

Discrete Geometry – Third Assignment

§1. LOWER ENVELOPES

LEMMA 1.1. Let f, g be monotonically increasing functions. Then $f \circ g$ is monotonically increasing.

PROOF OF LEMMA 1.1. Let $x \leq y$. As g is monotonically increasing it follows that $g(x) \leq g(y)$. Moreover, as f is monotonically increasing it follows that $f(g(x)) \leq f(g(y))$. ■

PROBLEM 1. For every c there exists an n_0 such that for all $n \geq n_0$ we have that $A_k(n) \geq A_{k-1}^{(c)}(n)$.

SOLUTION FOR PROBLEM 1. It suffice to show that $A_{k-1}^{(c)}(n) = o(A_k(n))$. Note that

$$\begin{aligned} \lim_{n \rightarrow \infty} \frac{A_k(n)}{A_{k-1}^{(c)}(n)} &= \lim_{n \rightarrow \infty} \frac{A_{k-1}(A_k(n-1))}{A_{k-1}^{(c)}(n)} = \lim_{n \rightarrow \infty} \frac{A_{k-1}(A_{k-1}(A_k(n-2)))}{A_{k-1}^{(c)}(n)} = \dots = \lim_{n \rightarrow \infty} \frac{A_{k-1}^{(c)}(A_k(n-c))}{A_{k-1}^{(c)}(n)} \\ &\geq \lim_{n \rightarrow \infty} \frac{2^c \cdot A_k(n-c)}{2^c \cdot n} = \lim_{n \rightarrow \infty} \frac{A_k(n-c)}{n} \geq \lim_{n \rightarrow \infty} \frac{2(n-c)}{n} = \lim_{n \rightarrow \infty} n - \frac{2c}{n} = \infty, \end{aligned}$$

where the first and the second inequalities hold due to the monotonicity of Ackermann's function. Indeed, $A_1(n) = 2n$ and therefore monotonically increasing. In addition, Lemma 1.1 implies that the composition of monotonically increasing functions is also monotonically increasing. ■

PROBLEM 2. Let $A(n) = A_n(3)$. Prove that for every $k \in \mathbb{N}$ there exists $n_0 \in \mathbb{N}$ such that for every $n \geq n_0$ it holds that $A(n) \geq A_k(n)$.

SOLUTION FOR PROBLEM 2. Let n_0 satisfy $A_{n_0}(2) \geq n_0, k$. It then follows that for every $n \geq n_0$ it holds that

$$A_n(3) = A_{n-1}(A_n(2)) \geq A_{n-1}(n) \geq A_k(n),$$

Where both inequality follows from the chose of n_0 . Note that Corollary A.5 implies that $f_k := A_k(n)$ is unbounded and therefore such choice for n_0 is possible. ■

PROBLEM 3. In the lecture we defined $x_k(m) := e_k(m)/n_k(m)$. We claimed that $x_k(m) \geq k/2$ and the proof was by induction on k and for each k by induction on m .

1. Complete the base case where $k = 1$.
2. Complete the base case where $m = 1$.
3. Check the details.

SOLUTION FOR PROBLEM 3.

1. Applying to the definition of $S_1(m)$, we have that the length of the lower envelope is m . That is, $e_1(m) = m$. In addition, the number of fans is defined to be 1, i.e., $f_1(m) = 1$. It

follows that

$$x_1(m) = \frac{e_k(m)}{n_k(m)} = \frac{e_k(m)}{f_k(m) \cdot m} = \frac{m}{m} = 1 \geq \frac{1}{2} = \frac{k}{2}.$$

2. Following the recurrences relation that was stated in the class we have the following relations:

$$f_k(1) = 2f_{k-1}(2) \quad \text{and} \quad e_k(1) = e_{k-1}(2) + 2f_{k-1}(2).$$

Therefore,

$$x_k(1) = \frac{e_k(1)}{n_k(1)} = \frac{e_k(1)}{f_k(1)} = \frac{e_{k-1}(2) + 2f_{k-1}(2)}{2f_{k-1}(2)} = \frac{e_{k-1}(2)}{n_{k-1}(2)} + 1 = x_{k-1}(2) + 1.$$

The induction hypothesis implies that $x_k(1) \geq k/2$.

