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1 Template

```

1 // Racso programmed here
2 #include <bits/stdc++.h>
3 using namespace std;
4 typedef long long ll;
5 typedef long double ld;
6 typedef __int128 sll;
7 typedef pair<int,int> ii;
8 typedef pair<ll,ll> pll;
9 typedef vector<int> vi;
10 typedef vector<ll> vll;
11 typedef vector<ii> vii;
12 typedef vector<vi> vvi;
13 typedef vector<vii> vvii;
14 typedef vector<pll> vpll;
15 typedef unsigned int uint;
16 typedef unsigned long long ull;
17 #define fi first
18 #define se second
19 #define pb push_back
20 #define all(v) v.begin(),v.end()
21 #define rall(v) v.rbegin(),v.rend()
22 #define sz(a) (int)(a.size())
23 #define fori(i,a,n) for(int i = a; i < n; i++)
24 #define endl '\n'
25 const int MOD = 1e9+7;
26 const int INFTY = INT_MAX;

```

```

27 const long long LLINF = LLONG_MAX;
28 const double EPS = DBL_EPSILON;
29 void printVector( auto& v ){ fori(i,0,sz(v)) cout <<
    v[i] << " "; cout << endl; }
30 void fastIO() { ios_base::sync_with_stdio(0); cin.
    tie(0); cout.tie(0); }

```

Listing 1: Racso's template.

2 C++ Sintaxis

2.1 Compilation sentences

To compile and execute in C++:

```

g++-13 -Wall -o solucion.exe solucion.cpp
g++ -std=c++20 -Wall -o main a.cpp ./solucion.exe < input.txt > output.txt

```

To compile and execute in Java:

```

javac -Xlint Solucion.java
java Solucion < input.txt > output.txt

```

To execute in Python (two options):

```

python3 solucion.py < input.txt > output.txt
ppyy3 solucion.py < input.txt > output.txt

```

2.2 Custom comparators

```

1 // Using function
2 bool cmpFunction(const pair<int,int> &a, const pair<
    int,int> &b) {
3     return a.second < b.second;

```

```
4 }
5 sort(all(v), cmpFunction);
6 // Using functor
7 struct CmpFunctor {
8     bool operator()(const pair<int,int> &a, const pair
9         <int,int> &b) const {
10         return a.second < b.second;
11     }
12 };
13 sort(all(v), CmpFunctor());
14 // Using lambda function
15 sort(all(v), [](const pair<int,int> &a, const pair<
16     int,int> &b) {
17     return a.second < b.second;
18 });
```

Listing 2: Template for custom sorting, using as example ordering a vii ascending by second element.

2.3 *STL DS Usage

2.4 *Strings methods

2.5 *Pragmas

2.6 Bit Manipulation Cheat Sheet

Bitwise operators:

- & (AND): Sets each bit to 1 if both bits are 1.
- | (OR): Sets each bit to 1 if at least one bit is 1.
- ^ (XOR): Sets each bit to 1 if bits are different.
- ~ (NOT): Inverts all bits.
- << (Left shift): Shifts bits left, fills with 0.

- >> (Right shift): Shifts bits right.
- Get LSB: $n \& -n$

Basic Bit Tasks:

- Get bit: $(n \& (1 \ll i)) \neq 0$
- Set bit: $n \mid (1 \ll i)$
- Clear bit: $n \& \sim(1 \ll i)$
- Toggle bit: $n \hat{=} (1 \ll i)$
- Clear LSB: $n \& (n - 1)$

Set Operations:

- Subset check: $(A \& B) == B$
- Set union: $A \mid B$
- Set intersection: $A \& B$
- Set difference: $A \& \sim B$
- Toggle subset: $A \hat{=} B$

Gray code: $G = B \oplus (B \gg 1)$

2.7 Comparing Floats

```
1 long double a, b, EPS = 1e-9;
2 if( abs(a - b) < EPS ) {
3     // 'a' equals 'b'
4 }
```

Listing 3: Check if two real numbers are equal using an epsilon scope.

2.8 Ceil

$$\left\lceil \frac{a}{b} \right\rceil = \frac{a + b - 1}{b}$$

(Originally I thought $\left\lceil \frac{a}{b} \right\rceil = \frac{a-1}{b} + 1$, but this calculates the wrong way in the cases where $a = 0$)

```
1 int myCeil(long long a, long long b) {
2     return (a + b - 1)/b;
3 }
```

Listing 4: A way to do ceil operation between integers without making explicit conversions. `a` and `b` are `int`, but the operation `a+b-1` can cause an overflow, so they must be casted into `long long` to avoid this. The result must be `int` anyway.

3 Miscellaneous

3.1 Binary Search

```
1 int binary_search( vector<int>& list, int n, int
   target ) {
2     int x0 = 0, x1 = n-1, mid;
3     while( x0 <= x1 ) {
4         mid = (x0 + x1) / 2; // (x1 - x0) / 2 + x0;
5         if( list[mid] == target )
6             return mid;
7         list[mid] < target ? x0 = mid + 1 : x1 = mid
   - 1;
8     }
9     return -1;
10 }
```

Listing 5: Classic (Vanilla) implementation of Binary Search. Returns the index where `target` was found. Binary Search works on $O(\log_2(n))$, let n be the size of the container.

```
1 int binary_search( vector<int>& list, int n, int
   target ) {
2     int ans = 0;
```

```

3     function<bool(int)> check = [&](int idx)->bool {
4         return idx < n && list[idx] <= target;
5     };
6     for( int i = 31; i >= 0; i-- ) {
7         if( check( ans + (1 << i) ) )
8             ans += 1 << i;
9     }
10    return list[ans] == target ? ans : -1;
11 }

```

Listing 6: Logarithmic jumps implementation of Binary Search. Returns the index where `target` was found. Binary Search works on $O(\log_2(n))$, let n be the size of the container.

```

1 int binary_search( vector<int>& list, int n, int
   target ) {
2     int left = 1, right = n + 1, mid;
3     while(right - left >= 1) {
4         mid = left + (right - left) / 2;
5         if( list[mid] >= target && target > list[mid -
6             1]) {
7             //----- TODO logic here -----//
8             break;
9         }
10        else
11            if( list[mid] < target )
12                left = mid + 1;
13        else
14            right = mid;

```

```

14     }
15 }

```

Listing 7: Binary Search implementation for searching an element in the interval $(numbers_{i-1}, numbers_i]$. Originally used on Codeforces problem 474 - B (Worms). Returns the index where `target` was found. Binary Search works on $O(\log_2(n))$, let n be the size of the container.

```

1 function<bool(ll)> check = [&](ll t) -> bool {
2     ll products = 0;
3     for(ll machine : v) {
4         products += t / machine;
5         if( products >= target )
6             return true;
7     }
8     return products >= target;
9 };
10 for(int i = 0; i < 70; i++) {
11     mid = (x0 + x1) / 2;
12     check(mid) ? x1 = mid : x0 = mid + 1;
13 }

```

Listing 8: Implementation of Binary Search in the Answer. Originally used on CSES problem *Factory Machines*. Binary Search works on $O(\log_2(n))$, let n be the size of the container.

3.1.1 *Parallel Binary Search

3.1.2 *Ternary Search

3.2 Kadane's Algorithm

```

1 ll arns = v[0], maxSum = 0;
2 for(i,0,n)
3 {
4     maxSum += v[i];
5     arns = max(arns, maxSum);
6     maxSum = max(0LL,maxSum);
7 }

```

Listing 9: Uses Kadane's Algorithm to find maximum subarray sum in $O(n)$.

4 Queries

4.1 Prefix Sum 2D

```

1 for(int i = 1; i <= n; i++)
2     for(int j = 1; j <= n; j++)
3     {
4         prefix[i][j] = prefix[i][j-1] + prefix[i-1][j] -
5         prefix[i-1][j-1];
6         prefix[i][j] += forest[i-1][j-1] == '*' ? 1 : 0;
7     }
8 for(int i = 0; i < q; i++)
9 {
10     pair<int,int> p1, p2;
11     cin >> p1.fi >> p1.se >> p2.fi >> p2.se;
12     int arns = prefix[p2.fi][p2.se];
13     arns -= prefix[p2.fi][p1.se-1] + prefix[p1.fi-1][
14         p2.se];
15     arns += prefix[p1.fi-1][p1.se-1];

```

```

14     cout << arns << endl;
15 }

```

Listing 10: Construction and query of how many 1's are there in a matrix. Originally used on *Forest Queries* from CSES.

