Computational Intractability
P: class of problems that have poly-time algorithms
- model independent
We will work with decision problems Lo Equivalent to search versions
Search Decision
Given G(V,E), n,v, i 4 4,v > find _ Answer "yes" if
The shortest path the it bit in the
to visa 1
Languages LC {0,1}* are abstractions
of computational problems
LEP if there is a polytime algorithm A S. F \forall $\alpha \in \{0,1\}^{*}$ Sometime
if xeL, then A(n)=1 if x&L, then A(n)=0

Solving vs Verifying Variables of cps.+ q(2)=13 Boolean circuit * Trivial algorithm: Try all the 2" assignments * Venification: if q is satisfiable, then 7 Satisfying e a witness that can be assignment checked NP (Non-deterministic Poly-time) LENP if I poly-time algorithm A that takes two inputs s. I tac so, 13* * if xEL, then FyESO,15* st 14/= 1x1° & A(x,y)=1 # if x & L, Then Hye So, 15 * s.t ly 1= las A (a,y)=0 Ja poly-time verifier A-lhat can be convinced of membership iff a EL

Some examples * SAT= { \$ | \$ is satisfiable } witness: satisfying assignment * VC= { (G, k) | G has a VC of size < k} witness: vertix cover Not all problems are in NP! * PRIMES = {n | n is a prime number} Lowhat is a short writicate to convince primatily? - PRIMES E NP (not trivial to slow) - PRIMESEP (AKS 2002) - F easily ventiable certificate to show that in is not prime * Linear Programming & has a P-time algorithm - Chuk if f < K f = min \(\sum_{i} \arra{c}_{i} \arra{c}_{i} \) S.t $\sum a_{ij} x_{j} \ge b_{j}$ -Not clear why i=1 Should have a $\forall j \in \{1, 2, ..., m\}$ represe poly in input aij, bj, Ci ER