Turing Machines CSE 211 (Theory of Computation)

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Adapted from slides by Dr. Muhammad Masroor Ali



Programming Techniques for Turing Machines

- A TM is exactly as powerful as a conventional computer.
- We will extend the basic model of the TM by adding features
- None of these tricks extend the basic model of the TM they are only notational conveniences

- We can use the finite control not only to represent a position in the "program" of the Turing machine, but to hold a finite amount of data.
- Figure suggests this technique.
- This also suggests the idea of multiple tracks.

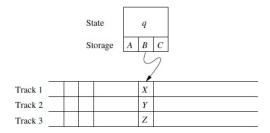
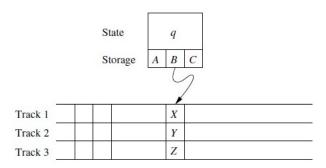
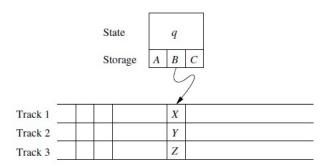


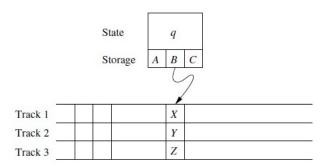
Figure 8.13: A Turing machine viewed as having finite-control storage and multiple tracks



- We see the finite control consisting of not only a "control" state q, but three data elements A, B, and C.
- The technique requires no extension to the TM model.



- We merely think of the state as a tuple.
- In the case of the diagram, we should think of the state as [q, A, B, C].



- Regarding states this way allows us to describe transitions in a more systematic way.
- This often makes the strategy behind the TM program more transparent.

Example

We shall design a TM

$$M = (Q, \{0, 1\}, \{0, 1, B\}, \delta, [q_0, B], \{[q_1, B]\})$$

that remembers in its finite control the first symbol (0 or 1) that it sees.

- And checks that it does not appear elsewhere on its input.
- Thus, M accepts the language $01^* + 10^*$.
- Accepting regular languages such as this one does not stress the ability of Turing machines.
- But it will serve as a simple demonstration.

- The set of states Q is $\{q_0, q_1\} \times \{0, 1, B\}$.
- That is, the states may be thought of as pairs with two components:
 - ① A control portion, q_0 or q_1
 - This remembers what the TM is doing.
 - Control state q₀ indicates that M has not yet read its first symbol.
 - q_1 indicates that it has read the symbol.
 - In q₁ it is checking that it does not appear elsewhere, by moving right and hoping to reach a blank cell.
 - A data portion.
 - This remembers the first symbol seen, which must be 0 or 1.
 - The symbol *B* in this component means that no symbol has been read.

- The set of states *Q* is $\{q_0, q_1\} \times \{0, 1, B\}$.
- That is, the states may be thought of as pairs with two components:
 - **a** A control portion, q_0 or q_1 .
 - This remembers what the TM is doing.
 - Control state q₀ indicates that M has not yet read its first symbol.
 - q₁ indicates that it has read the symbol.
 - In q₁ it is checking that it does not appear elsewhere, by moving right and hoping to reach a blank cell.
 - A data portion.
 - This remembers the first symbol seen, which must be 0 or 1.
 - The symbol B in this component means that no symbol has been read.

The transition function δ of M is as follows:

- **1** $\delta([q_0, B], a) = ([q_1, a], a, R)$ for a = 0 or a = 1.
 - Initially, q₀ is the control state, and the data portion of the state is B.
 - The symbol scanned is copied into the second component of the state.
 - And M moves right, entering control state q₁ as it does so.

The transition function δ of M is as follows:

- 2. $\delta([q_1, a], \overline{a}) = ([q_1, a], \overline{a}, R).$
 - Here \overline{a} is the "complement" of a.
 - That is, 0 if a = 1 and 1 if a = 0.
 - In state q₁ M skips over each symbol 0 or 1 that is different from the one it has stored in its state, and continues moving right.

The transition function δ of M is as follows:

- 3. $\delta([q_1, a], B) = ([q_1, B], B, R)$ for a = 0 or a = 1.
 - If M reaches the first blank, it enters the accepting state $[q_1, B]$.

- **1** $\delta([q_0, B], a) = ([q_1, a], a, R)$ for a = 0 or a = 1.
- $\delta ([q_1, a], \overline{a}) = ([q_1, a], \overline{a}, R).$
- **3** $\delta([q_1, a], B) = ([q_1, B], B, R)$ for a = 0 or a = 1.
 - Notice that M has no definition for δ ($[q_1, a], a$) for a = 0 or a = 1.
 - Thus, if M encounters a second occurrence of the symbol it stored initially in its finite control, it halts without having entered the accepting state.

$$M = (Q, \{0, 1\}, \{0, 1, B\}, \delta, [q_0, B], \{[q_1, B]\})$$

01111

- δ ($[q_0, B], a$) = ($[q_1, a], a, R$) for a = 0 or a = 1
- 2. $\delta([q_1, a], \overline{a}) = ([q_1, a], \overline{a}, R)$
- 3. $\delta([q_1, a], B) = ([q_1, B], B, R)$ for a = 0 or a = 1

[q₀, B] 01111

 $-0[q_1, 0]$

- 01[q_1 , 0]111

 $\vdash 011[q_1, 0]1$

 \vdash 0111[q_1 , 0]1

⊢ 01111[*q*₁, 0]*B*

 $\vdash 011111B[q_1, B]B$

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$[q_0, B]$ 01111 $\vdash 0[q_1, 0]$ 1111 $\vdash 01[q_1, 0]$ 111 $\vdash 011[q_1, 0]$ 11 $\vdash 0111[q_1, 0]$ 1 $\vdash 01111B[q_1, B]$ B



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```
[g_0, B] 01111

\vdash 0[q_1, 0] 1111

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\vdash 011[q_1, 0] 11

\vdash 0111[q_1, 0] 1

\vdash 01111[q_1, 0] B

\vdash 01111B[q_1, B] B
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```
[q_0, B] 01111

