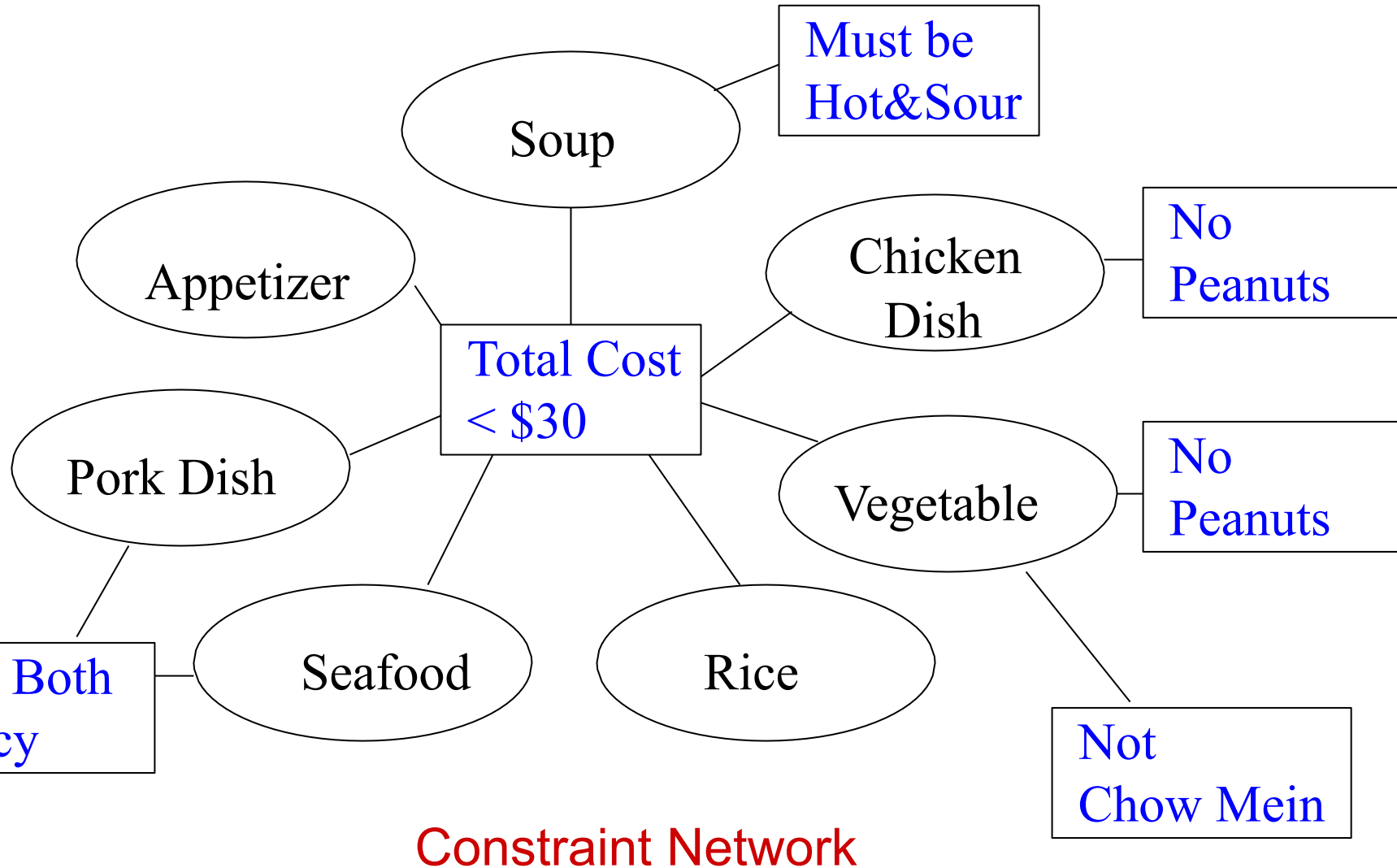


Constraint Satisfaction Problems



Formal Definition of CSP

- A constraint satisfaction problem (**CSP**) is a triple **(V, D, C)** where
 - **V** is a set of variables **X_1, \dots, X_n** .
 - **D** is the union of a set of domain sets **D_1, \dots, D_n** , where **D_i** is the domain of possible values for variable **X_i** .
 - **C** is a set of constraints on the values of the variables, which can be pairwise (simplest and most common) or **k** at a time.

