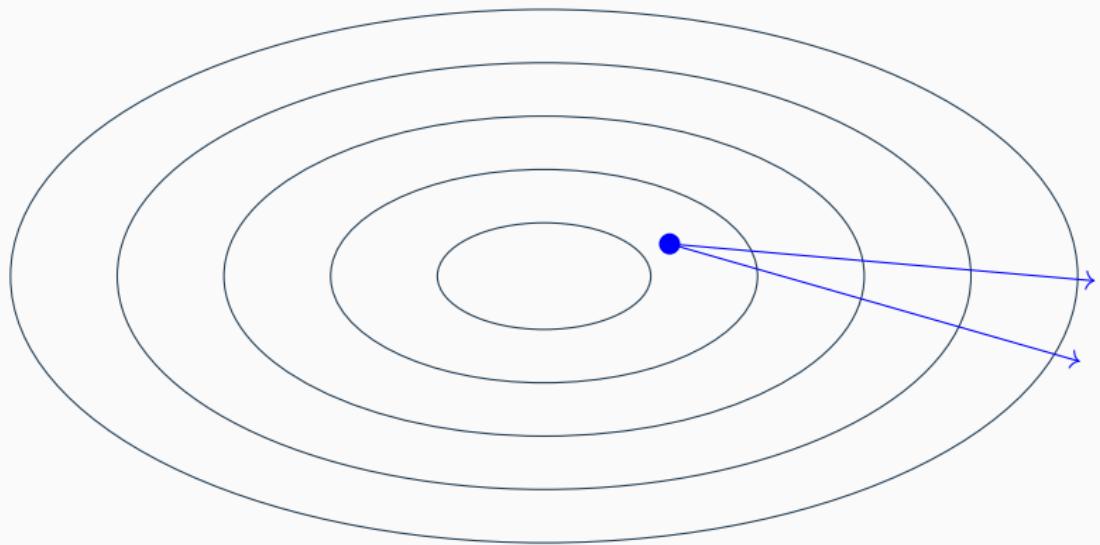


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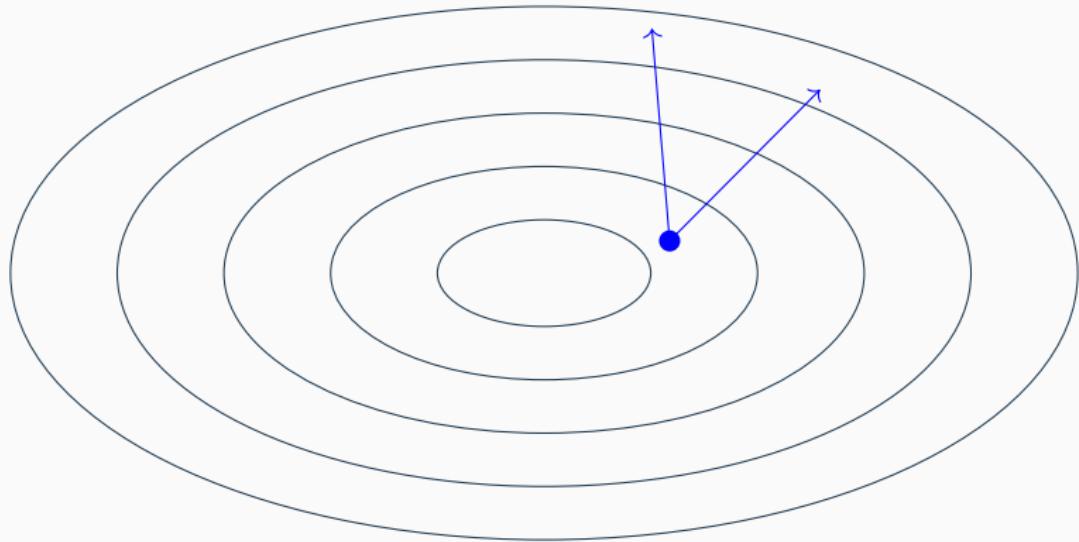
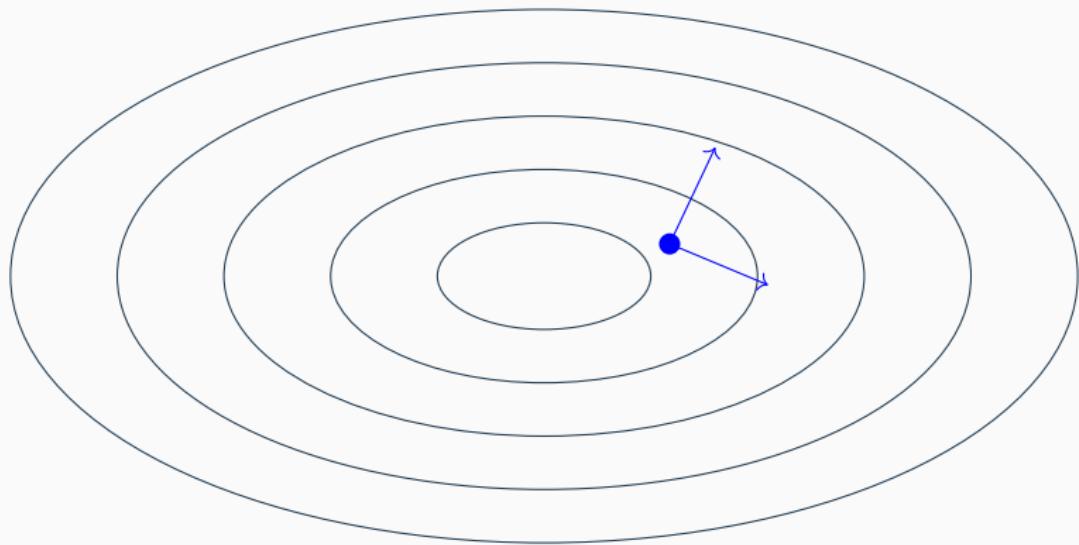


