CHAOS

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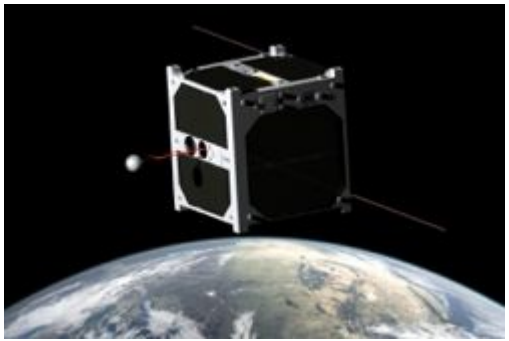
Presentation By: Gregor Peach

Motivation

Mixed Criticality Systems

- Hear me out, but what if have a system that does

TWO THINGS!!!!
(big wow)



Low and High Criticality

High Criticality

- Important tasks
- Highly tested code

Low Criticality

- Unimportant tasks
- Less tested code

State of the Art: Real-time Systems

Classic Real-time Operating Systems

- eg. FreeRTOS
- Pros:
  - Predictable execution (real-time)
  - Single core
  - High assurance
- Cons:
  - No multicore
  - No rich POSIX APIs

State of the Art: POSIX

LINUX

- eg. Ubuntu
- Pros:
  - Rich programming environment
  - Good hardware support
  - Real-time patch available
- Cons:
  - Really big and complex
  - Real-time performance terrible
  - Big surface area for attacks

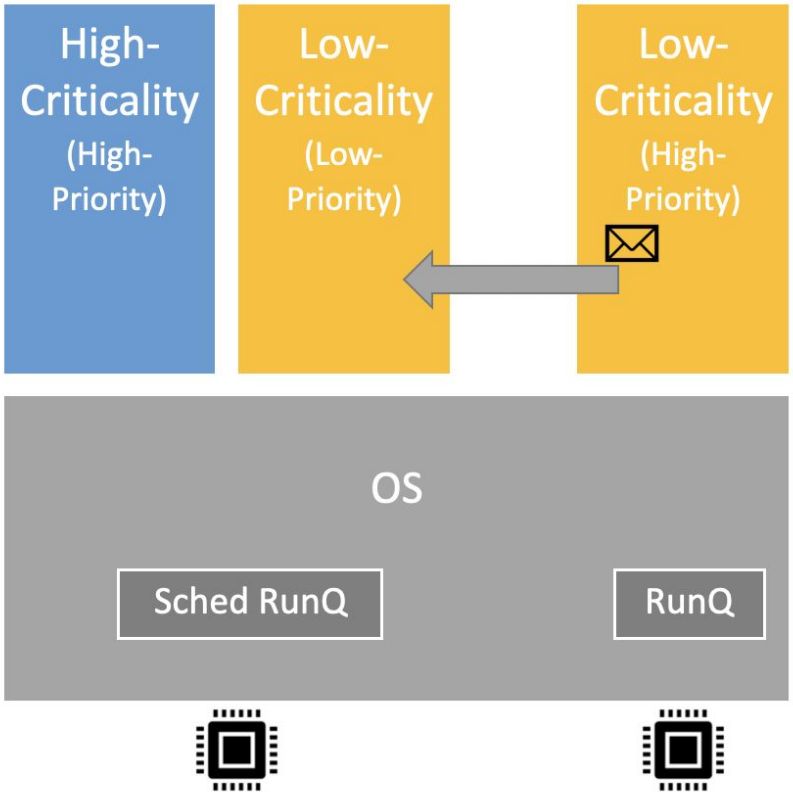
Mixed Criticality Software Lacks a Clear Home

If I have a mixed criticality system, what platform should I build it on?

What are the pros and cons of running it on FreeRTOS?

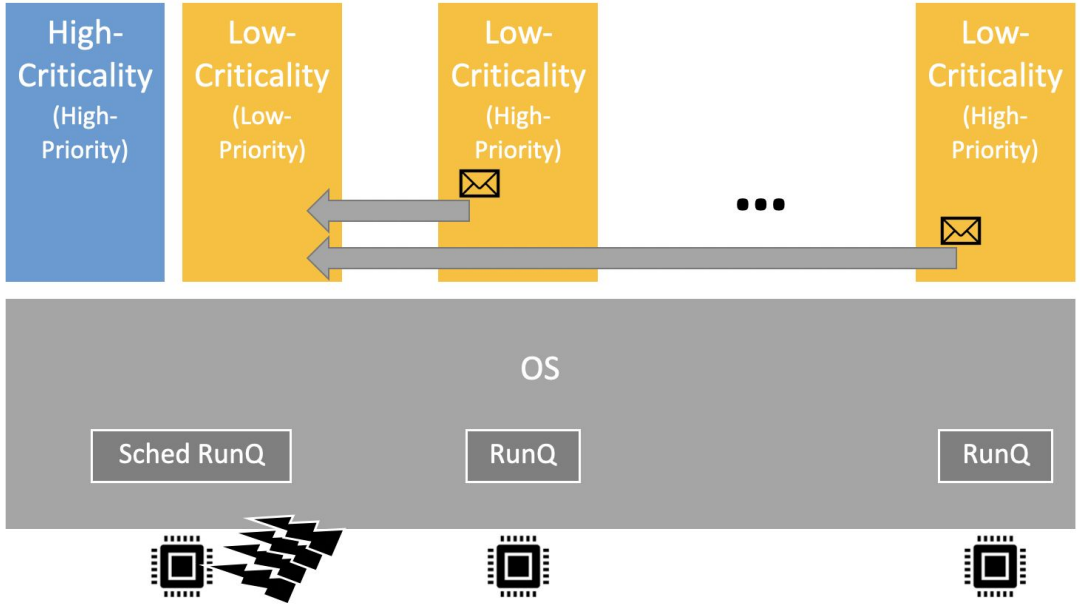
Linux?

