

ECE 236B - HW 5

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T4.26, A12.6, A3.11, A3.13, T4.43,

Q5.6

T4.26 ^{Verify} $x^T x \leq yz, y \geq 0, z \geq 0$

if and only if $\| \begin{bmatrix} 2x \\ y-z \end{bmatrix} \|_2 \leq y+z, y \geq 0, z \geq 0$.

Suppose $x^T x \leq yz$

$$\begin{aligned} \left\| \begin{bmatrix} 2x \\ y-z \end{bmatrix} \right\|_2^2 &= \begin{bmatrix} 2x^T & y-z \end{bmatrix} \begin{bmatrix} 2x \\ y-z \end{bmatrix} = 4x^T x + y^2 - yz - zy + z^2 \\ &= 4x^T x + y^2 - 2yz + z^2 \\ &= 4x^T x + y^2 + z^2 - 2yz \quad (y, z \text{ scalars}) \\ &\leq 4yz - 2yz + y^2 + z^2 \\ &= 2yz + y^2 + z^2 \\ &= (y+z)^2 \end{aligned}$$

Taking the square root of both sides,

$\left\| \begin{bmatrix} 2x \\ y-z \end{bmatrix} \right\|_2 \leq y+z$. Conversely, if this inequality holds, we have that:

$$\begin{aligned} 4x^T x + y^2 + z^2 - 2yz &\leq (y+z)^2 = y^2 + z^2 + 2yz \\ \Rightarrow 4x^T x &\leq 4yz \Rightarrow x^T x \leq yz. \end{aligned}$$

(a.) maximize $\left(\sum_{i=1}^m 1/(a_i^T x + b_i) \right)^{-1}$, with
domain $\{x \mid Ax \geq b\}$, where a_i^T is the i th row of A .

Recall: Given a proper cone K , $x \preceq_K y \Leftrightarrow y-x \in K$.

Def: SOCP

$$\min C^T x$$

$$\text{s.t. } -(A_i x + b_i, c_i^T x + d_i) \preceq_{K_i} 0, i=1, \dots, m$$

$$\text{where } K_i = \{(y, t) \in \mathbb{R}^{n_i+1} \mid \|y\|_2 \leq t\}$$

$$\text{maximize } \left(\sum_{i=1}^m 1/(a_i^T x + b_i) \right)^{-1} \quad (1)$$

is equivalent to:

$$\begin{aligned} \max \quad & 1^T t \\ \text{s.t.} \quad & 1/(a_i^T x + b_i) \leq t_i \quad \forall i \in [1, m] \\ & t \geq 0. \end{aligned} \quad (2)$$



$$\begin{aligned} \min \quad & 1^T t \\ \text{s.t.} \quad & t_i (a_i^T x + b_i) \geq 1, \quad \forall i \in [1, m] \\ & t \geq 0 \end{aligned}$$

Let $\alpha = 1$, $y_i = a_i^T x + b_i$, $z_i = t_i$. Then,
 $t_i (a_i^T x + b_i) \geq 1 \Leftrightarrow \|[a_i^T x + b_i - t_i]\|_2 \leq a_i^T x + b_i + t_i$,
 $a_i^T x + b_i \geq 0$, $t_i \geq 0$, $\forall i \in [1, m]$.

Let $K_i = \{ y \mid \|y\|_2 \leq a_i^T x + b_i + t_i \}$.
 This gives: $-[a_i^T x + b_i - t_i] \leq_{K_i} 0$.

In summary, we have:

$$\begin{aligned} \min \quad & 1^T t \\ \text{s.t.} \quad & \|[a_i^T x + b_i - t_i]\|_2 \leq a_i^T x + b_i + t_i, \\ & a_i^T x + b_i \geq 0, \quad t_i \geq 0, \quad \forall i \in [1, m]. \end{aligned}$$

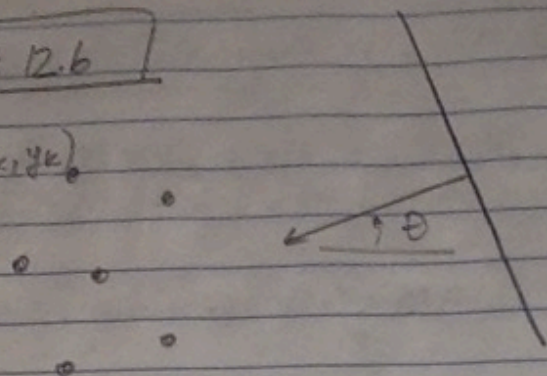


$$\begin{aligned} \min \quad & 1^T t \\ \text{s.t.} \quad & -[a_i^T x + b_i - t_i] \leq_{K_i} 0 \\ & a_i^T x + b_i \geq 0, \quad t_i \geq 0, \quad \forall i \in [1, m], \\ \text{where} \quad & K_i = \{ y \mid \|y\|_2 \leq a_i^T x + b_i + t_i \} \end{aligned}$$

③

A 12.6

(x_k, y_k)



$$G(\theta) = \sum_{k=1}^n w_k e^{i(x_k \cos \theta + y_k \sin \theta)}$$

$$= \sum_{k=1}^n [w_{r,k} \cos(\gamma_k(\theta)) - w_{i,k} \sin(\gamma_k(\theta))] + i [w_{r,k} \sin(\gamma_k(\theta)) + w_{i,k} \cos(\gamma_k(\theta))],$$

$$\gamma_k(\theta) = x_k \cos \theta + y_k \sin \theta, w_k = w_{r,k} + i w_{i,k}$$

$$\min_{w_k} \left(\max_n \{ |G(\theta_n)| \mid |\theta_n - \theta^{\text{tar}}| \geq \Delta \} \right).$$

$$(a.) \min t \quad -\Delta \leq \theta_n - \theta^{\text{tar}} \leq \Delta.$$

$$\text{s.t. } |G(\theta_n)| \leq t, \forall n.$$

$$\text{let } \gamma_{kn} = x_k \cos \theta_n + y_k \sin \theta_n.$$

$$|G(\theta_n)| = \left| \sum_k w_{r,k} \cos \gamma_{kn} - w_{i,k} \sin \gamma_{kn} + i \sum_k w_{r,k} \sin \gamma_{kn} + w_{i,k} \cos \gamma_{kn} \right| \leq t$$

$$\Leftrightarrow \left(\left(\sum_k w_{r,k} \cos \gamma_{kn} - w_{i,k} \sin \gamma_{kn} \right)^2 + \left(\sum_k w_{r,k} \sin \gamma_{kn} + w_{i,k} \cos \gamma_{kn} \right)^2 \right) \leq t^2$$

$$\Leftrightarrow \left\| \begin{pmatrix} \sum_k w_{r,k} \cos \gamma_{kn} - w_{i,k} \sin \gamma_{kn} \\ \sum_k w_{r,k} \sin \gamma_{kn} + w_{i,k} \cos \gamma_{kn} \end{pmatrix} \right\|_2 \leq t, \forall n.$$

12.6. Antenna array weight design.

We consider an array of n omnidirectional antennas in a plane, at positions (x_k, y_k) , $k = 1, \dots, n$.

