

A Composable Security Treatment of the Lightning Network

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Part 1

The Lightning Network



VISA

20,000 tx/s

 **bitcoin**

7 tx/s

Problem

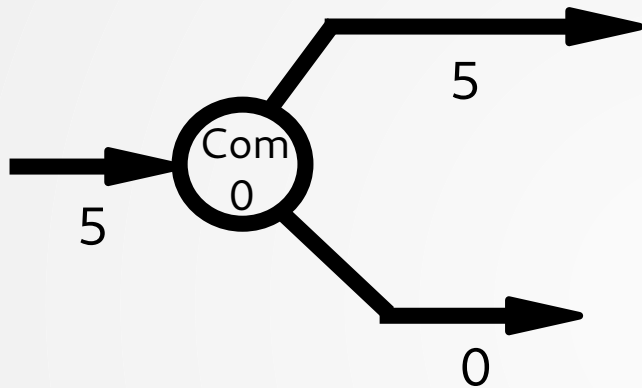
All txs validated by all wallets

Solution


- Move most txs off-chain
- Resolve disputes on-chain



Alice
→
5



Alice





Alice
& Bob



5

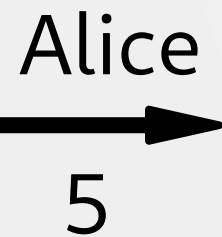
Com
0



5



0



5



Alice
& Bob

5

Com
0

5

Bob

0



Alice

5



Alice
& Bob

5



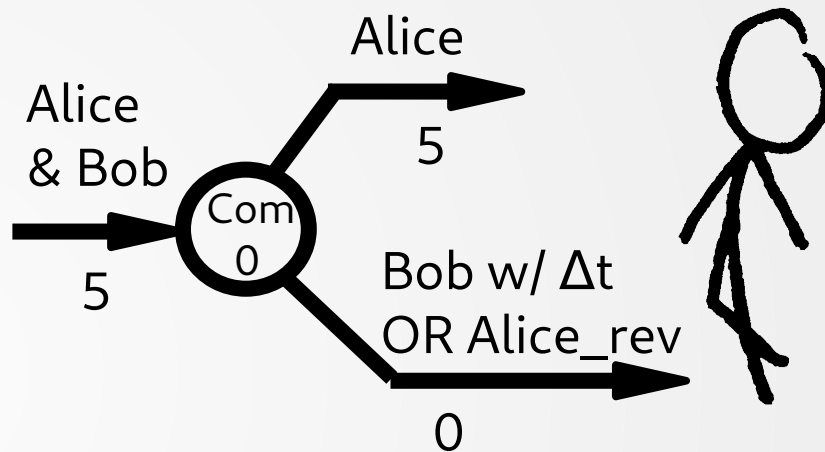
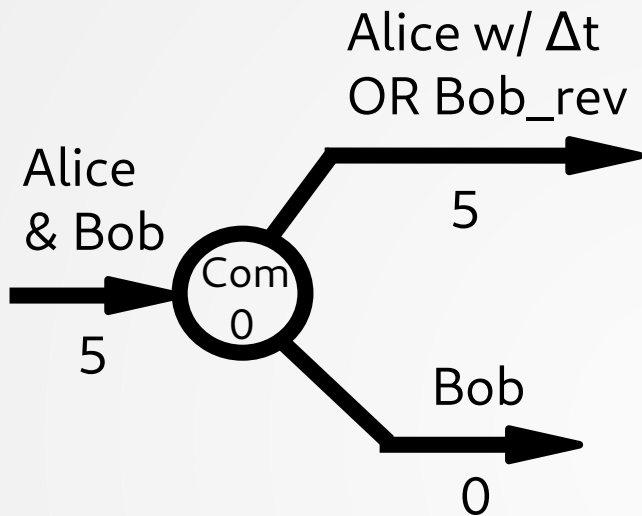
Alice w/ Δt
OR Bob_rev

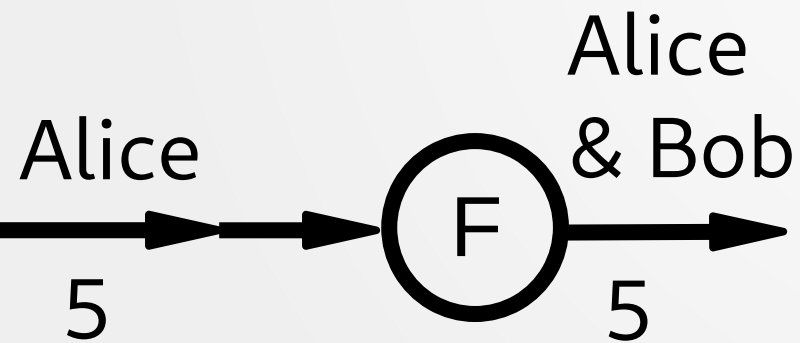
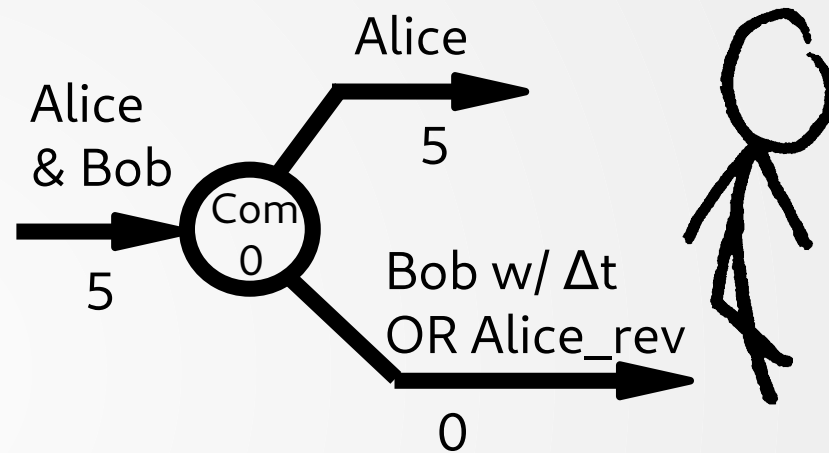
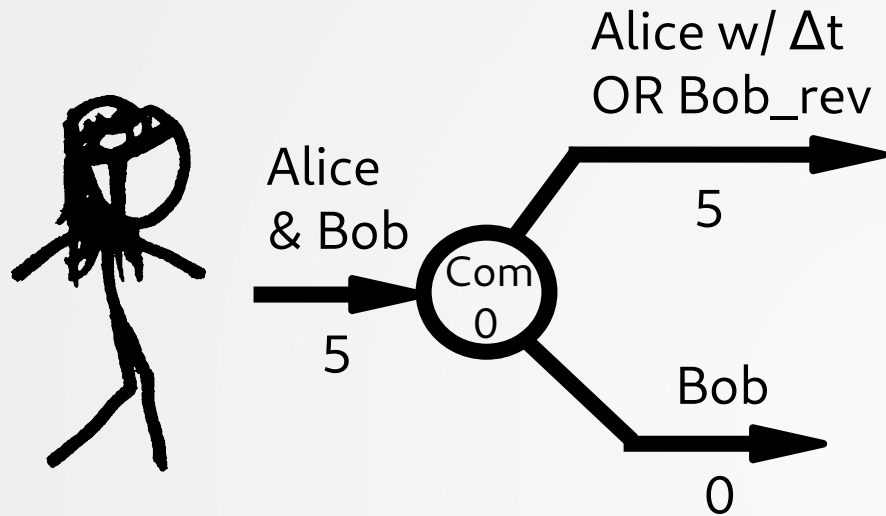
5

