

Abstract

More basic category theory. The main sources for this talk are the following.

- *nLab*.
- John Rognes's *Lecture Notes on Algebraic K-Theory*, Ch. 3.
- Peter Johnstone's lecture notes for "Category Theory" (Mathematical Tripos Part III, Michaelmas 2015), Ch. 1.

Definition 1. Let \mathcal{C} and \mathcal{D} be categories and $F, G : \mathcal{C} \rightarrow \mathcal{D}$ be functors. A *natural transformation* $\phi : F \Rightarrow G$ is a function $A \mapsto f_A$ from $\text{ob } \mathcal{C}$ to $\text{mor } \mathcal{D}$ such that $f_A : F(A) \rightarrow G(A)$ and the following diagram commutes for any morphism $f : A \rightarrow B$.

$$\begin{array}{ccc} FA & \xrightarrow{Ff} & FB \\ f_A \downarrow & & \downarrow f_B \\ GA & \xrightarrow{Gf} & GB \end{array}$$

In symbols, this may be written as $f_B f_* = f_* f_A$, where f_A and f_B are called the *components* of ϕ .

Remark 1. If every f_A is an isomorphism, then the $(f_A)^{-1}$ define a natural transformation between the same two functors.

Definition 2. Let $F, G, H : \mathcal{C} \rightarrow \mathcal{D}$ be functors. The *identity natural transformation* $\text{Id}_F : F \Rightarrow F$ is given by $A \mapsto \text{Id}_{F(A)}$. Moreover, given natural transformations $\phi : F \rightarrow G$ and $\psi : G \rightarrow H$, define the *composite natural transformation* $\psi \circ \phi$ by $A \mapsto (\psi \circ \phi)_A := \psi_A \circ \phi_A$.

Definition 3. If each f_A is an isomorphism, then we call $\phi : F \cong G$ a *natural isomorphism*.

Remark 2. If \mathcal{D} is a groupoid, then ϕ must be a natural isomorphism.

Lemma 1. A natural transformation $\phi : F \Rightarrow G$ is a natural isomorphism iff it has an inverse $\phi^{-1} : G \Rightarrow F$.

Proof. This follows from Remark 1 and the definition of composite natural transformation. \square

Example 4. Let R and S be commutative rings. Any ring homomorphism $f : R \rightarrow S$ induces a ring homomorphism $\text{GL}_n(f) : \text{GL}_n(R) \rightarrow \text{GL}_n(S)$ which satisfies $f(\det(A)) = \det(\text{GL}_n(f)(A))$. Viewing GL_n and $R \mapsto R^*$ as functors from **Rng** to **Grp** and $\det_R : \text{GL}_n(R) \rightarrow R^*$ as a morphism in **Grp**, we see that \det_R defines a natural transformation $\phi : \text{GL}_n \Rightarrow f^*$, where f^* denotes $f \upharpoonright_{R^*} R^* \rightarrow S^*$.

$$\begin{array}{ccc} \text{GL}_n(R) & \xrightarrow{\text{GL}_n(f)} & \text{GL}_n(S) \\ \det_R \downarrow & & \downarrow \det_S \\ R^* & \xrightarrow{f^*} & S^* \end{array}$$

Example 5. Recall the power set functor $P : \mathbf{Set} \rightarrow \mathbf{Set}$ given by $A \mapsto P(A)$ and $Pg(S) = g(S)$ where $g : A \rightarrow B$ is a function and $S \subset A$. Then the function $f_A : A \rightarrow P(A)$ given by $a \mapsto \{a\}$ defines a natural transformation $\phi : \text{Id}_{\mathbf{Set}} \Rightarrow P$.

Example 6. Set $\mathcal{C} = \mathcal{D} = \mathbf{Grp}$, $F = \text{Id}_{\mathcal{C}}$, and G equal to the abelianization functor. Then given a group H , the homomorphism $f : H \rightarrow H^{\text{ab}}$ defines a natural transformation $\phi : F \Rightarrow G$.

Example 7. Consider the preorders (P, \leq) and (Q, \leq) as small categories where functors $F, G : P \rightarrow Q$ are order-preserving functions. Then there is a unique natural transformation $\phi : F \Rightarrow G$ iff $F(x) \leq G(x)$ for every $x \in P$.

Example 8. The inversion isomorphism from a group G to G^{op} defines a natural transformation $\phi : \text{Id}_{\mathbf{Grp}} \Rightarrow (^{\text{op}} : \mathbf{Grp} \rightarrow \mathbf{Grp})$. In other words, G is naturally isomorphic to G^{op} .

Definition 9. Let \mathcal{C} and \mathcal{D} be categories with \mathcal{C} small. The *functor category* $\mathbf{Fun}(\mathcal{C}, \mathcal{D}) := \mathcal{D}^{\mathcal{C}}$ has functors $F : \mathcal{C} \rightarrow \mathcal{D}$ as objects and natural transformations as morphisms.

Remark 3. Given functors $F, G : \mathcal{C} \rightarrow \mathcal{D}$, why is the class of natural transformation $\phi : F \Rightarrow G$ necessarily a set? A G -Universe models ZFC, in particular **Replacement**.

