

Non-convergence Analysis of Probabilistic Direct Search

The 25th International Symposium on Mathematical Programming

Cunxin Huang

Joint work with Zaikun Zhang

Montréal, Canada July 26, 2024

The Hong Kong Polytechnic University

Non-convergence Analysis of Probabilistic Direct Search

The 25th International Symposium on Mathematical Programming

Zaikun Zhang (replacing Cunxin Huang)

Joint work with Cunxin Huang

Montréal, Canada July 26, 2024

The Hong Kong Polytechnic University

PRIMA and my gratitude



libprima.net

PRIMA is an acronym for

“Reference Implementation for Powell’s Methods
with Modernization and Amelioration”.

- Number of lines: > 100,000.
- The total time I spent on PRIMA:
 $\geq 3 \text{ years} \times 300 \text{ days per year} \times 10 \text{ hours per day} = 9,000 \text{ hours.}$

In the past years, due to the gap in my publication record while working on PRIMA, I needed a lot of support from the community. Thank you for the help and support, explicit or implicit, known or unknown to me.
Without your support, I would not have survived.

Non-convergence analysis: What?

Consider an algorithm

$$\mathcal{A} : \Xi \times \mathbb{F} \times \mathcal{X} \rightarrow \mathcal{X}^\infty, \quad (\xi, f, x_0) \mapsto \{x_k\}.$$

- ξ represents algorithmic parameters.
- f is the objective function.
- x_0 is the starting point.

When \mathcal{A} is deterministic:

- (Global) Convergence analysis: For all $(\xi, f, x_0) \in \hat{\Xi} \times \hat{\mathbb{F}} \times \mathcal{X}$, prove
 $\{x_k\}$ achieves stationarity asymptotically.
- Non-convergence analysis: For all $(\xi, f, x_0) \in \tilde{\Xi} \times \tilde{\mathbb{F}} \times \tilde{\mathcal{X}}$, prove
 $\{x_k\}$ fails to achieve stationarity asymptotically.

Non-convergence analysis: What?

Consider an algorithm

$$\mathcal{A} : \Xi \times \mathbb{F} \times \mathcal{X} \rightarrow \mathcal{X}^\infty, \quad (\xi, f, x_0) \mapsto \{x_k\}.$$

- ξ represents algorithmic parameters.
- f is the objective function.
- x_0 is the starting point.

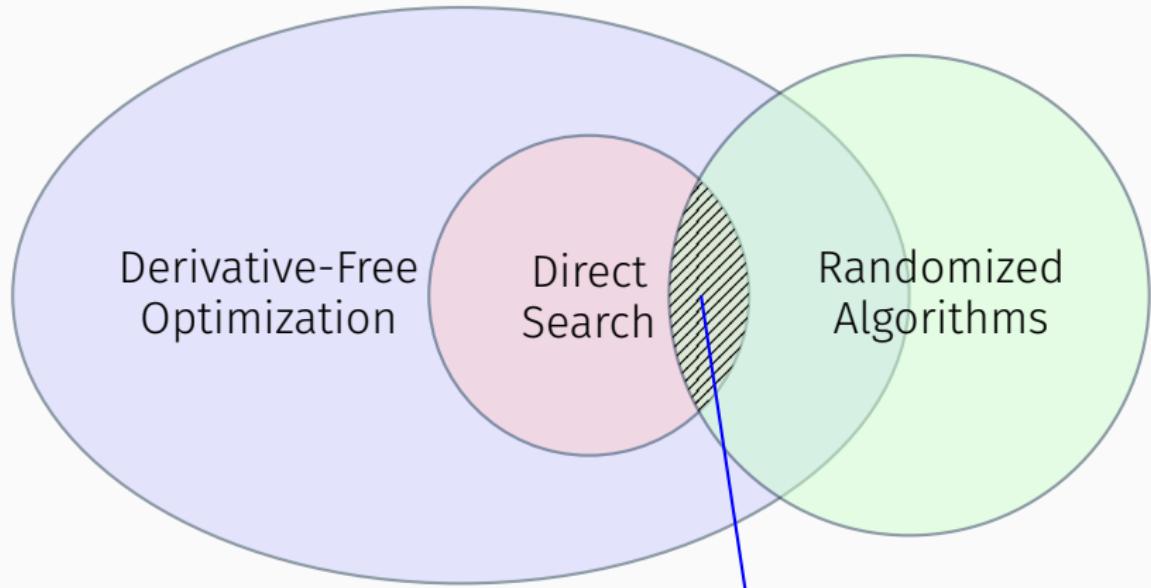
When \mathcal{A} is random:

- (Global) Convergence analysis: For all $(\xi, f, x_0) \in \hat{\Xi} \times \hat{\mathbb{F}} \times \mathcal{X}$, prove
 $\mathbb{P}(\{x_k\} \text{ achieves stationarity asymptotically}) = 1$.
- Non-convergence analysis: For all $(\xi, f, x_0) \in \tilde{\Xi} \times \tilde{\mathbb{F}} \times \tilde{\mathcal{X}}$, prove
 $\mathbb{P}(\{x_k\} \text{ fails to achieve stationarity asymptotically}) > 0$.

Non-convergence analysis: Why?

- Sharpen our knowledge about the algorithm.
- Deepen our understanding about the convergence analysis.
- Guide the selection of algorithmic parameters.
- Provide new perspectives on convergence analysis.

Probabilistic Direct Search (PDS)



The algorithm we consider in this talk:
Probabilistic Direct Search
(Gratton, Royer, Vicente, and Zhang 2015)

Derivative-Free Optimization (DFO)

Derivative-Free Optimization

- Do not use derivatives (first-order info.), only use function values
- Closely related: zeroth-order/black-box optimization ...

Derivatives are often **not available in applications**



Quantum Computing



Machine Learning

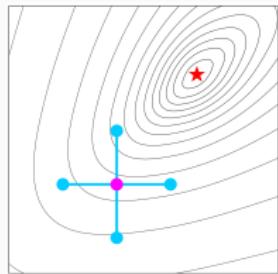


Circuit Design

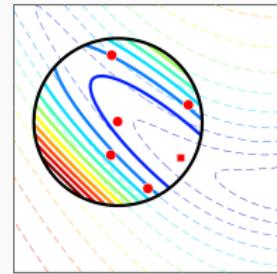
Direct-search methods and model-based methods

How to determine iterates?

- Direct-search methods: “simple” comparison of function values
- Model-based methods: build a surrogate of the objective function



Direct-search methods¹



Model-based methods²

¹Source: Kolda, Lewis, and Torczon 2003

²Source: Larson, Menickelly, and Wild 2019

PDS: a simplified framework

Algorithm 1: Direct Search based on **sufficient decrease**

PDS: a simplified framework

Algorithm 1: Direct Search based on **sufficient decrease**

Input: $x_0 \in \mathbb{R}^n$, $\alpha_0 \in (0, \infty)$, $c \in (0, \infty)$, $0 < \theta < 1 < \gamma$.

PDS: a simplified framework

Algorithm 1: Direct Search based on sufficient decrease

Input: $x_0 \in \mathbb{R}^n$, $\alpha_0 \in (0, \infty)$, $c \in (0, \infty)$, $0 < \theta < 1 < \gamma$.

for $k = 0, 1, \dots$ **do**

PDS: a simplified framework

Algorithm 1: Direct Search based on sufficient decrease

Input: $x_0 \in \mathbb{R}^n$, $\alpha_0 \in (0, \infty)$, $c \in (0, \infty)$, $0 < \theta < 1 < \gamma$.

for $k = 0, 1, \dots$ **do**

Select a finite set of directions $\mathcal{D}_k \subset \mathbb{R}^n$.

PDS: a simplified framework

Algorithm 1: Direct Search based on sufficient decrease

Input: $x_0 \in \mathbb{R}^n$, $\alpha_0 \in (0, \infty)$, $c \in (0, \infty)$, $0 < \theta < 1 < \gamma$.

for $k = 0, 1, \dots$ **do**

Select a finite set of directions $\mathcal{D}_k \subset \mathbb{R}^n$.

(In this talk, assume \mathcal{D}_k is a set of unit vectors for simplicity)

PDS: a simplified framework

Algorithm 1: Direct Search based on sufficient decrease

Input: $x_0 \in \mathbb{R}^n$, $\alpha_0 \in (0, \infty)$, $c \in (0, \infty)$, $0 < \theta < 1 < \gamma$.

for $k = 0, 1, \dots$ **do**

Select a finite set of directions $\mathcal{D}_k \subset \mathbb{R}^n$.

(In this talk, assume \mathcal{D}_k is a set of unit vectors for simplicity)

Set $d_k = \arg \min\{f(x_k + \alpha_k d) : d \in \mathcal{D}_k\}$.

(complete polling for simplicity)

PDS: a simplified framework

Algorithm 1: Direct Search based on sufficient decrease

Input: $x_0 \in \mathbb{R}^n$, $\alpha_0 \in (0, \infty)$, $c \in (0, \infty)$, $0 < \theta < 1 < \gamma$.

for $k = 0, 1, \dots$ **do**

Select a finite set of directions $\mathcal{D}_k \subset \mathbb{R}^n$.

(In this talk, assume \mathcal{D}_k is a set of unit vectors for simplicity)

Set $d_k = \arg \min\{f(x_k + \alpha_k d) : d \in \mathcal{D}_k\}$.

(complete polling for simplicity)

if $f(x_k + \alpha_k d_k) < f(x_k) - c\alpha_k^2$ **then**

PDS: a simplified framework

Algorithm 1: Direct Search based on sufficient decrease

Input: $x_0 \in \mathbb{R}^n$, $\alpha_0 \in (0, \infty)$, $c \in (0, \infty)$, $0 < \theta < 1 < \gamma$.

for $k = 0, 1, \dots$ **do**

 Select a finite set of directions $\mathcal{D}_k \subset \mathbb{R}^n$.

 (In this talk, assume \mathcal{D}_k is a set of unit vectors for simplicity)

 Set $d_k = \arg \min\{f(x_k + \alpha_k d) : d \in \mathcal{D}_k\}$.

