

# A\* algorithm

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The a\* algorithm combines features of uniform-cost search and pure heuristic search to efficiently compute optimal solutions. A\* algorithm is a best-first search algorithm in which the cost associated with a node is  $f(n) = g(n) + h(n)$ , where  $g(n)$  is the cost of the path from the initial state to node  $n$  and  $h(n)$  is the heuristic estimate or the cost of a path from node  $n$  to a goal. Thus,  $f(n)$  estimates the lowest total cost of any solution path going through node  $n$ . At each point a node with lowest  $f$  value is chosen for expansion. Ties among nodes of equal  $f$  value should be broken in favor of nodes with lower  $h$  values. The algorithm terminates when a goal is chosen for expansion.

A\* algorithm guides an optimal path to a goal if the heuristic function  $h(n)$  is admissible, meaning it never overestimates actual cost. For example, since airline distance never overestimates actual highway distance, and manhattan distance never overestimates actual moves in the gliding tile.

For puzzle, a\* algorithm, using these evaluation functions, can find optimal solutions to these problems. In addition, a\* makes the most efficient use of the given heuristic function in the following sense: among all shortest-path algorithms using the given heuristic function  $h(n)$ . A\* algorithm expands the fewest number of nodes.

The main **drawback** of a\* algorithm and indeed of any best-first search is its memory requirement. Since at least the entire open list must be saved, a\* algorithm is severely space-limited in practice, and is no more practical than best-first search algorithm on current machines. For example, while it can be run successfully on the eight puzzle, it exhausts available memory in a matter of minutes on the fifteen puzzle.