\vdash 0[q_1, 0] 1111

\vdash 01[q_1, 0] 111

\vdash 0111[q_1, 0] 11

\vdash 01111[q_1, 0] 1

\vdash 01111[q_1, 0] B

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[q_0, B] 01111

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\vdash 011[q_1, 0] 11

\vdash 0111[q_1, 0] 1

\vdash 01111[q_1, 0] B

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\vdash 011[q_1, 0] 11

\vdash 0111[q_1, 0] 1

\vdash 01111[q_1, 0] B

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\vdash 01111B[q_1, 0]B

\vdash 01111B[q_1, B]B
```

$$\textit{M} = (\textit{Q}, \{0,1\}, \{0,1,B\}, \delta \;, [\textit{q}_0,\textit{B}], \{[\textit{q}_1,\textit{B}]\})$$

We consider a string which does not belong to the language.

01101

- δ ($[q_0, B], a$) = ($[q_1, a], a, R$) for a = 0 or a = 1
- 2. $\delta([q_1, a], \overline{a}) = ([q_1, a], \overline{a}, R)$
- δ $([q_1, a], B) = ([q_1, B], B, R)$ for a = 0 or a = 1

[*q*₀, B] 01101

 $\vdash 0[q_1, 0] 1101$ $\vdash 01[q_1, 0] 101$

us *M* it halts without hav

$$\textit{M} = (\textit{Q}, \{0,1\}, \{0,1,B\}, \delta \;, [\textit{q}_0,\textit{B}], \{[\textit{q}_1,\textit{B}]\})$$

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$$[q_0, B]$$
 01101
 $\vdash 0[q_1, 0]$ 1101
 $\vdash 01[q_1, 0]$ 101
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 01101
 \vdash 0[$q_1, 0$] 1101
 \vdash 01[$q_1, 0$] 101
 \vdash 011[$q_1, 0$] 01

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We consider a string which does not belong to the language.

01101

- $\delta ([q_0, B], a) = ([q_1, a], a, R) \text{ for } a = 0 \text{ or } a = 1$
- $\delta ([q_1, a], \overline{a}) = ([q_1, a], \overline{a}, R)$
- δ $([q_1, a], B) = ([q_1, B], B, R)$ for a = 0 or a = 1

$$[q_0, B]$$
 01101
 \vdash 0[q_1 , 0] 1101
 \vdash 01[q_1 , 0] 101
 \vdash 011[q_1 , 0] 01

$$M = (Q, \{0, 1\}, \{0, 1, B\}, \delta, [q_0, B], \{[q_1, B]\})$$

We consider a string which does not belong to the language.

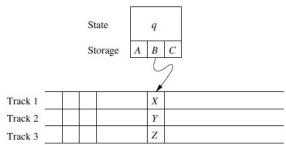
01101

- δ ([q_0 , B], a) = ([q_1 , a], a, R) for a = 0 or a = 1
- 2. $\delta([q_1, a], \overline{a}) = ([q_1, a], \overline{a}, R)$
- δ $([q_1, a], B) = ([q_1, B], B, R)$ for a = 0 or a = 1

$$[q_0, B]$$
 01101
 $\vdash 0[q_1, 0]$ 1101
 $\vdash 01[q_1, 0]$ 101
 $\vdash 011[q_1, 0]$ 01

Multiple Tracks

- Another useful "trick" is to think of the tape of a Turing machine as composed of several tracks.
- Each track can hold one symbol.
- The tape alphabet of the TM consists of tuples, with one component for each "track."
- Thus, for instance, the cell scanned by the tape head in figure contains the symbol [X, Y, Z].



Multiple Tracks

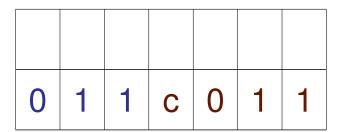
- Like the technique of storage in the finite control, using multiple tracks does not extend what the Turing machine can do.
- It is simply a way to view tape symbols and to imagine that they have a useful structure.

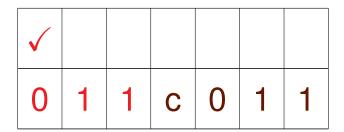
Example

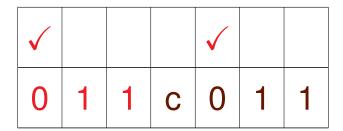
- A common use of multiple tracks is to treat one track as holding the data and a second track as holding a mark.
- We can check off each symbol as we "use" it.
- Or we can keep track of a small number of positions within the data by marking only those positions.
- Previous examples were instances of this technique.
- But in those examples we did not think explicitly of the tape as if it were composed of tracks.

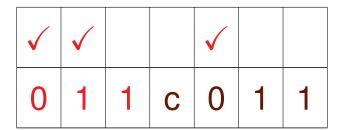
 In the present example, we shall use a second track explicitly to recognize the non-context-free language

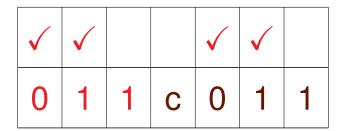
$$L_{wcw} = \{wcw \mid w \text{ is in } (0+1)^+\}$$

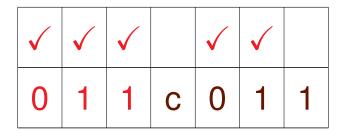


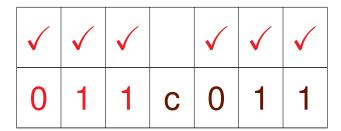












The Turing machine we shall design is:

$$M = (Q, \Sigma, \Gamma, \delta, [q_1, B], [B, B], \{[q_9, B]\})$$

- Q: The set of states is $\{q_1, q_2, ..., q_9\} \times \{0, 1, B\}$.
 - That is, pairs consisting of a control state q_i and a data component: 0, 1, or blank.
 - We again use the technique of storage in the finite control, as we allow the state to remember an input symbol 0 or 1.

The Turing machine we shall design is:

$$M = (Q, \Sigma, \Gamma, \delta, [q_1, B], [B, B], \{[q_9, B]\})$$

TM

- Γ : The set of tape symbols is $\{B,*\} \times \{0,1,c,B\}$.
- The first component, or track, can be either blank or "checked," represented by the symbols B and *, respectively.
- We use the * to check off symbols of the first and second groups of 0's and 1's.
- This eventually confirms that the string to the left of the center marker c is the same as the string to its right.
- The second component of the tape symbol is what we think of as the tape symbol itself.
- That is, we may think of the symbol [B, X] as if it were the tape symbol X, for X = 0, 1, c, B.

• The Turing machine we shall design is:

$$M = (Q, \Sigma, \Gamma, \delta, [q_1, B], [B, B], \{[q_9, B]\})$$

- Σ : The input symbols are [B, 0], [B, 1] and [B, c].