3. The general case where $k, m \geq 2$. The way that $S_k(m)$ was constructed implies that $f_k(m) = f_{k-1}(M) \cdot M$ where $M = f_k(m-1)$ is the number of fans in S' . We note that $M-1 \geq M/2$ as $M \geq 2$. Applying to the recurrence relation that was proven in the class we have

$$e_k(m) = e_k(m-1) \cdot f_{k-1}(M) + e_{k-1}(M) + (M-1) \cdot f_{k-1}(M).$$

To see this note that we have $f_{k-1}(M)$ copies of S' each contributes $e_k(m-1)$ to the length. In addition, we should consider the contribution of S^* , that is, $e_{k-1}(M)$. Last but not least, one should take into account the $M-1$ additional parts gained from this clever (yet tedious) construction at each fan. It follows that

$$\begin{aligned} e_k(m) &= e_{k-1}(M) + f_{k-1}(M) (e_k(m-1) + M-1) \\ &\geq e_{k-1}(M) + f_{k-1}(M) (e_k(m-1) + M/2). \end{aligned}$$

Dividing both sides by $n_k(m)$ yields:

$$\begin{aligned} x_k(m) &\geq \frac{e_{k-1}(M) + f_{k-1}(M) (e_k(m-1) + M/2)}{n_k} \\ &\geq \frac{e_{k-1}(M) + f_{k-1}(M) (e_k(m-1) + M/2)}{m \cdot f_k(m)} \\ &\geq \frac{e_{k-1}(M) + f_{k-1}(M) (e_k(m-1) + M/2)}{mM \cdot f_{k-1}(M)} \\ &\geq \frac{e_k(m-1) + M/2}{mM} + \frac{e_{k-1}(M)}{m \cdot n_k} \\ &\geq \frac{e_k(m-1)}{mM} \cdot \frac{n_k(m-1)}{n_k(m)} + \frac{1}{2m} + \frac{x_{k-1}(M)}{m} \\ &\geq \frac{e_k(m-1)}{mM} \cdot \frac{(m-1) \cdot M}{n_k(m-1)} + \frac{1}{m} \cdot \left(\frac{1}{2} + x_{k-1}(M) \right) \\ &\geq \frac{e_k(m-1)}{n_k(m-1)} \cdot \frac{m-1}{m} + \frac{1}{m} \cdot \left(\frac{1}{2} + x_{k-1}(M) \right) \\ &\geq \left(1 - \frac{1}{m} \right) \cdot x_k(m-1) + \frac{1}{m} \cdot \left(\frac{1}{2} + x_{k-1}(M) \right). \end{aligned}$$

The induction hypothesis now implies that $x_k(m-1) \geq k/2$ and that $x_{k-1}(M) \geq (k-1)/2$. Concluding that

$$\left(1 - \frac{1}{m}\right) \cdot \frac{k}{2} + \frac{1}{m} \left(\frac{1}{2} + \frac{k}{2} - \frac{1}{2}\right) \geq \frac{k}{2} \left(1 - \frac{1}{m} + \frac{1}{m}\right) = \frac{k}{2},$$

as claimed. ■

PROBLEM 4. Prove that if d, D are given in advance, and if r is chose large enough, then $x_{23} < x_{12}$.

SOLUTION FOR PROBLEM 4. Let ℓ_1, ℓ_2, ℓ_3 be the three segment from left to right as describe in Figure 1. Assume without loss of generality that the bottom left corner of the left segment (denoted by ℓ_1) is $p_1 = (0, 0)$ as describe in Figure 1.

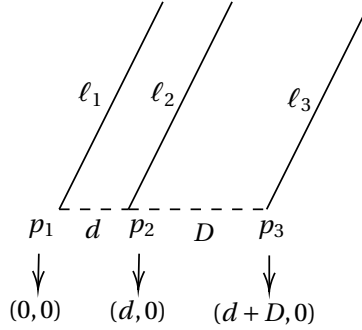


Figure 1: A schematic illustration.