CSPs vs. Standard Search Problems

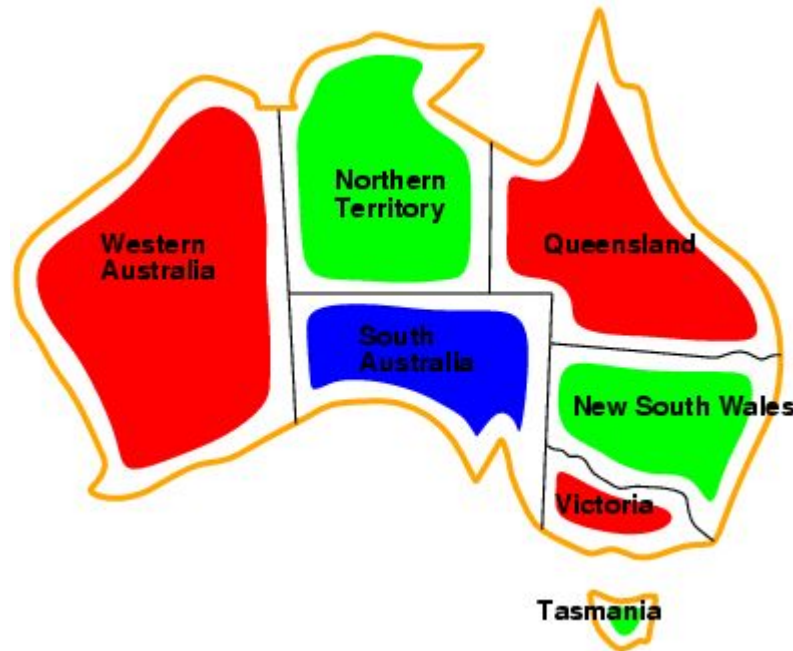
- Standard search problem:
 - **state** is a "black box" - any data structure that supports successor function, heuristic function, and goal test
- CSP:
 - **state** is defined by **variables** X_i with **values** from **domain** D_i
 - **goal test** is a set of **constraints** specifying allowable combinations of values for subsets of variables
- Simple example of a **formal representation language**
- Allows useful **general-purpose** algorithms with more power than standard search algorithms

Example: Map-Coloring



- **Variables** WA, NT, Q, NSW, V, SA, T
- **Domains** $D = \{\text{red}, \text{green}, \text{blue}\}$
- **Constraints** _{i} : adjacent regions must have different colors
- e.g., $WA \neq NT$, or $(WA, NT) \in \{(\text{red}, \text{green}), (\text{red}, \text{blue}), (\text{green}, \text{red}), (\text{green}, \text{blue}), (\text{blue}, \text{red}), (\text{blue}, \text{green})\}$

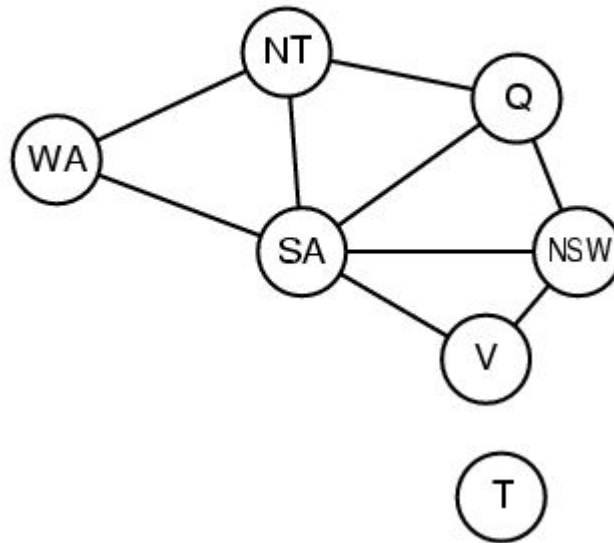
Example: Map-Coloring



- Solutions are **complete** and **consistent** assignments, e.g., WA = red, NT = green, Q = red, NSW = green, V = red, SA = blue, T = green

