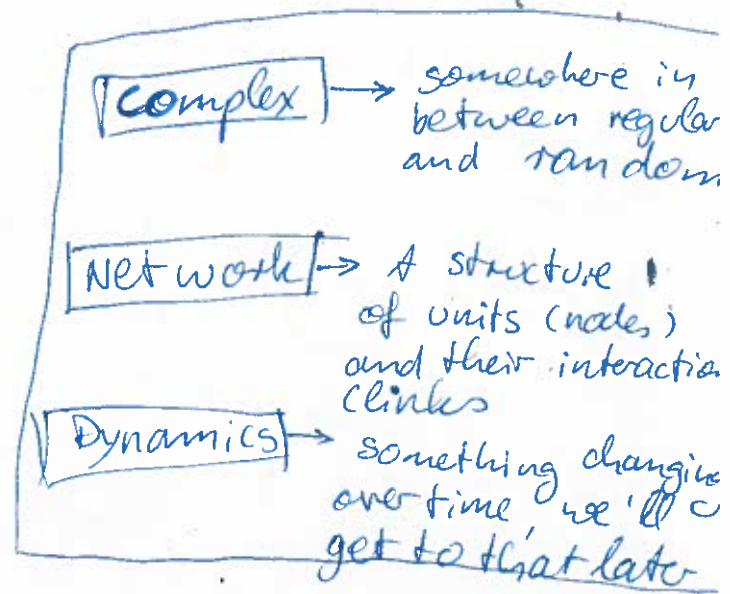


# Complex Network <sup>①</sup>

## DYNAMICS

### Motivation

- networks are everywhere!
- Game: Find a topic that is not a network and we'll try to prove you wrong!
- There is a lot of fun historical and pop-sci motivation on youtube and we have a bunch of links ... somewhere



- Networks sort of make the opposite simplification to almost everywhere else in physics:

### Physics:

Figure out how a thing behaves and how 2 things interact

### Network science:

Never mind what the things and their interactions are, can we already say something by the structure of interactions

Put them together and explain everything but Hyper-graphs

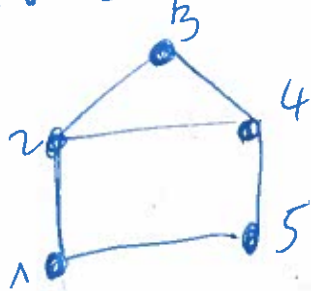
### Reading Material

- Newman Book or review paper
- Statistical Mechanics of Complex Networks
- Barabási
- Spatial Networks
- Erdős-Rényi

Show my funny plot proving that Networks are still growing and still cool

### ③ WALKS, PATHS & CONNECTIVITY

walk : sequence of edges



i.e.

1-2-4-3-2-1-2

closed walk → starts & ends on the same node

Path : walk without repeats of vertices or edges

5-1-2-4-3

cycle : closed path

4-3-2-4

Trail : vertices may repeat but edges not

4-2-3-4-5-1

Just so you know the lingo, mostly we'll stick to paths and leave the other terms to the Graph theorists.

### Connectivity

an undirected Graph  $G$  is said to be **CONNECTED** if every pair  $u, v$  of vertices is connected, that is if  $G$  contains a path from  $u$  to  $v$



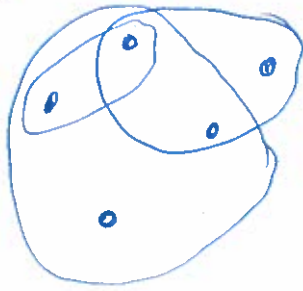
"disconnected"  
not connected



connected

Hyper-graph: Graph where edges "hyper-edges" are types of several vertices (5)