A unit plane wave with frequency ω is incident from an angle θ . This incident wave induces in the k th antenna element a (complex) signal $\exp(i(x_k \cos \theta + y_k \sin \theta - \omega t))$, where $i = \sqrt{-1}$. (For simplicity we assume that the spatial units are normalized so that the wave number is one, i.e., the wavelength is $\lambda = 2\pi$.) This signal is demodulated, i.e., multiplied by $e^{i\omega t}$, to obtain the baseband signal (complex number) $\exp(i(x_k \cos \theta + y_k \sin \theta))$. The baseband signals of the n antennas are combined linearly to form the output of the antenna array

The complex weights in the linear combination, $w_k = w_{re,k} + iw_{im,k}$, $k=1, \dots, n$, are called the antenna array coefficients or shading coefficients, and will be the design variables in the problem. For a given set of weights, the combined output $G(\theta)$ is a function of the angle of arrival θ of the plane wave. The design problem is to select weights w_i that achieve a desired directional pattern $G(\theta)$.

We now describe a basic weight design problem. We require unit gain in a target direction θ_{tar} , i.e., $G(\theta_{tar}) = 1$. We want $|G(\theta)|$ small for $|\theta - \theta_{tar}| \geq \Delta$, where 2Δ is our beamwidth. To do this, we can minimize $\max_{|\theta - \theta_{tar}| \geq \Delta} |G(\theta)|$, where the maximum is over all $\theta \in [-\pi, \pi]$ with $|\theta - \theta_{tar}| \geq \Delta$. This number is called the sidelobe level for the array; our goal is to minimize the sidelobe level. If we achieve a small sidelobe level, then the array is relatively insensitive to signals arriving from directions more than Δ away from the target direction. This results in the optimization problem minimize $\max_{|\theta - \theta_{tar}| \geq \Delta} |G(\theta)|$ subject to $G(\theta_{tar}) = 1$, with $w \in \mathbb{C}^n$ as variables.

The objective function can be approximated by discretizing the angle of arrival with (say) N values (say, uniformly spaced) $\theta_1, \dots, \theta_N$ over the interval $[-\pi, \pi]$, and replacing the objective with $\max\{|G(\theta_k)| \mid |\theta_k - \theta_{tar}| \geq \Delta\}$

```
In [1]: import numpy as np
import cvxpy as cp
import matplotlib.pyplot as plt
```

```
In [73]: n = 40
N = 400
theta_tar = 15*(np.pi/180)
Delta = 15*(np.pi/180)

X = 30*np.random.uniform(low=0,high=1,size=n)
Y = 30*np.random.uniform(low=0,high=1,size=n)

# Theta = np.linspace(start=-np.pi,stop=np.pi,num=N)
# Exclude theta's less than delta from theta_tar:
ratio = np.round(np.abs(-180-15)/(np.abs(180-15)+np.abs(-180-15))*1000)/1000 # (Ratio of number of points)
Theta1 = np.linspace(start=-np.pi,stop=theta_tar-Delta,num=round(N*ratio))
Theta2 = np.linspace(start=theta_tar+Delta,stop=np.pi,num=round(N*(1-ratio)))
Theta = np.concatenate((Theta1,Theta2))
```

```

In [18]: ## Failed attempts:

# Gamma_kn = np.outer(X,np.cos(Theta)) + np.outer(Y,np.sin(Theta))
# gamma_tar = X*np.cos(theta_tar) + Y*np.sin(theta_tar)

# # gamma = np.matrix(Gamma_kn[:,1])
# # Gamma_kn_cos = np.cos(Gamma_kn)
# # Gamma_kn_sin = np.sin(Gamma_kn)

# # np.shape(np.cos(gamma))
# # np.dot(w,np.cos(gamma)) + np.dot(w,np.sin(gamma))
# # cos_gamma = np.matrix(Gamma_kn_cos[:,1])
# # np.shape(cos_gamma)

# w_re = cp.Variable(n)
# w_im = cp.Variable(n)
# t = cp.Variable(1)

# constr = []
# for i in np.arange(N):
#     theta = Theta[i]
#     gamma = X*np.cos(theta) + Y*np.sin(theta)
#     term1 = cp.sum(np.cos(gamma)*w_re - np.sin(gamma)*w_im)
#     term2 = cp.sum(np.sin(gamma)*w_re + np.cos(gamma)*w_im)
#     constr += [cp.norm(term1,term2) <= t]
#     constr += [theta - theta_tar <= Delta, -(theta - theta_tar) <= Delta]

# term1 = cp.sum(np.cos(gamma_tar)*w_re - np.sin(gamma_tar)*w_im)
# term2 = cp.sum(np.sin(gamma_tar)*w_re + np.cos(gamma_tar)*w_im)
# constr += [term1 == 1,term2 == 0]
# problem = cp.Problem(cp.Minimize(t), constr)
# problem.solve()

## -----

# w_re = cp.Variable(n)
# w_im = cp.Variable(n)
# t = cp.Variable(1)

# constr = []
# for i in np.arange(N):

```

```
# theta = Theta[i]
# gamma = np.matrix(Gamma_kn[:,i])
# cos_gamma = Gamma_kn_cos[:,i].T
# sin_gamma = Gamma_kn_sin[:,i].T
# term1 = cp.sum(cp.multiply(cos_gamma,w_re) - cp.multiply(sin_gamma,w_im))
# term2 = cp.sum(cp.multiply(sin_gamma,w_re) + cp.multiply(cos_gamma,w_im))

# constr += [cp.sum(cp.multiply(term1,term1)) + cp.sum(cp.multiply(term2,term2)) <= t**2]
# constr += [theta - theta_tar <= Delta]
# constr += [-(theta - theta_tar) <= Delta]

# term1 = cp.sum(cp.multiply(np.cos(gamma_tar),w_re) - cp.multiply(np.sin(gamma_tar),w_im))
# term2 = cp.sum(cp.multiply(np.sin(gamma_tar),w_re) + cp.multiply(np.cos(gamma_tar),w_im))
# constr += [term1 == 1,term2 == 0]
# problem = cp.Problem(cp.Minimize(t), constr)
# problem.solve(solver=cp.ECOS)
```