 - This, as mentioned above, we identify with 0, 1 and c, respectively.

• The Turing machine we shall design is:

$$\textit{M} = (\textit{Q}, \Sigma, \Gamma, \delta \;, [\textit{q}_{1}, \textit{B}], [\textit{B}, \textit{B}], \{[\textit{q}_{9}, \textit{B}]\})$$

TM

 δ : The transition function.

- 1. $\delta([q_1, B], [B, a]) = ([q_2, a], [*, a], R).$
 - In the initial state,
 - M picks up the symbol a (which can be either 0 or 1),
 - stores it in its finite control,
 - goes to control state q₂,
 - "checks off" the symbol it just scanned,
 - and moves right.
 - Notice that by changing the first component of the tape symbol from B to *, it performs the check-off.

- 2. $\delta([q_2, a], [B, b]) = ([q_2, a], [B, b], R)$.
 - M moves right, looking for the symbol c.
 - Remember that a and b can each be either 0 or 1, independently, but can not be c.

- 3. $\delta([q_2, a], [B, c]) = ([q_3, a], [B, c], R).$
 - When M finds the c, it continues to move right, but changes to control state q₃.

- 4. $\delta([q_3, a], [*, b]) = ([q_3, a], [*, b], R).$
 - In state q_3 , M continues past all checked symbols.

- 5. $\delta([q_3, a], [B, a]) = ([q_4, B], [*, a], L).$
 - If the first unchecked symbol that M finds is the same as the symbol in its finite control, it checks this symbol.
 - It has matched the corresponding symbol from the first block of 0's and 1's.
 - M goes to control state q₄, dropping the symbol from its finite control, and starts moving left.

- 6. $\delta([q_4, B], [*, a]) = ([q_4, B], [*, a], L).$
 - M moves left over checked symbols.

- 7. $\delta([q_4, B], [B, c]) = ([q_5, B], [B, c], L).$
 - When M encounters the symbol c, it switches to state q₅ and continues left.
 - In state q₅, M must make a decision, depending on whether or not the symbol immediately to the left of the c is checked or unchecked.
 - If checked, then we have already considered the entire first block of 0's and 1's – those to the left of the c.
 - We must make sure that all the 0's and 1's to the right of the c are also checked.



7 (cont).
$$\delta([q_4, B], [B, c]) = ([q_5, B], [B, c], L)$$
.

- We accept if no unchecked symbols remain to the right of the c.
- If the symbol immediately to the left of the c is unchecked,
 - we find the leftmost unchecked symbol,
 - pick it up, and
 - start the cycle that began in state q_1 .

- 8. $\delta([q_5, B], [B, a]) = ([q_6, B], [B, a], L).$
 - This branch covers the case where the symbol to the left of c is unchecked.
 - M goes to state q₆ and continues left, looking for a checked symbol.

- 9. $\delta([q_6, B], [B, a]) = ([q_6, B], [B, a], L).$
 - As long as symbols are unchecked, M remains in state q₆ and proceeds left.

 δ is defined by the following rules, in which a and b each may stand for either 0 or 1.

10.
$$\delta([q_6, B], [*, a]) = ([q_1, B], [*, a], R).$$

 When the checked symbol is found, M enters state q₁ and moves right to pick up the first unchecked symbol.

11.
$$\delta([q_5, B], [*, a]) = ([q_7, B], [*, a], R).$$

- Now, let us pick up the branch from state q_5 .
- In q₅ we have just moved left from the c and find a checked symbol.
- We start moving right again, entering state q_7 .

12.
$$\delta([q_7, B], [B, c]) = ([q_8, B], [B, c], R).$$

- In state q_7 we shall surely see the c.
- We enter state q_8 as we do so, and proceed right.

- 13. $\delta([q_8, B], [*, a]) = ([q_8, B], [*, a], R).$
 - M moves right in state q₈, skipping over any checked 0's or 1's that it finds.

- 14. $\delta([q_8, B], [B, B]) = ([q_9, B], [B, B], R).$
 - If M reaches a blank cell in state q₈ without encountering any unchecked 0 or 1, then M accepts.
 - If M first finds an unchecked 0 or 1, then the blocks before and after the c do not match, and M halts without accepting.

$$\delta ([q_1, B], [B, a]) = ([q_2, a], [*, a], R)$$

 $\delta ([q_2, a], [B, b]) =$

$$\delta ([q_2, a], [B, c]) = ([q_3, a], [B, c], R)$$

$$\delta ([q_4, B], [z, a]) = ([q_4, B], [*, a], z)$$

$$\delta ([q_4, B], [*, a]) = ([q_4, B], [*, a], L)$$

8.
$$\delta([q_5, B], [B, a]) = ([q_6, B], [B, a], L)$$

$$([q_6, B], [B, a], L)$$

$$\delta([q_6, B], [*, a]) - ([q_4, B], [*, a], B)$$

$$0 \delta([q_5, B], [*, a]) = ([q_7, B], [*, a], R$$

Before State: $[q_1, B]$

В	В	В	В	В	В	В
0	1	1	С	0	1	1

 $\it M$ picks up the symbol 0, stores it in its finite control, goes to control state $\it q_2$, "checks off" the symbol just scanned, and moves right.

After State: $[q_2, 0]$

- $\delta ([q_2, a], [B, b]) = ([q_2, a], [B, b], R)$
- $([q_3, a], [B, c], R)$
- 4. $\delta([q_3, a], [*, b]) = ([q_3, a], [*, b], R$
- $\delta ([q_3, a], [B, a]) = ([q_4, B], [*, a], L$
- $\delta ([q_4, B], [*, a]) = ([q_4, B], [*, a], L)$
- δ ([q₄, B], [B, c]) = ([q₅, B], [B, c], L)
- $\delta([q_5, B], [B, a]) = ([q_6, B], [B, a], L)$
- 9. $\delta([q_6, B], [B, a]) = ([q_6, B], [B, a], L)$
- $0 \delta([q_6, B], [*, a]) = ([q_1, B], [*, a], R$

Before State: $[q_2, 0]$

*	В	В	В	В	В	В
0	1	1	O	0	1	1

 ${\it M}$ moves right, looking for the symbol ${\it c}$.

