

# First Year Review

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**Abstract.** This research proposal concerns the study of economic trust in the decentralised, asynchronous setting and its applications on online commerce and blockchains. In particular, the aim of this project twofold. The first goal is to understand whether it is possible to achieve robust online commerce with minimal risk of fraud without trusted third parties and if so at what cost. The second goal is to provide a general model of off-chain payments and explain how a relaxation of the trustless element of current solutions can facilitate the creation of a scalable blockchain. Tools from the fields of Cryptography and Game Theory will be used to achieve the aforementioned targets.

## 1 Research Topic

The advent of bitcoin [1] has promised to revolutionise the way online commerce takes place, only to meet multiple hurdles in the way. Today there exists an entire ecosystem of cryptocurrencies, e.g. [2,3,4], each with its own benefits and drawbacks. Nevertheless, there are several fundamental questions that still remain unanswered. Some of the prominent issues are whether cryptocurrencies can become the global unit of account, whether it is possible to regulate and fairly distribute them, what are the exact benefits they provide when compared to fiat currencies, in which respects the latter outperform them and if it is possible to completely replace fiat currencies with decentralised cryptocurrencies for which the traditional measures of law enforcement are less applicable due to the global nature of blockchains.

This research proposal is crucially motivated by the following observation: The single most advertised advantage of permissionless blockchains is the lack of need for trust between the interacting individuals *and* for a mediating trusted third party [1]. Unfortunately, this benefit is limited to a certain type of digital assets that admit a cryptographically secure

proof of ownership and can be provably transferred to another party. Additionally, the overhead for setting such a system in motion is too big for the entire economy to be realised in this way.

### 1.1 Decentralised E-Commerce

Going in more detail, blockchains provide us with novel capabilities. It is possible to immutably and truthfully record the fact that Alice paid Bob some cryptographic coins on a blockchain; it is possible for Bob to provably and irrefutably give her a digital asset in return (e.g. an ICO [5], or a sword in a future blockchain-backed game) and record the fact on a blockchain as well; additionally, it is possible to make the exchange atomic so that no trust between the two parties is needed. Unfortunately, all the above are impossible to accomplish if the asset is physical. Alice may pay first and then Bob can refuse to give the asset; or Bob may indeed hand in the asset, but Alice may then claim that it was not as advertised and demand her money back. The possibility of fraud is exacerbated if the exchange takes place online between parties who do not know each other personally. Crucially, an external observer cannot always resolutely decide which of the parties is right. Even if she can, it seems impossible for her to prove that she is not colluding with either player and that she is completely fair and impartial.

For comparison, if the payment is done in fiat currency, then there exists a mediator (usually a bank, often a court) who is charged with deciding what is the fair way to settle the dispute. Additionally, it is in the mediator's direct benefit to strive for the objectively fairest solution. Obviously the cryptographic security of this state of affairs is very low: it is entirely possible that the bank employee or the judge who has the final say is a friend of the fraudster. Nevertheless, it seems that currently the incentives are much better aligned when relying on the traditional legal system. As a result, transacting in fiat currency is, for practical and everyday purchases of physical goods and especially in the online setting, safer and more robust than using cryptocurrencies.

The cryptographic achievements of blockchains are currently realised only in a closed system, where everything can be proven mathematically. Nevertheless, the economy is not such a system. In order for decentralised blockchains to substitute the traditional fiat currencies, a mechanism for decentralised e-commerce with guarantees of minimal trust is needed. We propose to attempt the creation of such a mechanism or give plausible arguments why such a mechanism is unattainable. In case we manage to create such a system, it should ideally be able to constitute the backend of

a decentralised ecommerce platform, similar in usage to existing platforms such as OpenBazaar [6].

In essence, such a mechanism would provide a decentralised reputation system. In current online marketplaces (e.g. ebay.com) there exists a central authority that aggregates ratings and reviews from past interactions between users and serves them to participants of future transactions with the aim of minimising fraud. We aspire to replace this central authority with a distributed protocol or prove that such an endeavour is impossible.