As we saw in the lecture, we can assume that the slopes of the lines are $1 + \varepsilon m_1, 1 + \varepsilon m_2, 1 + \varepsilon m_3$, where $\varepsilon > 0$. We start by finding the lines:

$$\ell_1: y_1(x) = (1 + \varepsilon m_1)x, \quad \ell_2: y_2(x) = (1 + \varepsilon m_2) \cdot (x - d), \quad \ell_3: y_3(x) = (1 + \varepsilon m_3) \cdot (x - (d + D)).$$

Now, we find the intersection between the points that satisfies ℓ_1 and ℓ_2 :

$$\begin{aligned} y_1(x) &= y_2(x) \\ (1 + \varepsilon m_1)x &= (1 + \varepsilon m_2) \cdot (x - d) \\ x(1 + \varepsilon m_2 - (1 + \varepsilon m_1)) &= d \cdot (1 + \varepsilon m_2) \\ x\varepsilon(m_2 - m_1) &= d \cdot (1 + \varepsilon m_2) \end{aligned}$$

$$x_{12} = \frac{d \cdot (1 + \varepsilon m_2)}{\varepsilon \cdot (m_2 - m_1)}$$

Similarly we find x_{23} , namely, the intersection of ℓ_2, ℓ_3 :

$$\begin{aligned}
 (1 + \varepsilon m_2) \cdot (x - d) &= (1 + \varepsilon m_3) \cdot (x - (d + D)) \\
 x(1 + \varepsilon m_2 - (1 + \varepsilon m_3)) &= d \cdot (1 + \varepsilon m_2 - (1 + \varepsilon m_3)) - D \cdot (1 + \varepsilon m_3) \\
 \varepsilon x(m_2 - m_3) &= \varepsilon d(m_2 - m_3) - D(1 + \varepsilon m_3) \\
 \boxed{x_{23} = d + D \cdot \frac{(1 + \varepsilon m_3)}{\varepsilon(m_3 - m_2)}}
 \end{aligned}$$

As the question ask for satisfying $x_{23} < x_{12}$ we have the constrain:

$$\begin{aligned}
 x_{23} = d + \frac{D \cdot (1 + \varepsilon m_3)}{\varepsilon(m_3 - m_2)} &< \frac{d \cdot (1 + \varepsilon m_2)}{\varepsilon \cdot (m_2 - m_1)} = x_{12} \\
 d(m_3 - m_2)[\varepsilon(m_2 - m_1) - (1 + \varepsilon m_2)] &< -D \cdot (1 + \varepsilon m_3)(m_2 - m_1) \\
 D \cdot (1 + \varepsilon m_3)(m_2 - m_1) &< d(m_3 - m_2)[(1 + \varepsilon m_2) - \varepsilon(m_2 - m_1)]
 \end{aligned}$$

Reordering,

$$\begin{aligned}
 \frac{D}{d} &< \frac{(m_3 - m_2)[(1 + \varepsilon m_2) - \varepsilon(m_2 - m_1)]}{(1 + \varepsilon m_3)(m_2 - m_1)} \\
 &= (m_3 - m_2) \cdot \left(\frac{1 + \varepsilon m_2}{(1 + \varepsilon m_3)(m_2 - m_1)} - \frac{\varepsilon}{1 + \varepsilon m_3} \right) \\
 &< \frac{m_3 - m_2}{m_2 - m_1} \cdot \underbrace{\frac{1 + \varepsilon m_2}{1 + \varepsilon m_3}}_{\leq 1} \\
 &\leq \frac{m_3 - m_2}{m_2 - m_1} \\
 &= r
 \end{aligned}$$

To see the last equality note that

$$\frac{m_3 - m_2}{m_2 - m_1} = \frac{\frac{m_3 - m_2}{m_2}}{\frac{m_2 - m_1}{m_2}} = \frac{r - 1}{1 - \frac{1}{r}} = \frac{r - 1}{\frac{r - 1}{r}} = r.$$

We established the following relation:

$$\boxed{r(d, D) \geq \frac{D}{d}}$$

■

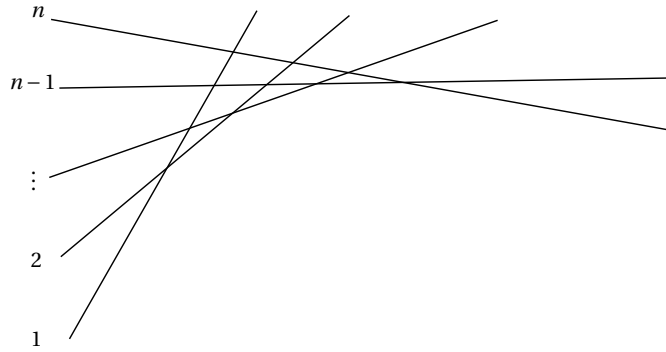
Let $\lambda_s(n)$ denote the maximum length of a Davenport-Schinzel sequence of order s .

