Master thesis

Work Distribution of a Heterogeneous Library Staff - A Personnel Task Scheduling Problem

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Abstract

Here is where you can write your abstract. It may be very long, or it may be very short, the reason you have an abstract is for people not to be forced to read lots of crap.

But still, they will have to read your abstract. After all, the abstract is what everyone reads. . .

Keywords: Keyword One, Chemostat, Another Key-Word, Key, Clé, Mot de cle, Nyckelhål, XBOX, Dagens viktigaste nyckelord, and Keywords.

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Nomenclature

Most of the reoccurring definitions, symbols and abbreviations are described here.

Definitions

Plocklista Text Library on wheels Text

Symbols

- Y_0 The amount of the variable Y inserted into a system.
- \hat{Y} The unit-dimension of the variable Y, for example $\hat{t} = 1s$.
- \bar{Y}_i A steady state (number i) value of Y.
- K_i Constants used in kinetic expressions, for example K_I .
- A The system matrix.

Abbreviations

 $\begin{array}{ccc} \operatorname{Exp} & \operatorname{Text} \\ \operatorname{Info} & \operatorname{Text} \\ \operatorname{PL} & \operatorname{Text} \\ \operatorname{PTSP} & \operatorname{Text} \\ \operatorname{SMPTSP} & \operatorname{Text} \end{array}$

CPI Competitive Product Inhibition (or Inhibited)
CSI Competitive Substrate Inhibition (or Inhibited)

CSTR Continuous Stirred Tank (bio)Reactor MMI Michaelis-Menten Inhibition (or Inhibited)

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Chapter 1

Introduction

1.1 Background

At a library absence can cause problems, both due to lack of personnel as well as due to the qualifications required to perform a task varies. If a worker were to be unavailable a day because of a meeting or being ill it would require for a stand-in to fill the vacancy. Therefore, it is of great interest to have a schedule with as many skilled stand-ins as possible to overcome such disturbances. Furthermore, the library personnel have certain demands and preferences as to how a satisfactory working schedule should be. For instance, it is neither preferable to work more than one evening each week nor work more weekends than required.

1.2 Problem description

1.2.1 Main objectives

Main objective is to create a schedule with as many stand-ins as possible, so that all days have a maximum amount of stand-ins. Diversity during the week and repetivity of the schedule each week is also desired.

1.2.2 Requirements on weekday activities

During weekdays, there is a worker demand at the stations Exp, Info, Plocklista and Library on wheels. ¡TABLE OF DEMAND¿. Plocklista is unique in the sense that its duration is longer than one shift. It is modeled as the three first shifts of a weekday.

In addition there is the Library on wheels, which has a different demand of workers depending on odd and even weeks.

1.2.3 Requirements on weekend activities

There are three different weekend stations: Exp, Info and Hageby. Working a weekend also means working friday unless you are scheduled to work in Hageby.

1.2.4 Diversity in skill

There are essentially two types of workers in a library. Librarians and assistants. The competence of the workers looks as follows: ¡TABLE OF DIFFERENT COMPETENCES;

A subset of the librarians can handle the Library on wheels and all are expected to take shifts at Hageby.

1.2.5 Personnel availability and task assignment limitations

The staff at the library have certain times of availability. Every day, a person can perform at maximum one task. Plocklistan is limited to once a week.

Personnel is only allowed to work one weekend out of five. The following week, the worker is free on the times requested by the worker. Also, some workers have specific

1.2.6 Inner work

The goal of this thesis is to distribute given tasks to the heterogeneous workforce at the library of Norrköping. Each task is either classified as an outer or an inner service where an outer service is when a librarian needs to interact with visitors. Inner services can in some rare cases require a predetermined person to be assigned to a specified time or day.

Demands and requests are to be fulfilled to the furthest extent possible. Weekends are included in the scheduling problem, which adds more constraints regarding the number of contiguous working days. However, the librarians are permitted a few exceptions from these laws regarding days of rest.

The main purpose of the thesis is to create a schedule robust enough to withstand absence, such that outer services always are assigned to a qualified and available worker. This is visualized as having a list of available stand-ins for each shift.