## 1.2 Off-chain payment channels

One additional impediment to the more widespread use of decentralised blockchains is the issue of scalability. Decentralised consensus protocols based on Proof-of-Work [7] or Proof-of-Stake [4] make the assumption that the majority of the hashing power or of the stake respectively is controlled by honest parties, each of which processes every single transaction. To oversimplify, one can think that such protocols function only as long as every member of a crucial portion of their users processes all transactions. A great duplication of effort takes place in order to avoid trusting a particular set of users. It seems intuitively obvious that, no matter how much the parameters are tuned, such systems can only handle a limited amount of transactions per second and thus cannot compete traditional centralised solutions such as VISA in their current state [8].

One solution to this issue is to ensure that most payments happen off-chain. A variety of mechanisms (discussed in more detail in Section 5) enable cryptocurrency owners to temporarily lock some funds in specially crafted channels with one on-chain transaction and subsequently perform any number of transactions between them with these funds without touching the blockchain. One more on-chain transaction unlocks the funds, ensuring each participant receives exactly the funds they owned in the latest state of the channel. Surprisingly, no trust is needed between the parties. The established term for this type of mechanism is “payment channels”.

As of today there exists no definitive systematisation of what payment channels can achieve and at what cost. Each proposal provides different benefits and incurs different costs; the language and formalism of each proposal is largely incompatible with the rest. Most of the proposals lack formal security definitions and proofs. More importantly, it is unclear what are the exact inherent tradeoffs and whether an ideal balance exists. One of the targets of this research is to provide a suitable conceptual

framework that encompasses the whole range of possibilities and limitations of such channels.

Two related issues that have not yet been tackled are the following. Firstly, currently all proposals offer trustless operation as a feature. Nevertheless, it seems plausible that a relaxation of this constraint may indeed allow for more efficient, faster and lighter payment channels. This is a new direction that this research will attempt to explore, ideally incorporating it in the general framework proposed above.

Lastly, current blockchains do not focus on how to facilitate payment channels, instead attempt to obviate their need. An interesting topic would be to provide a treatment on what properties should a blockchain have to aid the creation of lightweight payment channels.

## **2 Toolset and Strategy**

### **2.1 Simulation-based Security**

The two distinct research directions will both leverage tools and methods used in Cryptography and Game Theory. In order to formally model the interactions between marketplace participants, we will employ simulation-based security [9] and more specifically the Universally Composable Security framework [10]. The general strategy of simulation-based security is as follows: First we give some security definitions that formalise the desired properties our system should have. Subsequently we give an ideal-world functionality that satisfies these properties. This functionality acts as a centralised trusted party that executes the desired operation on behalf of the participants. Afterwards we describe a protocol that hopefully realises this functionality. Finally we prove that the proposed protocol indeed realises the functionality by showing that an external observer cannot distinguish an execution of the functionality from an execution of the actual protocol.

To make matters worse (i.e. closer to the real-world challenge the cryptographic protocol faces), the concept of the Adversary is introduced. This party is in control of all messages that the honest participants exchange and can change, delay or completely block them at will. A complete security proof of the fact that a given protocol realises a particular functionality states that for every efficient Adversary (polynomial algorithm) there exists a Simulator (polynomial algorithm) such that the distributions of the execution of the functionality as viewed from the perspective of each participant are statistically close to the respective distributions of the execution of the protocol.

The Universally Composable Security framework is a particular way of setting up and proving the security of a protocol. Achieving simulation of a protocol in the Universal Composition setting ensures that the protocol can be executed concurrently with other protocols that are secure in the Universal Composition setting. Consequently, managing to prove the security of our protocol in this setting ensures that it can coexist with other secure protocols in the same processing unit without suffering diminished security itself or damaging the security of the rest of the aforementioned protocols. Such an achievement would ensure that our protocol is more readily implementable whilst maintaining a high standard of security and would obviate the need of further security proofs for the case when it is executed in an interleaved fashion with other secure protocols.