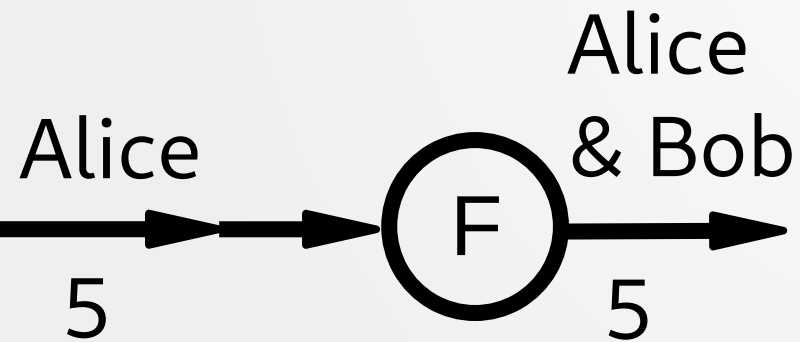
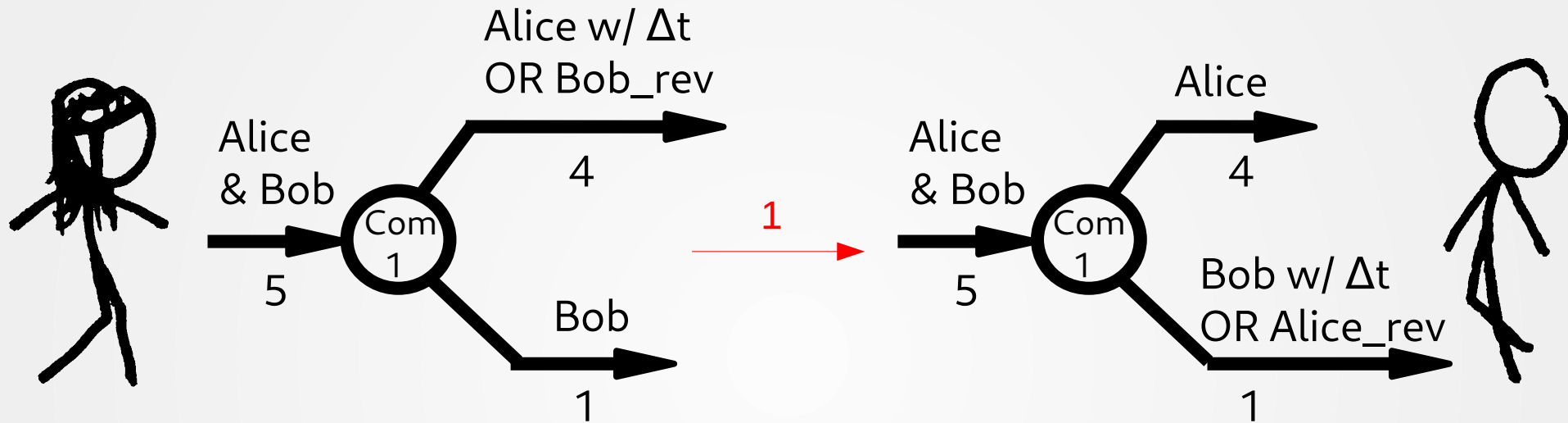
Bob