After

State: $[q_2, 0]$

*	В	В	В	В	В	В
0	1	1	С	0	1	1

- $\delta ([q_2, a], [B, b]) = ([q_2, a], [B, b], R)$
- $([q_3, a], [B, c], R)$
- 4. $\delta([q_3, a], [*, b]) = ([q_3, a], [*, b], R$
- $\delta ([q_3, a], [B, a]) = ([q_4, B], [*, a], L$
- $\delta ([q_4, B], [*, a]) = ([q_4, B], [*, a], L)$
- δ ([q₄, B], [B, c]) = ([q₅, B], [B, c], L)
- $\delta([q_5, B], [B, a]) = ([q_6, B], [B, a], L)$
- 9. $\delta([q_6, B], [B, a]) = ([q_6, B], [B, a], L)$
- $0 \delta([q_6, B], [*, a]) = ([q_1, B], [*, a], R$

Before State: $[q_2, 0]$

*	В	В	В	В	В	В
0	1	1	O	0	1	1

 ${\it M}$ moves right, looking for the symbol ${\it c}$.

After

State: $[q_2, 0]$

*	В	В	В	В	В	В
0	1	1	С	0	1	1

$$[.[*,a],R)$$

 $\delta ([q_2, a], [B, b]) = ([q_2, a], [B, b], R)$

$$\delta ([q_2, a], [B, c]) = ([q_3, a], [B, c], R)$$

- 4. $\delta([q_3, a], [*, b]) = ([q_3, a], [*, b], R)$
- 6. $\delta([q_3, a], [B, a]) = ([q_4, B], [*, a], L$
- $\delta ([q_4, B], [*, a]) = ([q_4, B], [*, a], L)$
- $\delta([q_4, B], [B, c]) = ([q_5, B], [B, c], L)$
- 8. $\delta([q_5, B], [B, a]) = ([q_6, B], [B, a], L)$
- $\delta ([q_6, B], [B, a]) = ([q_6, B], [B, a], L)$
 - $0 \quad \delta([q_6, B], [*, a]) = ([q_1, B], [*, a], R)$

011*c*011

Before State: $[q_2, 0]$

*	В	В	В	В	В	В
0	1	1	С	0	1	1

When M finds the c, it continues to move right, but changes to control state q_3

After

State: $[q_3, 0]$

*	В	В	В	В	В	В
0	1	1	С	0	1	1

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$$, [*, a], R)$$

$$\delta\left(\left[a_{2},a\right],\left[B,b\right]\right)=$$

$$([q_2, a], [B, b], R)$$

$$([q_2, a], [B, c])$$

$$\delta ([a_2, a], [*, b]) = ([a_2, a], [*, b], B)$$

5
$$\delta([q_3, a], [B, a]) = ([q_4, B], [*, a], L)$$

$$\delta ([q_4, B], [*, a]) = ([q_4, B], [*, a], L)$$

$$B, c]) =$$

, $c], L)$

$$[B, a]) = B, a], L)$$

$$([q_6, B], [B, a], L)$$

([q₆, B], [B, a]) =

$$\delta ([g_6, B], [B, a], E)$$

$$\delta ([g_6, B], [*, a]) = ([g_1, B], [*, a], R)$$

$$0 \delta([q_5, B], [*, a]) = ([q_7, B], [*, a], R$$

Before State: $[q_3, 0]$

*	В	В	В	В	В	В
0	1	1	С	0	1	1

Matched the corresponding symbol from first block. Checks this. Goes to control state q_4 . Drops the symbol from finite control. Starts moving left

After State: $[q_4, B]$

	*	В	В	В	*	В	В
ſ	0	4	4	0	0	4	4

$$\delta([q_1, B], [B, a]) = ([q_2, a], [*, a], R)$$

$$\delta ([q_2, a], [B, b]) = ([q_2, a], [B, b], R)$$

$$\delta([q_2, a], [B, c]) = ([q_3, a], [B, c], R)$$

4.
$$\delta([q_3, a], [*, b]) = ([q_3, a], [*, b], R)$$

5.
$$\delta([q_3, a], [B, a]) = ([q_4, B], [*, a], L)$$

6
$$\delta([q_4, B], [*, a]) = ([q_4, B], [*, a], L)$$

$$\delta ([q_4, B], [B, c]) = ([q_5, B], [B, c], L)$$

$$\delta ([q_5, B], [B, a]) = ([q_6, B], [B, a], L)$$

9.
$$\delta([q_6, B], [B, a]) = ([q_6, B], [B, a], L)$$

Before State: $[q_4, B]$

*	В	В	В	*	В	В
0	1	1	C	0	1	1

When M encounters the symbol c, it switches to state q_5 and continues left

After

State: $[q_5, B]$

*	В	В	В	*	В	В
0	1	1	С	0	1	1

$$\delta\left(\left[a_{0},a\right]\left[B,b\right]\right)=$$

$$([q_2, a], [B, b], R)$$

$$\delta ([q_3, a], [*, b]) = ([q_3, a], [*, b], R)$$

$$\delta ([q_3, a], [B, a]) = ([q_4, B], [*, a], L$$

$$\delta ([q_4, B], [*, a]) = ([q_4, B], [*, a], L)$$

$$, [B, c]) = B, c], L)$$

$$\delta([q_5, B], [B, a]) = ([q_6, B], [B, a], L)$$

$$\delta ([q_6, B], [B, a], E)$$

$$\delta ([q_6, B], [B, a]) =$$

$$\delta ([q_6, B], [*, a]) = ([q_1, B], [*, a], R)$$

Before State: [q₅, B]

*	В	В	В	*	В	В
0	1	1	С	0	1	1

This branch covers the case where the symbol to the left of c is unchecked. M goes to state q_6 and continues left, looking for a checked symbol

After State: $[a_6, B]$

* B B B * B B

0 11 14 C 0 0 1 1 1

$$\delta([q_1, B], [B, a]) = ([q_2, a], [*, a], R)$$

$$\delta ([q_2, a], [B, b]) = ([q_2, a], [B, b], R)$$

([
$$q_3, a$$
], [B, c]) = ([q_3, a], [B, c], R)

4.
$$\delta([q_3, a], [*, b]) = ([q_3, a], [*, b], R)$$

$$\delta ([q_3, a], [B, a]) = ([q_4, B], [*, a], L)$$

6.
$$\delta([q_4, B], [*, a]) = ([q_4, B], [*, a], L)$$

$$\delta([q_4, B], [B, c]) = ([q_5, B], [B, c], L)$$

8.
$$\delta([q_5, B], [B, a]) = ([q_6, B], [B, a], L)$$

$$\delta ([q_6, B], [B, a]) = ([q_6, B], [B, a], L)$$

Before State: [q_6 , B]

*	В	В	В	*	В	В
0	1	1	С	0	1	1

As long as symbols are unchecked, M remains in state q_6 and proceeds left

After

State: $[q_6, B]$

*	В	В	В	*	В	В
0	1	1	С	0	1	1



$$[q_1, B], [B, a]) = ([q_2, a], [*, a], R)$$

$$\delta ([q_2, a], [B, b]) = ([q_2, a], [B, b], R)$$

$$\delta([q_2, a], [B, c]) = ([q_3, a], [B, c], R)$$

$$\delta$$
 ([q₃, a], [*, b]) = ([q₃, a], [*, b], R)

6.
$$\delta([q_3, a], [B, a]) = ([q_4, B], [*, a], L$$

6
$$\delta([q_4, B], [*, a]) = ([q_4, B], [*, a], L)$$

$$\delta$$
 ([q_4 , B], [B , c]) = ([q_5 , B], [B , c], L)

$$\delta ([q_5, B], [B, a]) = ([q_6, B], [B, a], L)$$

$$\delta([q_6, B], [B, a]) = ([q_6, B], [B, a], L)$$

$$\delta ([q_5, B], [*, a]) = ([q_7, B], [*, a], R)$$

 $\delta([q_6, B], [*, a]) = ([q_1, B], [*, a], R)$