 (complete polling for simplicity)

if $f(x_k + \alpha_k d_k) < f(x_k) - c\alpha_k^2$ **then**

 Set $x_{k+1} = x_k + \alpha_k d_k$ and $\alpha_{k+1} = \gamma \alpha_k$.

 (Move and expand step size)

else

PDS: a simplified framework

Algorithm 1: Direct Search based on sufficient decrease

Input: $x_0 \in \mathbb{R}^n$, $\alpha_0 \in (0, \infty)$, $c \in (0, \infty)$, $0 < \theta < 1 < \gamma$.

for $k = 0, 1, \dots$ **do**

 Select a finite set of directions $\mathcal{D}_k \subset \mathbb{R}^n$.

 (In this talk, assume \mathcal{D}_k is a set of unit vectors for simplicity)

 Set $d_k = \arg \min\{f(x_k + \alpha_k d) : d \in \mathcal{D}_k\}$.

 (complete polling for simplicity)

if $f(x_k + \alpha_k d_k) < f(x_k) - c\alpha_k^2$ **then**

 Set $x_{k+1} = x_k + \alpha_k d_k$ and $\alpha_{k+1} = \gamma \alpha_k$.

 (Move and expand step size)

else

 Set $x_{k+1} = x_k$ and $\alpha_{k+1} = \theta \alpha_k$.

 (Stay and shrink step size)

PDS: a simplified framework

Algorithm 1: Probabilistic Direct Search based on sufficient decrease

Input: $x_0 \in \mathbb{R}^n$, $\alpha_0 \in (0, \infty)$, $c \in (0, \infty)$, $0 < \theta < 1 < \gamma$.

for $k = 0, 1, \dots$ **do**

Select a finite set of directions $\mathcal{D}_k \subset \mathbb{R}^n$ **randomly**.

(In this talk, assume \mathcal{D}_k is a set of unit vectors for simplicity)

Set $d_k = \arg \min\{f(x_k + \alpha_k d) : d \in \mathcal{D}_k\}$.

(complete polling for simplicity)

if $f(x_k + \alpha_k d_k) < f(x_k) - c\alpha_k^2$ **then**

Set $x_{k+1} = x_k + \alpha_k d_k$ and $\alpha_{k+1} = \gamma \alpha_k$.

(Move and expand step size)

else

Set $x_{k+1} = x_k$ and $\alpha_{k+1} = \theta \alpha_k$.

(Stay and shrink step size)

PDS: a simplified framework

Algorithm 1: Probabilistic Direct Search based on sufficient decrease

Input: $x_0 \in \mathbb{R}^n$, $\alpha_0 \in (0, \infty)$, $c \in (0, \infty)$, $0 < \theta < 1 < \gamma$.

for $k = 0, 1, \dots$ **do**

Select a finite set of directions $\mathcal{D}_k \subset \mathbb{R}^n$ **randomly**.

(In this talk, assume \mathcal{D}_k is a set of unit vectors for simplicity)

Set $d_k = \arg \min\{f(x_k + \alpha_k d) : d \in \mathcal{D}_k\}$.

(complete polling for simplicity)

if $f(x_k + \alpha_k d_k) < f(x_k) - c\alpha_k^2$ **then**

Set $x_{k+1} = x_k + \alpha_k d_k$ and $\alpha_{k+1} = \gamma \alpha_k$.

(Move and expand step size)

else

Set $x_{k+1} = x_k$ and $\alpha_{k+1} = \theta \alpha_k$.

(Stay and shrink step size)

Typical choice of $\{\mathcal{D}_k\}$ (GRVZ 2015): $\mathcal{D}_k = \{d_1, \dots, d_m\}$ with $d_\ell \stackrel{\text{i.i.d.}}{\sim} \mathcal{U}(S^{n-1})$

PDS: a simplified framework

Algorithm 1: Probabilistic Direct Search based on sufficient decrease

Input: $x_0 \in \mathbb{R}^n$, $\alpha_0 \in (0, \infty)$, $c \in (0, \infty)$, $0 < \theta < 1 < \gamma$.

for $k = 0, 1, \dots$ **do**

Select a finite set of directions $\mathcal{D}_k \subset \mathbb{R}^n$ **randomly**.

(In this talk, assume \mathcal{D}_k is a set of unit vectors for simplicity)

Set $d_k = \arg \min\{f(x_k + \alpha_k d) : d \in \mathcal{D}_k\}$.

(complete polling for simplicity)

if $f(x_k + \alpha_k d_k) < f(x_k) - c\alpha_k^2$ **then**

Set $x_{k+1} = x_k + \alpha_k d_k$ and $\alpha_{k+1} = \gamma \alpha_k$.

(Move and expand step size)

else

Set $x_{k+1} = x_k$ and $\alpha_{k+1} = \theta \alpha_k$.

(Stay and shrink step size)

Typical choice of $\{\mathcal{D}_k\}$ (GRVZ 2015): $\mathcal{D}_k = \{d_1, \dots, d_m\}$ with $d_\ell \stackrel{\text{i.i.d.}}{\sim} U(\mathcal{S}^{n-1})$

N.B.: typical choice in the deterministic case is $\{\pm e_i\}_{i=1}^n$, Coordinate Search (CS)

Illustration of how PDS works

$\mathcal{D}_k = \{d_1, d_2\}$, where $d_\ell \stackrel{\text{i.i.d.}}{\sim} \mathbf{U}(\mathcal{S}^1)$

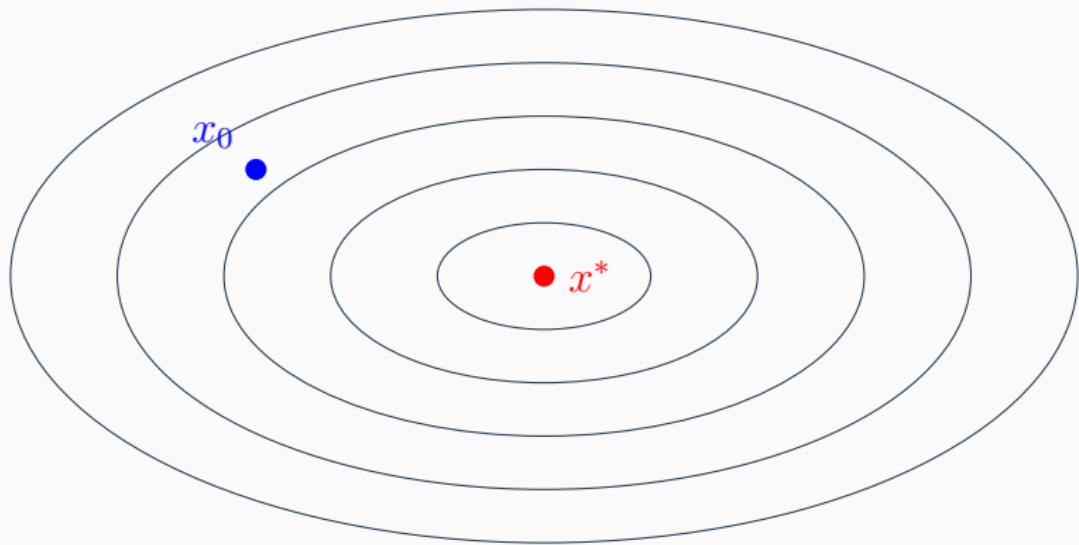


Illustration of how PDS works

$\mathcal{D}_k = \{d_1, d_2\}$, where $d_\ell \stackrel{\text{i.i.d.}}{\sim} \mathbf{U}(\mathcal{S}^1)$

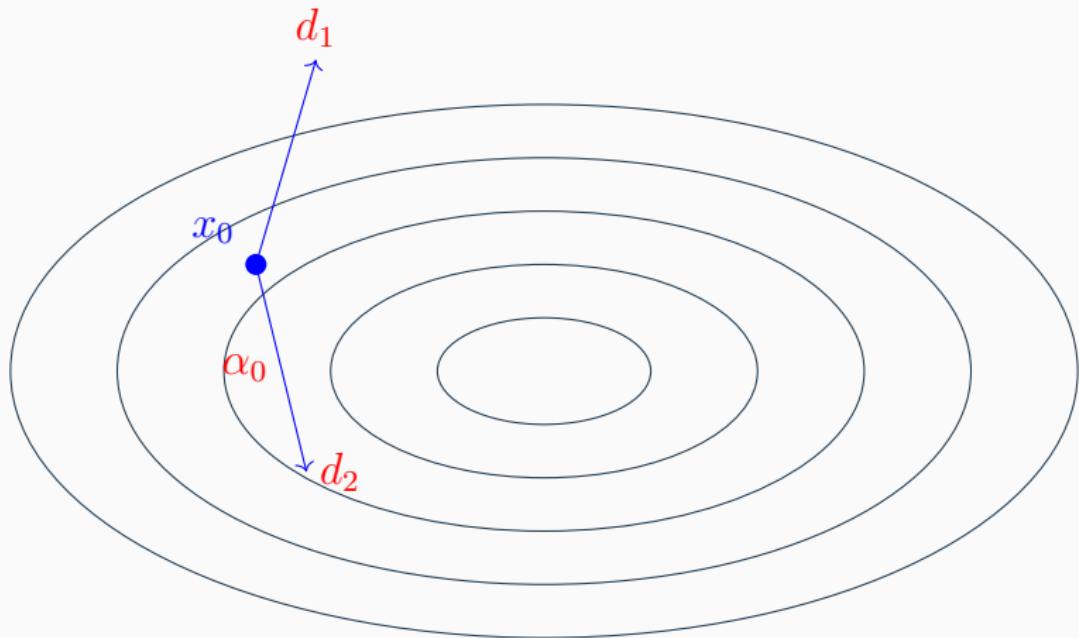


Illustration of how PDS works

$\mathcal{D}_k = \{d_1, d_2\}$, where $d_\ell \stackrel{\text{i.i.d.}}{\sim} \mathbf{U}(\mathcal{S}^1)$