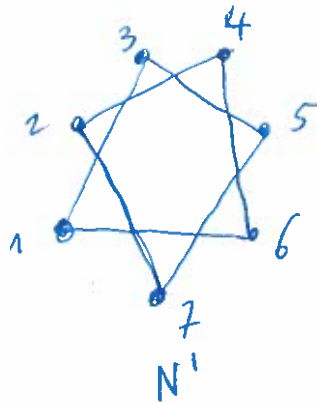
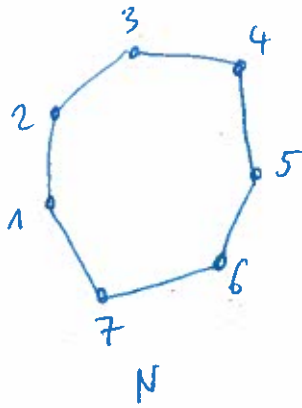
↳



haven't really found a use for this, but they get brought up occasionally

## Network MEASURES

Comparing several different networks can ~~be~~ quickly get out of hand as the network grows  $\Rightarrow$  we need aggregated quantities for the quantitative understanding



First idea: count the differences

Hamming distance:

$$H(N, N') = \sum_{i,j} |A_{ij}^N - A_{ij}^{N'}| = 14$$

$\Rightarrow$  However often we don't even care about the numbering of the nodes

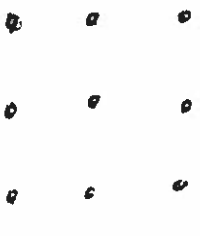
# Quantitatively understanding Networks

①

How can we describe and compare networks more simply?

What is the difference between?

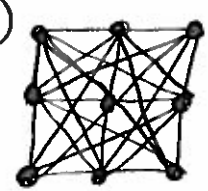
a)  $N=9$



$M=0$   
 $P=0$   
 $\langle k \rangle = 0$

and

b)  $N=9$



$M = \frac{N \cdot (N-1)}{2} = 9 \cdot 4 = 36$   
 $P = 1$   
 $\langle k \rangle = 8$

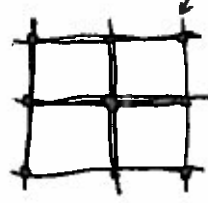
$\Rightarrow$

# links  $M$   
or link density  
 $P = \frac{M}{M_{\max}} = \frac{M}{N \cdot (N-1)/2}$

$\langle k \rangle = \frac{2M}{N}$

What is the difference between?


a)



$2M = N \cdot \langle k \rangle$   
 $= 4 \cdot 9 = 36$   
 $M = 18$

and

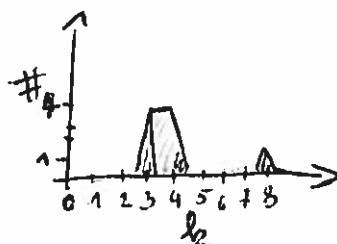
b)



$M = 18$

Peripatete rumbelungen

$\Rightarrow$  Degree distribution



there are

Global network measures Describing the network as a whole in a scalar or a vector (the adjacency matrix can be viewed as a matrix global network measure)



# Network Ensembles

①

Reasons why just one adjacency matrix doesn't tend to be of much use:

SYSTEMS in Nature tend to be

- \* fluctuating / evolving
- \* completely or partially unknown
- \* probabilistic / random in some way
- \* exist in many realizations

→ often we're not actually interested in the one true Network but rather in a TYPE OF NETWORKS

- Examples:
- Power Grid - like networks
  - friendship - like networks
  - protein - like networks

---

DEF: Network Ensemble

A set of possible networks together with a probability distribution of their occurrence

---

ERDŐS - Renyi - RANDOM Graphs

widely used as Null-Model, we will be computing a bunch of things about it analytically today.

a) a graph with  $N$  nodes and  $M$  edges is uniformly chosen from all such graphs

$G(N, M)$

b) each edge is realized with probability  $p$

$G(N, p)$

## Degree Distribution

the probability for a node to be connected to a specific set of  $k$  other nodes is