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Before State: [q_6 , B]

*	В	В	В	*	В	В
0	1	1	C	0	1	1

When the checked symbol is found, M enters state q_1 and moves right to pick up the first unchecked symbol

After

State: $[q_1, B]$

*	В	В	В	*	В	В
0	1	1	С	0	1	1

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$$\delta([q_1, B], [B, a]) = ([q_2, a], [*, a], R)$$

$$\delta ([q_2, a], [B, b], H)$$

$$\delta ([q_2, a], [B, c]) =$$

$$([q_3, a], [B, c], R)$$

$$\delta ([q_3, a], [p, a]) = ([q_4, B], [*, a], L$$

6
$$\delta([q_4, B], [*, a]) = ([q_4, B], [*, a], L)$$

7 $\delta([q_4, B], [B, c]) =$

$$([q_5, B], [B, c], L)$$

$$([q_6, B], [B, a], L)$$
① $\delta([q_6, B], [B, a]) =$

$$([q_6, B], [B, a], L)$$

$$\delta([q_6, B], [*, a]) = ([q_1, B], [*, a], R)$$

$$\delta([q_5, B], [*, a]) = ([q_7, B], [*, a], F_6)$$

Before State: $[q_1, B]$

*	В	В	В	*	В	В
0	1	1	C	0	1	1

M picks up the symbol 1, stores it in its finite control, goes to control state q_2 , "checks off" the symbol it just scanned, and moves right

After State: $[q_2, 1]$

- $\delta ([q_2, a], [B, b]) = ([q_2, a], [B, b], R)$
- $\delta ([q_2, a], [B, c]) = ([q_3, a], [B, c], R)$
- δ ([q₃, a], [*, b]) = ([q₃, a], [*, b], R)
- $\delta ([q_3, a], [B, a]) = ([q_4, B], [*, a], L$
- δ ([q_4 , B], [B, c]) = ([q_5 , B], [B, c], L)
- 8. $\delta([q_5, B], [B, a]) = ([q_6, B], [B, a], L)$
- 9. $\delta([q_6, B], [B, a]) = ([q_6, B], [B, a], L)$
- $0 \delta([q_5, B], [*, a]) = ([q_7, B], [*, a], R)$

Before State: $[q_2, 1]$

*	*	В	В	*	В	В
0	1	1	С	0	1	1

 $\it M$ moves right, looking for the symbol $\it c$

After

State: [*q*₂, 1]

*	*	В	В	*	В	В
0	1	1	С	0	1	1

$$[, [*, a], R)$$

$$\delta ([q_2, a], [B, b]) = ([q_2, a], [B, b], R)$$

$$\delta ([q_2, a], [B, c]) = ([q_3, a], [B, c], R)$$

4.
$$\delta([q_3, a], [*, b]) = ([q_3, a], [*, b], R)$$

$$\delta ([q_3, a], [B, a]) = ([q_4, B], [*, a], L$$

$$\delta ([q_4, B], [*, a]) = ([q_4, B], [*, a], L)$$

$$\delta$$
 ([q_4 , B], [B , c]) = ([q_5 , B], [B , c], L)

$$\delta ([q_5, B], [B, a]) = ([q_6, B], [B, a], L)$$

$$([q_6, B], [B, a], L)$$

$$0 \delta([q_6, B], [*, a]) = ([q_1, B], [*, a], R)$$

Before State: $[q_2, 1]$

*	*	В	В	*	В	В
0	1	1	С	0	1	1

When M finds the c, it continues to move right, but changes to control state q_3

After

State: [*q*₃, 1]

*	*	В	В	*	В	В
0	1	1	С	0	1	1

- $\delta ([q_2, a], [B, b]) = ([q_2, a], [B, b], R)$
- $\delta([q_2, a], [B, c]) = ([q_3, a], [B, c], R)$
- **4** $\delta([q_3, a], [*, b]) = ([q_3, a], [*, b], R)$
- $\delta ([q_3, a], [B, a]) = ([q_4, B], [*, a], L$
- $\delta ([q_4, B], [*, a]) = ([q_4, B], [*, a], L)$
- $\delta([q_4, B], [B, c]) = ([q_5, B], [B, c], L)$
- $\delta ([q_5, B], [B, a]) = ([q_6, B], [B, a], L)$
- 9. $\delta([q_6, B], [B, a]) = ([q_6, B], [B, a], L)$

Before State: $[q_3, 1]$

*	*	В	В	*	В	В
0	1	1	C	0	1	1

M continues past all checked symbols

After

State: $[q_3, 1]$

*	*	В	В	*	В	В
0	1	1	С	0	1	1

 $\delta ([q_1, B], [B, a]) = ([q_2, a], [*, a], R)$ $\delta ([q_2, a], [B, b]) =$

3.
$$\delta([q_2, a], [B, c])$$

 $([q_3, a], [B, c], R)$

$$(C_1, C_2)$$

$$[q_3, a], [B, a]) = ([q_4, B], [*, a], L)$$

$$\delta ([q_4, B], [*, a]) = ([q_4, B], [*, a], L)$$

$$\delta ([q_4, B], [B, c]) =$$

8.
$$\delta([q_5, B], [B, a]) = ([q_6, B], [B, a], L)$$

$$([q_6, B], [B, a], L)$$

$$\delta([q_6, B], [B, a], L) = ([q_6, B], [q_6, B], [q_6$$

①
$$\delta([q_5, B], [*, a]) = ([q_7, B], [*, a], F)$$

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Before State: $[q_3, 1]$

*	*	В	В	*	В	В
0	1	1	С	0	1	1

Matched the corresponding symbol from first block. Checks this. Goes to control state q_4 . Drops the symbol from finite control. Starts moving left

After State: $[q_4, B]$

- $\delta ([q_2, a], [B, b]) = ([q_2, a], [B, b], R)$
- $([q_2, a], [B, c]) = ([q_3, a], [B, c], R)$
- 4. $\delta([q_3, a], [*, b]) = ([q_3, a], [*, b], R)$
- 6. $\delta([q_3, a], [B, a]) = ([q_4, B], [*, a], L$
- 6 $\delta([q_4, B], [*, a]) = ([q_4, B], [*, a], L)$
- $\delta ([q_4, B], [B, c]) = ([q_5, B], [B, c], L)$
- 8. $\delta([q_5, B], [B, a]) = ([q_6, B], [B, a], L)$
- 9. $\delta([q_6, B], [B, a]) = ([q_6, B], [B, a], L)$

Before State: $[q_4, B]$

*	*	В	В	*	*	В
0	1	1	C	0	1	1

M moves left over checked symbols

After

State: $[q_4, B]$

*	*	В	В	*	*	В
0	1	1	С	0	1	1

$$=([q_2,a],[*,a],R)$$

$$\delta ([q_2, a], [B, b]) = ([q_2, a], [B, b], R)$$

$$0 ([q_2, a], [B, c]) = ([q_3, a], [B, c], R)$$

4.
$$\delta([q_3, a], [*, b]) = ([q_3, a], [*, b], R)$$

5.
$$\delta([q_3, a], [B, a]) = ([q_4, B], [*, a], L)$$

6
$$\delta([q_4, B], [*, a]) = ([q_4, B], [*, a], L)$$

$$\delta ([q_4, B], [B, c]) = ([q_5, B], [B, c], L)$$

$$b \delta([q_5, B], [B, a]) = ([q_6, B], [B, a], L)$$

$$\delta ([q_6, B], [B, a]) = ([q_6, B], [B, a], L)$$

$$0 \delta([q_6, B], [*, a]) = ([q_1, B], [*, a], R)$$

Before State: [q₄, B]

*	*	В	В	*	*	В
0	1	1	С	0	1	1

Encounters the symbol c, switches to state q_5 and continues left

After

State: $[q_5, B]$

*	*	В	В	*	*	В
0	1	1	С	0	1	1

- , [*, a], R)
- $\delta([q_1, 2], [2, 4]) ([q_2, 4], [1, 4], 11)$
