# Convex Optimization 10-725, Lecture 15: Sum of Square: The Ultimate form of Convex Optimization

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Today

#### So Far

- We have been mainly talking about convex optimization.
- Although all the algorithms we have learned can be applied to general settings (non-convex).
- ullet For example, people use SGD + M / Adam / Projected stochastic gradient descent in deep learning.
- However, we only have convergence guarantees of these algorithms in the convex setting.
- Question: Given a general optimization problem, can we make it convex?

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#### "Convexification"

- Question: Given a general optimization problem, can we make it convex?
- Meaning that given a general set  $\mathcal{D}$  and a general function f where we want to solve

$$\min_{x \in \mathcal{D}} f(x)$$

Can we write an equivalent optimization problem

$$\min_{y \in \mathcal{D}'} g(y)$$

- Such that  $\mathcal{D}'$  is a convex set, g is a convex function, and there is a KNOWN function r such that r(y) = x.
- After finding  $y^*$ , we can apply r and recover  $x^*$ .

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#### "Convexification"

• Example: non convex optimization over a non-convex constraint set:

$$\min_{x_1^4 + x_2^4 \le 1, x_1^2 - x_2^2 \ge 0} x_1^2 - 3x_2^2$$

This is equivalent to a convex optimization

$$\min_{y_1^2+y_2^2\leq 1, y_1-y_2\geq 0, y_1\geq 0, y_2\geq 0}y_1-3y_2$$

- Such that  $x_1^2 = y_1, x_2^2 = y_2$ .
- Question: How do we do it in general?
- For example if we have:

$$\min_{x_1^4 + x_2^4 + x_1 - x_2 \le 1, x_1^2 - x_1 x_2 - x_2^2 \ge 0} x_1^2 - 3x_2^2 + 4x_1$$

#### "Convexification"

- We will talk about the Sum of Squares approach, which can automatically turn any (analytic) optimization problem into a convex optimization problem.
- Sum of Squares approach is very simple, once we understand its spirit.

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- Motivation: Let us first consider un-constraint optimization:  $\min f(x)$  for general function f.
- Example of a very non-convex function:

$$f(x) = x_1^4 - 2x_1x_2 + x_2^4 - x_2^3 + 2x_2 - 3x_1^2$$

- How do we minimize this thing?
- How do we "convexifiy" this thing?

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 Let us start with an analytic function f, those functions can be written as:

$$f(x) = \sum_{\vec{\alpha}: \text{non-negative vectors in } \mathbb{Z}^d} C_{\vec{\alpha}} \prod_{i \in [d]} x_i^{\vec{\alpha}_i}$$

- Starting point: turn min f into a linear optimization problem by
- Creating a new variable  $X_{\vec{\alpha}}$  for each monomial  $\prod_{i \in [d]} x_i^{\vec{\alpha}_i}$ .
- Let  $X = \{X_{\vec{\alpha}}\}_{\vec{\alpha}: \text{non-negative vectors in } \mathbb{Z}^d$ , then we can write f(x) as:

$$G(X) = \sum_{\vec{\alpha}: \text{non-negative vectors in } \mathbb{Z}^d} C_{\vec{\alpha}} X_{\vec{\alpha}}$$

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• Let  $X = \{X_{\vec{\alpha}}\}_{\vec{\alpha}: \text{non-negative vectors in } \mathbb{Z}^d$ , then we can write f(x) as:

$$G(X) = \sum_{\vec{\alpha}: \text{non-negative vectors in } \mathbb{Z}^d} C_{\vec{\alpha}} X_{\vec{\alpha}}$$

• Example:  $f(x) = x_1^4 - 2x_1x_2 + x_2^2$ . Then we have:  $X = (X_{(4,0)}, X_{(1,1)}, X_{(0,2)})$  and

$$G(X) = X_{(4,0)} - 2X_{(1,1)} + X_{(0,2)}$$

 Observation: No matter how complicated (non-convex) f is, G is always linear (convex).

