# A Network-of-Networks Model for Electrical Infrastructure Networks

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#### Abstract

Modeling power transmission networks is an important area of research with applications such as vulnerability analysis, study of cascading failures, and location of measurement devices. Graph-theoretic approaches have been widely used to solve these problems, but are subject to several limitations. One of the limitations is the ability to model a heterogeneous system in a consistent manner using the standard graph-theoretic formulation. In this paper, we propose a *network-of-networks* approach for modeling power transmission networks in order to explicitly incorporate heterogeneity in the model. This model distinguishes between different components of the network that operate at different voltage ratings, and also captures the intra and inter-network connectivity patterns. By building the graph in this fashion we present a novel, and fundamentally different, perspective of power transmission networks. Consequently, this novel approach will have a significant impact on the graph-theoretic modeling of power grids that we believe will lead to a better understanding of transmission networks.

# 1. Introduction

The topological structure of electric power grids can be modeled as a graph. A graph G=(V,E) is a pair, where the vertex set V represents unique entities and the edge set E represents binary relationships on V. The study of topological and electrical structure of power transmission networks is an important area of research with several applications such as vulnerability analysis [1], locational marginal pricing [2], controlled islanding [3], and location of sensors [4]. In this paper, we focus on two interrelated problems: (i) graph-based modeling and characterization of power transmission networks, and (ii) random graph models for synthetic generation of transmission networks. These two problems have a significant impact on several aspects of our understanding and functioning of the power grid.

Graph-theoretic approaches have been studied extensively in the context of power grids (Section 6). However, there are several limitations that need to be addressed. For example, power grids have been characterized using graph-theoretic metrics to show conflicting features [5], and in the study of vulnerabilities to often misleading conclusions [1]. Graph based models and algorithms are effective when vertices represent homogeneous entities in a consistent manner. However, a transmission network is inherently heterogeneous. It consists of entities such as generators, loads, substations, transformers, and transmission lines that operate at different voltage ratings. Further, power transformers are not always treated consistently. They are generally represented as edges, but not always. Consequently, a purely graph-based model will misrepresent a transmission network when it is not constructed carefully. As an example, illustration of a model for the power grid in Poland is provided in Figure 1, where nodes operating at different voltages are shown in different colors and sizes - larger the size, higher the nominal voltage of a node. Thus, edges connecting two different types of vertices are transformers. In this paper, we address the problem of accurate representation of power transmission networks as a graph, and discuss the impact of this representation for graph-theoretic characterization and modeling using random graph models.

# A. Contributions

By accurately modeling transformers, and consequently decomposing a network into regions of different voltage

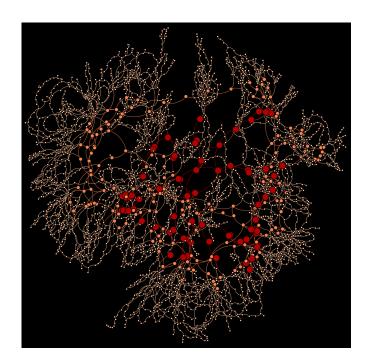


Fig. 1. Rendering of the Polish system with 3120 vertices and 3684 edges. The colors and sizes represent different voltage ratings. Larger the size, higher the voltage. Details are provided in Section 3

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ratings, we propose a novel network-of-networks model for power transmission networks. This simple idea has profound implications to the study of topological and electrical structure of power grids.

The contributions we make in this paper are: (i) a new decomposition method for power transmission networks using nominal voltage ratings of entities (Section 3); (ii) empirical evidence of the hierarchical nature of transmission networks based on the analysis of real-world data representing the north American and European grids (Section 3 and 4); (iii) characterization of the interconnection structure between networks of different voltages (Section 4); and (iv) presentation of random graph models – two methods for modeling the network at a specific voltage level, and one method for modeling the interconnection structure of networks of different voltage (Section 5).

#### 2. Network-of-Networks Model

An electric power transmission network consists of power generators, loads, transformers, substations, and branches connecting different components. While branches and transformers are modeled as edges of a graph, all the other components (buses) are modeled as vertices. While this model is fairly consistent, the inclusion of transformers as edges makes the graph representation inconsistent. An illustration of a symbolic power system is shown in Figure 2. Generation is shown in black, transmission is shown in blue, and distribution in green. A power transformer is used to step up or step down voltage

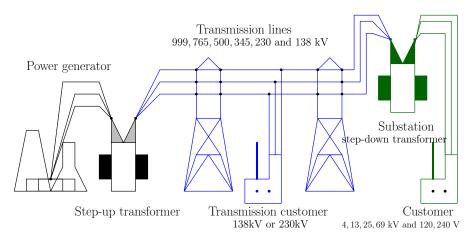


Fig. 2. An illustration of a power system. Power generation, shown in black, is connected to distribution, shown in green, via a transmission system, shown in blue. The system operates at different voltage ratings as shown in the figure.

between its two end-points, and therefore, connects two different regions of voltage magnitudes. Power transmitted at a certain voltage remains at that level unless stepped up or stepped down by another transformer connecting the network to a region of different voltage rating. Therefore, the component of the network transmitting power at a certain voltage rating can be treated as one network where all the vertices consistently represent entities operating at the same voltage rating, and edges represent transmission lines. The overall transmission network can thus be decomposed into a *network-of-networks* (defined further below), where each network operates at a certain voltage rating, and a pair of networks are connected via one or more transformers regulating the voltage levels. When expressed in this manner, a different graph structure emerges, which is the focus of our study.