In [*]: *# Attempt 3:*

```
w_re = cp.Variable(n)
w_im = cp.Variable(n)
t = cp.Variable(1,pos=True)

constr = []
for i in np.arange(N):
    theta = Theta[i]
    gamma = X*np.cos(theta) + Y*np.sin(theta)
    term1 = 0
    term2 = 0
    for j in np.arange(n):
        term1 += np.cos(gamma[j])*w_re[j] - np.sin(gamma[j])*w_im[j]
        term2 += np.sin(gamma[j])*w_re[j] + np.cos(gamma[j])*w_im[j]

    constr += [cp.SOC(term1 + term2,t)] # I think this should actually be reversed, but this doesn't wo

term1 = 0
term2 = 0
for j in np.arange(n):
    term1 += np.cos(gamma_tar[j])*w_re[j] - np.sin(gamma_tar[j])*w_im[j]
    term2 += np.sin(gamma_tar[j])*w_re[j] + np.cos(gamma_tar[j])*w_im[j]
constr += [term1 == 1,term2 == 0]

problem = cp.Problem(cp.Minimize(t), constr)
problem.solve()
```

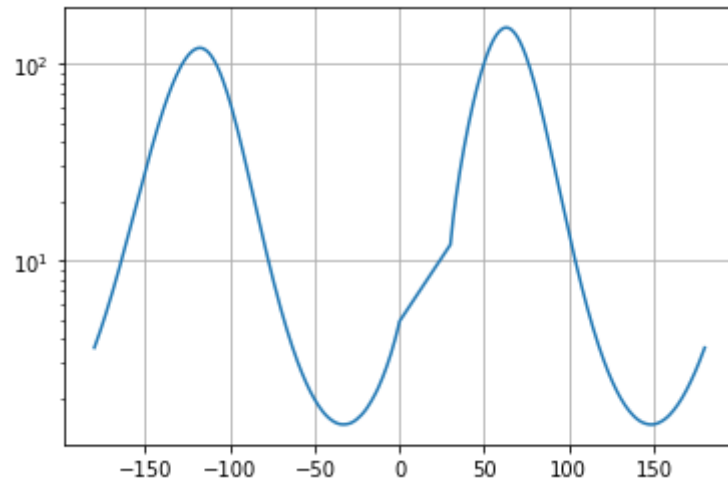
In [95]:

```
w_re_opt = w_re.value
w_im_opt = w_im.value

G = np.zeros(N)
for i in np.arange(N):
    theta = Theta[i]
    gamma = X*np.cos(theta) + Y*np.sin(theta)
    term1 = 0
    term2 = 0
    for j in np.arange(n):
        term1 += np.cos(gamma[j])*w_re_opt[j] - np.sin(gamma[j])*w_im_opt[j]
        term2 += np.sin(gamma[j])*w_re_opt[j] + np.cos(gamma[j])*w_im_opt[j]
    G[i] = np.sqrt(term1**2+term2**2)
```



```
In [99]: plt.plot(Theta*(180/np.pi),G)
plt.grid()
plt.yscale('log')
plt.show()
```



```
In [ ]: # (Obviously wrong)
```

④

Formulate as SDP:

A3.11 (a, b, c)

minimize $C^T x$

s.t. $x_1 F_1 + \dots + x_n F_n + G \leq 0$

$$F(x) = F_0 + x_1 F_1 + \dots + x_n F_n, \quad x \in \mathbb{R}^n, \quad F_i \in S^m.$$

$$\text{dom}(f) = \{x \in \mathbb{R}^n \mid F(x) \succ 0\}.$$

(a.) minimize $f(x) = C^T F(x)^{-1} C$, where $C \in \mathbb{R}^m$.

$$\begin{aligned} &\Downarrow \\ \min \quad &t \\ \text{s.t.} \quad &C^T F(x)^{-1} C \leq t \\ &F(x) \succ 0 \end{aligned}$$

$$\begin{aligned} &\Updownarrow \\ \min \quad &t \\ \text{s.t.} \quad &\begin{bmatrix} F(x) & C \\ C^T & t \end{bmatrix} \succeq 0 \\ &F(x) \succ 0 \end{aligned}$$

$$S = t - C^T F^{-1} C \geq 0$$

$$\Leftrightarrow C^T F^{-1} C \leq t.$$

$$-\begin{bmatrix} F(x) & C \\ C^T & t \end{bmatrix} = \begin{bmatrix} -F_0 - C & \\ -C^T & 0 \end{bmatrix} + x_1 \begin{bmatrix} -F_1 & 0 \\ 0 & 0 \end{bmatrix} + \dots + x_n \begin{bmatrix} -F_n & 0 \\ 0 & 0 \end{bmatrix} + t \begin{bmatrix} 0 & 0 \\ 0 & -1 \end{bmatrix} \leq 0$$

$$\begin{aligned} &\text{minimize } t \\ &\text{s.t.} \\ &\begin{bmatrix} -F_0 - C & \\ -C^T & 0 \end{bmatrix} + x_1 \begin{bmatrix} -F_1 & 0 \\ 0 & 0 \end{bmatrix} + \dots + x_n \begin{bmatrix} -F_n & 0 \\ 0 & 0 \end{bmatrix} + \begin{bmatrix} 0 & 0 \\ 0 & -1 \end{bmatrix} \leq 0 \\ &F(x) \succ 0 \end{aligned}$$

(5)

(b.) Minimize $f(x) = \max_{i=1, \dots, k} c_i^T F(x)^{-1} c_i$,where $c_i \in \mathbb{R}^m$, $i = 1, \dots, k$. \Updownarrow min t s.t. $\max_i c_i^T F(x)^{-1} c_i \leq t$ \Updownarrow min t s.t. $c_i^T F(x)^{-1} c_i \leq t, \forall i$ \Updownarrow min t s.t. $\begin{bmatrix} F(x) & c_i \\ c_i^T & t \end{bmatrix} \geq 0, \forall i$

Now $A \succ 0$ and $B \succ 0$ if and only if $\begin{bmatrix} A & 0 \\ 0 & B \end{bmatrix} \succ 0$. So we can write:

min t s.t. $\begin{bmatrix} A_i & 0 \\ 0 & t \end{bmatrix} \geq 0$ where $A_i = \begin{bmatrix} F(x) & c_i \\ c_i^T & t \end{bmatrix}$.

