

Master thesis

**Work Distribution for a Heterogeneous Library Staff -
A Personnel Task Scheduling Problem**

Claes Arvidson and Emelie Karlsson

LITH - MAT - EX - - 04 / 04 - - SE

Work Distribution for a Heterogeneous Library Staff - A Personnel Task Scheduling Problem

Optimeringslära, Linköpings Universitet

Claes Arvidson and Emelie Karlsson

LITH - MAT - EX - - 04 / 04 - - SE

Exam work: **30 hp**

Level: **A**

Supervisor: **Torbjörn Larsson**,
Optimeringslära, Linköpings Universitet

Examiner: **Elina Rönnberg**,
Optimeringslära, Linköpings Universitet

Linköping: **June 2016**

Abstract

The distribution of tasks to a heterogeneous work force at libraries and other service institutions is time consuming for manual schedulers, but well suited for optimization softwares. The problem studied concerns five different types of tasks, two types of workers and around 100 tasks to be scheduled per week. Also, there is a demand for workers during weekends and worker satisfaction is taken into account.

The main objective is to create a optimal ten week rotating schedule, in which the stand-in staff members are evenly distributed. That results in a robust schedule, as absence of workers is a common issue.

A mathematical model is formulated for the problem, which is solved using the commercial solver, CPLEX and by using two different LNS heuristic implementations. The first heuristic assigns week blocks to the workers, while the second distributes one task at a time. The latter heuristic works better than the former and achieves results comparable to those of the commercial solver. Our conclusion is that the second heuristic works better, since it focuses on finding a good weekend distribution before creating the rest of the schedule.

Keywords: Optimization, Scheduling, Task distribution, LNS, Weekend Scheduling, Heterogeneous workforce

URL for electronic version:

<http://urn.kb.se/resolve?urn=urn:nbn:se:liu:diva-77777>

Acknowledgements

The authors of this thesis would like to express our deepest gratitude to our supervisor Torbjörn Larsson at Linköping University, who has helped us find our way in moments of uncertainty and who has been guiding us through the whole project. Thank you for the time and effort you have spent on our thesis. We would also like to thank Elina Rönnerberg, who has provided encouraging words and valuable insights throughout the project.

Furthermore, we would like to thank Elisabeth Cserhalmi and Ingrid Loeld Rasch at Norrköpings Stadsbibliotek for providing us with an interesting thesis topic and for answering all our questions many times.

We thank our opponents, Akdas Hossain and Emma Miléus, for their comments and thoughts on this report.

Lastly, we would like to thank our families and loved ones, who have supported us until the end.

Nomenclature

Most of the reoccurring terms and abbreviations are described here.

Optimization terms

Heuristic An algorithm designed to find feasible, but not necessarily optimal, solutions.

Library terms

Fetch list The library task of collecting books from shelves, to be delivered elsewhere.
Library on wheels The task of driving a library bus with books to remote areas of town.

Abbreviations

Exp Service counter (sv. expeditionsdiskén)
Info Information counter (sv. informationsdiskén)
PL Fetch list (sv. plocklistan)
BokB Library on wheels (sv. bokbussen)
HB Hageby library

Contents

1	Introduction	1
1.1	Problem description	1
1.1.1	Description of the daily tasks at the library	2
1.1.2	Personnel attributes	4
1.1.3	Scheduling objectives: stand-in maximization and schedule variation	5
1.2	Aims and goals	6
1.3	Method	6
1.4	Outline	7
1.5	Contributions of the authors	7
2	Literature review	9
2.1	Personnel task scheduling problem	9
2.1.1	Applications	11
2.2	Shift minimization personnel task scheduling problem	11
2.3	Tour scheduling problem	13
2.4	Other similar problems	14
2.4.1	Fixed job schedule problem	14
2.4.2	Tactical fixed interval scheduling problem	14
2.4.3	Operational fixed interval scheduling problem	15
2.4.4	Stochastic job problems	15
2.5	Modelling soft constraints	16
2.6	Summary	17
2.7	Relevance to our problem	17
3	The mathematical model	19
3.1	Set and variable definitions	19
3.2	Objective function	21
3.3	Constraints	21
3.3.1	Demand and assignment constraints	21
3.3.2	Weekend and rotation constraints	22
3.3.3	Definitional constraints	24
3.3.4	Meeting constraints	25
4	Two heuristic solution methods	27
4.1	Week block scheduling approach	28
4.1.1	Block creation	29
4.1.2	Block filtering	31

4.1.3	Rotation assignment	32
4.1.4	Assignment of Library on Wheels	33
4.1.5	Initial solution	34
4.1.6	Costs	34
4.1.7	Destroy	36
4.1.8	Repair	36
4.1.9	Evaluation of solution	37
4.1.10	Final phase	37
4.2	Task distribution approach	38
4.2.1	Costs	38
4.2.2	Weekend phase	40
4.2.3	Weekday phase	41
5	Results and discussion	45
5.1	Results	45
5.1.1	AMPL implementation	45
5.1.2	Week block scheduling approach	45
5.1.3	Task distribution approach	48
5.2	Discussion	52
5.2.1	Week block scheduling approach	52
5.2.2	Task distribution approach	53
6	Concluding remarks	55
A	Problem definitions	59
A.1	Sets	59
A.2	Variables	60
A.3	Parameters	61
A.4	Objective function and constraints	61
B	Flow charts	65
B.1	Weekly scheduling approach	65
B.2	Task distribution approach	65
C	Week block table	69

List of Figures

4.1	Illustration of the differences between the model used in AMPL and the heuristic model	27
4.2	A flow chart of the implemented heuristic with week block scheduling	30
4.3	A flow chart of costs appearing when a PL is assigned a block on a Monday, third week relative to the library schedule.	35
5.1	Plot of the objective function value after destroy and repair iterations, one time without enough stand-ins when trying to enter the final phase.	47
5.2	Plot of the objective function value after destroy and repair iterations with enough stand-ins when trying to enter the final phase.	48
5.3	The weekend objective function for 1000 iterations	50
5.4	The weekend objective function minimum cost components for 1000 iterations	51
5.5	The weekend objective function average costs components for 1000 iterations	51
B.1	Algorithm for distributing weekends.	66
B.2	Algorithm for distributing weekday tasks.	67

List of Tables

1.1	Outer tasks can be performed either exclusively by librarians or by both librarians and assistants.	2
1.2	Staff demand during a week. PL is marked as one shift, but it is performed during the whole day.	3
1.3	Demand of staff for Library on Wheels	4

1.4	Availability schedule for a sample staff member. Yellow signifies that the person is available. In parenthesis, the weekend shift. . .	5
1.5	Example of a feasible week for a staff member.	5
2.1	PTSP variants	10
3.1	Resulting table of the function $\text{mod}_{10}(w - v + 10) + 1$	23
4.1	Showing changes in the model for the week block scheduling approach. The soft, relaxed and removed constraints are shown. . .	28
4.2	Showing changes in the model for the task distribution approach. The soft, relaxed and removed constraints are shown.	29
4.3	A general week block with all possible tasks	30
4.4	Illustration of one of the unique block appearances	31
4.5	Typical availability for a five-week period for a staff member. Yellow signifies that the staff member is available during the hours. In parenthesis, the weekend shift hours are given.	32
4.6	A weekend block with Desk tasks preventing any other tasks on Fridays.	32
4.7	An iteration in the destroy and repair iteration showing a swap of weekends when three staff members' week rotations are destroyed.	33
4.8	A block example to be evaluated using costs	34
4.9	Library demand at a shift and solution qualities.	36
4.10	List of all costs with description.	36
4.11	Cost evaluation for two blocks. The lower the total cost is, the more desired it is to insert in the repair.	37
4.12	Individual staff member costs and cost weights.	38
4.13	Library objective functions with their cost components and weights.	39
4.14	Staff member availability placing only weekends. Intensity of red indicates number of staff members.	41
4.15	Staff member availability after placing weekends as well as evening tasks and BokB for the same week.	41
4.16	Staff member demand during a week when entering weekday phase.	42
5.1	Results from solving the mathematical model with CPLEX solver.	45
5.2	List of all costs used in the week block scheduling approach and their respective values.	46
5.3	Number of successful runs	46
5.4	Number of stand-ins found in 420 successful runs.	47
5.5	Weights used in the implementation.	49
5.6	Results from the task distribution heuristic when varying It_{wend}	49
5.7	Results from the task distribution heuristic when varying It_{wday}	49
5.8	Results from the task distribution heuristic when varying weights. $It_{wend} = 1000$ and $It_{wday} = 20$ in all runs.	49
5.9	Pros and cons with the implemented week block scheduling approach	53
5.10	Pros and cons with the implemented task distribution approach	54
C.1	Number of assignable unique blocks for the workers based on their availability and qualification.	69

Chapter 1

Introduction

At a library, there exist a number of tasks which are to be distributed over all staff members. Poor work distribution among staff can cause problems and result in a shortage of staff at certain critical shifts. If a staff member becomes unavailable at such shifts, due to for example illness, a qualified stand-in is required to fill the vacancy. Therefore, it is of high priority to libraries and other similar institutions to create schedules with as many skilled stand-ins as possible in order to handle such unexpected disturbances.

The problem addressed in this thesis work concerns the library staff at the Central Library of Norrköping, Sweden. This library currently has 39 employees and the renown building from the 1950's is a central gathering point in Norrköping. The library is open weekdays from 8 a.m. to 8 p.m. and from 11 a.m. to 16 p.m. during weekends. The generous opening times also creates a scheduling challenge for the library as it requires a large pool of well coordinated personnel to keep the library open. In addition, the library also provides its services to one other smaller library, which adds further complexity to the problem of task distribution among the staff.

This chapter contains a problem description, where the scheduling problem at hand is described in detail, a section describing the goals and aims of the thesis work, a section about the methods used to solve the problem and an outline of this report.

1.1 Problem description

The main objective of the problem studied is to create a ten week schedule for all staff members and tasks at the library. Each staff member's ten week schedule should consist of two five-week schedules which should be as similar as possible, but not identical since some tasks are performed only even or odd weeks. The ten week schedule should be rotational.

The schedule should also be created so that the stand-ins are distributed evenly over all days. In order to do this, the library provided information about when staff members are available for task assignment, what tasks should be assigned when as well as the constraints for how these tasks should be assigned. These constraints are discussed in detail in this section.

1.1.1 Description of the daily tasks at the library

The highest priority at a library, as with any service institution, is to provide good service to visitors. This includes book loan services as well as being able to give visitors helpful information about the resources at disposal at the library. This type of work is referred to as "outer tasks" in this thesis. In addition to these, "inner tasks" such as sorting books, ordering new books and answering emails also exist and are a part of the every day library tasks as well. The problem of allocating the outer tasks, while taking inner tasks into consideration, is studied in this thesis.

Three main types of outer task can be identified at the library of Norrköping: working at the service counter (sv. expiditionsdisken), working at the information counter (sv. informationsdisken) and assembling books which are to be sent to other libraries according to the "fetch list" (sv. plocklista). The fetch list is a task for which the staff member is scheduled during a whole day, while the other tasks are scheduled for only one shift. The outer tasks can be performed by either librarians or both librarians and assistants, as described in Table 1.1.

Table 1.1: Outer tasks can be performed either exclusively by librarians or by both librarians and assistants.

Task	Description	Qualification
Service counter (Exp)	Administiring loans, library cards and the loaning machine	Librarian or assistant
Information counter (Info)	Handling questions about the library's resources.	Librarian
Fetch list (PL)	Fetching books that are to be sent to other libraries.	Librarian or assistant
Hageby (HB)	Handling librarian tasks at the filial Hageby during weekends.	Librarian
Library on Wheels (BokB)	Driving the Library on Wheels to different areas of town.	Librarian

As in the case with most libraries, the Central Library of Norrköping also has responsibilities that fall outside of its normal daily activities. One such responsibility is staffing a smaller library subsidiary in Hageby, situated in the suburban area of Norrköping, during weekends. It is decided that only librarians are qualified for this task. While some librarians prefer to work only at HB, others prefer never to.

Working during a weekend means working both Saturday and Sunday. When not working at HB, the Friday evening shift is also part of the weekend work. The weekend work is compensated with free days, often placed at the week following upon weekend work, in this thesis referred to as rest week.

Since the number of visitors at the library varies throughout the day and also during different days of the week, so does the demand for personnel for the outer tasks. This demand is illustrated in Table 1.2, which is constructed according to figures given by the library.

Another library task is the "Library on Wheels" (sv. bokbussen), for which only a handful of librarians are qualified. This work involves driving a library bus, which provides citizens in remoter areas of the city with books as well as other library services. The Library on Wheels only operates a few times a week

Table 1.2: Staff demand during a week. PL is marked as one shift, but it is performed during the whole day.

	Exp	Info	PL	HB
Monday				
Shift 1:	2	2	1	0
Shift 2:	3	3	0	0
Shift 3:	3	3	0	0
Shift 4:	3	3	0	0
Tuesday				
Shift 1:	2	2	1	0
Shift 2:	3	3	0	0
Shift 3:	3	3	0	0
Shift 4:	3	3	0	0
Wednesday				
Shift 1:	2	2	1	0
Shift 2:	3	3	0	0
Shift 3:	3	3	0	0
Shift 4:	3	3	0	0
Thursday				
Shift 1:	2	2	1	0
Shift 2:	3	3	0	0
Shift 3:	3	3	0	0
Shift 4:	3	3	0	0
Friday				
Shift 1:	2	2	1	0
Shift 2:	3	3	0	0
Shift 3:	3	3	0	0
Shift 4:	3	3	0	0
Saturday				
Shift 1:	3	3	0	1
Sunday				
Shift 1:	3	3	0	1

and the schedule differs between odd and even weeks in accordance with Table 1.3.

Apart from the outer tasks described above, there are also inner tasks at the library which sometimes need to be scheduled. One such type of inner task is meetings, which concerns a large number of staff members. There exist both library meetings, scheduled for the whole work force, and group meetings, scheduled only for a subset of the staff members. The library meetings are fixed in time and take place every fifth weeks on Monday mornings from 8:15 a.m. to 10:00 a.m. Also group meetings take place every fifth week but are flexible in time. For group meetings, all group members must be available. Only three groups have scheduled meetings: the child, adult and media group.

Table 1.3: Demand of staff for Library on Wheels

Odd Week	Mon	Tue	Wed	Thu	Fri
08:00-10:00	1	0	1	1	1
16:00-20:00	1	0	1	1	0
Even Week	Mon	Tue	Wed	Thu	Fri
08:00-10:00	1	0	1	1	0
16:00-20:00	0	0	1	1	0

1.1.2 Personnel attributes

The two types of library staff members that are to be assigned tasks are librarians and assistants. Librarians can perform all the tasks listed above, while assistants can perform a subset of these tasks.

The Central Library of Norrköping currently has 39 library staff members, 23 of which are librarians and 16 of which are assistants. These have different availabilities for performing tasks, depending on their working hours and the amount of inner work they have. In the standard case, each staff member is assigned one evening per week and once per five weeks he or she is assigned to work during the weekend. This is normally compensated the week after with two extra free days, placed according to the wishes of the staff member.

Let us consider a sample staff member who is a librarian, works full time and is also assigned to evening work on Wednesdays. The staff member is also assigned to work weekend on the fourth week and has chosen to take out its days off on Thursday and Friday during the week following the weekend. The availability for such a staff member is illustrated in Table 1.4. The schedule repeats itself after five weeks and illustrates only the availability for outer tasks, but does not show whether the staff member has been assigned any tasks during these weeks.

All staff members have a five week schedule in the same manner as the sample staff member. However, in order to meet the weekend demands, illustrated in Table 1.2, staff members have to be assigned for weekend work at different weekends. Thus, a staff member's schedule can be shifted in order to make sure that all weekend demand is fulfilled.

Considering again the sample staff member's schedule in Table 1.4, where the person is available throughout all days, five basic constraints regulate how tasks can be placed. Firstly, staff members are only allowed to take at a maximum one task per day. This guarantees that staff members also have time for inner tasks. Secondly, staff members are only allowed to perform a maximum of four tasks per week, in order for them to have a completely free day. Thirdly, staff members can only be assigned to PL at a most once a week. Also, there is a limit on how many PL tasks can be assigned to a staff member in total over all ten weeks which varies between different staff members, but which is at the most four. Lastly, a staff member should not have the same shift more than twice per week. This guarantees that there is fairness in the schedule so that no staff member has to work too many times at a shift which is considered to be "bad". An example of a feasible weekly schedule for the sample staff member is provided in Table 1.5.

The availability of all staff members is provided by the library. Some staff members work differently even and odd weeks while other never work evenings,

Table 1.4: Availability schedule for a sample staff member. Yellow signifies that the person is available. In parenthesis, the weekend shift.

Week 1	Mon	Tue	Wed	Thu	Fri	Sat	Sun
08:00-10:00 (11:00-16:00)							
10:00-13:00							
13:00-16:00							
16:00-20:00							
Week 2	Mon	Tue	Wed	Thu	Fri	Sat	Sun
08:00-10:00 (11:00-16:00)							
10:00-13:00							
13:00-16:00							
16:00-20:00							
Week 3	Mon	Tue	Wed	Thu	Fri	Sat	Sun
08:00-10:00 (11:00-16:00)							
10:00-13:00							
13:00-16:00							
16:00-20:00							
Week 4	Mon	Tue	Wed	Thu	Fri	Sat	Sun
08:00-10:00 (11:00-16:00)							
10:00-13:00							
13:00-16:00							
16:00-20:00							
Week 5	Mon	Tue	Wed	Thu	Fri	Sat	Sun
08:00-10:00 (11:00-16:00)							
10:00-13:00							
13:00-16:00							
16:00-20:00							

only work weekends or never work weekends. This data was used in order to solve the scheduling problem.

Table 1.5: Example of a feasible week for a staff member.

Week 1	Mon	Tue	Wed	Thu	Fri	Sat	Sun
08:00-10:00 (11:00-16:00)		Exp		PL			
10:00-13:00	Exp			PL			
13:00-16:00				PL			
16:00-20:00			Info				

1.1.3 Scheduling objectives: stand-in maximization and schedule variation

When we were first presented with the scheduling problem, it was explained to us that the main challenge for the library is to find a good number of stand-ins for all days. When we studied their schedule, we observed that only one stand-in was present at the worst day. Thus, the main objective of this thesis work is to create a feasible schedule, according to the constraints explained in this

section, with as many stand-ins as possible at the day with the least number of stand-ins.

