

資料結構

Data Structure

TZU-CHUN HSU¹

¹vm3y3rmp40719@gmail.com

¹Department of Computer Science, Zhejiang University



2021 年 1 月 16 日
Version 3.1

Disclaimer

本文「資料結構」為台灣研究所考試入學的「資料結構」考科使用，內容主要參考 Introduction to Algorithms[1]，以及 wjungle 網友在 PTT 論壇上提供的資料結構筆記 [2]。本文作者為 TZU-CHUN HSU，本文及其 L^AT_EX 相關程式碼採用 MIT 協議，更多內容請訪問作者之 GITHUB 分頁 [Oscarshu0719](#)。

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1 Overview

1. 本文頁碼標記依照 TKB 筆記 [2] 的頁碼。

2. TKB 筆記 [2] 章節頁碼：

Chapter	Page No.	Importance
1	3	***
2	259	*
3	52	***
4	259	*
5	82	*****
6	228	*****
7	180	*****
8	221	***
9	129	*****

3.

$$\log 2 = 0.3010$$

$$\log 3 = 0.4771$$

$$\log 5 = 0.6990$$

$$\log 7 = 0.8451$$

(1)

4. OBST 在「演算法」中，不再贅述。

Trees				
Tree	Insert x	Delete x	Search x	Remark
BST	$O(\log n) \sim O(n)$			Create: $O(n \log n) \sim O(n^2)$
AVL tree	$O(\log_m n)$			$F_{h+2} - 1 \leq n \leq 2^h - 1$
B tree				$1 + 2^{\frac{\lceil \frac{m}{2} \rceil^{h-1} - 1}{\lceil \frac{m}{2} \rceil - 1}} \leq n \leq \frac{m^h - 1}{m - 1}$
RBT				$h \leq 2 \log(n + 1)$
Splay tree				Worst: $O(n)$, Amortized: $O(\log n)$

Priority queues					
Operations	Max (Min)	Min-max & Deap & SMMH	Leftist	Binomial	Fibonacci
Insert x	$O(\log n)$	$O(\log n)$	$O(\log n)$	$O(\log n), O(1)^*$	$O(1)^*$
Delete max	$O(\log n)$	$O(\log n)$			
Delete min	$O(n)$	$O(\log n)$	$O(\log n)$	$O(\log n)$	$O(\log n)^*$
Delete x				$O(\log n)$	$O(\log n)^*$
Merge	$O(n)$		$O(\log n)$	$O(\log n)$	$O(1)^*$
Decrease key				$O(\log n)$	$O(1)^*$
Search x	$O(n)$				
Find max	$O(1)$	$O(1)$			
Find min		$O(1)$		$O(\log n)$	$O(1)$
Remark			$shortest(\text{root})$ $\leq \log(n+1) - 1$		

Sorting algorithms					
Method	Time complexity			Space complexity	Stable
	Best	Worst	Average		
Insertion	$O(n)$	$O(n^2)$		$O(1)$	✓
Selection	$O(n^2)$			$O(1)$	×
Bubble	$O(n)$	$O(n^2)$		$O(1)$	✓
Shell	$O(n^{1.5})$	$O(n^2)$		$O(1)$	×
Quick	$O(n \log n)$	$O(n^2)$	$O(n \log n)$	$O(n \log n) \sim O(n)$	×
Merge	$O(n \log n)$			$O(n)$	✓
Heap	$O(n \log n)$			$O(1)$	×
LSD Radix	$O(n \times k)$			$O(n + k)$	✓
Bucket/MSD Radix	$O(n)$	$O(n^2)$	$O(n + k)$	$O(n \times k)$	✓
Counting	$O(n + k)$				✓

2 Summary

1. Theorem (17) Permutation:

```
1: function PERM(list, i, n)
2:   if i == n then
3:     PRINT(list)
4:   else
5:     for j := i to n do
6:       SWAP(list, i, j)
7:       PERM(list, i + 1, n)
8:       SWAP(list, i, j)
9:     end for
10:  end if
11: end function
```

2. Theorem (87) 節點數:

$$n = \left(\sum_{i=1}^{\deg} i \times n_i \right) + 1 \quad (2)$$
$$n_0 = n_2 + 1 \text{ (二叉樹)}$$

3. Theorem (95, 97, 98)

- 可以確定二叉樹，其他則否：
 - Preorder 和 Inorder。
 - Postorder 和 Inorder。
 - Level-order 和 Inorder。
 - Complete 和任意排序。
- Preoder = Inoder: Empty、Root、Right-skewed tree。
- Postoder = Inoder: Empty、Root、Left-skewed tree。
- Preoder = Postoder: Empty、Root。