4.2 *Sparse Table

4.3 *Sqrt Decomposition

4.4 *Fenwick Tree

4.5 *Fenwick Tree 2D

4.6 Segment Tree

```

1 typedef long long ll;
2 typedef vector<ll> vll;
3 typedef vector<int> vi;
4 const int INF = INT_MAX;
5 class Segment_tree {
6     public: vll t;
7     Segment_tree( int n = 1e5+10 ) {
8         t.assign(n*4,INF);
9     }
10    void update(int node, int index, int tl, int tr,
11        int val) {
12        if( tr < index || tl > index ) return;
13        if( tr == tl ) t[node] = val;
14        else {
15            int mid = tl + ((tr-tl)>>1);
16            int lft = node << 1;

```



```

16         int right = lft + 1;
17         update(lft, index, tl, mid, val);
18         update(right, index, mid+1, tr, val);
19         t[node] = min(t[lft], t[right]);
20     }
21 }
22 ll query(int node, int l, int r, int tl, int tr)
23 {
24     if( tl > r || tr < l ) return INF;
25     if( tl >= l and tr <= r ) return t[node];
26     else {
27         int mid = tl + ((tr-tl)>>1);
28         int lft = node << 1;
29         int right = lft + 1;
30         ll q1 = query(lft, l, r, tl, mid);
31         ll q2 = query(right, l, r, mid+1, tr);
32         return min(q1, q2);
33     }
34 void build(vi &v, int node, int tl, int tr) {
35     if( tl == tr ) t[node] = v[tl];
36     else {
37         int mid = tl + ((tr-tl)>>1);
38         int lft = node << 1;
39         int right = lft + 1;
40         build(v, lft, tl, mid);
41         build(v, right, mid+1, tr);
42         t[node] = min(t[lft], t[right]);

```

```

43     }
44 }
45 };
46 Segment_tree st(n);
47 st.build(v, 1, 0, n-1);
48 st.update(1, a-1, 0, n-1, b);
49 st.query(1, a-1, b-1, 0, n-1));

```

Listing 11: Segment Tree for Dynamic Range **Minimum** Queries. Racso's Implementation.

```

1 vector<long long> v, sex;
2 int n;
3 void build(int node, int l, int r){
4     if(l == r) sex[node] = v[l];
5     else{
6         int mid = (l+r)/2;
7         build(2*node, l, mid);
8         build(2*node + 1, mid+1, r);
9         sex[node] = sex[2*node] + sex[2*node +1];
10    }
11 }
12 void update(int node, int l, int r, int idx, int val
13 ){
14     if(l == r){
15         v[idx] = val;
16         sex[node] = val;
17     }
18     else{

```

```

18     int mid = (l+r)/2;
19     if(l <= idx && idx <= mid) update(2*node, l,
mid, idx, val);
20     else update(2*node +1, mid+1, r, idx, val);
21     sex[node] = sex[2*node] + sex[2*node + 1];
22 }
23 }
24 int query(int node, int tl, int tr, int l, int r){
25     if(r < tl || tr < l) return 0;
26     if(l <= tl && tr <= r) return sex[node];
27     int tm = (tl+tr)/2;
28     return query(2*node, tl, tm, l, r) + query(2*
node +1, tm+1, tr, l, r);
29 }
30 v.resize(n);
31 sex.resize(4 * n);
32 build(1, 0, n - 1);
33 query(1, 0, n-1, l - 1, r - 1)

```

Listing 12: Segment Tree for Dynamic Range **Sum** Queries. Zum's Implementation.

4.6.1 *2D Segment Tree

4.6.2 *Persistent Segment Tree

4.6.3 *Lazy Propagation

4.7 Ordered Set

```

1 //<-- Header.
2 #include <bits/stdc++.h>

```

```

3 #include <ext/pb_ds/assoc_container.hpp>
4 #include <ext/pb_ds/tree_policy.hpp>
5 using namespace std;
6 using namespace __gnu_pbds;
7 template<typename T, typename Cmp = less<T>>
8 using ordered_set = tree<T,null_type,Cmp,rb_tree_tag
,tree_order_statistics_node_update>;
9 //<-- Declaration.
10 ordered_set<int> oset;
11 //<-- Methods usage.
12 // K-th element in a set (counting from zero).
13 ordered_set<int>::iterator it = oset.find_by_order(k
);
14 // Number of items strictly smaller than k.
15 ordered_set<int>::iterator it = oset.order_of_key(k)
;
16 // Every other std::set method.

```

Listing 13: Ordered set necessary includes in header, declaration of the object, and usage of its new methods.

4.7.1 Multi-Ordered Set

```

1 //<-- Header.
2 #include <bits/stdc++.h>
3 #include <ext/pb_ds/assoc_container.hpp>
4 #include <ext/pb_ds/tree_policy.hpp>
5 using namespace std;
6 using namespace __gnu_pbds;

```

```

7  template<typename T, typename Cmp = less<T>>
8  using ordered_set = tree<T,null_type,Cmp,rb_tree_tag
    ,tree_order_statistics_node_update>;
9  //<-- Use in main.
10 ordered_set<pair<int,int>> multi_oset;
11 map<int,int> cuenta;
12 function<void(int)> insertar = [&](int val) -> void
    {
13     multi_oset.insert({val,++cuenta[val]});
14 };
15 function<void(int)> eliminar = [&](int val) -> void
    {
16     multi_oset.erase({val,cuenta[val]--});
17 };

```

Listing 14: Declaration of multi-oset structure.

4.8 *Treap

4.9 *Trie

5 Graph Theory

5.1 Breadth-First Search (BFS)

```

1  vector<vector<int>> graph;
2  vector<bool> visited;
3  graph.assign(n, vector<int>() ); // <--- main
4  visited.assign(n, false); // <--- main
5
6  void bfs( int s ) {

```

```

7  queue<int> q;
8  q.push( s );
9  visited[ s ] = true;
10 while( ! q.empty() ) {
11     int u = q.front();
12     q.pop();
13     for( auto v : graph[ u ] ) {
14         if( ! visited[ u ] ) {
15             visited[ u ] = true;
16             q.push( v );
17             // --- ToDo logic here ---
18         }
19     }
20 }
21 return;
22 }

```

Listing 15: Iterative implementation of BFS graph traversal over a graph represented as a AdjacencyList on vector of vectors. BFS runs in $O(|V| + |E|)$.

```

1  int n, m;
2  string arns = "";
3  cin >> n >> m;
4  vector<vector<bool>> visited(n,vector<bool>(m,false)
    );
5  vector<string> path(n,string(m,'0'));
6  vector<string> grid(n);
7  vii dirs = {{0,1},{0,-1},{1,0},{-1,0}};
8  string commands = "LRUD";

```

```

9  queue<ii> q;
10 ii start, end, curr;
11 function<bool(int,int)> valid = [&](int i, int j) ->
    bool {
12     return ( i >= 0 && i < n && j >= 0 && j < m &&
        grid[i][j] != '#' && ! visited[i][j] );
13 };
14 fori(i,0,n) cin >> grid[i];
15 fori(i,0,n) {
16     fori(j,0,m)
17         if( grid[i][j] == 'A' ) {
18             visited[i][j] = true;
19             q.push( {i,j} );
20         }
21 }
22 while( ! q.empty() ) {
23     curr = q.front();
24     q.pop();
25     int i = curr.fi;
26     int j = curr.se;
27     if( grid[i][j] == 'B' ) {
28         end.fi = i;
29         end.se = j;
30         break;
31     }
32     int newI, newJ;
33     fori(I,0,4) {
34         newI = i + dirs[I].fi;

```

```

35         newJ = j + dirs[I].se;
36         if( valid(newI,newJ) ) {
37             visited[newI][newJ] = true;
38             q.push( {newI,newJ} );
39             path[newI][newJ] = commands[I];
40         }
41     }
42 }
43 while( path[ end.fi ][ end.se ] != '0' ) {
44     fori(i,0,4) {
45         if( path[ end.fi ][ end.se ] == commands[i] )
46         {
47             arns += i & 1 ? commands[i-1] : commands[i
48             +1];
49             end.fi -= dirs[i].fi;
50             end.se -= dirs[i].se;
51         }
52     }
53 reverse(all(arns));
54 if( arns == "" ) cout << "NO" << endl;
55 else cout << "YES" << endl << arns.size() << endl <<
    arns << endl;

```

Listing 16: BFS on Grid to find shortest path from an starting point *A* to an end *B*. Once the path is found, it reconstruct it with movements *LRUD*. Works in $O(n \cdot m)$. Originally used on problem *Labyrinth* from CSES.