### A. Preliminaries

A graph G=(V,E) is a pair, where the vertex set V represents unique entities and the edge set E represents binary relations on V. We define a network-of-networks model for power transmission as an undirected heterogenous graph  $G=(V,E,X,Y,\phi,\varphi)$  with a weight function  $w:E\to\mathbb{R}^+$ . The vertex set V represents unique entities such as substations, generators and loads; the edge set E represents binary relationships on V such that an edge  $e\in E$  represents either a power transmission line or a voltage transformer; a set X of values associated with V representing the nominal voltage rating at a particular vertex as defined in the mapping function  $\phi:V\to X$  where each vertex operates at a given voltage level  $\phi(v)\in X$ , and  $X\in\mathbb{R}^+$  is a positive real number. A set Y of values associated with E represent the type of an edge – a transmission line or a transformer – as defined in the mapping function  $\varphi:E\to Y$  where each edge is of a given type  $\varphi(e)\in Y$ , and  $Y\in\mathbb{Z}$  is an integer.

Traditionally, transmission networks are represented as graphs, making no distinction between entities operating at different voltage ratings and edges of different types. In our representation, the vertex and edge types (X,Y), and the mappings  $(\phi,\varphi)$ , allow us to make clear distinctions between different vertex and edge types. Given the voltage levels for each substation  $(X,\phi)$ , we can automatically detect the voltage transformers,  $E_T \subseteq E$ , where  $e_t(u,v) \in E_T$  such that  $\phi(u) \neq \phi(v)$ . By pruning the

transformer-edges from G, the remaining graph represents a collection of subgraphs where each subgraph has vertices operating at the same voltage level and edges that are transmission lines. The structure of this collection of subgraphs is of great interest in studying the transmission network and is the focus of this paper.

**Graph-theoretic Measures:** We use the following graph-theoretic measures in this paper. A path in a graph is a finite sequence of edges such that every edge in a path is traversed only once and any two consecutive edges in a path are adjacent to each other (share a vertex in common). The average shortest-path length is defined as the average of shortest paths between all possible pairs of vertices in a graph. The longest shortest-path is called the diameter of a graph. The local clustering coefficient of a vertex v is the ratio between the actual number of edges among the neighbors of v to the total possible number of edges among these neighbors. The average clustering coefficient of a graph is the sum of all local clustering coefficients of its vertices divided by the number of vertices.

## 3. CHARACTERIZATION OF REAL-WORLD NETWORKS

We now present details of the network-of-networks model by decomposing different datasets using the method described in Algorithm 1.

**Algorithm 1** Decompose a graph based on voltage ratings. **Input:** A graph  $G = (V, E, X, \phi)$ . **Output:** A graph  $G' = \{G_i, G_j, ...\}$ , which is a collection of disconnected components of G.

```
1: procedure DECOMPOSE-GRAPH(G = (V, E, X, \phi), G')
 2:
          Q \leftarrow V
          G' \leftarrow \emptyset

    A collection of graphs

 3:
          while Q \neq \emptyset do
 4:
               v \leftarrow Q \setminus \{v\} in some order
 5:
               Q_t \leftarrow \{v\}
 6:
               G_v \leftarrow \{v\}
                                                                                                                 > Represents a graph of same voltage level
 7:
               while Q_t \neq \emptyset do
 8:
                    x \leftarrow Q_t \setminus \{x\} in some order
 9:
                    Q \leftarrow Q \setminus \{x\}
10:
                    for each y \in adj(x) do
                                                                                                                                                     \triangleright Neighbors of x
11:
                         if X[x] = X[y] then

    Same voltage rating

12.
                              Q_t \leftarrow Q_t \cup \{y\}
13:
                              V(G_v) \leftarrow V(G_v) \cup \{y\}

    A new vertex

14:
                              E(G_v) \leftarrow E(G_v) \cup \{(x,y)\}

    A new edge

15:
               G' \leftarrow G_v
16:
```

Let  $G=(V,E,X,\phi)$  be the input graph, and G' be the decomposed graph that is computed as the output of Algorithm 1. The algorithm starts by adding the vertices in G to a queue Q (Line 2 in Algorithm 1) in an arbitrary order. The **while** loop on Line 4 iterates until all the vertices in the queue have been processed. Let us consider the process for an arbitrary vertex v chosen from Q (Line 5). Vertex v, which is added to a new queue  $Q_t$ , acts as the source of a new network (a connected component of G) consisting all the vertices connected to v through edges of the same voltage. The search for the network corresponding to v is managed by adding and removing vertices from  $Q_t$ , and persists until  $Q_t$  becomes empty (the **while** loop on Line 8). For each vertex v in v, we process all the neighbors of v (Line 11) to check if they are of the same voltage rating. The neighbors are added (Lines 14 and 15) to the network corresponding to vertex v if true (same voltage). Once all the connected vertices are added to v, the network is added to v (Line 16). The algorithm resumes from a new vertex that is chosen arbitrarily until all the vertices in v0 have been processed and added to corresponding networks in v1.

