

4月10日作业

1

1. Find

$$\min_{a,b \in \mathbb{R}} \max_{-1 \leq x \leq 2} |x^3 + ax + b|$$

without programming.

解: 先不妨猜测 $x_0 = -1$ 与 $x_2 = 2$ 两个端点是 Chebyshev 交错点,

设中间的交错点为 x_1 , 其满足

$$(x_1^3 + ax_1 + b)' = 3x_1^2 + a = 0 \quad (1)$$

由 Chebyshev 交错点的定义可知:

$$x_0^3 + ax_0 + b = x_2^3 + ax_2 + b$$

$$\text{即 } -1 - a + b = 8 + 2a + b \quad (2)$$

综合 (1) 与 (2) 得到: $a = -3$, $x_1 = 1$ (舍去)

再次利用交错点组的定义:

$$x_1^3 + ax_1 + b = -(x_0^3 + ax_0 + b)$$

$$\text{即 } 1 - 3 + b = -(-1 + 3 + b)$$

$$\Rightarrow b = 0$$

$$\text{令 } f(x) = x^3 - 3x$$

$$f(-1) = -1 + 3 = 2$$

$$f(2) = 8 - 6 = 2$$

$$\frac{1}{2} f'(x) = 3x^2 - 3 = 0 \Rightarrow x = \pm 1$$

$$\text{其中 } f(1) = 1 - 3 = -2$$

$\Rightarrow x_0 = -1, x_1 = 1, x_2 = 2$ 确实是 $x^3 - 3x$ 的交错点组

$\Rightarrow p(x) = 3x$ 是对 $f(x) = x^3$ 在 $[-1, 2]$ 上的最佳一致逼近

$$\Rightarrow \min_{a,b} \max_{-1 \leq x \leq 2} |x^3 + ax + b| = 2$$

$$\text{并且仅当 } a = -3, b = 0$$

□

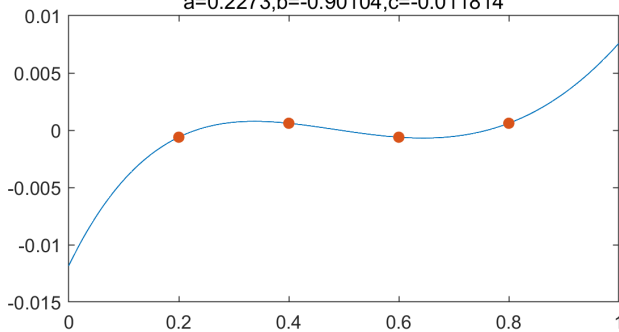
2. Find

$$\min_{a,b,c} \max_{0 \leq x \leq 1} |\ln(1+x) + ax^2 + bx + c|$$

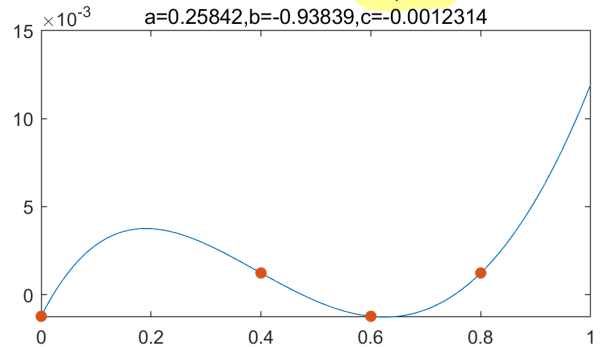
2

using Remez algorithm. Visualize the approximation error and the (nonuniform) alternating set by taking a few snapshots. (E.g., you may plot the error curve for the initial guess, two intermediate solutions, as well as the final solution.)

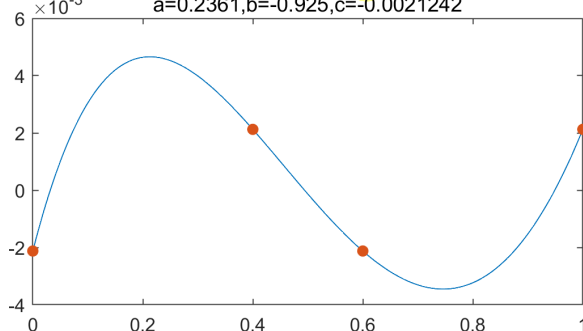
the initial alternating set and the error curve
a=0.2273, b=-0.90104, c=-0.011814



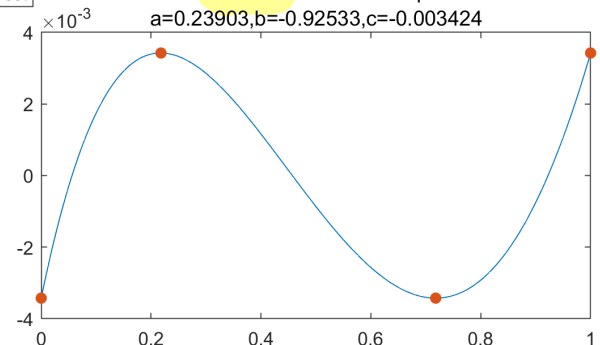
the error curve of loop=1
a=0.25842, b=-0.93839, c=-0.0012314



the error curve of loop=2
a=0.2361, b=-0.925, c=-0.0021242



the final error curve of loop=6
a=0.23903, b=-0.92533, c=-0.003424



□

3

3. Show that in each iterate of Remez algorithm, the linear system has a unique solution.

解: 线性系统可以表示为:

$$f(x_i) - \sum_{k=0}^n a_k \cdot x_i^k = (-1)^i \eta, \text{ 其中 } i=0, 1, \dots, n+1$$