There are a limited number of workers at the library and they make the resources that are to be distributed. Each individual has a set of *skills* and *competences*. Competences refer to the capability of being assigned the different outer services; Expedition, Norpan, Information desk, Library on wheels and Hageby as well as different inner services. The set of skills an individual can possess are described in Table 1.1. In total there are 39 workers available.

The outer services can be seen as assignments which requires available workers to be assigned to them. Each outer service is specified to a certain station, time and date. They also have a fix length and occur on a regular basis every ten weeks, which makes it possible to create a periodic schedule with a period of ten weeks.

Furthermore, outer and a few inner services can be characterized by different properties, which are represented in Table 1.2.

In addition to the properties mentioned above, there are several requirements that have to be met. These can be divided into job, robust and other requirements and are listed in Table 1.3 below.

Table 1.1: Personnel

| Skills | Description |
|--------------------|---|
| Work degree | 0-100 % |
| Type of employment | Librarian/Assistant |
| Competence | Inner and outer services the worker is qualified for |
| Weekly rest | Which days the worker has requested after working a weekend |
| Other requests | Does not work evenings etc. |

Table 1.2: Outer and inner services

| Outer service | Property | | | | |
|----------------|---|--|--|--|--|
| | Start time, end time, week and duration | | | | |
| | Station | | | | |
| | Number of qualified librarians demanded | | | | |
| | Number of qualified assistants demanded | | | | |
| Inner service | D 4 | | | | |
| illier service | Property | | | | |
| illier service | Start time, end time, week and duration | | | | |
| inner service | 1 0 | | | | |
| illier service | Start time, end time, week and duration | | | | |

Table 1.3: Requirements

| Job requirements | Description | | | | |
|---------------------|---|--|--|--|--|
| 1 | A maximum of one outer service is to be distributed to each person and day. | | | | |
| 2 | Remaining work time is individually distributed on assignments such as reshelving b | | | | |
| 3 | Weekend work are to be evenly distributed between the workers available on weekend | | | | |
| 4 | Working a weekend includes work on saturday and sunday the same week. | | | | |
| 5 | One evening shift on a weekday per person each week except when weekend work is | | | | |
| 6 | Every ten weeks the schedule is to be repeated. | | | | |
| 7 | It is recommended for each week to be as similar as possible. | | | | |
| Robust requirements | Description | | | | |
| 1 | Each outer service require at least one stand-in. | | | | |
| 2 | The stand-ins have to be qualified for the tasks they are stand-in for. | | | | |
| 3 | Focus is to maximize the lowest number of stand-ins of any task. | | | | |
| Other requirements | Description | | | | |
| 1 | Department and general meetings are to be held once per five weeks. | | | | |

There are also additional requirements of the resulting schedule made by the workers at the library. Two examples would be that a handful staff members are unable to work weekends as well as some personnel are unable to work in the evenings.

1.3 Problem categorization

The problem can be formulated as a personnel tasks scheduling problem for a heterogeneous workforce since the main objective is to distribute tasks to

workers during their available times. The workforce is heterogeneous as certain tasks can only be performed by librarians or is restricted to a certain subset of the personnel. Another aspect of the problem is the cyclic nature of the personnel schedules, which gives a degree of freedom in availability of personnel.

Chapter 2

Literature review

The scheduling problem is a mathematical optimization problem which has been studied since the 1950's with the objective of creating a feasible and satisfactory schedule for workers or machines performing tasks. Ernst et al. provide an overview of work in the area up to 2001. They state that, although the complexity of the scheduling problem has not increased in recent years, the mathematical models used to solve the scheduling problems have become more realistic and refined. Due to this as well as the development of more powerful computational methods, it is possible today to solve scheduling problems in a more satisfactory way than before. Such new models take into account softer values such as worker satisfaction and worker fatigue [1].