TABLE 1

DATASETS USED IN THIS STUDY. ASPL STANDS FOR AVERAGE SHORTEST PATH LENGTH.

| Dataset                   | (Vertices, Edges) | Clust. Coeff. | ASPL  | Diameter |
|---------------------------|-------------------|---------------|-------|----------|
| Polish                    | (2383, 8155)      | 0.012         | 12.58 | 30       |
| Western Interconnect (WI) | (15090, 18153)    | 0.027         | 24.33 | 66       |
| Eastern Interconnect (EI) | (49597, 62985)    | 0.071         | 35.80 | 96       |
| Texas Interconnect (TI)   | (4756, 5848)      | 0.020         | 17.06 | 40       |

In Table I we list the real-world datasets that are used in this study. The table also provides details of the network (without decomposition) in graph-theoretic metrics described in Section 2.1. We study the three major components of the North American power grid – the Eastern, Western and Texas Interconnects, and the Polish system. The Eastern Interconnect data is from a North American Electric Reliability Corporation (NERC) planning model for 2012. Data for Western and Texas Interconnects

are from the Federal Energy Regulatory Commission (FERC) Form 715 filings in 2005. The data for North American power grid is obtained through the U.S. Critical Energy Infrastructure Information (CEII) request process. Data for the Polish system is included with MATPOWER, an open source power system simulation package [6]. For all the datasets we perform several steps to clean the data for consistency. For example, we remove all the isolated vertices and edges – vertices with no edges incident on them, and edges disconnected from the graph.

For each input in the dataset, we provide details of the decomposed networks including information about the number of edges (power transformers) between networks of different voltage levels. Detailed graph measures are provided for the largest component in each level. In order to highlight the topological structure we also provide visual renderings of the networks at each level for all the inputs. With these details, we aim to support of observations that the topological structures are not only similar across different voltage levels, but also across different inputs.

## A. The Polish Network

In the network-of-networks model for the Polish system, there are three voltage levels: 400,200 and 100 kV. The number of components (a network at a specific voltage level) are one, one and five respectively for 400,200 and 100 kV. The size (|V|,|E|) of the largest components are (50,58),(135,174) and (2193,2485) respectively in the same order of voltage levels. For the largest component in each level, the average clustering coefficients are 0.141,0.032 and 0.008; the average path lengths are 6.484,7.899 and 37.572; and the diameters are 17,20 and 92.

Visualization of networks at different voltage levels of the Polish system is provided in Figure 3. We can visually observe from these renderings that the general topological structure consisting of chains and loops is consistent across different voltage levels. From an engineering perspective, loops provide alternative routes when one of the links in a chain fail, and chains help span large geographical areas.

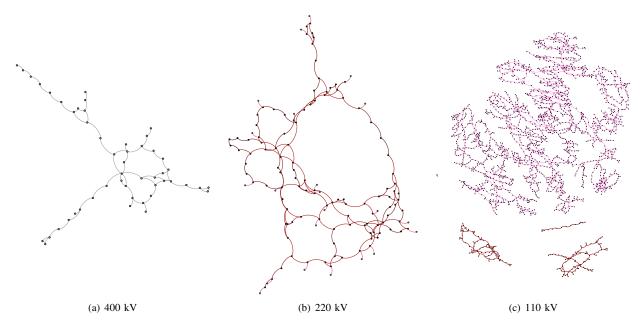


Fig. 3. Rendering of the Polish network at different nominal voltage levels.

In order to study the interconnection of different networks of the same voltage, we study the number of power transformers. The number of transformers (represented as edges) between voltage levels 400 and 200 kV is 15, and that between 400 and 100 kV is 31. There are 121 transformers between voltage levels 200 and 100 kV.

#### B. Western Interconnect Network

The Western Interconnect has four voltage levels: 500, 345, 230 and 138 kV. Key metrics from different levels are summarized in Table II. A large part of the network operates at 138 kV with 5815 vertices and 6371 edges. There are 104 connected components (networks) at the 138 kV, where the largest component has 550 vertices. The second largest portion of the network operates at 230 kV consisting of 20 components, with the largest component having an order of 1766 vertices. The third largest portion operates at 500 kV with 384 vertices organized in four components. The remaining part of the network operates at 345 kV with 305 vertices divided over 12 components. Different graph measures for the four voltage levels are proportional to the sizes of the corresponding graphs irrespective of the voltage ratings. Accordingly, we observe diameters of 64, 37, 72 and 135

TABLE II

DETAILS FROM DECOMPOSING THE WESTERN INTERCONNECT DATASET INTO FOUR VOLTAGE LEVELS. THE VALUES FOR AVERAGE CLUSTERING COEFFICIENT, AVERAGE SHORTEST PATH LENGTH AND DIAMETER ARE FOR THE LARGEST COMPONENT FOR A GIVEN VOLTAGE LEVEL.

| Voltage Levels             | 500 kV     | 345 kV     | 230 kV       | 138 kV       |
|----------------------------|------------|------------|--------------|--------------|
| Size $( V ,  E )$          | (384, 439) | (305, 341) | (1766, 2132) | (5815, 6371) |
| Num components             | 4          | 12         | 20           | 104          |
| V  of largest component    | 360        | 135        | 400          | 550          |
| Avg Clustering Coefficient | 0.02       | 0.047      | 0.055        | 0.022        |
| Avg Shortest Path Length   | 22.397     | 9.571      | 20.693       | 39.822       |
| Diameter                   | 64         | 37         | 72           | 135          |

for the largest components in different voltage levels presented in descending order. While the average clustering coefficients vary, the average shortest path lengths are proportional to the size.

