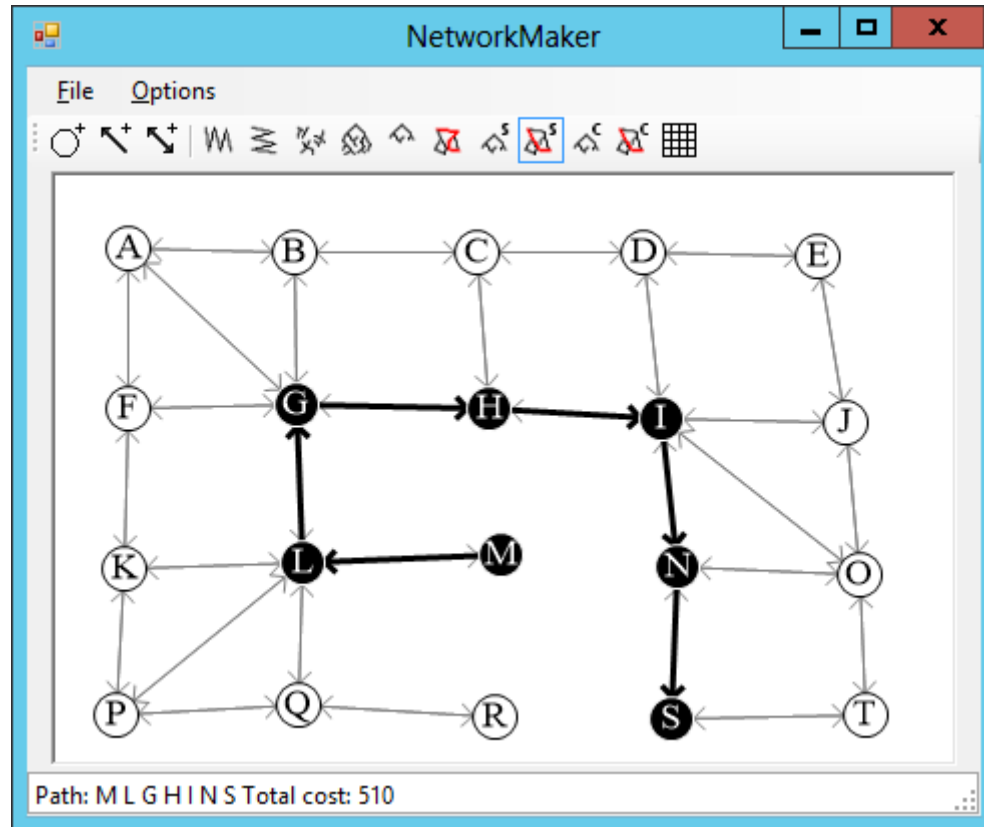


# Path-Finding Algorithms



# Agenda

- Strongly Connected Components
- Finding Paths
  - Finding Any Path
  - Label-Setting Shortest Paths
  - Label-Correcting Shortest Paths
  - All-Pairs Shortest Paths
- Transitivity
- Shortest Path Modifications
- Summary
- Exercises

# Strongly Connected Components

# Finding Paths

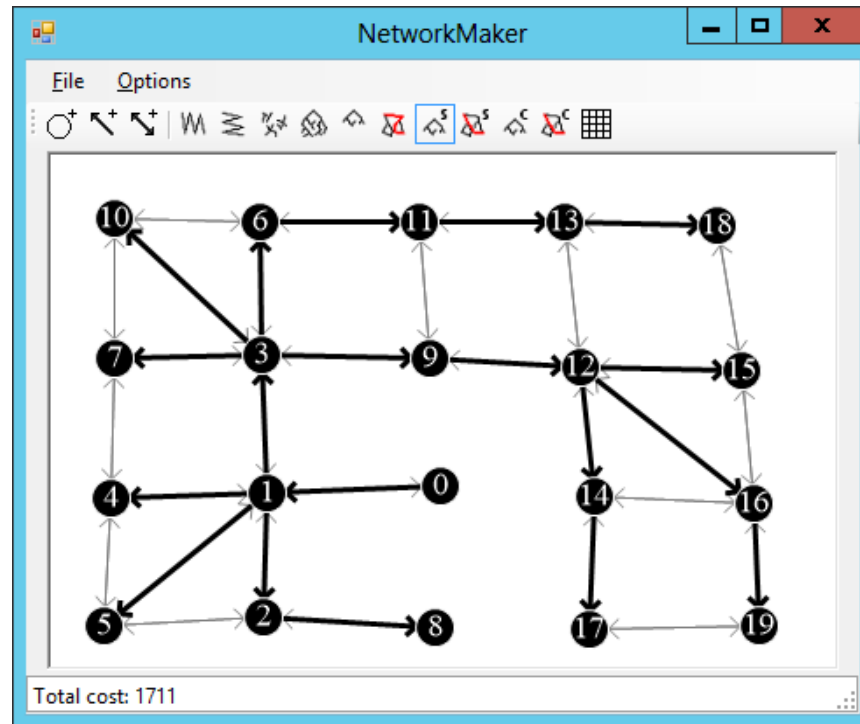
- Finding Any Path
- Label-Setting Shortest Paths
- Label-Correcting Shortest Paths
- All-Pairs Shortest Paths