In this section, the scheduling problem is classified into different subcategories which are areas related to the work of this paper. A few relevant areas for our work include Personnel Task Scheduling Problems (PTSP), Shift Minimization Task Scheduling Problems (SMTSP), Tour Scheduling Problems (TSP) and a few variations of these. Within these categories, the subproblem of task assignment, that is, the assignment of who does what is most relevant for our problem.

2.1 Personnel Task Scheduling Problem

In many practical instances production managers will face the Personnel Task Scheduling Problem (PTSP) while scheduling plant operations. It occurs when the rosterer or shift supervisor need to allocate tasks with specified start and end times to available personnel who have the required qualifications. Furthermore, it also occurs in situations where tasks of fixed times shall be assigned to machines. Decisions will then have to be made regarding the amount of maintenance workers needed and which machine the workers are assigned to look after. [2]

There are numerous variants to the PTSP. Studies on these have been made in article [2] by Krishnamoorthy et al. who gives a list of attributes that commonly appear in a PTSP and which are listed in Table 2.1 below. There are furthermore traits that always appear in a PTSP; tasks with fixed start and end time are to be distributed to staff members that possess certain skills, allowing them to perform only a subset of the available tasks. Start and end time of

their shifts are also predetermined for each day.

One variant, which also is the most simple, is mentioned in [2] and is called the *Feasibility Problem* where the aim is to just find a feasible solution. This requires that each task is allocated to a qualified and available worker. It is also required that a worker cannot be assigned more than one task simultaneously as well as tasks cannot be pre-empted, meaning that each task has to be completed by one and the same worker.

In Table 2.1 one can see attributes of PTSP variants. The nomenclature of the attributes T, S, Q, O refer to the *Task type*, *Shift type*, *Qualifications* and *Objective function* respectively.

Table 2.1: PTSP variants

| Attribute | Type | Explanation | | | |
|-----------|------------|---|--|--|--|
| Т | F | Fixed contiguous tasks | | | |
| | V | Variable task durations | | | |
| | S | Split (non-contiguous) tasks | | | |
| | $^{\rm C}$ | Changeover times between consecutive tasks | | | |
| S | F | Fixed, given shift lengths | | | |
| | I | Identical shifts which are effectively of infinite duration | | | |
| | D | Maximum duration without given start or end times | | | |
| | U | Unlimited number of shifts of each type available | | | |
| Q | I | Identical qualification for all staff (homogeneous workforce) | | | |
| | H | Heterogeneous workforce | | | |
| О | F | No objective, just find a feasible schedule | | | |
| | A | Minimise assignment cost | | | |
| | Т | Worktime costs including overtime | | | |
| | W | Minimise number of workers | | | |
| | U | Minimise unallocated tasks | | | |

Many of the most basic problems and a few more complex ones can be described with this definition of PTSP attributes. It is, however, not possible to describe all of the numerous types of PTSP using these nomenclatures [2].

By combining attributes it is possible to obtain more complex variants of the PTSP. An example would be the PTSP[F;F;H;A-T-W] mentioned in [2] where multiple objectives are used. This problem has fixed contiguous tasks, fixed shift lengths, heterogeneous workforce and three objective functions; A-T-W, which represent assignment costs, work time with overtime included and requirements to minimize the number of workers respectively. For this problem the objective function is then a linear combination with different parameters used to prioritize (weigh) them against each other.

Given the nomenclature above, our problem would be most related to the PTSP[F;F;H;F]. The difference is that the objective function is not empty. We are looking to maximize the number of qualified stand-ins each day as well as maximize employee satisfaction by meeting their recommendations. Since we have a fix number of workers, no costs and no unallocated tasks when a feasible solution is found, this cannot be described with the type of objective attributes given in Table 2.1. Therefore, none of the objective function types are relevant in our case.

Different variants of PTSP are given names in the literature. An example is

when the shifts and qualifications are identical (S=I and Q=I) and the objective function is to minimize the number of workers that are used (O=W). This variant, PTSP[F;I;I;W], has been published as the "fixed job schedule problem" and is described in Section 2.4 [2].

