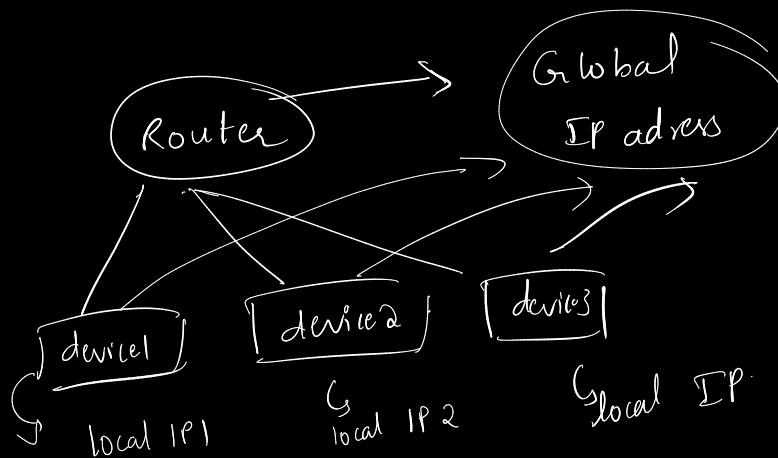


## IP Address:

- Router assigns local IP using DHCP (Dynamic Host Configuration protocol)

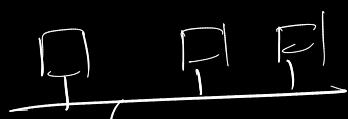


- Say its D1 made a request to google the router decides it by NAT (Network Access Translator)
- Which device is decided by IP address and which application to send is decided by port number.

- PORT number is a 16-bit number. Total ports possible are  $2^{16}$
- Web pages use HTTP protocol. *(Say some rules set by very cool people at Internet society.)*
- 0 - 1023  $\Rightarrow$  Reserved ports for HTTP stuff.
- 1024 - 49152  $\Rightarrow$  Applications

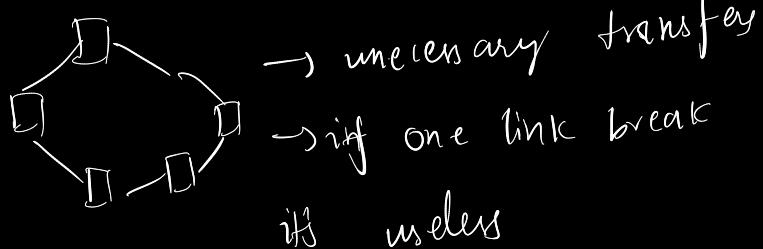
# Topology:

1) Bus Topology

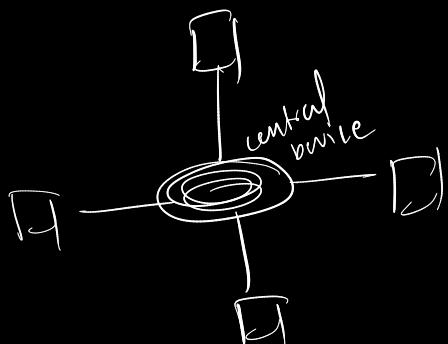


if link break all  
are disconnected

2) Ring Topology

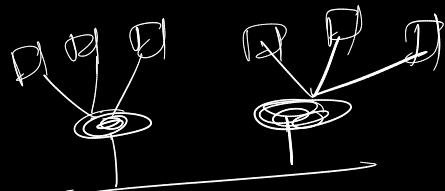


3) Star Topology



→ If central device  
stop the connection  
is gone.