Constraint graph

- **Binary CSP:** each constraint relates two variables
- **Constraint graph:** nodes are variables, arcs are constraints



Varieties of constraints

- **Unary** constraints involve a single variable,
 - e.g., $SA \neq \text{green}$
- **Binary** constraints involve pairs of variables,
 - e.g., $\text{value}(SA) \neq \text{value}(WA)$
- **Higher-order** constraints involve 3 or more variables,
 - e.g., cryptarithmic column constraints

Real-world CSPs

- **Assignment problems**
 - e.g., who teaches what class
- **Timetabling problems**
 - e.g., which class is offered when and where?
- **Transportation scheduling**
- **Factory scheduling**

Notice that many real-world problems involve real-valued variables

What Kinds of Algorithms are used for CSP?

- Backtracking Tree Search
- Tree Search with Forward Checking
- Tree Search with Discrete Relaxation (arc consistency, k-consistency)
- Many other variants
-

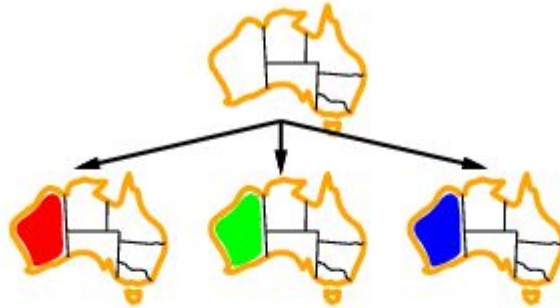
Backtracking Tree Search

- Variable assignments are **commutative**, i.e.,
[WA = red then NT = green] same as [NT = green then WA = red]
- Only need to consider assignments to a single variable at each node.
- Depth-first search for CSPs with single-variable assignments is called **backtracking** search.
- Backtracking search is the basic uninformed algorithm for CSPs.
- Can solve n -queens for $n \approx 25$.