In this context, a stand-in is defined as a staff member who is available for outer tasks during the first three shifts and who is not scheduled for any shift during that day. Both librarians and assistants can be stand-ins, but only librarians are qualified to be assigned to any task. Since the regular hour library activities are most crucial, there is no need for stand-ins during evenings and weekends. Similarly, there are no assigned stand-ins for the Library on Wheels or for Hageby.

Apart from maximizing the number of stand-ins for each day at the library, it is desirable for the sake of the personnel to create schedules that are repeated according to a five week pattern. Thus, the ten week schedule should ideally consist of two five week schedules.

1.2 Aims and goals

The goal of this thesis is to investigate how the ten week scheduling problem for the library staff members at the Central Library of Norrköping can be solved using mathematical optimization methods. This goal includes:

1. To investigate what has previously been done in the area of task distribution.
2. To find a mathematical model for the given problem.
3. To implement the model in AMPL.
4. To implement two heuristics for the problem using a large neighbourhood search framework, one using a week block scheduling approach and one using a task scheduling approach.
5. To evaluate and compare the models used in the three implementations as well as their performance.

1.3 Method

The method used in the thesis work for solving the problem described above was to first formulate a mathematical model for the problem, which was then solved using two different methods. The first method involved modeling the problem in the AMPL programming language and solving it by the commercial optimization solver CPLEX. The second method involved creating an independent solver in C++ and solving it using a heuristic. Worker availability data was imported from Excel sheets to data files used to solve the problem through Visual Basic.

Two different heuristics were tested in the second approach. The first heuristic distributes week blocks to staff members. All blocks are pre-generated and feasible according to the constraints described in the previous section. In the second heuristic, individual tasks are distributed to the staff members. In order to move through the solution space and find better solutions, both heuristics were based on a Large Neighbourhood Search (LNS) algorithm.

1.4 Outline

The thesis is divided into six different chapters. The first chapter provides a description of the problem studied as well as the context of the study. In Chapter 2, previous work in the area is presented. The mathematical model constructed for solving the problem is presented and explained in Chapter 3. In Chapter 4, the different implementations of the model are presented and the results for these implementations are given in Chapter 5. Lastly, conclusions from the work and suggestions for further development are presented in Chapter 6.

1.5 Contributions of the authors

The authors of this thesis have implemented some parts of the work individually and some together. In particular, the mathematical model was constructed and implemented by both authors, whereas the two heuristics were constructed and implemented separately. The week block scheduling heuristic was implemented by Claes while the task distribution heuristic was implemented by Emelie.

In the report, Chapter 1 was mainly written by Emelie. Chapter 2 was divided so that Claes wrote 2.1 and 2.2 Emelie wrote 2.3 and 2.5. Other parts in the chapter were written together. The third chapter was written by Claes and in Chapter 4 and 5, the authors wrote the parts which concerned their own heuristic. Chapter 6 was written by both authors.

Chapter 2

Literature review

Staff and machine scheduling problems are mathematical optimization problems which have been studied since the 1950s and concern creating feasible and satisfactory schedules for workers or machines performing tasks. One of the most extensive overviews of the area of staff scheduling is provided by Ernst et al. (2004). They state that, although the complexity of scheduling problems has not increased in recent years, the mathematical models used to solve scheduling problems have become more realistic and refined. Modern staff scheduling problems often concern the distribution of tasks and the creation of worker shifts, also taking softer values into account, such as worker satisfaction and worker fatigue. Due to this modelling refinement as well as the development of more powerful computational methods, it is possible today to solve staff scheduling problems in a more satisfactory way than before.

In this chapter, staff scheduling problems are classified into different subcategories which are related to the thesis work. The relevant areas for our work include Personnel Task Scheduling Problems (PTSP), Shift Minimization Task Scheduling Problems (SMTSP), Tour Scheduling Problems (TSP), and a few variations of these. Also problems featuring soft constraints will be studied in an additional section. Models and solution methods of these problems will be discussed under each headline and a summary will be provided at the end, together with a discussion about the relevance of the studied articles to our thesis work.

2.1 Personnel task scheduling problem

In many real life situations production managers will face the Personnel Task Scheduling Problem (PTSP) while scheduling service operations. Krishnamoorthy and Ernst (2001) write in their article that the PTSP occurs when the rosterer or shift supervisor need to allocate tasks with specified start and end times to available personnel who have the required qualifications. It also occurs in situations where tasks of fixed times shall be assigned to machines. Each machine have a certain amount of maintenance workers assigned to keep it in operation. Decisions will then have to be made both regarding which machine the workers are assigned to, as well as the optimal amount of maintenance workers needed at each machine.

There are numerous variants of the PTSP. Studies of these have been made by Krishnamoorthy and Ernst (2001), who give a list of attributes that commonly appear in a PTSP, which are listed in Table 2.1. Furthermore, there are traits that always appear in a PTSP; tasks with fixed start and end time are to be distributed to staff members that possess certain skills, allowing them to perform only a subset of the available tasks. Start and end time of their shifts are also predetermined for each day.

One variant, which also is the most simple, is mentioned in Krishnamoorthy and Ernst (2001) and is called the *Feasibility Problem*, where the aim is to merely find a feasible solution. This requires that each task is allocated to a qualified and available worker. It is also required that a worker cannot be assigned more than one task simultaneously, as well as tasks cannot be pre-empted, meaning that each task has to be completed by one and the same worker.

In Table 2.1 attributes of PTSP variants are shown. The nomenclature of the attributes T, S, Q and O refer to the *Task type*, *Shift type*, *Qualifications* and *Objective function* respectively.

Table 2.1: PTSP variants

Attribute	Type	Explanation
T	F	Fixed contiguous tasks
	V	Variable task durations
	S	Split (non-contiguous) tasks
	C	Changeover times between consecutive tasks
S	F	Fixed, given shift lengths
	I	Identical shifts which are effectively of infinite duration
	D	Maximum duration without given start or end times
	U	Unlimited number of shifts of each type available
Q	I	Identical qualification for all staff (homogeneous workforce)
	H	Heterogeneous workforce
O	F	No objective, just find a feasible schedule
	A	Minimize assignment cost
	T	Worktime costs including overtime
	W	Minimize number of workers
	U	Minimize unallocated tasks

Many of the most basic problems and a few more complex ones can be described with this definition of PTSP attributes. It is, however, not possible to describe all of the numerous types of PTSP using these nomenclatures according to Krishnamoorthy and Ernst (2001).

By combining attributes it is possible to obtain more complex variants of the PTSP. An example would be the PTSP[F;F;H;A-T-W] mentioned in Krishnamoorthy and Ernst (2001), where multiple objectives are used. This problem has fixed contiguous tasks, fixed shift lengths, heterogeneous workforce and three objective functions (A-T-W), which represent assignment costs, work time with overtime included and requirements to minimize the number of workers, respectively. The objective function for this problem is then a linear combination with parameters used to weigh (prioritize) them against each other.

Given the nomenclature above, our problem would be most related to the PTSP[F;F;H;F]. The difference is that the objective function is not empty. We

are looking to maximize the number of qualified stand-ins each day as well as maximize employee satisfaction by meeting their recommendations. We have a fixed number of workers, no costs and no unallocated tasks when a feasible solution is found. Therefore, none of the objective function types given in Table 2.1 are relevant in our case.

Some variants of the PTSP are given specific names in the literature, which is stated by Krishnamoorthy and Ernst (2001). An example is when the shifts and qualifications are identical ($S=I$, $Q=I$) and the objective function is to minimize the number of workers that are used ($O=W$). This variant, $PTSP[F;I;I;W]$, has been published as the Fixed Job Schedule Problem and is described in Section 2.4.

2.1.1 Applications

An example where the PTSP can be found is when developing a rostering solution for ground personnel at an airport. This is mentioned in the article by Krishnamoorthy and Ernst (2001). This problem can be dealt with by first assigning the workers to days in order to satisfy all the labour constraints, followed by assigning the tasks to the scheduled workers.

In the article mentioned above, three further problems of type PTSP, all related to airplanes, are mentioned. They occur when scheduling for either airport maintenance staff, planes to gates or staff that do not stay in one location, such as airline stewards. Scheduling for airport maintenance staff can lead to either $PTSP[F;I;H;U-A]$ or $PTSP[F;I-U;H;W]$, which are similar problems but given two different names: Operational Fixed Interval Scheduling Problem and Tactical Fixed Interval Scheduling Problem respectively. These are both described further in Section 2.4.

Another application, which has been frequently studied, is classroom assignments and is discussed in Krishnamoorthy and Ernst (2001). Based on specifications such as the amount of students in a class or the duration of a class, different classrooms have to be considered. Requirements of equipment, e.g. for a laboratory, may also greatly limit the available classrooms to choose from. A majority of the complications of this problem is due to the fact that lessons can span over multiple periods.

Worth noting for classroom assignment problems is that there are no start or end times for the shifts, as they represent the rooms. The aim in this problem would be to simply find a feasible assignment of classrooms. Therefore, the nomenclature of the problem would be $PTSP[S;I;H;F]$, with the possibility of adding preferences to the objective function. An example of a preference would be to assign the lessons as close to each other as possible on a day, preventing travel distances for teachers and students.

2.2 Shift minimization personnel task scheduling problem

A close relative to the PTSP is the Shift Minimization Personnel Task Scheduling Problem (SMPTSP), which is a variant where the aim is to minimize the cost occurring due to the number of personnel that are used. The same common traits for this problem are mentioned in the article Krishnamoorthy et al. (2012)

as for the PTSP: Workers with fixed work hours are to be assigned tasks with specified start and end times, for which they are qualified.

In the same article as above the authors "... concentrate mainly on a variant of the PTSP in which the number of personnel (shifts) required is to be minimized". In doing so, it is possible to determine the lowest number of staff and the mix of staff a company needs in order to be able to complete the tasks to be operational. They also presumed that the pool of workers is unlimited for each skill group, which is not the case in our problem due to the limitations of librarians and assistants.

Furthermore, it is said in Krishnamoorthy et al. (2012) that SMPTSP can be applied when there is a large number of workers available with different qualifications and it must be ensured that the tasks for that day are performed. The PTSP and SMPTSP are therefore useful day-to-day operational management tools and commonly occur in practical instances where tasks are allocated on a daily basis.

It is shown in Kroon et al. (1995) that SMPTSP is a complex problem even if the preemption constraint were to be removed. However, it is stated in Krishnamoorthy et al. (2012) that if the qualifications of the workers were identical it would be an easily solvable problem. The difference in publication year between the articles indicates that this problem has become less challenging in recent years.

During the last decade, a couple of heuristics have been implemented to deal with the SMPTSP. One method introduced by Krishnamoorthy et al. (2012) is a Lagrangean relaxation approach that combines two problem specific heuristics: Volume Algorithm (VA) and Wedelin's Algorithm (WA). These heuristics exploit the special structure of the SMPTSP by relaxing some of the harder constraints into the objective function, thus being a problem specific heuristic. What remains after the relaxation is a problem decomposed into several problems which can be solved independently. One way to solve these decomposed problems is discussed further in Krishnamoorthy et al. (2012).

Another way to solve the SMPTSP is to use a very large-scale neighbourhood search algorithm, as was presented by Smet and Vanden Berghe (2012). The main purpose of their implemented hybrid local search is to repeatedly fix and optimize to find neighbouring solutions. Using this method on 137 benchmark instances introduced by Krishnamoorthy et al. (2012), Smet and Vanden Berghe managed to find 81 optimal solutions compared to Krishnamoorthy et al. who only managed to find 67. Though, both methods found feasible solutions for 135 out of the 137 problem instances. This comparison is presented in Smet et al. (2014).

Smet et al. (2014) introduced a third and most effective method to solve the benchmark instances related to the SMPTSP. By using a versatile two-phase matheuristic approach, solutions to all 137 benchmark instances could be found for the first time. The procedure used in their implementation is to first generate an initial solution by using a constructive heuristic, followed by improving the solution using an improvement heuristic.

2.3 Tour scheduling problem

The Tour Scheduling Problem (TSP) is described by Loucks and Jacobs (1991) as a combination of shift scheduling and days-off scheduling. Shift scheduling refers to creating sets of contiguous hours during which a worker is assigned for work. The need for days off scheduling typically occurs when the time horizon for scheduling is weekly or more and when weekend staff is needed. Using the notation in Section 2.1, this would be classified as a variant of the PTSP[F;D;I;W] or PTSP[F;D;H;W], depending on if the workforce is homogeneous or heterogeneous.

According to Loucks and Jacobs (1991), the vast majority of all tour scheduling problems up to 1991 involve a homogeneous workforce, that is, any worker can perform any assigned task. One such early study of a tour scheduling problem is provided by Thompson (1988). The problem studied concerns only homogeneous workforces and the task assignment part is not considered.

In the article by Loucks and Jacobs (1991), the authors study a tour scheduling problem with a heterogeneous workforce. The problem both involves tour scheduling and task assignment, where the latter part is most interesting to us. The problem is studied in the context of fast food restaurants, where certain personnel is qualified only for certain stations in the restaurant. In such industries, the demand of staff differs between weekdays and times of the day. Two worker attributes are considered, their availability for work and their qualifications to perform different tasks. The problem concerns finding shifts for all workers which are to be assigned a length between a minimum and maximum number of hours per day.

The main problem studied in the article involves creating a one-week schedule for 40 workers in a fast food restaurant, available for eight different tasks with a seven-day and 128-hour workweek. Several synthetic problems are studied in the article, all, however, with a minimum shift length of three hours, a maximum shift length of eight hours and a maximum of five work days.

A similar problem to the one described by Loucks and Jacobs is studied by Choi et al. (2009). They focus on a particular fast food restaurant in Seoul, which is made a representative of fast food chains in general. In this study, only two types of workers are available, fulltime and part time workers, with no reference to difference in skill. The different shifts are already given by the restaurant managers and the task is to combine them into a tour. The task assignment aspect is lacking in this article.

In both articles the main objective is to minimize both overstaffing and understaffing, which will both have economic consequences for the fast food chain. This is done by reducing or increasing the workforce. For a problem with a fixed workforce, such as ours, this objective function is not relevant. In the example studied by Loucks and Jacobs there is also a goal to meet staff demand on total working hours. This is modeled as a secondary goal which is similar to ours and somehow models a "soft" value, which is of interest to us.

A more recent TSP concern monthly tour scheduling, as opposed to most literature which concerns only weekly scheduling. Such a study was done by Rong (2010). The main advantage of monthly scheduling over shorter time periods, as stated in the article, is the possibility to plan a schedule with respect to fairness and balance over a longer period of time. The problem concerns workers with different skills, where each worker also can possess multiple skills.

This is referred to as a mixed-skill problem. Thus, the problem is similar to our problem, where mixed-skill is also present. In the study, workers have individual weekend-off requirements. The problem does not involve task assignment, which makes it less relevant for us.

The solution methods used to solve the TSP differ greatly between the articles studied. In the older articles, such as Thompson (1988) and Loucks and Jacobs (1991), custom made algorithms very similar to the methods used in manual scheduling are proposed to solve the problem. These solution methods involve classifying staff and distributing them according to some rule (for example, the staff with the most scarce skill is assigned first). General commercial solvers are not proposed, due to their lack of efficiency during these times.

In Hojati and Patil (2011), the same model and data is used as in Loucks and Jacobs (1991). Also this article states that commercial solvers are insufficient for solving the problem. Instead two methods are proposed, one which decomposes the problem into two problems solvable by commercial solvers and one customized heuristic method based on a Lagrangian relaxation method. This method solves the problem with more satisfactory results than in Loucks and Jacobs (1991), as explained in the article. In Choi et al. (2009), a pure integer programming method is used.

2.4 Other similar problems

In this section a couple of other problems, similar to our own, will be described. The focus will not be on the variety of methods used to solve these problem; instead, the focus is to give clarity to how closely related many of these problems are.

2.4.1 Fixed job schedule problem

Variations of the task assignment problem relevant to our problem include for example the Fixed Job Schedule Problem (FJSP). The FJSP has been studied since the 1970s in the context of task assignment in processors. According to Krishnamoorthy et al. (2012), the problem concerns the distribution of tasks with fixed starting and ending times over a workforce with identical skills, such as processing units. Such problems have been solved by Gertsbakh and Stern (1978) and Fischetti et al. (1992).

In the article by Gertsbakh and Stern (1978) a situation is studied where n jobs need to be scheduled over an unlimited number of processors. The objective function of such a problem becomes to minimize the number of machines needed to perform all tasks. Fischetti solves a similar problem, but adds time constraints, saying that no processor is allowed to work for more than a fixed time T during a day as well as a constraint forcing tasks to spread out with time gaps S over a processor.

2.4.2 Tactical fixed interval scheduling problem

The Tactical Fixed Interval Scheduling Problem (TFISP) is a problem very closely related to the SMPTSP, with the sole difference being that the TFISP concerns workers which always are available, such as industrial machines or

processors. The problem is studied in Kroon et al. (1995). A typical TFISP can be expressed using the nomenclature in Table 2.1 and written as PTSP[F;I-U;H;W].

Opposed to the FJSP, the TFISP deals with a heterogeneous workforce. Two different contexts are studied by Kroon et al. (1995). One of them concerns the handling of arriving aircraft passengers at an airport. Two modes of transport from the airplane to the airport are investigated, directly by gate or by bus. The two transportation modes thus correspond to two processing units, which can only handle a number of jobs at the same time.

2.4.3 Operational fixed interval scheduling problem

The Operational Fixed Interval Scheduling Problem (OFISP) is a close relative to the TFISP. Both OFISP and TFISP are restricted by the following: Each machine (worker) cannot handle more than one job at a time, it can only handle a subset of the jobs and preemption of jobs is not allowed. The difference between OFISP and TFISP occurs in the objective function, as described in Kroon et al. (1995). TFISP tries to minimize the number of workers, while OFISP tries to minimize the operational costs and the number of unallocated tasks using priority indices. In the present nomenclature this would give rise to the problem PTSP[F;I;H;U-A].

Given the problem definition above, working shifts are to be created for the workers and tasks have to be allocated on a day-to-day basis. The OFISP, studied by Kroon et al. (1995), can therefore be seen both as a job scheduling problem and a task assignment problem.

2.4.4 Stochastic job problems

What differs mostly between the problem types described above and the problem studied in this thesis work, is the objective function. The main objective function is often to minimize staff for a fixed number of jobs, not taking stand-in assignment into account.

An area where the need for stand-in personnel appears is in the maintenance industry, where some jobs can be foreseen and other (emergency) jobs are of a stochastic nature, that is, there is a probability that such jobs will occur a certain hour. The problem combining both foreseen and stochastic maintenance worker scheduling was studied by Duffuaa and Al-Sultan (1999), as a continuation of the work by Roberts and Escudero (1983).