2.1.1 Applications

An example where PTSP can be found is when developing a rostering solution for ground personnel at an airport. Such a problem can be dealt with by first assigning the workers to days in order to satisfy all the labour constraints, followed by assigning the tasks to the scheduled workers [2].

Three problems of type PTSP related to airplanes can be found when scheduling for either airport maintenance staff, planes to gates or staff that do not stay in one location, such as airline stewards. Scheduling for airport maintenance staff can lead to either PTSP[F;I;H;U-A] or PTSP[F;I-U;H;W], which are similar problems but are given two different names; Operational Fixed Interval Scheduling Problem and Tactical Fixed Interval Scheduling Problem respectively. These are described further in Section 2.4 [2].

Another application, which has been frequently studied, is classroom assignments and is discussed in [2]. Based on specifications such as the amount of students in a class or the duration of a class, different classrooms have to be considered. Requirements of equipment, e.g. for a laboratory, may also greatly limit the available classrooms to choose from. A majority of the complications of this problem is due to the fact that lessons can span over multiple periods.

Worth noting for classroom assignment problems is that there are no start or end times for the shifts, as they represent the rooms. The aim in the present problem would be to simply find a feasible assignment of classrooms. Therefore the nomenclature of the problem would be PTSP[S;I;H;F], with the possibility of adding preferences to the objective function. An example of a preference would be to assign the lessons as close to each other as possible on a day, preventing traveling distances between classes for teachers and students [2].

2.2 Shift Minimisation Personnel Task Scheduling Problem

A close relative to the PTSP is the Shift Minimisation Personnel Task Scheduling Problem (SMPTSP) and is a special case in which the aim is to minimize the cost occuring due to the number of personnel (shifts) that are used. The same common traits are valid in this problem as in the PTSP; workers with fixed work hours are to be assigned tasks, with specified start and end times, that they are qualified for [3].

In article [3] they "... concentrate mainly on a variant of the PTSP in which the number of personnel (shifts) required is to be minimised.". In doing so, it is possible to determine the lowest number and mix of skilled staff a company should have to be able to complete the tasks and still be operational. They also presumed that the pool of workers are unlimited for either skill group, which is not the case in our problem due to the limitations on the amount of librarians and assistants available. SMPTSP can be applied when there are a large number of workers available with different qualifications and it is needed to ensure that the tasks for that day are performed. The PTSP and SMPTSP are therefore useful day-to-day management tools that commonly occurs in many practical instances where tasks are allocated on a daily basis [3].

It is shown in [4] that SMPTSP is a complex problem even if the preemption constraint were to be removed. However, if the qualifications of the workers were identical it would become an easily solvable problem [3].

SMPTSP is almost identical to another problem introduced by Kroon et al. which is called the Tactical Fixed Interval Scheduling Problem and is described in Section 2.4 below [3].

2.3 Tour Scheduling Problem with a heterogenous work force

The Tour Scheduling Problem (TSP) involves creating work shifts with days off for a work force. A shift here refers to a set of contiguous hours during which a worker is assigned for work. The need for days off occurs when there is weekend demand for staff and other free days need to be assigned instead.

According to Loucks and Jacobs, the vast majority of all tour scheduling problems up to 1991 involved a homogeneous workforce, that is, any worker can perform any assigned task [5]. One such early study of the our scheduling problem often mentioned in literature is provided by Thompson in 1988 [6]. The problem studied in this PhD thesis concern only homogeneous work forces and the task assignment part is lacking.

In the article by Loucks and Jacobs, the authors study a tour scheduling problem with a heterogeneous work force. The problem both involves tour scheduling and task assignment, where the latter part is most interesting to us. The problem is studied in the context of fast food restaurants, where certain personnel is qualified only for certain stations in the restaurant. In such industries, the demand of staff differs between different weekdays and different times of the day. Two worker attributes are considered; their availability for work and their qualification for performing different tasks. The problem concerns finding shifts for all workers which are to have a length between a minimum and maximum number of hours per day.

