

$$a) T(n) = 36T(n/6) + 2n$$

Master Method

$$a=36, b=6, f(n)=2n$$

$$n^{\log_b a} = n^{\log_6 36} = n^2$$

$$n < n^2$$

\therefore Case 1:

$$T(n) = (n^{\log_b a} - \epsilon) \quad \text{wherein this case } \epsilon=1$$

$$\therefore T(n) = O(n^2)$$

$$b) T(n) = 5T(n/3) + 17n^{1.2}$$

$$a=5, b=3, f(n)=17n^{1.2}$$

$$n^{\log_b a} = n^{\log_3 5} = n^{1.4}$$

$$n^{1.2} < n^{1.4}$$

Case 1:

$$T(n) = (n^{\log_b a} - \epsilon) \quad \text{where in this case } \epsilon=0.2$$

$$T(n) = O(n^{1.4})$$

$$c) T(n) = 12T(n/2) + n^2 \log n$$

$$a=12, b=2, f(n)=n^2 \log n$$

$$n^{\log_b a} = n^{\log_2 12} = n^{3.5}$$

$$n^2 \log n < n^{3.5}$$

Case 1:

$$T(n) = (n^{\log_b a} - \epsilon) \quad \text{where in this case } \epsilon=$$

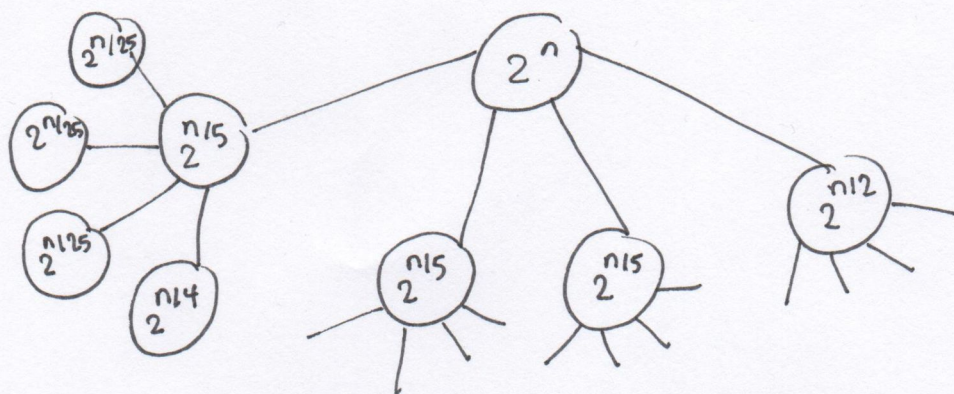
$$T(n) = O(n^{3.5})$$

Problem 2:

Algorithms

Osama Al-Wardi

d) $T(n) = 3T(n/5) + T(n/2) + 2^n$



$$3(2^{n/5}) + 2^{n/2} \quad \dots 1^{st} \text{ Level}$$

$$12(2^{n/125}) + 4(2^{n/4}) \quad \dots 2^{nd} \text{ Level}$$

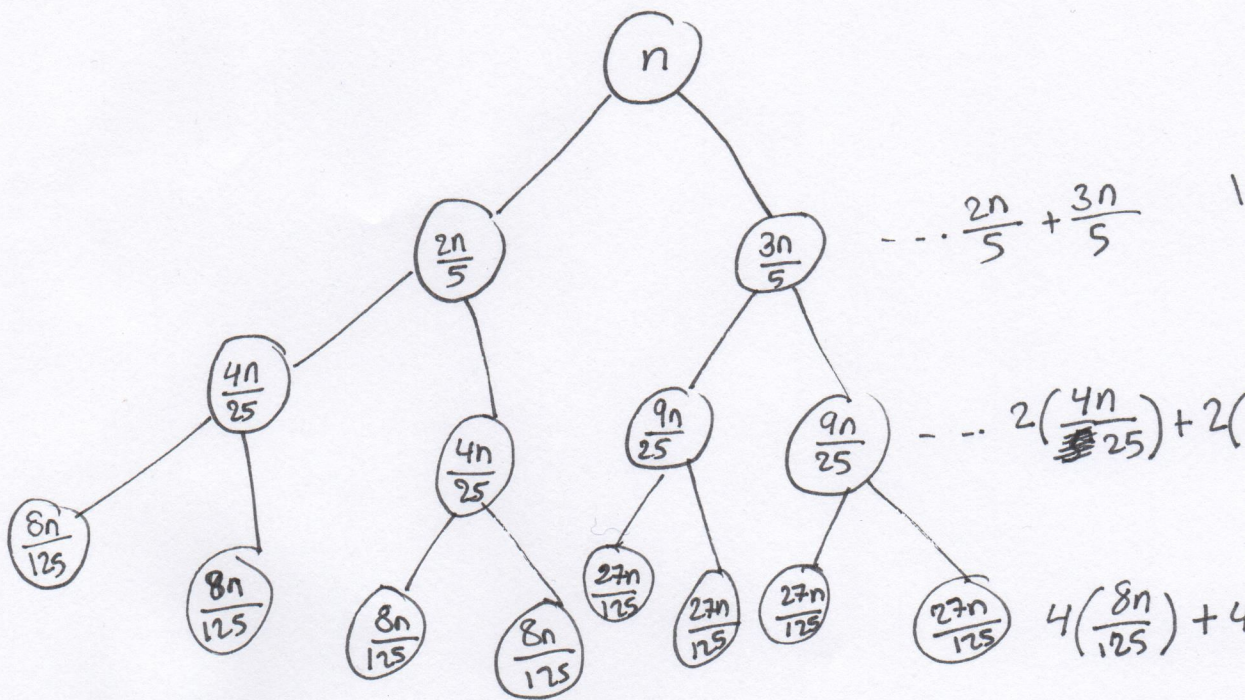
$$24(2^{n/125}) + 16(2^{n/8}) \quad \dots 3^{rd} \text{ Level}$$