This Setup is Unlikely to Work

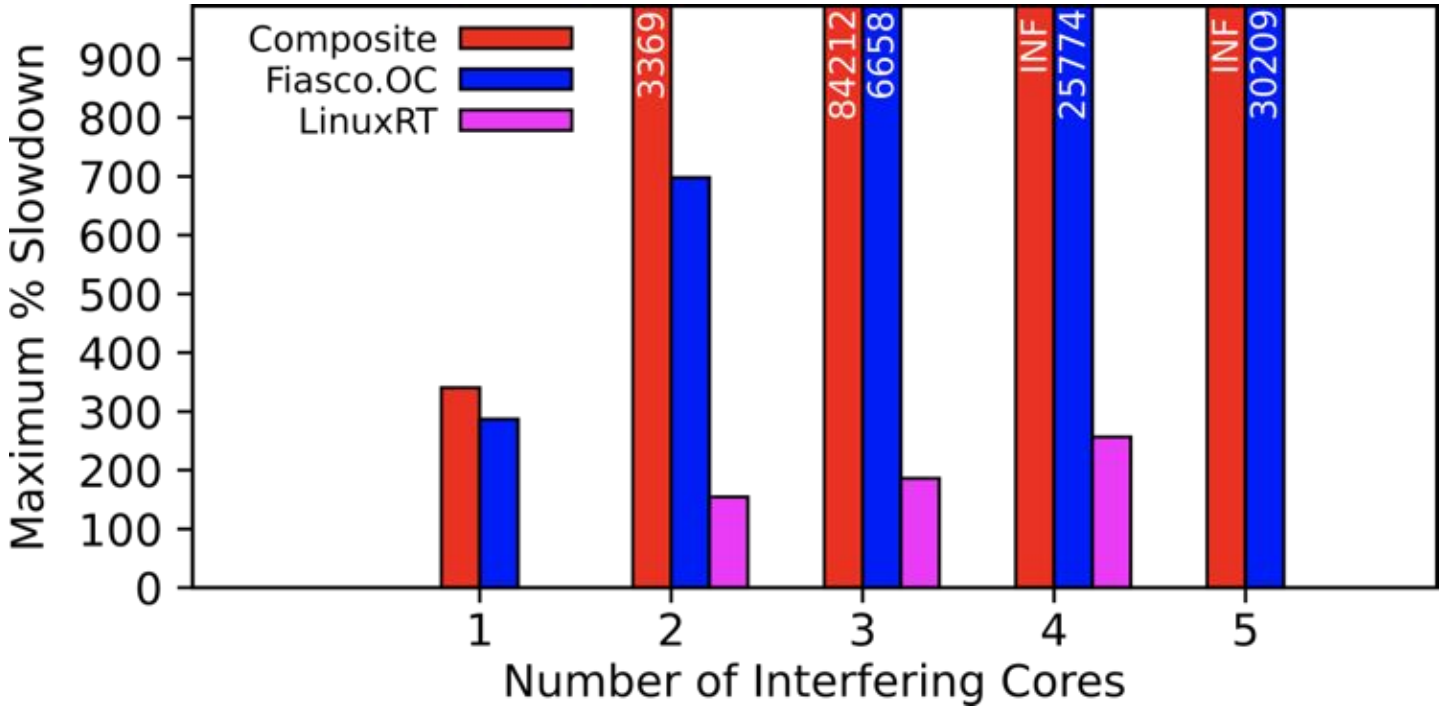


And Even If It Did...