The representative problem studied in the article involves creating a one-week schedule for 40 workers in a fast food restaurant, available for eight different tasks with a seven-day, 128-hour workweek. Several synthetic problems are studied in the article, all, however with minimum shift length three hours, maximum shift length eight hours and five maximum number of work days.

A similar problem to the one descibed by Loucks and Jacobs is studied by Choi et al. [7]. They focus on a particular fast food restaurant in Seoul, which is made a representative of fast food chains in general. In this study, only two types of workers are available; fulltime and part time workers, with no other reference to difference in skill. The different shifts are already given by the reastaurant managers and the task is to combine them into a tour. The task assignment aspect is lacking in this article.

In both articles the main objective is to minimize both overstaffing and

understaffing, which will both have economical consequences for the fast food chain. This is done by redusing or increasing the work force. For a problem with a fixed work force, such as ours, this objective is not relevant. In the example studied by Loucks and Jacobs there is also a goal to meet staff demand on total working hours. This is modeled as a secondary goal and is similar to our goal and somehow models a "soft" value, which is of interest to us.

A more recent tour scheduling problems concern monthly tour scheduling, as opposed to most literature which concerns only weekly scheduling. Such a study was done by Aiying Rong in 2010 [8]. The main advantage of monthly scheduling over shorter time periods, as stated in the article, is the possibility to plan a schedule with respect to fairness and balance over a longer period of time. The problem concerns workers with different skills, where each worker also can possess multiple skills. This is referred to as a mixed skill problem. Thus the problem is similar to our problem, where mixed skill is also present. In the study, workers have individual weekend-off requirements. The problem does not involve task assignment, which makes it less relevant for us.

2.4 Other similar problems

In this section a couple of problems similar to our own will be described in order to give clarity as to how closely related many of these problem types are.

2.4.1 Fixed Job Schedule Problem

Variations of the task assignment problem relevant for our problem include for example the Fixed Job Schedule Problem (FJSP). The FJSP has been studied since the 1970s in the context of task assignment in processors. The problem concerns the distribution of tasks with fixed starting and ending times over a workforce with identical skills, such as processing units [3]. Such problems have been solved by I. Gertsbakh, H.I. Stern [9] and Fischetti et al. [10].

In the article [9] by Gertsbakh, a situation where n jobs need to be scheduled over an unlimited number of procesors is studied. The objective function of such a problem becomes to minimize the number of machines needed to perform all tasks. Fischetti solves a similar problem, but adds time constraints, saying that no processor is allowed to work for more than a fixed time T during a day as well as a spread time constraint forcing tasks to spread out with time gap s over a processor.

2.4.2 Tactical Fixed Interval Scheduling Problem

Another type of problem is the Tactical Fixed Interval Scheduling Problem (TFISP). This is a problem very closely related to the SMPTSP problem with the only difference being that the TFISP concerns workers which always are available, such as industrial machines or processors. The problem is studied by for example Kroon et al. [4]. A typical TFISP can be expressed using the nomenclature in Table 2.1 and written as PTSP[F;I-U;H;W] [2].

As opposed to the FJSP, the TFISP deals with a heterogeneous workforce. Two different contexts are studied by Kroon et al. One of them concerns the handling of arriving aircraft passengers at an airport. Two modes of transport from the aeroplane to the airport are investigated; directly by gate or by bus. The two transportation modes thus correspond to two processing units which can only handle a number of jobs at the same time.

2.4.3 Operational Fixed Interval Scheduling Problem

The Operational Fixed Interval Scheduling Problem (OFISP) is a close relative to TFISP, where both types are restricted by the following; each machine (worker) cannot handle more than one job at a time, each machine can only handle a subset of the jobs and preemption is not allowed. The difference between them occurs in the objective function, as TFISP tries to minimize the number of workers while OFISP tries to minimize the operational costs and the number of unallocated tasks [11]. In the present nomenclature this would give rise to the problem PTSP[F;I;H;U-A] [2]. Given the problem definition above, working shifts are to be created for the workers and tasks are to be allocated on a day-to-day basis. OFISP can therefore be seen both as a job scheduling problem and a task assignment problem [11].