5.2 Deep-First Search (DFS)

```

1 vector<vector<int>> graph;
2 vector<bool> visited;
3 graph.assign(n, vector<int>() ); // <--- main
4 visited.assign(n, false); // <--- main
5
6 void dfs( int s ) {
7     if( visited[s] == true ) return;
8     visited[s] = true;
9     vector<int>::iterator i;
10    for( i = graph[s].begin(); i < graph[s].end();
11        ++i) {
12        if( ! visited[*i] ) {
13            // --- ToDo logic here ---
14            dfs(*i);
15        }
16    }
17 }

```

Listing 17: Recursive implementation of DFS graph traversal over a graph represented as a AdjacencyList on vector of vectors. DFS runs in $O(|V| + |E|)$.

```

1 vector<vector<int>> graph;
2 vector<bool> visited;
3 void dfs( int s ) {
4     stack<int> stk;
5     stk.push(s);
6     while ( !stk.empty() ) {

```

```

7         int u = stk.top();
8         stk.pop();
9         if ( visited[u] ) continue;
10        visited[u] = true;
11        // --- ToDo logic here ---
12        for(auto it = graph[u].rbegin(); it != graph
13            [u].rend(); ++it)
14            if (!visited[*it])
15                stk.push(*it);
16    }
17 }

```

Listing 18: Iterative implementation of DFS graph traversal over a graph represented as a AdjacencyList on vector of vectors. DFS runs in $O(|V| + |E|)$.

```

1 typedef long long ll;
2 vector<vector<ll>> adj;
3 vector<bool> visited;
4 function<void(ll)> dfs = [&](ll u) -> void {
5     if( visited[u] ) return;
6     visited[u] = true;
7     for( ll v : adj[u] )
8         dfs(v);
9 };
10 dfs(n);

```

Listing 19: DFS implementation with a lambda function (adjacency list and visited don't need to be passed thorough argument). DFS runs in $O(|V| + |E|)$.

```

1 typedef long long ll;
2 typedef vector<ll> vll;
3 map<ll,vll> adj;
4 set<ll> visited;
5 function<void(ll)> dfs = [&](ll u) -> void {
6     if( visited.count(u) ) return;
7     visited.insert(u);
8     for( ll v : adj[u] )
9         dfs(v);
10 };
11 dfs(n);

```

Listing 20: DFS implementation with a lambda function implemented with a map instead of vector of vectors, and a set to track visited nodes. DFS runs in $O(|V| + |E|)$.

5.3 Shortest Path

5.3.1 Dijkstra's Algorithm

```

1 typedef long long ll;
2 typedef pair<ll,ll> pll;
3
4 vector<vector<ll>> graph;
5 vector<ll> visited;
6 graph.assign(n, vector<ll>()); // <--- main
7 visited.assign(n, false); // <--- main
8
9 vector<ll> dijkstra( int n, int source, vector<
    vector<pll>> &graph ) {

```

```

10     vector<ll> dist( n, INFTY );
11     priority_queue<pll, vector<pll>, greater<pll>> pq;
12     dist[ source ] = 0;
13     pq.push( {0, source} );
14     while( ! pq.empty() ) {
15         ll d = pq.top().first;
16         ll u = pq.top().second;
17         pq.pop();
18         if( d > dist[ u ] ) continue;
19         for( auto &edge : graph[ u ] ) {
20             ll v = edge.first;
21             ll weight = edge.second;
22             if( dist[ u ] + weight < dist[ v ] ) {
23                 dist[ v ] = dist[ u ] + weight;
24                 pq.push( {dist[ v ], v} );
25             }
26         }
27     }
28     return dist;
29 }

```

Listing 21: Iterative implementation of Dijkstra's Algorithm for shortest path over a graph represented as a AdjacencyList on vector of vectors. Returns a vector with the shortest path to every other vertex in the graph. $O(|E| \times \log_2(|v|))$. doesn't work with negative weights.

```

1 vvp11 graph(n+1,vp11());
2 vector<bool> visited(n+1,false);
3 function<vll(int)> dijkstra = [&](int source) -> vll

```

```

    {
4   vll dist(n+1, INF);
5   priority_queue<pll, vpll, greater<pll>> pq;
6   dist[source] = 0;
7   pq.push({0, source});
8   while( ! pq.empty() ) {
9       ll d = pq.top().fi;
10      ll u = pq.top().se;
11      pq.pop();
12      if( d > dist[u] ) continue;
13      for(pll edge : graph[u]) {
14          ll v = edge.fi;
15          ll w = edge.se;
16          if( dist[u] + w < dist[v] ) {
17              dist[v] = dist[u] + w;
18              pq.push({dist[v], v});
19          }
20      }
21  }
22  return dist;
23 };

```

Listing 22: Iterative implementation of Dijkstra's Algorithm as a Lambda Function for shortest path over a graph represented as a AdjacencyList on vector of vectors. Returns a vector with the shortest path to every other vertex in the graph. $O(|E| \times \log_2(|V|))$. Doesn't work with negative weights.

5.3.2 *Floyd-Warshall's Algorithm

5.3.3 Bellman-Ford Algorithm

```

1  int V, E;
2  cin >> V >> E;
3  vvi edges(E, vi(3, 0));
4
5  for(int i = 1; i <= E; i++)
6      cin >> edges[i][0] >> edges[i][1] >> edges[i][2];
7
8  function<vi(int)> bellman_ford = [&](int src) -> vi
9  {
10     vi dist(V, INF);
11     dist[src] = 0;
12     for(int i = 0; i < V; i++)
13     {
14         for( vi edge : edges )
15         {
16             int u = edge[0];
17             int v = edge[1];
18             int w = edge[2];
19             if( dist[u] != INF and dist[u] + w < dist[v]
20         )
21         {
22             if( i == V-1 )
23                 return = {-1};
24             dist[v] = dist[u] + w;
25         }
26     }
27 }

```

```

25     }
26 }
27 return dist;
28 };

```

Listing 23: Finds the shortest route from a source vertex, to every other one in the graph. Works over a list of edges. Runs in $O(|V| \times |E|)$. Can be used to find negative cycles.

5.4 Minimum Spanning Tree (MST)

5.4.1 *Prim's Algorithm

5.4.2 *Kruskal's Algorithm

5.5 *Bipartite Checking

5.6 *Negative Cycles

5.7 Topological Sort

```

1 vi topo;
2 vvi graph(V+1,vi());
3 vector<bool> visited(V+1,false);
4 function<void(int)> dfs = [&](int u) -> void {
5     visited[u] = true;
6     for(int v : graph[u])
7         if( ! visited[v] )
8             dfs(v);
9     topo.pb(u);
10 };
11 function<void()> topological_sort = [&]() -> void {
12     for(int i = 1; i <= V; i++)

```

```

13         if( ! visited[i] )
14             dfs(i);
15         reverse(all(topo));
16 };
17 topological_sort();
18 for(int i = 0; i < V; i++) cout << topo[i] << " ";

```

Listing 24: Recursive toposort implementation for unweighted DAG through vvi with DFS with inverted postorder. Runs in $O(|V| \times |E|)$.

```

1 vi indegree(V+1,0);
2 vvi graph(V+1,vi());
3 vector<bool> visited(V+1,false);
4 fori(i,0,E) {
5     graph[u].pb(v);
6     indegree[v]++;
7 }
8 function<vi()> topological_sort = [&]() -> vi {
9     vi order, deg = indegree; // copy
10    queue<int> q;
11    for(int i = 1; i <= V; i++)
12        if( deg[i] == 0 )
13            q.push(i);
14    while( ! q.empty() ) {
15        int u = q.front(); q.pop();
16        order.pb(u);
17        for(int v : graph[u]) {
18            deg[v]--;
19            if(deg[v] == 0)

```



```

20         q.push(v);
21     }
22 }
23 return order;
24 };
25 vi topo = topological_sort();
26 if( (int)(topo.size()) != V ) cout << "IMPOSSIBLE"
    << endl;

```

Listing 25: Kahn's Algorithm for Topological Sorting using BFS and indegree vertex analysis (nodes in a cycle will never have indegree zero). Works over unweighted directed graphs containing cycles through vvi. Runs in $O(|V| \times |E|)$.