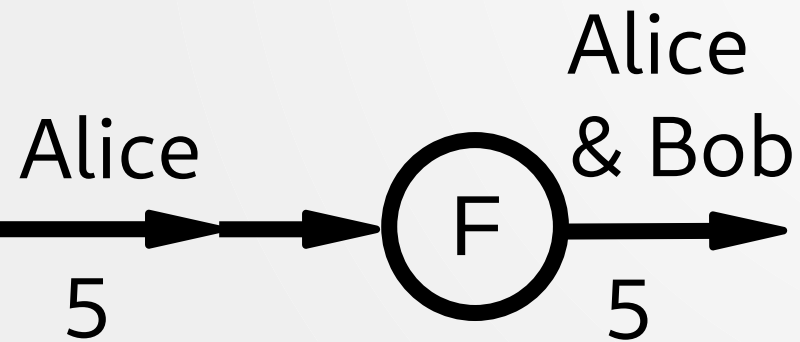
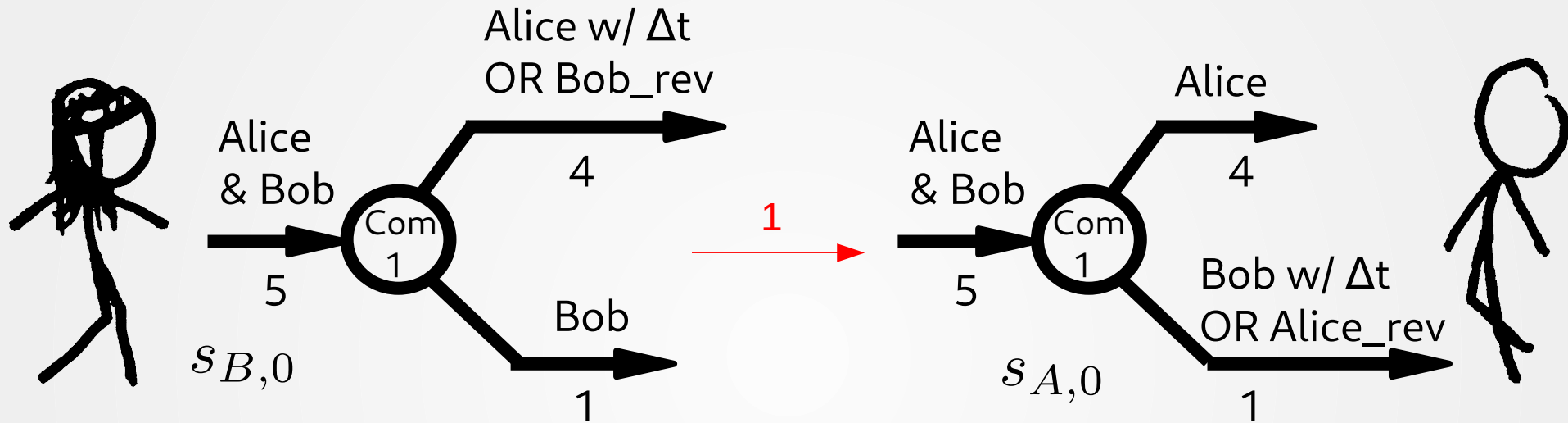
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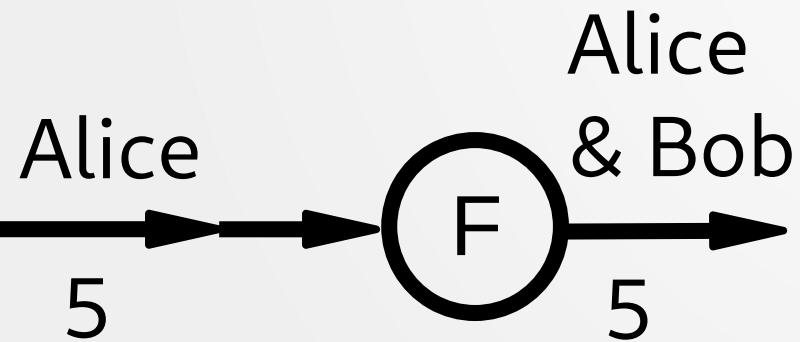
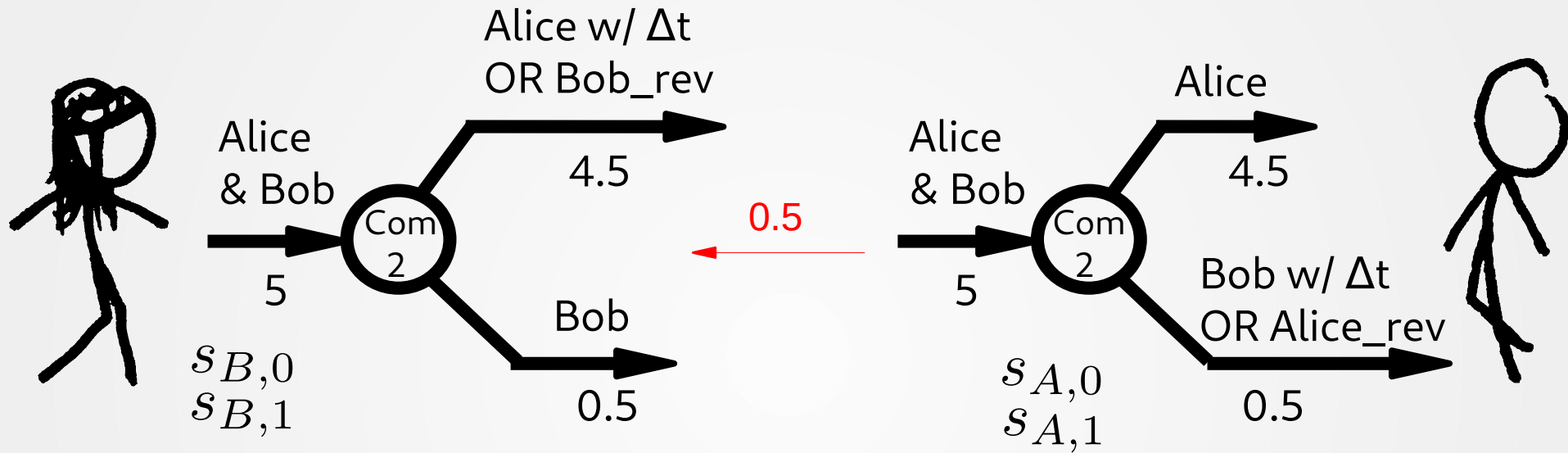