6

$$(c.) \text{ minimize } f(x) = \sup_{\|c\|_2 \leq 1} c^T F(x)^{-1} c$$

\Downarrow

$$\begin{aligned} \min t \\ \text{s.t. } \sup_{\|c\|_2 \leq 1} c^T F(x)^{-1} c \leq t \end{aligned}$$

\Downarrow

$$\begin{aligned} \min t \\ \text{s.t. } c^T F(x)^{-1} c \leq t \\ \text{for } \forall c \text{ s.t. } \|c\|_2 \leq 1. \end{aligned}$$

\Downarrow

$$\begin{aligned} \min t \\ \text{s.t. } \begin{bmatrix} F(x) & c \\ c^T & t \end{bmatrix} \geq 0 \\ \|c\|_2 \leq 1 \end{aligned}$$

$$\begin{aligned} \begin{bmatrix} F(x) & c \\ c^T & t \end{bmatrix} &= \begin{bmatrix} F_0 & 0 \\ 0 & 0 \end{bmatrix} + x_1 \begin{bmatrix} F_1 & 0 \\ 0 & 0 \end{bmatrix} + \dots + x_n \begin{bmatrix} F_n & 0 \\ 0 & 0 \end{bmatrix} + t \begin{bmatrix} 0 & 0 \\ 0 & 1 \end{bmatrix} \\ &+ c_1 \begin{bmatrix} 0 & \begin{bmatrix} 1 \\ 0 \\ \vdots \\ 0 \end{bmatrix} \\ \begin{bmatrix} 0 & \dots & 0 \end{bmatrix} & 0 \end{bmatrix} + \dots + c_n \begin{bmatrix} 0 & \begin{bmatrix} 0 \\ \vdots \\ 1 \end{bmatrix} \\ \begin{bmatrix} 0 & \dots & 0 \end{bmatrix} & 0 \end{bmatrix} \end{aligned}$$

$$\text{Let } X = \begin{bmatrix} x \\ c \\ t \end{bmatrix} \text{ and } C = \begin{bmatrix} 0 \\ 0 \\ 1 \end{bmatrix}, \text{ and } D = \begin{bmatrix} 0 \\ 1 \\ 0 \end{bmatrix}$$

We can then write:

$$\begin{aligned} \min C^T X \\ \text{s.t. } \begin{bmatrix} F(x) & c \\ c^T & t \end{bmatrix} \geq 0 \\ \|D^T X\|_2 \leq 1. \end{aligned}$$

But this isn't an SDP because of the norm constraint.

⑦

$$(c.) \text{ minimize } f(x) = \sup_{\|c\|_2 \leq 1} c^T F(x)^{-1} c$$

$$f(x) = \lambda_{\max}(F(x)^{-1})$$

$$\min t$$

$$\text{s.t. } \lambda_{\max}(F(x)^{-1}) \leq t$$

$$\lambda_{\max}(F(x)^{-1}) \leq t \Rightarrow (F(x)^{-1} - tI)c \leq 0 \quad \forall c$$

$$\Rightarrow c^T F(x)^{-1} c \leq t I c^T c \leq t \quad (\text{since } c^T c \leq 1).$$

$$(\text{since } F(x) > 0 \Rightarrow F(x)^{-1} > 0)$$

$$\Rightarrow c^T (F(x)^{-1} - tI) c \leq 0.$$

$$e_i^T F(x)^{-1} e_i \leq t \quad \forall i$$

$$\Rightarrow I^T F(x)^{-1} I \leq tI$$

$$\Rightarrow \begin{bmatrix} F(x) & I \\ I & tI \end{bmatrix} \succeq 0$$

We have:

$$\min t$$

$$\text{s.t. } \begin{bmatrix} F(x) & I \\ I & tI \end{bmatrix} \succeq 0$$

8

$$\boxed{A3.13} \quad H(A, B) = 2(A^{-1} + B^{-1})^{-1}.$$

Show that $X = \frac{1}{2} H(A, B)$ solves the SDP.

$$\max \operatorname{tr} X$$

$$\text{s.t. } \begin{bmatrix} X & X \\ X & X \end{bmatrix} \preceq \begin{bmatrix} A & 0 \\ 0 & B \end{bmatrix},$$

$$X \in S^n, A \in S_{++}^n, B \in S_{++}^n.$$

$$\text{Let } R = \begin{bmatrix} A^{-1} & I \\ B^{-1} & -I \end{bmatrix}.$$

$$A, B \in S_{++}^n \Rightarrow A^{-1}, B^{-1} \in S_{++}^n$$

$$\Rightarrow \operatorname{rank}(A^{-1}) = \operatorname{rank}(B^{-1}) = n.$$

We need to show that every row of $[A^{-1} \ I]$ and $[B^{-1} \ -I]$ are linearly independent.

$$\text{Suppose } v \in \operatorname{Null}(R), \quad v = \begin{bmatrix} v_1 \\ v_2 \end{bmatrix}.$$

$$\Rightarrow \begin{bmatrix} A^{-1} & I \\ B^{-1} & -I \end{bmatrix} \begin{bmatrix} v_1 \\ v_2 \end{bmatrix} = 0 \Rightarrow \begin{bmatrix} A^{-1}v_1 + Iv_2 \\ B^{-1}v_1 - Iv_2 \end{bmatrix} = \begin{bmatrix} 0 \\ 0 \end{bmatrix}.$$

Suppose row i of $[A^{-1} \ I]$ is linearly dependent on row i of $[B^{-1} \ -I]$.

$$\Rightarrow -[A^{-1} \ I]_i = [B^{-1} \ -I]_i \Rightarrow -A_i^{-1} = B_i^{-1}$$

$$\Rightarrow A_i^{-1} + B_i^{-1} = 0 \text{ for some } i.$$

$$\Rightarrow (v^T (A^{-1} + B^{-1}) v)_i = (v^T A^{-1} v)_i + (v^T B^{-1} v)_i = 0$$

$\Rightarrow A \leq 0$ or $B^{-1} \leq 0$, a contradiction.