5.8 *Disjoint Set Union (DSU)

5.9 *Condensation Graph

5.10 *Strongly Connected Components (SCC)

5.11 2-SAT

```

1 struct TwoSatSolver {
2     int n_vars;
3     int n_vertices;
4     vector<vector<int>> adj, adj_t;
5     vector<bool> used;
6     vector<int> order, comp;
7     vector<bool> assignment;
8
9     TwoSatSolver(int _n_vars) : n_vars(_n_vars),
    n_vertices(2 * n_vars), adj(n_vertices), adj_t(

```

```

    n_vertices), used(n_vertices), order(), comp(
    n_vertices, -1), assignment(n_vars) {
10         order.reserve(n_vertices);
11     }
12     void dfs1(int v) {
13         used[v] = true;
14         for (int u : adj[v]) {
15             if (!used[u])
16                 dfs1(u);
17         }
18         order.push_back(v);
19     }
20
21     void dfs2(int v, int c1) {
22         comp[v] = c1;
23         for (int u : adj_t[v]) {
24             if (comp[u] == -1)
25                 dfs2(u, c1);
26         }
27     }
28
29     bool solve_2SAT() {
30         order.clear();
31         used.assign(n_vertices, false);
32         for (int i = 0; i < n_vertices; ++i) {
33             if (!used[i])
34                 dfs1(i);
35         }

```

```

36
37     comp.assign(n_vertices, -1);
38     for (int i = 0, j = 0; i < n_vertices; ++i)
39     {
40         int v = order[n_vertices - i - 1];
41         if (comp[v] == -1)
42             dfs2(v, j++);
43     }
44     assignment.assign(n_vars, false);
45     for (int i = 0; i < n_vertices; i += 2) {
46         if (comp[i] == comp[i + 1])
47             return false;
48         assignment[i / 2] = comp[i] > comp[i +
49     1];
50     }
51     return true;
52
53     void add_disjunction(int a, bool na, int b, bool
54     nb) {
55         // na and nb signify whether a and b are to
56         be negated
57         a = 2 * a ^ na;
58         b = 2 * b ^ nb;
59         int neg_a = a ^ 1;
60         int neg_b = b ^ 1;
61         adj[neg_a].push_back(b);

```

```

60         adj[neg_b].push_back(a);
61         adj_t[b].push_back(neg_a);
62         adj_t[a].push_back(neg_b);
63     }
64
65     static void example_usage() {
66         TwoSatSolver solver(3); // a, b, c
67         solver.add_disjunction(0, false, 1, true);
68         //      a  v  not b
69         solver.add_disjunction(0, true, 1, true);
70         // not a  v  not b
71         solver.add_disjunction(1, false, 2, false);
72         //      b  v      c
73         solver.add_disjunction(0, false, 0, false);
74         //      a  v      a
75         assert(solver.solve_2SAT() == true);
76         auto expected = vector<bool>(True, False,
77         True);
78         assert(solver.assignment == expected);
79     }
80 };

```

Listing 26: 2-SAT implementation from `cp-algorithms.com`. Each component added is a expression of the form $a \vee b$, which is equivalent to $\neg a \Rightarrow b \wedge \neg b \Rightarrow a$ (if one of the variables is false, then the other one must be true). A directed graph is constructed based on these implication: For each x , there's two vertices v_x and $v_{\neg x}$. If there is an edge $a \Rightarrow b$, then there also is an edge $\neg b \Rightarrow \neg a$. For any x , if x is reachable from $\neg x$ and $\neg x$ is reachable from x , the problem has no solution. This means, each variable must be in a different SCC than their negative. This is verified by the method `solve_2SAT()`, which returns a boolean: `True` if it has a solution and `False` if it doesn't.

Giant Pizza How does a particular 2-SAT problem look like? Following is the statement for the problem CSES 1684 (Giant Pizza):

Uolevi's family is going to order a large pizza and eat it together. A total of n family members will join the order, and there are m possible toppings. The pizza may have any number of toppings. Each family member gives two wishes concerning the toppings of the pizza. The wishes are of the form "topping x is good/bad". Your task is to choose the toppings so that at least one wish from everybody becomes true (a good topping is included in the pizza or a bad topping is not included).

Input

The first input line has two integers n and m : the number of family members and toppings. The toppings are numbered $1, 2, \dots, m$. After this, there are n lines describing the wishes. Each line has two wishes of the form " $+x$ " (topping x is good) or " $-x$ " (topping x is bad).

Output

Print a line with m symbols: for each topping "+" if it is included and "-" if it is not included. You can print any valid solution. If there are no valid solutions, print "IMPOSSIBLE".

```

1 int main(){
2     fastIO();
3     int n = nxt(), m = nxt();
4     TwoSatSolver TwoSat(m);
5
6     for(i, 0, n){
7         char type1, type2;
8         int top1, top2;
9         cin >> type1 >> top1 >> type2 >> top2;
10
11         top1--; top2--;
12         TwoSat.add_disjunction(top1, type1 == '-',
top2, type2 == '-');
13     }
14
15     if(TwoSat.solve_2SAT()){
16         for(i, 0, m){
17             if(TwoSat.assignment[i])
18                 cout << "+ ";
19             else
20                 cout << "- ";
21         }
22     }
23     else cout << "IMPOSSIBLE";
24     return 0;
25 }
```

Listing 27: Main method for solving CSES 1684 Giant Pizza using 2-SAT template.

5.12 *Bridges and point articulation

5.13 Flood Fill

```

1 vector<string> grid(n);
2 vii dirs = {{0,1},{0,-1},{1,0},{-1,0}};
3 ii start;
4 int arns = 0;
5 function<void(int,int)> traverse = [&](int i, int j)
    -> void {
6     if( grid[i][j] == '#' ) return;
7     int newI, newJ;
8     if( grid[i][j] != '.' ) arns += grid[i][j] - '0';
9     grid[i][j] = '#';
10    for( ii move : dirs ) {
11        newI = i + move.fi; newJ = j + move.se;
12        if( newI >= 0 && newI < n && newJ >= 0 && newJ <
            m && grid[newI][newJ] == 'T' )
13            return;
14    }
15    for( ii move : dirs ) {
16        newI = i + move.fi; newJ = j + move.se;
17        if( newI >= 0 && newI < n && newJ >= 0 && newJ <
            m )
18            traverse(newI, newJ);
19    }
20 };
21 for(i,0,n)
22     cin >> grid[i];

```

```

23 for(i,0,n) {
24     for(j,0,m) {
25         if( grid[i][j] == 'S' ) {
26             grid[i][j] = '.';
27             start.fi = i;
28             start.se = j;
29         }
30     }
31 }
32 traverse(start.fi, start.se);
33 cout << arns << endl;

```

Listing 28: Traverse a matrix of 'n' x 'm' on grid representation. The matrix is composed of '.' for a valid space (empty), '#' for a wall, 'T' for a trap, and a number for a treasure. This implementation takes the sum of every treasure in the maze. The condition for moving to the next location is that there are no Traps nearby (up, down, left, right), so the player will never be killed while traversing. It also implements a way to read numerous test cases, but without knowing beforehand how many there are. Runs in $O(n \cdot m)$. Originally used on the problem *Treasures* from 2024-2025 ICPC Bolivia Pre-National Contest.

5.14 Lava Flow (Multi-source BFS)

```

1 typedef array<int,3> iii;
2
3 vii dirs = {{1,0},{0,1},{-1,0},{0,-1}};
4 map<int,string> path = {{0,"D"},{1,"R"},{2,"U"},{3,"
    L"}};
5 int n, m;

```

```

6  string arns = "";
7  bool escaped = false;
8  cin >> n >> m;
9  vector<string> grid(n);
10 vvi times(n,vi(m,INF)), prev(n,vi(m,-1));
11 vector<vector<bool>> visited(n,vector<bool>(m,false)
    );
12 queue<iii> q;
13 ii start, end;
14 for(int i = 0; i < n; i++) {
15     cin >> grid[i];
16     for(int j = 0; j < m; j++) {
17         if( grid[i][j] == 'M' ) {
18             q.push({i,j,0});
19             times[i][j] = 0;
20         }
21         else if( grid[i][j] == 'A' )
22             start = {i,j};
23     }
24 }
25 function<bool(int,int)> valid = [&](int I, int J) ->
    bool {
26     return (I >= 0) and (I < n) and (J >= 0) and (J <
        m) and (grid[I][J] != '#' ) and (times[I][J] ==
        INF);
27 };
28 function<bool(int,int)> valid_player = [&](int I,
    int J) -> bool {