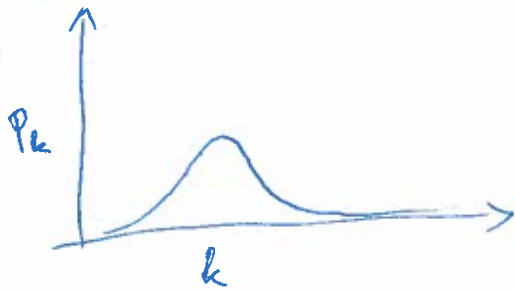
$$p^k (1-p)^{n-1-k}$$

# ways to select a set of  $k$  other nodes:

$$\binom{n-1}{k}$$

$$\Rightarrow \boxed{P_k = \binom{n-1}{k} p^k (1-p)^{n-1-k}}$$

Binomial degree distribution



---

## Clustering Coefficient

$C \hat{=}$  probability that two neighbours of a vertex are also neighbours of each other

$\Rightarrow C = p$  all node pairs have this very same connection probability

---

Other features

many other things can be analytically derived about ER-networks

- Size of the Giant component  
↳ percolation seminar topic
- Component size distribution

↳ see Newman & some new works from 2019

# The Configuration model

(5)

For some purposes you may not want everything about your network to be random a common thing to fix is the degree distribution or degree sequence.

- given a bunch of nodes with "stubs" or "half-edges", combine them randomly to obtain a realization of the configuration model  
(# stubs must be even)



Beware! this model has self-edges & multi-edges

Link probability between nodes  $i$  and  $j$

$$P_{ij} = \frac{k_i k_j}{2m-1}$$

say we start from node  $i$ . The probability that a specific one of the stubs of node  $i$  is connected to node  $j$  is  $\frac{\# \text{ stubs of } j}{\# \text{ all other stubs}} = \frac{k_j}{2m-1}$

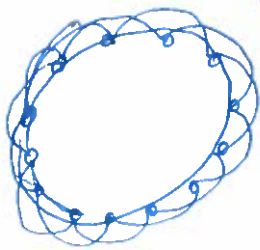
and we get  $k_i$  chances.

Random  $d$ -regular graph

configuration model where all nodes have degree  $d$

# Rewiring or adding links

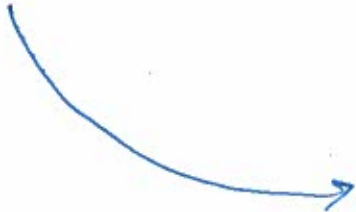
⑦



go through all edges, remove each with prob.  $P$  and add a random one instead

rewiring

preserves  $m$



just add random edges with probability  $P$  for each existing edge  
 $m \cdot P$  shortcuts are made  $\Rightarrow \langle k \rangle P$  per node

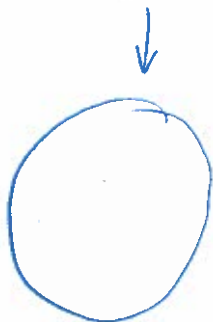
shortcuts

easier but increases  $m$

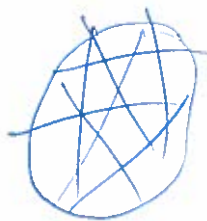
$P=0$

$P=1$

$0 < P < 1$



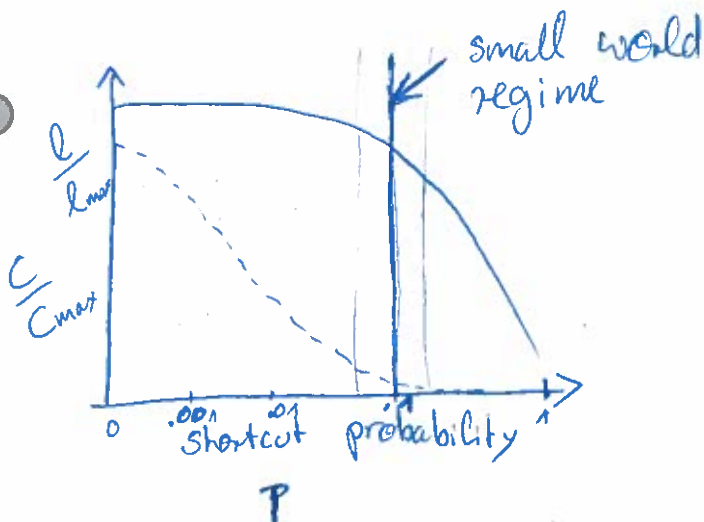
original circle



completely random



circle with some shortcuts



## Degree distribution

at  $P=0$  it's regular, all nodes have degree  $\langle k \rangle$   
 as we add shortcuts it becomes  $\langle k \rangle + \# \text{shortcuts}$   
 (Binomial distribution)