9

Thus, $[A^{-1} \ I]$ and $[B^{-1} \ -I]$ are linearly independent and hence $\text{rank}(R) = n$, so R is nonsingular. Thus,

$$R^T \begin{bmatrix} X & X \\ X & X \end{bmatrix} R \preceq R^T \begin{bmatrix} A & 0 \\ 0 & B \end{bmatrix} R$$

$$\begin{bmatrix} X & X \\ X & X \end{bmatrix} \begin{bmatrix} A^{-1} & I \\ B^{-1} & -I \end{bmatrix} = \begin{bmatrix} X(A^{-1} + B^{-1}) & 0 \\ X(A^{-1} + B^{-1}) & 0 \end{bmatrix}$$

Note: $B^{-T} = B^{-1}$.

$$X(A+B)^{-1} \begin{bmatrix} A^{-1} & B^{-1} \\ 0 & -I \end{bmatrix} \begin{bmatrix} I & 0 \\ I & 0 \end{bmatrix} = X(A^{-1} + B^{-1}) \begin{bmatrix} A^{-1} + B^{-1} & 0 \\ -I & 0 \end{bmatrix}$$

$$\begin{bmatrix} A & 0 \\ 0 & B \end{bmatrix} \begin{bmatrix} A^{-1} & I \\ B^{-1} & -I \end{bmatrix} = \begin{bmatrix} I & A \\ I & -B \end{bmatrix}$$

$$\begin{bmatrix} A^{-1} & B^{-1} \\ 0 & -I \end{bmatrix} \begin{bmatrix} I & A \\ I & -B \end{bmatrix} = \begin{bmatrix} A^{-1} + B^{-1} & 0 \\ -I & B \end{bmatrix}$$

So we have:

$$X(A^{-1} + B^{-1}) \begin{bmatrix} A^{-1} + B^{-1} & 0 \\ -I & 0 \end{bmatrix} \preceq \begin{bmatrix} A^{-1} + B^{-1} & 0 \\ -I & B \end{bmatrix}$$

Letting $X = \frac{1}{2} H(A, B) = (A^{-1} + B^{-1})^{-1}$, we have:

$$\begin{bmatrix} A^{-1} + B^{-1} & 0 \\ -I & 0 \end{bmatrix} \preceq \begin{bmatrix} A^{-1} + B^{-1} & 0 \\ -I & B \end{bmatrix},$$

which is true since $B \succeq 0$.

The constraint is satisfied with equality for all other elements, which implies that X is optimal.

(10)

Suppose $\exists Y$ s.t. $\text{tr}(Y) > \text{tr}(X)$
 and $\begin{bmatrix} Y & Y \\ Y & Y \end{bmatrix} \preceq \begin{bmatrix} A & 0 \\ 0 & B \end{bmatrix}$.

wlog, let $Y = X + Q_i$ where $Q_i = \begin{bmatrix} 0 & & \\ & \epsilon & \\ & & 0 \end{bmatrix}$,

That is, any matrix with $\text{tr}(Q_i) = \epsilon$
 where $\epsilon > 0$, so that $\text{tr}(Y) = \text{tr}(X) + \epsilon$.

Then, we have that

$$Y(A^{-1} + B^{-1}) = I + \epsilon(A^{-1} + B^{-1}) \Rightarrow$$

$$\begin{bmatrix} A^{-1} + B^{-1} & 0 \\ -I & 0 \end{bmatrix} + \epsilon \begin{bmatrix} (A^{-1} + B^{-1})^2 & 0 \\ -(A^{-1} + B^{-1}) & 0 \end{bmatrix} \preceq \begin{bmatrix} A^{-1} + B^{-1} & 0 \\ -I & B \end{bmatrix}$$

$$\Leftrightarrow \epsilon \begin{bmatrix} (A^{-1} + B^{-1})^2 & 0 \\ -(A^{-1} + B^{-1}) & 0 \end{bmatrix} \preceq \begin{bmatrix} 0 & 0 \\ 0 & B \end{bmatrix}$$

But $\epsilon(A^{-1} + B^{-1})^2 \not\preceq 0$ for $\forall \epsilon > 0$
 since $A^{-1}, B^{-1} \in S_{++}^n$, a contradiction.

Thus, X is the solution to the SDP.

It follows that $\text{tr}((A^{-1} + B^{-1})^{-1})$ is
 a concave function of (A, B)
 because $\text{tr}(X)$ is a concave function of
 X for $\forall X$ satisfying the given
 linear matrix inequality.

11

T 4.43 (b,c)

Suppose $A: \mathbb{R}^n \rightarrow \mathbb{S}^m$ is affine.

$$A(x) = A_0 + x_1 A_1 + \dots + x_n A_n, \quad A_i \in \mathbb{S}^m$$

Let $\lambda_1(x) \geq \lambda_2(x) \geq \dots \geq \lambda_m(x) = \text{eigen}(A)$.
Pose as SDP:

$$(b.) \quad \min (\lambda_1(x) - \lambda_m(x))$$

$$\lambda_{\max} = \sup_{\|C\|_2 \leq 1} C^T A(x) C \quad \text{But no simple expression for } \lambda_{\min}(x).$$

$$\begin{aligned} & \min_x (\lambda_1(x) - \lambda_m(x)) \\ &= \min_x \left(\max_{i,j} (\lambda_i(x) - \lambda_j(x)) \right) \end{aligned}$$

$$\begin{aligned} \Leftrightarrow & \min_{x,t} t \\ & \text{s.t. } \max_{i,j} (\lambda_i(x) - \lambda_j(x)) \leq t \end{aligned}$$

$$\begin{aligned} \Leftrightarrow & \min_{x,t} t \\ & \text{s.t. } \lambda_i(x) - \lambda_j(x) \leq t, \quad \forall i,j \end{aligned}$$

Let v_i be such that $A(x)v_i = \lambda_i v_i$

$$\lambda_i(x)v_i - \lambda_j(x)v_j = A(x)v_i - A(x)v_j = A(x)(v_i - v_j)$$

Consider $A(x) - I\lambda_m$. $(A(x) - I\lambda_m)v_m = 0$

$$\begin{aligned} & \min t \\ & \text{s.t. } \lambda_{\max} \leq t + \lambda_{\min} \end{aligned}$$

(12)

T 4.43 (b) (contd).