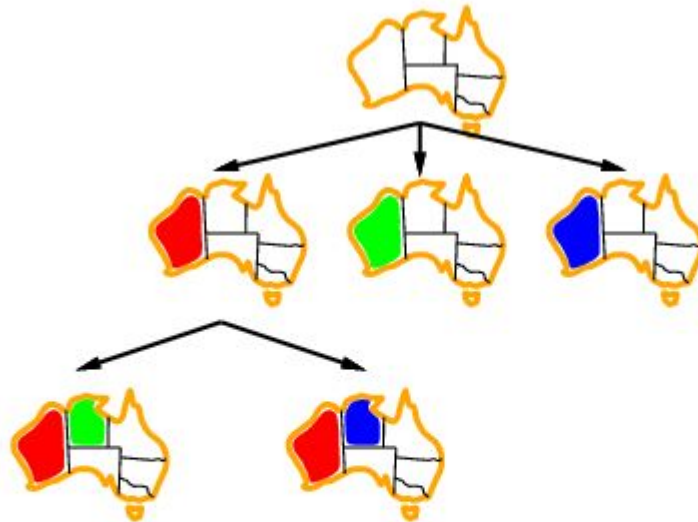
Backtracking Example



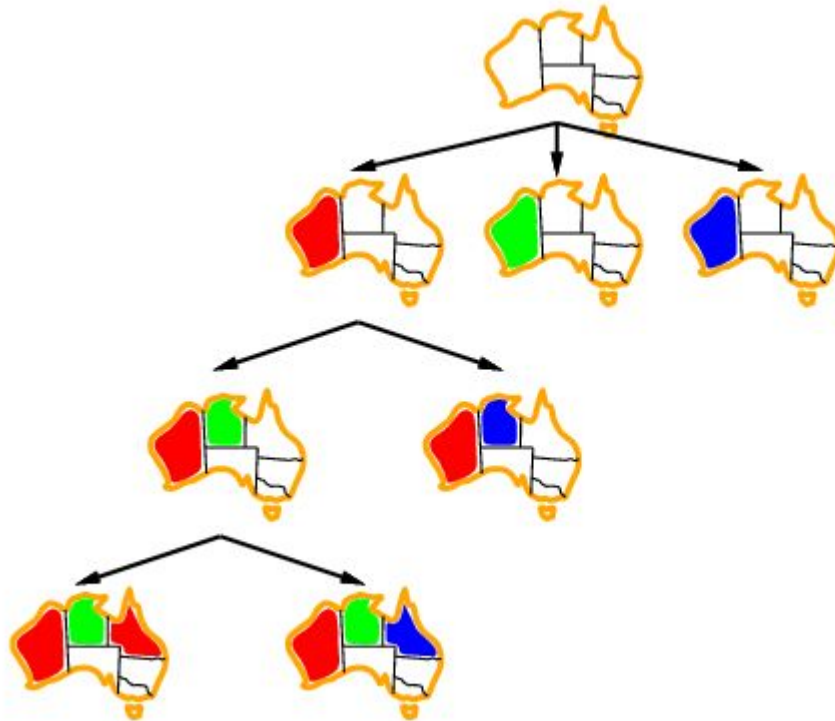
Backtracking Example



Backtracking Example

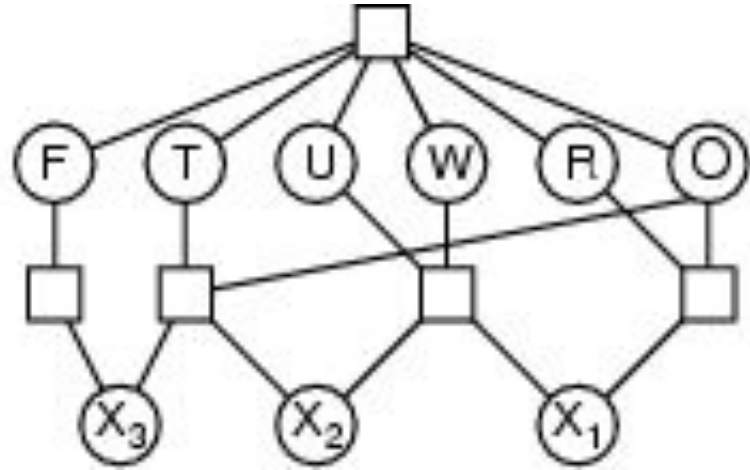


Backtracking Example



Example: Cryptarithmic

$$\begin{array}{r} \text{TWO} \\ + \text{TWO} \\ \hline \text{FOUR} \end{array}$$



- **Variables:**
 $\{F, T, U, W, R, O, X^1, X^2, X^3\}$
- **Domains:** $\{0, 1, 2, 3, 4, 5, 6, 7, 8, 9\}$
- **Constraints:** *Alldiff* (F, T, U, W, R, O)
 - $O + O = R + 10 \cdot X^1$
 - $X^1 + W + W = U + 10 \cdot X^2$
 - $X^2 + T + T = O + 10 \cdot X^3$
 - $X^3 = F, T \neq 0, F \neq 0$

Improving Backtracking Efficiency

- **General-purpose** methods can give huge gains in speed:
 - Which variable should be assigned next?
 - In what order should its values be tried?
 - Can we detect inevitable failure early?

