



# THE UNIVERSITY OF ARIZONA

## DEPARTMENT OF COMPUTER SCIENCE

# CSc 453: C-- Language Specification

[Lexical rules](#) | [Syntax rules](#) | [Semantic checking](#) | [Execution behavior](#)

**Notation:** The lexical and syntax rules given below use a notation called EBNF to specify patterns. EBNF notation characters are written in **magenta**. You can find more information about EBNF [here](#).

## 1. Lexical Rules

### 1.1. Notation

To reduce clutter, we use the following names in the patterns listed below:

*letter* ::= a | b | ... | z | A | B | ... | Z  
*digit* ::= 0 | 1 | ... | 9

### 1.2. Tokens

The tokens of the language are as follows:

<u>Name</u>	<u>Pattern</u>	<u>Comments</u>
ID	<i>letter</i> { <i>letter</i>   <i>digit</i>   _ }	identifier: e.g., x, abc, p_q_12
INTCON	<i>digit</i> { <i>digit</i> }	integer constant: e.g., 12345
LPAREN	(	Left parenthesis
RPAREN	)	Right parenthesis
LBRACE	{	Left curly brace
RBRACE	}	Right curly brace
COMMA	,	Comma
SEMI	;	Semicolon
kwINT	<b>int</b>	Keyword: <b>int</b>
kwIF	<b>if</b>	Keyword: <b>if</b>
kwELSE	<b>else</b>	Keyword: <b>else</b>
kwWHILE	<b>while</b>	Keyword: <b>while</b>
kwRETURN	<b>return</b>	Keyword: <b>return</b>
opASSG	=	Op: Assignment
opADD	+	Op: addition
opSUB	-	Op: subtraction
opMUL	*	Op: multiplication
opDIV	/	Op: division
opEQ	==	Op: equals
opNE	!=	Op: not-equals
opGT	>	Op: greater-than
opGE	>=	Op: greater than or equal
opLT	<	Op: less-than

opLE	<=	Op: less than or equal
opAND	&&	Op: logical and
opOR		Op: logical or

### 1.3. Comments

Comments are as in C, i.e. a sequence of characters preceded by `/*` and followed by `*/`, and not containing any occurrence of `*/`.

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## 2. Syntax Rules

Nonterminals are shown in lower-case italics; terminals are shown in boldface or upper-case. The symbol ‘ $\epsilon$ ’ (“epsilon”) denotes the empty sequence.

The grammar rules for full C-- are given below. Changes to the grammar relative to that for the [G2 subset of the language](#) are shown highlighted [here](#).

The start symbol of the grammar is *prog*.

### 2.1 Grammar Productions

```

prog           : func_defn prog
                | var_decl prog
                |  $\epsilon$ 

var_decl      : type id_list SEMI

id_list       : ID
                | ID COMMA id_list

type          : kwINT

func_defn     : type ID LPAREN opt_formals RPAREN LBRACE opt_var_decls opt_stmt_list
                | RBRACE

opt_formals   :  $\epsilon$ 
                | formals

formals       : type ID COMMA formals
                | type ID

opt_var_decls :  $\epsilon$ 
                | var_decl opt_var_decls

opt_stmt_list : stmt opt_stmt_list
                |  $\epsilon$ 

stmt          : fn_call SEMI
                | while_stmt
                | if_stmt
                | assg_stmt
                | return_stmt
                | LBRACE opt_stmt_list RBRACE
                | SEMI

```

```

if_stmt      : kwIF LPAREN bool_exp RPAREN stmt
               | kwIF LPAREN bool_exp RPAREN stmt kwELSE stmt

while_stmt   : kwWHILE LPAREN bool_exp RPAREN stmt

return_stmt  : kwRETURN SEMI
               : kwRETURN arith_exp SEMI

assg_stmt    : ID opASSG arith_exp SEMI

fn_call      : ID LPAREN opt_expr_list RPAREN

opt_expr_list :  $\epsilon$ 
               | expr_list

expr_list    : arith_exp COMMA expr_list
               | arith_exp

arith_exp    : ID
               | INTCON
               | arith_exp arithop arith_exp
               | LPAREN arith_exp RPAREN
               | opSUB arith_exp
               | fn_call

bool_exp     : arith_exp relop arith_exp
               | bool_exp logical_op bool_exp

arithop      : opADD
               | opSUB
               | opMUL
               | opDIV

relop        : opEQ
               | opNE
               | opLE
               | opLT
               | opGE
               | opGT

logical_op   : opAND
               | opOR

```

## 2.2. Operator Associativities and Precedences

The following table gives the associativities of various operators and their relative precedences. An operator with a higher precedence binds "tighter" than one with lower precedence. Precedences decrease as we go down the table.