# Finding Any Path

- To find a path from node A to node B:
  1. Find a spanning tree rooted at node A.
  2. Follow the reversed links in the spanning tree from node B to node A.
  3. Reverse the order in which the links were followed.

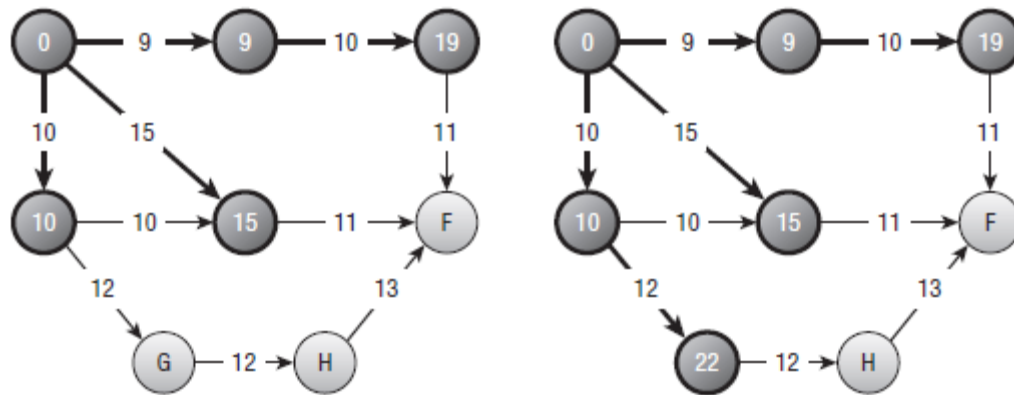
# Shortest Path Trees

- A shortest-path tree gives the shortest paths from the root node to any node in the network



# Label-Setting Shortest Paths

- Once set, a node's distance is never changed
- At each step, you add the link that gives the least total distance from the root to a node that is not in the tree



# Label-Setting Algorithm

1. Set the start node's distance to 0 and mark it as in the tree.
2. Add the start node's links to a *candidate list* of links.
3. While the candidate list is not empty, loop through the list.
  - a. If a link leads to a node in the tree, remove the link from the candidate list.
  - b. Suppose link  $L$  leads  $N_1 \rightarrow N_2$  and  $N_2$  is not yet in the tree. If  $D_1$  is the distance to node  $N_1$  in the tree and  $C_L$  is the cost of the link, let  $D_2 = D_1 + C_L$  be the possible distance for node  $N_2$  that uses this link. As you loop over the links in the candidate list, keep track of the link  $L_{\text{best}}$  and node  $N_{\text{best}}$  that give the smallest possible distance  $D_{\text{best}}$ .
  - c. Set the distance for  $N_{\text{best}}$  to  $D_{\text{best}}$  and mark  $N_{\text{best}}$  as part of the shortest path tree.
  - d. For all links  $L$  leaving node  $N_{\text{best}}$ , if  $L$  leads to a node that is not yet in the tree, add  $L$  to the candidate list.



# Label-Correcting Shortest Paths

- A node's distance may be updated later if a better path is found

# Label-Correcting Algorithm

1. Set the start node's distance to 0 and mark it as in the tree.
2. Add the start node's links to a candidate list of links.
3. While the candidate list is not empty:
  - a. Consider the first link in the candidate list.
  - b. Calculate the distance to the link's destination node:  
 $\text{<distance>} = \text{<source node distance>} + \text{<link cost>}$ .
  - c. If the new distance is better than the destination node's current distance:
    - i. Update the destination node's distance.
    - ii. Add all the destination node's links to the candidate list.