# The ensemble Zoo

①

More on Watts-Strogatz Networks

Graph Laplacian

$$L = -D + A - \begin{pmatrix} k_1 & & \\ & \ddots & \\ & & k_n \end{pmatrix} + \begin{pmatrix} 0 & 1 & 0 \\ 1 & 0 & 1 \\ 0 & 1 & 0 \end{pmatrix}$$

→ more on this in 2 to 4 weeks

• Diffusion, Synchronization etc.

↳ predict Eigenvalue spectrum of WS networks

first (and for our considerations main) step:  
Constructing the mean field Laplacian

$p=0$



original  
Network

rewiring  
fading

$0 < p < 1$



$p=1$



$$L^{MF} = \begin{pmatrix} c_0 & c_1 & c_2 & \dots & c_{N-1} \\ c_{N-1} & & & & \\ & c_2 & & & \\ & & \ddots & & \\ & & & c_1 & \\ c_1 & & & & c_0 \\ & & & c_{N-1} & \end{pmatrix}$$

this structure is called  
"Circulant Laplacian"

because the only thing that  
determines the entries of  $A^{MF}$   
and thus  $L^{MF}$  is the distance  
of two nodes along the  
circle

$$C_i^{p=0} = \begin{cases} -k & \text{if } i=0 \\ 1 & \text{if } i \in \{1, \dots, \frac{k}{2}, N-\frac{k}{2}, \dots, N-1\} := S_1 \text{ (ring)} \\ 0 & \text{if } i \in \{\frac{k}{2}+1, \dots, N-\frac{k}{2}-1\} := S_2 \text{ (off-ring)} \end{cases}$$

average total weight  $p \cdot k \cdot N/2 \stackrel{\text{expected}}{=} \# \text{rewired links}$

is subtracted UNIFORMLY from existing edges of the  
Ring and then distributed uniformly on ALL possible  
edges (on and off ring)

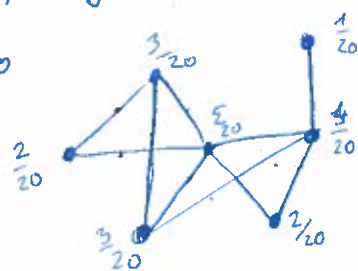
# BARABASI-ALBERT NETWORKS

(3)

Growth + Preferential attachment

(Earlier lesser known model by price for citation growth)

① Start from some network  $m_0$



② Add nodes with degree  $k = k_m$

③ Connect stumps to nodes in the existing network with

$$P_{k_i} = \frac{k_i}{\sum_j k_j} = \frac{k_i}{20}$$

degree distribution of B-A-Networks

1+2+3+4+5

Probability a new vertex attaches to a node of degree  $k$ :

$$\frac{k P_k}{\sum_k k P_k} = \frac{k P_k}{2m}$$

⇒ expected # degree  $k$  nodes to gain a link  
or  $\frac{k P_k}{2m}$

⇒ use this to formulate master equation

$$(n+1)P_{k,n+1} = \underbrace{n P_{k,n}}_{\text{expected \# nodes with degree } k \text{ at step } n} - \underbrace{\frac{1}{2} k P_{k,n}}_{\text{expected \# deg } k \text{ nodes gaining link}} + \underbrace{\frac{1}{2} (k-1) P_{k-1,n}}_{\text{expected \# deg } (k-1) \text{ nodes gaining a link}}$$

for  $k > m$

or

$$(n+1)P_{m,n+1} = n P_{m,n} - \frac{1}{2} m P_{m,n} + 1$$

$\langle k \rangle = Np \Rightarrow$  we can work out  $K(\langle k \rangle)$  or the radius for which we get a desired target mean degree