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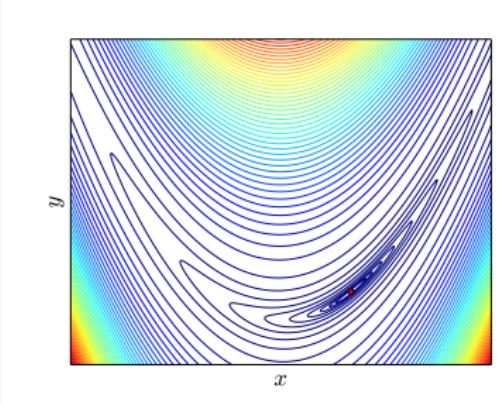
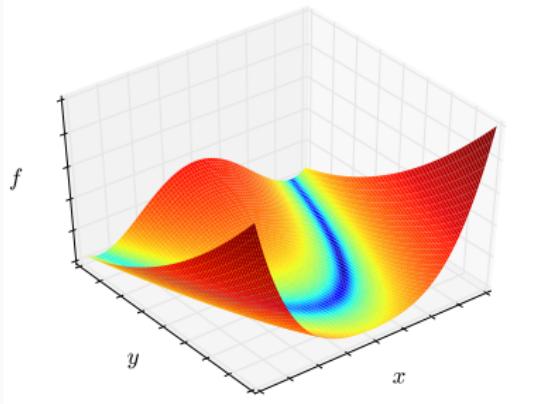
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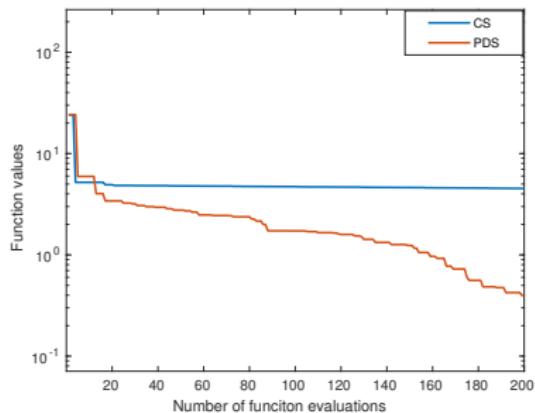
A numerical example: CS v.s. PDS with 2 directions

Rosenbrock “banana” function:

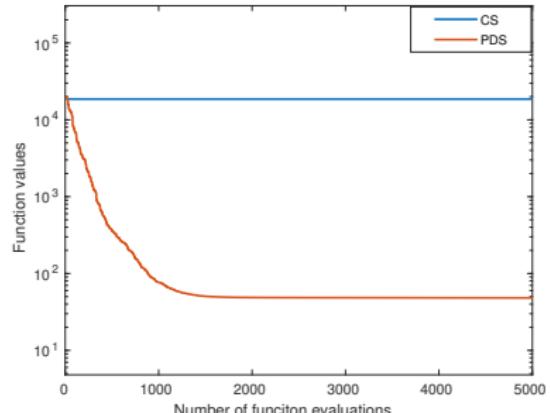
$$f(x) = \sum_{i=1}^{n-1} [(1 - x_i)^2 + 100(x_{i+1} - x_i^2)^2]$$



A numerical example: CS v.s. PDS with 2 directions



$$n = 2$$



$$n = 50$$

Function value v.s. number of function evaluations

Worst case complexity of function evaluations (GRVZ 2015)

$\mathcal{O}(n^2\epsilon^{-2})$ for CS while $\mathcal{O}(n\epsilon^{-2})$ for PDS

Cosine measure

Definition (Cosine measure w.r.t. a vector)

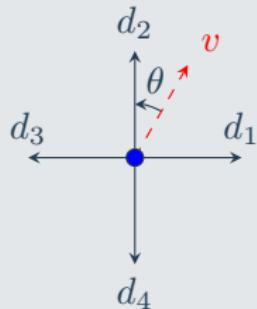
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$$\text{cm}(\mathcal{D}, v) = \max_{d \in \mathcal{D}} \frac{d^\top v}{\|d\| \|v\|},$$

which is the cosine measure of \mathcal{D} with respect to v .

Example

$$\text{cm}(\mathcal{D}, v) = \cos \theta$$



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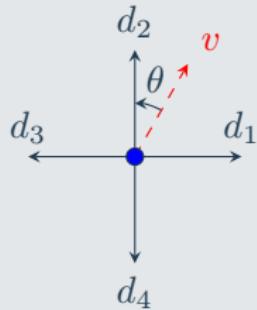
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$\text{cm}(\mathcal{D}, v)$ measures the ability of \mathcal{D} to “approximate” v

Convergence theory

Definition (p -probabilistically κ -descent)

$\{\mathcal{D}_k\}$ is said to be p -probabilistically κ -descent, if

$$\mathbb{P}(\text{cm}(\mathcal{D}_k, -g_k) \geq \kappa \mid \mathcal{D}_0, \dots, \mathcal{D}_{k-1}) \geq p \quad \text{for each } k \geq 0,$$

where $g_k = \nabla f(x_k)$.

Intuition

Each \mathcal{D}_k is “good enough” with probability at least p
no matter what has happened in the history

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Theorem (GRVZ, 2015)

If $\{\mathcal{D}_k\}$ is p_0 -probabilistically κ -descent with $\kappa > 0$ and

$$p_0 = \frac{\log \theta}{\log(\gamma^{-1}\theta)},$$

then PDS converges w.p.1 when f is L -smooth and lower-bounded.

Practical choice and natural question

Corollary (GRVZ, 2015)

If $\mathcal{D}_k = \{d_1, \dots, d_m\}$, where $d_\ell \stackrel{i.i.d.}{\sim} U(\mathcal{S}^{n-1})$, then PDS converges w.p.1 if

$$m > \log_2 \left(1 - \frac{\log \theta}{\log \gamma} \right).$$

Practical choice and natural question

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Moreover, are supermartingale-like assumptions essential?