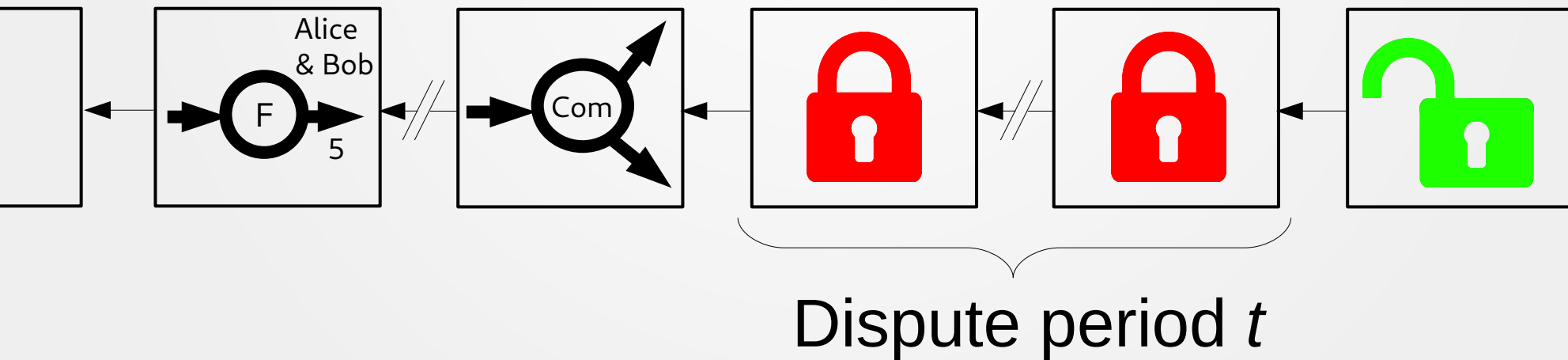
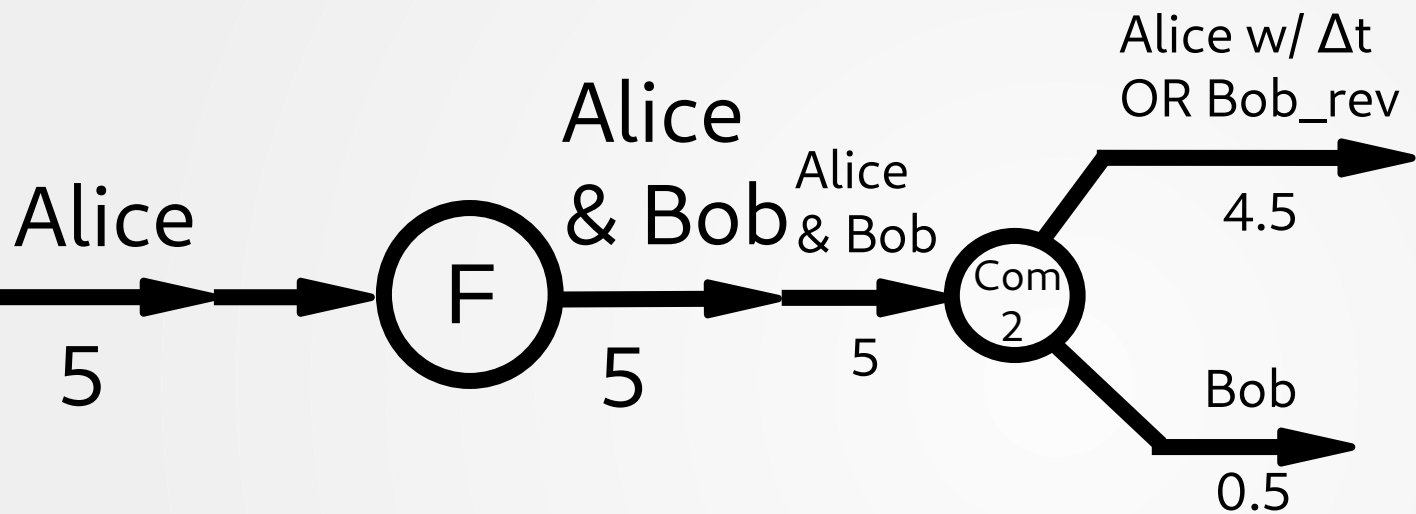