$p = \frac{\langle k \rangle}{N}$   
 $\downarrow$   
 Equate to above expression & solve for R

$$R = \frac{1}{\sqrt{\pi}} \left( \frac{\langle k \rangle}{N} \Gamma\left(\frac{d+2}{2}\right) \right)^{1/d}$$

## Random Geometric Graphs

- Null-model for embedded networks in space
- "randomized" alternative to grids
- many things can be computed about them
- model for peer to peer networks of devices based on short-range connections such as Wifi or blue tooth "ad-hoc-networks"

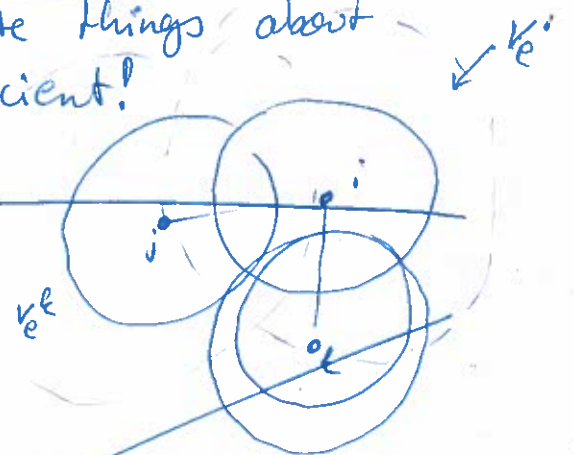
we can still analytically compute things about this! namely the clustering coefficient!

$$\langle C_d \rangle = \frac{1}{V_e} \int_{V_e} P_d(r) dV$$

Fraction of the excluded volume of  $i$

which is also in the excluded volume of  $j$

overlap  $P_d$  of the two spheres



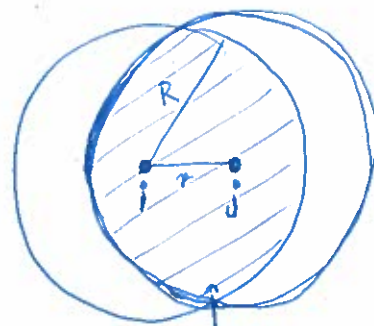
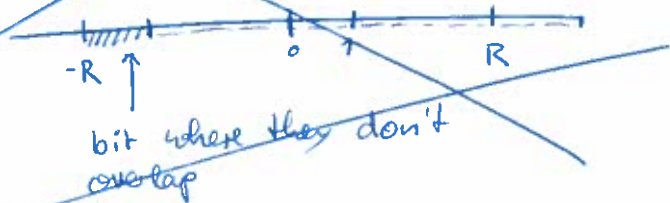
In 1d:

$$P_1(r) = \frac{(2R-r)}{2R}$$

$$= 1 - \frac{r}{2R}$$

$$\langle C_1 \rangle = \frac{1}{2R} \cdot \int_0^{2R} \left(1 - \frac{r}{2R}\right) dr$$

$$= \frac{1}{2R} \cdot R$$



$$C_d(r) = \frac{P(r)}{V_e}$$

Fractional Volume overlap =  $P_d(r) = \frac{\text{absolute overlap}}{\text{whole volume}}$



intuition behind this:

⑦

Small  $q$ : "You never have time for me"

lower-degree side feels neglected.

when choosing a new connection however, they still want a popular friend linearly more

large  $q$ : "you're not a sufficiently valuable connection"

higher-degree side thinks they can "do better"

### Results:

→ this process is repeated until an equilibrium in the degree distribution is reached

