```
task P_i.Join(R)
    wait for [S_j] ← R
    for wait for P_k ∈ [T_j] ← S_j do
        [hello] → P_k
        RTT_k = timeit wait for [hello] ← P_k with timeou
    N(P_i) = {P_k | RTT_k < timeout}
    N*(P_i) = {first K P_k ∈ N(P_i) | RTT_k < RTT_{k+1}}
task P_o.Leave(S_j)
    [goodbye] → S_j
    for P_k ∈ N(P_o do [goodbye] → P_k
    relay pending chunks
    wait for [goodbye] ← S_j with timeout
    while timeout do
        [goodbye] → S_j
        wait for [goodbye] ← S_j with timeout
task P_k.EchoHello()
    while True do
        wait for [hello] ← P_i
        [hello] → P_i
        N(P_k) = (NP_k) ∪ P_i
task P_k.ProcessGoodbye()
    while True do
        wait for [goodbye] ← P_o
        N(P_k) = N(P_k) \ P_o
task P_k.ShortestPathRouting()
    for all P_i ∈ N(P_k) do
        D[P_i] = ∞
    D[P_k] = 0
    for all P_i ∈ N*(P_k) do
        G[P_i] = P_i
    for all P_i ∈ N*(P_k) do
        [G[·]] → P_i
    while True do
        found_shorter_route = false
        [D*_{P_n}[·], G*_{P_n}[·]] ← from any P_n ∈ N*(P_k)
        for all D*_{P_n}[P_i] ∈ D*_{P_k}[·] do
            alternative_distance_P_i = D*_{P_n}[P_i] + D[P_n]
            if alternative_distance_P_i < D[P_j] then
                found_shorter_route = true
                D[P_i] = alternative_distance
                G[P_i] = P_n
        if found_shorter_route then
            for all P_i ∈ N*(P_k) do
```

$[D[\cdot], G[\cdot]] \rightarrow \mathcal{P}_i$

task $\mathcal{P}_k$.ChunkFlooding()

while True **do**

wait for $[x, \texttt{chunk}] \leftarrow \mathcal{S}_j$ **or** $[x, \mathcal{P}_i, \texttt{chunk}] \leftarrow \mathcal{P}$

if sender $= \mathcal{S}_j$ **then**

for all $\mathcal{P}_n \in N^*(\mathcal{P}_k)$ **do**

$[x, \mathcal{P}_k, \texttt{chunk}] \rightarrow \mathcal{P}_n$

else

for all $\mathcal{P}_n \in N^*(\mathcal{P}_k) \setminus \mathcal{P}_m$ **do**

if $\mathcal{P}_k \in G^*_{\mathcal{P}_n}[\mathcal{P}_i]$ **then**

$[x, \mathcal{P}_i, \texttt{chunk}] \rightarrow \mathcal{P}_n$

task $\mathcal{P}_k$.FreeRidingControl()

while True **do**

for all $\mathcal{P}_o \in N(\mathcal{P}_k)$ **do**

& **wait for** $[x, \mathcal{P}_i, \texttt{chunk}] \leftarrow \mathcal{P}_o$

with timeout **of** L rounds

on timeout $N(\mathcal{P}_k) = N(\mathcal{P}_k) \setminus \mathcal{P}_o$