Most Constrained Variable

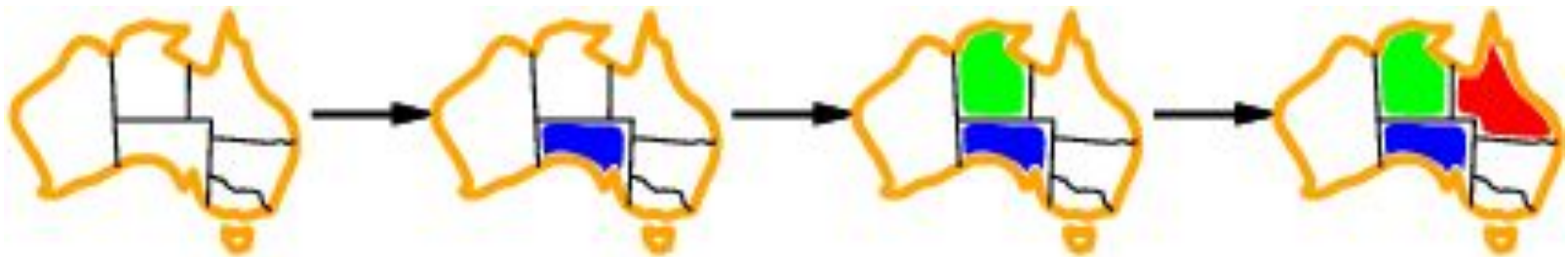
- Most constrained variable:
choose the variable with the fewest legal values



- a.k.a. **minimum remaining values (MRV)**
heuristic

Most Constraining Variable

- Tie-breaker among most constrained variables
- Most constraining variable:
 - choose the variable with the most constraints on remaining variables



Least Constraining Value

- Given a variable, choose the least constraining value:
 - the one that rules out the fewest values in the remaining variables

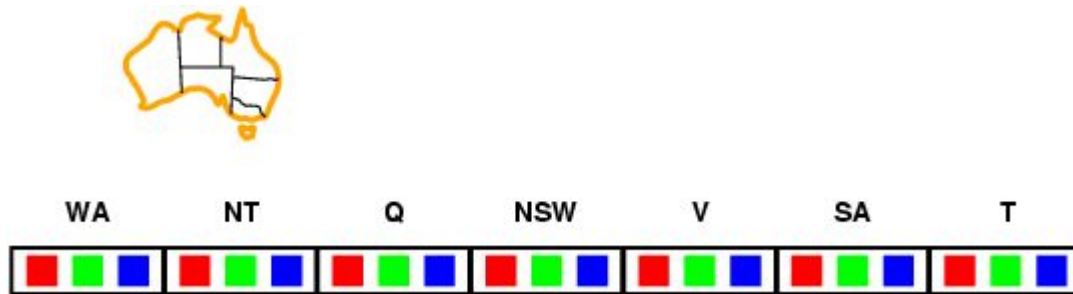


- Combining these heuristics makes 1000 queens feasible

How do we incorporate forward checking into a backtracking depth-first search?

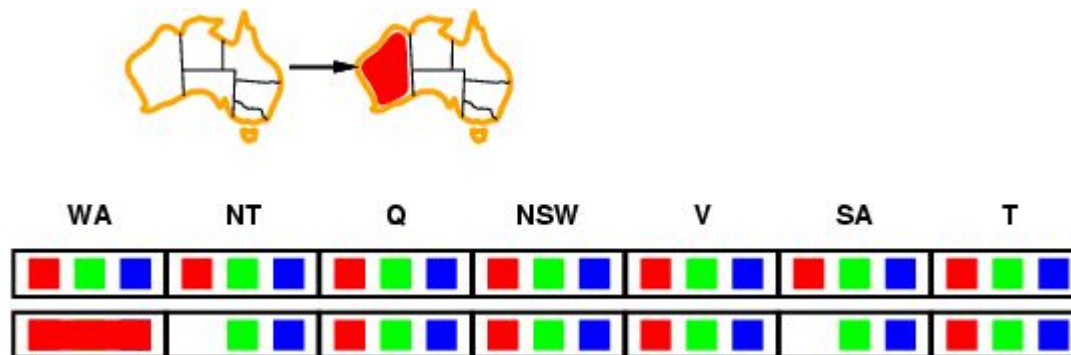
Book's Forward Checking Example

- Idea:
 - Keep track of remaining legal values for unassigned variables
 - Terminate search when any variable has no legal values



