Non-convergence of Probabilistic direct search

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DERIVATIVE-FREE OPTIMIZATION

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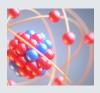
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Example (Typical case: f is a blackbox)

$$x \in \mathbb{R}^n \longrightarrow f: \mathbb{R}^n \to \mathbb{R}$$

Applications



Nuclear Physics



Machine Learning



Cosmology

BASIC ASSUMPTIONS

In this talk, we solve the unconstrained problem

$$\min_{x \in \mathbb{R}^n} f(x),$$

where we follow standard setup that

- ∇f is Lipschitz continuous with constant ν , cannot be evaluated,
- f is bounded below,
- the evaluation of f is expensive.

DIRECT-SEARCH METHODS

Direct-search methods

Sample points based on a finite direction set and make decisions by comparing values.

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$$D = \{e_1, -e_1, e_2, -e_2\} \qquad -e_1 \xrightarrow{e_2} x_k \xrightarrow{e_1} e_1$$

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Example (typical direction set in \mathbb{R}^2)

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In this talk, we assume direction set always be a set of unit vectors in \mathbb{R}^n .

Algorithm 1: Direct Search with sufficient decrease

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Everything almost clear, except "select D_k ".

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Positive Spanning Set (PSS)

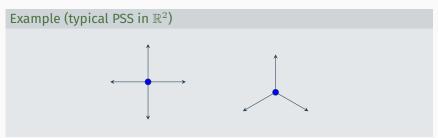
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Positive Spanning Set (PSS)



Question 3: how "good" is a PSS?

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Cosine measure

Define the cosine measure, $\operatorname{cm}(D)$, for a finite set of nonzero vectors, $D \in \mathbb{R}^n$, as

$$cm(D) = \min_{v \in \mathbb{R}^n \setminus \{0\}} cm(D, v) = \min_{v \in \mathbb{R}^n \setminus \{0\}} \max_{d \in D} \frac{d^{\top} v}{\|d\| \|v\|}.$$

Example

$$cm(D) = \cos\theta = \frac{\sqrt{2}}{2} \qquad \longleftrightarrow$$

CONVERGENCE OF DETERMINISTIC DIRECT SEARCH

Theorem

Under standard assumptions, if in Algorithm 1, there exists a positive constant κ such that

$$cm(D_k) \geq \kappa$$
 for all $k \geq 0$,

then we have

$$\liminf_{k \to \infty} \|g_k\| = 0.$$

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Each direction set D_k should be a "good enough" PSS.

Bad news: if D is a PSS in \mathbb{R}^n , then $|D| \ge n+1$. Complexity: $\mathcal{O}(n^2 \varepsilon^{-2})$

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All the randomness comes from \mathfrak{D}_k

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Notations for random variables or random vectors

$$D_k \Rightarrow \mathfrak{D}_k, \quad d \Rightarrow \mathfrak{d}, \quad x_k \Rightarrow X_k, \quad \alpha_k \Rightarrow A_k, \quad g_k \Rightarrow G_k$$

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Definition (p-probabilistically κ -descent)

 $(\mathfrak{D}_k)_{k\geq 0}$ is said to be $p ext{-probabilistically }\kappa ext{-descent, if}$

$$\mathbb{P}\left(\operatorname{cm}(\mathfrak{D}_k, -G_k) \geq \kappa \mid \mathfrak{D}_0, \dots, \mathfrak{D}_{k-1}\right) \geq p \quad \text{for each } k \geq 0.$$

Intuition: high probability that \mathfrak{D}_k is "always good enough", no matter what happened before.

CONVERGENCE OF PROBABILISTIC DIRECT SEARCH

Theorem (Gratton et al. 2015)

Under standard assumptions, if in Algorithm 2, there exists a positive constant κ such that $(\mathfrak{D}_k)_{k\geq 0}$ is p-probabilistically κ -descent with $p=\log \theta/\log(\gamma^{-1}\theta)$, then we have

$$\mathbb{P}\left(\liminf_{k\to\infty}\|G_k\|=0\right) = 1.$$

Complexity: $\mathcal{O}(mn\varepsilon^{-2})$ with extremely high probability, where $m \ll n$.

TABLE OF CONTENTS

- 1. Why and how of "non-convergence": ideas and main results
- 2. Tightness of non-convergence: counterexample
- 3. Conclusion

WHY AND HOW OF "NON-CONVERGENCE":

IDEAS AND MAIN RESULTS

NATURAL QUESTIONS

Recall "Converge or Diverge? A Story of Adam" by Prof. Luo

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Recall "Converge or Diverge? A Story of Adam" by Prof. Luo

What will happen if $(\mathfrak{D}_k)_{k\geq 0}$ fails to meet p-probabilistically κ -descent? Or worse \mathfrak{D}_k is always "bad"?

