資料結構和演算法 Data Structure and Algorithm

TZU-CHUN HSU^1

 $^{1}\mathrm{vm3y3rmp40719@gmail.com}$

¹Department of Computer Science, Zhejiang University



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Disclaimer

本文「資料結構與演算法」為「資料結構」和「演算法」筆記的總整理,內容主要參考 Introduction to Algorithms[2] 和洪捷先生的演算法參考書 [1],以及 wjungle 網友在 PTT 論壇上提供的資料結構筆記 [3][4]。

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1 Summary

Trees				
Tree	Insert x	Delete x	Search x	Remark
BST	$O(\log n) \sim O(n)$		(n)	Create: $O(n \log n) \sim O(n^2)$
AVL tree				$F_{n+2} - 1 \le n \le 2^h - 1$
B tree	$O(\log_m n)$			$1 + 2\frac{\lceil \frac{m}{2} \rceil^{h-1} - 1}{\lceil \frac{m}{2} \rceil - 1} \le n \le 2\lceil \frac{m}{2} \rceil^{h-1} - 1$
RBT				$h \le 2\log(n+1)$
Splay tree				Worst: $O(n)$

Disjoint set				
Combination	Union	Find		
Arbitrary Union &	O(1)	O(h)		
Simple find	0(1)	Worst: $O(n)$		
Union-by-height &	O(1)	$O(\log n)$		
Simple find	O(1)			
Union-by-height &	O(1)	$O(\alpha(m,n)) = O(\log^* n)$ close to $O(1)$		
Find with path compression	0(1)	close to $O(1)$		
		5		

Priority queues					
Operations	Max (Min)	Min-max & Deap & SMMH	Leftist	Binomial	Fibonacci
Insert x	$O(\log n)$	$O(\log n)$	$O(\log n)$	$O(\log n), O(1)^*$	$O(1)^*$
Delete max	$O(\log n)$	$O(\log n)$			
Delete min	O(n)	$O(\log n)$	$O(\log n)$	$O(\log n)$	$O(\log n)^*$
Delete x			MINIA	$O(\log n)$	$O(\log n)^*$
Merge	O(n)		$O(\log n)$	$O(\log n)$	$O(1)^*$
Decrease key				$O(\log n)$	$O(1)^*$
Search x	O(n)				
Find max	O(1)	O(1)			
Find min		O(1)		$O(\log n)$	O(1)
Remark	Create: $O(n)$		Merge faster than Max (Min) heap.	Find min can be down to $O(1)$.	Decrease key is faster than binomial heap

Sorting algorithms					
Method	Time complexity			Space complexity	Stable
Method	Best	Worst Average		Space complexity	Stable
Insertion	O(n)	$O(n^2)$		O(1)	
Selection		$O(n^2)$		O(1)	×
Bubble	O(n)	$O(n^2)$		O(1)	
Shell	$O(n^{1.5})$	$O(n^2)$		O(1)	×
Quick	$O(n \log n)$	$O(n^2)$	$O(n \log n)$	$O(n\log n) \sim O(n)$	×
Merge	$O(n \log n)$		O(n)		
Heap	$O(n \log n)$		O(1)	×	
LSD Radix	$O(n \times k)$		O(n+k)		
Bucket/MSD Radix	O(n)	$O(n^2)$	O(n+k)	$O(n \times k)$	
Counting $O(n+k)$					

	2			
Graph algorithms				
Problem	Time complexity	Remark		
Depth-First Search (DFS)	O(V + E)			
Kosaraju's	O(V + E)			
Kruskal's	$O(E \log V)$	7\\		
Prim's (Adjacency matrix)	$O(V ^2)$			
Prim's (Heap, Adjacency lists)	$O(E \log V)$			
Prim's (Fibonacci heap, Adjacency lists)	$O(E + V \log V)$			
Sollin's (Borůvka's)	$O(E \log V)$			
Dijkstra's (Min-heap)	$\Theta((\lceil V \rceil + E)\log V)$	Greedy, no negative		
Dijkstra's (Fibonacci-heap)	$\Theta(E + V \log V)$	edges or cycles		
Bellman-Ford	O(V E)	DP, no negative cycles		
Floyd-Warshall	$\Theta([V]^3)$	DP, no negative cycles		
Johnson's	$\Theta(V E + V ^2 \log V)$	No negative cycles		
Ford-Fulkerson	$O(E f^*)$	Greedy, f^* 為最大流		
Edmond-Karp	$O(V E ^2)$			

Dynamic Programming algorithms				
Problem	Time complexity	Space complexity	Remark	
Making change	O(kn)	O(n)		
Fractional Knapsack problem	$\Theta(n \log n)$	O(n)	Greedy	
0/1 Knapsack problem (DP)	$O(n2^{\log W})$	$O(n2^{\log W})$		
0/1 Knapsack problem (Branch-and-Bound)	$O(2^n)$			
Longest Common Subsequence (LCS)	O(mn)	O(mn)	不必連續	
Longest Increasing Subsequence (LIS)	$O(n^2)$	$O(n^2)$		
Longest Common Substring	O(mn)	O(mn)	必須連續	
Minimum Edit Distance	O(mn)	O(mn)		
Matrix-chain Multiplication	$O(n^3)$	$O(n^2)$		
Traveling Salesperson problem	$\Theta(n^2 2^n)$	$O(n2^n)$		
Optimal Binary Search Tree (OBST)	$\Theta(n^3)$	$\Theta(n^2)$		

Computational Ge	eometry algorithms
Problem	Time complexity
平面上點的 rank	$\Theta(n \log n)$
Maximal points	$\Theta(n \log n)$
Closest pair	$O(n \log n)$
Farthest pair	$O(n \log n)$
Graham scan	$\Theta(n \log n)$

References

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