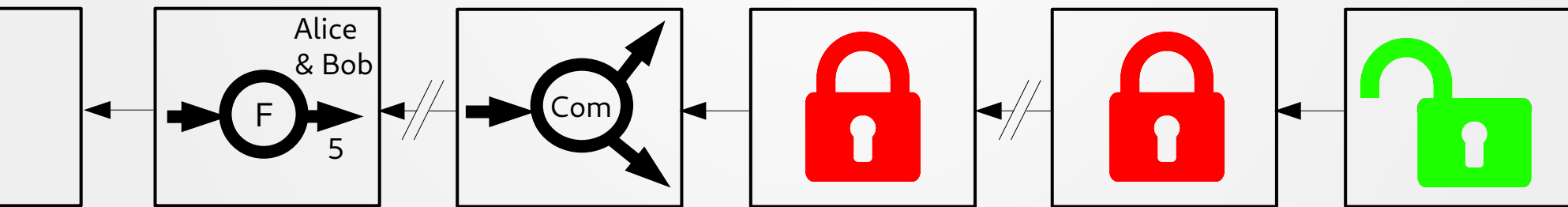
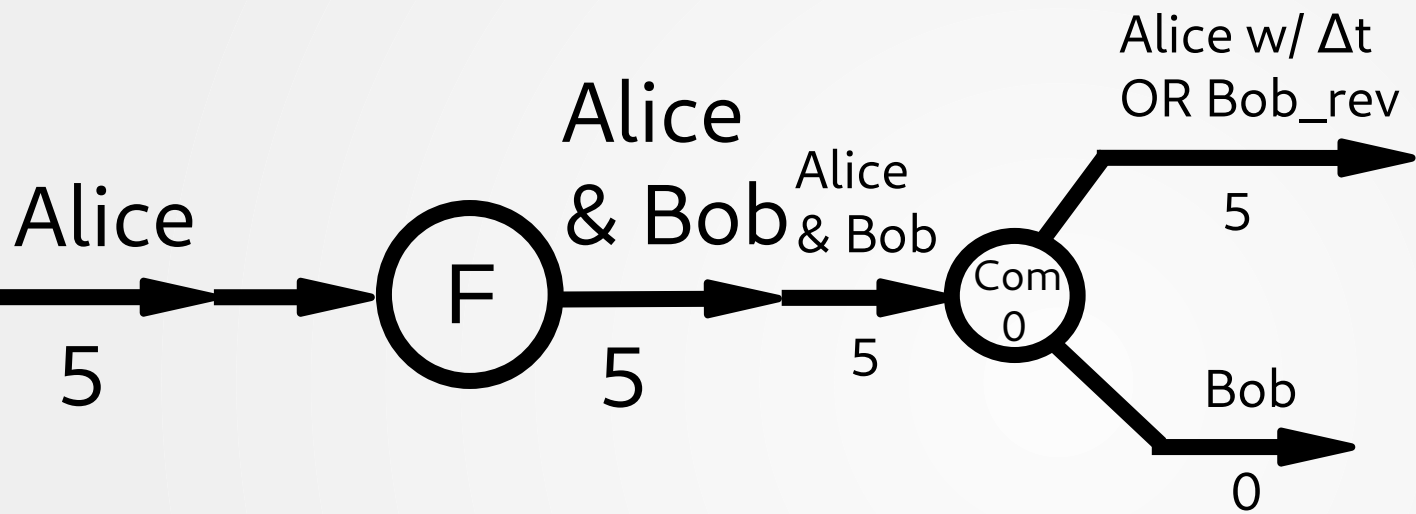




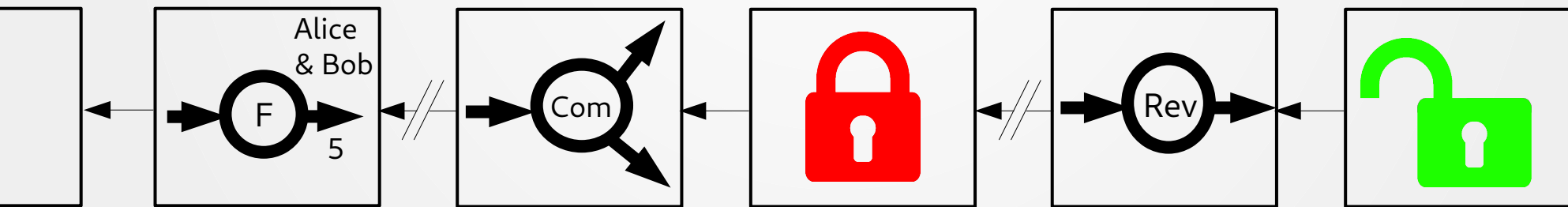
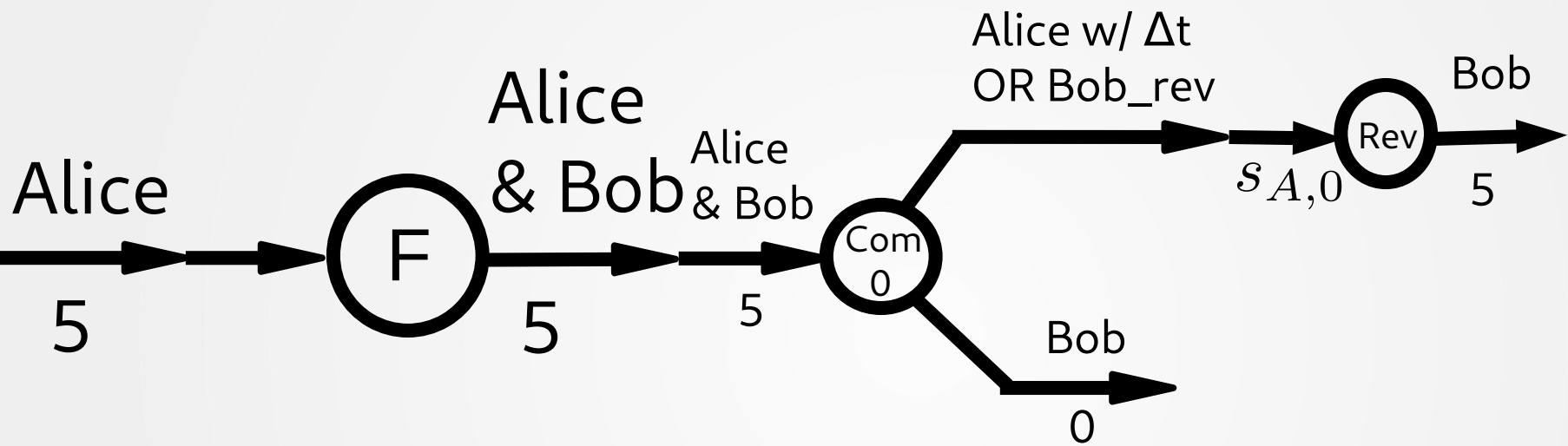