$$\mathbb{P}(\text{some event} \mid \mathcal{F}) \geq p$$

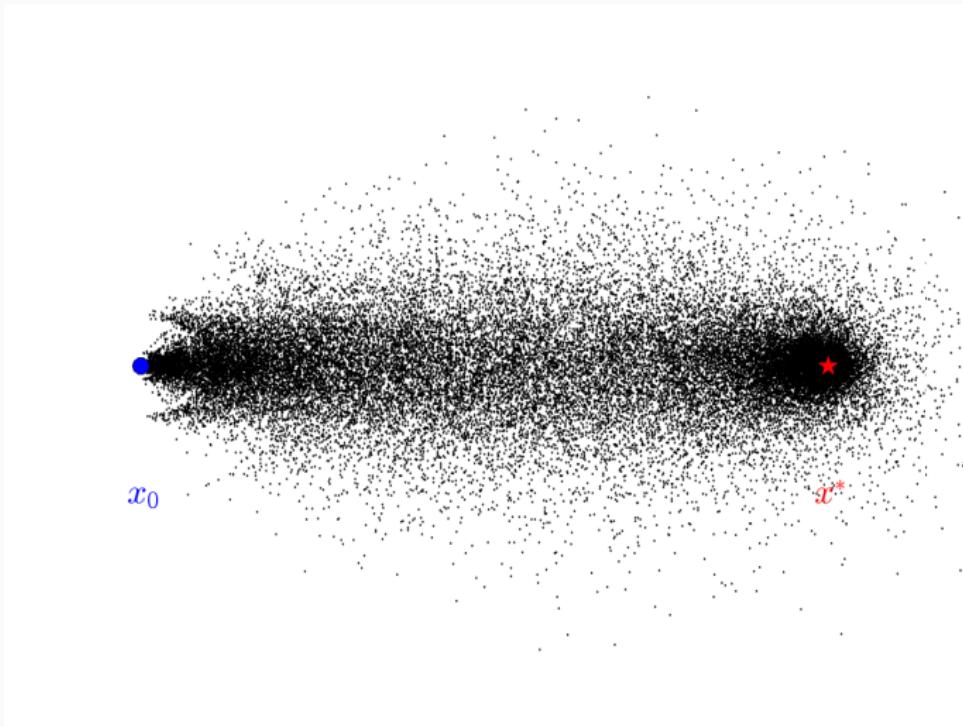
Related talks: Coralia, Kwassi Joseph, Matt, Anne, Warren, Sara, Lindon

A simple test

- Objective function: $f(x) = \|x\|^2/2$
- Initial point: $x_0 = (-10, 0)^\top$
- Stopping criterion: $\alpha_k \leq \text{machine epsilon}$
- Number of experiments: 100,000
- Parameters of PDS: $\alpha_0 = 1, \theta = 0.25, \gamma = 1.5, m = 2$

$$m = 2 < 2.143 \approx \log_2 \left(1 - \frac{\log \theta}{\log \gamma} \right)$$

A simple test (Cont'd)



Note: each **black dot** represents the **output point** of one run of PDS.

Non-convergence study is not rare

Many well-known algorithms have non-convergence [examples](#)

- Powell, On search directions for minimization algorithms, 1973.
- Yuan, An example of non-convergence of [trust region](#) algorithms, 1998.
- Reddi, Kale, and Kumar, On the convergence of [Adam](#) and beyond, 2018.
- Chen, He, Ye, and Yuan, The direct extension of [ADMM](#) for multi-block convex minimization problems is not necessarily convergent, 2016.
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Instead of finding a non-convergence example,
can we develop a theorem?

An overview of our theory

We assume that f is smooth and **convex** (explained later).

We denote the optimal solution set of f by \mathcal{S}^* .

We will establish the following.

Under **some assumption on $\{\mathcal{D}_k\}$** and algorithmic parameters, there exist **choices of x_0** such that

$$\mathbb{P} \left(\liminf_{k \rightarrow \infty} \text{dist}(x_k, \mathcal{S}^*) > 0 \right) > 0.$$

Differences from a non-convergence example

one function	v.s.	some function class
special parameters	v.s.	conditions for parameters
a specific initial point	v.s.	a region for initial points

Assumption on $\{\mathcal{D}_k\}$: probabilistic ascent

Recall p -probabilistically κ -descent

$$\mathbb{P}(\text{cm}(\mathcal{D}_k, -g_k) \geq \kappa \mid \mathcal{D}_0, \dots, \mathcal{D}_{k-1}) \geq p \quad \text{for each } k \geq 0.$$

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q -probabilistically **ascent**

$$\mathbb{P}(\text{cm}(\mathcal{D}_k, -g_k) \leq 0 \mid \mathcal{D}_0, \dots, \mathcal{D}_{k-1}) \geq q \quad \text{for each } k \geq 0.$$