In the former article, a fixed heterogeneous workforce consisting of electricians, plumbers and mechanics is studied. The shifts of the personnel are predetermined by their given work times and thus the problem becomes a pure task assignment problem. The goal is to maximize the number of planned and unplanned jobs performed by the workers, by taking into account the probability of unplanned work to occur. Thus, certain workers will be left at the station as stand-ins in the case an unplanned job arises. The commercial solver LINDO was used to solve the problem.

2.5 Modelling soft constraints

For most scheduling problems, the main objective is to minimize worker-related costs by reducing the number of workers needed to perform a task, or by reducing the working time for part-time employees. Recently, however, many studies have started to focus more on softer goals such as worker satisfaction as an objective function. Such goals are usually considered when scheduling is done manually, but have been forgotten or set aside in mathematical modelling.

In an article by Akbari et al. (2013) a scheduling problem for part-time workers with different preferences, seniority level and productivity is investigated. In this article, these aspects are reflected in the objective function and weighted against each other. A similar problem was also studied by Mohan (2008), but for a workforce of only part-time workers.

Other factors which may affect worker satisfaction, and in the long run efficiency and presence at work are fatigue, fairness and boredom. These are discussed by Eiselt and Marianov (2008). Repetitiveness of a job as well as the level of challenge can cause boredom of workers. The variance in the schedule is increased by Eiselt and Marianov (2008) through providing an upper bound of how many tasks can be performed in a given time span. The article suggests a sort of measurement of the distance between the task requirements and the worker abilities used. This will then be minimized in the objective function.

Another modelling method which is relevant specifically for scheduling problems featuring soft constraints is fuzzy goal programming. The method is discussed by Shahnazari-Shahrezaei et al. (2013), who model soft constraints as "fuzzy goals". These goals can become contradictory, for example could a preference of high seniority level workers come in conflict with a preference in working hours by an employee. The article uses fuzzy set theory, and a solution approach involving Li's two-phase method (Li, 1990). Soft constraints are modeled as trapezoid functions and an optimal solution with the best average value of all functions is found.

The solution methods proposed by Akbari et al. (2013) to solve a scheduling problem with soft constraints are two metaheuristics: Simulated Annealing (SA) and Variable Neighbourhood Search (VNS). According to Akbari et al. (2013), SA has been studied as a solution method for the scheduling problem since the 1990s and many studies have shown that it is capable of providing near-optimal solutions in a short time compared to optimizing integer programming models, for a variety of problems. An exponential cooling time was used for the algorithm and it was concluded that it was faster than the commercial solver LINGO in finding a solution.

VNS is the other proposed method by Akbari et al. (2013). A big difference between this and other methods is that VNS requires very little parameter tuning, while often providing good solutions. Larger movement sizes are used at higher temperatures, and smaller at lower temperatures. The method uses both a random and a systematic phase. In the random phase, worker schedules are regenerated randomly and better solutions are saved. In the systematic phase, two shifts are swapped. In the article, it was concluded that VNS could solve the problem faster than the implemented SA heuristic.

In the article by Eiselt and Marianov (2008), a commercial solver was used. This was also the case in Mohan (2008), as the article compares these commercial solver results with the results obtained from a branch-and-cut algorithm. Also

Shahnazari-Shahrezaei et al. (2013) finds results using a commercial solver.

2.6 Summary

In order to get an overview over the problems discussed in this chapter it is a good idea to look at general historical trends among the problems. One clear trend is the shift from problems concerning homogeneous to heterogeneous workforces. As stated by Krishnamoorthy et al. (2012), the problem of heterogeneous workforces was trivial at the publishing year, while in Loucks and Jacobs (1991), it is introduced as a rather new and challenging concept. Furthermore, the articles concerning the different PTSPs are mainly from the 1980s and 1990s and were challenging in their structure during these times. Some of them, such as FJSP, are related to the scheduling in processors, which was a hot topic at the start of the computer era. All the articles about modelling of soft constraints are presented after the year 2000, probably as a result of better computational powers.

The various solution methods found in the studied articles include commercial solvers, heuristics and matheuristics. Heuristics which have been studied in the articles include SA, VNS and Lagrangian relaxation, with and without VA. Also some local search-based heuristics have been used. Commercial solvers are used increasingly in newer articles such as Hojati and Patil (2011) and Mohan (2008), probably due to the improvement in performance in such solvers. Similarly, matheuristics are also discussed mostly in more recent articles, such as Akbari et al. (2013).

2.7 Relevance to our problem

As described in Section 2.1, many different types of personnel tasks scheduling problems exist. Only a few of them have been discussed in this chapter, of which some are closer related to our problem than others. In order to get a better understanding for how they are related to our problem, a few connections to them will be presented in this section.

Using Table 2.1, the closest classification of the problem studied in this thesis work would be the PTPS[F;F;H;F]. This describes a Personnel Task Scheduling Problem with fixed tasks, fixed shift lengths, a heterogeneous work force and an empty objective function. As stated in Section 2.1, the main difference lies in the objective function, as ours is not empty. Thus, many problem types discussed in this chapter concern problems which are similar to ours but with a different objective function. This includes the shift minimization personnel task scheduling problem, the tour scheduling problem, the fixed job schedule problem, the tactical fixed interval scheduling problem and the operational fixed interval scheduling problem. The first two are relevant to our problem as they concern task assignment and, in the case of tour scheduling problem, days off scheduling. However, both involve shift scheduling while we have fixed shifts. The other three problems concern only task assignment, but are otherwise further from our problem since the tasks types are different from ours.

As stated in Section 2.1, our problem has an objective function in which the minimal number of stand-in personnel is to be maximized. Such problems

often arise in the maintenance job scheduling problem, as is discussed in Section 2.4.4. As a secondary objective, we are also interested in making the schedule varied and with a repeating pattern over a cycle of five weeks. Modelling of such constraints is discussed in Section 2.5.

Chapter 3

The mathematical model

In this chapter the mathematical model implemented to solve the ten week scheduling problem described in Section 1.1 will be presented. Prior to the objective function and constraints, the most significant sets and variables will be provided. Section 3.2 presents the objective function and gives a short description of what it represents. In Section 3.3 the essential constraints will be presented and explained. A complete model with all definitions and the full set of constraints can be found in Appendix A.

3.1 Set and variable definitions

To solve the problem many sets and variables had to be declared. As mentioned before, there are many unique and individual requirements given by the library that have to be met. An example is that some staff members want a day free from outer tasks to be able to attend meetings or perform inner tasks. Another example is that some staff members have two alternating schedules for odd and even weeks. These specific cases have to be modelled and result in a variety of sets and variable definitions. Below, we focus on the common and general cases in order to keep the presentation simple and easily accessible. A complete list of the definitions can be found in Appendix A.

I	Set of staff members
I_{lib}	Set of librarians ($I_{lib} \subseteq I$)
I_{ass}	Set of assistants ($I_{ass} \subseteq I$)
W	Set of all ten weeks
W_5	Set of first five weeks
D	Set of all days in a week
D_5	Set of all five weekdays
S_d	Set of shifts available day d
S_3	Set of first three shifts on a weekday
J_d	Set of task types available day d

In order to further define the problem we introduce the following variables.

$$x_{i w d s j} = \begin{cases} 1, & \text{if staff member } i \text{ is assigned to a task } j \text{ in week } w, \text{ day } d, \text{ shift } s \\ 0, & \text{otherwise} \end{cases} \quad (3.1)$$

$$H_{i w h} = \begin{cases} 1, & \text{if staff member } i \text{ works weekend } h \text{ in week } w \\ 0, & \text{otherwise} \end{cases} \quad (3.2)$$

$$r_{i w} = \begin{cases} 1, & \text{if staff member } i \text{ has its schedule rotated } w-1 \text{ steps forwards} \\ 0, & \text{otherwise} \end{cases} \quad (3.3)$$

$$l_{i w d} = \begin{cases} 1, & \text{if librarian } i \text{ is a stand-in week } w, \text{ day } d \\ 0, & \text{otherwise} \end{cases} \quad (3.4)$$

$$a_{i w d} = \begin{cases} 1, & \text{if assistant } i \text{ is a stand-in week } w, \text{ day } d \\ 0, & \text{otherwise} \end{cases} \quad (3.5)$$

$$y_{i w d s} = \begin{cases} 1, & \text{if staff member } i \text{ works at task type E, I or P in week } w, \text{ day } d, \text{ shift } s \\ 0, & \text{otherwise} \end{cases} \quad (3.6)$$

$$W_{i w d} = \begin{cases} 1, & \text{if a staff member } i \text{ is working a shift in week } w, \text{ day } d \\ 0, & \text{otherwise} \end{cases} \quad (3.7)$$

$$b_{i w} = \begin{cases} 1, & \text{if staff member } i \text{ works at HB in week } w \\ 0, & \text{otherwise} \end{cases} \quad (3.8)$$

$$f_{i w} = \begin{cases} 1, & \text{if staff member } i \text{ is assigned to work Friday evening week } w \\ 0, & \text{otherwise} \end{cases} \quad (3.9)$$

$$M_{w d s} = \begin{cases} 1, & \text{if a big meeting is placed in week } w, \text{ day } d, \text{ shift } s \\ 0, & \text{otherwise} \end{cases} \quad (3.10)$$

$$m_{w d s D} = \begin{cases} 1, & \text{if a meeting is placed at group } D \text{ in week } w, \text{ day } d, \text{ shift } s \\ 0, & \text{otherwise} \end{cases} \quad (3.11)$$

$$d_{i w d s} = \begin{cases} 1, & \text{if there is a difference for a staff member } i \text{ in assignment of tasks at} \\ & \text{shift } s, \text{ day } d \text{ between weeks } w \text{ and } w+5 \\ 0, & \text{otherwise} \end{cases} \quad (3.12)$$

$$l^{min} = \text{lowest number of stand-in librarians found (integer)} \quad (3.13)$$

$$a^{min} = \text{lowest number of stand-in assistants found (integer)} \quad (3.14)$$

$$s^{min} = \text{weighted sum of numbers of stand-in librarians and assistants} \quad (3.15)$$

$$\delta = \text{total number of shifts that differ for all staff members} \quad (3.16)$$

Based on the variables defined above it has been possible to solve our scheduling problem. The variables l^{min} and a^{min} are the ones of most significance as they represent the number of stand-ins found after a run.

3.2 Objective function

As the problem contains multiple objective functions it has been necessary to weigh them against each other using parameters. These are shown in Equation 3.17 below, as M and N .

$$\text{maximize } M \cdot s^{min} - N \cdot \delta \quad (3.17)$$

The first part of the objective function represent the lowest amount of stand-in librarians and assistants found, while the second part is a preference from the library that two weeks with a five-week interval should be as similar as possible (e.g. week 1 and 6, 2 and 7 etc.). The parameter N prioritizes the similarity of weeks compared to the number of stand-ins. Based on the information given by the library it is a much higher priority to have many stand-ins. Hence, $M \gg N$ holds in our case.

3.3 Constraints

To model the problem it has been of relevance to divide many of the constraints into weekend and weekday constraints. Several auxiliary constraints and variables have also been added to avoid multiplication of two variables, which would have made the problem non-linear. These auxiliary constraints have been left out of this chapter for simplicity reasons. Instead, these can be seen in Appendix A.

3.3.1 Demand and assignment constraints

The most crucial constraint is to ensure that the demand of staff members is met each day. This can be modelled as

$$\sum_{i \in I} x_{i w d s j} = \text{demand}_{w d s j}, \quad w \in W, d \in D, s \in S, j \in J_d. \quad (3.18)$$

Here, $\text{demand}_{w d s j}$ is an integer representing the number of staff members required week w , day d , shift s for a task j .

The following constraint says whether or not a staff member is assigned a task during a weekday, where the evening shifts and task type Library on Wheels are excluded.

$$y_{i w d s} = \sum_{j \in J_d \setminus \{B\}} x_{i w d s j}, \quad i \in I, w \in W, d \in D_5, s \in S_3 \quad (3.19)$$

The variable assignment is used to simplify a couple of constraints later on.

To ensure that no staff member is assigned more than one task the following constraint is implemented.

$$\sum_{s \in S} \sum_{j \in J_d} x_{i w d s j} \leq 1, \quad i \in I, w \in W, d \in D \quad (3.20)$$

However, if we allow a person to have two shifts at the Library on Wheels on a day, Equation 3.20 has to be slightly modified, which is left out of this chapter.

It is preferred to allow only one PL per week and a maximum of three PL per ten weeks. These are easily modelled with the following constraints.

$$\sum_{s \in S_d} \sum_{d \in D} x_{iwsP} \leq 1, \quad i \in I, w \in W \quad (3.21)$$

$$\sum_{w \in W} \sum_{s \in S_d} \sum_{d \in D} x_{iwsP} \leq 3, \quad i \in I \quad (3.22)$$

The long duration of a PL (in the equations named "P") is the reason for these preferences, as some staff members are required to have some time free from outer tasks to be able to perform their inner tasks.

Another preference is to have varying start times of the assigned tasks each week, so that the more and less desired shifts are fairly distributed. The following equation is implemented to meet such requirements.

$$\sum_{d \in D_5} y_{iws} \leq 2, \quad i \in I, w \in W, s \in S_3 \quad (3.23)$$

Equation 3.23 allows a staff member to have at most two tasks starting at the same hour every week. Worth noting is that task type Library on Wheels is disregarded in this constraint as the variable y_{iws} is used; see Equation 3.19 for the definition.

It is desirable to avoid assigning too many tasks to a staff member. The reason for this, as mentioned previously, is to let the staff member have some time to allot for inner services during weekdays. The equation below models this preference.

$$\sum_{d \in D_5} \sum_{s \in S_d} \sum_{j \in J_d} x_{iwsj} \leq 4, \quad i \in I, w \in W \quad (3.24)$$

Equation 3.24 allows a staff member at most four weekday shifts per week. The exception, which is one of the BokB staff members, is left out of this equation for simplicity reasons. To model this completely a new subset of I needs to be created where that staff member is left out.

3.3.2 Weekend and rotation constraints

Every staff member's schedule is able to rotate up to nine times, where the decision variable r_{iv} decides the rotation. Therefore, the parameter $qualavail_{iwsj}$ has to align with the rotation so that all staff members are assigned tasks only when available. This is provided in the following equation.

$$x_{iwsj} \leq \sum_{v \in V} r_{iv} * qualavail_{i(mod_{10}(w-v+10)+1)dsj}, \quad i \in I, w \in W, d \in D, s \in S_d, j \in J_d \quad (3.25)$$

Table 3.1 gives a better understanding of how the modulus function in the parameter $qualavail_{iwsj}$ is used. Here, w represents the current week in the initial schedule and v represents $v-1$ rotations to the right from the initial schedule. With initial schedule we refer to the schedule where all staff members' weekends

occur at the first week. The resulting week, after rotations can be seen in the cells.

Table 3.1: Resulting table of the function $\text{mod}_{10}(w - v + 10) + 1$

		w =									
		1	2	3	4	5	6	7	8	9	10
v =	1	1	2	3	4	5	6	7	8	9	10
	2	10	1	2	3	4	5	6	7	8	9
	3	9	10	1	2	3	4	5	6	7	8
	4	8	9	10	1	2	3	4	5	6	7
	5	7	8	9	10	1	2	3	4	5	6
	6	6	7	8	9	10	1	2	3	4	5
	7	5	6	7	8	9	10	1	2	3	4
	8	4	5	6	7	8	9	10	1	2	3
	9	3	4	5	6	7	8	9	10	1	2
	10	2	3	4	5	6	7	8	9	10	1

An example to better understand the function: imagine that we are looking at an unrotated initial schedule where everyone is assumed to work weekend the first week. Say that we look at a staff member's first week, $w = 1$. If this schedule is rotated two times to the right, then $v = 3$, then the first week in the new rotated schedule represents the previous ninth week, that is $w = 9$ from the old schedule. The availability to look at in the old schedule is, therefore, week nine.

Initially, all staff members are assigned weekend work the first week. Then, a week rotation variable is introduced that rotates the staff members' availability matrices so that the demand constraints on evenings can be met. The most basic constraints regarding which week a staff member's weekend work occurs and its week rotation can be seen below.

$$\sum_{w \in W} r_{iw} = 1, i \in I \quad (3.26)$$

$$\sum_{w \in W} H_{iwh} \leq 1, i \in I, h = 1, 2 \quad (3.27)$$

$$r_{iw} \geq H_{iw1}, i \in I, w \in W \quad (3.28)$$

Equation 3.26 provides all staff members with a rotation of their schedule regardless if they are working weekends or not. Equation 3.27 allows a staff member a maximum of two weekends, for $h = 1$ and $h = 2$, per ten weeks. As mentioned in previous section all weekends occur the first week initially. Once a staff member is assigned a new weekend, the rotation has to align with this new assignment in order to be able to correctly rotate the availability matrix. This is done using Equation 3.26. In case a staff member is never due for weekend work, i.e. $H_{iw1} = 0, i \in I, w \in W$, the rotation r_{iw} is free as its availability matrix then is identical for all weeks.

A staff member is supposed to work weekends with a five-week interval. However, in case there are enough staff members to satisfy the demand on weekends, it may be enough for some to work one weekend per ten weeks. To

avoid problems with the rotation when only the second weekend is assigned to a staff member, the following equation is used.

$$r_{i(mod_{10}(w+4)+1)} \geq H_{iw2}, i \in I, w \in W \quad (3.29)$$

Worth noting is that if a staff member is assigned both weekends, then Equation 3.29 provides the same information as Equation 3.28.

A staff member is supposed to work with the same task both Saturday and Sunday when working a weekend. This can be modelled with the following constraints.

$$\sum_{j \in J_d} x_{iw61j} + \sum_{j \in J_d} x_{iw71j} = 2 * \sum_{h=1}^2 H_{iwh}, i \in I, w \in W \quad (3.30)$$

$$x_{iw61j} = x_{iw71j}, i \in I, w \in W, j \in J_d \quad (3.31)$$

Friday evening is also included in an extension to working Saturday and Sunday, unless the staff member is assigned to HB. It is, however, not a necessity to perform the same task Friday evening as during the weekend, thus Fridays are not included in Equation 3.31. Equation 3.32 below adds Fridays to the weekend.

$$\sum_{j \in J \setminus \{L\}} x_{iw54j} = f_{iw}, i \in I, w \in W \quad (3.32)$$

Here, "L" refers to the task type Library on Wheels. The variable f_{iw} is, as mentioned in the variable declaration, equal to one if and only if a staff member is working weekend as well as not being assigned to HB.

