# 面成流经营住.

1. Clarification | Edge cases | input type + input size and range | output type | time + Space complicity | give examples | 多華的日

Objectses:

(b) 负数 10, 已数

② 数的 type

② 极偏值

③ empty

2. Think out loud: speak whatever you are thinking 实在 3-5%钟超元进来,要hint

ifs
3. The before coding: Write idea and steps: else
Edge cases.
Run an example
な 記想 Test case. i) 面低電 对元对
な 智治等 → 把所有特別 用commont 写出

4. Test in real time: 1 Edge cases
example
in-out momory
large input

5. Refrotor the code: Make it more efficient.

Frequently opposition -> make it a value

dutestructure 店屋文化 男(第一つ onestep systep Communitate Code Imp System design

## Leet Code \$15 -> Amay

1. 数组有序? Yes: binary Search: 是否有相同值 No: O(n): 双指针, Hosh table, DP No: >O(n): 先 Sort, 后用 binary Search

id: Peut - Binay Search

- 2. 如果主体导位 > O(nlyn), 那么可以先 SOIT 共看是否会简单些
- 2 红B+Oon): Hosh tuble, 双稻计 乱为 +Ocn + Space Oci : 双档针
- 4. In place 操作: { · index 作为值: g arr [index] = arr [arrinks];

  Ou) 例用 Swap 季版
  用取点式 特殊值 标记

  S. 问 [Min 的种类数 OP Greedy
- 6. 不要求 space os 耐候: temp almy 也许全大大简化过程
- 7. 双指针: 可以相互追逐 武 两编同时扫 (追击
- 8. 对于找知单题,一定要运营问题特征,一等一步发现,试索
- 9. 旧台问题: | Back tracking | Bit Minpulation
- 10. Range 可能 step= current + arr [current]
- 11. 相同模式问题: Reconsion

1 digit = step Result 3 10;

- 12. Carry 问题: Step Result = digit[]+digit[]]+carry; | Carry = Step Result /10;
- 1). 超字种类有限: Counting Sort
- 4 有时间差接作无需移动 数因元李
- 1. Majorty: Moor Voting Algo

1. 字符串 dp:

Mecaning: dp[i][j] text(0.i) pattorn(0,j)

init: ...." a b c

J.

function: 填dp[i][j] 找 你律

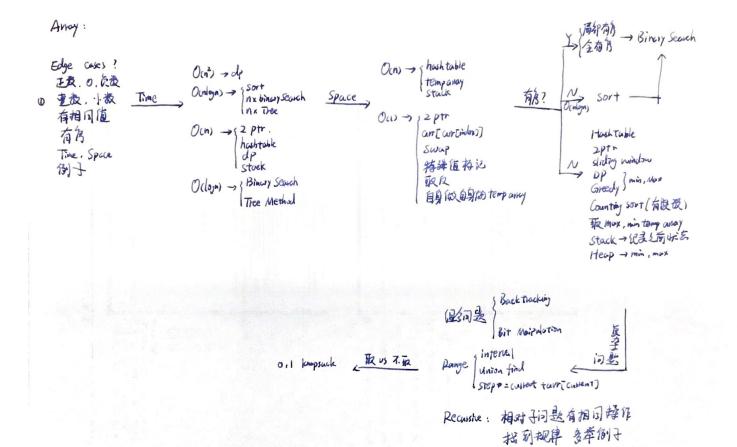
#### lost Code & ts → Linked List

- 1. Fake Head
- 2. (fost pointer
- 3. Revenue Linked List: 当需要逆向走时,先Housek,多Theuerse可悔其变为双向征丧.
- 4. 链表现的版 (值 以色别目的)
- 5. Linked List len ffort, slow 指电、闭截点, One pass. 花林俊, 再接作 > thu passes.
- 6. two or three pointous.
- 7. 做匙对一定要画图.
- 8. 用Hash May 存起 也许名简化挣印
- 9. Insert Node

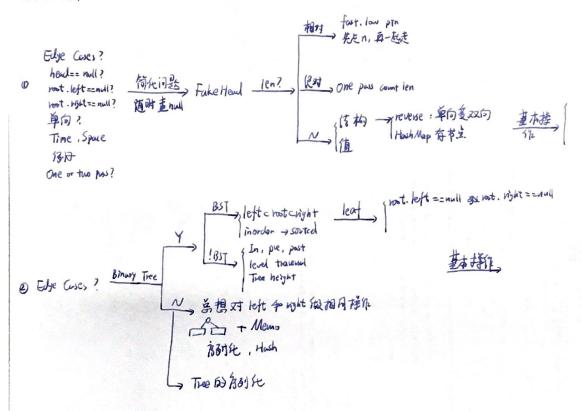
### Leet Code & 15 - Tree

- 1. In order, pie-order, past-order Trainsal
- 2. Level order Queue + count
- 3. Tree Depth
- 4. 节点、上有值的, 色入 sum, 使其为 sum-value 进行更新
- 5. The Node Insert
- 6. 不住 root,且版图的格尔 { left = Nath. max(0, helper(noot. right));
  result = Nath. max(0, helper(root. right));
  result = Math. max(result, left + right + root. red);
  return Nath. max(left, right) + root. val;
- 7. BST + inonde = 有序数组
- 8. balenced BST = 取中立、主动vost 「中主-left = recursion() 中主、right = recursion()
- 9- Lonest Common Amoster of BST:是像在一个Subtree 中.
- 10 Leaf: root.lest == null 8,8 root.right == null
- 11. Same Tree?
- 12. 从 n=1 → n= 4 举例 → 找 Base case → 找 送归模式 → 确定参数.
- 13. 序列化, 查重 serialize
- 14. 将树 Hash 成一个值

int t = dfs ( not left);
int r = dfs ( nort . lfyht);
int north trs ((f 1253 x31) + r 1235) x31 + nort. val;
return nort v ;



#### Linked List



Graph.

DFS:
BFS: shortest porth
Topo Sort: Pre requisits
Union find: Sets marge
は、佐京
i可是

プレートを存在版相同接続
Param 和 Meturn 励送程

1. DES 模板, Basic Idea, param pais

void offs (intests) map, int row, int col, intests visited) public

if (row 20 11 col co 11 row > map. length 11 col > rooms. length 11 visited [row] [col] == true) return !

visited trow] [col] = true;

dfs(10w-1, col);

ofs ( now+1, col);

dts ( pow, col-1);

dfs (vow, culti);

2 Topological sort.

- ① 图 HashMap < Key, Set < child>> Graph.
- ②入後 HashMap (Key, Integer>
- 3 BFS

3. Union find.

- ① Make Set : Q 自己指历自己。
- @ Union: path compression + always union with parent mode
- @ find : walk along parent pointer

Data structure: { int rank

lint data

Node present

4. BPS.