Can we obtain some kind of "non-convergence"?

"BAD" DIRECTION SET

p-probabilistically ascent

 $(\mathfrak{D}_k)_{k\in\mathbb{N}}$ is said to be p-probabilistically ascent if it satisfies

$$\mathbb{P}\left(\operatorname{cm}\left(\mathfrak{D}_{k},-G_{k}\right)\leq0\mid\mathfrak{D}_{0},\ldots,\mathfrak{D}_{k-1}\right)\geq\,p\quad\text{for each }k\geq0.$$

Observation

For any convex function, if $\operatorname{cm}(D_k, -g_k) \leq 0$, then no descent direction in D_k , leading to shrinking of step size.

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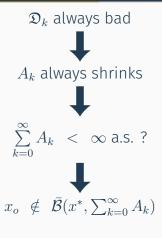
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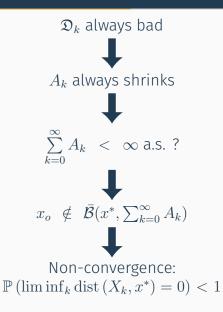
If we show non-convergence for convex functions, then we cannot expect convergence for others.

 \mathfrak{D}_k always bad

 \mathfrak{D}_k always bad $lackbox{4}{}$ A_k always shrinks







• Define indicator function $Y_k = \mathbb{1}_{\{\operatorname{cm}(\mathfrak{D}_k, -G_k) > 0\}}$ Indicator for "good" event

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- $A_k \leq \alpha_0 \prod_{i=0}^{k-1} \gamma^{Y_i} \theta^{1-Y_i} =: \alpha_0 U_k$
- $\sum_{k=1}^{\infty} A_k \le \alpha_0 \sum_{k=1}^{\infty} U_k$

QUANTITATIVE ANALYSIS

Assumption

$$\mathbb{P}\left(Y_k=0\mid Y_0,\dots,Y_{k-1}\right)\ \geq\ p\ >\ p_0\quad\text{for each }k\geq 0,$$
 where $p_0=\log\gamma/\log(\theta^{-1}\gamma).$

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Result

1.

$$\mathbb{P}\left(\sum_{k=1}^{\infty} U_k < \infty\right) = 1$$

2.

$$\mathbb{P}\left(\sum_{k=1}^{\infty} U_k < \zeta\right) > 0 \iff \zeta > \frac{\theta}{1-\theta}$$

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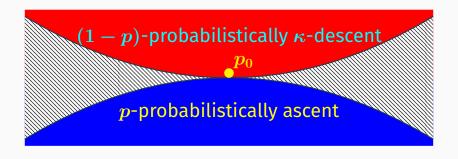
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Observation

- Easy to achieve non-convergence: $\theta/(1-\theta)$ is fixed number.
- Cannot further improve: $\sum_{k=1}^{\infty} U_k \ge \theta/(1-\theta)$.

GRAPHICAL EXPLANATION

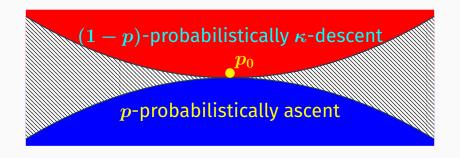


Space of $(\mathfrak{D}_k)_{k\geq 0}$

 \mathfrak{D}_k is measurable map from Ω to $2^{\mathbb{R}^n}$

Red: Convergence Blue: Non-convergence

GRAPHICAL EXPLANATION



Space of $(\mathfrak{D}_k)_{k\geq 0}$

 \mathfrak{D}_k is measurable map from Ω to $2^{\mathbb{R}^n}$

Red: Convergence Blue: Non-convergence Too much "shadow area", not satisfied!

QUALITATIVE ANALYSIS

Question: what if we only need $\sum_{k=1}^{\infty} U_k < \infty$?

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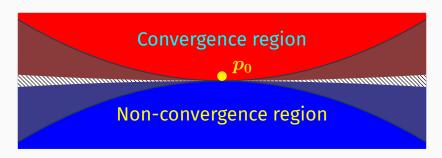
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GRAPHICAL EXPLANATION

Also use "liminf"-type assumption in convergence theory

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TIGHTNESS OF NON-CONVERGENCE: COUNTEREXAMPLE

CAN p_0 BE INCLUDED?

Natural question: can p_0 be included in the non-convergence region?