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• Example:  $f(x) = x_1^4 - 2x_1x_2 + x_2^2$ . Then we have:  $X = (X_{(4,0)}, X_{(1,1)}, X_{(0,2)})$  and

$$G(X) = X_{(4,0)} - 2X_{(1,1)} + X_{(0,2)}$$

- Observation: No matter how complicated (non-convex) f is, G is always linear (convex).
- But clearly, minimizing G w.r.t. X can not recover the minimizer of f, because the variables  $X_{\vec{\alpha}}$  are constraint and dependent.
- For example,  $x_2^2 \ge 0$ , so clearly we need to add constraint  $X_{(0,2)} \ge 0$
- But more importantly,

$$x_1^4 + 1/4 \ge x_1^2$$
,  $x_1^2 + 2x_1x_2 + x_2^2 \ge 0$ 

So we should also add constraints like

$$X_{(4,0)} + 1/4 \geq X_{(2,0)}, \quad X_{(2,0)} + 2X_{(1,1)} + X_{(0,2)} \geq 0$$

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But more importantly,

$$x_1^4 + 1/4 \ge x_1^2$$
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So we should also add constraints like

$$X_{(4,0)} + 1/4 \geq X_{(2,0)}, \quad X_{(2,0)} + 2X_{(1,1)} + X_{(0,2)} \geq 0$$

• Main Question: What are the sets of constraints we need to add?

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• Main observation: For every polynomial h(x), we must have

$$h(x)^2 \ge 0$$

.

• Let C(h) be the coefficient of h in monomial basis, so

$$h(x) = \sum_{\vec{\alpha}: \text{non-negative vectors in } \mathbb{Z}^d} C(h)_{\vec{\alpha}} \prod_{i \in [d]} x_i^{\vec{\alpha}_i}$$

So we need to add the constraint:

$$\sum_{\vec{\alpha}: \text{non-negative vectors in } \mathbb{Z}^d} C(h^2)_{\vec{\alpha}} X_{\vec{\alpha}} \geq 0$$

- For every polynomial h.
- Example:  $(x_1 + 1)^2 \ge 0$ , so we need to add  $X_{(2,0)} + 2X_{(1,0)} + 1 \ge 0$ .

To minimize the general function

$$f(x) = \sum_{\vec{\alpha}: \text{non-negative vectors in } \mathbb{Z}^d} C_{\vec{\alpha}} \prod_{i \in [d]} x_i^{\vec{\alpha}_i}$$

We want to alternatively minimize

$$G(X) = \sum_{\vec{\alpha}: \text{non-negative vectors in } \mathbb{Z}^d} C_{\vec{\alpha}} X_{\vec{\alpha}}$$

With constraints:

$$\sum_{\vec{\alpha}: \text{non-negative vectors in } \mathbb{Z}^d} C \big( \, h^2 \big)_{\vec{\alpha}} X_{\vec{\alpha}} \geq 0$$

- For every polynomial *h*.
- Main observation: This is convex optimization over convex constraint set (polytope)!

#### Theorem,

The new minimization problem is equivalent to the original problem, in the sense that

- (1).  $\min_{X \text{ satisfies all constraints}} G(X) = \min_{X} f(x)$ .
- (2). As long as f has a unique minimizer  $x^*$ , then G has a unique minimizer  $X^*$  and

$$x^* = (X^*_{(1,0,0,...,0)}, X^*_{(0,1,0,...,0)}, ..., X^*_{(0,0,0,...,0,1)})$$

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- The new minimization problem is convex, and it is equivalent to the original problem.
- Here, the new variables  $X_{\vec{\alpha}}$  is called the Pseudo-Expectation.
- ullet For notation simplicity, for a polynomial g, we use  $ilde{\mathbb{E}}_X g$  to denote

$$\tilde{\mathbb{E}}_X g = \sum_{\vec{\alpha}} C(g)_{\vec{\alpha}} X_{\vec{\alpha}}$$

• We can alternatively write the minimization as:

 $\min \tilde{\mathbb{E}}_X f$ , s.t. for all polynomial  $h: \tilde{\mathbb{E}}_X h^2 \geq 0$ 

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• Now we need to worry about the constraints:

$$\sum_{\vec{\alpha}: \text{non-negative vectors in } \mathbb{Z}^d} C(h^2)_{\vec{\alpha}} X_{\vec{\alpha}} \geq 0$$

- For every polynomial h.
- Clearly, this is still inefficient, since there are infinitely mainly polynomials h.
- Idea 1: Restrict the degree of  $h^2$  to be no more than D:  $\deg(h^2) \leq D$ . Meaning that we only consider h such that for every  $\vec{\alpha}$  with  $\|\vec{\alpha}\|_1 > D$ ,  $C(h^2)_{\vec{\alpha}} = 0$ .