# All-Pairs Shortest Paths, Part 1

- Floyd–Warshall algorithm
- Two-dimensional array named Distance
- `Distance[start_node, end_node]` is the shortest distance between nodes `start_node` and `end_node`

# All-Pairs Shortest Paths, Part 2

- Node via can improve a path if:
  - start\_node  $\rightarrow$  end\_node
- Becomes:
  - start\_node  $\rightarrow$  via\_node  $\rightarrow$  end\_node

# All-Pairs Shortest Paths Algorithm

- Initialize the Distance array
- For  $\text{via\_node} = 0, 1, 2, \dots, N - 1$ :
  - For every pair of nodes  $A \rightarrow B$ :
    - See if you can use  $\text{via\_node}$  to improve the path to:  
 $A \rightarrow \text{via\_node} \rightarrow B$

# All-Pairs Shortest Paths Code, Part 1

1. Initialize the Distance array:
  - a. Set  $\text{Distance}[i, j] = \text{infinity}$  for all entries.
  - b. Set  $\text{Distance}[i, i] = 0$  for all  $i = 1$  to  $N - 1$ .
  - c. If nodes  $i$  and  $j$  are connected by a link  $i \rightarrow j$ , set  $\text{Distance}[i, j]$  to the cost of that link.

# All-Pairs Shortest Paths Code, Part 2

## 2. Initialize the Via array:

- a. For all  $i$  and  $j$ :
  - i. If  $\text{Distance}[i, j] < \text{infinity}$ , set  $\text{Via}[i, j]$  to  $j$  to indicate that the path from  $i$  to  $j$  goes via node  $j$ .
  - ii. Otherwise, set  $\text{Via}[i, j]$  to  $-1$  to indicate that there is no path from node  $i$  to node  $j$ .

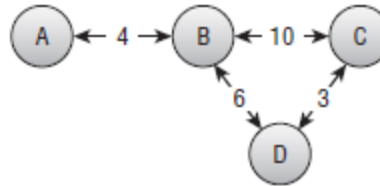
# All-Pairs Shortest Paths Code, Part 3

## 2. Execute the following nested loops to find improvements:

```
For via_node = 0 To N - 1
  For from_node = 0 To N - 1
    For to_node = 0 To N - 1
      Integer: new_dist =
        Distance[from_node, via_node] + Distance[via_node, to_node]
      If (new_dist < Distance[from_node, to_node]) Then
        // This is an improved path. Update it.
        Distance[from_node, to_node] = new_dist
        Via[from_node, to_node] = via_node
      End If
    Next to_node
  Next from_node
Next via_node
```



# All-Pairs Shortest Paths Example



	A	B	C	D
A	0	4	$\infty$	$\infty$
B	4	0	10	6
C	$\infty$	10	0	3
D	$\infty$	6	3	0

	A	B	C	D
A	0	4	$\infty$	$\infty$
B	4	0	10	6
C	$\infty$	10	0	3
D	$\infty$	6	3	0

	A	B	C	D
A	0	4	14	10
B	4	0	10	6
C	14	10	0	3
D	10	6	3	0

	A	B	C	D
A	0	4	14	10
B	4	0	10	6
C	14	10	0	3
D	10	6	3	0

	A	B	C	D
A	0	4	13	10
B	4	0	9	6
C	13	9	0	3
D	10	6	3	0

	A	B	C	D
A	A	B		
B	A	B	C	D
C		B	C	D
D		B	C	D

	A	B	C	D
A	A	B		
B	A	B	C	D
C		B	C	D
D		B	C	D

	A	B	C	D
A	A	B	B	B
B	A	B	C	D
C	B	B	C	D
D	B	B	C	D

	A	B	C	D
A	A	B	B	B
B	A	B	C	D
C	B	B	C	D
D	B	B	C	D

	A	B	C	D
A	A	B	D	B
B	A	B	D	D
C	D	D	C	D
D	B	B	C	D

Improve with node A

Improve with node B

Improve with node C

Improve with node D

# Transitivity

- Transitive Reduction
- Transitive Closure

# Shortest Path Modifications

- Shape Points
- Early Stopping
- Bidirectional Search
- Best-First Search
- Turn Penalties and Prohibitions

# Summary

- Finding Paths
  - Finding Any Path
  - Label-Setting Shortest Paths
  - Label-Correcting Shortest Paths
  - All-Pairs Shortest Paths
- Transitivity
- Shortest Path Modifications

# Exercises

- Chapter 13 Exercises 8 – 21, 23 – 26.
- Read *Essential Algorithms, 2e* Chapter 14 pages 351 – 470. (Stop before the section “Network Cloning.”)