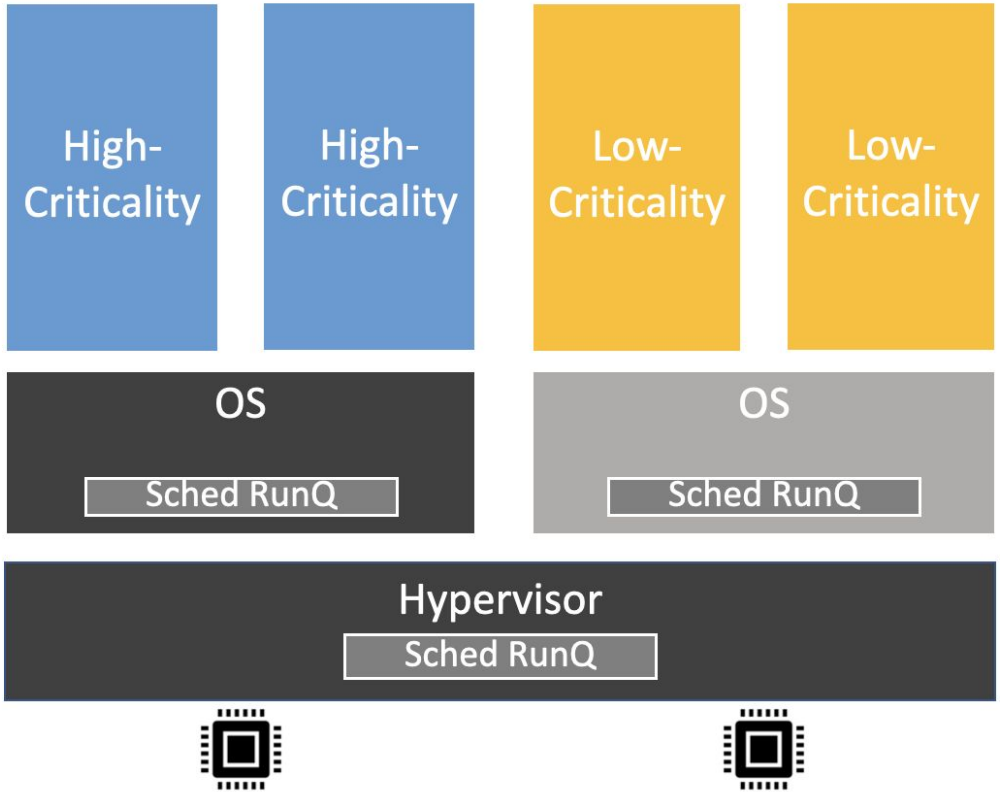
- Naive core sharing has a problem
  - Cross core interrupt denial of service attack



Performance Problem!

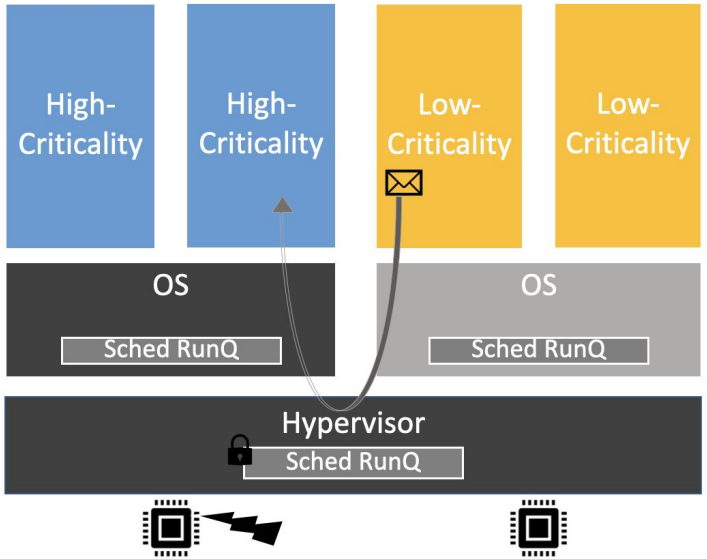


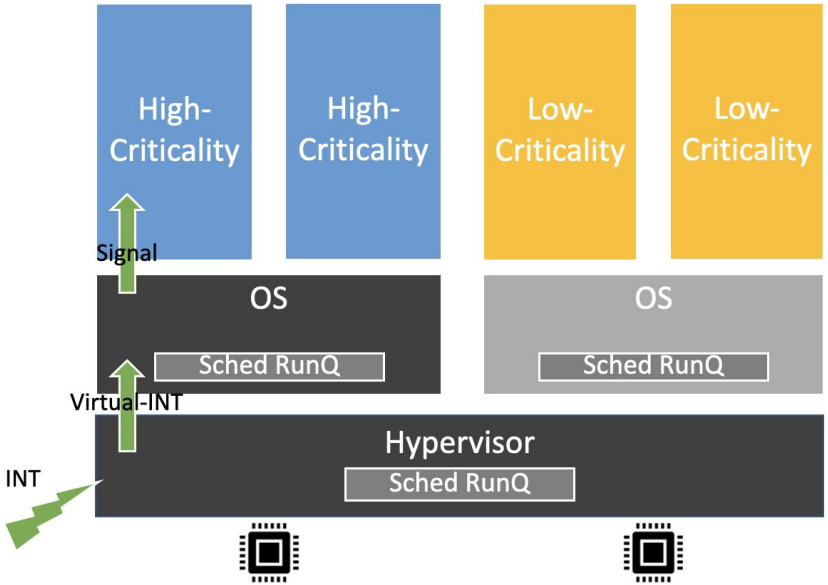
Solution in Practice: Per Core Virtualization



Solution in Practice: Per Core Virtualization

Interrupts handled properly...