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So we need to add the constraint:

$$\sum_{\vec{\alpha}: \text{non-negative vectors in } \mathbb{Z}^d} C(h^2)_{\vec{\alpha}} X_{\vec{\alpha}} \geq 0$$

- For every polynomial h.
- Clearly, this is inefficient, since there are infinitely mainly polynomials
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- Idea 1: Restrict the degree of  $h^2$  to be no more than D:  $\deg(h^2) \leq D$ . Meaning that we only consider h such that for every  $\vec{\alpha}$  with  $\|\vec{\alpha}\|_1 > d$ ,  $C(h^2)_{\vec{\alpha}} = 0$ .
- So we only add constraints

$$\sum_{\vec{\alpha}: \mathsf{non-negative} \ \mathsf{vectors} \ \mathsf{in} \ \mathbb{Z}^d} C(\mathit{h}^2)_{\vec{\alpha}} X_{\vec{\alpha}} \geq 0$$

- For every polynomial h with  $deg(h^2) \le D$ .
- *D* is some hyper-parameter you will choose in practice.

So we only add constraints

$$\sum_{\vec{\alpha}: \text{non-negative vectors in } \mathbb{Z}^d} C(h^2)_{\vec{\alpha}} X_{\vec{\alpha}} \geq 0$$

- For every polynomial h with  $deg(h^2) \le D$ .
- But still, there are infinitely many polynomials h with  $deg(h^2) \le D$ .

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So we only add constraints

$$\sum_{\vec{\alpha}: \text{non-negative vectors in } \mathbb{Z}^d} C(h^2)_{\vec{\alpha}} X_{\vec{\alpha}} \geq 0$$

- For every polynomial h with  $deg(h^2) \le D$ .
- Main observation: Let  $M_D(X)$  be the matrix such that

$$[M_D(X)]_{\vec{\alpha},\vec{\alpha'}} = X_{\vec{\alpha}+\vec{\alpha'}}$$

for all  $\vec{\alpha}, \vec{\alpha'}$ : non-negative vectors in  $\mathbb{Z}^d$  with  $\|\vec{\alpha}\|_1 \leq D/2, \|\vec{\alpha'}\|_1 \leq D/2$ 

 Then the original set of constraints is equivalent to the following constraint:

$$M_D(X) \geq 0$$

• This is because  $\sum_{\vec{\alpha}} C(h^2)_{\vec{\alpha}} X_{\vec{\alpha}} = c^{\mathsf{T}} M_D(X) c$  for  $c = (C(h)_{\vec{\alpha}})_{\vec{\alpha}}$ .

#### Degree D Sum of Squares

- Degree D Sum of Squares for un-constraint optimization: To minimize min f(x):
- It is equivalent to minimizing the new problem:

$$G(X) = \sum_{\vec{\alpha}: \text{non-negative vectors in } \mathbb{Z}^d} C_{\vec{\alpha}}(f) X_{\vec{\alpha}}$$

With constraints that form a convex set:

$$M_D(X) \geq 0$$

• Where for all  $\vec{\alpha}$ ,  $\vec{\alpha'}$ : non-negative vectors in  $\mathbb{Z}^d$  with  $\|\vec{\alpha}\|_1 \le D/2, \|\vec{\alpha}'\|_1 \le D/2$  ":

$$[M_D(X)]_{\vec{\alpha},\vec{\alpha'}} = X_{\vec{\alpha}+\vec{\alpha'}}$$

- Then  $x^* = (X_{(1,0,\ldots,0)}, X_{(0,1,\ldots,0)}, X_{(0,0,\ldots,1)}).$
- This is the Sum of Squares program, where  $X_{\vec{\alpha}}$  are called the "pseudo-expectations".
- We denote  $\tilde{E}(f) := \sum_{\vec{c}} C_{\vec{c}}(f) X_{\vec{c}}$

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• To solve  $\min_{X:M_D(X)\geq 0}G(X)$ , we need to use the interior point method, the self-concordant barrier function for this PSD cone is

$$R(X) = -\log \det(M_D(X))$$

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- For constraint minimization min f(x) such that  $\{h_i(x) \ge 0\}_{i \in [m]}$ :
- Choice 1: Use duality to turn it into unconstraint optimization over *x*.
- Choice 2: We still minimize

$$G(X) = \sum_{\vec{\alpha}: \text{non-negative vectors in } \mathbb{Z}^d} C_{\vec{\alpha}} X_{\vec{\alpha}}$$