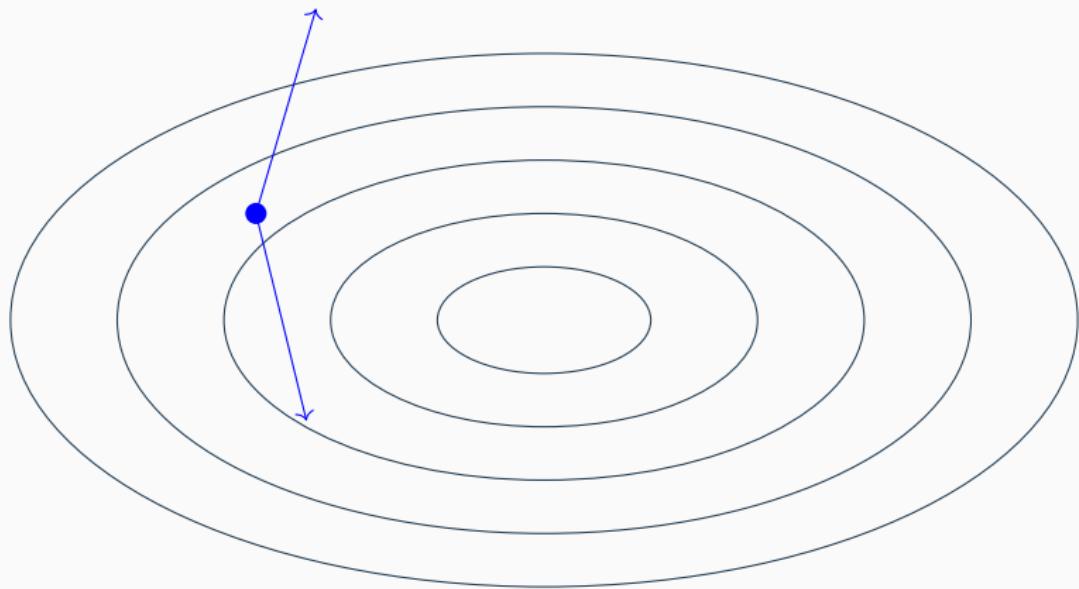


Illustration of how PDS works

$\mathcal{D}_k = \{d_1, d_2\}$, where $d_\ell \stackrel{\text{i.i.d.}}{\sim} \mathbf{U}(\mathcal{S}^1)$

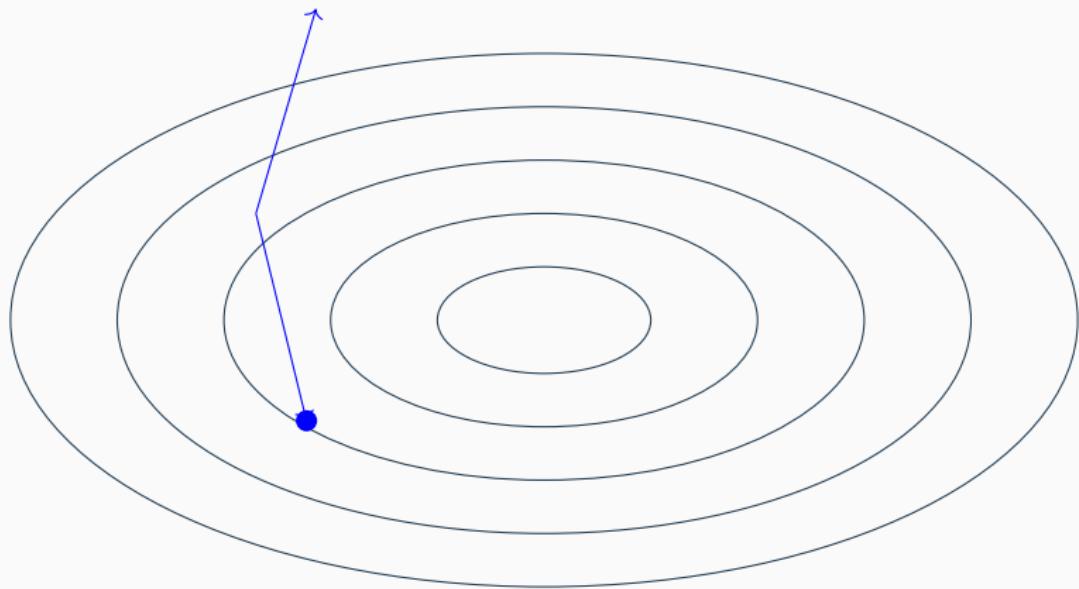


Illustration of how PDS works

$\mathcal{D}_k = \{d_1, d_2\}$, where $d_\ell \stackrel{\text{i.i.d.}}{\sim} \mathbf{U}(\mathcal{S}^1)$

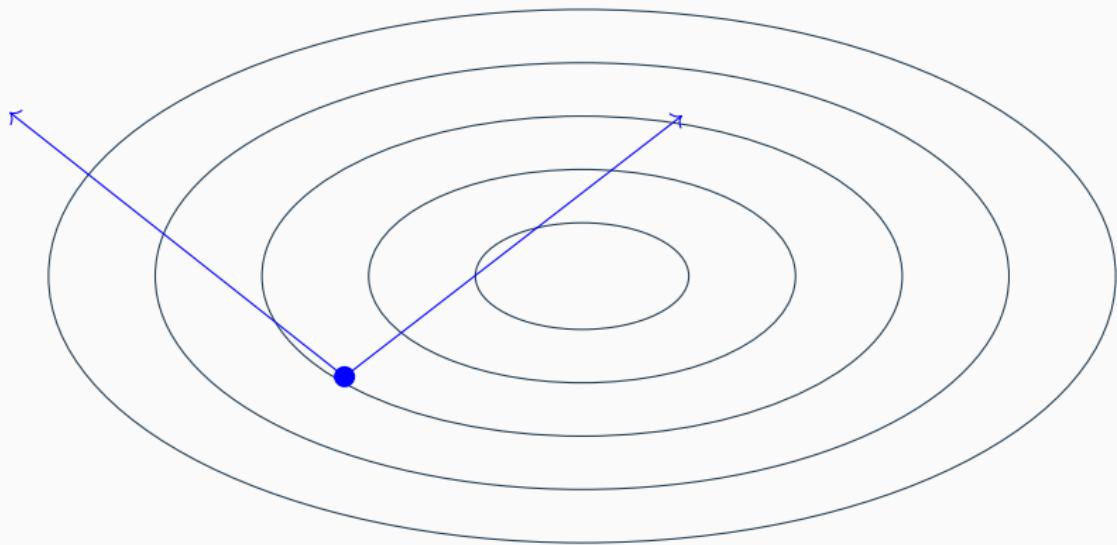


Illustration of how PDS works

$\mathcal{D}_k = \{d_1, d_2\}$, where $d_\ell \stackrel{\text{i.i.d.}}{\sim} \mathbf{U}(\mathcal{S}^1)$

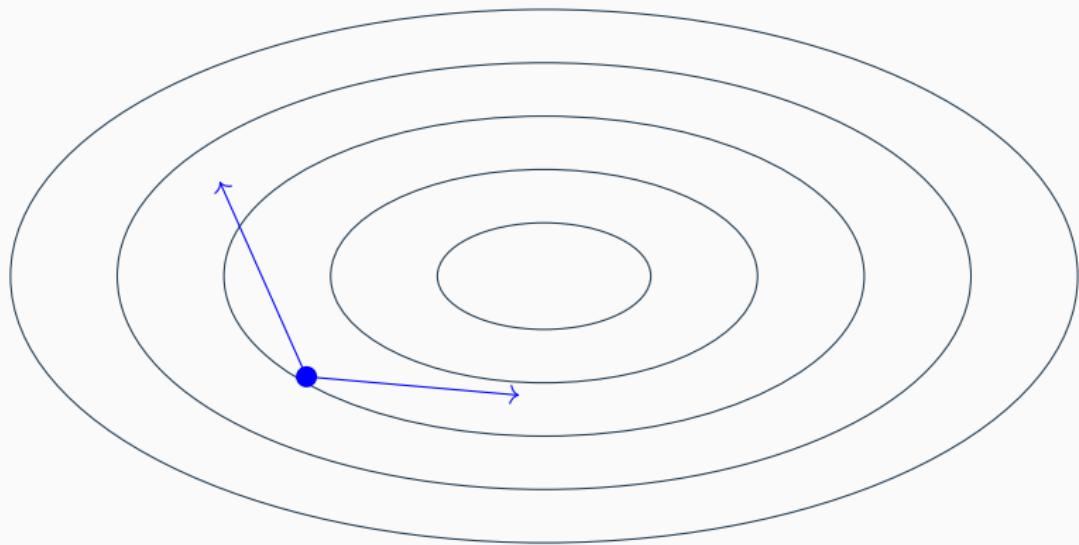


Illustration of how PDS works

$\mathcal{D}_k = \{d_1, d_2\}$, where $d_\ell \stackrel{\text{i.i.d.}}{\sim} \mathbf{U}(\mathcal{S}^1)$