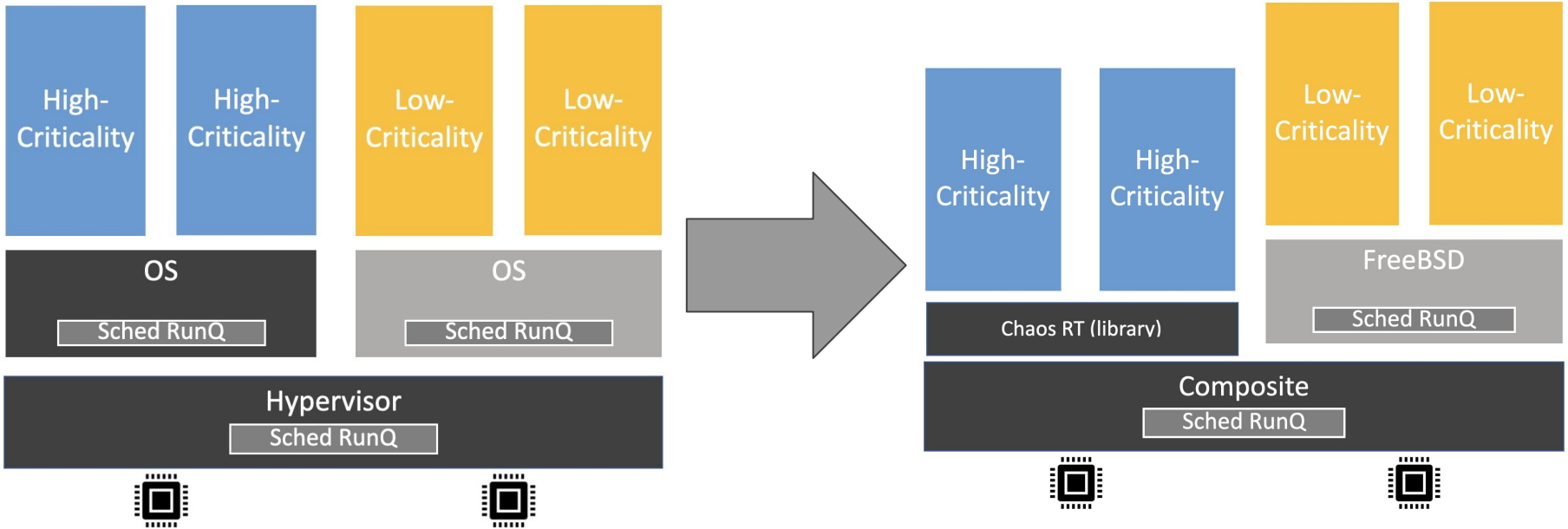


Problems with Per Core Virtualization

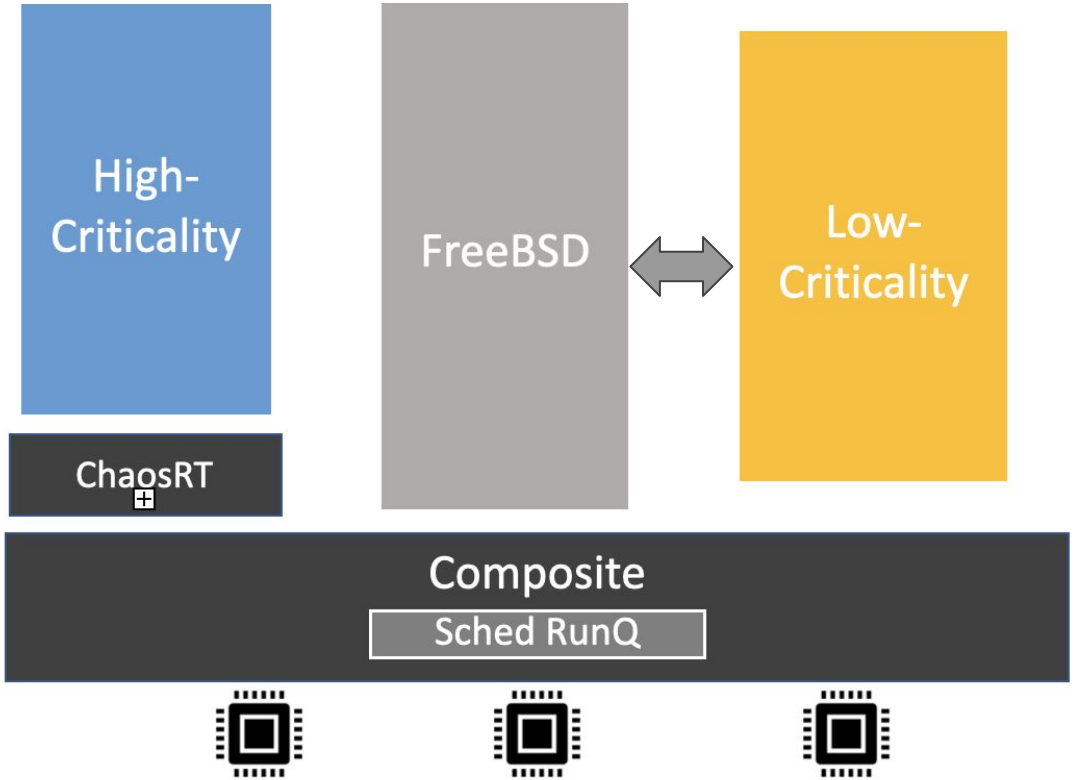
- It’s super slow, big overheads for normal computation
  - Real-time worst case bounds are subject to variance (Hierarchical Overheads)
- Slower to coordinate between high and low assurance

What is CHAOS?