Visualization of networks at voltage levels 500 and 345 kV is provided in Figure 4, and for 230 and 138 kV in Figure 5. Similar to networks in the Polish system, we observe that the pattern of loops and links is consistent across different voltage levels. Further, these patterns are also similar to the patterns observed in the Polish system.

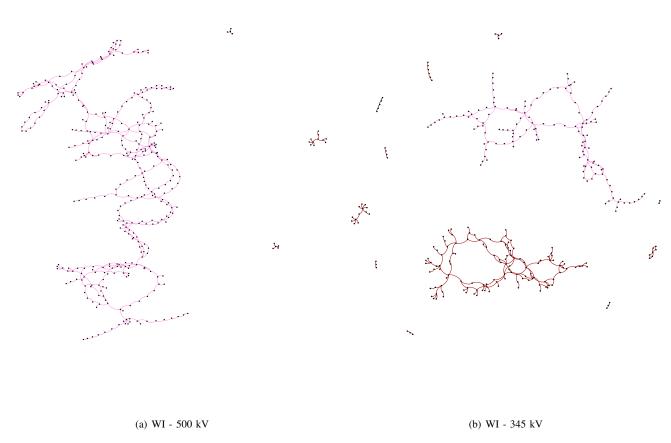


Fig. 4. Rendering of the Western Interconnect networks at 500 and 345 kV.

The number of transformers (represented as edges) between different voltage levels for the Western Interconnect system is summarized in Table III. We observe that the maximum number of connections (549) exist between levels 138 and 230 kV; followed by 123 edges between levels 138 and 345 kV; and 115 edges between levels 230 and 500 kV. There are 10 transformers between 500 and 345 kV. The manner in which these transformers are connected provides a rich interaction network. Details and visualization of this interconnection structure is presented in Section 4.

## C. Eastern Interconnect Network

The Eastern Interconnect is the largest instance in our dataset and has five voltage levels: 765, 500, 345, 230 and 138 kV. Key details of the system are summarized in Table IV. The largest portion of the network operates at 138 kV with 16082

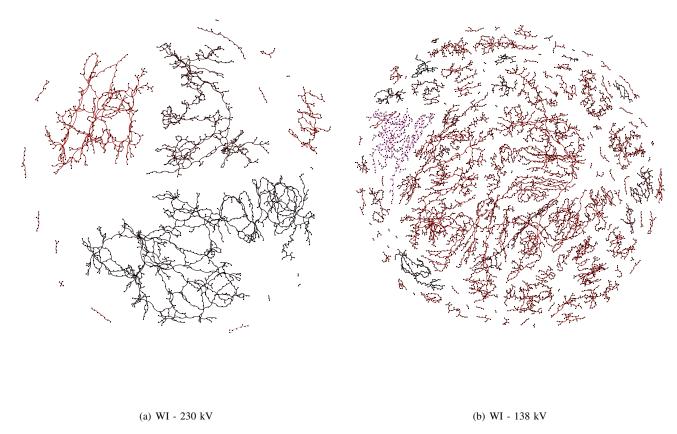


Fig. 5. Rendering of the Western Interconnect networks at 230 and 138 kV.

 $\begin{tabular}{ll} TABLE III \\ TRANSFORMERS IN THE WESTERN INTERCONNECT SYSTEM \\ \end{tabular}$ 

|        | 500 kV | 345 kV | 230 kV | 138 kV |
|--------|--------|--------|--------|--------|
| 500 kV | -      | 10     | 115    | 18     |
| 345 kV | 10     | -      | 54     | 123    |
| 230 kV | 115    | 54     | _      | 549    |
| 138 kV | 18     | 123    | 549    | _      |

vertices and 19110 edges. At this level, there are 121 connected components with the largest component comprising of 12997 vertices. As a result, a large fraction of components at this level are small in size. A similar pattern also emerges at 230 kV with 41 components (Figure 7(b)).

Similar to the Western Interconnect, different graph measures for the five voltage levels are proportional to the sizes of the corresponding graphs irrespective of the voltage ratings. We observe diameters of 386,211,78,48 and 16 for the largest components in different voltage levels in descending order of kV. Again, the average clustering coefficients vary, but the average shortest path lengths are proportional to the size of the largest component. Rendering of the networks at voltage levels 765 and 500 kV is provided in Figure 6, and for 345 and 230 kV in Figure 7.

The number of transformers (represented as edges) between different voltage levels for the Eastern Interconnect system is summarized in Table V. Similar to the Western Interconnect, the maximum number of edges (1093) exist between the two lowest levels of voltage -138 and 230 kV. Visualization of the interconnection structure is discussed in Section 4.