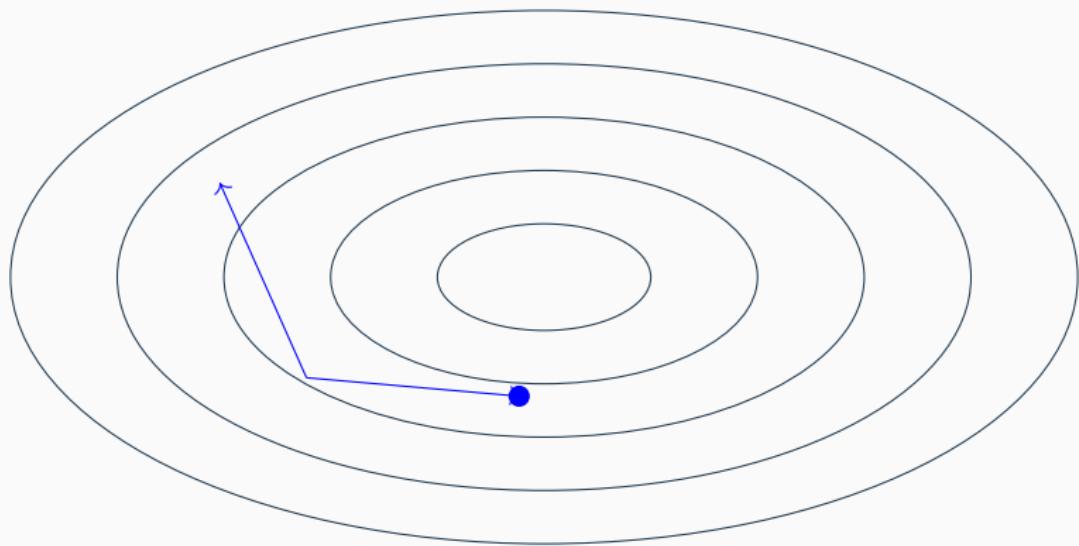


Illustration of how PDS works

$\mathcal{D}_k = \{d_1, d_2\}$, where $d_\ell \stackrel{\text{i.i.d.}}{\sim} \mathbf{U}(\mathcal{S}^1)$

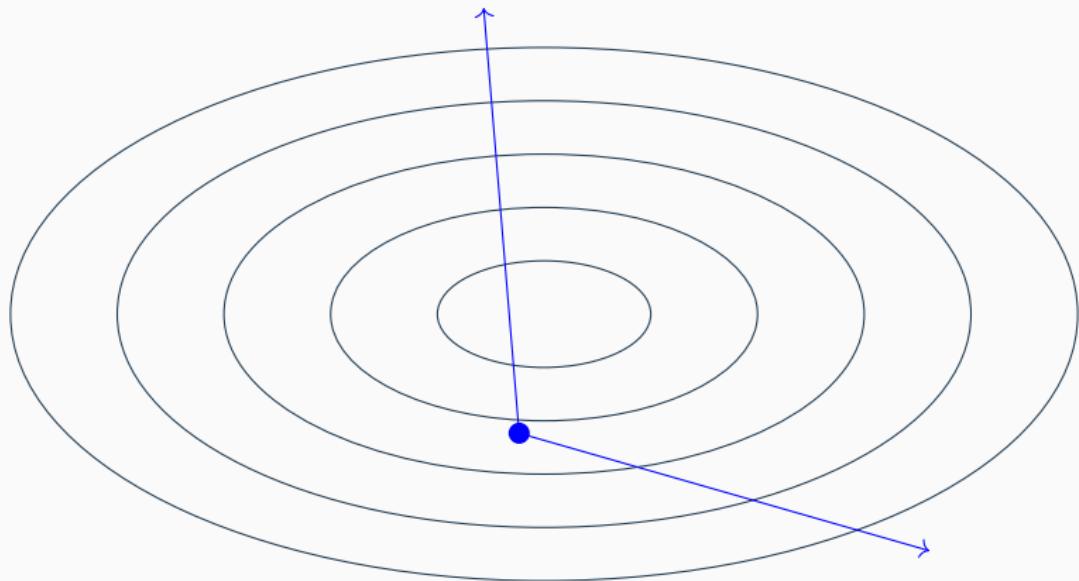


Illustration of how PDS works

$\mathcal{D}_k = \{d_1, d_2\}$, where $d_\ell \stackrel{\text{i.i.d.}}{\sim} \mathbf{U}(\mathcal{S}^1)$

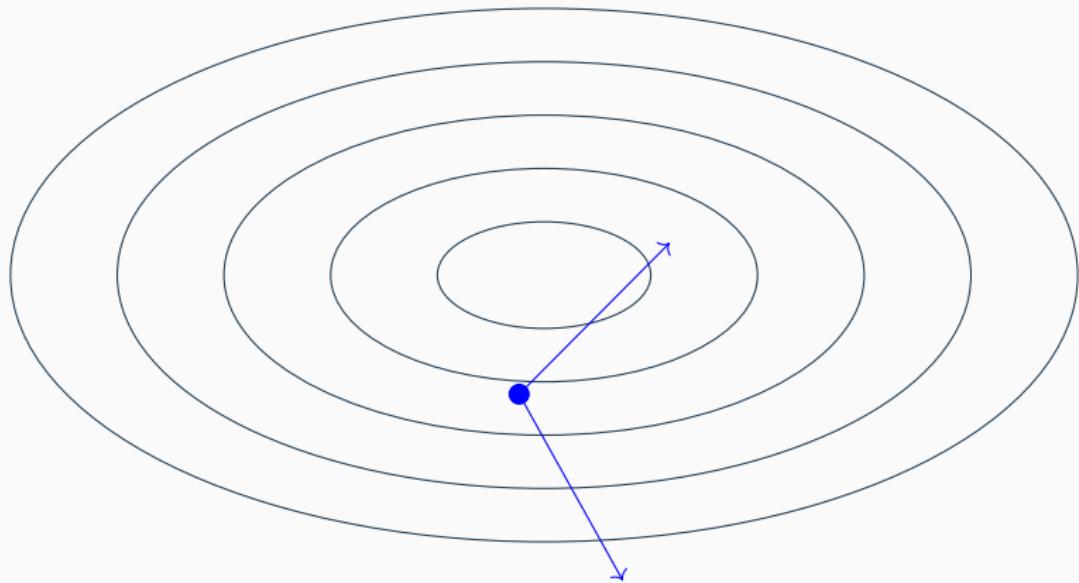


Illustration of how PDS works

$\mathcal{D}_k = \{d_1, d_2\}$, where $d_\ell \stackrel{\text{i.i.d.}}{\sim} \mathbf{U}(\mathcal{S}^1)$

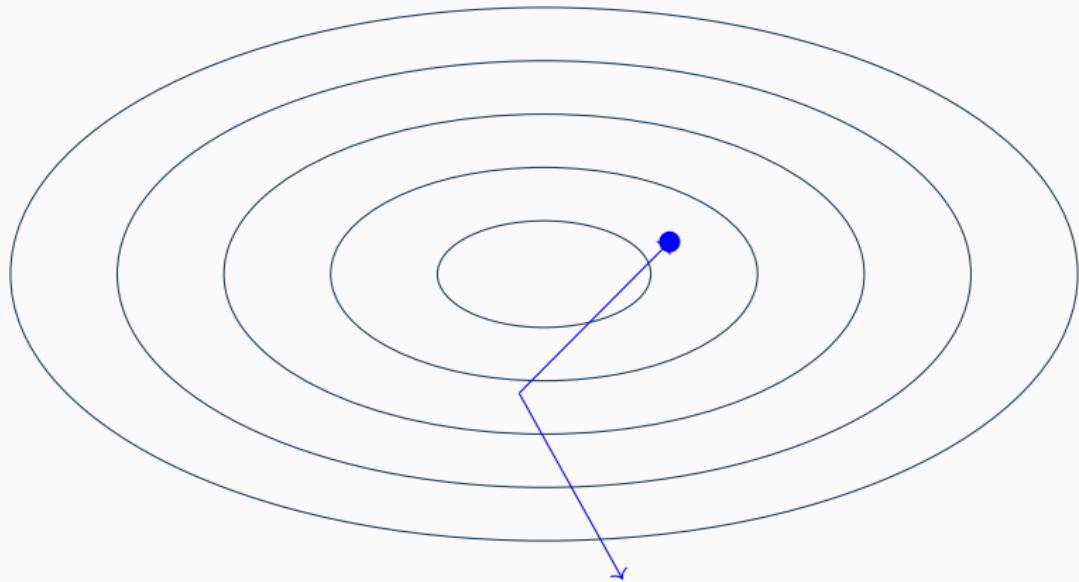


Illustration of how PDS works

$\mathcal{D}_k = \{d_1, d_2\}$, where $d_\ell \stackrel{\text{i.i.d.}}{\sim} \mathbf{U}(\mathcal{S}^1)$

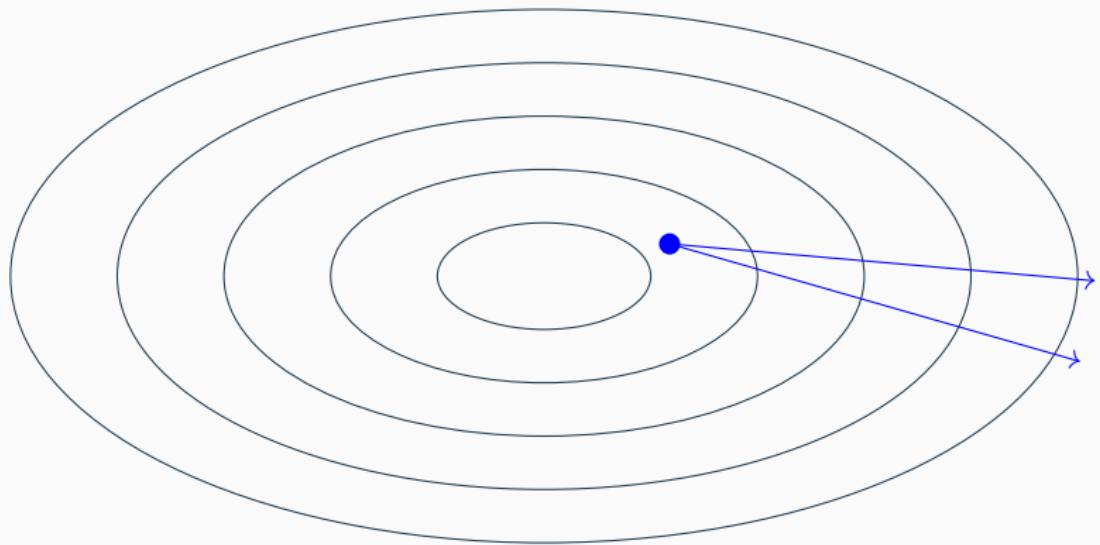


Illustration of how PDS works

$\mathcal{D}_k = \{d_1, d_2\}$, where $d_\ell \stackrel{\text{i.i.d.}}{\sim} \mathbf{U}(\mathcal{S}^1)$