## 2.2 Game Theoretic approach

Before the advent of blockchains, most cryptographic protocols solved problems where there were clear-cut lines between the aims of honest parties and the machinations of the Adversary. For example, in the case of public key encryption [11], the honest parties should be able to exchange messages encrypted in a fashion that only those who know the private key can decrypt, even though the adversary has the ability to employ any method conceivable (within some computational limits) to read the original contents of the message. Any protocol that claims to realise secure encryption should satisfy (the formal version of) this requirement. It is intuitively obvious that it is in the benefit of the honest parties to follow such a protocol.

On the other hand, blockchains made the state of affairs much more complex. Indeed, bitcoin was a concrete solution to an ill-defined problem; several years passed before a formal definition of the problem blockchains solve was formulated [12]. Nevertheless, this definition does not necessarily coincide with the desires of the participants to the protocol. The definitions demand that a blockchain realises an immutable ledger that is inexorably extended, however the participants' goals span a great variety, such as maximising their share of coins or preventing certain transactions from entering the blockchain. In this case, assuming that there exist only honest parties that execute the protocol exactly as described and an Adversary that employs every conceivable technique in order to break the security definitions is an oversimplification. It ignores slight but crucial deviations from the protocol (e.g. [13]) that may provide advantages not considered in the security definitions. Therefore the assumption of an honest majority can be contested.

When considering decentralised e-commerce, the problem is even more exacerbated. Formulating a simple definition for the characteristics of a healthy marketplace is highly non-trivial. The individual desires of parties should employ a central role in the analysis and not be necessarily restricted to a clear-cut protocol. Assuming that a portion of honest participants exists may be wrong, especially if other participants follow counter-strategies that make honest behaviour unprofitable or if honest players have incentive to profitably deviate from their strategy. Conversely, permitting arbitrary actions to the Adversary could complicate the analysis more than necessary, since it may be reasonable to assume that certain “unreasonable” strategies will never be followed by the Adversary.

This is where rational analysis and mechanism design [14] come into play. Game Theory [15] provides a series of tools and concepts that greatly support the study of multi-agent rational systems. The concept of Nash Equilibrium [16], fundamental to the study of Game Theory, will be used throughout this research. Indeed, our aim will be to describe a mechanism that resembles the real e-commerce setting as closely as possible—while abstracting away irrelevant complications—and find strategies that constitute Nash Equilibria.

Fortunately, we can leverage an arsenal of cryptographic tools to help us rule out some categories of undesired behaviours and thus design a system that achieves the goals of modern e-commerce whilst minimising the trust towards third parties. The guarantees given will not be necessarily cryptographic, but rather rational; that is, participants will refrain from acting badly not because they do not have the ability, but rather because it will be detrimental to their social standing, access to goods and services or other desirable attributes.

### 3 Progress to date

#### 3.1 Modelling for Decentralised E-Commerce

The main aim of this research is to construct a model that closely follows the dynamics of a marketplace. Following the Universal Composability framework, we define  $n$  interactive Turing Machines (ITMs) that represent the players that take part in the marketplace game. These ITMs comprise the set  $\mathcal{P}$ . Furthermore, we define an ITM named Environment and represented by  $\mathcal{E}$ .

This ITM is the first to start functioning. Initially, its responsibility is to allocate a utility function to each player, drawn from a distribution on legal utility functions that is common knowledge [?] to the players. We

will elaborate on the nature of the utility functions later on. Additionally,  $\mathcal{E}$  provides players with an initial “endowment” of assets and money. Contrary to the utility functions, these are not sent only to the players, but additionally to the  $\mathcal{F}_{\text{Assets}}$  and  $\mathcal{F}_{\text{Ledger}}$  functionalities respectively.

After providing the utility functions and the endowments, the main part of the game begins.  $\mathcal{E}$  can send a message to a player  $P \in \mathcal{P}$  that contains a desire  $d$  to be satisfied. To satisfy this desire,  $P$  has to buy it from another player.

### 3.2 Research on Payment Channels

## 4 Plan

## 5 Literature Review

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