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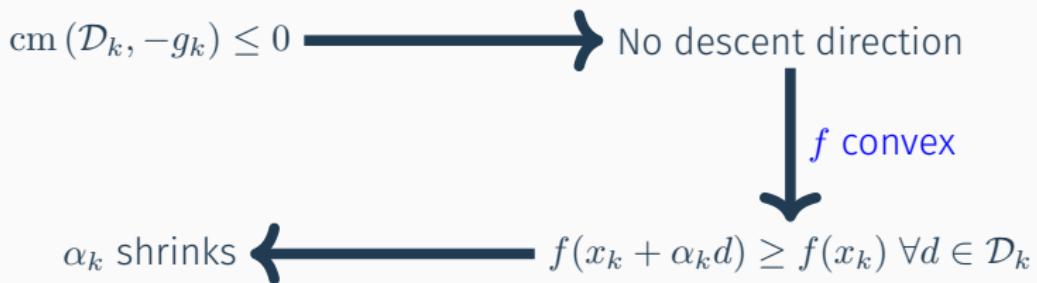
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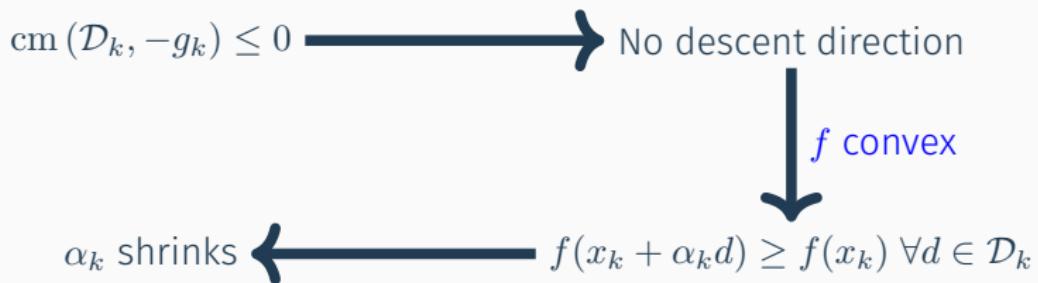
Note

If $\text{cm}(\mathcal{D}_k, -g_k) \leq 0$, then \mathcal{D}_k is “bad” (no descent direction).

Why assuming convexity?

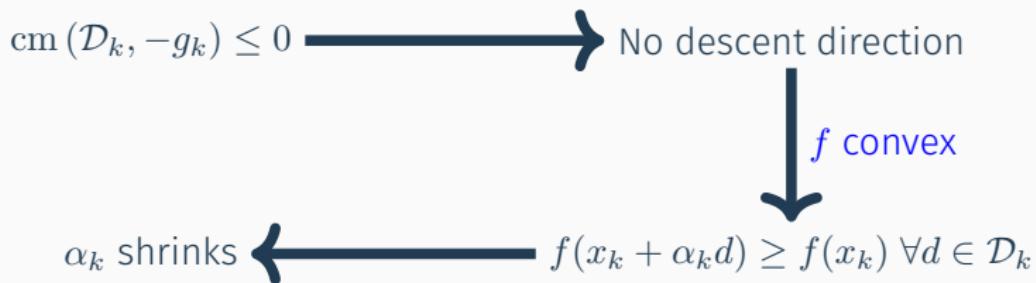


Why assuming convexity?



- Convexity connects $\text{cm}(\mathcal{D}_k, -g_k) \leq 0$ and shrinking of step size

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- Convexity connects $\text{cm}(\mathcal{D}_k, -g_k) \leq 0$ and shrinking of step size
- $\{\mathcal{D}_k\}$ is probabilistic ascent implies α_k “often” shrinks

From probabilistic ascent to non-convergence: How?

$\{\mathcal{D}_k\}$ is probabilistically ascent

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$$\mathbb{P} \left(\sum_{k=0}^{\infty} \alpha_k \text{ is “bounded”} \right) > 0 ?$$

From probabilistic ascent to non-convergence: How?

$\{\mathcal{D}_k\}$ is probabilistically ascent



α_k “often” shrinks



$$\mathbb{P} \left(\sum_{k=0}^{\infty} \alpha_k \text{ is “bounded”} \right) > 0 ?$$



$\mathbb{P}(\text{non-convergence}) > 0$ if $\text{dist}(x_0, \mathcal{S}^*)$ is “large”?