4) Tree (Star + Bus)



5) Mesh → every single comp connected to every other  
comp

⇒ expensive  
⇒ less scalability

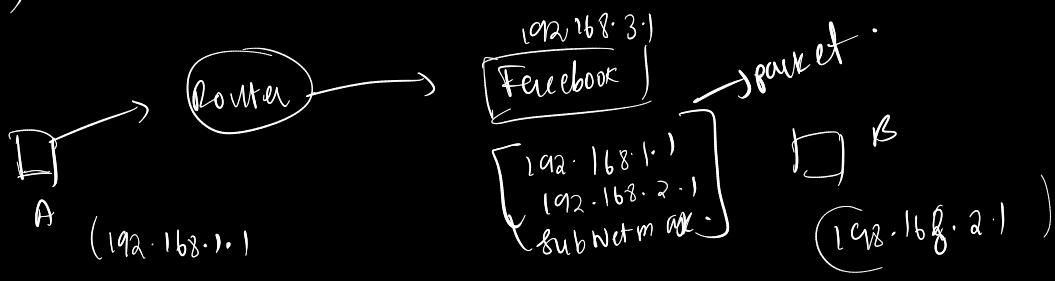
# Structure of the Network:

## → OSI MODEL (Open System Inter Connection)

- There are 7 layers

- Example: A whatsapp msg.

- 1) WhatsApp see msg sending (Application layer)
- 2) Application layer → Presentation layer
- 3) Presentation Layer converts ASCII to something else (Encryption + compression + translation)
- 4) Session layer ⇒ Authentication, Authorisation stuff. It just establishes & terminates sessions
- 5) Transport Layer ⇒
  - i) Segments (divides into small parts)
  - ii)
- 6) Network Layer ⇒ sends to computer on other networks say the router www in the layer. IP addressing of senders & receivers and forms IP packet and also routing
- 7) Data Link ⇒ receives data packet from network layer



8) Handles logical addressing & physical addressing

TCP/IP Model : (Practically more used in real world)  
It contains lesser steps compared to

OSI Model -

1. Application
2. Transport
3. Network
4. Datalink
5. Physical

## 1) Application Layers:

- i) User interaction
- ii) Where: devices
- iii) Protocols
- iv) Client- Server Architecture

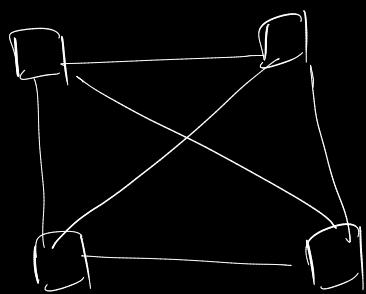
⇒ Collection of servers are called data centers

⇒ Client- Server Architecture

$$\boxed{\text{Client}} \geq \boxed{\text{Server}}$$

⇒ Peer to Peer (P to P) Architecture

Eg: Bit torrent



Key advantage

- 1) Scalable
- 2) Decentralized
- 3) Every Comp acts as Client & Server

## ⇒ Protocols:

### \* TCP / IP :

- \* HTTP
- \* DHCP
- \* FTP (File Transfer Protocol)
- \* SMTP (Simple Mail Transfer Protocol) (for sending EMAIL)
- \* POP3 & IMAP (Receiving EMAIL)
- \* SSH (To login to someone's comp using terminal)
- \* UDP (. stateless  $\Rightarrow$  Data maybe lost)

\* Sockets : Used to send msg's from one comp  $\rightarrow$  other.

\* Ports : Tells which application

say we know to send data chrome but to which tab will be determined by "Ephemeral Ports". It basically assigns its own random ports.

\* HTTP : It is a client server protocol. It tells how to request data and how to send data. This is basically a application Layer protocol. HTTP uses TCP (Transmission Control Protocol) inside it.



(at transport layer)