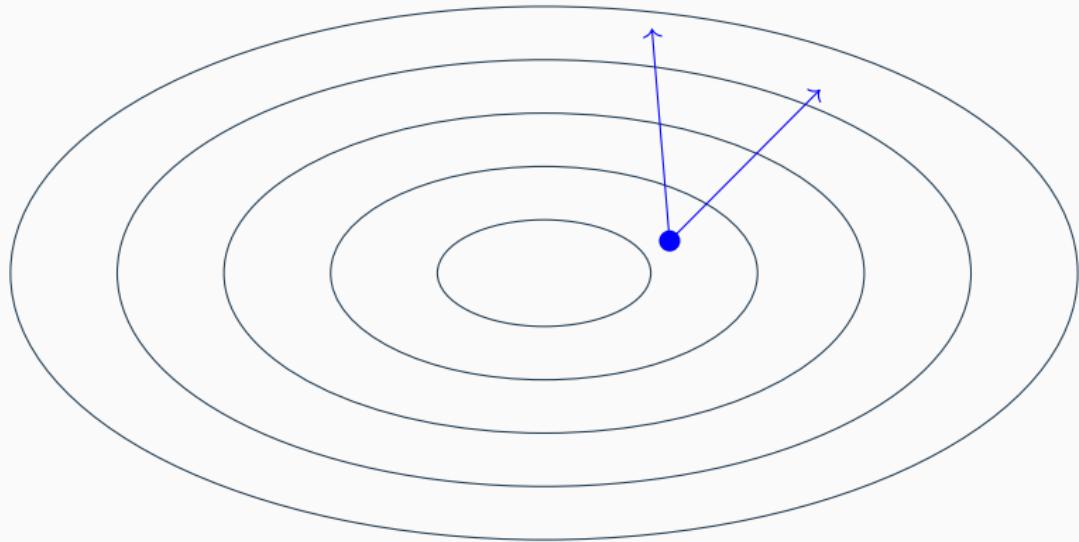
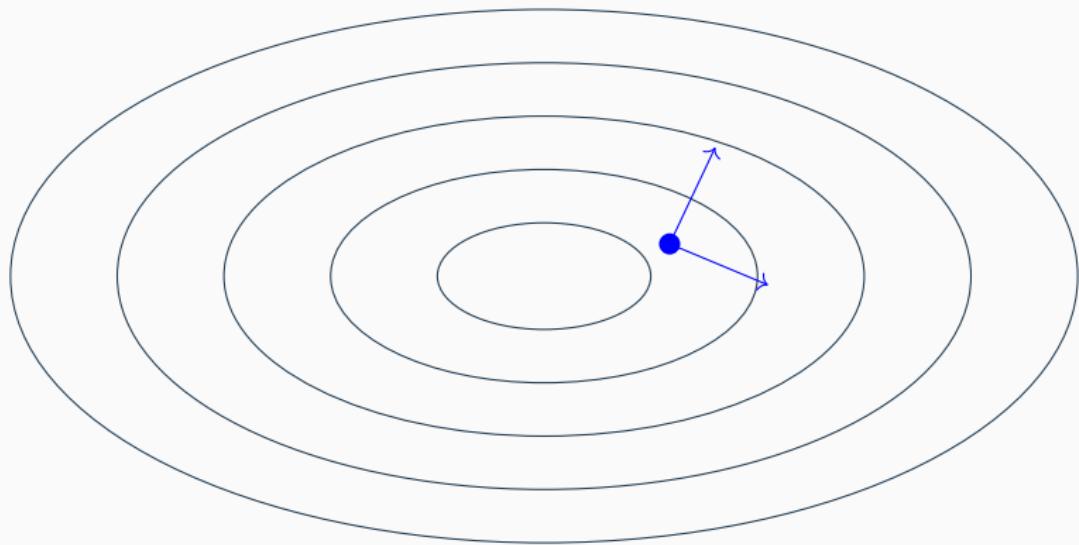


Illustration of how PDS works

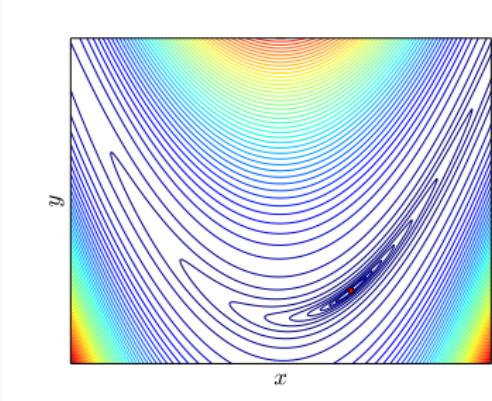
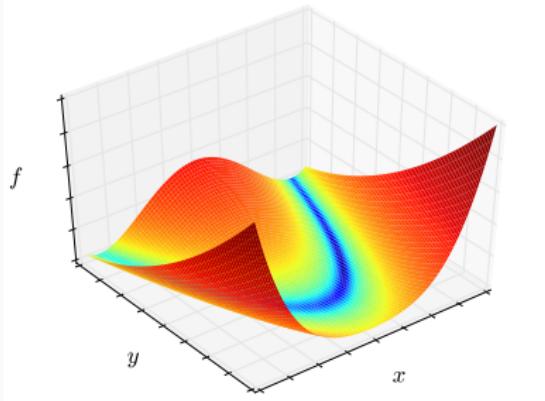
$\mathcal{D}_k = \{d_1, d_2\}$, where $d_\ell \stackrel{\text{i.i.d.}}{\sim} \mathbf{U}(\mathcal{S}^1)$



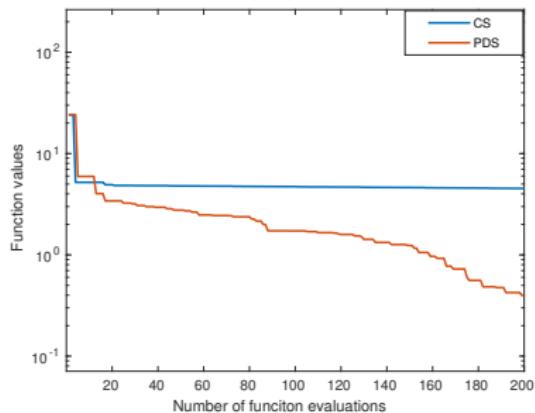
A numerical example: CS v.s. PDS with 2 directions

Rosenbrock “banana” function:

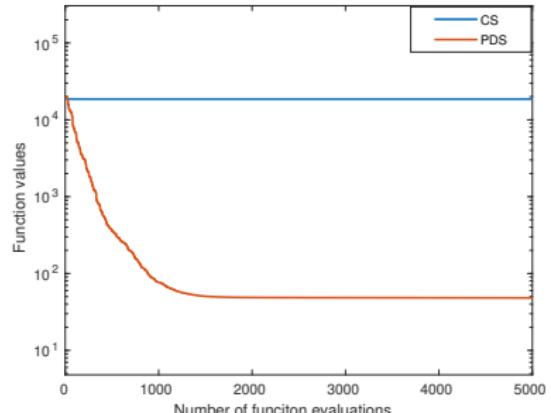
$$f(x) = \sum_{i=1}^{n-1} [(1 - x_i)^2 + 100(x_{i+1} - x_i^2)^2]$$



A numerical example: CS v.s. PDS with 2 directions



$$n = 2$$



$$n = 50$$

Function value v.s. number of function evaluations

Worst case complexity of function evaluations (GRVZ 2015)

$\mathcal{O}(n^2\epsilon^{-2})$ for CS while $\mathcal{O}(n\epsilon^{-2})$ for PDS

Cosine measure

Definition (Cosine measure w.r.t. a vector)

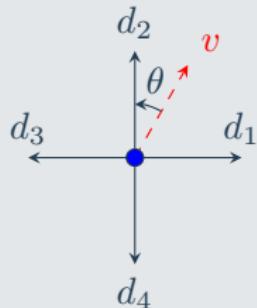
Given a finite set $\mathcal{D} \subseteq \mathbb{R}^n \setminus \{0\}$ and a vector $v \in \mathbb{R}^n \setminus \{0\}$, define

$$\text{cm}(\mathcal{D}, v) = \max_{d \in \mathcal{D}} \frac{d^\top v}{\|d\| \|v\|},$$

which is the cosine measure of \mathcal{D} with respect to v .

Example

$$\text{cm}(\mathcal{D}, v) = \cos \theta$$



Cosine measure

Definition (Cosine measure w.r.t. a vector)

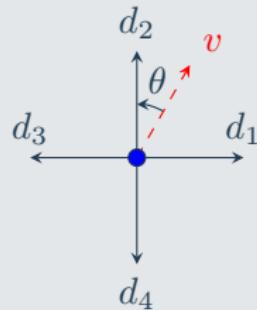
Given a finite set $\mathcal{D} \subseteq \mathbb{R}^n \setminus \{0\}$ and a vector $v \in \mathbb{R}^n \setminus \{0\}$, define

$$\text{cm}(\mathcal{D}, v) = \max_{d \in \mathcal{D}} \frac{d^\top v}{\|d\| \|v\|},$$

which is the cosine measure of \mathcal{D} with respect to v .

Example

$$\text{cm}(\mathcal{D}, v) = \cos \theta$$



$\text{cm}(\mathcal{D}, v)$ measures the ability of \mathcal{D} to “approximate” v

Convergence theory

Definition (p -probabilistically κ -descent)

$\{\mathcal{D}_k\}$ is said to be p -probabilistically κ -descent, if

$$\mathbb{P}(\text{cm}(\mathcal{D}_k, -g_k) \geq \kappa \mid \mathcal{D}_0, \dots, \mathcal{D}_{k-1}) \geq p \quad \text{for each } k \geq 0,$$

where $g_k = \nabla f(x_k)$.