It is of interest to implement a constraint to prevent staff members from being assigned to HB more than once every ten weeks in order to avoid unfairness. This constraint can be modelled as.

$$\sum_{w \in W} \sum_{d=6}^7 x_{iwd1B} \leq 2, i \in I. \quad (3.33)$$

Here, "B" refers to HB. Equation 3.31 together with Equation 3.33 ensure that once a staff member is assigned to HB it will be for two consecutive weekend days, which is the amount of days the last mentioned equation allows every ten weeks. Why this is preferable is described in Section 1.1.1.

3.3.3 Definitional constraints

The variable δ in the objective function is defined by the following equation.

$$\delta = \sum_{i \in I} \sum_{w \in 1..5} \sum_{d \in 1..5} \sum_{s \in 1..3} d_{iws} \quad (3.34)$$

This equation gives an integer value, representing the total number of shift differences that occur for all workers through all weekdays in the library. The variable d_{iws} , defined in Equation 3.36 below, compares the task assignment between two five-week separated weeks.

To calculate the variable s^{min} used in the objective function, 3.17, the following equation is used.

$$s^{min} \leq L \cdot \sum_{i \in I_{lib}} l_{iwd} + A \cdot \sum_{i \in I_{ass}} a_{iwd}, \quad w \in W, d \in D_5 \quad (3.35)$$

Here, l_{iwd} and a_{iwd} are binary variables stating whether a staff member is a stand-in or not on a specific day. Since s^{min} is being maximized in the objective function, it will assume the lowest value of the sum of stand-in librarians and assistants for any day during the ten weeks. Just as for previous equations, auxiliary constraints have been left out for simplicity reasons. Here it is left out how l_{iwd} and a_{iwd} are determined.

If $L < A$ holds in the model, the solver would prioritize assistants over librarians as stand-ins. Librarians are, however, more desired as stand-ins due to their ability to perform all types of tasks. Therefore, it is desired to let $L \geq A$.

As stated in Section 3.2, two weeks with a five-week interval should be as similar as possible. Equation 3.36 together with Equation 3.34 and the objective function equation, 3.17, provides this preference.

$$d_{iwd} = |y_{iwd} - y_{i(w+5)d}|, \quad i \in I, w \in W_5, d \in D_5, s \in S_3 \quad (3.36)$$

The decision variable d_{iwd} states if there is a difference in assignment between two tasks at the same hour and day for the two weeks, w and $w + 5$. The differences occur if, say, an Exp task is assigned on Monday week one at 08:00-10:00 and no task is assigned the same shift Monday week six. Thus, a variable that is minimized in the objective function.

3.3.4 Meeting constraints

At the library there are both library meetings and group meetings, which both occur with a five-week interval. Library meetings are set to take place from 08:00-10:00 on Mondays, whereas group meetings are more freely distributed. A few staff members are not assigned library meetings, since they are needed at the stations to keep the library running. The set I_{big} in the equation below consists of all staff members who are to be assigned library meetings. The constraints modelling library meetings are as follows.

$$\sum_{w \in W_5} M_{w11} = 1 \quad (3.37)$$

$$M_{(w+5)11} = M_{w11}, \quad w \in W_5 \quad (3.38)$$

$$\sum_{s=1}^3 \sum_{j \in J_1 \setminus \{L\}} x_{iwsj} \leq 1 - M_{w11}, \quad i \in I \setminus I_{big}, w \in W \quad (3.39)$$

Equation 3.37 and 3.38 assign two library meetings with a five-week interval during the ten-week scheduling period. Equation 3.39 makes sure that the staff members that do not attend library meetings are not assigned any other task the day of the meeting.

The constraints added to model the group meetings are somewhat similar to the library meeting constraints. Just as for library meetings they take place

two times during the ten weeks with a five-week interval, which is described by Equations 3.40 and 3.41.

$$\sum_{w \in W_5} \sum_{d \in D_5} \sum_{s \in S_3} m_{wdsD} = 1 \quad (3.40)$$

$$m_{(w+5)dsD} = m_{wdsD}, \quad D = 1, \dots, 3, w \in W_5, d \in D_5, s \in S_3 \quad (3.41)$$

$$m_{wdsD} + x_{iwsdj} \leq 1, \quad D = 1, \dots, 3, i \in I_D, w \in W, d \in D_5, s \in S_3, j \in J_d \quad (3.42)$$

$$m_{wdsD} \leq \sum_{v \in V} r_{iv} * qualavail_{i(mod_{10}(w-v+10)+1)dsE}, \quad D = 1, \dots, 3, i \in I_D, w \in W, d \in D_5, s \in S_3 \quad (3.43)$$

Equation 3.42 prohibits a staff member from multitasking, that is, to attend a meeting and work with a task simultaneously. Lastly, Equation 3.43 enables group meetings only when everyone in that group is available, The rotation is also taken into the account. "E" in Equation 3.43 stands for Exp and is used due to it is the only task type everyone is available for and *qualavail* is dependent on the variable *j*.

Chapter 4

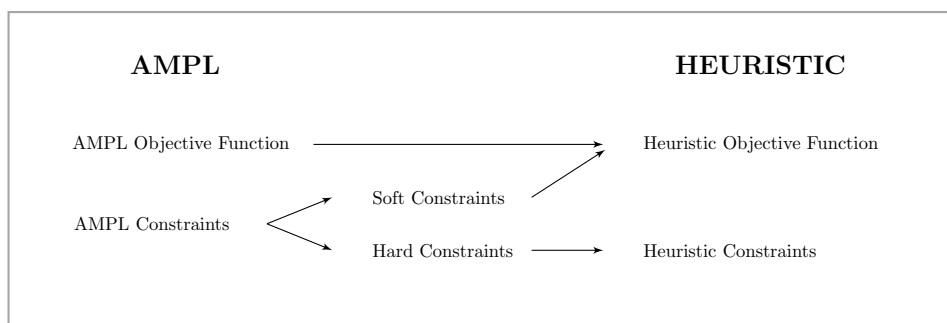
Two heuristic solution methods

The first heuristic is focusing on creating a large pool of unique week appearances. In this pool, there should be one block for every existing combination of tasks for any of the staff members. After the pool creation, each staff member is assigned one block from this pool for every week that is scheduled for.

The second heuristic initially assigns the week rotations to the staff members. Afterwards, the tasks are allocated to the staff members using a greedy heuristic until a feasible solution is found.

In order to implement the heuristics, some of the mathematical model's constraints are softened, relaxed or removed. It is due to lack of time a few constraints have been removed as they would require time to implement which we did not have. Figure 4.1 illustrates the process where some constraints which are considered to be hard in the original model are softened when the heuristics are applied. The changes of the model for the two heuristics are described in Table 4.1 and 4.2.

Figure 4.1: Illustration of the differences between the model used in AMPL and the heuristic model



It was natural to remove the meetings and shorten the scheduling period to five weeks instead of ten in the first implementation. Meetings are not in the

Table 4.1: Showing changes in the model for the week block scheduling approach. The soft, relaxed and removed constraints are shown.

Soft Constraints	
Affected constraints	Constraint description
3.18	The number of staff members needed for every shift and task type in the library.
3.22	Maximum number of PL per staff member and week cycle.
Relaxed Constraints	
Affected constraints	Constraint relaxation
Several. $W = W_5$ in all constraints.	Five week scheduling instead of ten week scheduling.
BokB-constraints	BokB manually assigned, due to low degree of freedom.
Availability data	Even and odd week staff members have availability at each shift according to the stricter of the two sets.
Removed Constraints	
Affected constraints	Constraint description
3.36	Any two weeks w and $w+5$ shall be as similar as possible.
3.37 - 3.43	Meetings are not implemented.

library's current schedule, as it was more of a request if we could take them into account when creating a schedule. A ten week cycle is, more or less, a copy of five week scheduling cycles. The difference occur for those who have varying work hours during even and odd weeks. Therefore, the scheduling period was shortened to five weeks.

Whenever constraint violations occur in the soft constraints, mentioned in Table 4.1 and 4.2, a cost is added to the objective function. Therefore, a feasible solution is the same as having no constraint violations in the soft constraints.

4.1 Week block scheduling approach

The first heuristic approach creates initially a large pool of week block appearances. These are then filtered for every staff member based in its qualities, such as availability and qualification. Then, a large neighbourhood search is applied by alternating between a destroy and repair on a given number of staff members. A critical step during a destroy and repair phase is to be able to swap rotations between staff members, so that they are scheduled for weekend work another week.

The implemented heuristic can be seen in Figure 4.2. The LNS framework is shown in the middle of the figure.

If a run does not find a feasible solution after a given number of iterations, the run is discarded and a new iteration is executed. It is shown on the right side of Figure 4.2. In case a run is close enough to enter the final phase, a check is made whether there are enough stand-ins to be assigned the remaining tasks.

Table 4.2: Showing changes in the model for the task distribution approach. The soft, relaxed and removed constraints are shown.

Soft Constraints	
Affected constraints	Constraint description
3.20	At most one task per day.
3.24	At most four tasks per week.
3.21	At most one PL per week.
3.22	Maximal number of PL per ten weeks.
3.23	Not more than two tasks at the same shift in a week restriction.
Relaxed Constraints	
Affected constraints	Constraint relaxation
Several. $W = W_5$ in all constraints.	Five week scheduling instead of ten week scheduling.
BokB-constraints	BokB placed according to constraints, but the schedule is fixed.
Availability data	Even and odd week staff members have availability at each shift according to the stricter of the two sets.
Removed Constraints	
Affected constraints	Constraint description
3.36	Any two weeks w and $w+5$ shall be as similar as possible.
3.37 - 3.43	Meetings are not implemented.

If not, all staff members will be assigned empty week blocks, so that a new try of finding a feasible solution can be made.

Every staff member's information such as availability and qualification is inserted into an Excel table. It is then written to a text file using a Visual Basic code, which in turn is read and used by the heuristic.

4.1.1 Block creation

A major part of this heuristic is to create the large pool of unique week appearances. These are then filtered for each of the staff member based on its availability. The staff members' availability are generalized into three categories: weekend week, week rest week and weekday week, where weekday week occurs three times during a five week period, see Table 1.4.

One week block contains seven days and up to four shifts where each shift can contain at most three different tasks. Table 4.3 below is a representation of a general week block with all possible tasks for each day d and shift s . In the table, I represents that "No task" is assigned that day, D represents "Desk task" meaning either Exp or Info, PL represents "Fetch list" and HB represents "Hageby".

Every day must contain exactly *one* task from either of the four shifts when creating a week block. The tasks I and PL are ranging over more than one shift. The duration of a PL is three shift and I refers to the entire day. They are both placed in the first shift to simplify the complete task representation. When

Figure 4.2: A flow chart of the implemented heuristic with week block scheduling

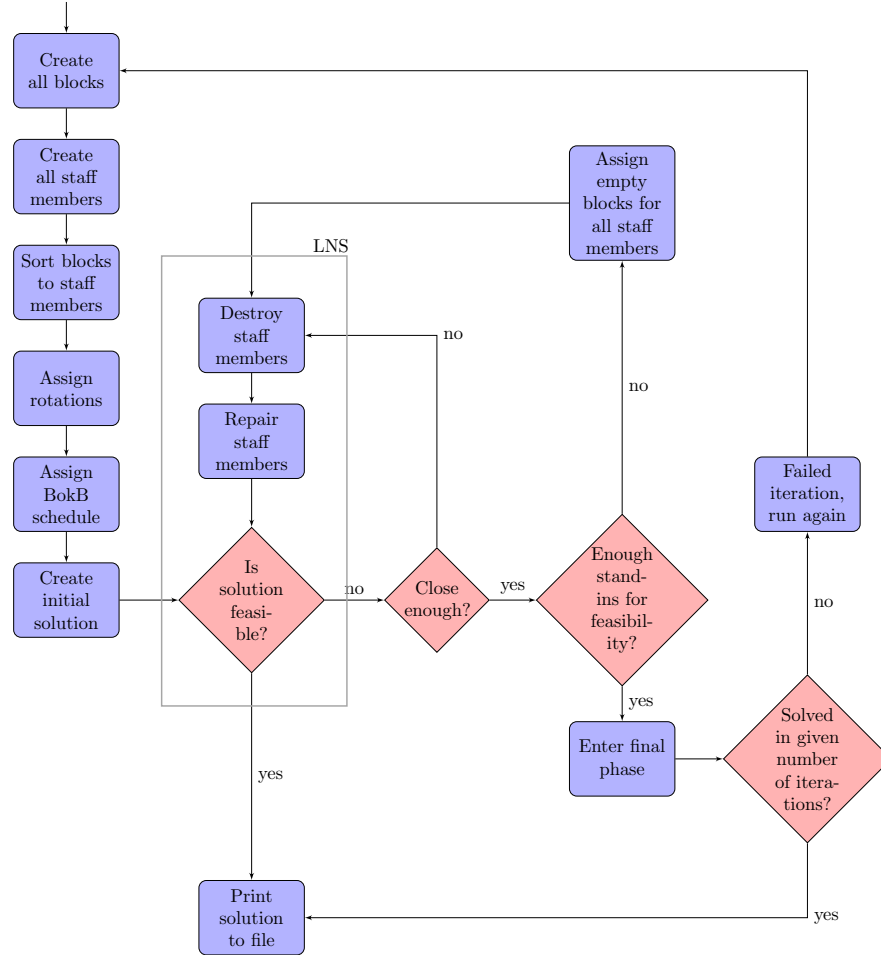


Table 4.3: A general week block with all possible tasks

	Mon	Tue	Wed	Thu	Fri	Sat	Sun
08:00-10:00	I,D,PL	I,D,PL	I,D,PL	I,D,PL	I,D,PL	I,D,HB	I,D,HB
10:00-13:00	D	D	D	D	D		
13:00-16:00	D	D	D	D	D		
16:00-20:00	D	D	D	D	D		

creating the combinations of block appearances there are additional conditions that have to be met. These are:

1. At most two tasks per week can be assigned the same shift.
2. At most one evening is allowed per week, if Friday evenings are excluded.
3. At most one PL is allowed in a week block.
4. Saturday and Sunday shall always contain the same type of task.

5. If Saturday and Sunday contain Desk tasks, then so shall Friday afternoon (fourth shift).
6. No more than four tasks are allowed during the weekdays, leaving at least one day free from tasks.

Too illustrate the growth of the number of blocks when more tasks are added, consider the following example. If items 1, 2, 3, 5 and 6 are disregarded there exists $6^5 * 3 = 23,328$ unique week blocks. In contrast, if Exp and Info were to be considered separately, instead of the combination of the two, the possible combinations would be $10^5 * 4 = 400,000$. By applying all conditions above the total number of unique block appearances for this implementation are only 4,175.

An illustration of one of the 4,175 existing blocks can be seen in Table 4.4 below.

Table 4.4: Illustration of one of the unique block appearances

	Mon	Tue	Wed	Thu	Fri	Sat	Sun
08:00-10:00		PL	I		I	HB	HB
10:00-13:00							
13:00-16:00				D			
16:00-20:00	D						

This block contains five tasks; two of them are weekend tasks and three are weekday tasks. Which weeks that this block can be assigned to is dependent on the staff member's rotation. Since Hageby is assigned to the week block, one can conclude that only a librarian can have this block assigned to itself. Due to this fact, the Desk tasks can mean either Exp or Info desk work, as librarians are qualified for both.

4.1.2 Block filtering

After creating all existing week appearances they are filtered for each of the staff members based on their availability in each of the three categories mentioned in Section 4.1.1. Table ?? in Appendix ?? shows the results with respect to the number of existing blocks after this filtering has been made for all staff members.

All of the values in the table are fractions of the total number of 4,175 blocks. Knowing the structure of the problem, one can deduce that the number of available week rest blocks are always less than or equal to the available weekday blocks, since the only difference between the two mentioned block categories is when a staff member is free from work due to its week rest. Hence, they are equally many when the staff member is never working weekends.

Looking at a typical staff member's availability, shown in Table 4.5, one might think that there shall be more weekend blocks available than weekday blocks, as the availability, almost in all cases, is higher for weekend blocks. This is not the case, however, as all blocks without weekend tasks are filtered out from the weekend block category. This means that all combinations with "No task" on weekends are excluded in that category. Furthermore, the case shown in Table 4.6 decreases the number of existing weekend blocks further. This case

Table 4.5: Typical availability for a five-week period for a staff member. Yellow signifies that the staff member is available during the hours. In parenthesis, the weekend shift hours are given.

Weekend week	Mon	Tue	Wed	Thu	Fri	Sat	Sun
08:00-10:00 (11:00-16:00)							
10:00-13:00							
13:00-16:00							
16:00-20:00							
Week rest week	Mon	Tue	Wed	Thu	Fri	Sat	Sun
08:00-10:00 (11:00-16:00)							
10:00-13:00							
13:00-16:00							
16:00-20:00							
Weekday week	Mon	Tue	Wed	Thu	Fri	Sat	Sun
08:00-10:00 (11:00-16:00)							
10:00-13:00							
13:00-16:00							
16:00-20:00							
Weekday week	Mon	Tue	Wed	Thu	Fri	Sat	Sun
08:00-10:00 (11:00-16:00)							
10:00-13:00							
13:00-16:00							
16:00-20:00							
Weekday week	Mon	Tue	Wed	Thu	Fri	Sat	Sun
08:00-10:00 (11:00-16:00)							
10:00-13:00							
13:00-16:00							
16:00-20:00							

occurs when a staff member is assigned weekend Desk tasks and, therefore, can not be assigned any other task that Friday, as only one task is allowed per day.

Table 4.6: A weekend block with Desk tasks preventing any other tasks on Fridays.

	Mon	Tue	Wed	Thu	Fri	Sat	Sun
08:00-10:00	X	X	X	X		D	D
10:00-13:00							
13:00-16:00							
16:00-20:00					D		

The notion "X" in Table 4.6 can represent any task, and shifts colored in black mean that no other task can be assigned that shift.

4.1.3 Rotation assignment

The rotation assignment is an early step in the destroy and repair process. Every worker have a rotation assigned to them stating which week the weekend work

occurs. In the rotation assignment this rotation is assigned to a worker based on a few constraints. After a destroy and repair iteration the three randomly chosen workers can have had their rotation swapped. The constraints and an illustration of a swap will be presented here.