```

```

29     return (I >= 0) and (I < n) and (J >= 0) and (J <
        m) and (!visited[I][J]) and (grid[I][J] != '#');
30 };
31 function<bool(int,int)> is_border = [&](int I, int J
    ) -> bool {
32     return I == 0 || I == n-1 || J == 0 || J == m-1;
33 };
34 // Corner cases
35 if( is_border(start.fi,start.se) ) {
36     cout << "YES" << endl << "0" << endl;
37     return 0;
38 }
39 // Multi-Source BFS
40 while( ! q.empty() ) {
41     iii u = q.front();
42     q.pop();
43     for(ii dir : dirs) {
44         int newI = u[0] + dir.fi;
45         int newJ = u[1] + dir.se;
46         int w = u[2] + 1;
47         if( valid(newI,newJ) ) {
48             times[newI][newJ] = w;
49             q.push({newI,newJ,w});
50         }
51     }
52 }
53 // Player BFS
54 q.push({start.fi,start.se,0});

```

```

55 visited[start.fi][start.se] = true;
56 while( ! q.empty() and !escaped ) {
57     iii u = q.front();
58     q.pop();
59     for(int i = 0; i < 4; i++) {
60         int newI = u[0] + dirs[i].fi;
61         int newJ = u[1] + dirs[i].se;
62         int w = u[2] + 1;
63         if( valid_player(newI,newJ) and w < times[newI][
newJ] ) {
64             visited[newI][newJ] = true;
65             prev[newI][newJ] = i;
66             q.push({newI,newJ,w});
67             if( is_border(newI,newJ) ) {
68                 end.fi = newI;
69                 end.se = newJ;
70                 escaped = true;
71                 break;
72             }
73         }
74     }
75 }
76 if( !escaped ) {
77     cout << "NO" << endl;
78     return 0;
79 }
80 // Path reconstruction
81 cout << "YES" << endl;

```

```

82 int i = end.fi;
83 int j = end.se;
84 while( prev[i][j] != -1 ) {
85     int oldI = i;
86     arns += path[ prev[i][j] ];
87     i -= dirs[prev[i][j]].fi;
88     j -= dirs[prev[oldI][j]].se;
89 }
90 reverse(all(arns));
91 cout << sz(arns) << endl << arns << endl;

```

Listing 29: Classic Lava Flow problem implementation, where the timer from the starting point A needs to be less than every other in the MS-BFS starting in M places. Once one edge is reached, the path is reconstructed from the output. Runs in BFS complexity $O(|V| + |E|)$. Originally used in the CSES problem *Monsters*.

5.15 MaxFlow

5.15.1 Dinic's Algorithm

```

1 const ll INF = 1e17;
2 /**
3  * @brief Represents a directed edge in a flow
4  *         network.
5  * @details Stores the edge's source, destination,
6  *         capacity, and current flow.
7  *         Used in max-flow algorithms like Dinic
8  *         or Ford-Fulkerson. */
9 struct flowEdge {
10     int u; // Source node

```

```

8   int v; // Destination node
9   ll cap; // Maximum flow capacity of the edge
10  ll flow = 0; // Current flow through the edge (
    initially 0)
11  flowEdge( int u, int v, ll cap ) : u(u), v(v), cap
    (cap) {};
12 };
13 /**
14  * @brief Implementation of Dinic's max-flow
    algorithm.
15  * @details Manages a flow network with BFS (Level
    Graph) and DFS (Blocking Flow) optimizations. */
16 struct Dinic {
17     vector<flowEdge> edges; // All edges in the flow
    network (including reverse edges)
18     vector<vi> adj;
19     int n; // Total number of nodes in the graph
20     int s; // Source node
21     int t; // Sink node (destination of flow)
22     int id = 0; // Counter for edge indexing
23     vi level; // Stores the level (distance from 's')
    of each node during BFS
24     vi next; // Optimization for DFS: tracks the next
    edge to explore for each node
25     queue<int> q; // Queue for BFS traversal
26     /**
27     * @brief Constructs a Dinic solver for a flow
    network.

```

```

28     * @param n Number of nodes.
29     * @param s Source node.
30     * @param t Sink node. */
31 Dinic( int n, int s, int t ) : n(n), s(s), t(t) {
32     adj.resize(n); // Initialize adjacency list for
    'n' nodes.
33     level.resize(n); // Prepare level array for BFS.
34     next.resize(n); // Prepare next-edge array for
    DFS.
35     fill(all(level),-1); // Mark all levels as
    unvisited (-1).
36     level[s] = 0; // The source has level 0.
37     q.push(s); // Start BFS from the source.
38 }
39 /**
40  * @brief Adds a directed edge and its residual
    reverse edge to the flow network. */
41 void addEdge( int u, int v, ll cap ) {
42     edges.emplace_back(u,v,cap); // Original edge: u
    -> v
43     edges.emplace_back(v,u,0); // Residual edge: v
    -> u
44     adj[u].pb(id++);
45     adj[v].pb(id++);
46 }
47 /**
48     * @brief Performs BFS to construct the level
    graph (Layered Network) from source 's' to sink '

```

```

t'.
49  * @details Assigns levels (minimum distances
    from 's') to all nodes and checks if 't' is
    reachable.
50  *      Levels are used to guide the DFS
    phase in Dinic's algorithm.
51  * @return bool True if the sink 't' is reachable
    (i.e., there exists an augmenting path), false
    otherwise. */
52 bool bfs() {
53     while( ! q.empty() ) {
54         int curr = q.front();
55         q.pop();
56         for( auto e : adj[curr] ) {
57             if( edges[e].cap - edges[e].flow < 1 ) //
    Skip saturated edges (no residual capacity).
58                 continue;
59             if( level[ edges[e].v ] != -1 ) // Skip
    already visited nodes (level assigned).
60                 continue;
61             // Assign level to the neighbor node.
62             level[ edges[e].v ] = level[ edges[e].u ] +
    1; // Next level = current + 1.
63             q.push( edges[e].v ); // Add neighbor to the
    queue for further BFS.
64         }
65     }

```

```

66     return level[t] != -1; // Return whether the
    sink 't' was reached (level[t] != -1).
67 }
68 /**
69  * @brief Finds a blocking flow using DFS in the
    level graph constructed by BFS.
70  * @param u Current node being processed.
71  * @param flow Maximum flow that can be sent from
    'u' to the sink 't'.
72  * @return ll The amount of flow successfully
    sent to 't'. */
73 ll dfs( int u, ll flow ) {
74     if( flow == 0 ) // No remaining flow to send.
75         return 0;
76     if( u == t ) // Reached the sink; return
    accumulated flow.
77         return flow;
78     // Explore edges from 'u' using 'next[u]' to
    avoid revisiting processed edges.
79     for( int& cid = next[u]; cid < sz(adj[u]); cid++
    ) {
80         int e = adj[u][cid]; // Index of the edge in '
    edges'.
81         int v = edges[e].v; // Destination node of
    the edge.
82         // Skip invalid edges:
83         // 1. Not in the level graph (level[u] + 1 !=
    level[v]). Just edges in exactly one leve ahead (

```



```

ensures shortest paths).
84     // 2. No residual capacity (cap - flow < 1).
85     if( level[edges[e].u] + 1 != level[v] || edges
[e].cap - edges[e].flow < 1 )
86         continue;
87     ll f = dfs( v, min(flow, edges[e].cap - edges[
e].flow ) ); // Recursively send flow to 'v'.
88     if( f == 0 ) // No flow could be sent via this
edge.
89         continue;
90     // Update residual capacities:
91     edges[e].flow += f;          // Increase flow in
the original edge.
92     edges[ e ^ 1 ].flow -= f; // Decrease flow in
the reverse edge. (All reverse edges have
distinct parity)
93     return f;                  // Return the flow
sent.
94 }
95 return 0; // No augmenting path found from 'u'.
96 }
97 /**
98  * @brief Computes the maximum flow from source '
s' to sink 't' using Dinic's algorithm.
99  * @details Iterates through BFS and DFS phases
to find the maximum flow.
100     Accumulates flow while there exists
augmentation paths in the residual graph.

```

```

101     Restart auxiliary structures for every new
phase.
102     * @return ll The maximum flow value. */
103 ll maxFlow() {
104     ll flow = 0; // Tracks the total flow sent.
105     while( bfs() ) { // While there are augmenting
paths:
106         fill(all(next),0); // Reset 'next' for DFS.
107         for( ll f = dfs(s,INF); f != 0; f = dfs(s,
INF) ) // Send blocking flow in the level graph:
108             flow += f;
109         // Reset for next BFS phase:
110         fill(all(level),-1);
111         level[s] = 0;
112         q.push(s);
113     }
114     return flow;
115 }
116 /**
117  * @brief Finds edges belonging to the minimum
cut after maxFlow().
118  * @details First, it marks all the reachable
nodes from 's' with an augmentation path after
obtained the max flow
119         and all the saturated edges coming out from
any of the nodes who belong to the min-cut.
120         For 'minCut()' to work, 'maxFlow()' must be
first executed to get the min-cut.