Intuitive meaning of p -probabilistically κ -descent

Each \mathcal{D}_k is “good enough” with probability at least p
no matter what has happened in the history

Convergence theory

Definition (p -probabilistically κ -descent)

$\{\mathcal{D}_k\}$ is said to be p -probabilistically κ -descent, if

$$\mathbb{P}(\text{cm}(\mathcal{D}_k, -g_k) \geq \kappa \mid \mathcal{D}_0, \dots, \mathcal{D}_{k-1}) \geq p \quad \text{for each } k \geq 0,$$

where $g_k = \nabla f(x_k)$.

Intuitive meaning of p -probabilistically κ -descent

Each \mathcal{D}_k is “good enough” with probability at least p
no matter what has happened in the history

Theorem (GRVZ 2015)

If $\{\mathcal{D}_k\}$ is p_0 -probabilistically κ -descent with $\kappa > 0$ and

$$p_0 = \frac{\log \theta}{\log(\gamma^{-1}\theta)},$$

then PDS converges w.p.1 when f is L -smooth and lower-bounded.

Practical choice and natural questions

Corollary (GRVZ 2015)

If $\mathcal{D}_k = \{d_1, \dots, d_m\}$, where $d_\ell \stackrel{i.i.d.}{\sim} U(\mathcal{S}^{n-1})$, then PDS converges w.p.1 if

$$m > \log_2 \left(1 - \frac{\log \theta}{\log \gamma} \right).$$

Practical choice and natural questions

Corollary (GRVZ 2015)

If $\mathcal{D}_k = \{d_1, \dots, d_m\}$, where $d_\ell \stackrel{i.i.d.}{\sim} U(\mathcal{S}^{n-1})$, then PDS converges w.p.1 if

$$m > \log_2 \left(1 - \frac{\log \theta}{\log \gamma} \right).$$

Questions:

- Is p_0 -probabilistically κ -descent an **essential** assumption or a **technical** one? (Such **supermartingale-like assumptions** are ubiquitous in the convergence analysis of randomized methods!)
- What will happen if

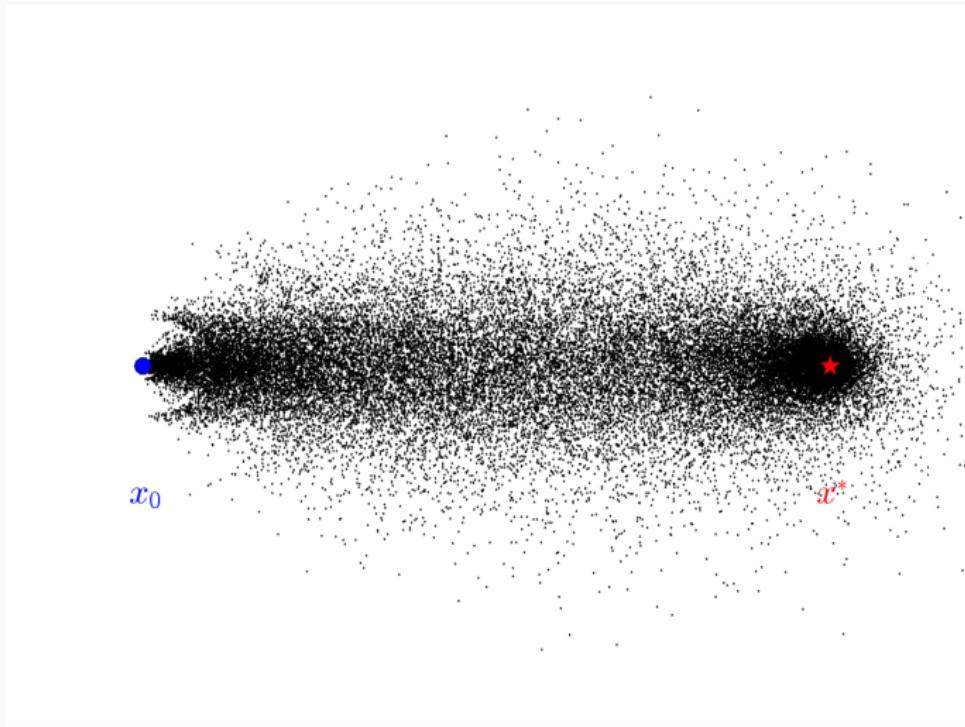
$$m \leq \log_2 \left(1 - \frac{\log \theta}{\log \gamma} \right)?$$

A simple test

- Objective function: $f(x) = x^T x / 2$
- Initial point: $x_0 = (-10, 0)^T$
- Stopping criterion: $\alpha_k \leq \text{machine epsilon}$
- Number of experiments: 100,000
- Parameters of PDS: $\alpha_0 = 1, \theta = 0.25, \gamma = 1.5, m = 2$, which render

$$m = 2 < 2.143 \approx \log_2 \left(1 - \frac{\log \theta}{\log \gamma} \right)$$

A simple test (Cont'd)



Note: each **black dot** represents the **output point** of one run of PDS.

Non-convergence study is not rare

Many well-known algorithms have non-convergence [examples](#)

- Powell, On search directions for minimization algorithms, 1973.
- Yuan, An example of non-convergence of [trust region](#) algorithms, 1998.
- Reddi, Kale, and Kumar, On the convergence of [Adam](#) and beyond, 2018.
- Chen, He, Ye, and Yuan, The direct extension of [ADMM](#) for multi-block convex minimization problems is not necessarily convergent, 2016.
- Dai, A perfect example for the [BFGS](#) method, 2013.
- Mascarenhas, The divergence of the [BFGS](#) and [Gauss Newton](#) methods, 2014.

Non-convergence study is not rare

Many well-known algorithms have non-convergence examples

- Powell, On search directions for minimization algorithms, 1973.
- Yuan, An example of non-convergence of trust region algorithms, 1998.
- Reddi, Kale, and Kumar, On the convergence of Adam and beyond, 2018.
- Chen, He, Ye, and Yuan, The direct extension of ADMM for multi-block convex minimization problems is not necessarily convergent, 2016.
- Dai, A perfect example for the BFGS method, 2013.
- Mascarenhas, The divergence of the BFGS and Gauss Newton methods, 2014.

Instead of finding a non-convergence example,
can we develop a theorem?

An overview of our theory

We assume that f is smooth and **convex** (explained later).

We denote the optimal solution set of f by \mathcal{S}^* .

We will establish the following.

Under **some assumption on $\{\mathcal{D}_k\}$** and algorithmic parameters, there exist **choices of x_0** such that

$$\mathbb{P} \left(\liminf_{k \rightarrow \infty} \text{dist}(x_k, \mathcal{S}^*) > 0 \right) > 0.$$

Differences from a non-convergence example:

one function	v.s.	some function class
special parameters	v.s.	conditions for parameters
a specific initial point	v.s.	a region for initial points

Assumption on $\{\mathcal{D}_k\}$: probabilistic ascent

Recall p -probabilistically κ -descent

$$\mathbb{P}(\text{cm}(\mathcal{D}_k, -g_k) \geq \kappa \mid \mathcal{D}_0, \dots, \mathcal{D}_{k-1}) \geq p \quad \text{for each } k \geq 0.$$

Assumption on $\{\mathcal{D}_k\}$: probabilistic ascent

Recall p -probabilistically κ -descent

$$\mathbb{P}(\text{cm}(\mathcal{D}_k, -g_k) \geq \kappa \mid \mathcal{D}_0, \dots, \mathcal{D}_{k-1}) \geq p \quad \text{for each } k \geq 0.$$

q -probabilistically **ascent**

$$\mathbb{P}(\text{cm}(\mathcal{D}_k, -g_k) \leq 0 \mid \mathcal{D}_0, \dots, \mathcal{D}_{k-1}) \geq q \quad \text{for each } k \geq 0.$$

Assumption on $\{\mathcal{D}_k\}$: probabilistic ascent

Recall p -probabilistically κ -descent

$$\mathbb{P}(\text{cm}(\mathcal{D}_k, -g_k) \geq \kappa \mid \mathcal{D}_0, \dots, \mathcal{D}_{k-1}) \geq p \quad \text{for each } k \geq 0.$$

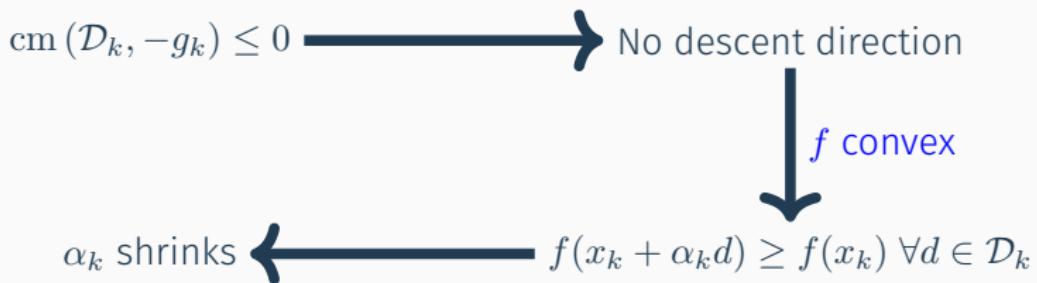
q -probabilistically **ascent**

$$\mathbb{P}(\text{cm}(\mathcal{D}_k, -g_k) \leq 0 \mid \mathcal{D}_0, \dots, \mathcal{D}_{k-1}) \geq q \quad \text{for each } k \geq 0.$$

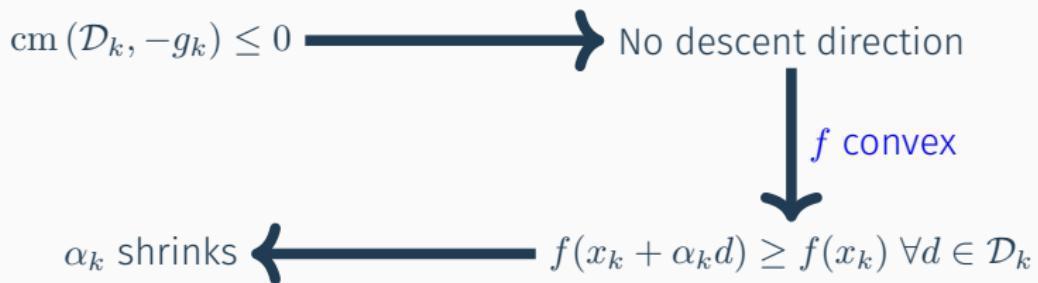
Note

If $\text{cm}(\mathcal{D}_k, -g_k) \leq 0$, then \mathcal{D}_k is “bad” (no descent direction).