## Cookies

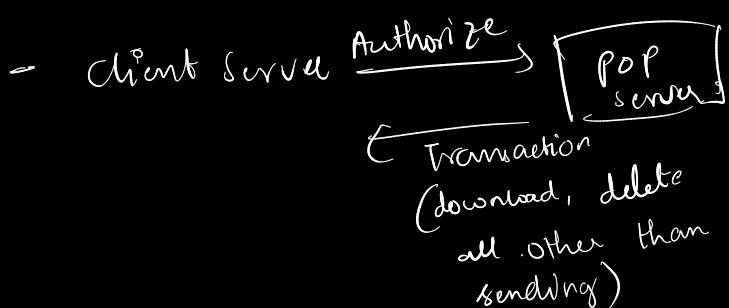
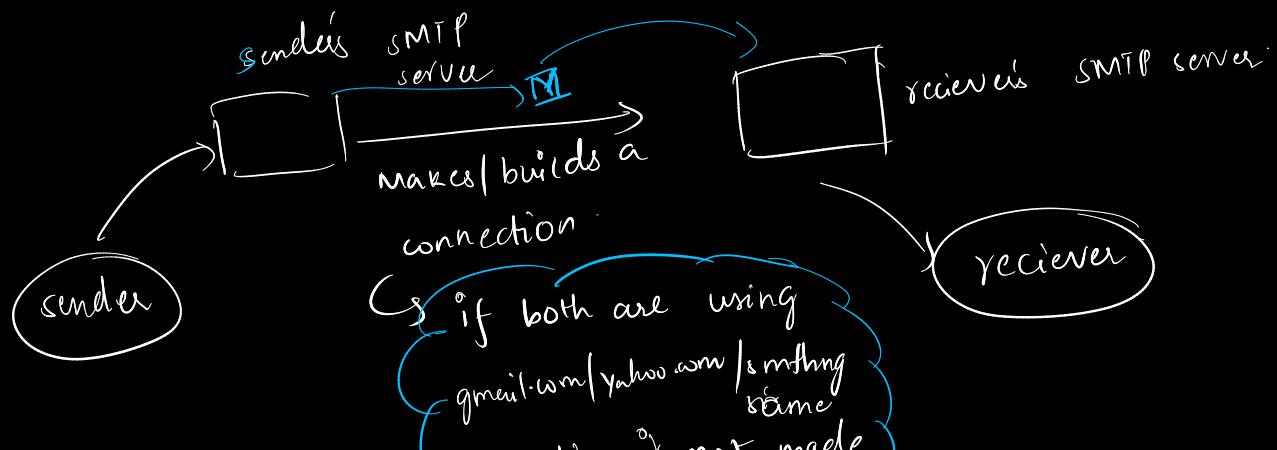
A unique string stored in browser. Login data is stored in cookies and then it is sent in headers every time it makes a request to backend.

## How Email Works:

Application Layer : SMTP, POP3 → how to send & receive

Transfer : TCP protocol

Ex:

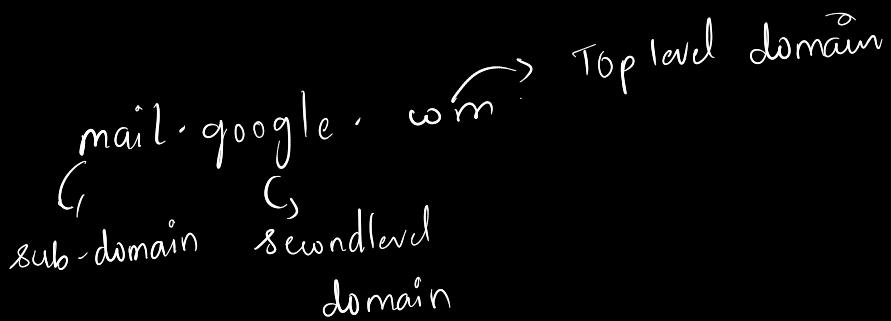


So basically gmail has two separate servers SMTP & POP

⇒ IMAP : Allows to use emails on multiple devices so the syncing thing on diff devices

## ⇒ DNS (Domain Name System):

Used to find the IP address associated with a URL.