Key ingredients of the analysis

- Define the indicator function for “bad \mathcal{D}_k ”

$$Y_k = \mathbb{1}(\text{cm}(\mathcal{D}_k, -g_k) \leq 0)$$

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- Note the following inequality between step sizes (f is convex)

$$\alpha_{k+1} \leq \begin{cases} \gamma\alpha_k, & \text{if } Y_k = 0 \\ \theta\alpha_k, & \text{if } Y_k = 1 \end{cases} = \gamma^{1-Y_k} \theta^{Y_k} \alpha_k$$

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$$\alpha_k \leq \alpha_0 \prod_{\ell=0}^{k-1} \gamma^{1-Y_\ell}\theta^{Y_\ell}$$

- Get an upper bound of series of step sizes

$$\sum_{k=1}^{\infty} \alpha_k \leq \alpha_0 \sum_{k=1}^{\infty} \prod_{\ell=0}^{k-1} \gamma^{1-Y_\ell}\theta^{Y_\ell} =: \alpha_0 S$$

- Analyze the behavior of the random series S

A closer look at the random series S

Recall that

$$S = \sum_{k=1}^{\infty} \prod_{\ell=0}^{k-1} \gamma^{1-Y_\ell} \theta^{Y_\ell},$$

where $Y_\ell = \mathbb{1}(\text{cm}(\mathcal{D}_\ell, -g_\ell) \leq 0)$.

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Two questions

- (Q1) Does there exist a constant ζ such that

$$\mathbb{P}(S < \zeta) > 0?$$

- (Q2) Moreover, can we specify the value of ζ ?

Answer to Q1 and Q2

Proposition

If $\{\mathcal{D}_k\}$ is q -probabilistically ascent with $q > q_0$, where

$$q_0 = 1 - p_0 = \frac{\log \gamma}{\log(\theta^{-1}\gamma)},$$

then

1.

$$\mathbb{P}(S < \infty) = 1,$$

2.

$$\mathbb{P}(S < \zeta) > 0 \iff \zeta > \frac{\theta}{1 - \theta}.$$

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Note

- $\mathbb{P}(S < \infty) = 1$ implies the existence of a ζ but not its value.
- The lower bound in 2 is tight, as $S = \sum_{k=1}^{\infty} \prod_{\ell=0}^{k-1} \gamma^{1-Y_\ell} \theta^{Y_\ell} \geq \frac{\theta}{1 - \theta}$.

Non-convergence of PDS

Theorem

Under aforementioned assumptions on f , if the sequence $\{\mathcal{D}_k\}$ in PDS is q -probabilistically ascent with $q > q_0$, then

$$\mathbb{P} \left(\liminf_{k \rightarrow \infty} \text{dist}(x_k, \mathcal{S}^*) > 0 \right) > 0,$$

provided that $\text{dist}(x_0, \mathcal{S}^*) > \alpha_0 / (1 - \theta)$.

Weaker assumption than probabilistic ascent

Denote $\mathbb{P}(\text{cm}(\mathcal{D}_k, -g_k) \leq 0 \mid \mathcal{D}_0, \dots, \mathcal{D}_{k-1})$ by P_k .

Recall that $\{\mathcal{D}_k\}$ is q -probabilistically ascent if $P_k \geq q$ for each $k \geq 0$.

Note that $\{P_k\}$ are random variables.

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For the latter, we can relax the assumption

$$\text{from } P_k \geq q > q_0 \quad \text{to} \quad \mathbb{P}\left(\liminf_{k \rightarrow \infty} P_k > q_0\right) > 0.$$

What happens in the typical implementation of PDS?

Let $\mathcal{D}_k = \{d_1, \dots, d_m\}$, where $d_\ell \stackrel{\text{i.i.d.}}{\sim} \mathcal{U}(\mathcal{S}^{n-1})$.

Recall that PDS is convergent if

$$m > \log_2 \left(1 - \frac{\log \theta}{\log \gamma} \right).$$

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With our non-convergence analysis, PDS is non-convergent if

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or, equivalently,

$$m < \log_2 \left(1 - \frac{\log \theta}{\log \gamma} \right).$$

Assumptions for convergence and non-convergence are essential.

Tightness of our assumption on $\{\mathcal{D}_k\}$

Our assumption on $\{\mathcal{D}_k\}$:

q -probabilistically ascent with $q > q_0$.

Natural question:

Is it sufficient to require $q \geq q_0$?

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Answer: NO!

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Natural question:

Is it sufficient to require $q \geq q_0$?

Answer: NO!

Example

We assume

- $\theta = 1/2$ and $\gamma = 2$, which implies $q_0 = 1/2$;
- $\mathcal{D}_k = \{g_k/\|g_k\|\}$ or $\{-g_k/\|g_k\|\}$ with probability $1/2$, respectively.

Then PDS converges w.p.1.