Sum of Level 1

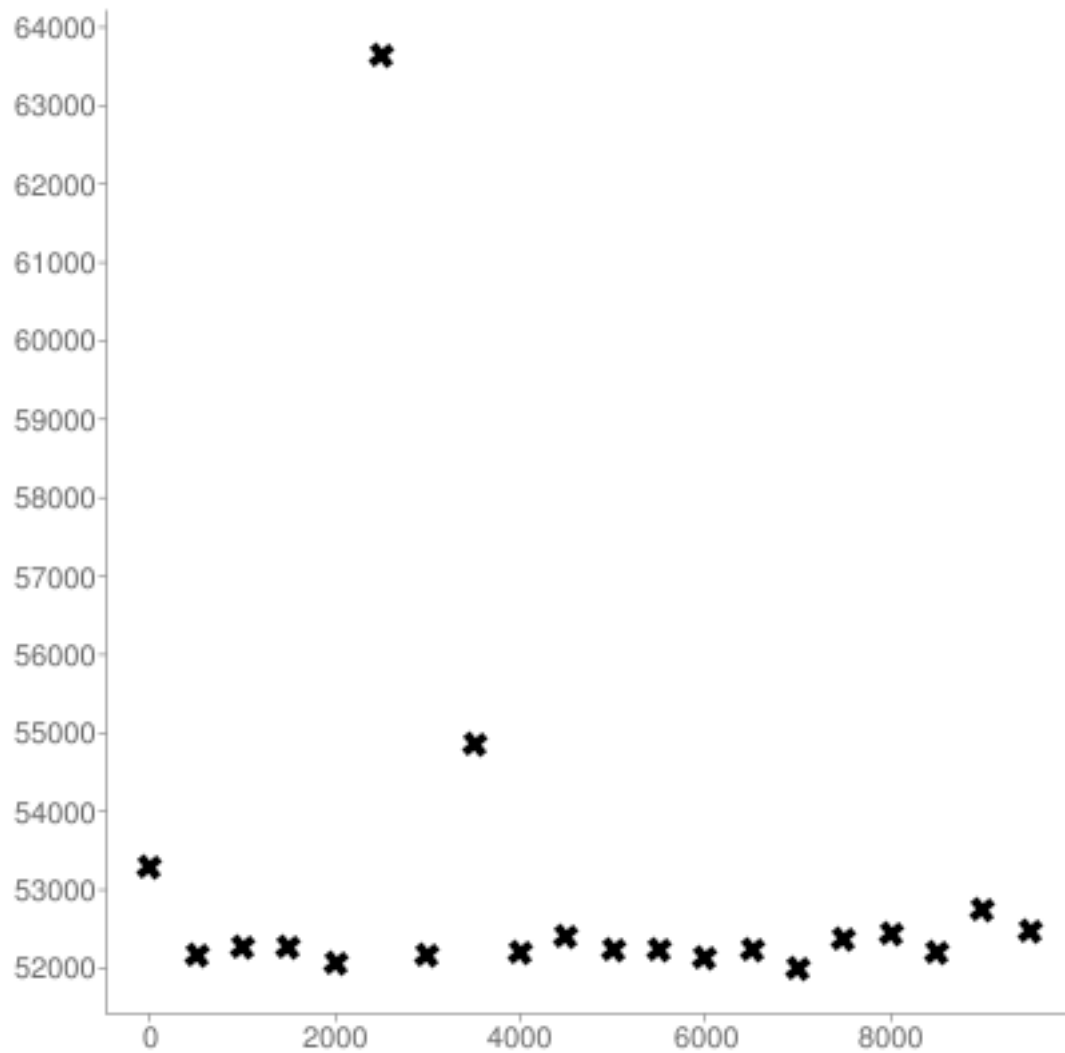
$$\text{Total} = (3(2^{n/5}) + 2^{n/4} + 12(2^{n/125}) + 4(2^{n/4}) + 24(2^{n/125}) + 16(2^{n/8}))$$