其中 a_0, \dots, a_n, η 是 $n+2$ 个未知数, 写成矩阵形式:

$$\begin{bmatrix} x_0^n & \dots & x_0, & 1, & 1 \\ \vdots & & & & \\ x_{n+1}^n & \dots & x_{n+1}, & 1, & (-1)^{n+1} \end{bmatrix}_{(n+2)} \begin{bmatrix} a_0 \\ \vdots \\ a_n \\ \eta \end{bmatrix} = \begin{bmatrix} f(x_0) \\ \vdots \\ f(x_{n+1}) \end{bmatrix}$$

(记矩阵为 A)

注意到沿着第 $n+2$ 列 (最后一列) 进行展开时, 所有的余子式都是 Vandermonde 行列式

$$\text{记 } F(x_k) = \prod_{\substack{0 \leq i < j \leq n+1 \\ i, j \neq k}} (x_j - x_i) > 0$$

$$\Rightarrow \det A$$

$$= (-1)^{n+1} \cdot \left[1 \cdot F(x_0) + (-1)^{n+1} F(x_1) + \dots + (-1)^{n+1+n+1} F(x_{n+1}) \right]$$

$$= (-1)^{n+1} \cdot \sum_{k=0}^{n+1} F(x_k) \neq 0 \Rightarrow A \text{ 可逆} \Rightarrow \text{有唯一解}$$

□

4. Let $p(x) = \sum_{k=0}^n a_k x^k$ be a real polynomial such that

$$\max_{-1 \leq x \leq 1} |p(x)| \leq 1.$$

4

Show that $|a_n| \leq 2^{n-1}$.

Remark: There was a mistake regarding the location of Chebyshev alternating set in today's lecture. You will be able to fix the mistake by solving this exercise.

解:
$$\begin{aligned} p(x) &= a_n x^n + \dots + a_0 \\ &= a_n \cdot \left(x^n + \frac{a_{n-1}}{a_n} x^{n-1} + \dots + \frac{a_0}{a_n} \right) \\ &= a_n \cdot (x^n + b_{n-1} x^{n-1} + \dots + b_0) \end{aligned}$$

其中 $b_i = \frac{a_i}{a_n}$, $a_n \neq 0$, $i = 0, 1, 2, \dots, n-1$

记 $\tilde{p}(x) = x^n + b_{n-1} x^{n-1} + \dots + b_0$

$\tilde{p}(x)$ 是 n 阶首一多项式

$$\Rightarrow \max_{-1 \leq x \leq 1} |p(x)| = |a_n| \cdot \max_{-1 \leq x \leq 1} |\tilde{p}(x)|$$

则只需证明 $\max_{-1 \leq x \leq 1} |\tilde{p}(x)| \geq 2^{1-n}$ 即可

$$\begin{aligned} \tilde{p}(x) &= x^n + (b_{n-1} x^{n-1} + \dots + b_0) \\ &= x^n - p_{n-1}(x) \end{aligned}$$

其中 $p_{n-1}(x)$ 是不超过 $n-1$ 次的多项式

那么只需证明:

$$\min_{p_{n-1}} \max_{-1 \leq x \leq 1} |x^n - p_{n-1}(x)| \geq 2^{1-n}$$

也就是 $\|x^n - p_{n-1}^*(x)\|_\infty \geq 2^{1-n}$

其中 $p_{n-1}^*(x)$ 是对 x^n 的 $n-1$ 次最佳一致逼近多项式

记 n 次 Chebyshev 多项式为 $T_n(x)$, 下面来说明:

$$x^n - p_{n-1}^*(x) = 2^{1-n} \cdot T_n(x) \text{ 时}$$

$p_{n-1}^*(x)$ 恰是 $n-1$ 次的最佳一致逼近

$$\text{注意到 } T_n(x) = \cos(n \cdot \arccos x)$$

$\Rightarrow T_n(x)$ 在 $[-1, 1]$ 上有 $n+1$ 个最值点:

$$x = \cos\left(\frac{k\pi}{n}\right), \quad k=0, 1, \dots, n$$

且相邻最值点依次取 $1, -1, 1, -1, \dots$

此外, $2^{1-n} \cdot T_n(x)$ 是 n -次多项式 ($T_n(x)$ 最高项系数为 2^{n-1})

$$\text{则 } p_{n-1}^*(x) = x^n - 2^{1-n} \cdot T_n(x) \text{ 不超过 } n-1 \text{ 次}$$

则由 Chebyshev 定理, $x^n - p_{n-1}^*(x) = 2^{1-n} T_n(x)$ 时,

$$\text{也都 } p_{n-1}^*(x) = x^n - 2^{1-n} \cdot T_n(x) \text{ 时}$$

$p_{n-1}^*(x)$ 是 x^n 的不超过 $n-1$ 次最佳一致逼近

$$\Rightarrow \|x^n - p_{n-1}^*(x)\|_\infty = \|2^{1-n} T_n(x)\|_\infty = 2^{1-n}$$

$$\Rightarrow |a_n| \leq \frac{1}{\|\tilde{p}(x)\|_\infty} \leq \frac{1}{\|x^n - p_{n-1}^*(x)\|_\infty} = 2^{n-1}$$

□