```

```

121     If only is needed the value, is enough
    returning the value of 'maxFlow()'.
122     * @return vii List of edges (u, v) in the min-
    cut. Its size is the minimum number of 'roads' to
    close. */
123 vii minCut() {
124     vii ans;
125     fill(all(level),-1); // Reset levels.
126     level[s] = 0;        // Mark source as reachable
    .
127     q.push(s);
128     while( ! q.empty() ) { // BFS to mark nodes
    reachable from 's' in the residual graph.
129         int curr = q.front();
130         q.pop();
131         for( int id = 0; id < sz(adj[curr]); id++ ) {
    // For every edge going out from 'curr'.
132             int e = adj[curr][id];
133             // If 'v' is has not been visited yet, and
    the edge have residual capacity.
134             if( level[edges[e].v] == -1 && edges[e].cap
    - edges[e].flow > 0 ) {
135                 q.push(edges[e].v);
136                 level[edges[e].v] = level[edges[e].u] + 1;
137             }
138         }
139     }
140     for( int i = 0; i < sz(level); i++ ) {

```

```

141     if( level[i] != -1 ) {
142         for( int id = 0; id < sz(adj[i]); id++ ) {
143             int e = adj[i][id];
144             if( level[edges[e].v] == -1 && edges[e].
    cap - edges[e].flow == 0 )
145                 ans.emplace_back(edges[e].u,edges[e].v);
146         }
147     }
148 }
149 return ans;
150 }
151 /**
152     * @brief Reconstructs the maximum bipartite
    matching after running 'maxFlow()'.
153     * @details Every edge that belong to the
    original graph and have flow greater than zero,
    belongs to the matching.
154     For 'maximumMatching()' to work, 'maxFlow()'
    '.
155     * @return vii List of matched pairs (boy, girl).
    */
156 vii maximumMatching() {
157     vii ans;
158     fill(all(level),-1); // Reset levels.
159     level[s] = 0;        // Mark source as reachable
    .
160     q.push(s);

```

```

161     while( ! q.empty() ) { // BFS to mark nodes
reachable via saturated edges with flow greater
than zero.
162         int curr = q.front();
163         q.pop();
164         for( int id = 0; id < sz(adj[curr]); id++ ) {
165             int e = adj[curr][id];
166             // If 'v' has not been visited yet, the edge
is saturated and have flow greater than zero.
167             if( level[edges[e].v] == -1 && edges[e].cap
- edges[e].flow == 0 && edges[e].flow != 0ll ) {
168                 q.push(edges[e].v);
169                 level[edges[e].v] = level[edges[e].u] + 1;
170             }
171         }
172     }
173     for( int i = 0; i < sz(level); i++ ) { //
Collect original edges (boy -> girl) that are
saturated and have flow > 0.
174         if( level[i] != -1 ) {
175             for( int id = 0; id < sz(adj[i]); id++ ) {
176                 int e = adj[i][id];
177                 if( edges[e].u != s && edges[e].v != t
&& edges[e].cap - edges[e].flow == 0 && edges[e].
flow != 0ll )
178                     ans.emplace_back(edges[e].u, edges[e].v
);
179             }

```

```

180     }
181 }
182 return ans;
183 }
184 };

```

Listing 30: Commented template for solving MaxFlow problems with Dinic's algorithm. Works in complexity $O(|V|^2 \times |E|)$. In bipartite graphs and graphs with unitary max capacity the complexity turns $O(|E| \times \sqrt{|V|})$.

Download Speed How does a particular flow problem look like? Following is the statement for the problem CSES 1694 (Download Speed):

Consider a network consisting of n computers and m connections. Each connection specifies how fast a computer can send data to another computer.

Kotivalo wants to download some data from a server. What is the maximum speed he can do this, using the connections in the network?

Input

The first input line has two integers n and m : the number of computers and connections. The computers are numbered $1, 2, \dots, n$. Computer 1 is the server and computer n is Kotivalo's computer.

After this, there are m lines describing the connections. Each line has three integers a, b , and c : computer a can send data to computer b at speed c .

Output

Print one integer: the maximum speed Kotivalo can download data.

```

1 int main()
2 {
3     fastIO();
4
5     int n, m, u, v, w;

```

```

6
7     cin >> n >> m;
8
9     Dinic flow(n+1,1,n); // size n+1 to fix 0-
indexed indexes, 1 is the source (server), 'n' is
the sink (Kotivalo)
10
11     fori(i,0,m)
12     {
13         cin >> u >> v >> w;
14         flow.addEdge(u,v,w);
15     }
16
17     cout << flow.maxFlow() << endl;
18
19     return 0;
20 }

```

Listing 31: Main method for solving CSES 1697 Download Speed using MaxFlow template.

Police Chase Max Flow-Min Cut Theorem: $\text{MaxFlow} = \text{MinCut}$.

Following is the statement for the problem CSES 1695 (Police Chase):

Kaaleppi has just robbed a bank and is now heading to the harbor. However, the police wants to stop him by closing some streets of the city.

What is the minimum number of streets that should be closed so that there is no route between the bank and the harbor?

Input

The first input line has two integers n and m : the number of crossings and streets. The crossings are numbered $1, 2, \dots, n$. The bank is located at crossing 1, and the harbor is located at crossing n .

After this, there are m lines that describing the streets. Each line has two integers a and b : there is a street between crossings a and b . All streets are two-way streets, and there is at most one street between two crossings.

Output

First print an integer k : the minimum number of streets that should be closed. After this, print k lines describing the streets. You can print any valid solution.

```

1 int main()
2 {
3     fastIO();
4
5     int n, m, u, v;
6     vii minCut;
7
8     cin >> n >> m;
9
10    Dinic flow(n+1,1,n); // size n+1 to fix 0-
indexed indexes, 1 is the source (bank), 'n' is
the sink (harbor)
11
12    fori(i,0,m)
13    {
14        cin >> u >> v;
15        flow.addEdge(u,v,1);
16        flow.addEdge(v,u,1);
17    }
18

```

```

19     flow.maxFlow();
20     minCut = flow.minCut();
21
22     cout << (sz(minCut)/2) << endl;
23     for(int i = 0; i < sz(minCut); i += 2)
24         cout << minCut[i].fi << " " << minCut[i].se
25         << endl;
26     return 0;
27 }

```

Listing 32: Main method for solving CSES 1695 Police Chase using MaxFlow template.

School Dance MaxFlow = MinCut = MaxMatching.

Following is the statement for the problem CSES 1696 (School Dance):

There are n boys and m girls in a school. Next week a school dance will be organized. A dance pair consists of a boy and a girl, and there are k potential pairs.

Your task is to find out the maximum number of dance pairs and show how this number can be achieved.

Input

The first input line has three integers n , m and k : the number of boys, girls, and potential pairs. The boys are numbered $1, 2, \dots, n$, and the girls are numbered $1, 2, \dots, m$.

After this, there are k lines describing the potential pairs. Each line has two integers a and b : boy a and girl b are willing to dance together.

Output

First print one integer r : the maximum number of dance pairs. After this, print r lines describing the pairs. You can print any valid solution.

```

1  int main()
2  {
3      fastIO();
4
5      int n, m, k, a, b;
6      ll maxPairs;
7      vii pairs;
8
9      cin >> n >> m >> k;
10
11     Dinic flow(n+m+2,0,n+m+1);
12
13     fori(boy,0,n+1)
14         flow.addEdge(0,boy,1);
15
16     fori(girl,n+1,n+m+1)
17         flow.addEdge(girl,n+m+1,1);
18
19     fori(i,0,k)
20     {
21         cin >> a >> b;
22         flow.addEdge(a,n+b,1);
23     }
24
25     maxPairs = flow.maxFlow();
26     pairs = flow.maximumMatching();
27

```

```

28     cout << maxPairs << endl;
29     fori(i,0,sz(pairs))
30         cout << pairs[i].fi << " " << (pairs[i].se -
        n) << endl;;
31
32     return 0;
33 }

```

Listing 33: Main method for solving CSES 1696 School Dancing using MaxFlow template.