4. Theorem (116)

```

1: function CREATEMINHEAP(Tree  $s$ , size  $n$ )
2:   for  $i := n/2$  to 1 do                                ▷ Start from parent of the last node.
3:      $tmp := s[i]$ 
4:      $j := 2 \times i$                                          ▷ Left child of  $i$ .
5:     while  $j \leq n$  do                                     ▷ There is a child.
6:       if  $j < n$  then                                     ▷ Right child exists.
7:         if  $s[j] > s[j+1]$  then                             ▷ Choose the smaller child.
8:            $j := j + 1$ 
9:         end if
10:      end if
11:      if  $tmp \leq s[j]$  then
12:        Break.
13:      else                                                ▷ Percolate one level.
14:         $s[j/2] := s[j]$ 
15:         $j := j \times 2$ 
16:      end if
17:    end while
18:     $s[j/2] := tmp$ 
19:  end for
20: end function

```

5. Theorem (195) Quick sorting:

```

1: function QUICKSORT(Array  $A$ , index  $p, r$ )                ▷ Sorting from  $A[p]$  to  $A[r]$ 
2:   if  $p < r$  then
3:      $q := \text{PARTITION}(A, p, r)$ 
4:     QUICKSORT( $A, p, q - 1$ )
5:     QUICKSORT( $A, q + 1, r$ )
6:   end if
7: end function

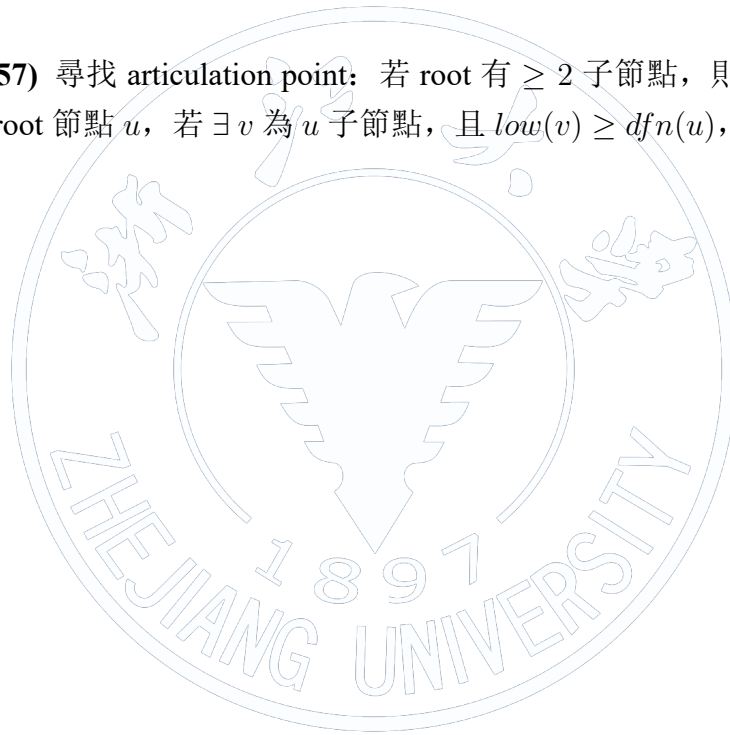
```

```

1: function PARTITION(Array  $A$ , index  $p$ ,  $r$ )
2:    $x := A[r]$  ▷ Pivot.
3:    $i := p - 1$ 
4:   for  $j := p$  to  $r - 1$  do
5:     if  $A[j] \leq x$  then
6:        $i := i + 1$ 
7:       SWAP( $A[i]$ ,  $A[j]$ )
8:     end if
9:   end for
10:  SWAP( $A[r]$ ,  $A[i + 1]$ )
11:  return  $i + 1$ 
12: end function

```

6. **Theorem (257)** 尋找 articulation point: 若 root 有 ≥ 2 子節點, 則 root 為 articulation point; \exists 非 root 節點 u , 若 $\exists v$ 為 u 子節點, 且 $low(v) \geq dfn(u)$, 則 u 為 articulation point。



References

- [1] Thomas H. Cormen, Charles E. Leiserson, Ronald L. Rivest, and Clifford Stein. *Introduction to Algorithms, Third Edition*. The MIT Press, 3 edition, 2009.
- [2] wjungle@ptt. 資料結構 @tkb 筆記. <https://drive.google.com/file/d/0B8-2o6L73Q2VeFpGejlYRk1WeFk/view?usp=sharing>, 2017.