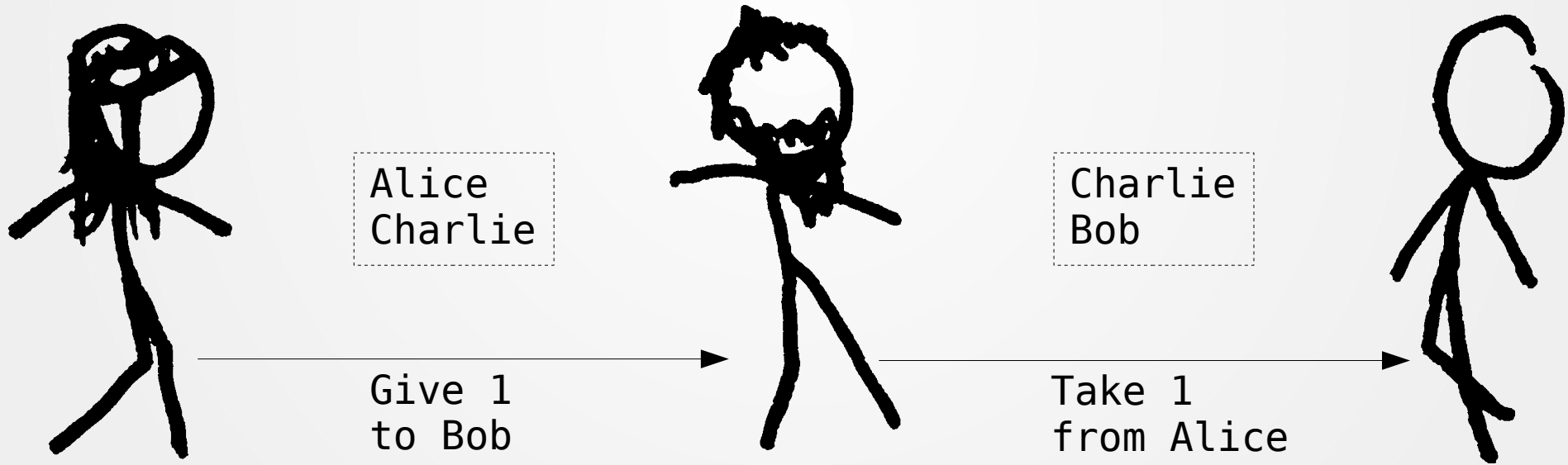


Dispute period t



Dispute period t

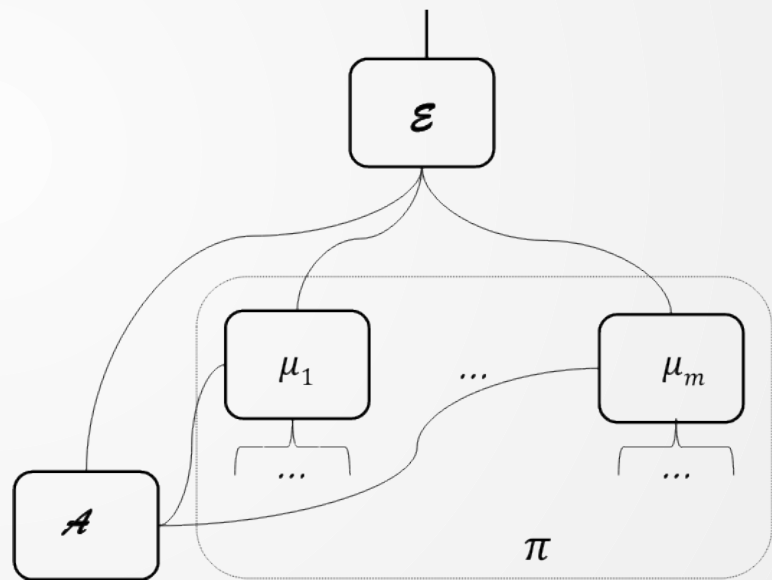
Multi-hop payments



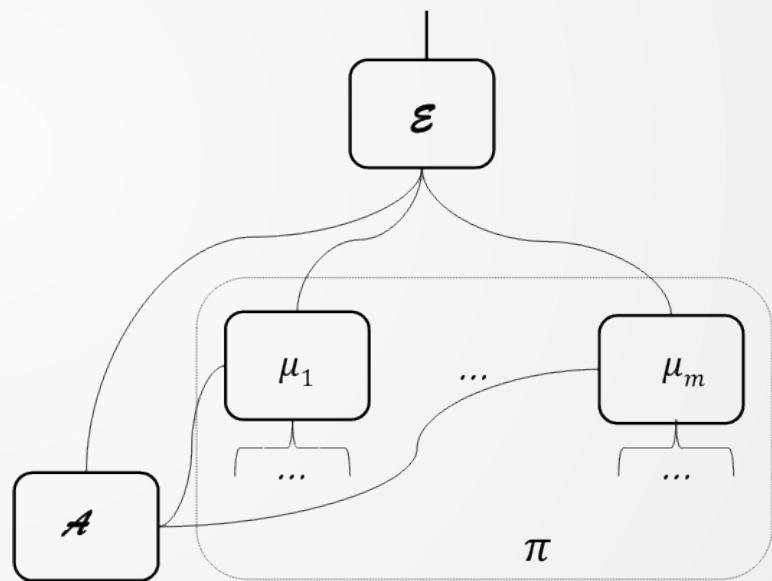
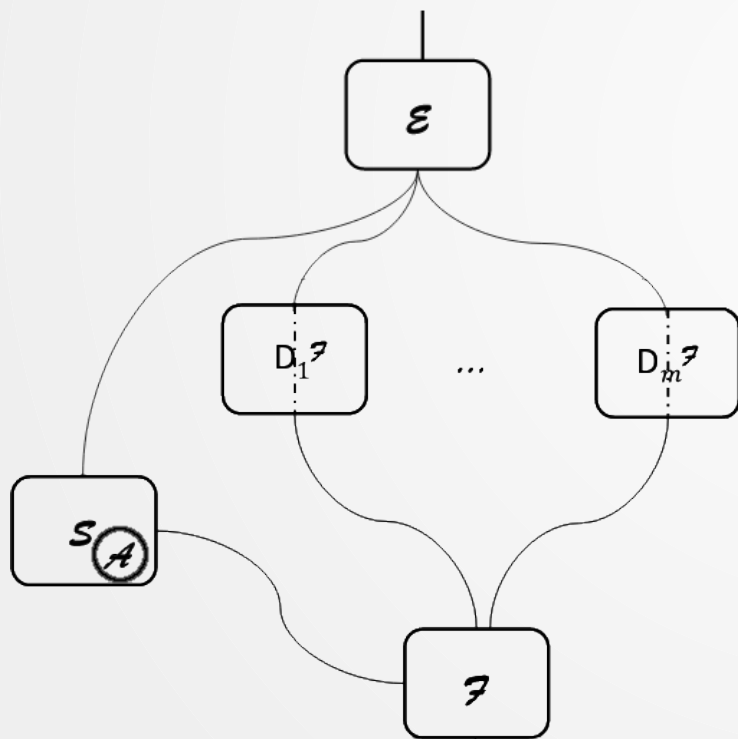
Part 2