PROBLEM 5. Prove that:

1. $\lambda_1(n) = n$.
 2. $\lambda_2(n) = 2n - 1$.
- **HINT (Upper-Bound).** Let z be the last symbol to appear for the first time. Show that z can appear only one. Delete z and use induction on n .

SOLUTION FOR PROBLEM 5.

1. We start by showing that $\lambda_1(n) \geq n$. Consider the following Davenport-Schinzel sequence of order 1 of length n : $\langle 1, 2, 3, \dots, n \rangle$.



The upper bound is trivial. The appearance of any additional segment would form an $a \dots b \dots a$ for some $a, b \in \Sigma = \{1, \dots, n\}$, and hence negates the definition of Davenport-Schinzel sequence of order 1.

2. **Upper-Bound.** As the hint suggests, we prove the claim by induction on n . For $n = 1$ then there is only one element, and indeed $\lambda_2(1) = 2 \cdot 1 - 1 = 1$. Assuming $\lambda_2(n-1) = 2n-3$ and let A be a Davenport-Schinzel sequence of order 2. Let z be the last symbol to appear for the first time and let y be the successive element in the sequence. We show that z appears only once in A . Assume towards a contradiction that the claim is false, then A cannot be a Davenport-Schinzel sequence of order 2 as it forms (see Figure 2) the configuration $(\dots y \dots zy \dots y \dots)$ which violates the definition of Davenport-Schinzel sequence of order 2.

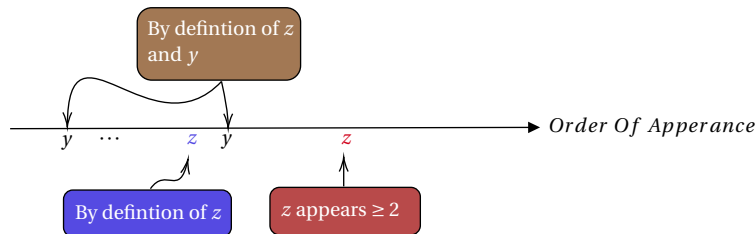


Figure 2: The assumption that z at least twice yields the configuration $y \rightarrow z \rightarrow y \rightarrow z$.

Removing the single appearance of z from U can yield a violation of the definition of Davenport-Schinzel sequence if the two neighbours of z share the same symbol. In this

case, choose one of the neighbours arbitrarily and remove also the chosen neighbour. The resulting sequence is a Davenport-Schinzel sequence of order 2 with $n - 1$ segments (cause we remove the segment z). It follows by the induction hypothesis that the maximum length of A' is $2(n - 1)$, and by the way A' was formed, we have that A' is shorter by at most 2 elements from A . Hence, the maximum length of A is at most $2n - 1$, as claimed.

Lower-Bound. The sequence $(1, 2, \dots, n - 1, n, n - 1, \dots, 2, 1)$ is a Davenport-Schinzel sequence of order 2 with n segments of length $2n - 1$.

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§2. TVERBERG'S THEOREM

PROBLEM 6. Write that last part of the proof formally (and not just with an example).

Proof. We adopt the notation from the slides and note that for every $1 \leq i \leq N + 1$ it holds that $\sum_{j=1}^r \hat{q}_{i,j} = \hat{0}$. As we want to show that $\hat{0} \in \text{conv}(q_{i,r}, \dots, q_{i,r})$, we normalize the last equation by dividing by r and driving that there exists a convex combination of $\hat{q}_{i,r}, \dots, \hat{q}_{i,r}$ which forms $\hat{0}$. The colorful Carathéodory Theorem implies that there exists a sequence j_1, \dots, j_{N+1} such that $\hat{0} \in \text{conv}(\hat{q}_{1,j_1}, \dots, \hat{q}_{N+1,j_{N+1}})$. Hence there exists scalars $\alpha_1, \dots, \alpha_{N+1} \geq 0$ such that $\sum_{i=1}^{N+1} \alpha_i = 1$ and $\hat{0} = \sum_{i=1}^{N+1} \alpha_i \hat{q}_{i,j_i}$. As vector are equal iff the entries are element-wise equals we have

$$0 = \sum_{i: j_i=k} \alpha_i q_i + \sum_{i: j_i=r} \alpha_i \cdot (-q_i) \iff \sum_{i: j_i=r} \alpha_i q_i = \sum_{i: j_i=k} \alpha_i q_i,$$

for every $1 \leq k \leq r - 1$. Recall that each $q_i = (p_i, 1)$ and hence the sum of the coefficient at each side of the equation sums to the same quantity, say λ . Hence, dividing all equations by λ yields a point $\rho := \frac{1}{\lambda} \cdot \sum_{i: j_i=r} \alpha_i q_i$ such that ρ can be expressed as r convex combinations, say $\text{conv}(\mathcal{G}_1), \dots, \text{conv}(\mathcal{G}_r)$ where $\cup_{i=1}^r \mathcal{G}_i = \mathcal{U} := \{p_1, \dots, p_{N+1}\}$, i.e. $\mathcal{G}_1, \dots, \mathcal{G}_r$ is a partition of \mathcal{U} .