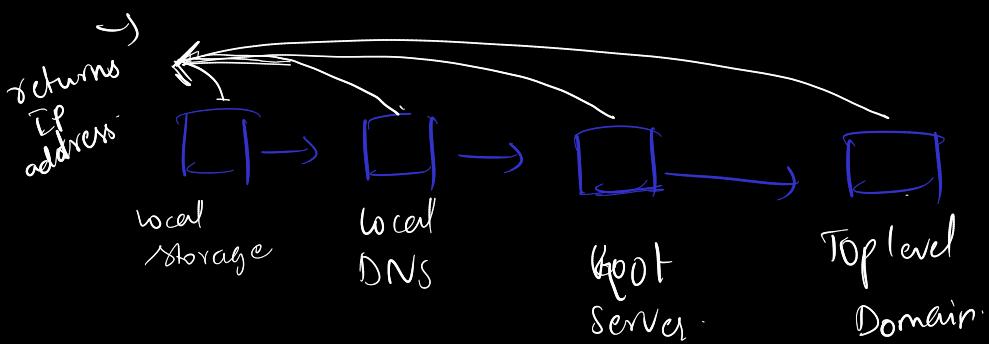


### \* Root DNS Servers (Top level domains)

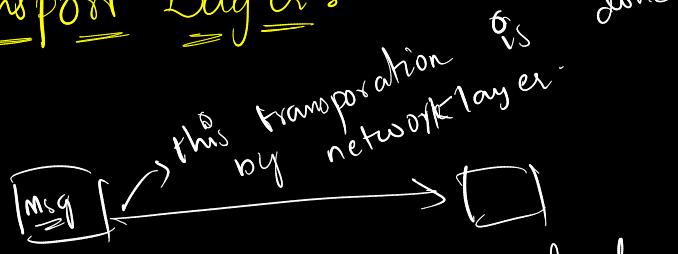


### IP Searching Process:

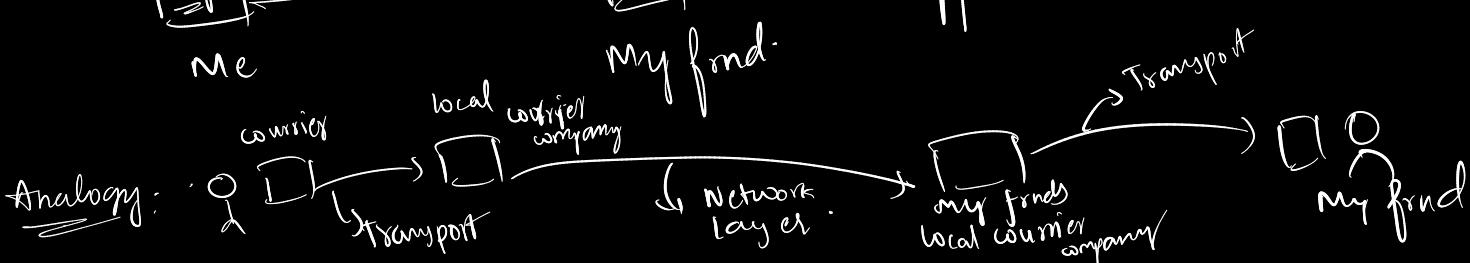
- First checks in the device storage
- Look it up in the local DNS server (TSP, Router)



## 2) Transport Layer:



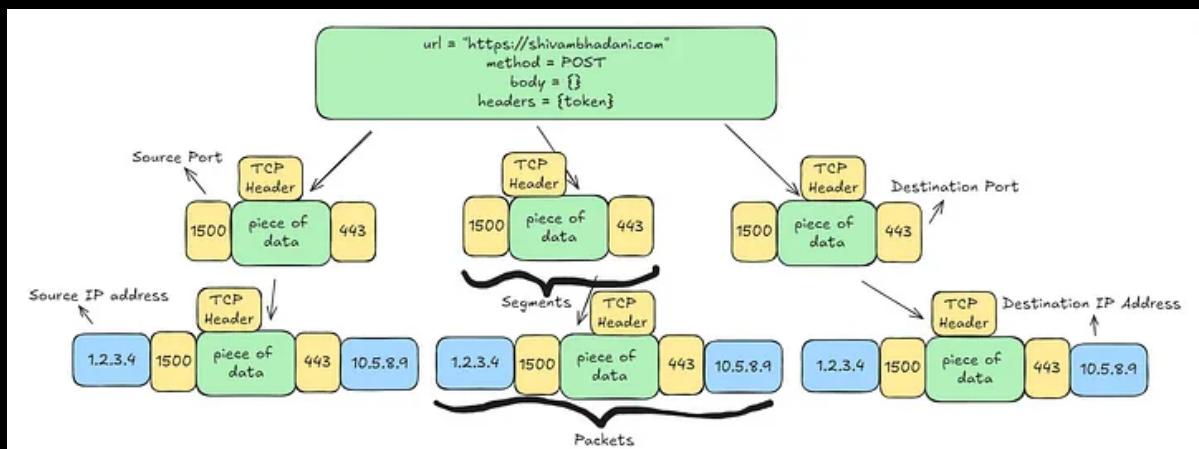
⇒ Transport layer just transports data from network layer to application within the computer