- $\delta ([q_2, a], [B, c]) =$
- $([q_3, a], [B, c], R)$
- $\delta ([a_0, b_1, [a_1, b_2]) ([a_1, b_1], [a_1, b_2]) = ([a_1, b_1], [a_1, b_2])$
 - $0 ([q_4, D], [*, d]) = ([q_4, D], [*, d], L)$
- $([q_5, B], [B, c], L)$
- $([q_6, B], [B, a], L)$
- $([q_6, B], [B, a], L)$

Before State: $[q_5, B]$

*	*	В	В	*	*	В
0	1	1	C	0	1	1

.....

After

State: [q₂, 1]

*	*	В	В	*	*	В
0	1	1	С	0	1	1

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[*, a], R)

c], R)

 $\delta ([q_0, q_1, [z, q_1)] - ([q_1, p_1, [z, q_1, z]))$

[B, c]) =

[B,c],L)

 $([q_6, B], [B, a], L)$

 $([q_6, B], [B, a], L)$ $\delta([q_6, B], [*, a]) = ([q_1, B], [*, a], R)$

1 $\delta([q_5, B], [*, a]) = ([q_7, B], [*, a], F$

011*c*011

Before State: $[q_2, 1]$

*	*	*	В	*	*	В
0	1	1	С	0	1	1

.

After State: $[q_3, 1]$

* * * B * B 0 1 1 c 0 1 1

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 $\delta\left(\left[\alpha_{2},a\right],\left[B,b\right]\right)=$

$$([q_2, a], [B, b], R)$$

$$0 ([q_2, a], [B, c])$$

 $([q_3, a], [B, c], R)$

$$\delta([a, a], [a, b]) - ([a, a] + b] R$$

$$\delta ([q_3, a], [*, b]) = ([q_3, a], [*, b], h)$$

$$\delta ([q_3, a], [B, a]) = ([q_4, B], [*, a], L)$$

$$\delta ([q_4, B], [*, a]) = ([q_4, B], [*, a], L)$$

$$[B,c]) = (c,c,L)$$

$$([q_6, B], [B, a], L)$$

$$([q_6, B], [B, a], L)$$

$$\emptyset \quad \delta([q_6, B] [* a]) = ([q_4, B] [* a], B)$$

$$0 \delta([q_5, B], [*, a]) = ([q_7, B], [*, a], F$$

011*c*011

Before State: $[q_3, 1]$

*	*	*	В	*	*	В
0	1	1	С	0	1	1

Matched the corresponding symbol from first block. Checks this. Goes to control state q_4 . Drops the symbol from finite control. Starts moving left

After State: $[q_4, B]$

*	*	*	В	*	*	*
0	.1.	1	G.	O=	. 1 ₌	1, .

[*, a], R)

 $\delta ([q_2, a], [B, b]) =$

 $([q_2, a], [B, b], R)$

 $([q_3, a], [B, c], R)$

6 $\delta([q_3, a], [*, b]) = ([q_3, a], [*, b], R$

 $\delta ([q_4, B], [*, a]) = ([q_4, B], [*, a], L)$

 $\delta([q_4, B], [B, c]) = ([q_5, B], [B, c], I)$

([q_6, B], [B, a], B) =

 $([q_6, B], [B, a], L)$ $\delta([q_6, B], [*, a]) = ([q_1, B], [*, a], R)$

 $\delta ([q_5, B], [*, a]) = ([q_7, B], [*, a], F_6)$

011*c*011

Before State: $[q_4, B]$

* * * B * * * 0 1 1 c 0 1 1

.....

After State: $[q_4, B]$

*	*	*	В	*	*	*
0	1	1	С	0	1	1

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- $\delta ([q_2, a], [B, b]) = ([q_2, a], [B, b], R)$
- $([q_2, a], [B, c]) = ([q_3, a], [B, c], R)$
- 4 $\delta([q_3, a], [*, b]) = ([q_3, a], [*, b], R$
- 6 $\delta([q_3, a], [B, a]) = ([q_4, B], [*, a], L$
- $\delta ([q_4, B], [*, a]) = ([q_4, B], [*, a], L)$
- $\delta ([q_4, B], [B, c]) = ([q_5, B], [B, c], L)$
- $\delta([q_5, B], [B, a]) = ([q_6, B], [B, a], L)$
- 9. $\delta([q_6, B], [B, a]) = ([q_6, B], [B, a], L)$
- $\emptyset \quad \delta([q_5, B], [*, a]) = ([q_7, B], [*, a], R)$

Before State: [q₄, B]

	*	*	*	В	*	*	*
Ì	0	1	1	С	0	1	1

Encounters c, switches to q_5 and continues left

After

State: $[q_5, B]$

*	*	*	В	*	*	*
0	1	1	С	0	1	1

$$\delta([q_1, B], [B, a]) = ([q_2, a], [*, a], R)$$

2
$$\delta([q_2, a], [B, b]) = ([q_2, a], [B, b], R)$$

$$\delta([q_2, a], [B, c]) = ([q_3, a], [B, c], R)$$

$$\delta$$
 ([q_3 , a], [*, b]) = ([q_3 , a], [*, b], R)

$$\delta ([q_3, a], [B, a]) = ([q_4, B], [*, a], L)$$

$$\delta ([q_4, B], [*, a]) = ([q_4, B], [*, a], L)$$

$$\delta ([q_4, B], [B, c]) = ([q_5, B], [B, c], L)$$

$$\delta ([q_5, B], [B, a]) = ([q_6, B], [B, a], L)$$

①
$$\delta([q_6, B], [B, a]) = ([q_6, B], [B, a], L)$$

Before State: [q₅, B]

*	*	*	В	*	*	*
0	1	1	С	0	1	1

Just moved left from c and find a checked symbol. Have already considered the entire block of 0's and 1's to the left of the c. Must make sure that all the 0's and 1's to the right of c are also checked. Start moving right again, entering state q_7

After

State: [*q*₇, *B*]



- $\delta ([q_2, a], [B, b]) = ([q_2, a], [B, b], B)$
- ([q_2, a], [B, c]) =
- 4. $\delta([q_3, a], [*, b]) = ([q_3, a], [*, b], R)$
- $\delta ([q_3, a], [B, a]) = ([q_4, B], [*, a], L$
- 6. $\delta([q_4, B], [*, a]) = ([q_4, B], [*, a], L)$
- δ ([q₄, B], [B, c]) = ([q₅, B], [B, c], L)
- 8. $\delta([q_5, B], [B, a]) = ([q_6, B], [B, a], L)$
- $\delta ([q_6, B], [B, a]) = ([q_6, B], [B, a], L)$

 Ω $S([\alpha D][D \alpha])$

- $\delta ([q_5, B], [*, a]) = ([q_7, B], [*, a], R)$

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Before State: $[q_7, B]$

*	*	*	В	*	*	*
0	1	1	С	0	1	1

In state q_7 we shall surely see the c. We enter state q_8 as we do so, and proceed right

After State: $[q_8, B]$

*	*	*	В	*	*	*
0	1	1	С	0	1	1

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- $\delta ([q_2, a], [B, b]) = ([q_2, a], [B, b], B)$
- $([q_2, a], [B, c]) = ([q_3, a], [B, c], R)$
- $\delta([q_3, a], [*, b]) = ([q_3, a], [*, b], R)$
- $\delta ([q_3, a], [B, a]) = ([q_4, B], [*, a], L)$
- 6. $\delta([q_4, B], [*, a]) = ([q_4, B], [*, a], L)$
- δ ([q_4 , B], [B, c]) = ([q_5 , B], [B, c], L)
- $\delta([q_5, B], [B, a]) = ([q_6, B], [B, a], L)$