#### D. Texas Interconnect Network

The Texas Interconnect has three significant voltage levels: 345, 138 and 69 kV. Key details of the system are summarized in Table VI. The largest portion of the network operates at 138 kV with 2772 vertices and 3274 edges. At this level, there are 3 connected components with the largest component comprising of 2770 vertices. The second largest portion of the network

TABLE IV

DETAILS FROM DECOMPOSING THE EASTERN INTERCONNECT DATASET INTO FIVE VOLTAGE LEVELS. THE VALUES FOR AVERAGE CLUSTERING COEFFICIENT, AVERAGE SHORTEST PATH LENGTH AND DIAMETER ARE FOR THE LARGEST COMPONENT FOR A GIVEN VOLTAGE LEVEL.

| Voltage Levels             | 765 kV     | 500 kV     | 345 kV       | 230 kV       | 138 kV         |
|----------------------------|------------|------------|--------------|--------------|----------------|
| Size $( V ,  E )$          | (115, 152) | (370, 441) | (1224, 1521) | (6251, 7435) | (16082, 19110) |
| V  of largest component    | 79         | 310        | 1065         | 5688         | 12997          |
| Num components             | 2          | 4          | 18           | 41           | 121            |
| Avg Clustering Coefficient | 0.028      | 0.086      | 0.073        | 0.035        | 0.047          |
| Avg Shortest Path Length   | 5.93       | 18.68      | 24.098       | 73.453       | 134.84         |
| Diameter                   | 16         | 48         | 78           | 211          | 386            |

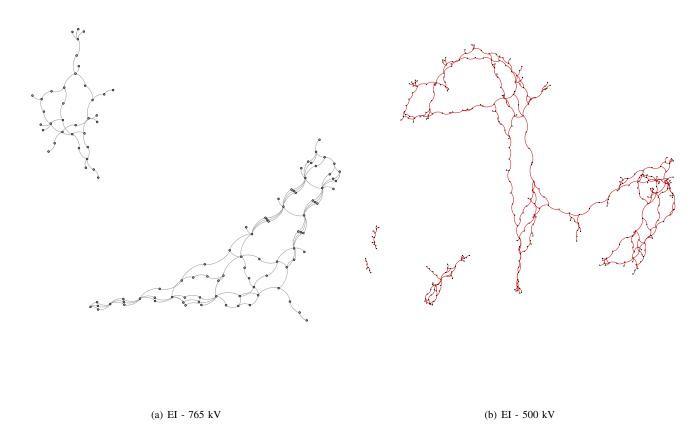


Fig. 6. Rendering of the Eastern Interconnect networks at 765 and 500 kV.

 $\label{table V} TRANSFORMERS IN THE EASTERN INTERCONNECT SYSTEM$ 

|        | 765 kV | 500 kV | 345 kV | 230 kV | 138 kV |
|--------|--------|--------|--------|--------|--------|
| 765 kV | _      | 5      | 38     | 13     | 18     |
| 500 kV | 5      | _      | 9      | 194    | 33     |
| 345 kV | 38     | 9      | _      | 166    | 602    |
| 230 kV | 13     | 194    | 166    | _      | 1093   |
| 138 kV | 18     | 33     | 602    | 1093   | ı      |

operates at 69 kV, with 1749 vertices and 1845 edges. Unlike the network at 138 kV, the network at 69 kV has significantly larger number of connected components – 44 components. Consequently, the existence of a large number of small components are visible in Figure 8(c). Rendering of the networks at voltage levels 345, 138 and 69 kV is provided in Figure 8.

We conclude this section by highlighting the existence of a similar topological structure not only across different voltage levels, but also across different systems. The structure of a network at a given voltage level, decomposed using Algorithm 1, provides insight about the existence of large diameter (and average shortest path length) and low clustering coefficient in power

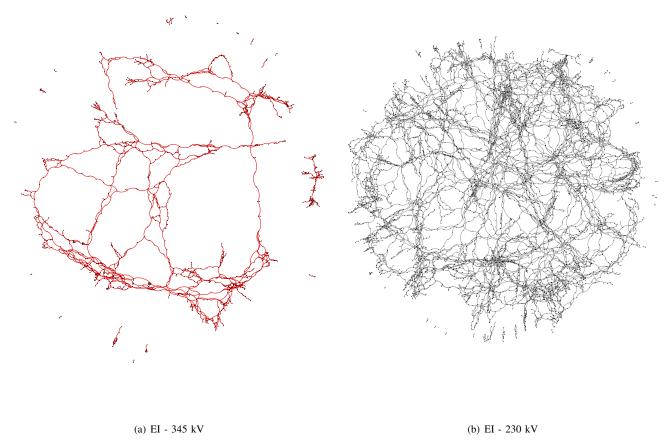


Fig. 7. Rendering of the Eastern Interconnect networks at 345 and 230 kV.

TABLE VI

DETAILS FROM DECOMPOSING THE TEXAS INTERCONNECT DATASET INTO THREE VOLTAGE LEVELS. THE VALUES FOR AVERAGE CLUSTERING
COEFFICIENT, AVERAGE SHORTEST PATH LENGTH AND DIAMETER ARE FOR THE LARGEST COMPONENT FOR A GIVEN VOLTAGE LEVEL.

| Voltage Levels             | 345 kV     | 138 kV       | 69 kV        |
|----------------------------|------------|--------------|--------------|
| Size $( V ,  E )$          | (210, 291) | (2772, 3274) | (1749, 1845) |
| Order of largest component | 208        | 2770         | 1100         |
| Num components             | 2          | 3            | 44           |
| Avg Clustering Coefficient | 0.062      | 0.019        | 0.02         |
| Avg Path Length            | 8.32       | 32.466       | 47.22        |
| Diameter                   | 22         | 84           | 134          |

transmission networks.