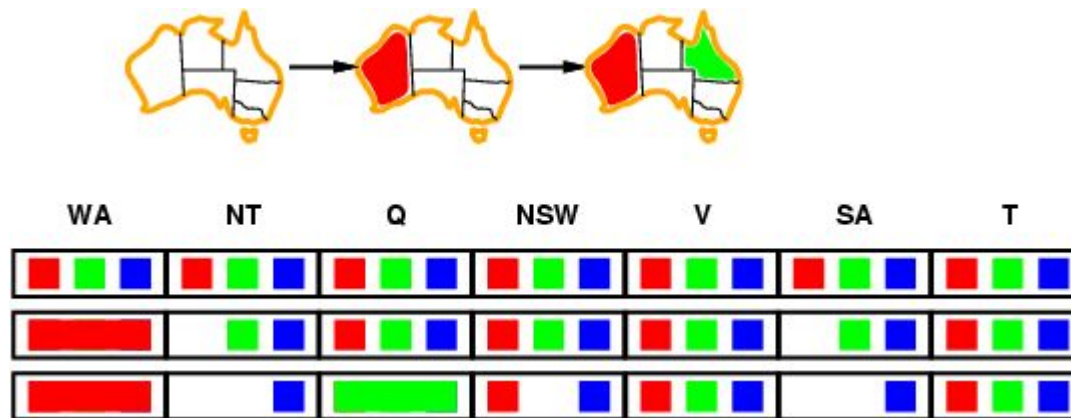
Forward Checking

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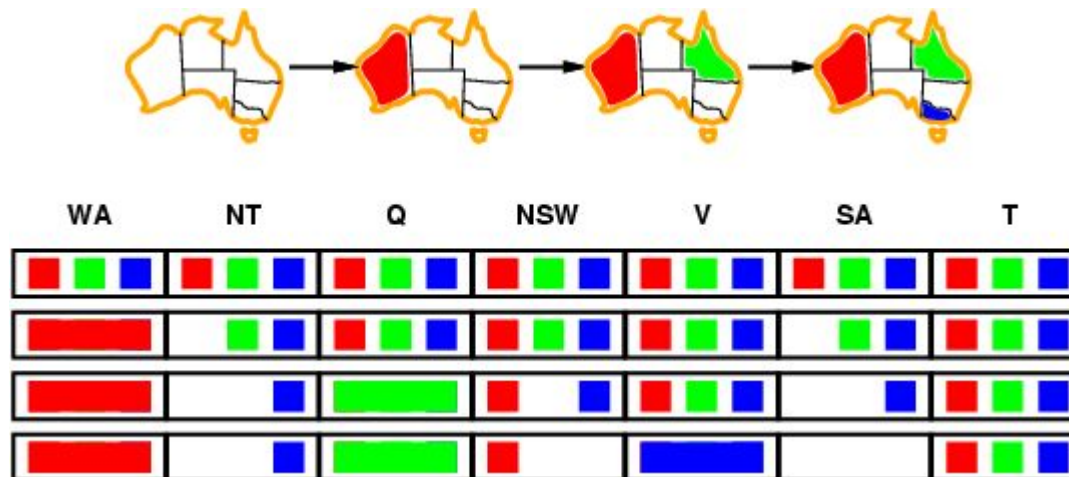
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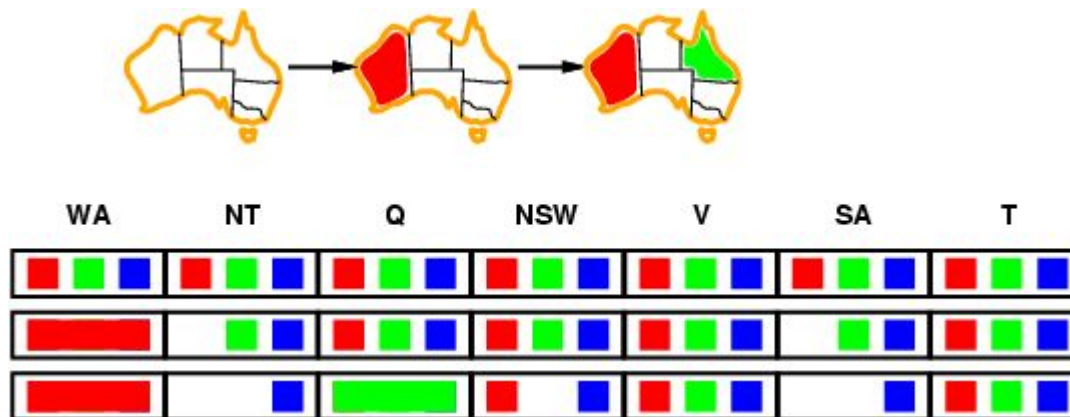
Forward Checking