CHAOS: Presenting Devirtualization!!



CHAOS: Presenting Devirtualization!!

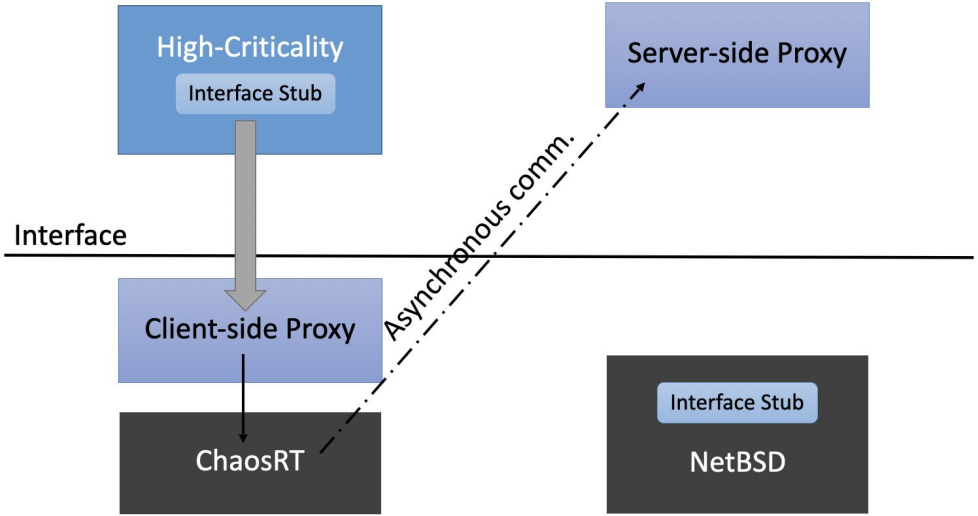


Devirtualization: Get Rid of the Slow VM Layer

- Per Core Virtualization Slow
  - Too many layers of isolation
- Removing it is a win for predictability and performance!
- Works for both low and high criticality code!

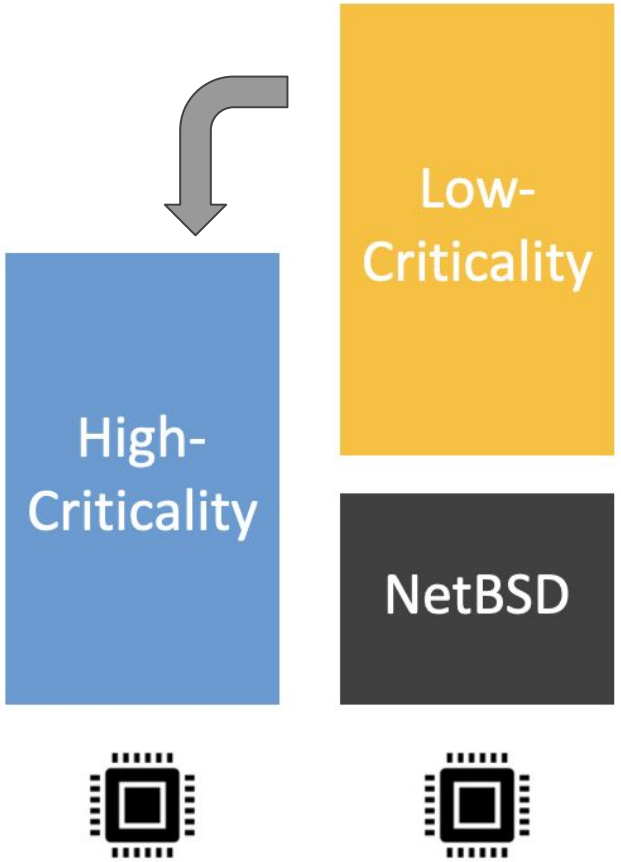
Extension: High Assurance Code uses UNIX APIS

Proxies = mechanism for letting High-Criticality and Low-Criticality processes communicate