- 9. $\delta([q_6, B], [B, a]) = ([q_6, B], [B, a], L)$
- $0 \delta([q_6, B], [*, a]) = ([q_1, B], [*, a], R$

Before State: [q₈, B]

	*	*	*	В	*	*	*
Ì	0	1	1	С	0	1	1

M moves right in state q_8 , skipping over any checked 0's or 1's that it finds

After

State: $[q_8, B]$

*	*	*	В	*	*	*
0	1	1	С	0	1	1

- $\delta ([q_2, a], [B, b]) = ([q_2, a], [B, b], B)$
- $([q_2, a], [B, c]) = ([q_3, a], [B, c], R)$
- $\delta([q_3, a], [*, b]) = ([q_3, a], [*, b], R)$
- $\delta ([q_3, a], [B, a]) = ([q_4, B], [*, a], L)$
- 6. $\delta([q_4, B], [*, a]) = ([q_4, B], [*, a], L)$
- δ ([q_4 , B], [B, c]) = ([q_5 , B], [B, c], L)
- $\delta([q_5, B], [B, a]) = ([q_6, B], [B, a], L)$
- 9. $\delta([q_6, B], [B, a]) = ([q_6, B], [B, a], L)$
- $0 \delta([q_6, B], [*, a]) = ([q_1, B], [*, a], R$

Before State: [q₈, B]

	*	*	*	В	*	*	*
Ì	0	1	1	С	0	1	1

M moves right in state q_8 , skipping over any checked 0's or 1's that it finds

After

State: $[q_8, B]$

*	*	*	В	*	*	*
0	1	1	С	0	1	1

$$\delta\left([a_0,a],[B,b]\right) =$$

$$([q_2, a], [B, b], R)$$

$$([q_3, a], [B, c], R)$$

4.
$$\delta([q_3, a], [*, b]) = ([q_3, a], [*, b], R$$

$$\delta ([q_3, a], [B, a]) = ([q_4, B], [*, a], L)$$

$$\delta ([q_4, B], [*, a]) = ([q_4, B], [*, a], L)$$

$$[B,C]) = B,C],L)$$

$$([q_6, B], [B, a], L)$$

$$([q_6, B], [B, a], L)$$

Before State: $[q_8, B]$

*	*	*	В	*	*	*
0	1	1	С	0	1	1

M moves right in state q_8 , skipping over any checked 0's or 1's that it finds

After State: $[q_8, B]$

*	*	*	В	*	*	*	В
0	1	1	С	0	1	1	В

$$\delta\left([q_2,a],[B,b]\right) = \\ ([q_2,a],[B,b],R)$$

$$\delta\left([q_2,a],[B,c]\right) = \\ \delta\left([q_2,a],[B,c]\right) = \\ \text{unchecked 0 or 1, it accepts}$$

$$\delta ([q_4, B], [B, c]) = ([q_5, B], [B, c], L)$$

$$\delta ([q_5, B], [B, a]) = ([q_6, B], [B, a], L)$$

$$\delta ([q_6, B], [B, a]) = ([q_6, B], [B, a], L)$$

$$\delta ([q_6, B], [B, a], L)$$

M reaches a blank cell in state q₈ without encountering any

Example

 Let's design a TM to add two binary numbers. The tape initially contains

$$W_1 C W_2 C$$

- w₁ and w₂ are two binary number (assumed to be of equal length)
- The leftmost bits are LSBs for convenience
- The *c*'s are separators
- We will write the sum to the right of the second c

1	0	1	С	0	1	1	С	В	В	В	В

\checkmark											
1	0	1	С	0	1	1	С	В	В	В	В

\checkmark				√							
1	0	1	С	0	1	1	С	1	В	В	В

	\checkmark			√								
1	0	1	С	0	1	1	С	1	В	В	В	

\checkmark	*			•	√						
1	0	1	С	0	1	1	С	1	1	В	В

		-		*							
1	0	1	С	0	1	1	С	1	1	В	В

					√						
1	0	1	С	0	1	1	С	1	1	0	В

√											
1	0	1	С	0	1	1	С	1	1	0	1

- We will store two bits and the carry in the state
- States will be 4-tuples of the form [q, a, b, c_i]
- We start at [q₀, B, B, 0]
- After reading the bit from the first number, we switch to q_1 , store the bit in the state and go right
- We switch to q₂ after seeing the first c
- Once we find an unchecked bit, we switch to q₃ and store the bit in the state
- We keep going right until we see a blank. We then write the sum of the bits and the carry stored in the state



- We then switch to q₄ and keep going left
- We switch to q₅ when see the first c and then to q₆ when see the second c
- When we find the check mark we go to q₀ and move right
- Then either the cycle repeats
- Or if we see a c while at q_0 , we are done with the numbers
- In that case we switch to q₇ and keep going right until we find a blank
- We then write the carry and switch to q₈ and halt

$$[q_0,B,B,0]$$

1	0	1	С	0	1	1	С	В	В	В	В

$$[q_1, 1, B, 0]$$

\checkmark											
1	0	1	С	0	1	1	С	В	В	В	В

$$[q_2, 1, B, 0]$$

\checkmark											
1	0	1	С	0	1	1	С	В	В	В	В

$$[q_3, 1, 0, 0]$$

\checkmark				\checkmark							
1	0	1	С	0	1	1	С	В	В	В	В

$$[q_4, B, B, 0]$$

\checkmark				√							
1	0	1	С	0	1	1	С	1	В	В	В

$$[q_5, B, B, 0]$$

\checkmark				√							
1	0	1	С	0	1	1	С	1	В	В	В

$$[q_6, B, B, 0]$$

\checkmark				√							
1	0	1	С	0	1	1	С	1	В	В	В

$$[q_0, B, B, 0]$$

V					√							
	1	0	1	С	0	1	1	С	1	В	В	В

$$[q_0, 0, B, 0]$$

\checkmark	√			√							
1	0	1	С	0	1	1	С	1	В	В	В

$$[q_3, 0, 1, 0]$$

√	•	✓			√	√						
1	(0	1	С	0	1	1	С	1	В	В	В

$$[q_4, B, B, 0]$$

	√			√	*						
1	0	1	С	0	1	1	С	1	1	В	В

$$[q_1, 1, B, 0]$$

1	0	1	С	0	1	1	С	1	1	В	В

$$[q_3, 1, 1, 0]$$

\checkmark	√	√		√	√	√					
1	0	1	С	0	1	1	С	1	1	В	В

$$[q_4, B, B, 1]$$

\checkmark	\checkmark	\checkmark		\checkmark	√	√					
1	0	1	С	0	1	1	С	1	1	0	В

$$[q_0, B, B, 1]$$

1	0	1	С	0	1	1	С	1	1	0	В

$$[q_7, B, B, 1]$$

		√									
1	0	1	С	0	1	1	С	1	1	0	В

$$[q_8, B, B, B]$$

\checkmark	√	√		√	√	√					
1	0	1	С	0	1	1	С	1	1	0	1

The Turing machine we shall design is:

$$M = (Q, \Sigma, \Gamma, \delta, [q_0, B, B, 0], [B, B], \{[q_7, B, B, B]\})$$

where:

- Q: The set of states is $\{q_0, q_2, ..., q_7\} \times \{0, 1, B\} \times \{0, 1, B\} \times \{0, 1, B\}.$
 - That is, pairs consisting of a control state q_i and three data components: 0, 1, or blank
 - The data components store two bits from two numbers and a carry.