Consider $\lambda I - A(x)$. Let $A(x) = U^T \Lambda U$, $U^T U = I$
 $\Rightarrow \lambda I - A = U^T (\lambda I - \Lambda) U$, $\lambda I - A \geq 0 \Leftrightarrow \lambda I - \Lambda \geq 0$
 $\lambda I - \Lambda \geq 0 \Leftrightarrow \lambda \geq \lambda_{\max}(\Lambda) = \lambda_{\max}(A)$
 $\Rightarrow \lambda_{\max}(A) = \inf_{\lambda I - A \geq 0} \lambda \Leftrightarrow \lambda_{\max}(A) = -\sup_{\lambda I - A \leq 0} (-\lambda)$

Consider $A(x) - \lambda' I = U^T (\Lambda - \lambda' I) U$
 $\Lambda - \lambda' I \geq 0 \Leftrightarrow \lambda_{\min}(A) = \lambda_{\min}(\Lambda) \geq \lambda'$
 $\Rightarrow \lambda_{\min} = \sup_{A - \lambda' I \geq 0} \lambda' \Leftrightarrow \lambda_{\min} = -\inf_{\lambda' I - A \leq 0} (-\lambda')$

Thus, $\min_x (\lambda_{\max}(A) - \lambda_{\min}(A)) = \min_x \inf_{\lambda, \lambda'} \lambda - \lambda'$
 $\lambda' I - A \leq 0 \leq \lambda I - A$

$$= \min_{\lambda, \lambda'} \lambda - \lambda' \quad \text{s.t.} \quad \begin{cases} \lambda' I - A(x) \leq 0 \\ \lambda I - A(x) \geq 0 \end{cases}$$

(c) $\kappa(A(x)) = \lambda_1(x) / \lambda_m(x)$, $\text{dom}\{x \mid A(x) \succ 0\}$.
 $A(x) \succ 0$ for at least one x .

Make change of variables $y = x/\gamma$, $t = \lambda/\gamma$, $s = 1/\gamma$. ($\Rightarrow y = sx$)
 $\min \lambda_1(x) / \lambda_m(x) \quad \text{s.t.} \quad A(x) \geq 0$

$$\min \lambda / \gamma \quad \Leftrightarrow \quad \min t$$

s.t. $\gamma I - A \leq 0 \leq \lambda I - A$ s.t. $I - sA \leq 0 \leq tI - sA$

$$\Leftrightarrow \min t \quad \Leftrightarrow \min t$$

s.t. $I \leq sA \leq tI$, $s > 0$ s.t. $I \leq s(A_0 + x_1 A_1 + \dots + x_n A_n) \leq tI$, $s > 0$

$$\Leftrightarrow \min t$$

s.t. $I \leq sA_0 + y_1 A_1 + \dots + y_n A_n \leq tI$, $-I \leq 0$

(13)

There are 3 constraints, so we can rewrite them using 3×3 matrices.

$$s \begin{bmatrix} -I & 0 & 0 \\ 0 & A_0 & 0 \\ 0 & 0 & -A_0 \end{bmatrix} + y_1 \begin{bmatrix} 0 & A_1 & 0 \\ 0 & 0 & 0 \\ 0 & 0 & -A_1 \end{bmatrix} + \dots + y_n \begin{bmatrix} 0 & A_n & 0 \\ 0 & 0 & 0 \\ 0 & 0 & -A_n \end{bmatrix} \\ + \begin{bmatrix} 0 & 0 & 0 \\ 0 & 0 & 0 \\ 0 & 0 & I \end{bmatrix} + t \begin{bmatrix} 0 & 0 & 0 \\ 0 & -I & 0 \\ 0 & 0 & 0 \end{bmatrix} \leq \underline{0}$$

The problem is now an SDP.

(14)

6. GP: $\min f_0(x)$ s.t. $f_i(x) \leq 1, h_i(x) = 1$,
 f_0, \dots, f_m posynomial & h_1, \dots, h_p monomial.
 That is, $h_i(x) = c x_1^{a_1} x_2^{a_2} \dots x_n^{a_n}$, $c > 0, a_i \in \mathbb{R}$,
 $f_i(x) = \sum c_k x_1^{a_{1k}} \dots x_n^{a_{nk}}$, $c_k > 0$.

$$\text{minimize } (A + \mu T = \sum w_i + \mu \max(T_i))$$

$$\text{s.t. } T_1 = \rho k_0 \left[(w_3 + c_{e1}) \left(\frac{1}{w_1} + \frac{1}{w_2} + \frac{1}{w_3} \right) + w_2 \left(\frac{1}{w_1} + \frac{1}{w_2} \right) + (w_1 + w_4 + w_5 + w_6 + c_{e2} + c_{e3}) \frac{1}{w_1} \right]$$

$$T_2 = \rho k_0 \left[(w_5 + c_{e2}) \left(\frac{1}{w_1} + \frac{1}{w_4} + \frac{1}{w_5} \right) + w_4 \left(\frac{1}{w_1} + \frac{1}{w_4} \right) + (w_6 + c_{e3}) \left(\frac{1}{w_1} + \frac{1}{w_1} \right) + (w_1 + w_2 + w_3 + c_{e1}) \frac{1}{w_1} \right]$$

$$T_3 = \rho k_0 \left[(w_6 + c_{e2}) \left(\frac{1}{w_1} + \frac{1}{w_4} + \frac{1}{w_6} \right) + w_4 \left(\frac{1}{w_1} + \frac{1}{w_4} \right) + (w_1 + w_2 + w_3 + c_{e1}) \frac{1}{w_1} + (w_5 + c_{e2}) \left(\frac{1}{w_1} + \frac{1}{w_4} \right) \right]$$

$$T_i = \sum_j C_i R_{ij} = \rho k_0 \sum_j w_i R_{ij}$$

$$\min (\sum w_i + \mu \max(T_i))$$

$$\Leftrightarrow \min (\sum w_i + \mu t) \quad \Leftrightarrow \min (\sum w_i + \mu t)$$

$$\text{s.t. } \max(T_i) \leq t \quad \text{s.t. } -T_i \leq t, \forall i.$$

$$T_i \leq t \Leftrightarrow \frac{1}{t} T_i \leq 1 \text{ since } t > 0.$$

Each T_i restraint is clearly a posynomial inequality of the form $f_i(x) \leq 1$. The objective is also a posynomial. Thus, the scalarized problem is a GP.

6. Interconnect sizing. We consider the problem of sizing the interconnecting wires of the simple circuit shown below, in which one voltage source drives three different capacitive loads Cload1, Cload2, and Cload3.

