

Elements of AI - Assignment 0

Soln 1. The abstraction used in the program with $N \times N$ chessboard and using N pieces has been summarized as follows:-

i) Initial state :-

In the beginning the initial board is a matrix of 0's (representation of empty squares). Thus, the initial state could be said to be an empty chessboard with none of the pieces placed on it. Eg. for a board, $N=2$, the initial state is a list of lists: $[[0,0], [0,0]]$ shown as:

— —
— —

ii) Goal state :-

The goal for the problem is the state with all the N pieces placed at such a position that none of the pieces attack one another. For the n rooks problem, it is such that none of the n rooks are in the same row and column. There can be multiple solutions (goal states). One eg :-
for $N=2$.

R —
— R

iii) Successor Function: The

The successor function generates the list of next valid states and appends them to the fringe. In the original program, it just adds a piece to all the squares on the board that is passed to it. If the new board isn't a goal state then it is appended to the fringe. A point to note is that if the square already has a piece we still go through the process of putting a piece there and return technically the same board. The Depth First search paradigm has been used to implement the ~~fringe~~ search with the help of stacks data structure.

Eg. let's assume a board of $_ _ R _ _$ and the successor function will give :-

$\begin{array}{ccccccc} R & R & ; & _ & R & ; & _ & R & ; & _ & R \\ _ & _ & & _ & _ & & R & _ & & _ & R \end{array}$
 (this will give)
 $\begin{array}{ccccccc} R & R & ; & R & R & ; & R & R & ; & R & R \\ _ & _ & & _ & _ & & R & _ & & _ & R \end{array}$

3 rooks can also exist on a 2×2 board with this function.

iv) Cost function :-

The cost function represents the cost of going from one state to another state. Since we are just placing a rook at a square (if it already has one, we still go through the entire process), the cost function for going from one state to the next is a constant factor of 1. (The cost of placing a rook anywhere on the board is the same)

v) Valid states :-

A valid state can be said to be state that is possible / attainable by the successor function. A valid state for this problem could be a state with rooks placed randomly anywhere on it. The total number of rooks can exceed N also. We also don't take into account the attacking constraint while generating valid states.

eg of valid states : $_ _ _ ; \begin{matrix} R \\ R \end{matrix} _ _ ;$

$RR ; R _ ; R R$, etc.
 $R _ ; _ R ; _ _$

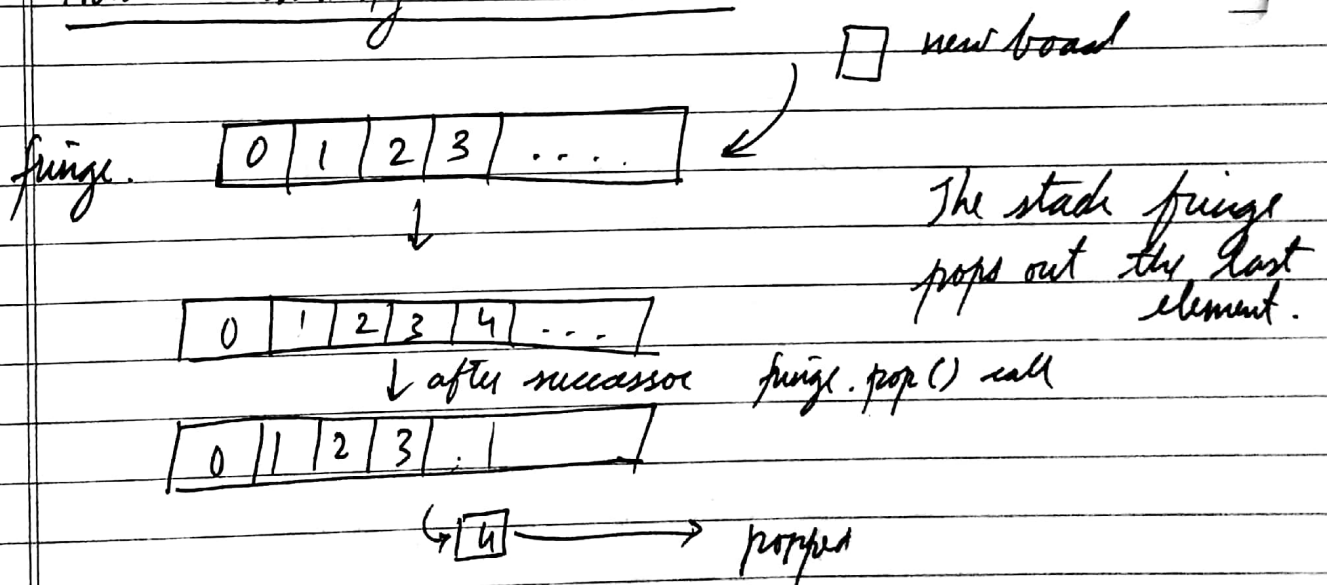
Soln 2. The code is converted to BFS by using (fringe.pop(0)) which pops the elements from the beginning of the fringe i.e. first in first out (FIFO) paradigm.

A tabular description of N versus run-time.

Type \ N	1	2	3	4
DFS	runs	runs	no output	no output
BFS	runs	runs	runs	runs.

For $N > 2$ DFS doesn't give output whereas BFS gives an output quickly. DFS goes on an infinite loop checking the same board again & again.

How to modify to BFS?



So convert to a BFS we must remove the board that came first into the list. We can do this by using:

`fringe.pop(0)`

↑ specify the location to pop from.

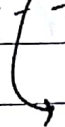
Now all the boards will be taken from the 0th position and we get a BFS implementation instead of a DFS implementation.

0	1	2	3	
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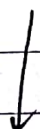
`fringe.pop(0)`

1	2	3		
---	---	---	--	--



0

→ popped



add board '4'

1	2	3	4	
---	---	---	---	--



added at the end as before.

Soln 3. Successors 2 is implemented as follows (given in code file). The following logic is added:

- Only append to ^{successors list} fringe, if new-board generated from add-price is not same as original board.
- Only return ^{append} the board to successors if the total number of pieces on board are less than equal to N .

⇒ The add-price function returns a board with a piece at (r, c) passed to it.

- We must now check if this new board is equal to original board or not and
- check if there are more than N pieces on the new board.

If the new-board passes these tests we append it to a list of successors that is returned to solve at the end.

⇒ Now both BFS & DFS run but still they both are very slow. DFS runs because now fringe doesn't have duplicate boards in it and doesn't go on infinite loop.

Soln 4. Successors 3 has been implemented in code file.

The logic used has been:-

- We specify a list empty - r $[]$ with $[1, 2, 3, \dots, N]$
- Now iterate over all the squares in the board passed to successors. If the square has a piece, remove the row from empty - r $[]$ ie. if board $[3][4]$ has a piece $(=1)$ then remove "3" from empty - r .
- Now specify ~~row~~ a variable "col" equal to the number of pieces on the board. ie. it'll be "4" if board already has 4 rooks placed & so on.
- Now simply iterate over all the elements in empty - r $[]$ place a piece at the values of empty - r $[]$ & col. (fixed).
- Use the limiting condition defined in successors2 as is and append to successors function if it is valid.
- Return successors list.
- Fringe uses DFS (DFS is very slow)

Tabular representation of my successors 3 with time:-

N	Time taken in seconds
5	0.00014
20	0.00454
30	0.0162
40	0.0437
50	0.095
100	1.214
150	5.214
200	15.683
250	37.542
300	75.367

The program runs for $N=280$ in 58 seconds.
So, I guess $N=280$ will be the limit for running under 60 seconds.