Our contribution

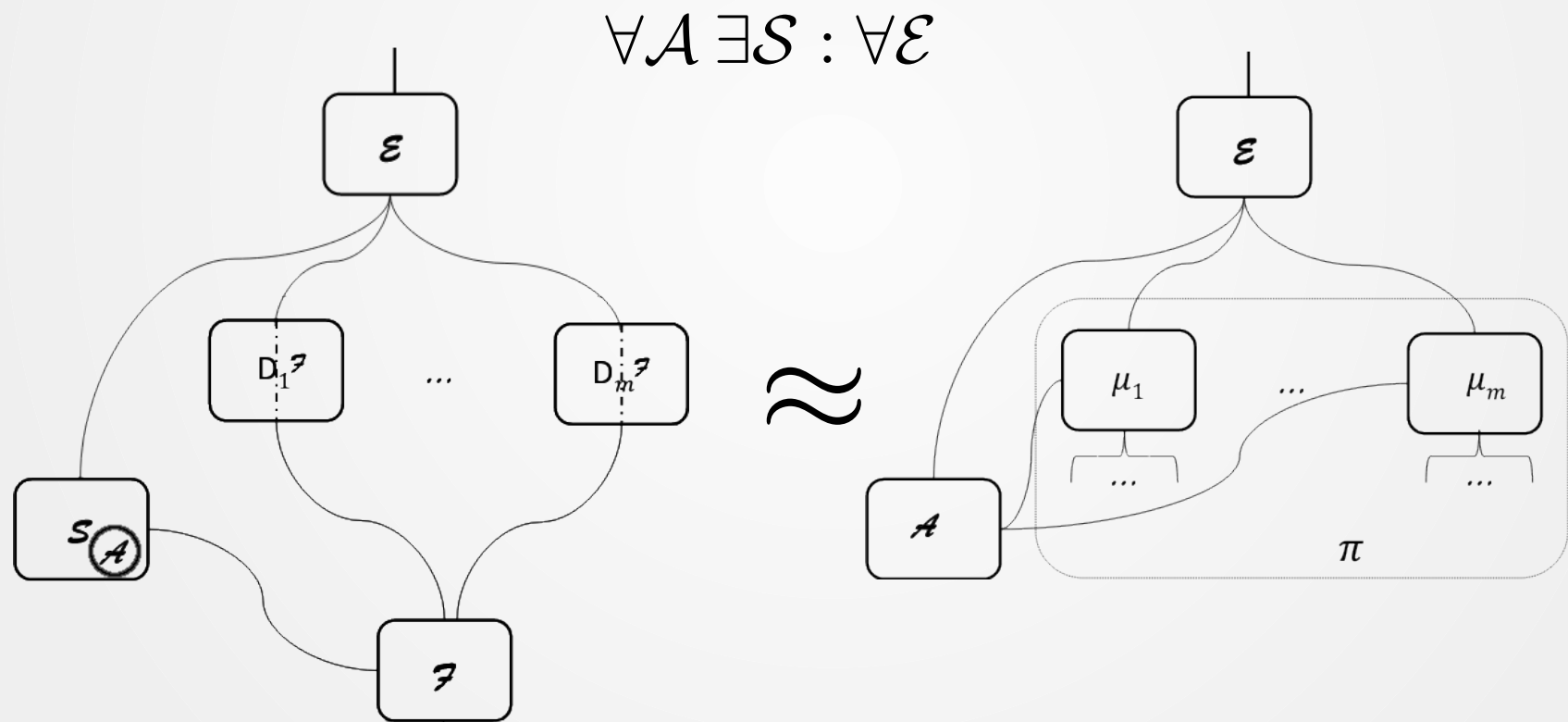
Simulation-based Security



Simulation-based Security



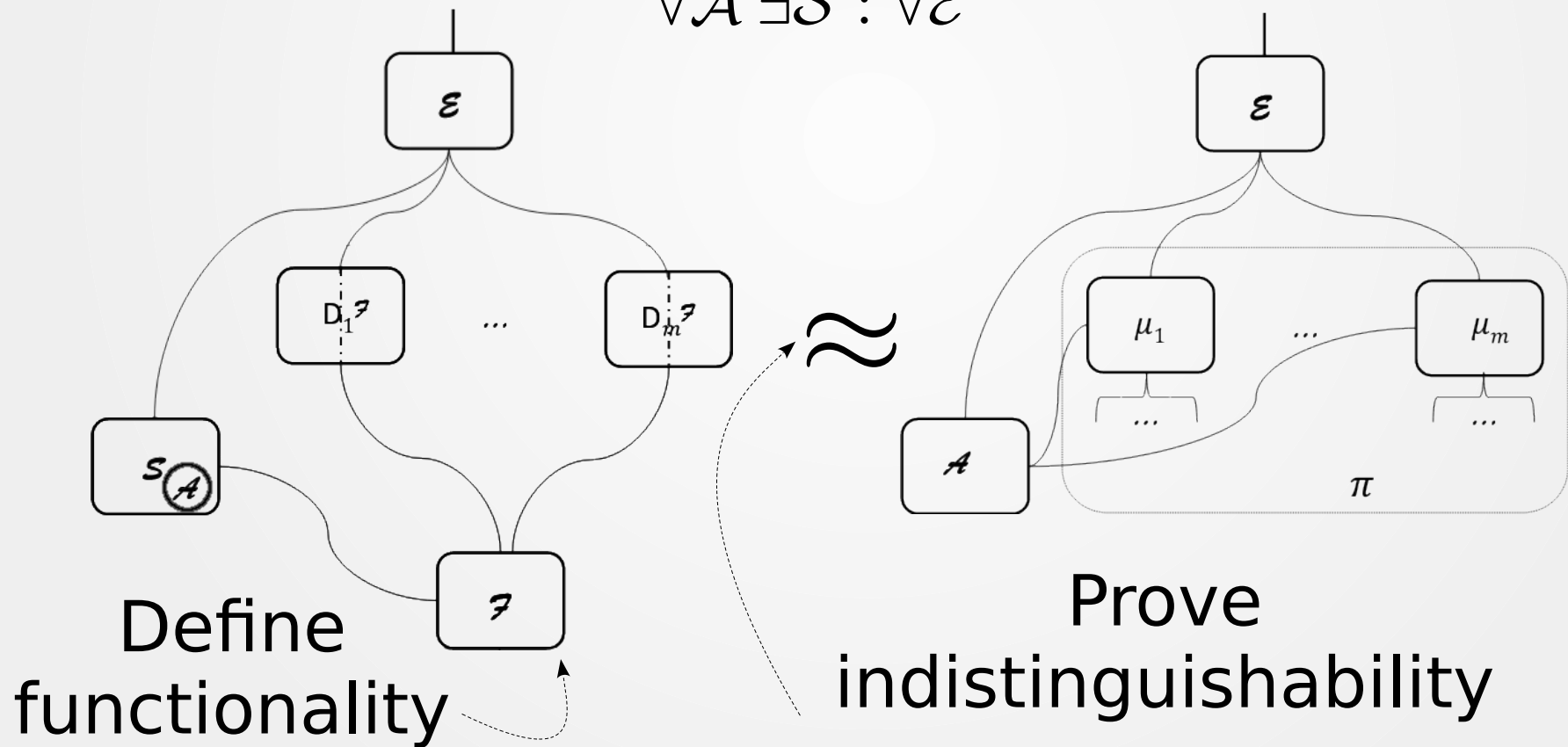
Simulation-based Security



Credits: "Universally Composable Security", Ran Canetti
<https://eprint.iacr.org/2000/067>

Our paper

$$\exists A : SE A \forall$$



Universal Composition

If functionality F is UC-realized by protocol P , then a protocol R that internally uses F is indistinguishable from R' that is like R but internally uses P

Universal Composition

$\mathcal{G}_{\text{ledger}}$ [BMTZ'17, BGKRZ'18]

Functionality $\mathcal{F}_{\text{PayNet}}$ – interface

- from \mathcal{E} :
 - (REGISTER, delay, relayDelay)
 - (TOPPEDUP)
 - (OPENCHANNEL, *Alice*, *Bob*, *x*, *tid*)
 - (CHECKFORNEW, *Alice*, *Bob*, *tid*)
 - (PAY, *Bob*, *x*, $\overrightarrow{\text{path}}$, receipt)
 - (CLOSECHANNEL, receipt, *pchid*)
 - (FORCECLOSECHANNEL, receipt, *pchid*)
 - (POLL)
 - (PUSHFULFILL, *pchid*)
 - (PUSHADD, *pchid*)
 - (COMMIT, *pchid*)
 - (FULFILLONCHAIN)
 - (GETNEWS)
- to \mathcal{E} :
 - (REGISTER, *Alice*, delay(*Alice*), relayDelay(*Alice*), pubKey)