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PROBLEM 7. Let $p_1, \dots, p_{d+1} \in \mathbb{R}^d$ form the vertices of a simplex. Let q be the center of the simplex (i.e., $q = (p_1 + \dots + p_{d+1})/(d + 1)$). Replace each p_i by a tiny cloud of $r - 1$ points. What are **all** the Tverberg partitions of the set? And how many Tverberg partitions are there?

Proof. Note that the number of points is $(r - 1) \cdot (d + 1) + 1$, that is a Tverberg's number. For the first set in the partition there are $(r - 1)^{d+1}$ options, since in each cloud there are $(r - 1)$ option and there are $(d + 1)$ clouds. Once it has been chosen, the second set in the partition has $(r - 2)^{d+1}$ as at each cloud one vertex has been discarded cause repetition is not allowed in a partition. Finally, the last set has $1^{d+1} = 1$ options. The element q is also taken as a set in the partition, i.e., the set $\{q\}$ is a set in the partition. Therefore, we have

$$(r - 1)^{d+1} \cdot (r - 2)^{d+1} \cdot \dots \cdot 2^{d+1} \cdot 1^{d+1} = (1 \cdot 2 \cdot \dots \cdot r - 1)^{d+1} = ((r - 1)!)^{d+1} \text{ Options.}$$

■

§A. THE ACKERMANN'S HIERARCHY

LEMMA A.1. *Prove that $A_k(n+1) \geq A_k(n) + 2$ whenever $k, n \in \mathbb{N}$.*

COROLLARY A.2. *$A_k(n) \geq 2n$ whenever $k, n \in \mathbb{N}$.*

PROOF OF LEMMA A.1. We prove it using induction over $(k, n) \in \mathbb{N} \times \mathbb{N}$, where $(a, b) < (c, d)$ if either $a < c$, or $a = c$ and $b < d$. Observe that the claim holds for $k, n = 0$ as $A_0(1) = 2 \geq 0 + 2 = A_0(0) + 2$. Assume the claim holds for every $(k', n') < (k, n)$ and show for (k, n) . Note that

$$A_k(n+1) = A_{k-1}(A_k(n)) \geq 2 \cdot A_k(n) \geq A_k(n) + 2n \geq A_k(n) + 2.$$

■

COROLLARY A.3. *Prove that $A_k(n+1) > A_k(n)$ whenever $k, n \in \mathbb{N}$.*

LEMMA A.4. *Prove that $A_{k+1}(n) \geq A_k(n)$ whenever $m, n \in \mathbb{N}$.*

PROOF OF LEMMA A.4. Note that $A_1(0) = 0 = A_0(0)$. Assume the claim holds for all $(k', n') < (k, n)$ and show for (k, n) . To that end, note

$$A_{k+1}(n) = A_k(A_{k+1}(n-1)) \geq A_k(2(n-1)) \geq A_k(n),$$

where the first inequality is due to Lemma A.1. ■

COROLLARY A.5. *Fix $n \in \mathbb{N}$. Then $f_k = A_k(n)$ is unbounded.*

PROOF OF COROLLARY A.5. In the proof of Lemma A.4 we showed that:

$$A_{k+1}(n) = A_k(A_{k+1}(n-1)) \geq A_k(2(n-1)).$$

Continue this fashion one obtains:

$$A_{k+1}(n) \geq A_k(2(n-1)) \geq A_{k-1}(2^2(n-1)) \geq \dots \geq A_{k-k'}(2^{k'}(n-1)) \geq A_1(2^k(n-1)) = 2^{k+1}(n-1).$$

Clearly, as the last expression in the right hand side tends to infinity as k tends to infinity the claim follows. ■

REFERENCES

- [1] J. Matoušek. *Lectures on discrete geometry*, volume 212 of *Graduate Texts in Mathematics*. Springer-Verlag, New York, 2002.