Definition 10. Given a category \mathcal{C} , the *arrow category* $\text{Ar}(\mathcal{C})$ of \mathcal{C} has as objects morphisms $f : X_0 \rightarrow X_1$ in \mathcal{C} and as morphisms $M : (f : X_0 \rightarrow X_1) \rightarrow (g : Y_0 \rightarrow Y_1)$ the pairs (M_0, M_1) of morphisms $M_0 : X_0 \rightarrow Y_0$ and $M_1 : X_1 \rightarrow Y_1$ such that

$$\begin{array}{ccc} X_0 & \xrightarrow{f} & X_1 \\ M_0 \downarrow & & \downarrow M_1 \\ Y_0 & \xrightarrow{g} & Y_1 \end{array}$$

commutes.

Note that $\text{Ar}(\mathcal{C}) \cong \mathbf{Fun}([1], \mathcal{C})$.

Lemma 2. $\mathbf{Fun}(\mathcal{C} \times \mathcal{D}, \mathcal{E}) \cong \mathbf{Fun}(\mathcal{C}, \mathbf{Fun}(\mathcal{D}, \mathcal{E}))$ via currying.

Definition 11. A functor $F : \mathcal{C} \rightarrow \mathcal{D}$ is an *equivalence* if there is a functor $G : \mathcal{D} \rightarrow \mathcal{C}$ such that $F \circ G \cong \text{Id}_{\mathcal{D}}$ and $G \circ F \cong \text{Id}_{\mathcal{C}}$. In this case, we say that F and G are *equivalent categories*. Moreover, we say that a property of \mathcal{C} is *categorical* if it is invariant under such equivalence.

Example 12. Let k be a field. Let the category \mathbf{Mat}_k have natural numbers as objects and morphisms $n \rightarrow p$ given by $p \times n$ matrices over k . Let \mathbf{fdMod} denote the category of finite-dimensional vector spaces with linear maps as morphisms. These two categories are equivalent. Send $\text{nat } n$ to k^n in one direction and the space V to $\dim V$ in the other direction.

Definition 13. A functor $F : \mathcal{C} \rightarrow \mathcal{D}$ is *essentially surjective* if for each object Z of \mathcal{D} , there is some object Y of \mathcal{C} such that $F(Y) \cong Z$.

Theorem 3. A functor is an equivalence iff it is full, faithful, and essentially surjective.

Proof. See Rognes, Theorem 3.2.10. □

Definition 14. A *skeleton* of \mathcal{C} is a full subcategory $\mathcal{C}' \subset \mathcal{C}$ such that each element of $\text{ob } \mathcal{C}$ is isomorphic to exactly one element of $\text{ob } \mathcal{C}'$.

Lemma 4. With notation as before, \mathcal{C}' and \mathcal{C} are equivalent categories via the inclusion functor.

Proof. Apply Theorem 3. □

Lemma 5. Any two skeleta $\mathcal{C}', \mathcal{C}'' \subset \mathcal{C}$ are isomorphic.

Proof. Define $F : \mathcal{C}' \rightarrow \mathcal{C}''$ by $F(X) = Y$ where $h_X : X \cong Y$ and $F(f) = h_Y \circ f \circ (h_X)^{-1}$ for $f \in \mathcal{C}(X, Y)$. To get F^{-1} , similarly define $G : \mathcal{C}'' \rightarrow \mathcal{C}'$ by choosing $(h_X)^{-1}$. □

Remark 4. The previous two lemmas are equivalent to the axiom of choice, as is the statement that every category admits a skeleton.

Definition 15. Fix $X \in \text{ob } \mathcal{C}$. Define the functor $\mathcal{Y}^X : \mathcal{C} \rightarrow \mathbf{Set}$ by $Y \mapsto \mathcal{C}(X, Y)$ and mapping each morphism $g : Y \rightarrow Z$ to $g_* : \mathcal{C}(X, Y) \rightarrow \mathcal{C}(X, Z)$ given by $f \mapsto gf$. We call $\mathcal{C}(X, -) := \mathcal{Y}^X$ the set-valued functor *corepresented* by X in \mathcal{C} .

Definition 16. Fix $Z \in \text{ob } \mathcal{C}$. Define the contravariant functor $\mathcal{Y}_Z : \mathcal{C}^{\text{op}} \rightarrow \mathbf{Set}$ by $Y \mapsto \mathcal{C}(Y, Z)$ and mapping each morphism $f : X \rightarrow Y$ in \mathcal{C} to $f^* : \mathcal{C}(Y, Z) \rightarrow \mathcal{C}(X, Z)$ given by $g \mapsto gf$. We call $\mathcal{C}(-, Z) := \mathcal{Y}^Z$ the set-valued functor *represented* by Z in \mathcal{C} .

Definition 17. A functor $F : \mathcal{C} \times \mathcal{C}' \rightarrow \mathcal{D}$ is also called a *bifunctor*.