5.15.2 *Ford-Fulkerson Algorithm

5.15.3 *Goldber-Tarjan Algorithm

6 Trees

6.1 Counting Childrens

```

1 vi childrens(n+1,0);
2 vvi graph(n+1);
3 vector<bool> visited(n+1, false);
4 fori(i,2,n+1) {
5     cin >> tmp;
6     graph[tmp].pb(i);
7     graph[i].pb(tmp);
8 }
9 function<int(int)> dfs = [&](int u) -> int {
10     visited[u] = true;
11     for(int v : graph[u]) {
12         if( !visited[v] )

```

```

13         childrens[u] += dfs(v);
14     }
15     return childrens[u] + 1;
16 };
17 dfs(1);

```

Listing 34: Algorithm that counts how many childrens does every node have, from 2..n in a rooted tree (root = 1).

6.2 *Tree Diameter

6.3 *Centroid Decomposition

6.4 *Euler Tour

6.5 *Lowest Common Ancestor (LCA)

6.6 *Heavy-Light Decomposition (HLD)

7 Strings

7.1 Knuth-Morris-Pratt Algorithm (KMP)

```

1 // Longest Prefix-Suffix
2 vi compute_LPS(string s) {
3     size_t len = 0, i = 1, sz = s.size();
4     vi lps(sz,0);
5     while( i < sz ) {
6         if( s[i] == s[len] )
7             lps[i++] = ++len;
8         else
9             if( len != 0 )
10                 len = lps[len-1];

```

```

11     else
12         lps[i++] = 0;
13     }
14     return lps;
15 }
16 // Get number of occurrences of a pattern p in a
    string s
17 int kmp(string s, string p) {
18     vi lps = compute_LPS(p);
19     size_t n = s.size(), m = p.size(), i = 0, j = 0;
20     int cnt = 0;
21     while( i < n ) {
22         if( p[j] == s[i] ) {
23             j++; i++;
24         }
25         if( j == m ) { // Full match
26             cnt++;
27             j = lps[ j - 1 ];
28         }
29         else if( i < n and p[j] != s[i] ) { // Mismatch
    after j matches
30             if( j != 0 )
31                 j = lps[ j - 1 ];
32             else
33                 i++;
34         }
35     }
36     return cnt;

```

```

37 }

```

Listing 35: KMP algorithm for counting how many times a pattern appear into a string. Runs in $O(n + m)$.

7.2 *Suffix Array

7.3 *Rolling Hashing

7.4 *Z Function

7.5 *Aho-Corasick Algorithm

8 Dynamic Programming

8.1 *Coins

8.2 *Longest Increasing Subsequence (LIS)

8.3 *Edit Distance

8.4 *Knapsack

8.5 *SOS DP

8.6 *Digit DP

8.7 *Bitmask DP

9 Mathematics

9.1 Number Theory

9.1.1 Greatest Common Divisor (GCD)

```

1 int gcd(int a, int b) {
2     if (a == 0) return b;
3     if (b == 0) return a;
4     if (a == b) return a;

```

```

5  if (a > b)
6      return gcd(a - b, b);
7  return gcd(a, b - a);
8  }

```

Listing 36: Implementation of handmade GCD, because using `gcd()` runs slow with long long, also `_gcd()`.

9.1.2 Gauss Sum

The sum of the first n natural numbers in $O(1)$.

$$S = \frac{n(n+1)}{2}$$

$$n = \sqrt{2S + \frac{1}{4}} - \frac{1}{2}$$

```

1  int S = (1LL * n * (1LL * n + 1LL))/2;
2  int n = (int)( sqrt( 2 * S + 0.25 ) - 0.5 )

```

Listing 37: Implementation of the Gauss Sum.

9.1.3 *Modular Theory

9.1.4 *Modulo Inverse

9.1.5 *Fermat's Little Theorem

9.1.6 *Chinese Remainder Theorem

9.1.7 Bnpow

```

1  const int MOD = 1e9+7;
2  int bnpow( long long a, long long b ) { // a^b
3      long long sol = 1;

```

```

4      a %= MOD;
5      while( b > 0 ) {
6          if( b & 1 )
7              sol = ( 1LL * sol * a ) % MOD;
8          a = ( 1LL * a * a ) % MOD;
9          b >>= 1;
10     }
11     return sol % MOD;
12 }

```

Listing 38: Applying binary exponentiation to a problem requiring $a^b \bmod (10^9 + 7)$ in $O(\log_2(b))$.

(2) 9.1.8 Matrix Exponentiation (Linear Recurrency)

```

1  template <typename T> void matmul(vector<vector<T>>
    &a, const vector<vector<T>>& b) {
2      size_t n = a.size(), m = a[0].size(), p = b[0].
        size();
3      assert(m == b.size());
4      vector<vector<T>> c(n, vector<T>(p));
5      for(size_t i = 0; i < n; i++)
6          for(size_t j = 0; j < p; j++)
7              for(size_t k = 0; k < m; k++)
8                  c[i][j] = (c[i][j] + a[i][k] * b[k][j])
9                      % MOD;
10     a = c;
11 }
12 template <typename T> struct Matrix {

```



```

12  vector<vector<T>> mat;
13  Matrix() {}
14  Matrix(vector<vector<T>> a) { mat = a; }
15  Matrix(int n, int m) {
16      mat.resize(n);
17      for(int i = 0; i < n; i++) {mat[i].resize(m);
18  }
19  int rows() const { return mat.size(); }
20  int cols() const { return mat[0].size(); }
21  void makeIden() {
22      for(int i = 0; i < rows(); i++)
23          for(int j = 0; j < cols(); j++)
24              mat[i][j] = (i == j ? 1 : 0);
25  }
26  Matrix operator*=(const Matrix &b) {
27      matmul(mat, b.mat);
28      return *this;
29  }
30  void print() {
31      for(int i = 0; i < rows(); i++) {
32          for(int j = 0; j < cols(); j++)
33              cout << mat[i][j] << " ";
34          cout << endl;
35      }
36  }
37  Matrix operator*(const Matrix &b) { return Matrix
    (*this) *= b; }

```

```

38  };
39  int main() {
40      Matrix<ll> A( {{1,1},{1,0}} );
41      Matrix<ll> ini(2,1);
42      ini.mat[0][0] = 0;
43      ini.mat[1][0] = 1;
44      Matrix<ll> iden(2,2);
45      iden.makeIden();
46      ll n;
47      cin >> n;
48      while(n > 0) {
49          if( n & 1 ) iden *= A;
50          A *= A;
51          n >>= 1;
52      }
53      Matrix<ll> res = iden * ini;
54      cout << res.mat[0][0] << endl;
55      return 0;
56  }

```

Listing 39: Template to pow a matrix of size n to a certain exponent with logarithmic time (using binpow), and multiply it to another matrix, with modulo operation, as well as how to use it. Full implementation for calculating n -th Fibonacci term with linear recurrency.

9.1.9 Prime checking

```

1  bool prime( int n ){
2      if( n == 2 )

```

```

3     return true;
4     if( n % 2 == 0 || n <= 1 )
5         return false;
6     for( int i = 3; i * i <= n; i += 2 )
7         if( ( n % i ) == 0 )
8             return false;
9     return true;
10 }

```

Listing 40: Returns if n is a prime number in $O(\sqrt{n})$. Avoids overflow $\forall n \leq 10^6$ ($\approx INT_MAX$).

9.1.10 Prime factorization

```

1 void prime_factorization(vll& factorization, ll n) {
2     for(long long d = 2; d*d <= n; d++) {
3         while(n % d == 0) {
4             factorization.push_back(d);
5             n /= d;
6         }
7     }
8     if( n > 1 )
9         factorization.push_back(n);
10 }

```

Listing 41: Returns prime factorization of the number n using *trial division*, simplest way. Runs in $O(\sqrt{n})$. e.g. for 12 the result is 2x2x3.

9.1.11 Sieve of Eratosthenes

```

1 void sieve_of_eratosthenes(vector<bool>& is_prime,
2     int n) {
3     is_prime.assign(n+1,true);
4     is_prime[0] = is_prime[1] = false;
5     for(int i = 2; i <= n; i++) {
6         if( is_prime[i] && (long long)i * i <= n ) {
7             for(int j = i*i; j <= n; j += i)
8                 is_prime[j] = false;
9         }
10    }

```

Listing 42: Calculates every prime number up to n with sieve of eratosthenes in a boolean 1-indexed vector. Runs in $O(n \log \log n)$.