- Idea:
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Constraint Propagation

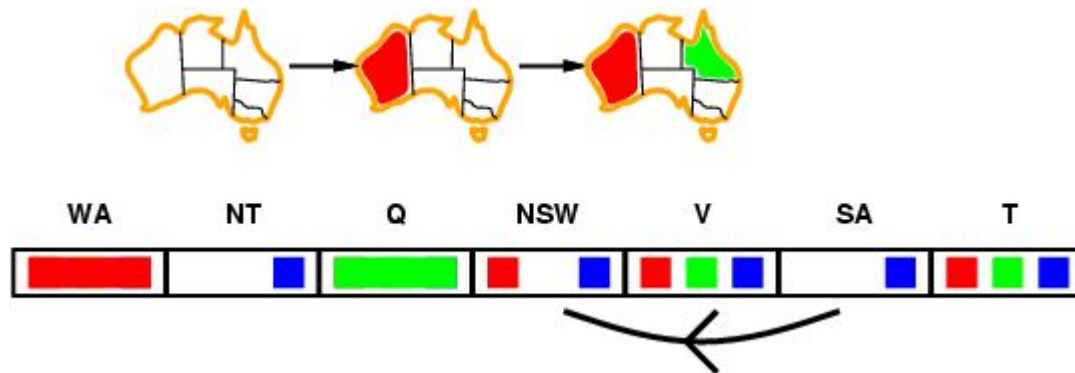
- Forward checking propagates information from assigned to unassigned variables, but doesn't provide early detection for all failures:



- NT and SA cannot both be blue!
- Constraint propagation repeatedly enforces constraints locally

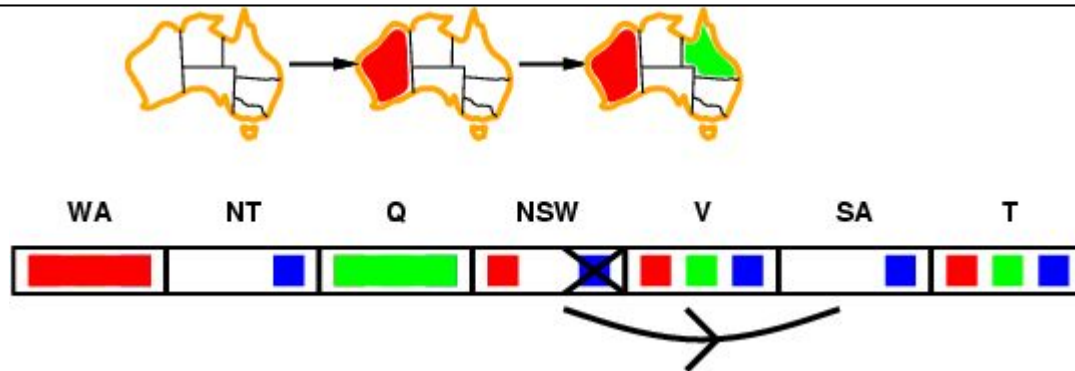
Arc Consistency

- Simplest form of propagation makes each arc **consistent**
- $X \bowtie Y$ is consistent iff
for **every** value x of X there is **some** allowed value y of Y