# 4. Interconnection of Networks

Given a network-of-networks (decomposed) model of transmission network, we compute the interconnection structure of the networks operating at different voltage levels using Algorithm 2. Let  $G' = \{G_i, G_j, ...\}$  be the input (decomposed) graph as computed from Algorithm 1. Let  $G'' = (V, E, X, \phi)$  be the output graph representing the interconnection structure of the decomposed graph G'.

Algorithm 2 proceeds by assigning a new vertex to each component  $G_i \in G'$ . The voltage rating of  $G_i$  is assigned to the representing vertex. An edge is added between two vertices if the corresponding components are interconnected via a power transformer. Consequently, note that in G'', each vertex represents a connected component of the same voltage rating, and each edge represents a transformer in the original graph G. Thus, the graph G'' represents how different regions of similar voltage are interconnected with each other.

We now present the details of the interconnection structure of Western and Eastern Interconnects computed using Algorithm 2. Key details are summrized in Table VII. We note that the interconnection graphs of the Polish and Texas Interconnect are

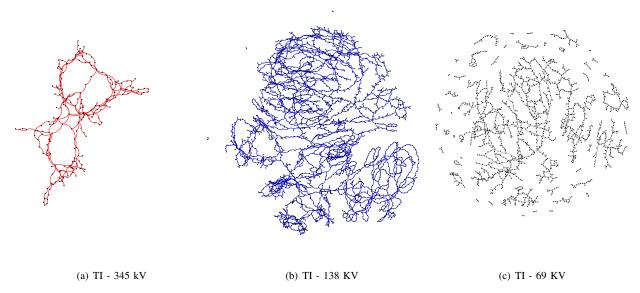


Fig. 8. Rendering of the Texas Interconnect network at 345, 138 and 69 kV.

```
Algorithm 2 Interconnection structure
 1: procedure BUILD-INTERCONNECT-GRAPH(G', G'')
          for each G_i \in G' do
                                                                                                                                      \triangleright Add the vertices to G''
 2:
               V(G'') \leftarrow V \cup \{v_i\}
                                                                                                          \triangleright Add a vertex v_i representing component G_i
 3:
               \phi(G'') \leftarrow \phi \cup \{v_i \rightarrow X(G_i)\}\
 4:
                                                                                                                          \triangleright Voltage rating of component G_i
          for each G_i \in G' do
                                                                                                                                         \triangleright Add the edges to G''
 5:
               for each G_j \in G' such that i \neq j do E(G'') \leftarrow E \cup \{(v_i, v_j)\}
 6:
                                                                              \triangleright Vertices v_i and v_j represent components G_i and G_j respectively
 7:
```

TABLE VII INTERCONNECTION NETWORKS. ASPL STANDS FOR AVERAGE SHORTEST PATH LENGTH.

| Dataset | (Vertics, Edges) | Clust. Coeff. | ASPL | Diameter |
|---------|------------------|---------------|------|----------|
| Western | (497, 565)       | 0.619         | 3.27 | 7        |
| Eastern | (575, 619)       | 0.508         | 2.91 | 7        |

small and unsuitable for analysis. The topological properties of the interconnection networks is substantially different from the properties of the decomposed networks presented in Section 3. While the average clustering coefficient of these networks are higher, the average shortest path and diameter are significantly smaller (relative to decomposed networks). The fundamental differences in the topological structure is also evident from the visualization of the two systems presented in Figure 9.

From the interconnection graphs we observe a hierarchical nature of the manner in which different networks at different voltages interconnect with each other. The lowest voltage levels are used for distribution purposes and generally form degree-one vertices in the interconnection networks. The remaining vertices (networks of higher voltage) form a hierarchical structure spanning the network. In order to highlight this feature, we present a 2-core of the interconnection network of the Eastern Interconnect in Figure 10. The k-core of a graph is a subgraph of the graph such that each vertex has degree k or more.

We will conclude this section by observing that the topological structure of the interconnection network is significantly different from the topological structure of the decomposed networks operating at the same voltage level. Further, there exists a hierarchical structure based on the voltage levels to the manner in which different vertices are linked to one another.

## 5. RANDOM GRAPH MODELS

An important research topic in the study of electric power grids is the ability to synthetically generate graphs that match the characteristics of real-world power grids. We refer the reader to [5] for a detailed analysis of different random graph models and their comparison to real-world datasets. We summarize key findings from previous work in Section 6. In this section, we present two random graph models to generate graphs that match the characteristics of graphs at a given voltage level that were presented in Section 3. We then present graph models that match the characteristics of the interconnection structure that we

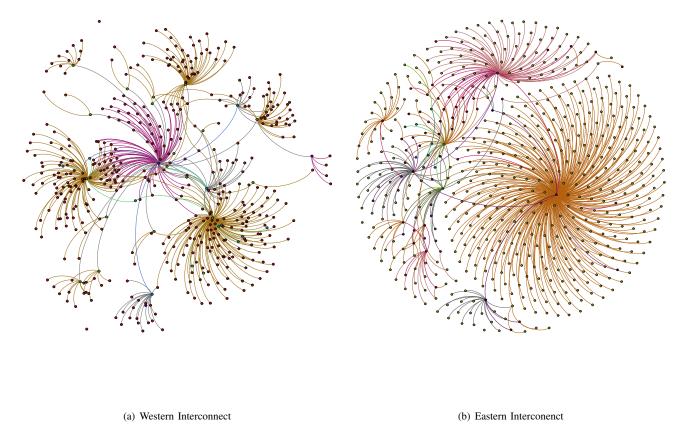


Fig. 9. Rendering of the interconnection network for the Western and Eastern interconnects. For the Western Interconnect, vertices in red color represent networks at 138 kV, green at 230 kV, blue at 345 kV and purple at 500 kV. For the Eastern Interconnect, the color code is as follows: 138 – yellow green, 230 – red, 345 – purple, 500 – blue, and 765 – green.

presented in Section 4. By combining the two separate models, we propose that synthetic graphs can be generated to match the overall characteristics of real-world power grids.