There are 35 weekend staff members available, out of which 21 are librarians and 14 are assistants. The demand for weekend staff members each week is seven, that is, the demand for five weeks is exactly $7 * 5 = 35$ staff members. Another requirement, in addition to the seven staff members each weekend, is that at least four of them have to be librarians, due to three librarians are needed in the Information desks and one in Hageby. Because of the exact balance between supply and demand, it seems reasonable to swap rotations between staff members in the destroy and repair loop, so that there always are enough qualified weekend workers every week.

The swaps are done using a random generator in the assignment of rotations such that the above described weekend staffing requirements are always met. However, all of the BokB-staff members have fixed weekends and hence are not given new rotations in the destroy and repair loop. This is more thoroughly described in Section 4.1.4 below.

Table 4.7 shows a destroy and repair iteration regarding rotation assignments. What happens is that the three workers will be able to be assigned a new weekend week, as long as the constraints above are met.

Table 4.7: An iteration in the destroy and repair iteration showing a swap of weekends when three staff members' week rotations are destroyed.

Initial assignment:					
Week	1	2	3	4	5
Librarians	4	4	5	4	4
Assistants	3	3	2	3	3
After destroy:					
Week	1	2	3	4	5
Librarians	3	4	4	4	4
Assistants	3	2	2	3	3
After repair:					
Week	1	2	3	4	5
Librarians	4	5	4	4	4
Assistants	3	2	3	3	3

Yellow indicates that a staff member, either librarian or assistant, has had its week rotation destroyed that week and green indicates that a staff member's week rotation has been repaired. Comparing the "Initial assignment" with "After repair", a swap can be seen between week 2 and 3. Worth noting is that a swap can occur, even if the number of qualified staff members remains the same after a repair. To understand this, consider the case when two librarians with different rotations have them destroyed. Then two cases can occur; either they are assigned the same rotation as before or they swap.

4.1.4 Assignment of Library on Wheels

In order to avoid creating enormous number of block combinations, BokB tasks are assigned manually. This also leads to a fix week rotation for the five BokB-

staff members, and it slightly reduces the degrees of freedom of the problem as the BokB-staff members' rotation will remain unchanged. However, as there are only a couple of different sets of feasible BokB assignments it shall not be the deciding factor for the quality of the solution.

Without a manual assignment of BokB, the number of possible tasks in a week block would increase significantly. The number of tasks would increase from 36 to 43, as there are seven BokB tasks during a week, which would result in a lot more than 4,175 unique week appearances.

4.1.5 Initial solution

The initial solution is created in a similar fashion as the repair function of this heuristic. Based on a greedy heuristic the best week block for a random staff member and week is found and inserted. This is done until every staff member has one weekend block, one week rest block and three weekday blocks assigned.

In order to find the best week block, several costs are introduced to measure whether a block is good or bad to assign to a staff member. Say, if the library demands two librarians at the Information desk on Monday at 08:00 and currently there is only one assigned, then it would be good to assign another one. Such an assignment will, therefore, be rewarded using a cost contribution. Good assignments will provide negative costs and bad assignments will provide positive costs to the objective function value.

Table 4.8 together with Figure 4.3 show an incremental change where different cost contributions have to be considered when evaluating a block before assigning it to a staff member. In order to calculate demand costs for the PL, assume that the block which is being evaluated is to be inserted for week three.

Table 4.8: A block example to be evaluated using costs

	Mon	Tue	Wed	Thu	Fri	Sat	Sun
08:00-10:00	PL	I		I		I	I
10:00-13:00					D		
13:00-16:00			D				
16:00-20:00							

Task to be evaluated

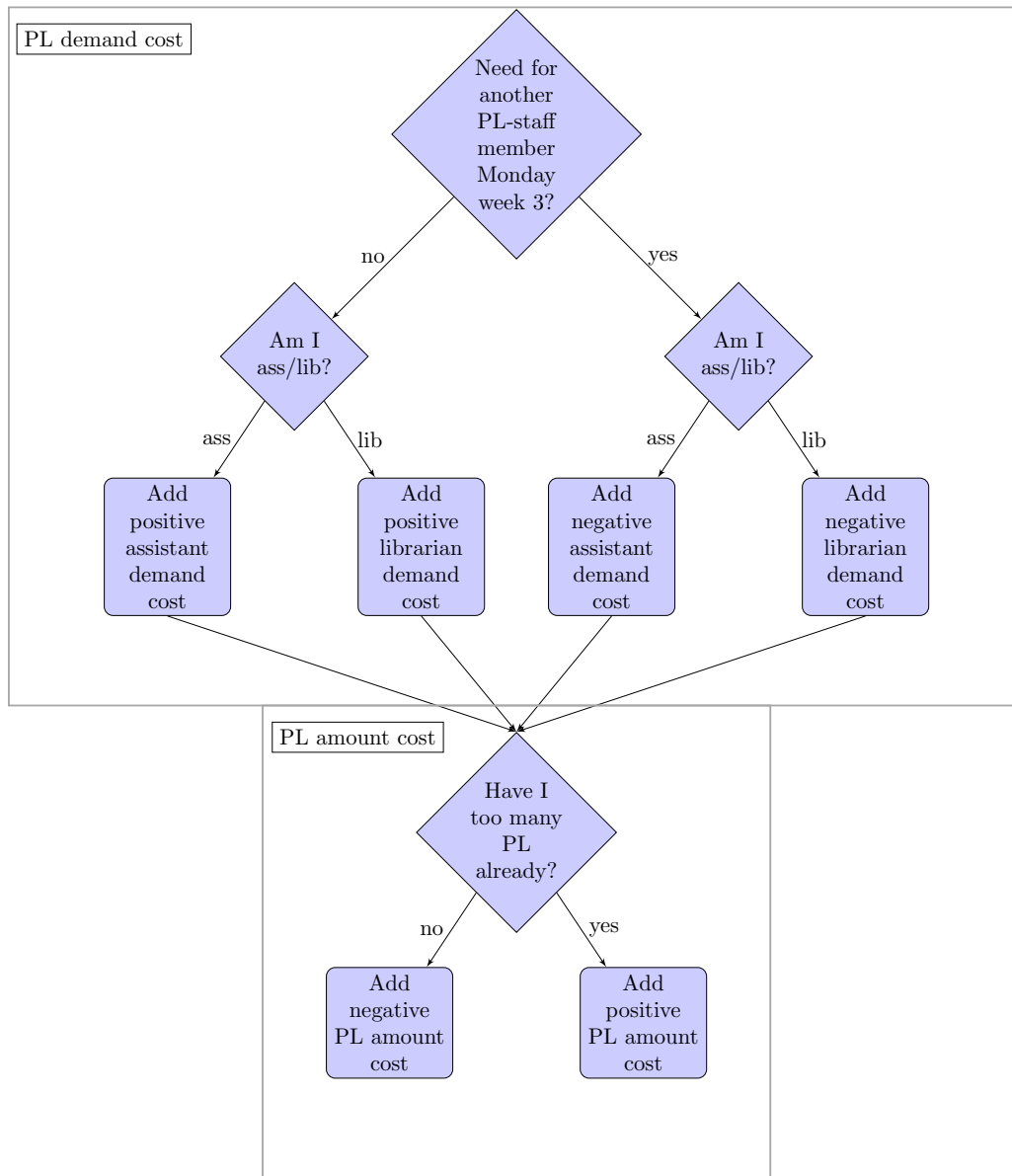
	Mon
08:00-10:00	PL
10:00-13:00	
13:00-16:00	

4.1.6 Costs

In order to find a feasible solution the implemented costs must be carefully chosen. Presently, there are 16 different costs, where some of them are correlated with each other.

Six of the costs are presented in Figure 4.3. They have all been assigned unique values so that, for example, the positive assistant demand cost differs from the negative assistant demand cost in absolute value. To explain why, consider a case where four staff members are demanded, out of which two have

Figure 4.3: A flow chart of costs appearing when a PL is assigned a block on a Monday, third week relative to the library schedule.



to be librarians, see Table 4.9. The first case, where three assistants and one librarian have been assigned is infeasible, as there is a need of two librarians at the Information desks. The second case, where one assistant and three librarians have been assigned is feasible, however, not optimal due to the exceeding use of librarians. The optimal solution is, therefore, Case 3.

The complete list of the 16 costs with descriptions can be seen in Table 4.10.

Whenever a new week block is to be inserted in a repair all available blocks for that staff member have to be given costs. The cost will be based on the

Table 4.9: Library demand at a shift and solution qualities.

Demand:	4 staff members (≥ 2 librarians)	
Case 1:	3 assistants, 1 librarian	(infeasible)
Case 2:	1 assistant, 3 librarians	(feasible)
Case 3:	2 assistants, 2 librarians	(optimal)

Table 4.10: List of all costs with description.

Demand costs	
Cost name	Description
Demand_few_ass	In need of more assistants to fill the demand.
Demand_few_lib	In need of more librarians to fill the demand.
Demand_many_ass	More assistants assigned than needed.
Demand_many_lib	More librarians assigned than optimal.
Demand_few_total	Incorrect number of staff members assigned a task.
Demand_many_total	Incorrect number of staff members assigned a task.
Demand_evening_cost	Incorrect number of staff members assigned an evening task.
Demand_PL_good_ass	Assigning a PL to an assistant, when not already assigned to another.
Demand_PL_good_lib	Assigning a PL to a librarian, when not already assigned to another.
Demand_PL_bad_ass	Assigning a PL to an assistant, when already assigned to another.
Demand_PL_bad_lib	Assigning a PL to a librarian, when already assigned to another.
PL costs	
Cost name	Description
PL_good_amount	Assigning a PL to a staff member still in need of more for feasibility.
PL_violate_amount	Assigning more PL than allowed to that staff member.
Weekend costs	
Cost name	Description
HB_amount	Either no or more than one HB staff member assigned the same weekend.
No_weekend	No weekend blocks available for assignment.
Stand-in costs	
Cost name	Description
Stand_in_cost	Occuring when a possible stand-in is ruined due to an assignment.

current number of staff members assigned to the tasks in the library. The block with the lowest total cost will be inserted in the repair. An illustrative example can be seen in Table 4.11.

4.1.7 Destroy

The destroy consists of two contiguous steps: Choosing three staff members and then destroying their assigned blocks. The BokB tasks remain unchanged even if a staff member assigned to those tasks is destroyed.

4.1.8 Repair

Initially in the repair function the destroyed staff members are assigned new rotations, see Table 4.7, followed by five new week blocks. Which order the blocks are inserted in is randomly generated. For each iteration, one of the three

Table 4.11: Cost evaluation for two blocks. The lower the total cost is, the more desired it is to insert in the repair.

	Mon	Tue	Wed	Thu	Fri	Sat	Sun
08:00-10:00	PL	I		I		I	I
10:00-13:00					D		
13:00-16:00			D				
16:00-20:00							
$Totalcost = -Demand_PL_good_lib + PL_violate_amount +$ $Demand_many_lib - Demand_few_lib = 400$							
	Mon	Tue	Wed	Thu	Fri	Sat	Sun
08:00-10:00	I	D	I		D	I	I
10:00-13:00							
13:00-16:00							
16:00-20:00				D			
$Totalcost = -Demand_few_lib - Demand_evening_cost$ $+Demand_many_lib = -20,260$							

staff members is randomly chosen, followed by one block that has not already been reassigned, until all 15 destroyed weeks have been repaired. Which week block that is chosen is based on the costs listed in Table 4.10. An illustrative example is shown in Table 4.11.

4.1.9 Evaluation of solution

After each destroy and repair loop the new solution is evaluated. It is common for the new solution to be worse than the previous one. However, all new solutions are accepted regardless of this.

All costs are regarded in the evaluation of the solution. If there are any demand violations in the library or in the number of PL assigned to a staff member these are given the respective costs. Furthermore, in case there are stand-ins these are provided with a negative cost.

4.1.10 Final phase

At the end of a run there are mostly a few infeasibilities left that do not get solved. To fix this, a final phase was implemented. Whenever a run is "relatively close" and it is likely to be able to find a feasible solution the final phase is initiated. "Relatively close" means in this case that the only infeasibility left is the lack of staff members at a handful of weekday shifts. It is likely to be able to find a solution when there are more stand-ins available than demand violations for all days.

The final phase destroys one random week for a random staff member and immediately repairs it, resulting in the same or better solution. This is repeated until a feasible solution is found.

With well chosen cost parameter values the run will always eventually get relatively close to feasibility and reach the final phase. If there are fewer stand-ins available a day where infeasibility occurs, then the solution will be discarded in the final phase.

4.2 Task distribution approach

The second heuristic approach tested distributes individual tasks to staff members and greatly resembles the process of manual scheduling. The implemented algorithm undergoes two phases. In the first phase, weekends, as well as some weekday tasks, are distributed and in the second phase the rest of the weekday tasks are distributed.

The primary method used in this approach is a Large Neighbourhood Search (LNS) together with a Simulated Annealing (SA) accept function. Destroying and repairing the solution, as is customary in LNS, helps leading the solution out of local optima or plateaus. Similarly, SA is used in order to allow the solution to move in a less favorable directions to move out of local optima regions.

4.2.1 Costs

The search algorithm is guided using costs, both for staff member schedules and the whole library schedule. Each staff member has a number of costs, as displayed in Table 4.12. Each of these costs represents a badly placed task, resulting in an infeasible schedule or unused stand-in potential. The costs have different weights, as illustrated in the table. How the weights should be set is discussed in Chapter 5.

For the library schedule, costs are weighed together in three different ways. These are referred to as the three library objective functions, as is illustrated in Table 4.13. These costs partly overlap with the individual staff member costs.

Table 4.12: Individual staff member costs and cost weights.

Stand-in cost		
Cost	Weight	Cost description
C_{SI}	W_{w_SI}	Cost from having a task on a day where the staff member can be a stand-in.
Infeasibility costs		
Cost	Weight	Cost description
C_{Task_D}	$W_{w_Task_D}$	Cost from all tasks exceeding the daily limit.
C_{Task_W}	$W_{w_Task_W}$	Cost from all tasks exceeding the weekly limit.
C_{PL_D}	$W_{w_PL_W}$	Cost from all PL exceeding the weekly limit.
C_{PL_Tot}	$W_{w_PL_Tot}$	Cost from all PL exceeding the total limit.
C_{SShift_W}	$W_{w_SShift_W}$	Cost from the number of task performed at the same shift in a week exceeding the weekly limit.

The weekend objective function, as referred to in the table, is associated with the weekend distribution phase of the problem. There are three types of costs in this objective function, each measuring a certain aspect of a good weekend schedule. Since we want to increase the number of stand-ins at the most critical days, the worst shift or day of these aspects is to be maximized. This is referred to as the min cost. Also, the average of each aspect is considered, in order to distinguish solutions for which the worst shift or day is identical. All costs in this objective function are calculated using the formula:

Table 4.13: Library objective functions with their cost components and weights.

Weekend Objective Function		
Cost	Weight	Cost description
C_{SI_m}	W_{SI_m}	Cost of day with minimum number of stand-ins.
C_{S_m}	W_{S_m}	Cost of shift with minimum number of staff members.
C_{D_m}	W_{D_m}	Cost of day with minimum number of staff members.
C_{SI_a}	W_{SI_a}	Average number of stand-ins at a day.
C_{S_a}	W_{S_a}	Average number of staff members at a shift.
C_{D_a}	W_{D_a}	Average number of staff members at a day.
Staff member Objective Function		
Cost	Weight	Cost description
C_{Task_D}	W_{Task_D}	Cost from all tasks exceeding the daily limit for all staff members.
C_{Task_W}	W_{Task_W}	Cost from all tasks exceeding the weekly limit for all staff members.
C_{PL_W}	W_{PL_W}	Cost from all PL exceeding the weekly limit for all staff members.
C_{PL_Tot}	W_{PL_Tot}	Cost from all PL exceeding the total limit for all staff members.
C_{SShift_W}	W_{SShift_W}	Cost from the number of task performed at the same shift in a week exceeding the weekly limit for all staff members.
Weekday Objective Function		
Cost	Weight	Cost description
C_{SI_m}	-	Cost of day with minimum number of stand-ins.

$$C_{type} = W_{lib} * num_of_lib + W_{ass} * num_of_ass \quad (4.1)$$

The weights W_{lib} and W_{ass} are constants used in all weekend objective function costs. Typically W_{lib} is larger than W_{ass} since librarians can perform a greater number of tasks. The weekend objective function becomes the sum of the weighted costs displayed in Table 4.13.

After a certain number of iteration the weekend phase is finished and the weekday distribution phase is entered. Here, the staff member objective function is used to identify infeasibilities in the staff member schedules and is simply a sum of the staff member costs for all staff members. When this cost is zero, the schedule is feasible and the weekday objective function value is calculated. This objective function contains the minimum stand-in cost, which is calculated in the same way as in the weekend objective function. The weekday objective function corresponds to the objective function in the Mathematical Model and is to be maximized.

4.2.2 Weekend phase

The weekend phase is the first phase the algorithm enters and its flow is described also in Appendix B Figure B.1. The LNS component is part of the "destroy and repair loop", while the simulated annealing step is marked as the decision to accept a solution which is worse than the current one.

The reason for implementing the weekend phase separate from the distribution of other tasks is because of the fact that the number of stand-ins in the final schedule depends to a large extent on the weekend-staff member constellation. In particular, it is the week rest following upon weekend work which reduces the number of staff members at a particular shift or day.

The outer loop in the flow chart is performed a specified number of iterations. This loop is guided by the weekend objective function described in the previous subsection. Better solutions are always accepted and the globally best solution is saved. The SA component is implemented using exponential cooling, so that a solution which is worse than the current one is accepted with a probability P :

$$P = \exp(-\Delta E/T) \quad (4.2)$$

where

$$\Delta E = \text{Wend}_t - \text{Wend}_{t-1} \quad (4.3)$$

and

$$T = T_0 \alpha^t \quad (4.4)$$

T is the temperature of the SA accept function, which has an initial value of T_0 and which cools down with the iteration count t at a rate α ($0 < \alpha < 1$). The objective function value at iteration t is written Wend_t .

The inner loop of the algorithm destroys and repairs the solution. The loop checks if the produced solution is infeasible, that is if the difference between the number of available staff members at each shift and the number of required staff members is negative. Such schedules are discarded and a new destroy and repair is performed. The number of available staff members is calculated per shift and is the sum of all staff members who are available at that shift and who do not have any tasks scheduled at that day.

The algorithms for distributing, destroying and repairing weekends is simply a random function since it is very hard to predict what is a good or bad placement of a specific weekend without placing the whole schedule. Since the weekend objective function is only a measurement of the current number of stand-ins and available staff members, this number will be greatly reduced when placing the weekday tasks.