 - We again use the technique of storage in the finite control, as we allow the state to remember an input symbol 0 or 1.

TM

Example-continued

• The Turing machine we shall design is:

$$M = (Q, \Sigma, \Gamma, \delta, [q_0, B, B, 0], [B, B], \{[q_7, B, B, B]\})$$

where:

- Γ : The set of tape symbols is $\{B,*\} \times \{0,1,c,B\}$.
- The first component, or track, can be either blank or "checked," represented by the symbols B and *, respectively.
- We use the * to check off symbols of the first and second groups of 0's and 1's.
- The second component of the tape symbol is what we think of as the tape symbol itself.
- That is, we may think of the symbol [B, X] as if it were the tape symbol X, for X = 0, 1, c, B.

The Turing machine we shall design is:

$$M = (Q, \Sigma, \Gamma, \delta, [q_0, B, B, 0], [B, B], \{[q_7, B, B, B]\})$$

where:

- Σ : The input symbols are [B, 0], [B, 1] and [B, c].
 - This, as mentioned above, we identify with 0, 1 and c, respectively.

• The Turing machine we shall design is:

$$M = (Q, \Sigma, \Gamma, \delta, [q_0, B, B, 0], [B, B], \{[q_7, B, B, B]\})$$

TM

where:

 δ : The transition function.

- 1. $\delta([q_0, B, B, c_i], [B, a]) = ([q_1, a, B, c_i], [*, a], R).$
 - In the initial state,
 - M picks up the symbol a (which can be either 0 or 1),
 - stores it in its finite control,
 - goes to control state q₁,
 - "checks off" the symbol it just scanned,
 - and moves right.
 - Notice that by changing the first component of the tape symbol from B to *, it performs the check-off.

- 2. $\delta([q_1, a, B, c_i], [B, 0/1]) = ([q_1, a, B, c_i], [B, 0/1], R).$
 - M moves right, looking for the symbol c.
 - Remember that a and b can each be either 0 or 1, independently, but can not be c.

- 3. $\delta([q_1, a, B, c_i], [B, c]) = ([q_2, a, B, c_i], [B, c], R).$
 - When M finds the c, it continues to move right, but changes to control state q₂.

- 4. $\delta([q_2, a, B, c_i], [*, 0/1]) = ([q_2, a, B, c_i], [*, 0/1], R).$
 - In state q_2 , M continues past all checked symbols.

- 5. $\delta([q_2, a, B, c_i], [B, b]) = ([q_3, a, b, c_i], [*, b], R).$
 - When it sees the first unchecked symbol, it saves the bit in the state.
 - M goes to control state q₃ and keeps moving right.

- 6. $\delta([q_3, a, b, c_i], [B, 0/1/c]) = ([q_3, a, b, c_i], [B, 0/1/c], R).$
 - It moves right over all symbols until it finds a blank

- 7. $\delta([q_3, a, b, c_i], [B, B]) = ([q_4, B, B, c_o], [B, s], L).$
 - When it finds a blank, it writes the sum s to the tape and saves the carry out co in the state
 - s and c_o are given by

а	b	C_i	S	c_0
0	0	0	0	0
0	0	1	1	0
0	1	0	1	0
0	1	1	0	1
1	0	0	1	0
1	0	1	0	1
1	1	0	0	1
1	1	1	1	1

- 8. $\delta([q_4, B, B, c_i], [B, 0/1]) = ([q_4, B, B, c_i], [B, 0/1], L).$
 - Moves over symbols until it finds c

- 9. $\delta([q_4, B, B, c_i], [B, c]) = ([q_5, B, B, c_i], [B, c], L).$
 - Switches to q₅ when it sees c

- 10. $\delta([q_5, B, B, c_i], [B/*, 0/1]) = ([q_5, B, B, c_i], [B/*, 0/1], L).$
 - Moves over symbols until it finds c

 δ is defined by the following rules, in which a, b and c_i each may stand for either 0 or 1.

11.
$$\delta([q_5, B, B, c_i], [B, c]) = ([q_6, B, B, c_i], [B, c], L).$$

Switches to q₆ when it sees c

- 12. $\delta([q_6, B, B, c_i], [B, 0/1]) = ([q_6, B, B, c_i], [B, 0/1], L).$
 - Moves over unchecked symbols

- 13. $\delta([q_6, B, B, c_i], [*, 0/1]) = ([q_0, B, B, c_i], [*, 0/1], R).$
 - Switches to q_0 and starts moving right when it sees a checked symbol

 δ is defined by the following rules, in which a, b and c_i each may stand for either 0 or 1.

14.
$$\delta([q_0, B, B, c_i], [B, c]) = ([q_7, B, B, c_i], [B, c], R).$$

• If it sees c while at q_0 , switches to q_7

 δ is defined by the following rules, in which a, b and c_i each may stand for either 0 or 1.

15.
$$\delta([q_7, B, B, c_i], [B/*, 0/1/c]) = ([q_7, B, B, c_i], [B/*, 0/1/c], R).$$

Moves over everything until it sees a blank

- 16. $\delta([q_7, B, B, c_i], [B, B]) = ([q_8, B, B, B], [B, c_i], R).$
 - When it finds a blank, writes carry to tape, switches to q₈ and halts

Subroutines

- As with programs in general, it helps to think of Turing machines as built from a collection of interacting components, or "subroutines."
- A Turing machine subroutine is a set of states that perform some useful process.
- This set of states includes a start state and another state that temporarily has no moves, and that serves as the "return" state to pass control to whatever other set of states called the subroutine.

Subroutines

- The "call" of a subroutine occurs whenever there is a transition to its initial state.
- We note that the TM has no mechanism for remembering a "return address," that is, a state to go to after it finishes.
- We should our design of a TM call for one subroutine to be called from several states, we can make copies of the subroutine, using a new set of states for each copy.
- The "calls" are made to the start states of different copies of the subroutine, and each copy "returns" to a different state.