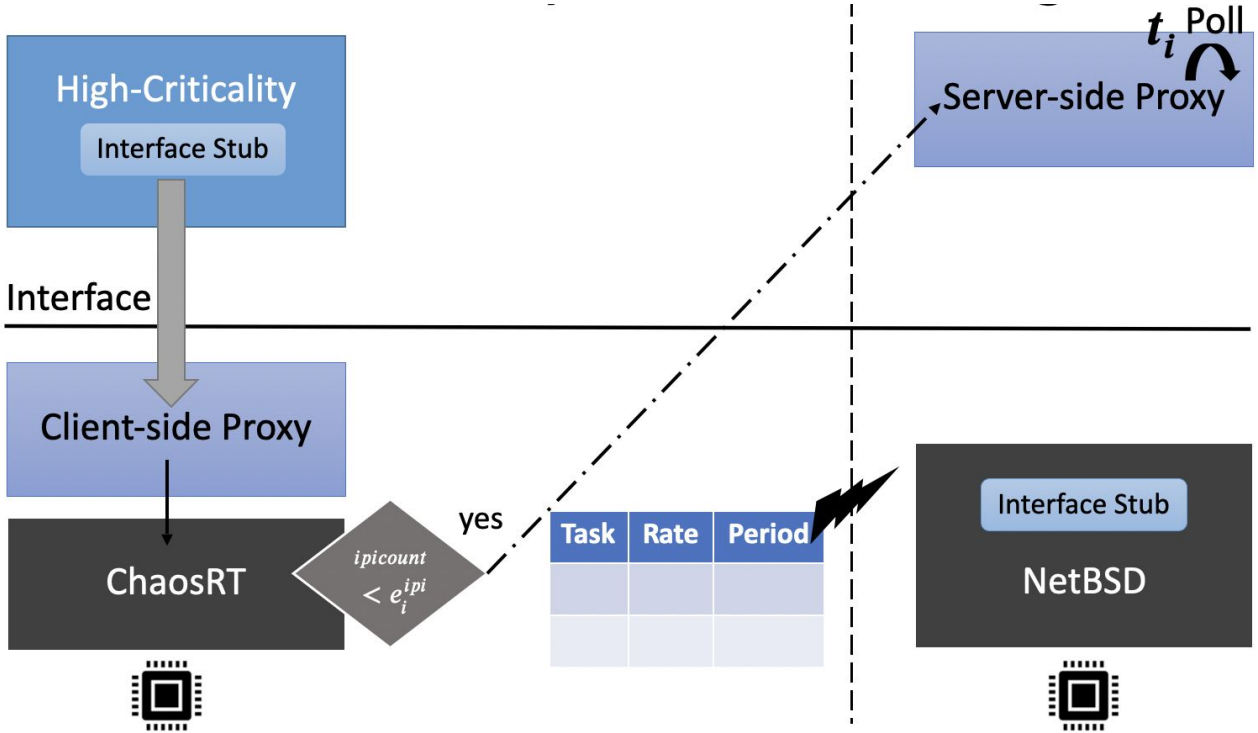


Why Not Talk Directly?

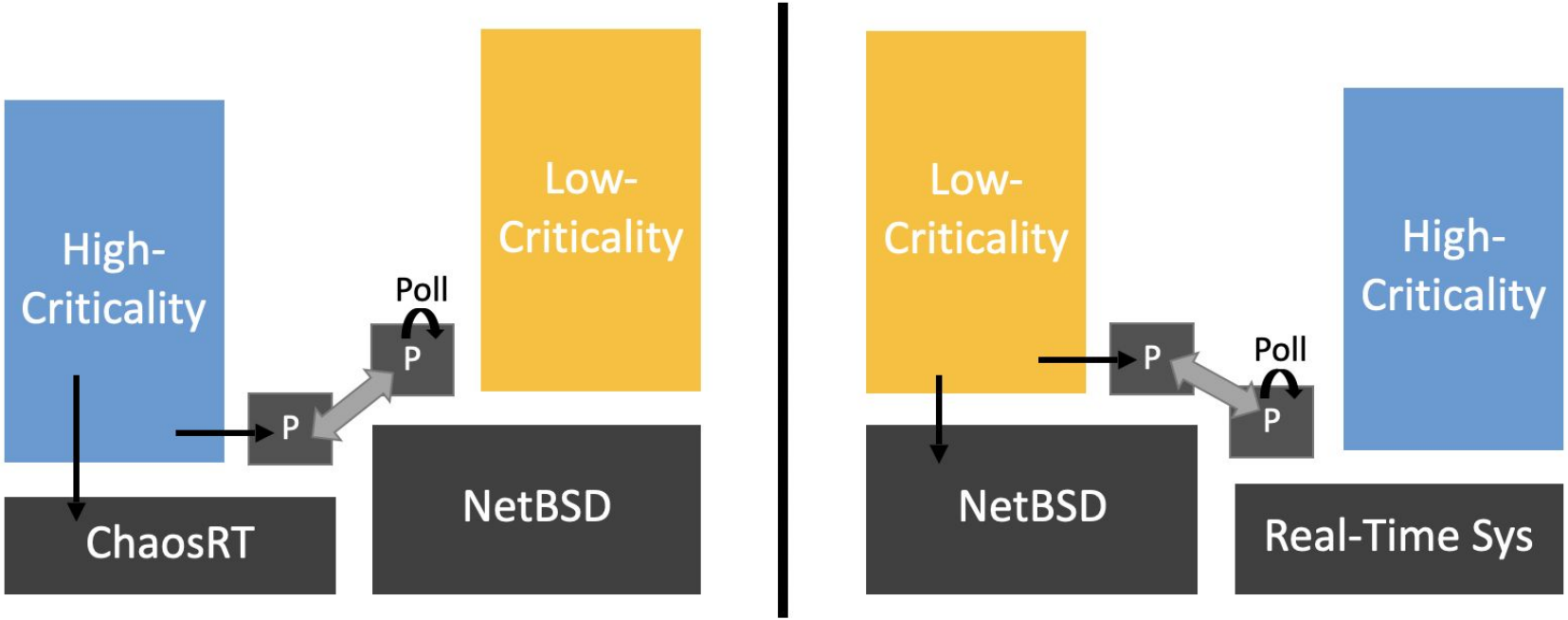
- What might happen here?
- (Recall earlier)



IPI Interference!