2.4.4 Maintenance scheduling

What differs mostly between the problem types described above and the problem studied in this thesis work, is the difference in objective. The main objective for the librarian scheduling problem, after assigning all tasks to available personnel, is to maximize the number of stand-in staff. A similar problem arises in the maintenance industry, where some jobs can be forseen and other jobs are of a stochastic nature, that is, there is a probability that a maintenance job will occur a certain hour. The problem combing both unplanned and planned maintenance worker scheduling was studied in 1999 by Duffuaa and Al-Sultan, as a continuation of Robers' and Escudero's work in 1983 [12] [13].

In the article, a fixed, heterogeneous work force consisting of electricians, plumbers and mechanics is studied. The shifts of the personnel are predetermined by their given work times and thus the problem becomes a pure task assignment problem. An objective function is used where the goal is to maximize the number of planned and unplanned jobs performed by the workers, by taking into account the probability of unplanned work to occur. Thus, certain workers will be left at the station as stand-ins in the case an unplanned job arrives.

2.5 Modeling soft constraints

For most scheduling problems, the main objective is to minimize worker-related costs by reducing the number of workers needed to perform a task, or by reducing the working time for part-time employees. Equivalently, the goal in production industries is to reduce the number of machines needed. Recently, however, many studies have started to focus more or softer values such as worker satisfaction as an objective. Such values are usually considered when scheduling is done manually, but have been forgotten or set aside in mathematical modeling.

In an article by Akbari from 2012 a scheduling problem for part-time workers with different preferences, seniority level and productivity is investigated. In this

2.6. Solution Methods 11

article, these aspects are reflected in the objective function and weighted against each other. [14]. A similar problem was also studied by Mohan in 2008, but for a work force of only part-time workers [15].

Other factors which may affect worker satisfaction, and in the long run efficiency and presence at work are fagique, fairness and boredom. These are discussed by Eiselt and Marianov [16]. Repetitiveness of a job as well as the level of challenge can cause bordom in workers. Increasing variance is done by Eiselt and Marianov through providing an upper bound of how many tasks can be performed in a given time span. The article suggests some sort of measurement of the distance between the task requirements and the worker abilities is used. This will then be minimized in the objective function.

2.6 Solution Methods

In many real life situations, the scheduling method used to create worker schedules is a simple matching algoritm between two can do what and when. The process is most often left in the hands of experienced and knowledgable schedulers, who konw the capacity of the work force and how to maximize productivity by meeting task demands as well as employee demands and individual personality traits. This is referred to as the "art" of scheduling [13]. However, when personnel forces grow large and there are regulations, task skill requirements or several personnel preferences to take into account, the problem becomes too large to solve manually in a satisfactory manner.

The first computational methods used for solving scheduling problems were in many cases simple heuristics resembling the scheduling process as performed in a manual way. One example of this is the heuristic presented by Loucks and Jacobs which assigns workers to tasks, following certain rules, until all tasks are assigned [5]. An overview of solution methods is given by Ernst et al, where almost 30 different methods are presented and it is not uncommon that special purpose algorithms are used to suite a specific problem [?]. Some of the more interesting solution methods with respect to the probelm studied in this thesis are discussed in this section. These include solving with commercial solvers, matheuristic methods such as simulated anneahling and variable neighbourhood search, pure heuristic methods, goal programming and fuzzy goal programming.

2.6.1 Commersial software

Commersially available scheduling programs.

2.6.2 Mathematical Programming

Formulating a mathematical model. Objective function and constraints. Solving using commercial solver such as CPLEX or Guroby.

Stochastic and non-stochastic

2.6.3 Simulated Annealing

Simulated Annealing (SA) has been studied as a solution method to the scheduling problem by researchers such as Brusco and Jacobs in the early 1990's. The

method is a metaheuristic method which has the advantage over local search methods that it does not easily get stuck in local optima. The method is a random optimization method designed to find a global optimum solution. The method allows bad moves according to a function According to Akbari, Simulated Anneahling

2.6.4 Variable Neighbourhood Search

Also avoids local optima.