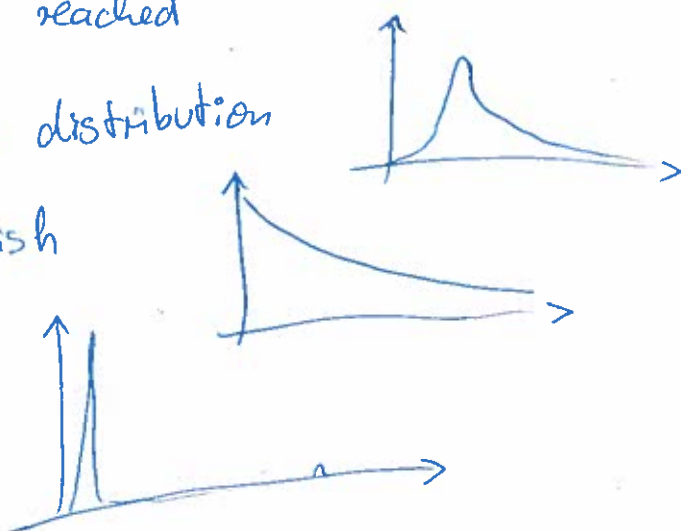
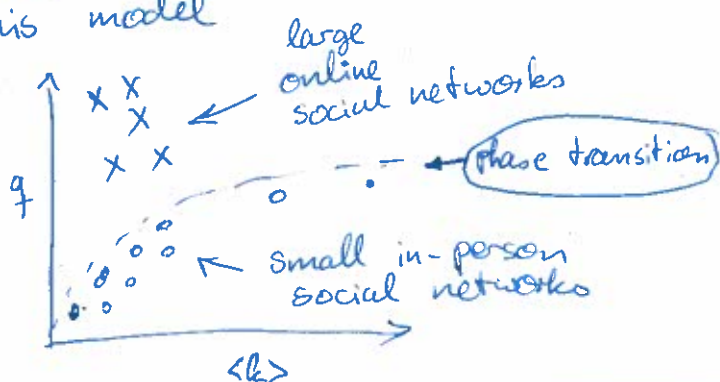
Small  $q$ : peaked degree distribution

large  $q$ : power law-ish

Very large  $q$ : two peaks

$q > 0.98$

real networks fitted with this model



↑  
sort of leader-follower network

## NON-Overlapping Chain

Proteins are chains of amino acids. Amino acids are physical objects that can't overlap indefinitely



# Exponential Random Graphs OR Canonical Network Ensembles

(1)

Defining a network ensemble through its network measure.

→ sort of the opposite of what we did before where we defined it through the formation mechanisms and then computed expected network measures

→ more of a "Statistical mechanics approach"

- fix the desired average value of a network measure

- draw from an ensemble that tends to be close to that value

⇒ this approach first came from social sciences, where you had measurements

Let  $\mathcal{G}$  be the set of all graphs with  $n$  nodes

and  $\sum_{G \in \mathcal{G}} P(G) = 1$

and let  $x_1, x_2, \dots$  be network measures

$$\Rightarrow \langle x_i \rangle = \sum_{G \in \mathcal{G}} P(G) x_i(G)$$

↑  
now we fix this

Now under these constraints we want to MAXIMIZE ENTROPY  
so we are looking for the most generic ensemble satisfying the constraints.

⇒ Use Lagrange Multipliers to maximize Gibbs Entropy under network measure constraints

Lagrange multipliers:

$$\mathcal{L}(x, \lambda) = f(x) - \lambda g(x)$$

← constraint  $g(x) = 0$ :  $g(x) = \left( \sum_{G \in \mathcal{G}} P(G) - 1 \right)$

and  $g(x) = \left( \sum_{G \in \mathcal{G}} P(G) x_i(G) - \langle x_i \rangle \right)$

Gibbs Entropy  $S = - \sum_{G \in \mathcal{G}} P(G) \ln(P(G))$

Simplest Example:  
Erdős-Rényi; again (sort of)

(3)

$H = \beta M$  as  $M$  is what we want to fix

$$\Rightarrow P(G) = \frac{e^{\beta M}}{Z} \quad \text{where } Z = \sum_G e^{\beta M}$$