We first present two random graph models to generate graphs at a given voltage. These two models can be seen as a variant of the random geometric graphs [7]. A d-dimensional random geometric graph, G(n,r), is a graph formed by randomly placing n vertices in a d-dimensional space and adding an edge between pairs of vertices whose Euclidean distance is less than or equal to r. Sparsity and structure of the generated graph can be controlled by varying the distance parameter r. We note here that we experimented with several random graph models and found that the random geometric graph models to be the most promising [5]. Here, we present two such models. The first model is given in Algorithm 3. The algorithm takes the size of the graph (number of vertices and edges) as the input, and generates a random graph as the output. The algorithm starts generating n two-dimensional points uniformly at random (Line 3 in Algorithm 3). For each vertex i, a set of  $\lfloor \frac{m}{n} \rfloor$  neighbors are chosen by minimizing the Euclidean distance as follows:

$$\min_{j} \quad (x_i - x_j)^2 + (y_i - y_j)^2 
\text{s.t.} \quad j \notin adj(i),$$
(1)

where adj(i) represents the adjacency set of vertex i. An additional edge is generated with with probability  $P = (\frac{m}{n} - \lfloor \frac{m}{n} \rfloor)$  (Line 6 in Algorithm 3).

Graphs generated by Algorithm 3 do not match all the desired properties of power grids. Therefore, we consider an adapted version of this algorithm that modifies how vertices get connected using the distance function given by Equation 1. The modification is driven by a bisection cost  $c_b$ . Unlike the previous algorithm, new nodes can either be connected to an existing node through a new edge, or it can be placed near an existing edge and become a bisecting node for that edge. This change is motivated by economical considerations in the construction of power transmission networks. For example, consider the construction of a new town that needs to connect to the power grid. The town can either be connected via a transmission

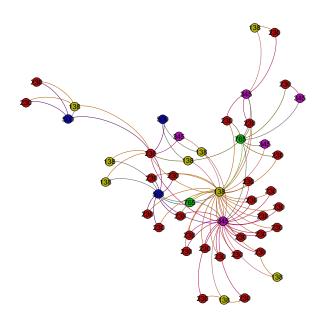


Fig. 10. Rendering of the 2-core of the Eastern interconnection network. Color code is as follows: 138 – yellow green, 230 – red, 345 – purple, 500 – blue, and 765 – green.

**Algorithm 3** Simple minimum-distance algorithm. **Input:** Parameters for the number of vertices and edges (n and m). **Output:** A random graph G.

```
1: procedure SIMPLE-MINDIST-GRAPH((n, m), G)
```

2: **for** i = 1 : n **do** 

3:

- Randomly generate planar coordinates for i ( $x_i$ ,  $y_i$ ) from a uniform distribution
- 4: **for** i = 1 : n **do**
- 5: Generate  $\lfloor \frac{m}{n} \rfloor$  edges (i, j) by iteratively selecting j to minimize the Euclidean distance between i and j using Equation 1
- 6: Generate an additional edge with probability  $P = (\frac{m}{n} \lfloor \frac{m}{n} \rfloor)$  using Equation 1

line or feeder to an existing substation, or a new substation could be built, bisecting an existing transmission line. With this modification, the min-dist selection criteria (Eq. 1) is modified as follows:

min 
$$(C_1, C_2)$$
  
s.t.  $C_1 = \min_{j \notin adj(i)} (x_i - x_j)^2 + (y_i - y_j)^2$   
 $C_2 = \min_{k \in \{1...m\}} d(i \to e_k) + c_j,$  (2)

where  $d(i \to e_k)$  is the distance between point i and the nearest point along line segment  $e_k$ , and  $e_j$  is an exogenously selected bisection cost. If  $C_1$  is smaller than  $C_2$ , then the algorithm creates a new edge (i, j). Otherwise, the new node bisects the existing edge  $e_k$ .

The second step in generating synthetic transmission networks is to generate random graphs to match the characteristics of the interconnection graphs. We propose the use of preferential attachment (PA) model for this purpose. We generate the

preferential attachment (PA) graphs using the algorithm proposed in [8]. Since we aim to generate graphs of a given size (number of vertices, n, and edges m), we modify the algorithm in [8] to allow for a fractional average degree < k >. For each new vertex i,  $\lfloor \frac{m}{n} \rfloor$  edges are added by linking i to an existing vertex j, where j is chosen randomly from the probability distribution  $P = k_j / \sum_{c=1}^n k_c$ , where  $k_c$  is the degree of vertex c. Further, an additional edge (i,j) is added with probability  $\frac{m}{n} - 1$  to generate a preferential attachment graph with n vertices and approximately m edges. We generate Erdős-Rényi (ER) random graphs by randomly adding edges between a pair of vertices until there are exactly m edges in the graph. Further details are provided in [5]. In Figure 11, we provide the cumulative degree distribution of the interconnection graphs of the Western and Eastern Interconnects (described in Section 4) along with degree distributions from random graphs – preferential attachment (PA) and Erdős-Rényi (ER) of the same sizes. We observe that preferential attachment provides a better fit to the degree distribution of interconnection graphs.

