

APSP for unweighted undirected graphs
(dense graphs)

- Repeated BFS: $O(nm + n^2) = O(n^3)$ for dense graphs

Seidel's algorithm

$A \rightarrow$ adjacency matrix of G

$G^2 = (V, E^2)$

$(u, v) \in E^2$ if \exists path of length ≤ 2 in G

\rightarrow Computing the adjacency A_2 matrix of G^2 : Compute A^2 & set

$$A_2(i, j) = 1 \text{ if } A^2(i, j) > 0 \\ i \neq j \text{ or } A(i, j) > 0$$

\rightarrow what is the relation between the shortest paths in G^2 & G

Claim: If d_{uv} & D_{uv} are the shortest path distances in G & G^2 resp.

$$\text{then } D_{uv} = \left\lceil \frac{d_{uv}}{2} \right\rceil$$

Proof: $u \rightarrow a_1 \rightarrow b_1 \rightarrow a_2 \rightarrow b_2 \rightarrow \dots \rightarrow a_k \rightarrow b_k \rightarrow v$
even $\hookrightarrow G$

$$\left\{ \begin{array}{l} u \rightarrow a_1 \rightarrow b_1 \rightarrow a_2 \rightarrow b_2 \rightarrow \dots \rightarrow a_k \rightarrow b_k \rightarrow v \\ \quad \quad \quad \hookrightarrow G \\ u \rightarrow b_1 \rightarrow b_2 \rightarrow \dots \rightarrow b_k \rightarrow v \longrightarrow G^2 \end{array} \right.$$
$$\left\{ \begin{array}{l} u \rightarrow a_1 \rightarrow b_1 \rightarrow a_2 \rightarrow b_2 \rightarrow \dots \rightarrow a_k \rightarrow b_k \rightarrow a_{k+1} \rightarrow v \\ \qquad \qquad \qquad \qquad \qquad \qquad \qquad \qquad \qquad \qquad \hookrightarrow G \\ u \rightarrow b_1 \rightarrow b_2 \rightarrow \dots \rightarrow b_k \rightarrow v \end{array} \right.$$

$$D_{uv} \leq \left\lceil \frac{d_{uv}}{2} \right\rceil$$

If \exists path of length $l \leq \lceil \frac{d_{uv}}{2} \rceil$

Since each edge corr. to ≤ 2 edges in G

\exists path of length $2d < d_{uv}$

Corollary: $d_{uv} \in \{2D_{uv}, 2D_{uv}-1\}$

Lemma: ① If $d_{uv} = 2D_{uv}$, then $\forall w \in N_G(v)$

$$D_{nw} \geq D_{nr}$$

(2) If $d_{uv} = 2D_{uv} - 1$, then $\forall w \in N_G(v)$

$$D_{nw} \subseteq D_{nr}, \text{ and}$$

$$\exists z \in N_G(v) \text{ s.t. } D_{uz} < D_{uv}$$

Proof: ① Suppose $\exists w \in N_G(v)$ s.t

$$D_{uw} < D_{uv} \Rightarrow D_{uw} \leq D_{uv} - 1$$

$$2D_{uw} < 2D_{uv} - 1$$

Consider the shortest path from u to w

+ (w, v)

$$\text{has length} \leq 2D_{uw} + 1 < 2D_{uv} = d_{uv}$$

② For any $w \in N_G(v)$

Consider shortest path from u to v + (v, w)

$$d_{uw} \leq d_{uv} + 1$$

$$= 2D_{uv} \Rightarrow D_{uv} \geq \frac{d_{uw}}{2}$$

$$\Rightarrow D_{uv} \geq D_{uw}$$

Consider z in the shortest path from u to v

s.t $z \in N_G(v)$

$$d_{uz} = d_{uv} - 1 = 2D_{uv} - 2$$

$$\frac{d_{uz}}{2} = D_{uv} - 1 \Rightarrow D_{uz} < D_{uv}$$

Corollary: $d_{uv} = 2D_{uv} \Leftrightarrow \frac{\sum_{w \in N_G(v)} D_{uw}}{\deg(v)} \geq D_{uv}$

Consider the matrix $D \cdot A$

$$(D \cdot A)(u, v) = \sum_{w \in V} D(u, w) \cdot A(w, v)$$

$$= \sum_{w \in N_a(v)} D(u, w)$$