

Error Control Coding - Low-Density Parity-Check Codes

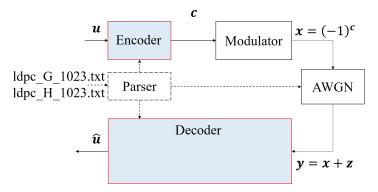
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- Consider the (1023,781) low-density parity-check (LDPC) code with block length n=1023, dimension k=781, row weight $\rho=32$, and column weight $\gamma=32$.
- Block Diagram:





Encoding of Low-Density Parity-Check Codes

Use the recursion

$$u_{l+6} = u_{l+1} \oplus u_l, \quad \text{for } l \ge 0$$

with the initial conditions

$$u_0 = 1, u_1 = u_2 = u_3 = u_4 = u_5 = 0$$

to generate k = 781 information bits.

• The generated sequence is 100000100001... with period 63.



Additive white Gaussian noise (AWGN) channel:

• For a binary code of rate R with BPSK modulation, the noise variance σ^2 is given by

$$\sigma^2 = \left(2R\frac{E_b}{N_0}\right)^{-1}$$

where E_b/N_0 is the bit signal-to-noise ratio (SNR).

• The code rate is R = 781/1023 here.



• Please use the following pseudo code:

```
#define IA 16807
#define IM 2147483647
#define AM (1.0/IM)
#define IQ 127773
#define IR 2836
#define NTAB 32
#define NDIV (1+(IM-1)/NTAB)
#define EPS 1.2e-7
#define RNMX (1.0-EPS)
main()
   long *idum;
   idum = (long *)malloc(sizeof(long));
   *idum = SEED; //SEED must be a negative integer.
```



Please use **normal()** to output two independent normal random

```
variables. n_1 and n_2.
normal(n_1, n_2, \sigma)
      do{
          x_1 = ran1(idum);
           x_2 = ran1(idum);
          x_1 = 2x_1 - 1:
          x_2 = 2x_2 - 1:
           s = x_1^2 + x_2^2;
      } while (s \ge 1.0)
      n_1 = \sigma x_1 \sqrt{-2 \ln s/s};
      n_2 = \sigma x_2 \sqrt{-2 \ln s/s};
```

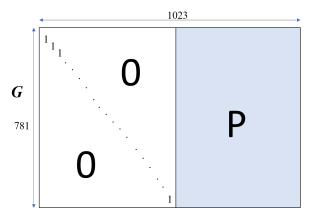


 Please use ran1() to generate a random variable uniformly distributed in the interval (0,1).

```
ran1(long *idum)
   int j;
   long k;
   static long iy=0;
   static long iv[NTAB];
   double temp;
   if (*idum <= 0 || !iv){
       if (-(*idum) < 1) *idum=1;
       else *idum = -(*idum);
       for (j=NTAB+7; j>=0; j--){
           k=(*idum)/IQ:
           *idum=IA*(*idum-k*IQ)-IR*k:
           if (*idum < 0) *idum += IM:
           if (j < NTAB) iv[j] = *idum;
       iy=iv[0];
   k=(*idum)/IQ:
   *idum=IA*(*idum-k*IO)-IR*k:
   if (*idum < 0) *idum += IM;
   j=iv/NDIV;
   iv=iv[j];
   iv[i] = *idum:
   if ((temp=AM*iv) > RNMX) return RNMX;
   else return temp;
```

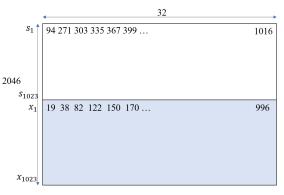


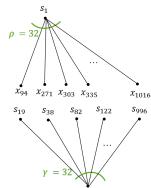
Input Idpc_G_1023.txt:





Input ldpc_H_1023.txt:







Decoding of Low-Density Parity-Check Codes

程式DEMO方式

- 沒有DEMO (不需跟助教比對)。
- 5/31 開始進行期末Project。

上台報告

- Bonus
- 已經完成的同學做簡報分享,需要先跟我報名登記(限量)。
- 6/7 上課時間。

報告繳交

- 6/14 12:00前私訊檔案給"助教_高子傑"。
- 檔案名稱格式: 學號_姓名.zip 例如: E94081042_許博士.zip
- 壓縮檔內包含所有的.cpp (必須要有註釋) 以及 一個報告.pdf。
- 報告檔名格式: Project2_學號_姓名.pdf 例如: Project2_E94081042_許博士.pdf

Decoding of Low-Density Parity-Check Codes

報告繳交内容

- 報告必要内容 (Baseline):
 - 解碼採用Sum-Product Algorithm;
 - 系統架構圖 (細部blcok diagram);
 - 程式流程解釋 (流程或是Pseudo code);
 - 參數設定: number of maximum iterations = 100 等其他參數;
 - 模擬數據 (表格): SNR (dB) \ No. decoded bits \ No. decoded bit errors \ bit error rate (BER);
 - 模擬圖: 至少到 BER = 10^{-4} .
- 報告加分内容 (Bonus):
 - 不同 iteration個數的比較,收斂的狀況與分析:
 - 其他演算法: bit-flipping algorithm、min-sum algorithm、或其他文獻 上的做法。
 - 錯誤率能夠越低越好,但是要數值穩定的曲線。
 - BER之外,提供Block error rate (BLER)的模擬效能圖。
 - 其他觀察或比較。



Decoding of Low-Density Parity-Check Codes

參考數值:

- Sum-Product Algorithm (Maximum 100 iterations):
 - BER = 4.0×10^{-2} at $E_b/N_0 = 2.2$ dB;
 - BER = 1.0×10^{-3} at $E_b/N_0 = 3.0$ dB.