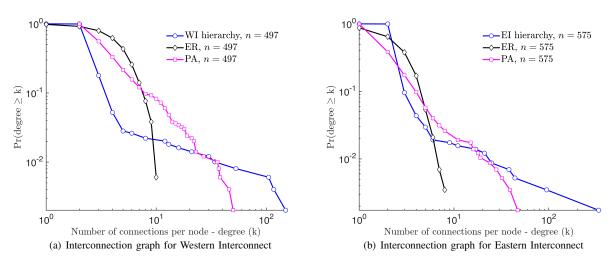


Fig. 11. Cumulative degree distribution of the interconnection graphs of the Western and Eastern Interconnects, along with degree distributions from random graphs – Erdős-Rényi (ER) and preferential attachment (PA).

We conclude this section by noting that accurate models for synthetic recreation of transmission networks should include several aspects such as geography, cost of adding vertices and edges, size and distribution of power transformers, and the inherent hierarchy of networks of different voltage ratings. While we present a multi-step process involving different random graph models, careful construction and validation is part of our work in the near future.

## 6. RELATED WORK

Many researchers have explored the use of random graphs to model power grid [5], [9], [10], [11], [12], [13], but in each case the resulting random graph does not fully model the structure of interest. In most cases, authors attempt to match graph characteristics from power grid networks, such as degree distribution, average path length, and clustering coefficient, in their random models.

In [9] the author examines how vulnerability tests on the Western Interconnect (WECC) and the Nordic grid, and then compare with the same tests on an Erdős-Rényi (ER) random graph and the scale-free model of Barabási and Albert [14]. They first note that the clustering coefficient and average path length of the real graphs are significantly larger than that of both types of random graphs. When comparing vulnerability to failure of generators (i.e., removal of vertices) they see that the real networks show more susceptibility than the random networks.

The authors of [10] also show that the topological properties of the Eastern United States power grid and the IEEE 300 system differ significantly from random networks created using preferential-attachment, small-world, and ER models. They provide a new model, the minimum distance graph, which more closely matches measured topological characteristics like degree distribution, clustering coefficient, diameter, and assortativity of the real networks. We build on these ideas in our work. Extensions of the minimum-distance model are discussed in Section 5.

In a more recent paper by some of the same authors [10] they look more specifically at both the electrical and topological connectivity of three North American power infrastructures. They again compare these with random, preferential-attachment, and small-world networks and see that these random networks differ greatly from the real power networks. They propose to represent electrical connectivity of power systems using electrical distances rather than physical connectivity and geographic connections. In particular, they propose a distance based on sensitivity between active power transfers and nodal phase angle differences. Electrical distance is calculated as a positive value for all pairs of vertices which then yields a complete weighted graph. Therefore, to make it more comparable to a geographic network the authors use a threshold value and keep those edges which are below the threshold.

Pagani and Aiello, in [11], present a survey of the most relevant work using complex network analysis to study the power grid. In this survey they remark that most of the networks studied were High Voltage from either North America, Europe, or China. They remark that most of the surveyed work show that degree distribution is exponential. The geography of the country seems to be important as the results differ somewhat between countries with differing geographies. One point of agreement between all studies is that the power grid is resilient to random failures but extremely vulnerable to attacks targeting nodes with high degree or high betweenness scores.

Finally, we cite the work in [13] in which the authors characterize many graph measures in the context of power grid graphs. As an example, they show that power grids are sparsely connected, the degree distribution has an exponential tail, and the line impedance has a heavy-tailed distribution. Based on their findings they propose an algorithm to generate random power grids which features the same topology and electrical characteristics discovered from real data. They, like us, take a hierarchical approach to generating synthetic power networks. However, their approach differs by looking at geographic zones, whereas our approach is to break up the network by voltage level. Our approach requires no geographic knowledge of the system and leads to a systematic approach for annotating the nodes and edges with different electrical properties such as voltage ratings.

# 7. CONCLUSIONS AND FUTURE WORK

Graph-theoretic analysis of power grids has often resulted in misleading conclusions. In this paper, we hypothesized that one of the reasons for such conclusions is the inability to account for heterogeneity in a power grid. We therefore developed a new method for decomposition based on nominal voltage rating, and using power transformers to generate a network-of-networks model of power transmission networks. While the individual networks operating at the same voltage level are characterized by exponential degree distribution, large values for average shortest-path length and diameter, and low clustering coefficient; the interconnection structure representing the interconnection of these networks is characterized by non-exponential degree distribution, small values for average shortest-path length and diameter, and relatively higher clustering coefficient.

Consequently, we proposed two random graph methods for synthetic generation of power transmission networks. Our approach not only models the topology, but also provides a method for annotating the network with voltage ratings so that the electrical properties of a grid can be modeled correctly. To the best of our knowledge, this approach is novel.

The resulting characterization of power transmission networks and the ability to synthetically recreate networks has important implications for studying the vulnerability of power systems, evolution of power networks, and transmission expansion planning. Thus, the ideas proposed in this paper hold the potential to significantly improve our understanding of several aspects of electric infrastructure networks, as well as other critical infrastructure networks.

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