Now we need to add new constraints:

$$\sum_{\vec{\alpha}: \mathsf{non-negative} \ \mathsf{vectors} \ \mathsf{in} \ \mathbb{Z}^d} C \big( h_i h^2 \big)_{\vec{\alpha}} X_{\vec{\alpha}} \geq 0$$

- For every  $i \in [m]$  and every polynomial h.
- Theorem: With these new constraints, the optimization is also equivalent to min f(x) such that  $\{h_i(x) \ge 0\}_{i \in [m]}$  (not going to prove here)

- Main Application: Robust linear regression.
- Suppose we are given N data points  $(x^{(i)}, y^{(i)})$ ,  $1 \tau$  fraction of them satisfies  $(y \langle w^*, x \rangle)^2 \le \varepsilon^2$ , and  $\tau$  fraction of the data is arbitrarily corrupted (both x and y).
- Solving linear regression

$$\min_{w} \sum_{i \in [N]} \left( \langle x^{(i)}, w \rangle - y^{(i)} \right)^{2}$$

will not work.

We want to solve the Robust Linear Regression

$$\min_{\alpha_i} \min_{w} \sum_{i \in [N]} \alpha_i \left( \langle x^{(i)}, w \rangle - y^{(i)} \right)^2$$

Such that  $\alpha_i^2 = \alpha_i$ ,  $\frac{1}{N} \sum_i \alpha_i \ge 1 - \tau$ .

- Main Application: Robust linear regression.
- We want to solve the Robust Linear Regression

$$\min_{\alpha_i} \min_{w} \sum_{i \in [N]} \alpha_i \left( \langle x^{(i)}, w \rangle - y^{(i)} \right)^2$$

Such that  $\alpha_i^2 = \alpha_i$ ,  $\frac{1}{N} \sum_i \alpha_i \ge 1 - \tau$ .

- This is non-convex optimization.
- Main Claim: We can use degree 8 sum of squares (over both  $\alpha$  and w) to solve this.

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We want to solve the Robust Linear Regression

$$\min_{\alpha_i} \min_{w} \sum_{i \in [N]} \alpha_i \left( \langle x^{(i)}, w \rangle - y^{(i)} \right)^2$$

Such that  $\alpha_i^2 = \alpha_i$ ,  $\frac{1}{N} \sum_i \alpha_i \ge 1 - \tau$ .

• Let  $\tilde{x}^{(i)}, \tilde{y}^{(i)}$  be the uncorrupted data, such that for every i,

$$(\tilde{y}^{(i)} - \langle w^*, \tilde{x}^{(i)} \rangle)^2 \le \varepsilon^2$$

- Let  $\tilde{\alpha}_i = 0$  if i is corrupted, and  $\tilde{\alpha}_i = 1$  otherwise (we do not know  $\tilde{\alpha}_i = 0$ , they are just for proof purpose).
- Then

$$1 = \alpha_i \tilde{\alpha}_i + \alpha_i (1 - \tilde{\alpha}_i) + (1 - \alpha_i)$$

Since

$$1 = \alpha_i \tilde{\alpha}_i + \alpha_i (1 - \tilde{\alpha}_i) + (1 - \alpha_i)$$

So we have:

$$\sum_{i \in [N]} \left( \langle \tilde{x}^{(i)}, w \rangle - \tilde{y}^{(i)} \right)^{2}$$

$$= \sum_{i \in [N]} \alpha_{i} \tilde{\alpha}_{i} \left( \langle \tilde{x}^{(i)}, w \rangle - \tilde{y}^{(i)} \right)^{2}$$

$$+ \sum_{i \in [N]} \alpha_{i} (1 - \tilde{\alpha}_{i}) \left( \langle \tilde{x}^{(i)}, w \rangle - \tilde{y}^{(i)} \right)^{2}$$

$$+ \sum_{i \in [N]} (1 - \alpha_{i}) \left( \langle \tilde{x}^{(i)}, w \rangle - \tilde{y}^{(i)} \right)^{2}$$