Convergence result inspired by non-convergence analysis

Define a series

$$S(\kappa) = \sum_{k=1}^{\infty} \prod_{\ell=0}^{k-1} \gamma^{Z_\ell(\kappa)} \theta^{1-Z_\ell(\kappa)},$$

where $Z_\ell(\kappa) = \mathbb{1}(\text{cm}(\mathcal{D}_\ell, -g_\ell) \geq \kappa)$.

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If there exists a $\kappa > 0$ such that $S(\kappa) = \infty$, then DS converges.

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If there exists a $\kappa > 0$ such that $S(\kappa) = \infty$, then DS converges.

Relation with existing result (GRVZ, 2015)

p_0 -probabilistically κ -descent $\implies S(\kappa) = \infty$ w.p.1

Take away

In this talk, we

- theoretically explain the non-convergence phenomenon of PDS,
- find out the behavior of PDS is closely related to the random series

$$S = \sum_{k=1}^{\infty} \prod_{\ell=0}^{k-1} \gamma^{1-Y_\ell} \theta^{Y_\ell}.$$

Non-convergence analysis can

- verify whether your assumption for convergence is essential,
- deepen our understanding of mathematical tools we use,
- provide new perspectives on convergence analysis,
- guide the choice of algorithmic parameters,

One more thing: OptiProfiler

OptiProfiler (joint work with Tom M. Ragonneau and Zaikun Zhang) is
a [benchmarking platform](#) for DFO solvers.

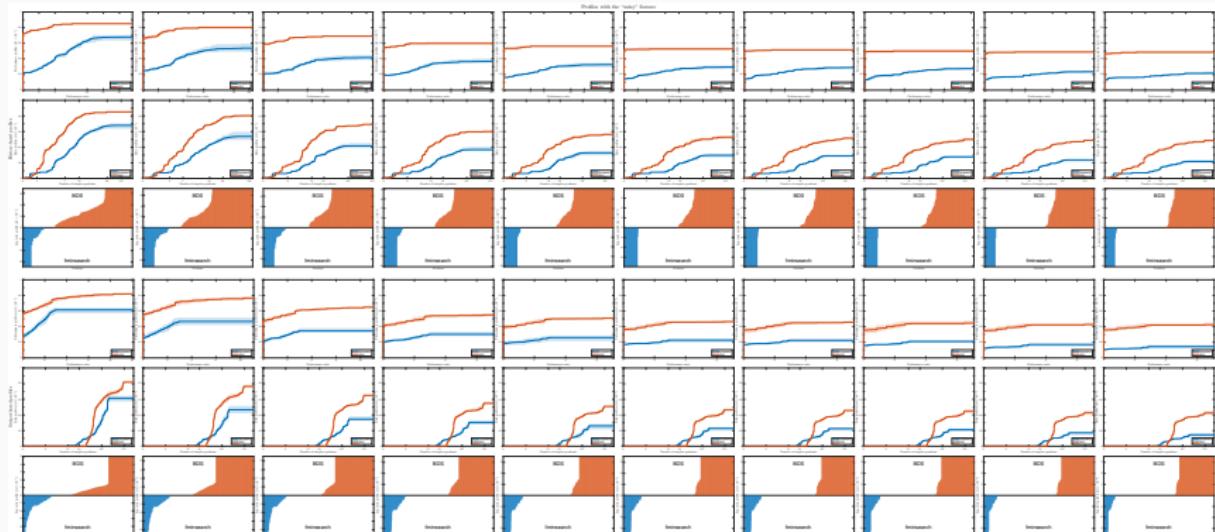
Our goal: [fair](#), [convenient](#), and [uniform](#) benchmarking.

- Creating [performance profiles](#), [data profiles](#), and [log-ratio profiles](#).
[Moré, Wild, 2009; Shi, Xuan, Oztoprak, and Nocedal, 2023]
Thanks for Nikolaus's nice talk: runtime distributions and COCO!
- Providing multiple types of tests
noisy function, unrelaxable constraints, randomized initial point...
- Implemented in Python and MATLAB

One more thing: OptiProfiler

Just one line MATLAB code:

```
benchmark({@bds, @fminsearch}, "noisy")
```



GitHub repository: <https://github.com/optiprofiler>

Acknowledgement

- Thanks to the organizers!
- Thanks to all the speakers!
- Thanks Giampaolo and Geovani for saving my life!
- Thanks Zaikun for giving me this great opportunity!



Grazie mille!

Photo taken by Lindon.

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