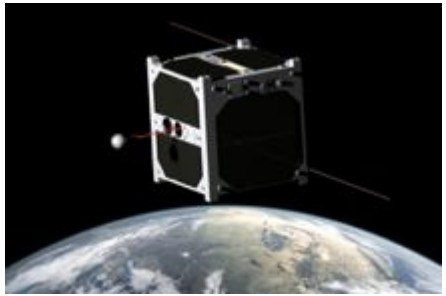


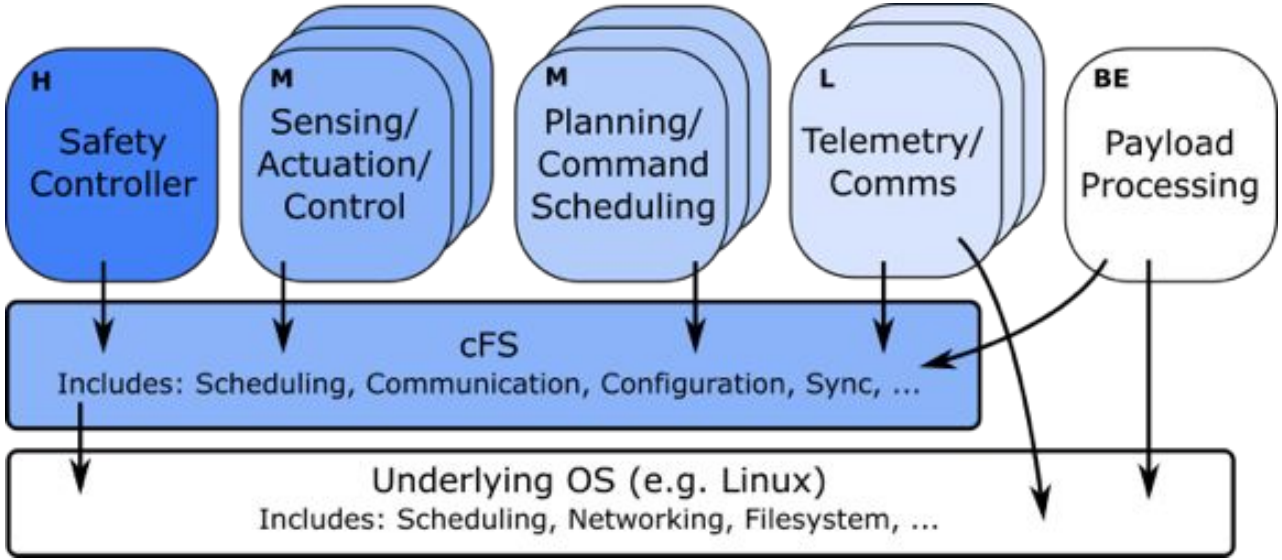
Proxies: Limiting IPI Interference



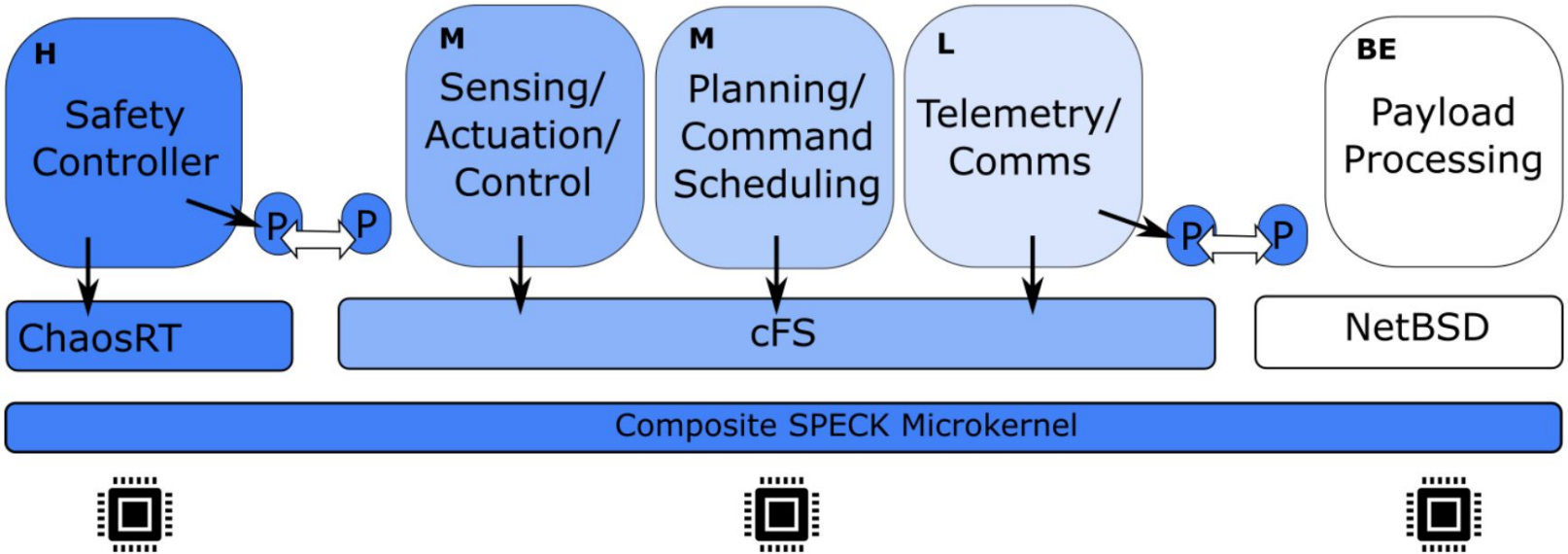
Evaluation

Use Case: The cFE (aka cFS)

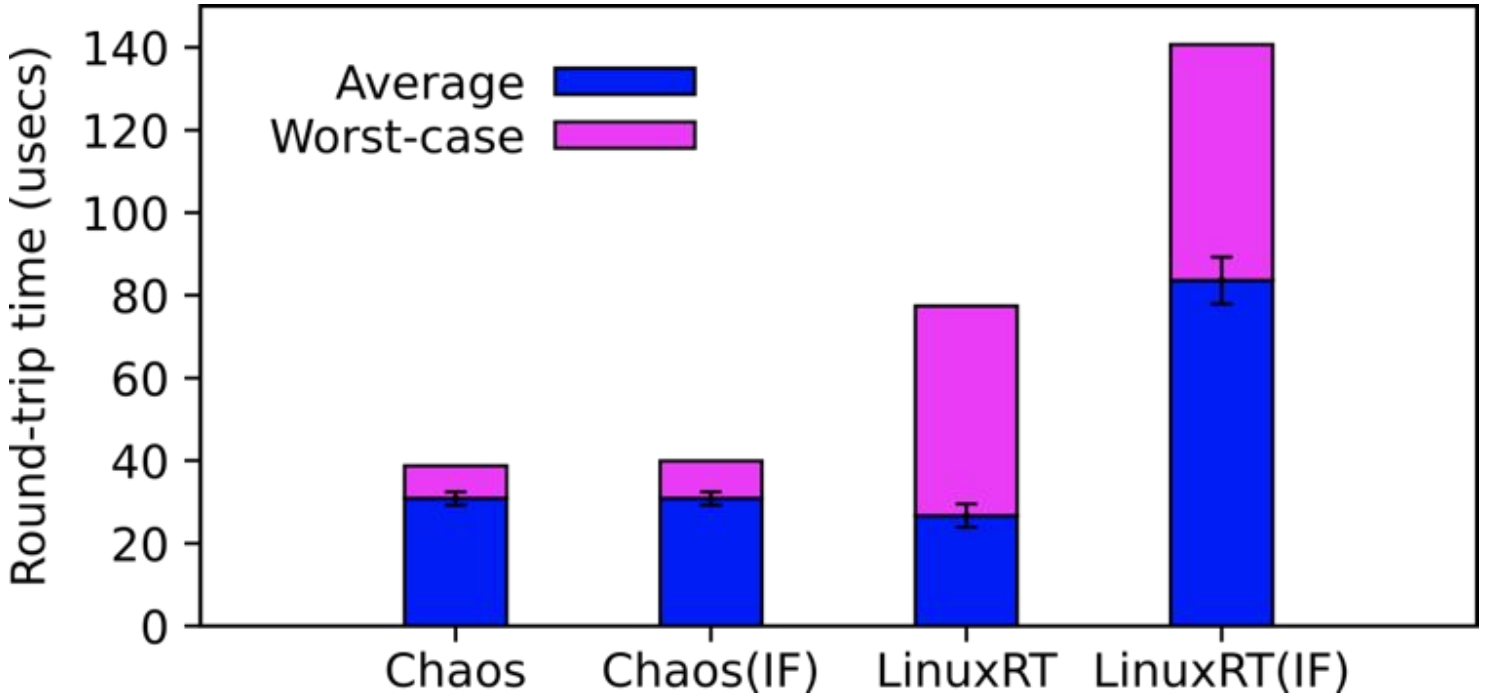




cFE under CHAOS



Worst Case Communication Latency



Summary of Features

Problem	Composite OS	Chaos Solution
Shared-memory interference	Lockless kernel	
IPI interference	Explicit IPIs	Bound IPI interference
Configurable, isolated sched	User-level sched	
High-assurance, predictable execution		Minimal runtime
Inter-assurance-level coordination		Efficient, latency-bounded comm.
Strong inter-assurance isolation	Capability-based	Devirtualization

Critiques

Questions from the Issue

@pcodes: Could this scale to multiple assurance levels (does it even make sense to)?

@anguyen0204: How do TCAPS work?

@rebeccc: IPIs vs shared memory: why not try and resolve the issues of shared memory?

Issues from the Issue

@pcodes: The JpS (jargon per sentence) was too high
    @others: The paper was too dense

@rebeccc: Devirtualization--unclear what this meant in the context of the paper

@Others: The evaluation graphs kinda suck