In order to get a better measurement of the number of available staff members during this phase, evening tasks and BokB tasks are distributed after each weekend repair. The evening tasks are distributed using the same method as is described in Section 4.2.3 while BokB tasks are placed according to a fixed schedule. The effect of this distribution is visible in Tables 4.14 and 4.15. High intensity of red in a cell indicates a higher number of staff members. It can be seen that the first table is generally redder and thus there are fewer available staff members in the second table. This state reflects better the real number of available staff members for a certain weekend constellation.

Table 4.14: Staff member availability placing only weekends. Intensity of red indicates number of staff members.

Num available assistants							
	Mo	Tu	We	Th	Fr	Sa	Su
Shift 1:	8	10	10	10	6	0	0
Shift 2:	8	9	9	9	6	0	0
Shift 3:	9	8	9	7	5	0	0
Shift 4:	3	2	2	3	0	0	0
Num available librarians							
	Mo	Tu	We	Th	Fr	Sa	Su
Shift 1:	16	15	16	13	12	0	0
Shift 2:	18	15	17	14	13	0	0
Shift 3:	17	14	18	18	13	0	0
Shift 4:	3	4	4	3	1	0	0
Num available BokB-librarians							
	Mo	Tu	We	Th	Fr	Sa	Su
Shift 1:	2	0	1	1	1	0	0
Shift 2:	0	0	0	0	0	0	0
Shift 3:	0	0	0	0	0	0	0
Shift 4:	1	0	2	2	0	0	0

Table 4.15: Staff member availability after placing weekends as well as evening tasks and BokB for the same week.

Num available assistants							
	Mo	Tu	We	Th	Fr	Sa	Su
Shift 1:	7	10	9	9	8	0	0
Shift 2:	7	9	8	8	8	0	0
Shift 3:	7	7	8	6	6	0	0
Shift 4:	0	0	0	0	0	0	0
Num available librarians							
	Mo	Tu	We	Th	Fr	Sa	Su
Shift 1:	14	11	13	11	12	0	0
Shift 2:	14	11	14	11	12	0	0
Shift 3:	13	11	14	13	13	0	0
Shift 4:	0	0	0	1	1	0	0
Num available BokB-librarians							
	Mo	Tu	We	Th	Fr	Sa	Su
Shift 1:	0	0	0	0	0	0	0
Shift 2:	0	0	0	0	0	0	0
Shift 3:	0	0	0	0	0	0	0
Shift 4:	0	0	1	0	0	0	0

4.2.3 Weekday phase

When entering the weekday phase, all weekends, evenings and BokB tasks are already placed. In this phase, the rest of the tasks, referred to as "weekday"

tasks are to be placed. The demand for a week at this point is illustrated in Table 4.16 and is identical for all weeks. The algorithm for distributing weekday tasks is found in Appendix B, Figure B.2, and greatly resembles the algorithm for distributing weekends.

Table 4.16: Staff member demand during a week when entering weekday phase.

	Exp	Info	PL	HB	BokB
Monday					
Shift 1:	2	2	1	0	0
Shift 2:	3	3	0	0	0
Shift 3:	3	3	0	0	0
Shift 4:	0	0	0	0	0
Tuesday					
Shift 1:	2	2	1	0	0
Shift 2:	3	3	0	0	0
Shift 3:	3	3	0	0	0
Shift 4:	0	0	0	0	0
Wednesday					
Shift 1:	2	2	1	0	0
Shift 2:	3	3	0	0	0
Shift 3:	3	3	0	0	0
Shift 4:	0	0	0	0	0
Thursday					
Shift 1:	2	2	1	0	0
Shift 2:	3	3	0	0	0
Shift 3:	3	3	0	0	0
Shift 4:	0	0	0	0	0
Friday					
Shift 1:	2	2	1	0	0
Shift 2:	3	3	0	0	0
Shift 3:	3	3	0	0	0
Shift 4:	0	0	0	0	0
Saturday					
Shift 1:	0	0	0	0	0
Sunday					
Shift 1:	0	0	0	0	0

The outer loop of the weekday phase, or the "weekday task distribution loop", is performed a specified number of iterations and calculates the weekday objective function for each iteration. Better solutions are saved, others are discarded.

The inner "destroy and repair loop" destroys and repairs the solution until the staff member objective function is zero, that is until all staff member schedules are feasible. If the loop cannot find such a schedule during a maximum number of iterations, the whole solution is discarded as infeasible, and the outer loop is entered.

When destroying and repairing tasks, the tasks are first sorted according to qualification requirement, meaning that librarian tasks are placed first and

assistant tasks second. This guarantees that not all librarians will be used up for assistant tasks. Secondly, in each subgroup, the tasks are sorted according to the difference between the number of available staff members and the demand at each task. This makes sure that the most critical tasks are placed first.

The tasks are then staffed one at a time. The process starts with temporarily placing the task at all available staff members. This will generate a cost for each staff member according to Table 4.12. The cheapest staff members are then chosen and permanently placed on the task.

The destroy function identifies the staff member with the highest cost. The worst week of this staff member is then identified and destroyed. Also, a few random staff members are chosen, their number depending on the destroy amount specified, for which the same week is destroyed. This makes it possible for the first staff member to get a better worst week by changing tasks with the other destroyed staff members in the repair function.

Chapter 5

Results and discussion

5.1 Results

In the following sections, results from testing the AMPL implementation and both heuristics will be presented. The parameter values used in the runs will also be given.

5.1.1 AMPL implementation

Implementing the model in AMPL and solving it with CPLEX 12.5 gives the values displayed in Table 5.1. The weights in Equation 3.17 and 3.35 were set to $L = 2$, $A = 1$, $M = 100$ and $N = 1$, thus making librarians twice as valuable as assistants and stand ins 100 times more valuable than no shift changes.

Table 5.1: Results from solving the mathematical model with CPLEX solver.

	Min num stand-ins (lib, ass)	Stand-in cost	Solution time
Result	(3,0)	6	19 min

Only the stand-in part of the objective function is presented, since this value can be used as a benchmark in measuring the performance of the heuristics. The number of shift changes, which is the second part of the objective function in Equation 3.17, is zero.

5.1.2 Week block scheduling approach

As most of the costs are correlated with each other some parameter tuning was required. The final result of the parameter tuning can be seen in Table 5.2. Their respective descriptions can be seen in Table 4.10.

Worth noting is that these values are most certainly not the best possible, as they have mostly been assigned using intuition. However, a few relations were determined by either running some tests or using reasoning. An example is if $PL_good_cost > Demand_PL_good_ass$ or $PL_good_cost > Demand_PL_good_lib$. If one of those holds, all staff members that is able to be assigned another PL will be, leaving the library overstaffed with staff members assigned to PL. This

Table 5.2: List of all costs used in the week block scheduling approach and their respective values.

Demand costs	
Cost name	Cost value
Demand_few_ass	350
Demand_few_lib	300
Demand_many_ass	200
Demand_many_lib	40
Demand_few_total	800
Demand_many_total	700
Demand_evening_cost	20,000
Demand_PL_good_ass	1,200
Demand_PL_good_lib	800
Demand_PL_bad_ass	1,200
Demand_PL_bad_lib	1,600
PL costs	
Cost name	Cost value
PL_good_amount	1,000
PL_violate_amount	1,500
Weekend costs	
Cost name	Cost value
HB_amount	15,000
No_weekend	5,000
Stand-in costs	
Cost name	Cost value
Stand_in_cost	5

is because the net of $-PL_good_cost + Demand_PL_good_ass/lib$ will then be negative, meaning it is preferable to assign yet another PL to a staff member regardless of how many there already are. This leads to an overstaffed library on some days, and lack of staff members on other days.

Table 5.3 shows the success rate of finding a solution during a script execution. A run is considered to fail if no feasible solution is found during the final phase. This script execution is using the cost parameters shown in Table 5.2.

Table 5.3: Number of successful runs

Successful runs	420
Runs (total)	638
%	~66

Table 5.4 shows statistics of how many stand-ins were found after 420 successful runs.

Figure 5.1 shows a plot of the objective function value after destroy and repair iterations. When the value on the y-axis reaches zero, a feasible solution is found.

Every time the value skyrockets, a solution has been discarded in the final phase, meaning empty blocks are assigned to all staff members and once again

Table 5.4: Number of stand-ins found in 420 successful runs.

Stand-in spread	Times
No stand-ins	351
1 assistant	33
1 librarian	36

entering the destroy and repair loop. Destroy and repair iterations are being executed until a solution is likely to be found, that is, there are enough stand-ins available to take care of the demand violations.

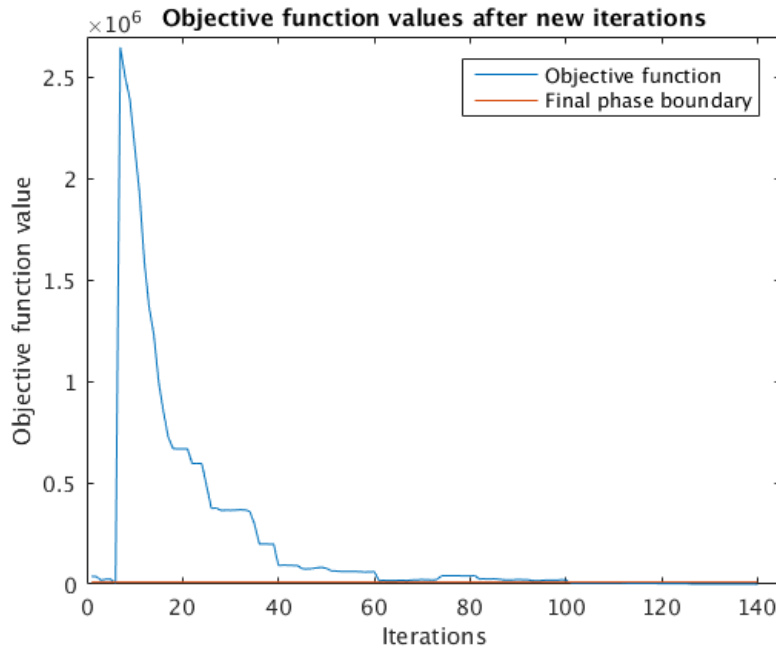


Figure 5.1: Plot of the objective function value after destroy and repair iterations, one time without enough stand-ins when trying to enter the final phase.

As seen in Figure 5.1 there are not enough stand-ins one time during the run. What happens is that the objective function gets below the final phase boundary, but does not meet the stand-in criterion and, therefore, destroys every workers' schedule.

Figure 5.2 illustrates a run when there are enough stand-ins to enter the final phase. Whenever an iteration is below the final phase boundary, the entering criterion is checked. If the final phase is entered, the few demand violations will be taken care of, one by one, by assigning the stand-ins where the violations occur.

To illustrate why the solution gets worse after a destroy and repair iteration; consider the case when two librarians with the same rotation are being repaired, where one can only work at HB during weekends and the other one can work at either HB or the Desk task. If the more versatile librarian is repaired first and being assigned to HB due to the cost parameters, then it will be impossible for

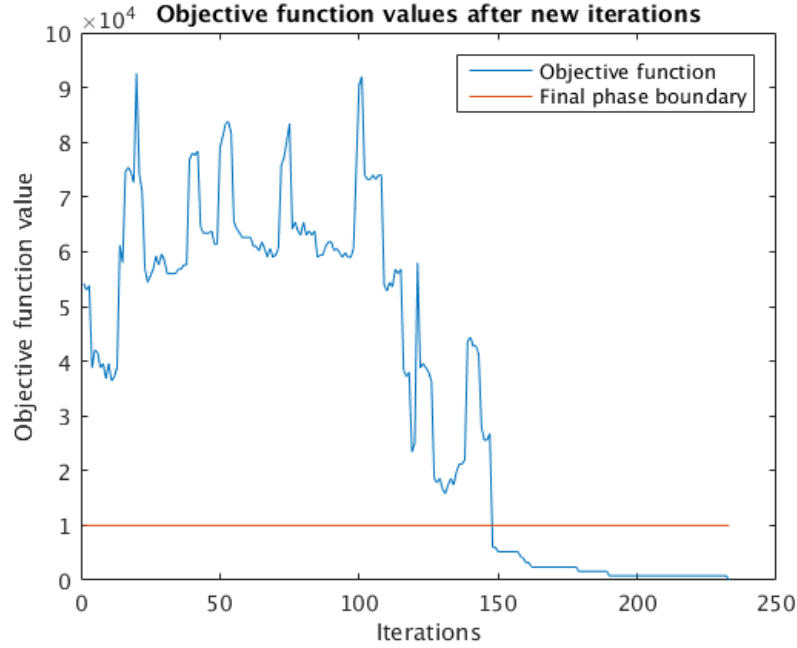


Figure 5.2: Plot of the objective function value after destroy and repair iterations with enough stand-ins when trying to enter the final phase.

the second librarian to be assigned any weekend tasks. This results in a high cost contribution to the objective function.

5.1.3 Task distribution approach

In order to find a good value for the weights described in Section 4.2.1 in the previous chapter, different tuning tests were performed. These tests resulted in the weights presented in Table 5.5. These weights are used in all test in this section unless stated otherwise.

The first six weights, belonging to the weekend objective function, were most experimented with and will be discussed further in this section. The first three, the min weights, are scaled relatively to each other, while the latter three are simple calculated as one tenth of their equivalent min weights. The weights in Equation 4.1, W_{lib} and W_{ass} were set to 2 and 1 respectively since it can be argued that one librarian can perform twice as many tasks as one assistant.

The individual staff member cost and the staff member objective function costs measuring the same aspect of the schedule have been given the same weights, as is illustrated in Table 5.5. This is due to the fact that the weights, relative to each other, decide what constraints are most strict. It is, for example more severe to break the rule of one task per day, than to break the rule of having too many shifts at the same time in a week, which is reflected in the weights. The stand-in weight W_{w_SI} is twice as high for librarians as for assistant.

The two phases of the implementation are run a specified number of iterations, referred to as It_{wend} and It_{wday} . These parameters were also trimmed

Table 5.5: Weights used in the implementation.

Weight	Weight value
Weekend Objective Function Weights	
W_{SI_m}	0.1
W_{S_m}	0.1
W_{D_m}	10
W_{SI_a}	0.01
W_{S_a}	0.01
W_{D_a}	1
staff member/staff member Objective Function Weights	
W_{w_SI}	2 if librarian 1 if assistant
$W_{Task_D} / W_{w_Task_D}$	100
$W_{Task_W} / W_{w_Task_W}$	10
$W_{PL_W} / W_{w_PL_W}$	5
$W_{PL_Tot} / W_{w_PL_Tot}$	5
$W_{SShift_W} / W_{w_SShift_W}$	4

and the results from different iteration values are displayed in Tables 5.6 and 5.7. The results for weight tuning are displayed in Table 5.8.

Table 5.6: Results from the task distribution heuristic when varying It_{wend}

It_{wend}/It_{wday}	Heuristic cost	AMPL cost
10/10	4.40	4.78
50/10	5.00	5.34
100/10	5.26	5.52
500/10	5.66	5.78
1000/10	5.54	5.66

Table 5.7: Results from the task distribution heuristic when varying It_{wday}

It_{wend}/It_{wday}	Heuristic cost	AMPL cost
500/5	5.47	5.74
500/10	5.66	5.78
500/15	5.6	5.74
500/20	5.57	5.74

Table 5.8: Results from the task distribution heuristic when varying weights. $It_{wend} = 1000$ and $It_{wday} = 20$ in all runs.

$W_{SI_m}/W_{S_m}/W_{D_m}$	Heuristic cost	AMPL cost
0/0/0	2.85	3.66
1/1/1	5.43	5.54
10/0.1/0.1	5.11	5.36
0.1/10/0.1	5.43	5.60
0.1/0.1/10	5.57	5.80

The tables display two costs for each test, the Heuristic cost and the AMPL cost. These are calculated as the average result of 100 runs. Each run takes between a few seconds and a few minutes to perform. The heuristic cost is the actual result of the heuristic, measuring how well the heuristic solves the problem. The AMPL cost is the result obtained when solving the problem to optimality in CPLEX, using the weekends found in the heuristic. Thus, this value measures how good the weekend phase works. Both figures should be compared to the optimal stand-in value for the problem, which is 6 (Table 5.1).

Studying Table 5.6, it can be seen that the AMPL cost increases with more weekend iterations. However, after 500 iterations, the cost stops to improve and it is actually deteriorating slightly. This is probably indicating that the weekend objective function value can only give a rough measure of a good schedule, and thus the result does not only depend on the number of iterations but also on chance.

Similarly, in Table 5.7, the Heuristic cost increases until a certain threshold, when the solutions are not improving any more. This might indicate that the second phase is not stochastic enough, so the same schedules are found multiple times. This might be solved through implementing varying weights or introducing more randomness in task placement.

Table 5.8 shows a few extreme weight combinations. It is clear from the table that, although the performance does not vary to any larger extent when shifting the focus in the weights, using a completely random search with no weights produces bad results. The last combination of weights seems to be performing slightly better than the others and thus, these weights were chosen in the implementation.

The weekend objective function value during a typical run is shown in Figure 5.3. The effects of the SA accept function are visible in the plot as smaller or larger dives in the objective function value. In the measuring, the SA-parameters $T_0 = 0.4$ and $\alpha = 0.985$. $It_{wend} = 1000$ were used. The graph suggests that a high weekend objective function value is found already within the 200 first iterations, which is consistent with the results in Table 5.6.

