

# Dynamic Programming

## \* Text-segmentation

Given a string  $S[1, 2, \dots, n]$  &  
access to a function  $\text{IsWord}(i, j)$   
that checks if  $S[i, \dots, j]$  is a valid word  
Check if  $S$  can be partitioned into  
valid words

Recursion 1:

$$\text{split}(i) = \begin{cases} 1 & \text{if } S[1, \dots, i] \text{ can be split} \\ 0 & \text{o/w} \end{cases}$$

$$\text{split}(i) = \bigvee_{j=1}^i (\text{IsWord}(1, j) \wedge \text{split}(j)) \quad ?$$

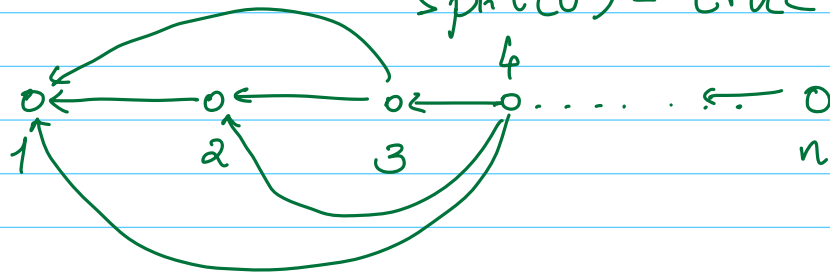
$$\text{split}(i) = \bigvee_{j=1}^i (\text{split}(j) \wedge \text{IsWord}(j+1, i))$$

\* Compute  $\text{split}(n)$

① How does the dependency digraph look?

② What are the base cases?

$\hookrightarrow \text{Split}(1) = \text{IsWord}(1, 1)$   
 $\text{Split}(0) = \text{true}$



### Memoized Version

Segment(j)

if  $j=0$  return true

if  $\text{Split}(j)$  defined

return  $\text{Split}(j)$

for  $1 \leq k \leq j$

if ( $\text{Segment}(k-1) \wedge \text{IsWord}(k, j)$ )

$\text{Split}(j) = \text{true} \rightarrow \text{Memoization}$

return

$\text{Split}(j) = \text{false}$

return

## Bottom-up version

Segment (n)

split(0) = true

for  $1 \leq i \leq n$   
split(i) = false  
for  $1 \leq k \leq i$

if (split(k-1)  $\wedge$  IsWord(k, i) )

split(i) = true (store split(i)=k)

$O(n^2)$

DFS on  
the  
digraph

- Understand the recursive structure and the dependencies to create the bottom-up version of the DP.
- You can construct the actual split of the string into segments as well.

(How?)