Why assuming convexity?

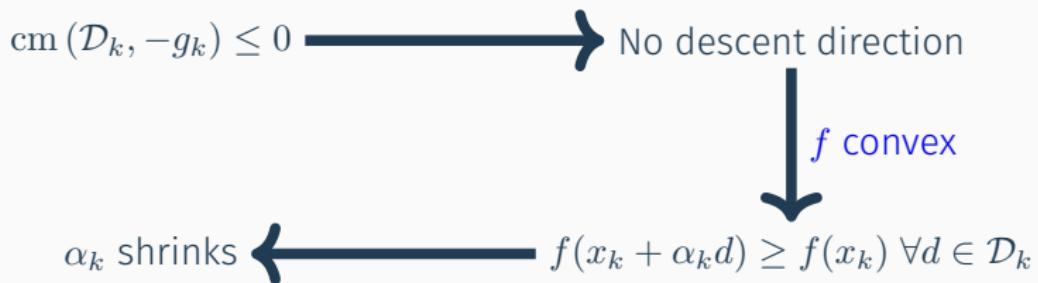


Why assuming convexity?



- Convexity connects $\text{cm}(\mathcal{D}_k, -g_k) \leq 0$ and shrinking of step size

Why assuming convexity?



- Convexity connects $\text{cm}(\mathcal{D}_k, -g_k) \leq 0$ and shrinking of step size
- $\{\mathcal{D}_k\}$ is probabilistic ascent implies α_k “often” shrinks

From probabilistic ascent to non-convergence: How?

$\{\mathcal{D}_k\}$ is probabilistically ascent

From probabilistic ascent to non-convergence: How?

$\{\mathcal{D}_k\}$ is probabilistically ascent



α_k “often” shrinks

From probabilistic ascent to non-convergence: How?

$\{\mathcal{D}_k\}$ is probabilistically ascent



α_k “often” shrinks



$$\mathbb{P} \left(\sum_{k=0}^{\infty} \alpha_k \text{ is “bounded”} \right) > 0 ?$$

From probabilistic ascent to non-convergence: How?

$\{\mathcal{D}_k\}$ is probabilistically ascent



α_k “often” shrinks



$$\mathbb{P} \left(\sum_{k=0}^{\infty} \alpha_k \text{ is “bounded”} \right) > 0 ?$$



$\mathbb{P}(\text{non-convergence}) > 0$ if $\text{dist}(x_0, \mathcal{S}^*)$ is “large”?

Key ingredients of the analysis

- Define the indicator function for “bad \mathcal{D}_k ”:

$$Y_k = \mathbb{1}(\text{cm}(\mathcal{D}_k, -g_k) \leq 0)$$

Key ingredients of the analysis

- Define the indicator function for “bad \mathcal{D}_k ”:

$$Y_k = \mathbb{1}(\text{cm}(\mathcal{D}_k, -g_k) \leq 0)$$

- Note the following inequality between step sizes (f is convex):

$$\alpha_{k+1} \leq \begin{cases} \gamma\alpha_k, & \text{if } Y_k = 0 \\ \theta\alpha_k, & \text{if } Y_k = 1 \end{cases} = \gamma^{1-Y_k} \theta^{Y_k} \alpha_k$$

Key ingredients of the analysis

- Define the indicator function for “bad \mathcal{D}_k ”:

$$Y_k = \mathbb{1}(\text{cm}(\mathcal{D}_k, -g_k) \leq 0)$$

- Note the following inequality between step sizes (f is convex):

$$\alpha_{k+1} \leq \begin{cases} \gamma\alpha_k, & \text{if } Y_k = 0 \\ \theta\alpha_k, & \text{if } Y_k = 1 \end{cases} = \gamma^{1-Y_k}\theta^{Y_k}\alpha_k$$

- Use the above inequality iteratively:

$$\alpha_k \leq \alpha_0 \prod_{\ell=0}^{k-1} \gamma^{1-Y_\ell}\theta^{Y_\ell}$$

Key ingredients of the analysis

- Define the indicator function for “bad \mathcal{D}_k ”:

$$Y_k = \mathbb{1}(\text{cm}(\mathcal{D}_k, -g_k) \leq 0)$$

- Note the following inequality between step sizes (f is convex):

$$\alpha_{k+1} \leq \begin{cases} \gamma\alpha_k, & \text{if } Y_k = 0 \\ \theta\alpha_k, & \text{if } Y_k = 1 \end{cases} = \gamma^{1-Y_k}\theta^{Y_k}\alpha_k$$

- Use the above inequality iteratively:

$$\alpha_k \leq \alpha_0 \prod_{\ell=0}^{k-1} \gamma^{1-Y_\ell}\theta^{Y_\ell}$$

- Get an upper bound of the series of step sizes:

$$\sum_{k=1}^{\infty} \alpha_k \leq \alpha_0 \sum_{k=1}^{\infty} \prod_{\ell=0}^{k-1} \gamma^{1-Y_\ell}\theta^{Y_\ell} =: \alpha_0 S$$

- Analyze the behavior of the random series S

A closer look at the random series S

Recall that

$$S = \sum_{k=1}^{\infty} \prod_{\ell=0}^{k-1} \gamma^{1-Y_\ell} \theta^{Y_\ell},$$

where $Y_\ell = \mathbb{1}(\text{cm}(\mathcal{D}_\ell, -g_\ell) \leq 0)$.

A closer look at the random series S

Recall that

$$S = \sum_{k=1}^{\infty} \prod_{\ell=0}^{k-1} \gamma^{1-Y_\ell} \theta^{Y_\ell},$$

where $Y_\ell = \mathbb{1}(\text{cm}(\mathcal{D}_\ell, -g_\ell) \leq 0)$.

Two questions

- (Q1) Does there exist a constant ζ such that

$$\mathbb{P}(S < \zeta) > 0?$$

- (Q2) Moreover, can we specify the value of ζ ?

Answer to Q1 and Q2

Proposition

If $\{\mathcal{D}_k\}$ is q -probabilistically ascent with $q > q_0$, where

$$q_0 = 1 - p_0 = \frac{\log \gamma}{\log(\theta^{-1}\gamma)},$$

then

1.

$$\mathbb{P}(S < \infty) = 1,$$

2.

$$\mathbb{P}(S < \zeta) > 0 \iff \zeta > \frac{\theta}{1 - \theta}.$$

Answer to Q1 and Q2

Proposition

If $\{\mathcal{D}_k\}$ is q -probabilistically ascent with $q > q_0$, where

$$q_0 = 1 - p_0 = \frac{\log \gamma}{\log(\theta^{-1}\gamma)},$$

then

1.

$$\mathbb{P}(S < \infty) = 1,$$

2.

$$\mathbb{P}(S < \zeta) > 0 \iff \zeta > \frac{\theta}{1 - \theta}.$$

Note

- $\mathbb{P}(S < \infty) = 1$ already implies the existence of a ζ but not its value.
- The lower bound in 2 is tight, as $S = \sum_{k=1}^{\infty} \prod_{\ell=0}^{k-1} \gamma^{1-Y_\ell} \theta^{Y_\ell} \geq \frac{\theta}{1 - \theta}$.

Non-convergence of PDS

Theorem

Under aforementioned assumptions on f , if the sequence $\{\mathcal{D}_k\}$ in PDS is q -probabilistically ascent with $q > q_0$, then

$$\mathbb{P} \left(\liminf_{k \rightarrow \infty} \text{dist}(x_k, \mathcal{S}^*) > 0 \right) > 0,$$

provided that $\text{dist}(x_0, \mathcal{S}^*) > \alpha_0 / (1 - \theta)$.

What happens in the typical implementation of PDS?

Let $\mathcal{D}_k = \{d_1, \dots, d_m\}$, where $d_\ell \stackrel{\text{i.i.d.}}{\sim} \mathsf{U}(\mathcal{S}^{n-1})$.

Recall that PDS is convergent if

$$m > \log_2 \left(1 - \frac{\log \theta}{\log \gamma} \right).$$

What happens in the typical implementation of PDS?

Let $\mathcal{D}_k = \{d_1, \dots, d_m\}$, where $d_\ell \stackrel{\text{i.i.d.}}{\sim} \mathcal{U}(\mathcal{S}^{n-1})$.

Recall that PDS is convergent if

$$m > \log_2 \left(1 - \frac{\log \theta}{\log \gamma} \right).$$

With our non-convergence analysis, PDS is non-convergent if

$$\mathbb{P} (\text{cm}(\mathcal{D}_k, -g_k) \leq 0 \mid \mathcal{D}_0, \dots, \mathcal{D}_{k-1}) > q_0,$$

What happens in the typical implementation of PDS?

Let $\mathcal{D}_k = \{d_1, \dots, d_m\}$, where $d_\ell \stackrel{\text{i.i.d.}}{\sim} \mathcal{U}(\mathcal{S}^{n-1})$.

Recall that PDS is convergent if

$$m > \log_2 \left(1 - \frac{\log \theta}{\log \gamma} \right).$$

With our non-convergence analysis, PDS is non-convergent if

$$\mathbb{P} (\text{cm}(\mathcal{D}_k, -g_k) \leq 0 \mid \mathcal{D}_0, \dots, \mathcal{D}_{k-1}) > q_0,$$

which is equivalent to

$$\left(\frac{1}{2} \right)^m > q_0 = \frac{\log \gamma}{\log(\theta^{-1}\gamma)},$$

What happens in the typical implementation of PDS?