Example 18. Let \mathcal{C} be a category. Define $\mathcal{C}(-, -) : \mathcal{C}^{\text{op}} \times \mathcal{C} \rightarrow \mathbf{Set}$ by $(X, X') \mapsto \mathcal{C}(X, X')$ and mapping each morphism $(f, f') : (X, X') \rightarrow (Y, Y')$ to $\mathcal{C}(f, f') : \mathcal{C}(X, X') \rightarrow \mathcal{C}(Y, Y')$ given by $g \mapsto f'gf$.

Definition 19. This is due to Dan Kan. Let \mathcal{C} and \mathcal{D} be categories and $F : \mathcal{C} \rightarrow \mathcal{D}$ and $G : \mathcal{D} \rightarrow \mathcal{C}$ be functors. Consider the set-valued bifunctors $\mathcal{D}(F(-), -), \mathcal{C}(-, G(-)) : \mathcal{C}^{\text{op}} \times \mathcal{D} \rightarrow \mathbf{Set}$. An *adjunction* between F and G is a natural isomorphism $\phi : \mathcal{D}(F(-), -) \Rightarrow \mathcal{C}(-, G(-))$. If such ϕ exists, then we say that (F, G) is an *adjoint pair* of functors. We also call F the *left adjoint* to G and G the *right adjoint* to F .

Remark 5. For each $c : X' \rightarrow X$ and $d : Y \rightarrow Y'$, the following commutes.

$$\begin{array}{ccc} \mathcal{D}(F(X), Y) & \xrightarrow{\phi_{X,Y}} & \mathcal{C}(X, G(Y)) \\ c^* d_* \downarrow & & \downarrow c^* d_* \\ \mathcal{D}(F(X'), Y') & \xrightarrow{\phi_{X',Y'}} & \mathcal{C}(X', G(Y')) \end{array}$$

Example 20. The forgetful functor $U : \mathbf{Grp} \rightarrow \mathbf{Set}$ admits a left adjoint $F : \mathbf{Set} \rightarrow \mathbf{Grp}$ which maps a set to the free group generated by A . The adjunction is the natural bijection $\mathbf{Set}(A, U(G)) \cong \mathbf{Grp}(F(A), G)$.

Example 21. Let R be a ring. The forgetful functor $U : R\text{-}\mathbf{Mod} \rightarrow \mathbf{Set}$ admits a left adjoint $R(-)$ sending a set S to $\bigoplus_{s \in S} R$, the free R -module generated by S . The adjunction is the natural bijection $\mathbf{Set}(S, U(M)) \cong R\text{-}\mathbf{Mod}(R(S), M)$.

Remark 6. U does not admit a right adjoint in either of the previous two examples.

Example 22. The forgetful functor $U : \mathbf{Top} \rightarrow \mathbf{Set}$ has left adjoint that sends a set to the same set equipped with the discrete topology. It also has a right adjoint via the functor sending a set to the same set equipped with the indiscrete topology.

Example 23. Let \mathbf{CMon} be the category of commutative monoids. Given $M \in \text{ob } \mathbf{CMon}$, we can construct the completion, or Grothendieck group, $G(M)$ on $M \times M$ as follows. Define addition on $M \times M$ component-wise and say that $(m_1, m_2) \sim (n_1, n_2)$ if $m_1 + m_2 + k = m_2 + n_1 + k$ for some $k \in M$. Set $G(M)$ as $(M \times M / \sim, +)$.

Then $G : \mathbf{CMon} \rightarrow \mathbf{Ab}$ is a functor. This is left adjoint to the forgetful functor $U : \mathbf{Ab} \rightarrow \mathbf{CMon}$.

Remark 7. Read Rognes, Definition 3.4.8, where he constructs the group completion $K(M)$ of non-commutative monoids M . It turns out that $K(M)$ is realized as the fundamental group of an important classifying space.

Definition 24. A subcategory $\mathcal{C} \subset \mathcal{D}$ is *reflective* if the inclusion functor is a right adjoint and is *coreflective* if the inclusion functor is a left adjoint.

Example 25.

1. $\mathbf{Ab} \subset \mathbf{CMon}$ is reflective by Example 8.
2. $\mathbf{Ab} \subset \mathbf{Grp}$ is reflective.
3. Let $\mathbf{Ab}_T \subset \mathbf{Ab}$ denote the category of torsion groups. This is coreflective via the functor sending an abelian group to its torsion subgroup because any homomorphism $f : A \rightarrow B$ where A is torsion has $f(A) \subset B_T$.

Definition 26. Given an adjunction $\phi : \mathcal{D}(F(-), -) \Rightarrow \mathcal{C}(-, G(-))$, define the *unit morphism*

$$\eta_X = \phi_{X, F(X)}(\text{Id}_{F(X)})$$

and the *counit morphism*

$$\epsilon_Y = \phi_{G(Y), Y}^{-1}(\text{Id}_{G(Y)}).$$

Lemma 6. *The unit morphisms η_X define a natural transformation $\eta : \text{Id}_{\mathcal{C}} \Rightarrow GF$, and the counit morphisms ϵ_Y define a natural transformation $\epsilon : FG \Rightarrow \text{Id}_{\mathcal{D}}$.*