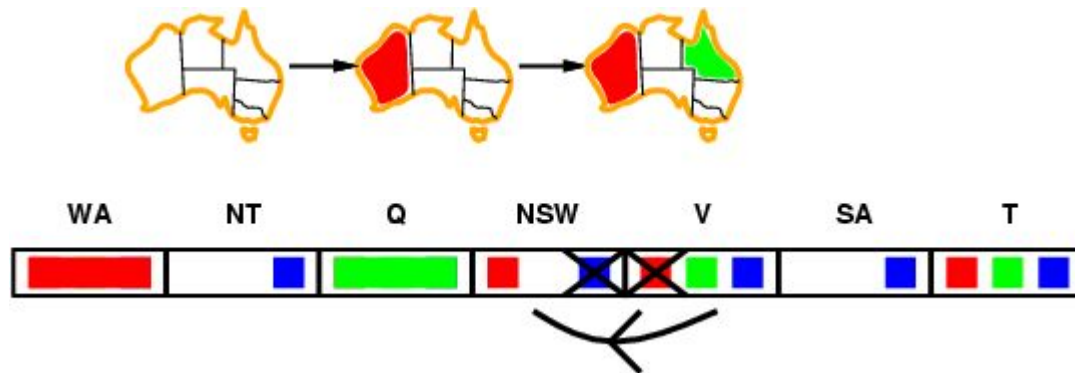
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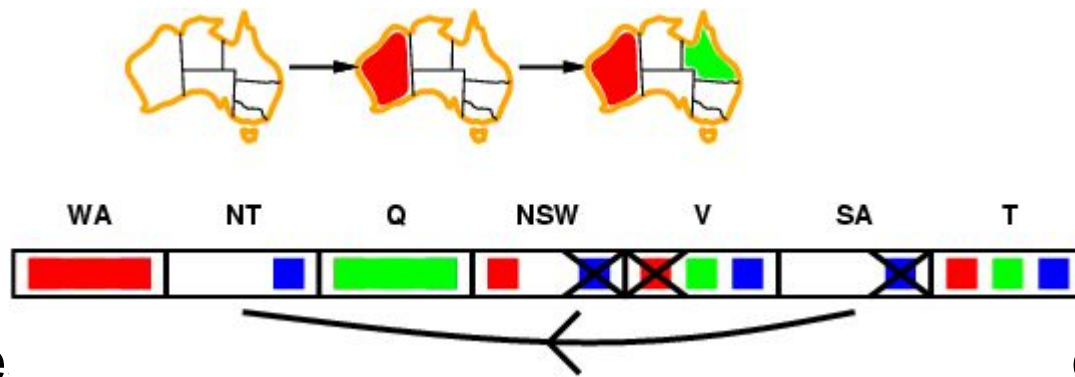
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- If X loses a value, neighbors of X need to be rechecked

Arc consistency

- Simplest form of propagation makes each arc **consistent**
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for **every** value x of X there is **some** allowed value y of Y



- If X loses a value, neighbors must be rechecked
- Arc consistency detects failure earlier than forward checking
- Can be run as a preprocessor or after each assignment

Comparison of Methods

- Backtracking tree search is a blind search.
- Forward checking checks constraints between the current variable and all future ones.
- Arc consistency then checks constraints between all pairs of future (unassigned) variables.

- What is the complexity of a backtracking tree search?
- How do forward checking and arc consistency affect that?