INF-2700: Assignment 3

INF-2700: Database Systems

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November 5, 2023

1 Introduction

In this assignment we were tasked with implementing natural join using two different algorithm. The first was a nested loop join, the second was a block nested loop join. Below the design and implementation of these two algorithms for implementing the natural join operation will be explained.

In addition, the performance of each algorithm measured up against each other will be discussed in the experiments and results section and at the end there will be a section discussing if there is a different algorithm that would be better suited if the database was stored in a B+ tree organized file and how that algorithm would compare to the block nested loop join.

In order to run the program you must navigate to the correct directory (where you have downloaded the repository), run the "make" command, run ./run_front, then type the commands "test", "test2" and "select * from test_table natural join test_table2" and you will perform the natural join operation. By default the nested loop join is run. If you wish you test the block nested loop join, you must comment out line 1000 to 1025, and then uncomment line 1028 to line 1064. You must then save, run make cleanall and repeat the previous stated line of commands.

2 Design

[1, 2] The design of the natural join was largely influenced by the precode and the implementation of the previous assignments. In order to join to tables, these tables had to exist, so they were created and the tables needed schemas and records, so all this needed to be created, in addition to result records, schema and table that could hold the results. Since this had also been done in the previous assignments, the same recipe could be followed for this assignment. Finding the common field between the two was the first challenge for this implementation. Luckily it was specified that there only needed to be one common field, so the functions that created the test tables could be amended to have a common int field and a unique string field and in the function for the natural join. It would then be possible to loop through the fields in the tables to find the common one, then create a join table and add the fields from the other tables to that table.

In the book [1] page 704-707 details how to implement a nested-loop join and a block nested-loop join. For a nested-loop join, one simply uses a for loop to iterate over all the records in the left table, then another for loop to iterate over all the records in the right table and then check if the int field record value of the left table record is equal to the int field record value of the right table, if it is, the record can be added to the result table.

For the block nested-loop join, the design was intended to follow the book where now, instead of two four loops, the book describes four for loops. The first iterates over all the blocks in the left table, then the next

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for loop iterates over all the blocks in the right table, then the third for loop iterates over all the records in the blocks in the left table and the fourth for loop iterates over all the records in the right table.

3 Implementation

[1, 2] The table_natural_join function takes two table pointers as arguments. It performs a natural join operation between the two tables. The function first finds a common integer field between the two schemas. If no common integer field is found, an error message is printed. The function then creates a result schema that is a copy of the left schema and adds fields from the right schema that are not already in the result schema.

For the nested-loop join, while loops are utilized (the book features for loops, but iterating over the records in the table using a while loop and the get_record function was easier). The function then iterates over the records of both tables. If a record from the left table has the same value for the common field as a record from the right table, a new record is created in the result schema.

The block nested-loop first checks if both tables exist. If not, an error message is printed. The function then declares variables for iterating over the blocks and pages of the tables. It sets the position of both tables to the beginning and iterates over the records of both tables. If a record from the left table has the same value for a field as a record from the right table, a new record is created in the result schema.

4 Experiments and Results

[1] In order to test the performance of the two algorithms against each other, the number of records in each of the two tables were set to 1000. The nested loop join had the following number of disk seeks/reads/writes/IOs: 60/56/105/161. In comparison, the block nested loop join had the following number of disk seeks/reads/writes/IOs: 4/2/56/58. The block nested loop join has significantly fewer disk seeks/reads/writes/IOs than the nested loop join. This indicates that the block nested loop join is more efficient in terms of I/O operations and will likely have better performance than the nested loop join.

Fewer I/O operations mean that less time is spent reading from or writing to disk, which can significantly speed up queries and other operations. That is why database systems often use techniques like indexing, partitioning, and in-memory processing to minimize I/O operations.

5 B+ Tree organized database

[1, 3] A possible algorithm that could be used is a sort-merge join, which is a general technique for joining two sorted lists of data. One way to use a sort-merge join could be by doing the following. First, you would sort both tables by the join column using the B+ tree index. This would be done by traversing the leaf nodes of the B+ tree in order. Then, merge the two sorted tables by comparing the values in the join column and outputting the matching rows. This would be done by using two pointers that scan the tables from left to right and advance according to the comparison result.

Now as for comparing this to a block nested-loop join, there are advantages and disadvantages for both, depending on the size, order and index of the tables. A sort-merge join algorithm is more efficient for larger tables, especially if they are already sorted or have an index on the join column. A block nested-loop join algorithm is more suitable for small tables or when one table is much smaller than the other. A sort-merge join algorithm can produce an ordered output, while block nested-loop join algorithm might not give an ordered output. Sort-merge also requires more memory than a block nested-loop, which could be an issue depending on system specs.

6 Sources

References

- [1] silberschatz Silberschatz, A., Korth, H. F.,, Sudarshan, S. (2010). Database system concepts. New York: McGraw-Hill.
- [2] nested loop join Nested-loop join algorithm. Retrieved 16:45, October 29, 2023, from https://dev.mysql.com/doc/refman/8.0/en/nested-loop-joins.html
- [3] sort-merge join Sort-merge join algorithm. Retrieved 18:30, November 5th, 2023, from https://en.wikipedia.org/wiki/Sort-merge_join