"A variable neighborhood search based matheuristic for nurse rostering problems" Della Croce et Salassa. "VNS outoperforms exact commercial general purpose solvers" Matheuristic approach!

Early work by: Hansen and Mladenovic, 2001, Mladenovic and Hansen, 1997

2.6.5 Tabu Search

Commonly used meta-heuristic.

2.6.6 Goal programming and Fuzzy Goal Programming

GP: Used for multiple goals.

FGP: Used for contradictory goals.

Bellman and Zadeh's max-min operator!

Fuzzy goal programming. "Fuzzy goals" = soft constraints. Fuzzy set theory. The basic idea of FGP is to present some of the model parameters as imprecise numbers. Goal programming: good when combining soft and hard constraints.

Using an average value approach with goals that are contradictory makes it possible to maximize the amount of "goodness" in the solution, by priotritizing one constraint over another, which in total generates the most good.

Chapter 3

The mathematical model

In this chapter the mathematical model implemented to solve this problem will be presented. Variables and sets will be stated prior to the objective function and constraints. A list of all the parameters can be found in Appendix B.

3.1 Set and variable definitions

To solve the problem many sets and definitions had to be declared as there are many unique requirements that have to be met. An example is that some want a day free from outer tasks so that they can focus on other assignments or attend meetings. Another one is that some personnel have different schedules for odd and even weeks. These specific cases have to be modeled and result in a variety of set and variable definitions. Hence, a bigger part of the definitions are instead listed in Appendix A.

Table 3.1: Significant sets and variables

| I | Set of workers |
|-------------|--|
| W | Set of weeks |
| D | Set of days in a week |
| S_d | Set of shifts day d |
| J_d | Set of task types day d |
| x_{iwdsj} | 1 if worker i is assigned task j in week w , day d , shift s , 0 otherwise |
| $h1_{iw}$ | 1 if worker i works its first weekend week w , 0 otherwise |
| $h2_{iw}$ | 1 if worker i works its second weekend week w , 0 otherwise |

3.2 Objective function

Due to multiple objective functions, it has been necessary to weigh them against each other to prioritize using parameters. As can be seen in equation ... below

$$maxN1l*lowest_stand_in_amount_lib + N1a*lowest_stand_in_amount_ass$$
 (3.1)

$$\begin{split} maxN1l*lowest_stand_in_amount_lib \\ + N1a*lowest_stand_in_amount_ass \\ - N2*\sum_{n=1}^{\infty} 2^{-n} = 1 \end{split}$$

3.3 Constraints

To model this problem it has been of relevance to divide many of the constraints into weekend- and weekday constraints.

Chapter 4

The ideal CSTR: the chemostat

In this chapter we study exponential growth, the logistic. . . .

4.1 Some simple models of biological growth

4.1.1 Exponential growth

If $\mu = \text{constant} > 0$, we get $X(t) = X_0 e^{\mu t}$.

4.1.2 The logistic equation

Let us assume that $\frac{dX}{dt} = \mu \cdot X$, with $\mu = \mu(S) = k \cdot S \dots$

$$\begin{cases} \frac{dX}{dt} = kSX & (a) \\ \frac{dS}{dt} = -\alpha kSX & (b) \end{cases}$$

$$\frac{dX}{dt} = r(1 - \frac{X}{B})X \tag{4.1}$$

An explicit solution to (4.1) is: $X(t) = \frac{X_0B}{X_0 + (B - X_0)e^{-rt}}$, if $0 < X_0 < B$. It can be found by separating variables in equation (4.1)

4.2 The chemostat

A chemostat is made of two main parts; a nutrient reservoir, and a growth-chamber, reactor, in which the bacteria reproduces.

$$\begin{cases}
\frac{dX}{dt} = \mu(S)X - X \frac{F}{V} \\
\frac{dS}{dt} = -\alpha\mu(S)X - S \frac{F}{V} + S_0 \frac{F}{V}
\end{cases}$$
(4.2)

$$\mathbf{A} = \left(\begin{array}{cc} 0 & \sigma \alpha_1 \\ -\frac{1}{\alpha_1} & -\sigma - 1 \end{array} \right)$$