Now all this summing over  $G$  can be done using adjacency matrices:

$$M = \sum_{i < j} A_{ij}$$

$$Z = \sum_{\{A_{ij}\}} e^{\beta \sum_{i < j} A_{ij}} = \sum_{\{A_{ij}\}} \prod_{i < j} e^{\beta A_{ij}} = \prod_{i < j} \sum_{A_{ij}=0,1} e^{\beta A_{ij}} = \prod_{i < j} (1 + e^{\beta})$$

all allowed values of  $A_{ij}$   $\Rightarrow (1 + e^{\beta})^{\binom{n}{2}}$

$\Rightarrow$  compute free Energy  $F = \ln(Z) = \binom{n}{2} \ln(1 + e^{\beta})$

$$\langle M \rangle = \frac{\partial F}{\partial \beta} = \frac{e^{\beta}}{1 + e^{\beta}} \binom{n}{2} = \binom{n}{2} \frac{1}{e^{-\beta} + 1}$$

if we have a particular value  $\langle M \rangle$  in mind we should fix  $\beta$  to

$$\beta = \ln \left( \frac{\langle M \rangle}{\binom{n}{2} - \langle M \rangle} \right)$$

probability for node-pair to be connected:

$$P_{ij} = \langle A_{ij} \rangle = \frac{1}{Z} \sum_{\{A_{ij}\}} A_{ij} e^{\beta \sum_{i < j} A_{ij}} = \frac{\sum_{A_{ij}=0,1} A_{ij} e^{\beta A_{ij}}}{\sum_{A_{ij}=0,1} e^{\beta A_{ij}}}$$

$$= \frac{(0 + e^{\beta})}{(1 + e^{\beta})} = \frac{1}{1 + e^{-\beta}} = \frac{\langle M \rangle}{\binom{n}{2}}$$

just like ER

# RELATIVE CANONICAL NETWORK ENSEMBLES

(5)

so far we've maximized Entropy, however, is that reasonable?  
why would Entropy always be maximal?

⇒ Real world networks may not actually be completely maximized in Entropy but relatively maximized to something we may know about the formation process, which is the

## Background Ensemble

- Essentially the Background ensemble is what you get if you run the Ensemble for infinite temperature.  
(It's practically very convenient to define it like this!)

$$P_{R,q}^R = \frac{e^{-\beta R(G)}}{Z_{R,q}(\beta)} q(G)$$

$R$ : Property function

$\beta$ : inverse temperature/genericity

$q$ : Background ensemble

## RCNE's as a tool for analyzing network measures

- while Canonical network models are a good way of constructing network ensembles based on network measures and Background ensembles, their realism can be limited and the amount of things you can actually compute can be ~~very~~ constrained by computational complexity
- But they can actually be used as a way to analyze the NETWORK MEASURES THEMSELVES

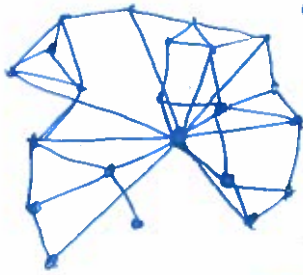
what about another measure?

(7)

$$R_{WL} = \frac{W \cdot L}{N} \leftarrow \text{Average shortest path length}$$

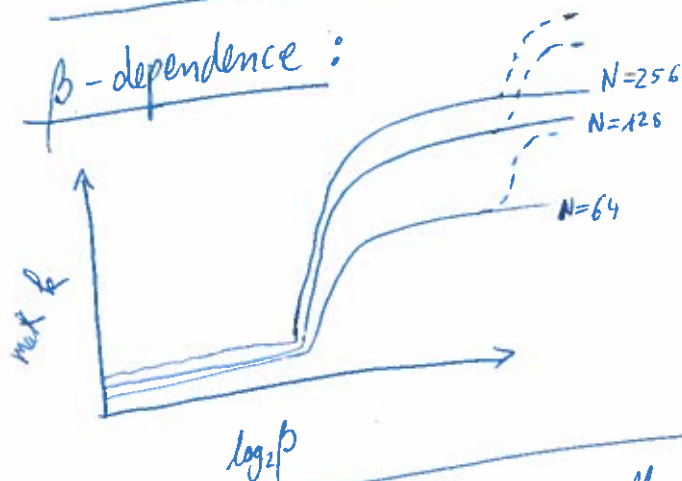
total Euclidean length of all links

→ Better but still no



one Hub with long links, short links otherwise

$\beta$ -dependence:



phase transition  
at some value  
of  $\beta$   
(in all 4 cases)

→ sometimes actually 2 phase transitions  
one for clique and one for hub-formation

Problems:  
15.3