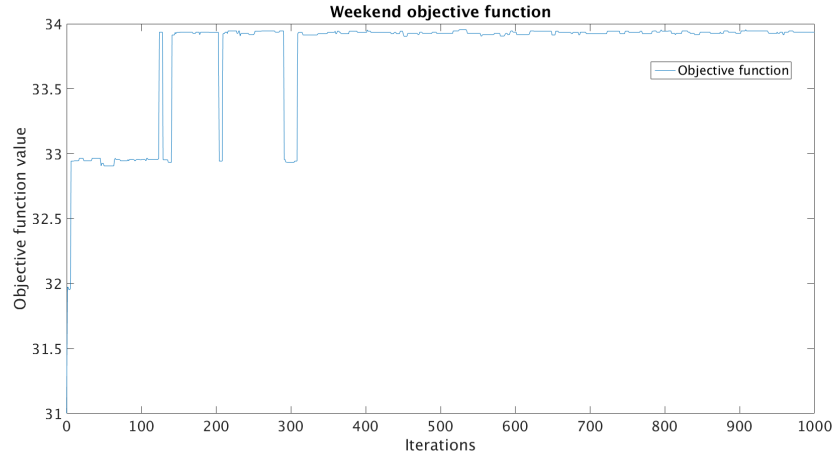


Figure 5.3: The weekend objective function for 1000 iterations

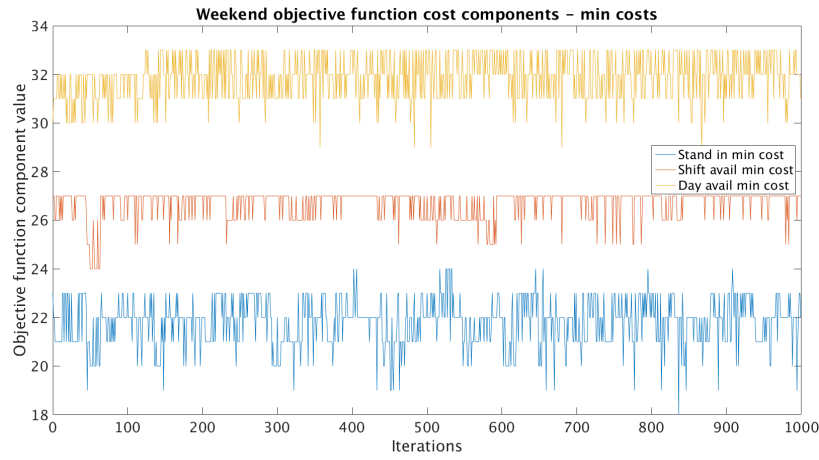


Figure 5.4: The weekend objective function minimum cost components for 1000 iterations

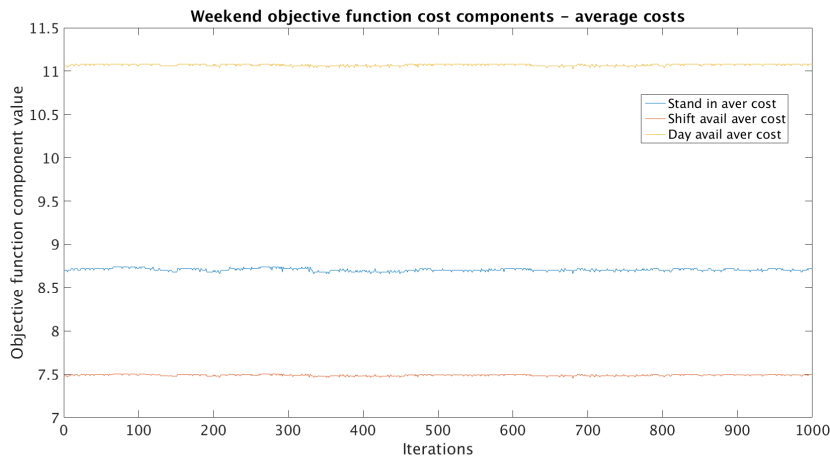


Figure 5.5: The weekend objective function average costs components for 1000 iterations

In Figure 5.4 and Figure 5.5, the three different cost components from the weekend objective function for the same run are displayed. In the first plot, the minimum costs are displayed. These seem to move between a few discrete values, although a small improvement over the iterations can be seen at least in the day avail cost component. The average values are almost constant throughout the iterations, as can be seen in the second plot. This is probably due to the fact that the total number of available staff members barely changes when shifting their weeks.

Attempts were made to correlating the min components of the weekend objective function value. However, the results were fluctuating between different tests so no clear correlations could be identified. This suggests that they are measuring three independent aspects.

5.2 Discussion

In this section, the different results presented in Section 5.1 will be discussed. Pros and cons for both methods are given so that they can be more easily compared.

5.2.1 Week block scheduling approach

Table 5.9 lists pros and cons with the implemented week block scheduling approach. Some pros and cons were considered before the heuristic was chosen. Mainly the con concerning the exponential growth of week blocks was taken into consideration, and an estimate of the upper limit of the problem size was done.

The pro regarding the same amount of week blocks was taken into consideration, as the heuristic was thought of having firstly, a block construction phase and secondly, an assignment phase. However, the most significant pro of this approach is that many of the constraints are already met, implicitly, when the blocks are created. Constraints like maximum of one task per day, one evening per week, one PL per week, two shifts at the same hour per week and more are already taken care of.

Cons like rotations, solution time and costs are no major issues as they can be avoided by a few smarter implementations. For instance, rotations can be improved by assigning a chosen value whenever a staff member is available as stand-in on a day. From the sum of those values an even distribution of possible stand-ins can be acquired by, for instance, always letting the difference between the lowest and the highest sum of stand-ins, through the days, be as small as possible. This improvement is done instead of randomly generating the new rotations in the destroy and repair iteration.

The improvement mentioned above should also implicitly decrease the solution time, as less iterations would be required to find a feasible solution.

In contrast to the cons mentioned above, a major issue is to always be able to create a pool of possible week block appearances regardless of the problem size. Just by adding meetings and the assignment of the library on wheels tasks would make the number of possible week block appearances grow considerably.

Table 5.9: Pros and cons with the implemented week block scheduling approach

Pros	Cons
Many constraints are already met, implicitly, when the blocks are created. The same number of week block appearances will exist for five and ten weeks.	Rotations need to be assigned in a systematic way in order to achieve reasonable results regarding lowest number of stand-ins through the days.
Quick iterations when destroying and repairing.	The number of unique block appearances grows exponentially in case more task types are added, such as meetings.
	The solution time can vary considerably as several random generators have been used.
	More cost parameters are needed in this approach (some correlated), where each of them affects the solution procedure.

5.2.2 Task distribution approach

The results from the previous section point to the fact that finding a good weekend distribution is essential in order to get a good final schedule. When It_{wend} is increased the results improve, up to a limit of around 500 iterations. The results are comparable to those of the AMPL implementation. This suggests that the weekend objective function used gives an indication of a good schedule. However, it only gives an indication, not an optimal solution, as the results seem to converge to a value slightly lower than the AMPL result.

Table 5.8 suggests that the weights used in the weekend objective function are important for how well the implementation works, but only to a certain degree. More testing could be done on the weights in order to see if the result can be further improved. However, since the objective function is an incomplete measure of a good schedule the results will never be completely optimal. This is listed as the main drawback in the cons column in Table 5.10.

Table 5.10: Pros and cons with the implemented task distribution approach

Pros	Cons
<p>In most cases, the method finds the optimal solution.</p> <p>The method is very fast compared to the CPLEX solver. Only a few iterations in each phase gives good results.</p> <p>The method can be used to find a good weekend schedule. The rest of the schedule can then be placed using another method such as CPLEX.</p>	<p>The weekend objective function is not an exact measurement of a good weekend schedule.</p>

Chapter 6

Concluding remarks

In this thesis work, a mathematical model for the staff scheduling problem at the Library of Norrköping was developed. The overall goal was to find schedules that are robust with respect to disturbances by evenly distributing stand-ins, that is, staff members that are not assigned to any tasks during a day. The model was solved by the commercial optimization solver CPLEX, which generated an optimal schedule with three librarian stand-ins and zero assistant stand-ins at the worst day. This is better than the current schedule at the library, where only one staff member was assigned as stand-in on the worst day.

When developing the two heuristics for solving the problem, the conclusion was reached that the weekend distribution is essential for the stand-in distribution in the final schedule. The heuristic which focused on making weekend scheduling separately from the rest of the problem thus performed better.

Although good results were obtained both by using the AMPL model and the second heuristic, there is room for further development in the modelling of the problem. One of the things which could be done in a more robust way is the modelling of exceptions and preferences among the staff members. In the current model, these are simply modelled as hard constraints. If these could be modelled in a more general way or separately from the core problem, the model would not be so sensitive to changes in exceptions or the staff preferences.

Several components of the model could be simplified. For example, the library on wheels could be modelled separately from the rest of the problem. Another suggestion is to give the same availability for all weeks despite if the staff member is working different hours during odd and even weeks. This would make it easier to model as the staff member's ten-week schedule would consist of two five-week schedules instead. Also, a more general way of describing meetings would simplify their distribution. One way of doing this could be to produce schedules which are not entirely feasible according to the problem description, but which can be used as a basis for manual adjustments.

Regarding the implementations, there are several components missing in order for the heuristics to fully correspond to the mathematical model. Although these features are not crucial, they could be implemented in order to make the heuristics correspond to the complete model. Due to lack of time, these missing components were not implemented. Furthermore, a more systematic framework for testing and evaluating the heuristics could be developed, in order to be able to measure the results more fairly.

Bibliography

- A. T. Ernst, H. Jiang, M. Krishnamoorthy, B. Owens, and D. Sier. An annotated bibliography of personnel scheduling and rostering. *Annals of Operations Research*, 127(1-4):21–144, 2004.
- M. Krishnamoorthy and A. T. Ernst. The personnel task scheduling problem. In *Optimization Methods and Applications*, pages 343–368. Springer, 2001.
- M. Krishnamoorthy, A. T. Ernst, and D. Baatar. Algorithms for large scale shift minimisation personnel task scheduling problems. *European Journal of Operational Research*, 219(1):34–48, 2012.
- L. G. Kroon, M. Salomon, and L. N. Van Wassenhove. Exact and approximation algorithms for the operational fixed interval scheduling problem. *European Journal of Operational Research*, 82(1):190–205, 1995.
- P. Smet and G. Vanden Berghe. A matheuristic approach to the shift minimisation personnel task scheduling problem. *Practice and Theory of Automated Timetabling*, pages 145–151, 2012.
- P. Smet, T. Wauters, M. Mihaylov, and G. Vanden Berghe. The shift minimisation personnel task scheduling problem: A new hybrid approach and computational insights. *Omega*, 46:64–73, 2014.
- J. S. Loucks and F. R. Jacobs. Tour scheduling and task assignment of a heterogeneous work force: A heuristic approach. *Decision Sciences*, 22(4):719–738, 1991.
- G. M. Thompson. *A Comparison of Techniques for Scheduling Non-Homogeneous Employees in a Service Environment Subject to Non-Cyclical Demand Volume I, Chapters 1-7*. PhD thesis, The Florida State University, 1988.
- K. Choi, J. Hwang, and M. Park. Scheduling restaurant workers to minimize labor cost and meet service standards. *Cornell Hospitality Quarterly*, 50(2):155–167, 2009.
- A. Rong. Monthly tour scheduling models with mixed skills considering weekend off requirements. *Computers & Industrial Engineering*, 59(2):334–343, 2010.
- M. Hojati and A. S. Patil. An integer linear programming-based heuristic for scheduling heterogeneous, part-time service employees. *European Journal of Operational Research*, 209(1):37–50, 2011.

- I. Gertsbakh and H. I. Stern. Minimal resources for fixed and variable job schedules. *Operations Research*, 26(1):68–85, 1978.
- M. Fischetti, S. Martello, and P. Toth. Approximation algorithms for fixed job schedule problems. *Operations Research*, 40(1-supplement-1):S96–S108, 1992.
- S. O. Duffuaa and K. S. Al-Sultan. A stochastic programming model for scheduling maintenance personnel. *Applied Mathematical Modelling*, 23(5):385–397, 1999.
- S. M. Roberts and L. F. Escudero. Scheduling of plant maintenance personnel. *Journal of Optimization Theory and Applications*, 39(3):323–343, 1983.
- M. Akbari, M. Zandieh, and B. Dorri. Scheduling part-time and mixed-skilled workers to maximize employee satisfaction. *The International Journal of Advanced Manufacturing Technology*, 64(5-8):1017–1027, 2013.
- S. Mohan. Scheduling part-time personnel with availability restrictions and preferences to maximize employee satisfaction. *Mathematical and Computer Modelling*, 48(11):1806–1813, 2008.
- H. A. Eiselt and V. Marianov. Employee positioning and workload allocation. *Computers & Operations Research*, 35(2):513–524, 2008.
- P. Shahnazari-Shahrezaei, R. Tavakkoli-Moghaddam, and H. Kazemipoor. Solving a multi-objective multi-skilled manpower scheduling model by a fuzzy goal programming approach. *Applied Mathematical Modelling*, 37(7):5424–5443, 2013.

Appendix A

Problem definitions

A.1 Sets

I	Set of workers
I_{lib}	Set of librarians ($I_{lib} \subseteq I$)
I_{ass}	Set of assistants ($I_{ass} \subseteq I$)
W	Set of weeks
W_5	Set of first five weeks
W_9	Set of first nine weeks
D	Set of days in a week
D_4	Set of first four weekdays
D_5	Set of all five weekdays
D_w	Set of weekend days, Saturday and Sunday
S_d	Set of shifts day d
S_3	Set of first three shifts on a weekday
J_d	Set of task types day d
I_{LOW}	Set of librarians available to work in library on wheels
I_{free_day}	Set of workers that shall be assigned a free weekday per week
I_{odd_even}	Set of all workers with odd or even weeks
$I_{weekend_avail}$	Set of workers available for weekend work
$I_{big_meeting}$	Set of all workers that attend big meetings
I_{no_PL}	Set of workers not working on PL
$I_{weekend_avail}$	Set of workers working 3-4 times PL in 10 weeks
Dep	Set of departments
$I_{dep\{Dep\}}$	Set of workers in departments
V	Set of possible week rotations (shift the week by 1-10 steps)

A.2 Variables

$$x_{iwdsj} = \begin{cases} 1, & \text{if worker } i \text{ is assigned in week } w, \text{ day } d, \text{ shift } s \text{ to a task } j \\ 0, & \text{otherwise} \end{cases} \quad (\text{A.1})$$

$$H_{iwh} = \begin{cases} 1, & \text{if worker } i \text{ works weekend } h (= 1, 2) \text{ in week } w \\ 0, & \text{otherwise} \end{cases} \quad (\text{A.2})$$

$$r_{iw} = \begin{cases} 1, & \text{if worker } i \text{ has its schedule rotated } w-1 \text{ steps} \\ 0, & \text{otherwise} \end{cases} \quad (\text{A.3})$$

$$l_{iwd} = \begin{cases} 1, & \text{if librarian } i \text{ is a stand-in week } w, \text{ day } d \\ 0, & \text{otherwise} \end{cases} \quad (\text{A.4})$$

$$a_{iwd} = \begin{cases} 1, & \text{if assistant } i \text{ is a stand-in week } w, \text{ day } d \\ 0, & \text{otherwise} \end{cases} \quad (\text{A.5})$$

$$y_{iwd} = \begin{cases} 1, & \text{if worker } i \text{ works week } w, \text{ day } d, \text{ shift } s \text{ at task type E, I or P} \\ 0, & \text{otherwise} \end{cases} \quad (\text{A.6})$$

$$W_{iwd} = \begin{cases} 1, & \text{if a worker } i \text{ is working a shift week } w, \text{ day } d \\ 0, & \text{otherwise} \end{cases} \quad (\text{A.7})$$

$$b_{iw} = \begin{cases} 1, & \text{if worker } i \text{ works at HB week } w \\ 0, & \text{otherwise} \end{cases} \quad (\text{A.8})$$

$$f_{iw} = \begin{cases} 1, & \text{if worker } i \text{ is assigned to work friday evening week } w \\ 0, & \text{otherwise} \end{cases} \quad (\text{A.9})$$

$$M_{wds} = \begin{cases} 1, & \text{if a big meeting is placed on week } w, \text{ day } d, \text{ shift } s \\ 0, & \text{otherwise} \end{cases} \quad (\text{A.10})$$

$$m_{wdsD} = \begin{cases} 1, & \text{if a meeting is placed on week } w, \text{ day } d, \text{ shift } s \text{ at department } D \\ 0, & \text{otherwise} \end{cases} \quad (\text{A.11})$$

$$d_{iwd} = \begin{cases} 1, & \text{if there is a difference in assignment of tasks at a shift } s, \text{ for a worker } i, \text{ day } d \text{ between week } w \text{ and } w+5 \\ 0, & \text{otherwise} \end{cases} \quad (\text{A.12})$$

$$l^{min} = \text{lowest number of stand-in librarians found (integer)} \quad (\text{A.13})$$

$$a^{min} = \text{lowest number of stand-in assistants found (integer)} \quad (\text{A.14})$$

$$s^{min} = \text{weighted sum with number of stand-in librarians and assistants} \quad (\text{A.15})$$

$$\delta = \text{total number of shifts that differ for all staff members} \quad (\text{A.16})$$

A.3 Parameters

$$N1l = \text{a value to prioritize the amount of stand-in librarians} \quad (\text{A.17})$$

$$N1a = \text{a value to prioritize the amount of stand-in assistants} \quad (\text{A.18})$$

$$N1 = \text{a value to prioritize total number of stand ins} \quad (\text{A.19})$$

$$N2 = \text{a value to prioritize similar weeks} \quad (\text{A.20})$$

$$avail_day_{iwd} = \begin{cases} 1, & \text{if worker } i \text{ is available for work week } w, \text{ day } d \\ 0, & \text{otherwise} \end{cases} \quad (\text{A.21})$$

$$demand_{dsj} = \text{number of workers required day } d, \text{ shift } s \text{ for task type } j \quad (\text{A.22})$$

$$qualavail_{iwsj} = \begin{cases} 1, & \text{if worker } i \text{ is qualified and available week } w, \text{ day } d, \text{ shift } s \text{ for task type } j \\ 0, & \text{otherwise} \end{cases} \quad (\text{A.23})$$

$$LOW_demand_{wds} = \text{number of workers required day } d, \text{ shift } s \text{ at the library on wheels} \quad (\text{A.24})$$