The invariant line: conclusions

| Model | Monods Chemostat | CSI-CSTR |
|-----------------|-------------------------------------|---|
| μ | $\frac{S}{1+S}$ | $\frac{S}{1+S+\frac{S^2}{K_I}}$ |
| $\frac{dX}{dt}$ | $\alpha_1 \frac{S}{1+S} X - X$ | $\alpha_1 \frac{S}{1 + S + \frac{S^2}{K_I}} X - X$ |
| $\frac{dS}{dt}$ | $-\frac{S}{1+S}X - S + \alpha_2$ | $-\frac{S}{1+S+\frac{S^2}{K_I}}X - S + \alpha_2$ |
| XNC | $S = \frac{1}{\alpha_1 - 1}$ | $S = \frac{K_I(\alpha_1 - 1)}{2} \pm \sqrt{\left(\frac{K_I(\alpha_1 - 1)}{2}\right)^2 - K_I}$ |
| SNC | $X = \frac{(\alpha_2 - S)(1+S)}{S}$ | $X = \frac{(\alpha_2 - S)(1 + S + \frac{S^2}{K_I})}{S}$ |
| limit | _ | $K_I 	o \infty$ |

The other three models, the chemostat, the MMI-CSTR and the CPI-CSTR are quite similar in comparison to the CSI-CSTR.

Monods chemostat does not "feel" this inhibition and does not care...

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Appendix A

Set and variable list

 $D=1\ 2\ 3\ 4\ 5\ 6\ 7$

• • •

Appendix B

Parameter list

 $\begin{array}{c} A=1 \\ B=2 \end{array}$

Appendix C

The Linearized stability

C.1 The Linearization

F(x), a one-variable function of x can be Taylor-expanded around a fix X. We get $F(X+x)=F(X)+F'(X)x+O(x^2)$. For small perturbations of x around X we get the linearization: $F(X+x)\approx F(X)+F'(X)x$, containing only the constant and the linear terms.

For functions of two variables F(X + x, S + s) and G(X + x, S + s):

```
 \left\{ \begin{array}{l} F(X+x,S+s) = F(X,S) + F_X'(X,S)x + F_S'(X,S)s + O((x+s)^2) \\ G(X+x,S+s) = G(X,S) + G_X'(X,S)x + G_S'(X,S)s + O((x+s)^2) \end{array} \right.
```

```
function chemostat_inhibited(alpha1, alpha2, xp0, sp0, xc)
%chemostat_inhibited Displays a phaseportrait, nullclines
    and an Euler-path of an inhibited Chemostat.
    chemostat_inhibited(alfa1, alfa2, np0, cp0, nc) will run if
    alpha1 > 1/xc, thus there is a reproduction.
    alpha2 > 1/(xc*alpha1-1), thus there is sufficient stock-nutrition. xp0 > 0 , you can not have a nonpositive population.
    sp0 > 0 , you can not have a nonpositive concentration.
    The blue arrows represent the vectorfield.
    The black lines are two of the three nullclines.
    The black dotted line is the invariant line (no solution crosses it). The red line is an Eulerpath, starting in + and ending in \ast.
    chemostat_inhibited(5, 3, 0.2, 0.3, 6)
    by Per Erik Strandberg, 2003-2004.
% Start-condition:
if ((alpha1>1) & (alpha2>0) & (sp0>0) & (xp0>0) & xc>0),
    if (alpha2<1/(alpha1-1)),
         disp(' ')
disp (' (HINT: Only the trivial steady state, alpha2 is too small...)')
         disp (' (HINT: Two steady states, alpha2 is quite large...)')
```

```
% The illegal indata case:
%-----
else
    disp(' chemostat_inhibited.m by Per Erik Strandberg, 2003-2004.')
    disp(' Did not Finish OK. (You used illegal indata.)')
    disp(' ')
    disp(' For syntax help type: help chemostat_inhibited .')
    disp(' ')
end
```

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