We divide the wires into 6 segments of fixed length l_i ; the optimization variables in the problem will be the widths w_i of the segments. (The height of the wires is related to the particular integrated circuit technology process, and is fixed.) We take the lengths l_i to be one, for simplicity.

In the next figure each of the wire segments is modeled by a simple RC circuit, with the resistance inversely proportional to the width of the segment and the capacitance proportional to the width.

The capacitance and resistance of the i th segment is thus $C_i = k_0 w_i$, $R_i = \rho / w_i$, $i = 1, \dots, 6$, where k_0 and ρ are positive constants, which we take to be one for simplicity. We also take $C_{load1} = 1.5$, $C_{load2} = 1$, and $C_{load3} = 5$.

We are interested in the trade-off between area and delay. The total area used by the wires is the sum of the w_i 's

We use the Elmore delay to model the delay from the source to each of the loads. The Elmore delays to loads 1, 2, and 3 are:

$$T_1 = (C_3 + C_{load1})(R_1 + R_2 + R_3) + C_2(R_1 + R_2) + (C_1 + C_4 + C_5 + C_6 + C_{load2} + C_{load3})R_1$$

$$T_2 = (C_5 + C_{load2})(R_1 + R_4 + R_5) + C_4(R_1 + R_4) + (C_6 + C_{load3})(R_1 + R_4) + (C_1 + C_2 + C_3 + C_{load1})R_1$$

$$T_3 = (C_6 + C_{load3})(R_1 + R_4 + R_6) + C_4(R_1 + R_4) + (C_1 + C_2 + C_3 + C_{load1})R_1 + (C_5 + C_{load2})(R_1 + R_4)$$

(The general rule is as follows: the Elmore delay from the source to node j is given by Sum of $C_i R_{ij}$ all nodes i where C_i is the capacitance at node i and R_{ij} is the sum of the resistances on the intersection of the path from the source to node i and the path from the source to node j .) Our main interest is in the maximum of these delays, $T = \max \{T_1, T_2, T_3\}$.

We also impose minimum and maximum allowable values for the wire widths: $W_{min} \leq w_i \leq W_{max}$, $i = 1, \dots, 6$. For our specific problem, we take $W_{min} = 0.1$ and $W_{max} = 10$. We compare two choices of wire widths.

```
In [34]: import numpy as np
import cvxpy as cp
import matplotlib.pyplot as plt
```

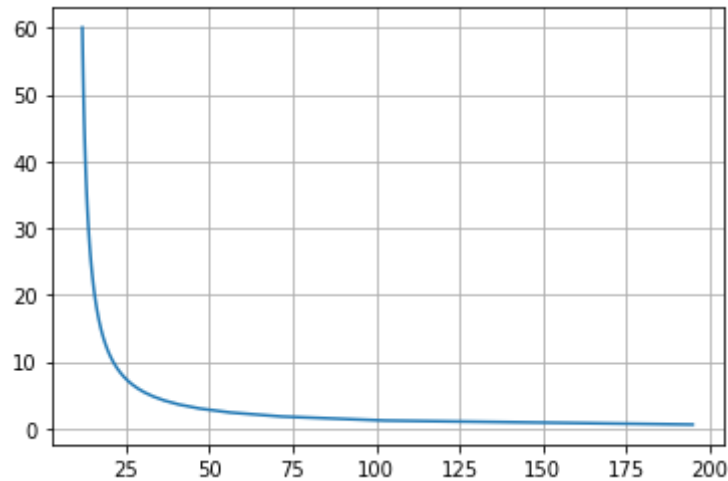
```
In [118]: Cload1 = 1.5
          Cload2 = 1
          Cload3 = 5
          W_min = 0.1
          W_max = 10
          N = 100
          w_v = np.linspace(start=W_min,stop=W_max,num=N)
          A_v = 6*w_v
```

(a) Equal wire widths. Plot the values of area A versus delay T , obtained if you take equal wire widths w_i (varying between W_{min} and W_{max}).

```
In [119]: T_v = np.zeros(N)
          for i in np.arange(N):
              w = w_v[i]
              C1 = C2 = C3 = C4 = C5 = C6 = w
              R1 = R2 = R3 = R4 = R5 = R6 = 1/w
              T1 = (C3 + Cload1)*(R1 + R2 + R3) + C2*(R1 + R2) + (C1 + C4 + C5 + C6 + Cload2 + Cload3)*R1
              T2 = (C5 + Cload2)*(R1 + R4 + R5) + C4*(R1 + R4) + (C6 + Cload3)*(R1 + R4) + (C1 + C2 + C3 + Cload1)
              T3 = (C6 + Cload3)*(R1 + R4 + R6) + C4*(R1 + R4) + (C1 + C2 + C3 + Cload1)*R1 + (C5 + Cload2)*(R1 +
              T = np.max([T1,T2,T3])
              T_v[i] = T
```

Area vs Delay:

```
In [120]: plt.plot(T_v,A_v)
plt.grid()
plt.show()
```



(b) Optimal wire widths. The optimal area-delay trade-off curve can be computed by scalarization, i.e., by minimizing $A + \mu T$, subject to the constraints on w , for a large number of different positive values of μ . Verify that the scalarized problem is a geometric program (GP). For the specific problem parameters given, compute the area-delay trade-off curve using CVX or CVXPY. You can choose the values of μ logarithmically spaced between 10^{-3} and 10^3 . Compare the optimal trade-off curve with the one obtained in part (a).

Consult chapter 7 of the CVX user guide for details on how to solve GPs. For reasons explained in the user guide, CVX is not very fast when solving GPs. If needed, you can limit the number of weights μ , for example, to 10 or 20.