A.4 Objective function and constraints

$$\text{maximize } M \cdot s^{min} - N \cdot \delta \quad (\text{A.25})$$

$$\sum_{i \in I} x_{iwsj} = (1 - M_{wds}) demand_{dsj}, \quad w \in W, d \in D_5, s \in S_d, j \in J_d \setminus \{L\} \quad (\text{A.26})$$

$$\sum_{i \in I} x_{iwsj} = demand_{dsj}, \quad w \in W, d \in D_w, s \in S_d, j \in J_d \quad (\text{A.27})$$

$$\sum_{i \in I} x_{iwsL} = LOW_demand_{wds}, \quad w \in W, d \in D_5, s \in S_d \quad (\text{A.28})$$

$$\sum_{d \in D_5} \sum_{s \in S_d} \sum_{j \in J_d} x_{iwsj} \leq 4, \quad i \in I \setminus \{36\}, w \in W \quad (\text{A.29})$$

$$\sum_{w \in W_5} M_{w11} = 1 \quad (\text{A.30})$$

$$M_{(w+5)11} = M_{w11} \quad (\text{A.31})$$

$$\sum_{w \in W} \sum_{d \in D_5} \sum_{s \in S_d} M_{wds} = 2 \quad (\text{A.32})$$

$$\sum_{s \in 2..3} \sum_{j \in J_d \setminus \{L\}} x_{iwsj} \leq 1 - M_{w11}, \quad i \in I \setminus I_big, w \in W \quad (\text{A.33})$$

$$\sum_{w \in W_5} \sum_{d \in D_5} \sum_{s \in S_3} m_{wdsD} = 1, \quad D \in Dep \quad (\text{A.34})$$

$$m_{(w+5)dsD} = m_{wdsD}, \quad D \in Dep, w \in W_5, d \in D_5, s \in S_3 \quad (\text{A.35})$$

$$m_{wdsD} + x_{iwsj} \leq 1, \quad D \in Dep, i \in I_dep\{Dep\}, w \in W, d \in D_5, s \in S_3, j \in J_d \quad (\text{A.36})$$

$$m_{wd2D} + m_{wd3D} + x_{iwd1P} \leq 1, \quad D \in Dep, i \in I_dep\{Dep\}, w \in W, d \in D_5, s \in S_3 \quad (\text{A.37})$$

$$m_{wdsD} \leq \sum_{v \in V} r_{iv} \cdot qualavail_{i(mod_{10}(w-v+10)+1)dsE}, \quad D \in Dep, i \in I_dep\{Dep\} \setminus \{39\}, w \in W, d \in D_5, s \in S_3 \quad (\text{A.38})$$

$$M_{w11} + \sum_{d \in D_5} \sum_{s \in S_3} m_{wdsD} \leq 1, \quad D \in Dep, i \in I_big \cup I_dep\{Dep\}, w \in W \quad (\text{A.39})$$

$$\sum_{j \in J_d} x_{iwsj} \leq 1, \quad i \in I, w \in W, d \in D, s \in S_d \quad (\text{A.40})$$

$$\sum_{s \in S_d} \sum_{j \in J_d \setminus \{L\}} x_{iwsj} \leq 1 - x_{iws_2L}, \quad i \in I, w \in W, d \in D_4, s_2 \in S_d \quad (\text{A.41})$$

$$\sum_{s \in S_5} \sum_{j \in J_d \setminus \{L\}} x_{iwsj} \leq 1, \quad i \in I, w \in W \quad (\text{A.42})$$

$$\sum_{j \in J_d \setminus \{L\}} x_{iwsj} \leq 1 - x_{iws_1L}, \quad i \in I, w \in W, s \in S_3 \quad (\text{A.43})$$

$$\sum_{s \in S_d} \sum_{j \in J_d} x_{iwsj} \leq 1, \quad i \in I, w \in W, d \in D_w \quad (\text{A.44})$$

$$\sum_{d \in D_4} \sum_{j \in J_d} x_{iwd4j} \leq 1, \quad i \in I \setminus \{36\}, w \in W \quad (\text{A.45})$$

$$\sum_{d \in D_5} x_{iwd1P} \leq 1, \quad i \in I, w \in W \quad (\text{A.46})$$

$$x_{iwsP} = 0, \quad i \in I_{noP}, w \in W, d \in D_5, s \in S_d \quad (\text{A.47})$$

$$\sum_{w \in W} \sum_{d \in D_5} x_{iwd1P} \leq 4, \quad i \in I_{manyP} \quad (\text{A.48})$$

$$\sum_{w \in W} \sum_{d \in D_5} x_{iwd1P} \leq 3, \quad i \in I_{noP} \setminus I_{manyP} \quad (\text{A.49})$$

$$\sum_{w \in W} r_{iw} = 1, \quad i \in I \quad (\text{A.50})$$

$$\sum_{w \in W} H_{iwh} \leq 1, \quad i \in I_{weekendavail}, h \in \{1, 2\} \quad (\text{A.51})$$

$$r_{iw} \geq H_{iw1}, \quad i \in I, w \in W \quad (\text{A.52})$$

$$r_{i(2 \cdot w)} = 0, \quad i \in I_{OddEven}, w \in W_5 \quad (\text{A.53})$$

$$r_{i(mod_{10}(w+4)+1)} \geq H_{iw2}, \quad i \in I, w \in W \quad (\text{A.54})$$

$$\sum_{s \in S_6} \sum_{j \in J_6} x_{iw6sj} + \sum_{s \in S_7} \sum_{j \in J_7} x_{iw7sj} = 2 \cdot \sum_{h \in \{1, 2\}} H_{iwh}, \quad i \in I_{weekendavail}, w \in W \quad (\text{A.55})$$

$$\sum_{s \in S_6} x_{iw6sj} = \sum_{s \in S_7} x_{iw7sj}, \quad i \in I_{weekendavail}, w \in W, j \in J_7 \quad (\text{A.56})$$

$$\sum_{j \in J_d \setminus \{L\}} x_{iw54j} = f_{iw}, \quad i \in I_{weekendavail}, w \in W \quad (\text{A.57})$$

$$b_{iw} = x_{iw61B}, \quad i \in I_{weekendavail}, w \in W \quad (\text{A.58})$$

$$f_{iw} \geq H_{iwh} + (1 - b_{iw}) - 1, \quad i \in I_{weekendavail}, w \in W, h \in \{1, 2\} \quad (\text{A.59})$$

$$f_{iw} \leq \sum_{h \in \{1, 2\}} H_{iwh}, \quad i \in I_{weekendavail}, w \in W \quad (\text{A.60})$$

$$f_{iw} \leq 1 - b_{iw}, \quad i \in I_{weekendavail}, w \in W \quad (\text{A.61})$$

$$\sum_{w \in W} \sum_{d \in \{6, 7\}} x_{iwd1B} \leq 2, \quad i \in I_{lib} \setminus \{23\} \quad (\text{A.62})$$

$$\sum_{d \in \{6, 7\}} \sum_{j \in \{E, I\}} x_{23wd1j} = 0, \quad w \in W \quad (\text{A.63})$$

$$s^{min} \leq N1l \cdot \sum_{i \in I_{lib}} l_{iwd} + N1a \cdot \sum_{i \in I_{ass}} a_{iwd}, \quad w \in W, d \in D_5 \quad (\text{A.64})$$

$$W_{iwd} \geq \sum_{s \in S_3} \sum_{j \in J_d} x_{iwdsj}, \quad i \in I, w \in W, d \in D_5 \quad (\text{A.65})$$

$$W_{iwd} \geq \sum_{j \in J_d} x_{iwd4j}, \quad i \in I, w \in W, d \in D_5 \quad (\text{A.66})$$

$$l_{iwd} \geq \sum_{v \in V} (r_{iv} \cdot \text{avail_day}_{i(\text{mod}_{10}(w-v+10)+1)d}) + (1 - W_{iwd}) - 1, \quad i \in I_{lib}, w \in W, d \in D_5 \quad (\text{A.67})$$

$$l_{iwd} \leq \sum_{v \in V} (r_{iv} \cdot \text{avail_day}_{i(\text{mod}_{10}(w-v+10)+1)d}), \quad i \in I_{lib}, w \in W, d \in D_5 \quad (\text{A.68})$$

$$l_{iwd} \leq 1 - W_{iwd}, \quad i \in I_{lib}, w \in W, d \in D_5 \quad (\text{A.69})$$

$$a_{iwd} \geq \sum_{v \in V} (r_{iv} \cdot \text{avail_day}_{i(\text{mod}_{10}(w-v+10)+1)d}) + (1 - W_{iwd}) - 1, \quad i \in I_{ass}, w \in W, d \in D_5 \quad (\text{A.70})$$

$$a_{iwd} \leq \sum_{v \in V} (r_{iv} \cdot \text{avail_day}_{i(\text{mod}_{10}(w-v+10)+1)d}), \quad i \in I_{ass}, w \in W, d \in D_5 \quad (\text{A.71})$$

$$a_{iwd} \leq 1 - W_{iwd}, \quad i \in I_{ass}, w \in W, d \in D_5 \quad (\text{A.72})$$

$$d_{iwd} \geq (y_{iwd} - y_{i(w+5)d}), \quad i \in I, w \in W_5, d \in D_5, s \in S_3 \quad (\text{A.73})$$

$$d_{iwd} \geq -(y_{iwd} - y_{i(w+5)d}), \quad i \in I, w \in W_5, d \in D_5, s \in S_3 \quad (\text{A.74})$$

$$\delta = \sum_{i \in I} \sum_{w \in W_5} \sum_{d \in D_5} \sum_{s \in S_3} d_{iwd} \quad (\text{A.75})$$

$$x_{iwd} \leq \sum_{v \in V} (r_{iv} \cdot \text{qualavail}_{i(\text{mod}_{10}(w-v+10)+1)ds}), \quad i \in I, w \in W, d \in D, s \in S_d, j \in J_d \quad (\text{A.76})$$

$$y_{iwd} = \sum_{j \in J_d \setminus \{L\}} x_{iwdj}, \quad i \in I, w \in W, d \in D_5, s \in S_3 \quad (\text{A.77})$$

$$\sum_{d \in D_5} y_{iwd} \leq 2, \quad i \in I, w \in W, s \in S_3 \quad (\text{A.78})$$

$$\sum_{w \in W} x_{25w44L} = 8 \quad (\text{A.79})$$

$$x_{36w14L} = 1 - x_{25w11L}, \quad w \in W \quad (\text{A.80})$$

$$\sum_{w \in W} x_{37w34L} = 5 \quad (\text{A.81})$$

$$x_{37w34L} = 1 - x_{37(w+1)34L}, \quad w \in W_9 \quad (\text{A.82})$$

$$x_{36wd4j} = 0, \quad w \in W, d \in D_4, j \in J_d \setminus \{L\} \quad (\text{A.83})$$

Appendix B

Flow charts

B.1 Weekly scheduling approach

B.2 Task distribution approach

Figure B.1: Algorithm for distributing weekends.

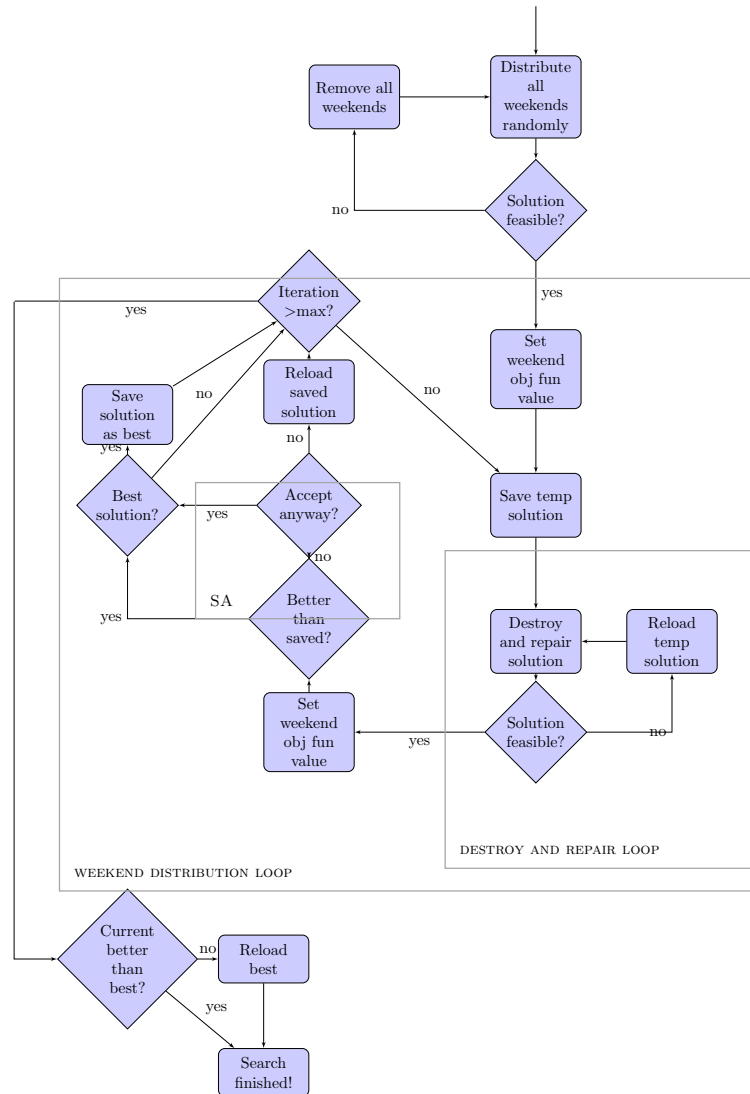
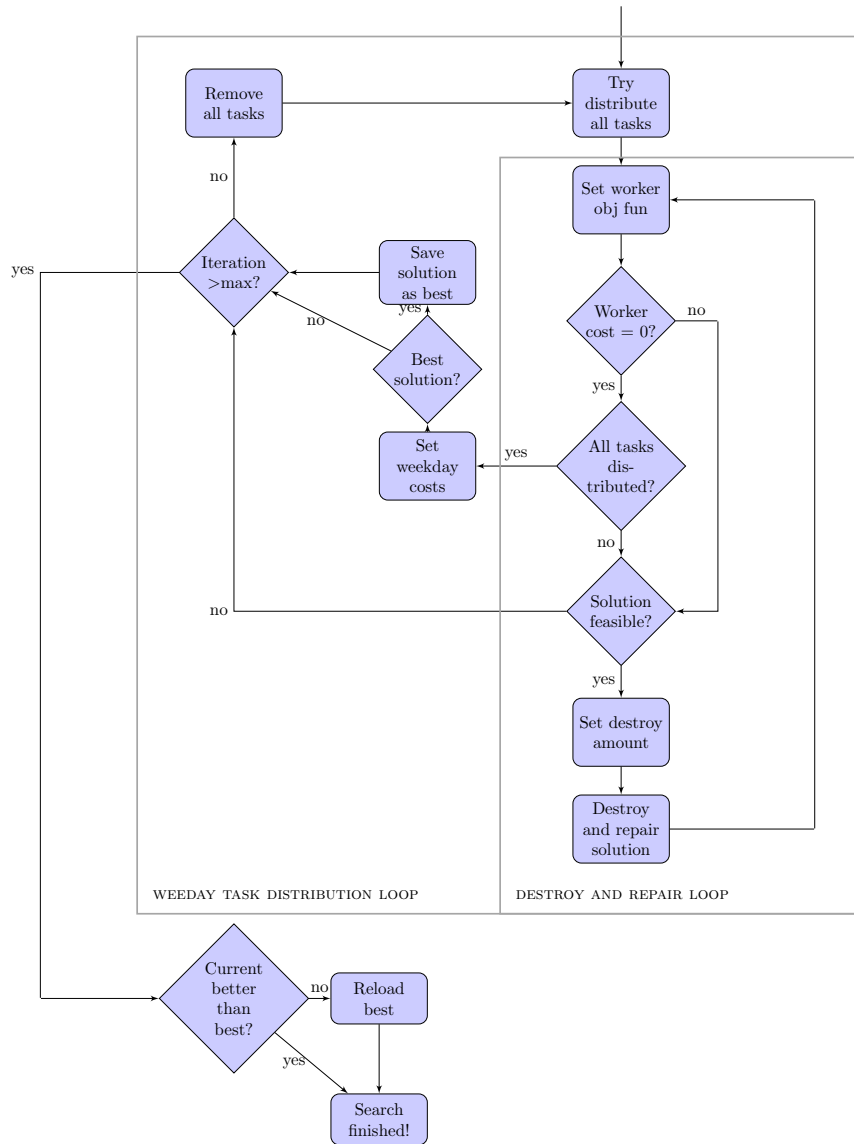


Figure B.2: Algorithm for distributing weekday tasks.



Appendix C

Week block table

Table C.1: Number of assignable unique blocks for the workers based on their availability and qualification.

Worker	Weekend	Weekrest	Weekday
1	532	347	1580
2	1580	1580	1580
3	1063	347	1580
4	557	165	779
5	261	130	531
6	532	130	1580
7	261	130	531
8	261	29	531
9	115	12	247
10	532	130	1580
11	9	8	8
12	1063	347	1580
13	771	92	1190
14	265	29	489
15	51	18	120
16	495	130	843
17	237	69	267
18	532	279	1580
19	495	47	843
20	532	130	1580
21	227	227	227
22	2	1	1
23	3	2	2
24	11	5	47
25	1063	279	1580
26	5	4	4
27	213	106	425
28	2	2	2
29	426	106	1281
30	127	126	455
31	495	130	843
32	261	29	531
33	72	45	306
34	425	425	425
35	91	20	221
36	55	27	101
37	1063	347	1580
38	3	1	1
39	2	1	1

Copyright

The publishers will keep this document online on the Internet - or its possible replacement - for a period of 25 years from the date of publication barring exceptional circumstances. The online availability of the document implies a permanent permission for anyone to read, to download, to print out single copies for your own use and to use it unchanged for any non-commercial research and educational purpose. Subsequent transfers of copyright cannot revoke this permission. All other uses of the document are conditional on the consent of the copyright owner. The publisher has taken technical and administrative measures to assure authenticity, security and accessibility. According to intellectual property law the author has the right to be mentioned when his/her work is accessed as described above and to be protected against infringement. For additional information about the Linköping University Electronic Press and its procedures for publication and for assurance of document integrity, please refer to its WWW home page: <http://www.ep.liu.se/>

Upphovsrätt

Detta dokument hålls tillgängligt på Internet - eller dess framtida ersättare - under 25 år från publiceringsdatum under förutsättning att inga extraordinära omständigheter uppstår. Tillgång till dokumentet innebär tillstånd för var och en att läsa, ladda ner, skriva ut enstaka kopior för enskilt bruk och att använda det oförändrat för ickekommersiell forskning och för undervisning. Överföring av upphovsrätten vid en senare tidpunkt kan inte upphäva detta tillstånd. All annan användning av dokumentet kräver upphovsmannens medgivande. För att garantera äktheten, säkerheten och tillgängligheten finns det lösningar av teknisk och administrativ art. Upphovsmannens ideella rätt innefattar rätt att bli nämnd som upphovsman i den omfattning som god sed kräver vid användning av dokumentet på ovan beskrivna sätt samt skydd mot att dokumentet ändras eller presenteras i sådan form eller i sådant sammanhang som är kränkande för upphovsmannens litterära eller konstnärliga anseende eller egenart. För ytterligare information om Linköping University Electronic Press se förlagets hemsida <http://www.ep.liu.se/>

© 2016, Claes Arvidson and Emelie Karlsson