So we have:

$$\begin{split} \sum_{i \in [N]} \left( \langle \tilde{x}^{(i)}, w \rangle - \tilde{y}^{(i)} \right)^2 \\ &= \sum_{i \in [N]} \alpha_i \tilde{\alpha}_i \left( \langle \tilde{x}^{(i)}, w \rangle - \tilde{y}^{(i)} \right)^2 + \sum_{i \in [N]} \alpha_i (1 - \tilde{\alpha}_i) \left( \langle \tilde{x}^{(i)}, w \rangle - \tilde{y}^{(i)} \right)^2 \\ &+ \sum_{i \in [N]} (1 - \alpha_i) \left( \langle \tilde{x}^{(i)}, w \rangle - \tilde{y}^{(i)} \right)^2 \end{split}$$

• This means: (since f = g implies  $\tilde{E}(f) = \tilde{E}(g)$ )

$$\sum_{i \in [N]} \tilde{E}\left(\langle \tilde{x}^{(i)}, w \rangle - \tilde{y}^{(i)}\right)^2$$

$$= \sum_{i \in [N]} \tilde{E} \alpha_i \tilde{\alpha}_i \left( \langle \tilde{x}^{(i)}, w \rangle - \tilde{y}^{(i)} \right)^2 + \sum_{i \in [N]} \tilde{E} \alpha_i (1 - \tilde{\alpha}_i) \left( \langle \tilde{x}^{(i)}, w \rangle - \tilde{y}^{(i)} \right)^2$$

$$+\sum \tilde{E}(1-lpha_i)\left(\langle \tilde{x}^{(i)},w \rangle - \tilde{y}^{(i)}\right)^2$$

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- So we have:
- First term:

$$\sum_{i \in [N]} \tilde{E} \alpha_i \tilde{\alpha}_i \left( \langle \tilde{x}^{(i)}, w \rangle - \tilde{y}^{(i)} \right)^2$$

$$= \sum_{i \in [N]} \tilde{E} \alpha_i \left( \langle \tilde{x}^{(i)}, w \rangle - \tilde{y}^{(i)} \right)^2$$

$$- \sum_{i \in [N]} \tilde{E} \alpha_i (1 - \tilde{\alpha}_i) \left( \langle \tilde{x}^{(i)}, w \rangle - \tilde{y}^{(i)} \right)^2$$

The last term is a (negative) sum of squares, so we know that

$$\sum_{i \in [N]} \tilde{E} \alpha_i \tilde{\alpha}_i \left( \langle \tilde{x}^{(i)}, w \rangle - \tilde{y}^{(i)} \right)^2$$

$$\leq \sum_{i \in [N]} \tilde{E} \alpha_i \left( \langle \tilde{x}^{(i)}, w \rangle - \tilde{y}^{(i)} \right)^2$$

Which is the objective value of the Sum of Squares program.

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Second term:

$$\sum_{i \in [N]} \tilde{E} \alpha_i (1 - \tilde{\alpha}_i) \tilde{\alpha}_i \left( \langle \tilde{x}^{(i)}, w \rangle - \tilde{y}^{(i)} \right)^2$$

• By Cauchy-Shwartz:  $\tilde{E}^2(fg) \leq \tilde{E}f^2\tilde{E}g^2$ ), this is because of  $\tilde{E}(f\tilde{E}(fg) - g\tilde{E}f^2)^2 \geq 0$ , so  $\tilde{E}f^2\tilde{E}^2(fg) - 2\tilde{E}^2(fg)\tilde{E}f^2 + (\tilde{E}f^2)^2\tilde{E}g^2 \geq 0$ . So

$$\left(\frac{1}{N}\sum_{i\in[N]}\tilde{E}\alpha_{i}(1-\tilde{\alpha}_{i})\left(\langle\tilde{x}^{(i)},w\rangle-\tilde{y}^{(i)}\right)^{2}\right)^{2}$$

$$\leq\left(\frac{1}{N}\sum_{i\in[N]}\tilde{E}(1-\tilde{\alpha}_{i})^{2}\right)\left(\frac{1}{N}\tilde{E}\sum_{i\in[N]}\alpha_{i}^{2}\left(\langle\tilde{x}^{(i)},w\rangle-\tilde{y}^{(i)}\right)^{4}\right)^{2}$$

$$\leq\tau\left(\frac{1}{N}\sum_{i\in[N]}\tilde{E}\alpha_{i}^{2}\left(\langle\tilde{x}^{(i)},w\rangle-\tilde{y}^{(i)}\right)^{4}\right)\leq\tau\left(\frac{1}{N}\sum_{i\in[N]}\tilde{E}\left(\langle\tilde{x}^{(i)},w\rangle-\tilde{y}^{(i)}\right)^{4}\right)$$