 - (REGISTERED)
 - (NEWS, newChannels, closedChannels, updatesToReport)
- from \mathcal{S} :
 - (REGISTERDONE, *Alice*, pubKey)
 - (CORRUPTED, *Alice*)
 - (CHANNELANNOUNCED, *Alice*, $p_{\text{Alice},F}$, $p_{\text{Bob},F}$, *fchid*, *pchid*, *tid*)
 - (UPDATE, receipt, *Alice*)
 - (CLOSEDCHANNEL, channel, *Alice*)
 - (RESOLVEPAYS, *payid*, charged)
- to \mathcal{S} :
 - (REGISTER, *Alice*, delay, relayDelay)
 - (OPENCHANNEL, *Alice*, *Bob*, *x*, *fchid*, *tid*)
 - (CHANNELOPENED, *Alice*, *fchid*)
 - (PAY, *Alice*, *Bob*, *x*, $\overrightarrow{\text{path}}$, receipt, *payid*)
 - (CONTINUE)
 - (CLOSECHANNEL, *fchid*, *Alice*)
 - (FORCECLOSECHANNEL, *fchid*, *Alice*)
 - (POLL, Σ_{Alice} , *Alice*)
 - (PUSHFULFILL, *pchid*, *Alice*)
 - (PUSHADD, *pchid*, *Alice*)
 - (COMMIT, *pchid*, *Alice*)
 - (FULFILLONCHAIN, *t*, *Alice*)

Our contributions

- Prove Lightning Network security in UC framework
- Derive exact time bounds for how often parties need to check the chain
- Prove that a ledger with instant finality is unrealizable

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Thank you!