

Assignment 2

1) Prove that if T is a tree with a vertex of degree k , then T has at least k leaves.

Proof:

Suppose T is a tree.

Let v be a vertex in T with degree k .

Thus v has k neighbors.

Since T is a tree each neighbor of v is either a leaf or has degree at least 2.

If all neighbors of v are leaves then T has at least k leaves.

If a neighbor, u , has degree at least 2, then by removing edge uv we have that the component containing u is still a tree and contains at least one leaf.

Therefore if T is a tree with a vertex of degree k , then T has at least k leaves.

QED

2) Prove that an edge e is a bridge in G if and only if e belongs to every spanning tree of G

Proof:

Suppose e is a bridge in G .

Assume G is a connected graph.

Therefore removing e will increase the number of components from 1 to 2.

Now consider every spanning tree of G .

Any spanning tree is a connected graph which contains all vertices of G .

Therefore, if e was removed from G we would gain 2 components and G would be

Therefore e must belong to every spanning tree of G to keep each component connected.

Suppose now, e is not a bridge.

Therefore removing e will not disconnect G .

Thus, some other path from the endpoints of e exists.

Therefore, e does not belong to every spanning tree.

Therefore, e is a bridge in G if and only if e belongs to every spanning tree of G .

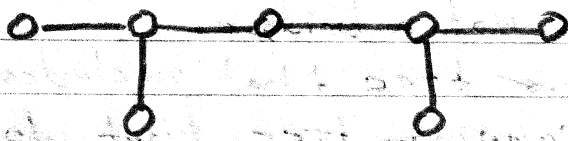
QED

3) Prove that there exists a pair of non-isomorphic graphs with the same degree sequence.

Proof:



Degree Sequence: 1 1 1 2 3 3



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4) Prove that there is only one positive integer k such that no graph contains exactly k spanning trees.

Proof:

K_1 and P_2 have only 1 spanning tree.

C_n has n spanning trees where $n \geq 3$

Toward a contradiction, suppose a graph G has 2 spanning trees.

Let u and v be two vertices in G

Assume edge uv is not a bridge.

Let T be the spanning tree that includes uv and T' be the spanning tree that does not include uv .

Since uv is not a bridge there exists some path from u to v in T' .

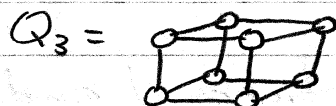
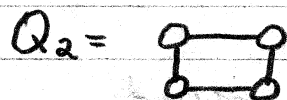
However, this implies the path from u to v in T' plus the original edge uv would form a cycle.

But C_n has n spanning trees where $n \geq 3$.

Therefore, there is only one positive integer k ($k=2$) such that no graph contains exactly k spanning trees.

QED

5/ How many automorphisms of Q_3 exist:-



Q_2 has 8 automorphisms:-

itself

itself rotated 90° 3 times

itself mirrored

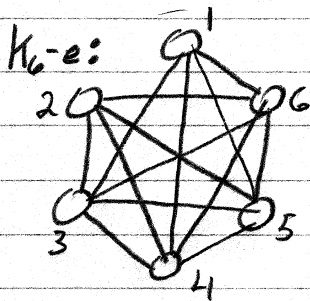
itself mirrored then rotated 90° 3 times

Therefore, each face has 8 automorphisms
and Q_3 has 6 faces, thus:

Q_3 has $6 \times 8 = 48$ automorphisms

6) Let $G = K_6 - e$.

How many automorphisms of G exist



Vertices 1 and 2 must stay in the same locations therefore vertices 3, 4, 5 and 6 can be moved around in any permutation thus 4! Therefore G has $2 \cdot 4! = 48$ automorphisms.

3, 4, 5 and 6 have $4!$ permutations for each 1 and 2 combo which is 2 possibilities