$$(2^{n/5}) + (2^{n/4}) \dots (2(2^{n/5}) + 2^{n/4} + 12(2^n$$

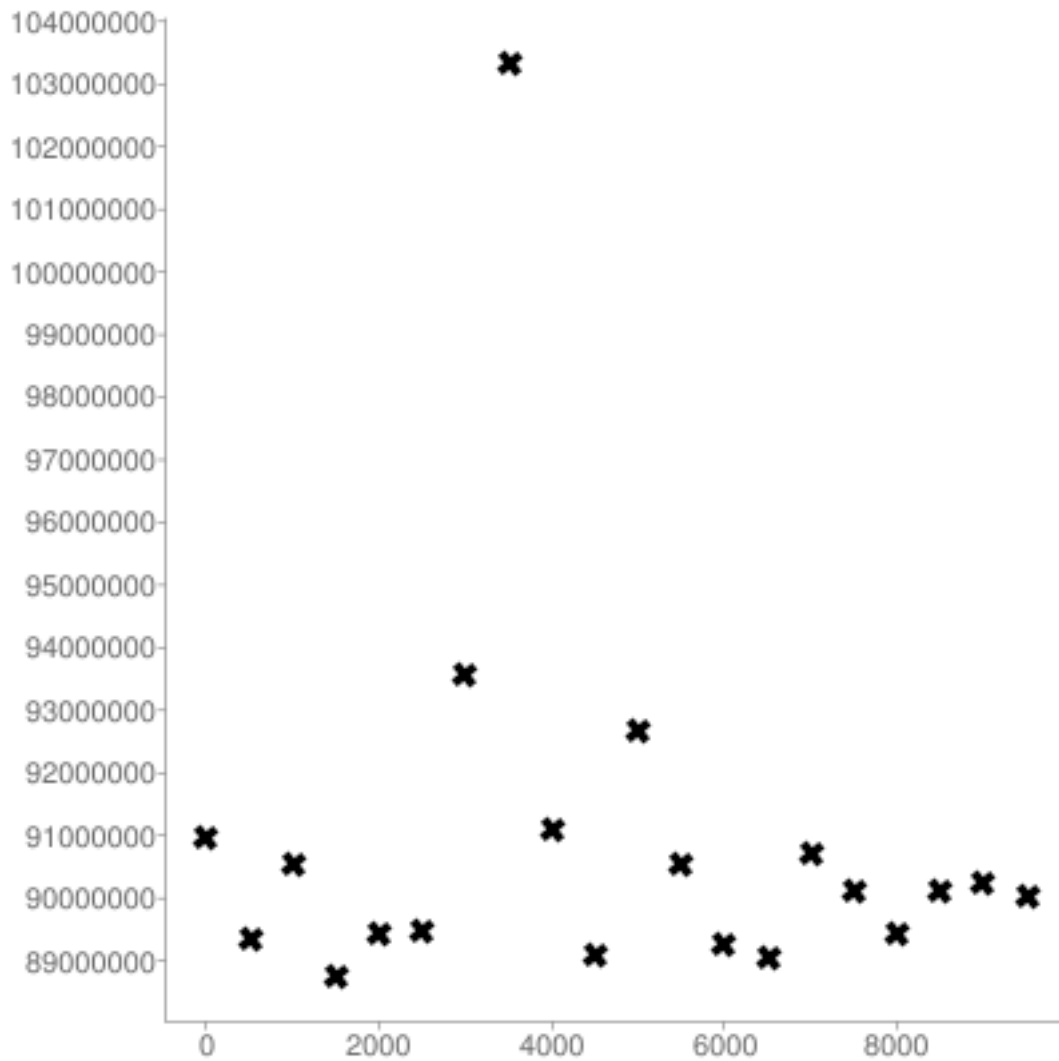
$$T(n) = 1(2n/5) + 1(3n/5) + O(1)$$



Best Case



Average Case



Worst Case