Let $\mathcal{D}_k = \{d_1, \dots, d_m\}$, where $d_\ell \stackrel{\text{i.i.d.}}{\sim} \mathcal{U}(\mathcal{S}^{n-1})$.

Recall that PDS is convergent if

$$m > \log_2 \left(1 - \frac{\log \theta}{\log \gamma} \right).$$

With our non-convergence analysis, PDS is non-convergent if

$$\mathbb{P} (\text{cm}(\mathcal{D}_k, -g_k) \leq 0 \mid \mathcal{D}_0, \dots, \mathcal{D}_{k-1}) > q_0,$$

which is equivalent to

$$\left(\frac{1}{2} \right)^m > q_0 = \frac{\log \gamma}{\log(\theta^{-1}\gamma)},$$

or, equivalently,

$$m < \log_2 \left(1 - \frac{\log \theta}{\log \gamma} \right).$$

Tightness of our assumption on $\{\mathcal{D}_k\}$

Our assumption on $\{\mathcal{D}_k\}$:

q -probabilistically ascent with $q > q_0$.

Natural question:

Is it sufficient to require $q \geq q_0$?

Tightness of our assumption on $\{\mathcal{D}_k\}$

Our assumption on $\{\mathcal{D}_k\}$:

q -probabilistically ascent with $q > q_0$.

Natural question:

Is it sufficient to require $q \geq q_0$?

Answer: NO!

Tightness of our assumption on $\{\mathcal{D}_k\}$

Our assumption on $\{\mathcal{D}_k\}$:

q -probabilistically ascent with $q > q_0$.

Natural question:

Is it sufficient to require $q \geq q_0$?

Answer: NO!

Example

We assume

- $\theta = 1/2$ and $\gamma = 2$, which implies $q_0 = 1/2$;
- $\mathcal{D}_k = \{g_k/\|g_k\|\}$ or $\{-g_k/\|g_k\|\}$ with probability $1/2$, respectively.

Then PDS converges w.p.1.

Convergence results inspired by non-convergence analysis

Consider the series

$$S(\kappa) = \sum_{k=1}^{\infty} \prod_{\ell=0}^{k-1} \gamma^{1-Y_\ell(\kappa)} \theta^{Y_\ell(\kappa)},$$

where $Y_\ell(\kappa) = \mathbb{1}(\text{cm}(\mathcal{D}_\ell, -g_\ell) \leq \kappa)$.

Convergence results inspired by non-convergence analysis

Consider the series

$$S(\kappa) = \sum_{k=1}^{\infty} \prod_{\ell=0}^{k-1} \gamma^{1-Y_\ell(\kappa)} \theta^{Y_\ell(\kappa)},$$

where $Y_\ell(\kappa) = \mathbb{1}(\text{cm}(\mathcal{D}_\ell, -g_\ell) \leq \kappa)$.

Roughly speaking, $S(0) < \infty$ implies non-convergence of PDS.

What can we say about convergence using $S(\kappa)$?

Convergence results inspired by non-convergence analysis

Consider the series

$$S(\kappa) = \sum_{k=1}^{\infty} \prod_{\ell=0}^{k-1} \gamma^{1-Y_\ell(\kappa)} \theta^{Y_\ell(\kappa)},$$

where $Y_\ell(\kappa) = \mathbb{1}(\text{cm}(\mathcal{D}_\ell, -g_\ell) \leq \kappa)$.

Roughly speaking, $S(0) < \infty$ implies non-convergence of PDS.

What can we say about convergence using $S(\kappa)$?

Theorem

If there exists a $\kappa > 0$ such that $S(\kappa) = \infty$, then DS converges.

Convergence results inspired by non-convergence analysis

Consider the series

$$S(\kappa) = \sum_{k=1}^{\infty} \prod_{\ell=0}^{k-1} \gamma^{1-Y_\ell(\kappa)} \theta^{Y_\ell(\kappa)},$$

where $Y_\ell(\kappa) = \mathbb{1}(\text{cm}(\mathcal{D}_\ell, -g_\ell) \leq \kappa)$.

Roughly speaking, $S(0) < \infty$ implies non-convergence of PDS.

What can we say about convergence using $S(\kappa)$?

Theorem

If there exists a $\kappa > 0$ such that $S(\kappa) = \infty$, then DS converges.

Relation with existing result in [GRVZ 2015]

p_0 -probabilistically κ -descent $\implies S(\kappa) = \infty$ w.p.1

Take away

In this talk, we

- theoretically explain the non-convergence phenomenon of PDS,
- find out the behavior of PDS is closely related to the random series

$$S = \sum_{k=1}^{\infty} \prod_{\ell=0}^{k-1} \gamma^{1-Y_\ell} \theta^{Y_\ell}.$$

Non-convergence analysis can

- sharpen our knowledge about the algorithm,
- deepen our understanding about the convergence analysis,
- guide the selection of algorithmic parameters, and
- provide new perspectives on convergence analysis.

Thank you!

One more thing: OptiProfiler



github.com/optiprofiler

OptiProfiler (joint work with Cunxin Huang and Tom M. Ragonneau) is
a [benchmarking platform](#) for DFO solvers.

Our goal: [fair](#), [convenient](#), and [uniform](#) benchmarking.

- Creating [performance profiles](#), [data profiles](#), and [log-ratio profiles](#) [Moré, Wild 2009; Shi, Xuan, Oztoprak, and Nocedal 2023]
- Providing multiple types of tests
noisy function, unrelaxable constraints, randomized initial point ...
- Implemented in Python and MATLAB
- Default problem set: S2MPJ [Gratton, Toint 2024]

One more thing: OptiProfiler

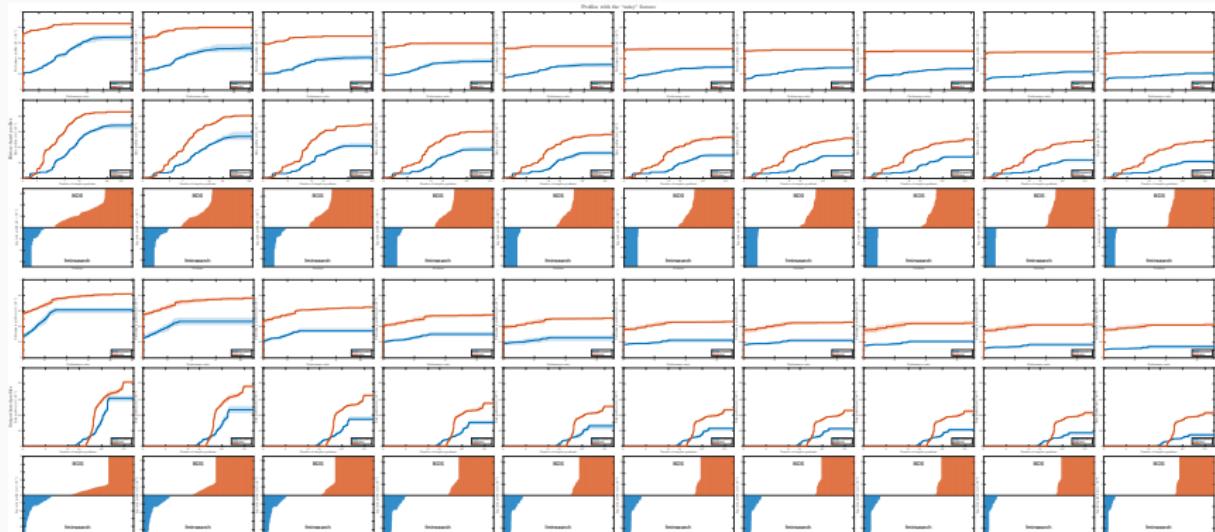
Example (MATLAB):

```
benchmark({@bds, @fminsearch}, "noisy")
```

One more thing: OptiProfiler

Example (MATLAB):

```
benchmark({@bds, @fminsearch}, "noisy")
```



N.B.: Separate profiles can also be generated.

References I

- ▶ Chen, C. et al. (2016). “The direct extension of ADMM for multi-block convex minimization problems is not necessarily convergent”. *Math. Program.* 155, pp. 57–79.
- ▶ Conn, A. R., Scheinberg, K., and Vicente, L. N. (2009). *Introduction to Derivative-Free Optimization*. Vol. 8. MOS-SIAM Ser. Optim. Philadelphia: SIAM.
- ▶ Durrett, R. (2010). *Probability: Theory and Examples*. Fourth. Camb. Ser. Stat. Probab. Math. Cambridge: Cambridge University Press.
- ▶ Fermi, E. and Metropolis, N. (1952). *Numerical solution of a minimum problem*. Tech. rep. Alamos National Laboratory, Los Alamos, USA.
- ▶ Ghanbari, H. and Scheinberg, K. (2017). “Black-box optimization in machine learning with trust region based derivative free algorithm”. *arXiv:1703.06925*.

References II

- ▶ Gratton, S. et al. (2015). “**Direct search based on probabilistic descent**”. *SIAM J. Optim.* 25, pp. 1515–1541.
- ▶ Kolda, T. G., Lewis, R. M., and Torczon, V. (2003). “**Optimization by direct search: New perspectives on some classical and modern methods**”. *SIAM Rev.* 45, pp. 385–482.
- ▶ Larson, J., Menickelly, M., and Wild, S. M. (2019). “**Derivative-free optimization methods**”. *Acta Numer.* 28, pp. 287–404.
- ▶ Mascarenhas, W. (2014). “**The divergence of the BFGS and Gauss Newton methods**”. *Math. Program.* 147, pp. 253–276.
- ▶ Powell, M. J. D. (1973). “**On search directions for minimization algorithms**”. *Math. Program.* 4, pp. 193–201.

References III

- Yuan, Y. (1998). “An example of non-convergence of trust region algorithms”. In: *Advances in Nonlinear Programming*. Ed. by Y. Yuan. Dordrecht: Kluwer Academic Publishers, pp. 205–215.