Recall  $\tilde{\alpha}_i = 0$  if i is corrupted, and  $\tilde{\alpha}_i = 1$  otherwise

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$$\tau\left(\frac{1}{N}\sum_{i\in[N]}\tilde{E}\alpha_{i}^{2}\left(\left\langle \tilde{x}^{(i)},w\right\rangle -\tilde{y}^{(i)}\right)^{4}\right)\leq\tau\left(\frac{1}{N}\sum_{i\in[N]}\tilde{E}\left(\left\langle \tilde{x}^{(i)},w\right\rangle -\tilde{y}^{(i)}\right)^{4}\right)$$

• This is because  $(\alpha_i^2 = \alpha_i)$  is in the Sum of Square constraint)

$$\tilde{E}\alpha_i^2 f^2 - \tilde{E}f^2 = \tilde{E}(\alpha_i^2 - 1)f^2$$
$$= \tilde{E}(-\alpha_i^2 + 2\alpha_i - 1)f^2$$
$$= -\tilde{E}(\alpha_i - 1)^2 f^2 \le 0$$

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• Third term can be done similarly, so we get:

$$\frac{1}{N} \sum_{i \in [N]} \tilde{E}\left(\langle \tilde{x}^{(i)}, w \rangle - \tilde{y}^{(i)}\right)^2 \leq \frac{1}{N} \mathsf{Sum of Square Objective}$$

$$+O(\tau)^{1/2}\left(\frac{1}{N}\sum_{i\in[N]}\tilde{E}\left(\langle \tilde{x}^{(i)},w\rangle-\tilde{y}^{(i)}\right)^4\right)^{1/2}$$

ullet If  $ilde{x}^{(i)}$ 's are Gaussian, then by hyper-contractivity, we have

$$\left(\frac{1}{N}\sum_{i\in[N]}\tilde{E}\left(\langle \tilde{x}^{(i)},w\rangle-\tilde{y}^{(i)}\right)^{4}\right)^{1/2}\leq O(1)\frac{1}{N}\sum_{i\in[N]}\tilde{E}\left(\langle \tilde{x}^{(i)},w\rangle-\tilde{y}^{(i)}\right)^{2}$$

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Therefore,

$$\frac{1}{N} \sum_{i \in [N]} \tilde{E} \left( \langle \tilde{x}^{(i)}, w \rangle - \tilde{y}^{(i)} \right)^2 \le \frac{1}{N} \mathsf{Sum of Square Objective}$$

$$+O(\tau)^{1/2}\frac{1}{N}\sum_{i\in[N]}\tilde{E}\left(\langle\tilde{x}^{(i)},w\rangle-\tilde{y}^{(i)}\right)^2$$

Which means

$$\frac{1}{N} \sum_{i \in [N]} \tilde{E} \left( \langle \tilde{x}^{(i)}, w \rangle - \tilde{y}^{(i)} \right)^2 \leq \frac{1 + O(1/\tau^{1/2})}{N} \text{Sum of Square Objective}$$

Which means

$$\frac{1}{N} \sum_{i \in [N]} \left( \langle \tilde{x}^{(i)}, \tilde{E}[w] \rangle - \tilde{y}^{(i)} \right)^2 \leq \frac{1 + O(1/\tau^{1/2})}{N} \text{Sum of Square Objective}$$

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• Last inequality we uses:

$$\tilde{E}(g(w))^{2} - (g(\tilde{E}[w]))^{2}$$

$$= \tilde{E}(g(w) - g(\tilde{E}[w]))^{2} \ge 0$$

Which means

$$\frac{1}{N} \sum_{i \in [N]} \left( \langle \tilde{x}^{(i)}, \tilde{E}[w] \rangle - \tilde{y}^{(i)} \right)^2 \leq \frac{1 + O(1/\tau^{1/2})}{N} \text{Sum of Square Objective}$$

$$\leq (1 + O(1/\tau^{1/2}))\varepsilon^2$$

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