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Example

We assume

- $f:\mathbb{R}^n \to \mathbb{R}$ be gradient Lipschitz and strongly convex,
- $\theta = 1/2$ and $\gamma = 2$,
- $\cdot \ \mathbb{P}(\mathfrak{d}_k = G_k/\|G_k\|) = \mathbb{P}(\mathfrak{d}_k = -G_k/\|G_k\|) = 1/2$,

then we have

$$\mathbb{P}\left(\lim_{k\to\infty}\|G_k\|=0\right) = 1.$$

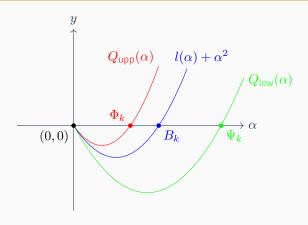
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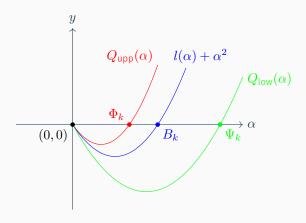
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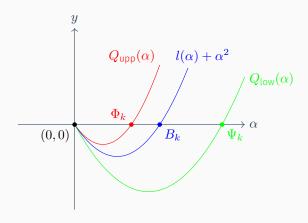
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- $l(\alpha) + \alpha^2$ is both upper and lower bounded by quadratic functions

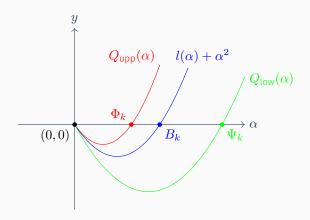




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- · Known that $A_k(\omega) \to 0 \ \forall \omega \in \Omega$, just to prove $\mathbb{P}(A_k \geq B_k \text{ i.o.}) = 1$

ANALYSIS OF A RANDOM WALK

Let $S_k = \log_2(A_k/\alpha_0)$ and $R_k = \log_2(B_k/\alpha_0)$. Then we abstract the following problem.

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Assume that $(S_k)_{k\geq 0}$ and $(R_k)_{k\geq 0}$ are two stochastic processes such that

•
$$S_0=0$$
 and $(R_k)_{k\geq 0}$ uniformly bounded,

•

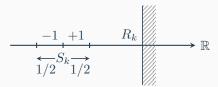
$$\mathbb{P}(S_{k+1} = S_k - 1 \mid S_1, \dots, S_k) = \frac{1}{2} \mathbb{1}_{\{S_k < R_k\}} + \mathbb{1}_{\{S_k \ge R_k\}},$$

$$\mathbb{P}(S_{k+1} = S_k + 1 \mid S_1, \dots, S_k) = \frac{1}{2} \mathbb{1}_{\{S_k < R_k\}}.$$

Then, $\mathbb{P}(S_k \geq R_k \text{ i.o.}) = 1$?

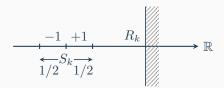
GRAPHICAL EXPLANATION

When $S_k < R_k$,

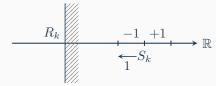


GRAPHICAL EXPLANATION

When $S_k < R_k$,

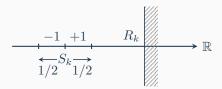


When $S_k \geq R_k$,

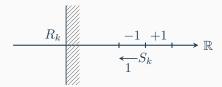


GRAPHICAL EXPLANATION

When $S_k < R_k$,



When $S_k \geq R_k$,



Does S_k go beyond the "wall" R_k i.o. with probability 1?

• Construct $\widetilde{S}_k=S_k+\sum_{i=0}^{k-1}T_i\mathbb{1}_{\{S_k\geq R_k\}}$, where T_0,T_1,\ldots are i.i.d., $\mathbb{P}(T_i=0)=\mathbb{P}(T_i=2)=1/2$

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- ${f \cdot}\ (\widetilde{S}_k)_{k\geq 0}$ is simple random walk $\Rightarrow \limsup_k \widetilde{S}_k = \infty$ w.p. 1

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- $(R_k)_{k\geq 0}$ uniformly bounded \Rightarrow so does $(S_k)_{k\geq 0}$
- $\sum_{i=0}^{\infty} T_i \mathbb{1}_{\{S_k \geq R_k\}} = \infty$ w.p. 1

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- establish non-convergence theory for probabilistic direct search,
- · distinguish convergence region and non-convergence region,
- construct one counterexample to show the tightness of non-convergence region boundary.

Future work:

- find estimation or lower bound for the probability of non-convergence,
- establish non-convergence theory for other models.

Thank you!

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