In [86]: *# Set up the problem:*

```

N = 21
mu_v = np.geomspace(start=1e-3, stop=1e3, num=N)
w = cp.Variable(shape=(6,), pos=True, name="w")
t = cp.Variable(1, pos=True)

T1 = (w[2] + Cload1)*(w[0]**-1 + w[1]**-1 + w[2]**-1) + w[1]*(w[0]**-1 + w[1]**-1) + (w[0]+w[3]+w[4]+w[5])
T2 = (w[4]+Cload2)*(w[0]+w[3]+w[4]) + w[3]*(w[0]**-1+w[3]**-1) + (w[5] + Cload3)*(w[0]**-1 + w[3]**-1) +
T3 = (w[5]+Cload3)*(w[0]**-1+w[3]**-1+w[5]**-1) + w[3]*(w[0]**-1+w[3]**-1) + (w[0]+w[1]+w[2]+Cload1)*w[0]
T1 /= t
T2 /= t
T3 /= t
print("T1:", T1.log_log_curvature)
print("T2:", T2.log_log_curvature)
print("T3:", T3.log_log_curvature)

```

T1: LOG-LOG CONVEX

T2: LOG-LOG CONVEX

T3: LOG-LOG CONVEX

In [87]: *# Solve a single instance:*

```
mu = 1
objective_fn = cp.sum(w) + mu*t
constraints = [T1 <= 1, T2 <= 1, T3 <= 1]
assert objective_fn.is_log_log_convex()
assert all(constraint.is_dgp() for constraint in constraints)
problem = cp.Problem(cp.Minimize(objective_fn), constraints)

print(problem)
print("Is this problem DGP?", problem.is_dgp())

minimize Sum(w, None, False) + 1.0 * var151138
subject to ((w[2] + 1.5) * (power(w[0], -1) + power(w[1], -1) + power(w[2], -1)) + w[1] * (power(w[0],
-1) + power(w[1], -1)) + (w[0] + w[3] + w[4] + w[5] + 1.0 + 5.0) * power(w[0], -1)) / var151138 <= 1.0
            ((w[4] + 1.0) * (w[0] + w[3] + w[4]) + w[3] * (power(w[0], -1) + power(w[3], -1)) + (w[5] +
5.0) * (power(w[0], -1) + power(w[3], -1)) + (w[0] + w[1] + w[2] + 1.5) * power(w[0], -1)) / var151138
<= 1.0
            ((w[5] + 5.0) * (power(w[0], -1) + power(w[3], -1) + power(w[5], -1)) + w[3] * (power(w[0], -
-1) + power(w[3], -1)) + (w[0] + w[1] + w[2] + 1.5) * power(w[0], -1) + (w[4] + 1.0) * (power(w[0], -
1) + power(w[3], -1))) / var151138 <= 1.0
Is this problem DGP? True
```

In [88]:

```
problem.solve(gp=True)
print("Optimal value:", problem.value)
print(w, ":", w.value)
print(t, ":", t.value)
print("Dual values: ", list(c.dual_value for c in constraints))
```

```
Optimal value: 21.486547050628886
w : [3.26297989e+00 5.36887644e-01 4.19543272e-01 2.03482520e+00
     3.65756051e-08 1.42861434e+00]
var151138 : [13.80369632]
Dual values: [array([0.12603657]), array([0.0702077]), array([0.44619011])]
```



```
In [102]: def f_T(w,Cload1,Cload2,Cload3):
    T1 = (w[2] + Cload1)*(w[0]**-1 + w[1]**-1 + w[2]**-1) + w[1]*(w[0]**-1 + w[1]**-1) + (w[0]+w[3]+w[4]
    T2 = (w[4]+Cload2)*(w[0]+w[3]+w[4]) + w[3]*(w[0]**-1+w[3]**-1) + (w[5] + Cload3)*(w[0]**-1 + w[3]**-
    T3 = (w[5]+Cload3)*(w[0]**-1+w[3]**-1+w[5]**-1) + w[3]*(w[0]**-1+w[3]**-1) + (w[0]+w[1]+w[2]+Cload1
    T = np.max([T1,T2,T3])
    return T
```

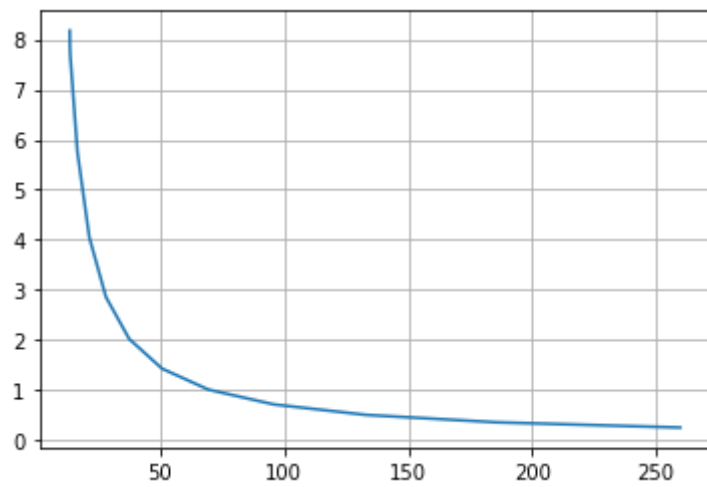
```
In [108]: # Solve for multiple values of mu:
opt_val = np.zeros(N)
opt_w = np.zeros((6,N))
opt_t = np.zeros(N)
opt_T = np.zeros(N)

for i in np.arange(N):
    w = cp.Variable(shape=(6,), pos=True, name="w")
    t = cp.Variable(1,pos=True)
    T1 = (w[2] + Cload1)*(w[0]**-1 + w[1]**-1 + w[2]**-1) + w[1]*(w[0]**-1 + w[1]**-1) + (w[0]+w[3]+w[4]
    T2 = (w[4]+Cload2)*(w[0]+w[3]+w[4]) + w[3]*(w[0]**-1+w[3]**-1) + (w[5] + Cload3)*(w[0]**-1 + w[3]**-
    T3 = (w[5]+Cload3)*(w[0]**-1+w[3]**-1+w[5]**-1) + w[3]*(w[0]**-1+w[3]**-1) + (w[0]+w[1]+w[2]+Cload1
    T1 /= t
    T2 /= t
    T3 /= t

    mu = mu_v[i]
    objective_fn = cp.sum(w) + mu*t
    constraints = [T1 <= 1, T2 <= 1, T3 <= 1]
    problem = cp.Problem(cp.Minimize(objective_fn), constraints)
    problem.solve(gp=True)
    opt_val[i] = problem.value
    opt_w[:,i] = w.value
    opt_t = t.value
    opt_T[i] = f_T(w.value,Cload1,Cload2,Cload3)
```

```
In [125]: A_opt = np.sum(opt_w,axis=0)
```

```
plt.plot(opt_T,A_opt)  
plt.grid()  
# plt.xscale('log')  
plt.show()
```



In [126]: *# Optimal solution and non-optimal solution from part a overlaid:*

```
plt.plot(opt_T,A_opt)
plt.plot(T_v,A_v,'--')
plt.grid()
# plt.xscale('log')
plt.show()
```