\* How `anios::post()` works under hood

- \* How `anios.post()` works ~~manually~~:
  - When we send the data along with URL it's first broken into small units → Y'cuz we can't transfer large data at one go through cables, wires or wifi
  - Each unit needs the receiver port and sender's port data for the transmission of data to be successful.

→ Each unit is called a Segment.



- Every packet sent by sender will be received. To make sure that happens we use acknowledgement. So kind of like replying I've received packet 1, packet 2, etc. and I have not received.

foo wox kbs

Note: TCP works on 3 way handshake

→ sending data over TCP has 3 steps:

i) Connection Establishment (3 way handshake)

ii) Transfer data

iii) Connection Termination (4 way handshake)

\* 3-way Handshake:

FLAG bits: SYN, ACK, SYN-ACK,  
FIN,

kind of switches like  
 $\rightarrow$  SYN = 1 (set)  
ACK = 0 (unset)

i) Client side sends SYN (synchronize) bit.

ii) Server side SYN-ACK (Acknowledge)

iii) Client side ACK

Analogy: Client: I want to establish a connection will initialise all my mugs to X.  
server: I acknowledge it. It will initialise all mug to Y.  
client: Accepted.

→ A random number (Seq) is generated at the start and from the messages are numbered by '+1'.

→ Y: random number → solves confusion on multiple connections.  
extra layer of security

\* 4 Way Handshake:

• client sends a FIN to indicate it wants to close connection

• server ⇒ ACK

• server ⇒ FIN

• Client ⇒ FIN

Congestion: → needed to learn: it might damage connections b/w comp  
When a network is overloaded by large data being transferred by multiple devices, then there might be packets loss, time delay, or the breakdown of entire network sometimes.

\* How does TCP handle it?

It uses some jargon (algo's) to figure out how much load can the network take and adjust the transmission speed accordingly

There are 4 phases in TCP congestion control:

- 1) Slow Phase
- 2) Congestion Avoidance
- 3) Fast retransmit
- 4) Fast Recovery

1) Slow Phase:

In this the sender sends the data by gradually increasing the amount of data until, he encounters congestion.

Congestion Window: TCP maintains a congestion window that tells the sender how much data he can send before ACK from receiver

- Initially cwnd is set to 1 MSS (Maximum Segment Size)
- It ↑es it for every ACK received.

## Congestion Avoidance:

- Once the cwnd reaches threshold value. It stops growing exponentially to avoid congestion. It increases linearly by a small amount.

## 2) Fast Retransmit:

- If the network  $\rightarrow$  congested. Then some packets will be lost. If the sender receives 3 ACKs (instead of 2 ACKs) that indicates a packet is lost.
- Instead of raising an error which would take more time, TCP retransmits the missing packet. This process  $\Rightarrow$  fast retransmit

## 3) Fast Recovery:

- To retransmit the missing packets the TCP does not go back to the slow start phase.
- Instead it cuts down the MSS to  $MSS/2 \Rightarrow$  Yes if linearly until senders start receiving normal ACKs.
- After that it goes back to normal congestion avoidance Mode.

## Sockets :

- socket is an endpoint for <sup>(sending/receiving)</sup> transferring data over the network. It is combined of
  - 1) IP address
  - 2) Port
  - 3) Transfer protocol (In web's it's TCP generally)
- what it does  $\Rightarrow$  attaching port (TCP segments) + IP addresses (packets)
  - + establishing 3 way handshake

Note: Socket  $\neq$  WebSocket