### 2.2.1. Arithmetic Operators

<u>Operator</u>	<u>Associativity</u>
opSUB (unary)	right to left
opMUL, opDIV	left to right
opADD, opSUB (binary)	left to right

### 2.2.2. Boolean Operators

**Operator    Associativity**

opAND        left to right

opOR         left to right

## 3. Semantic Rules

### 3.1. Scopes and types

There are two kinds of scope in the C-- subset of C--: (1) *global* scope; and (2) for each function in the program, the scope *local* to that function.

An identifier in C-- has one of two possible types in C--: (1) an **int** variable; or (2) a function.

Variables can be declared as globals or as locals. However, the grammar rules of C-- are such that functions can only be defined as globals. Thus, we can have the following possible combinations of scope and type:

	Global	Local
Variable	Yes	Yes
Function	Yes	No

### 3.2. Declarations

The scope of an identifier in a program is given as follows:

1. A variable declared outside a function definition has global scope.
2. A variable declared within a function definition has scope local to that function.
3. The formal parameters of a function have scope local to that function.

The type of an identifier in a program is given as follows:

1. An identifier declared as a function name has type *function*.
2. An identifier declared as a variable (i.e., not as a function) has type *variable*.

### 3.3. Uses

C-- follows the commonly used rule that a use of an identifier refers to the most deeply nested declaration enclosing that use. Since C-- has only global and local scopes, this translates to the following:

1. If an identifier *x* is declared as a local within a function *f* then uses of *x* within *f* refer to this local.
2. Otherwise, if *x* is declared as a global prior to the definition of the function *f* then uses of *x* within *f* refer to this global.
3. Otherwise, any use of *x* within *f* has no declaration to refer to.

A C-- program must satisfy the following requirements:

1. An identifier can be declared at most once as a global and at most once as a local within any particular function.

Note that an identifier can be declared as local in multiple functions (since each function has its own distinct local scope).

2. An identifier *x* that is used within a function body (i.e., which occurs in a statement or expression in the function body) must have been declared prior to the use. The corresponding declaration (see Section 3.3 above) of *x* must satisfy the following requirements:
  1. If the use of *x* is as a function call:
    - the corresponding declaration of *x* should be that of a function; and
    - the number of arguments in the call must be equal to the number of arguments in *x*'s declaration
  2. If the use of *x* is not a function call, the corresponding declaration of *x* should be that of a variable.

## 4. Execution behavior

The C-- language has the execution characteristics expected of a C-like block-structured language. The description below mentions only a few specific points that are likely to be of interest. For points not mentioned explicitly, you should consider the behavior of C-- to be as for C.

## 4.1. Data

An object of type `int` occupies 32 bits.

## 4.2. Order of Evaluation

1. The evaluation order of the operands of expressions with the following kinds of operators is unspecified: arithmetic operators (`opADD`, `opSUB`, `opMUL`, or `opDIV`) and relational operators (`opEQ`, `opNE`, `opGT`, `opGE`, `opLT`, `opLE`). This means that the operands of such expressions can be evaluated in any order.
2. Expressions involving the logical operators `opAND` and `opOR` must be evaluated using short circuit evaluation.
3. The order in which the actual parameters in a function call are evaluated is unspecified.

## 4.3. Program execution

1. Program execution begins at a function named `main()`.

**Note:** You are not required to check whether the program being compiled defines a function named `main()`. If a program does not define `main()`, SPIM will generate an error message.

2. Execution returns from a function if either an explicit **return** statement is executed, or if execution "falls off" the end of the function body. In the latter case, no value is returned.
3. Programs will use a function `println()` to print out integer values. Code for this function will not be part of the input program, but will be generated by your compiler as a hard-coded sequence of MIPS instructions as discussed in the document "*Translating Three-Address Code to MIPS Assembly Code*". This function will behave as though it was defined as

```
void println(int x) { printf("%d\n", x); }
```

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