9.1.12 Sum of Divisors

```

1 ll sum_of_divisors(ll n) {
2     ll sum = 1;
3     for (long long i = 2; i * i <= n; i++) {
4         if(n % i == 0) {
5             int e = 0;
6             do {
7                 e++;
8                 n /= i;
9             } while (n % i == 0);
10            ll s = 0, pow = 1;
11            do {

```

```

12     s += pow;
13     pow *= i;
14 } while (e-- > 0);
15     sum *= s;
16 }
17 }
18 if(n > 1)
19     sum *= (1 + n);
20 return sum;
21 }

```

Listing 43: Calculates the sum of all divisors of number n . e.g. $sum_of_divisors(12) = 18$. Runs in $O(\sqrt{n})$.

```

1 void sum_of_divisors_sieve( vll& sigma, int n ) {
2     sigma.assign(n+1,0);
3     for(int i = 1; i <= n; i++)
4         for(int j = i; j <= n; j+=i)
5             sigma[j] += i;
6 }

```

Listing 44: Calculates the sum of all divisors of all numbers from 1 to n . Runs in $O(n \log(n))$.

9.2 Combinatorics

9.2.1 Binomial Coefficients

```

1 const int MAXN = 1e6+1;
2 vll fact(MAXN+1), inv(MAXN+1);

```

```

3 int binpow( ll a, ll b ) { // a^b
4     ll sol = 1;
5     a %= MOD;
6     while( b > 0 ) {
7         if( b & 1 )
8             sol = ( 1LL * sol * a ) % MOD;
9         a = ( 1LL * a * a ) % MOD;
10        b >>= 1;
11    }
12    return sol % MOD;
13 }
14 void combi() {
15     fact[0] = inv[0] = 1;
16     fori(i,1,MAXN+1) {
17         fact[i] = fact[i-1] * i % MOD;
18         inv[i] = binpow( fact[i], MOD - 2 );
19     }
20 }
21 ll nCr( ll n, ll r ) {
22     return fact[n] * inv[r] % MOD * inv[n-r] % MOD;
23 }
24 combi();
25 nCr(a,b);

```

Listing 45: Template for calculating binomial coefficients $\binom{n}{k} = \frac{n!}{k!(n-k)!}$. Precalculate $fact$ and inv runs in $O(MAXN \cdot \log_2(MOD))$ ($\log_2(MOD) \approx 30$). So, in general case when $NMAX = 10^6$ and $MOD = 10^9 + 7$ can be generalized to $O(n \cdot \log(n))$, $n \leq 10^6$.

9.2.2 Common combinatorics formulas

$$\binom{n}{2} = \frac{n(n-1)}{2}$$

$$\sum_{k=0}^n \binom{n}{k} = 2^n$$

$$\sum_{k=0}^n \binom{n}{k} \binom{n}{n-k} = \binom{2n}{n}$$

$$\sum_{k=0}^n k \binom{n}{k} = n2^{n-1}$$

$$\sum_{k=0}^{\infty} \binom{2k}{k} \binom{2n-2k}{n-k} = 4^n$$

9.3 *Stars and Bars

9.4 Probability

9.5 Computational Geometry

9.5.1 *Cross Product

9.5.2 *Convex Hull

9.6 *Fast Fourier Transform (FFT)

10 Appendix

10.1 What to do against WA?

1. Have you done the correct complexity analysis?
2. Have you understood well the statement?
3. Have you corroborated yet the trivial test cases?
4. Have you checked all the corner cases?
5. Have you proposed a lot of non-trivial test cases?

6. Isn't there any possibility of overflow? (Multiplying two `int` needs to be fitted into a `long long`)

(3) 7. Have you done a desktop test?

(4) 8. Have you read all the variables? (`tc` variable on `main`)

(5) 9. Every part of your code works as it's meant to?

10.2 Primitive sizes

(6)	Data type	[B]	Minimum value it takes	Maximum value it takes
	bool	1	0	1
	signed char	1	0	255
(7)	unsigned char	1	-128	127
	signed int	4	-2,147,483,648 $\approx -2 \times 10^9$	2,147,483,647 $\approx 2 \times 10^9$
(8)	unsigned int	4	0	4,294,967,295 $\approx 4 \times 10^9$
	signed short	2	-32,768	32,767
	unsigned short	2	0	65,535
	signed long long int	8	-9,223,372,036,854,775,808 $\approx -9 \times 10^{18}$	9,223,372,036,854,775,807 $\approx 9 \times 10^{18}$
	unsigned long long int	8	0	18,446,744,073,709,551,615 $\approx 18 \times 10^{18}$
	float	4	1.1×10^{-38}	3.4×10^{38}
	double	8	2.2×10^{-308}	1.7×10^{308}
	long double	12	3.3×10^{-4932}	1.1×10^{4932}

Table 1: Capacity of primitive data types in C++.

10.3 *ASCII table

10.4 *Numbers bit representation

10.5 How a `vector<vector<pair<int,int>>>` looks like

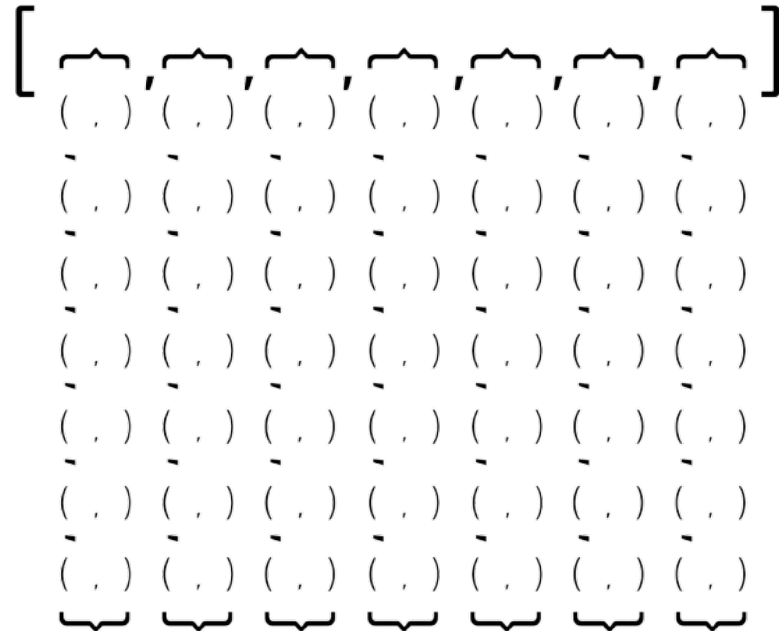


Figure 1: Visual representation of a vector of vector of pairs.

10.6 How all neighbours of a grid looks like

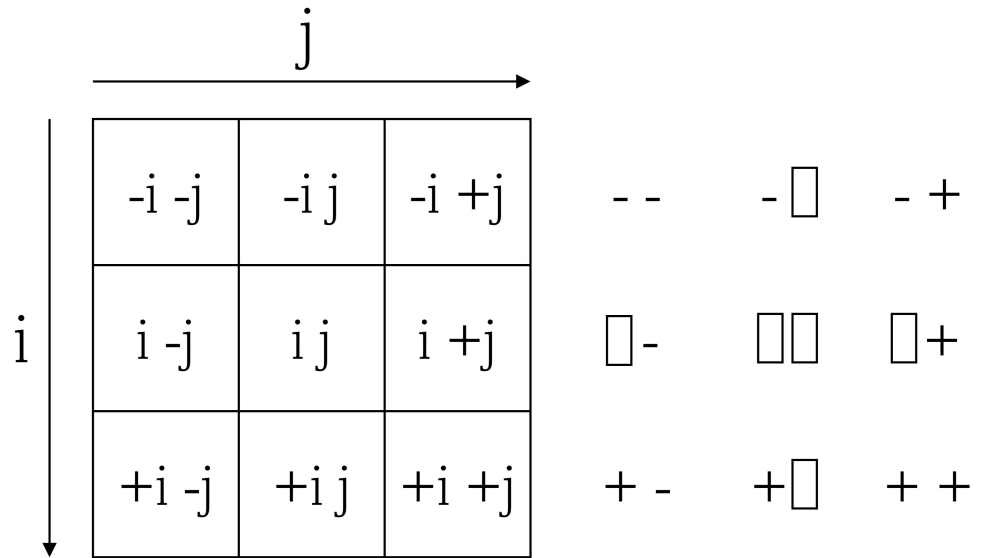


Figure 2: Visual representation of how all adjacent cells in a grid look like.