

**Grammar**

BASE TYPE		TYPE		LABEL	
$b$	$::=$	$\tau$	$::=$	$\ell$	$::=$
	BOOL		$b$		PUBLIC
	UINT <sub><math>n</math></sub>		ARR $\langle b, n \rangle$		SECRET
	INT <sub><math>n</math></sub>				
EXPRESSION					
		$e$	$::=$		
				TRUE	
				FALSE	
				$c$	integer literal
				$x$	variable
				$x[e]$	array get
				$\langle b, n, b_x \rangle x \Rightarrow e$	array comprehension
				VIEW( $x, e, n$ )	array view
				$\ominus e$	unary op
				$e_1 \oplus e_2$	binary op
				$e_1 ? e_2 : e_3$	ternary op
				REF $x$	mut ref
				$f(e_1, \dots, e_n)$	function call
STATEMENT					
		$s$	$::=$		
				$s_1; s_2$	sequence
				$\langle \tau, \ell, \sigma \rangle x := e$	variable declaration
				$x := e$	variable assignment
				$x[e_1] = e_2$	array assignment
				IF $e \{s_1\}$ ELSE $\{s_2\}$	conditional
				FOR $\langle b \rangle x$ FROM $e_1$ TO $e_2 \{s\}$	loop
				RETURN $e$	return
FUNCTION DEFINITION					
		$fdec$	$::=$		
				$\langle b, \ell \rangle f(\langle \tau_1, \ell_1, \sigma_1 \rangle x_1, \dots, \langle \tau_n, \ell_n, \sigma_n \rangle x_n) \{s\}$	

**Metavariables**

TYPE CONTEXT		VARIABLE TYPE STORE	
$\Gamma$	$::=$	$\mu$	$::=$
	$\emptyset$		$\emptyset$
	$\Gamma[e \mapsto \langle \tau, \ell, \sigma \rangle]$		$\mu[x \mapsto \langle \tau, \ell, \sigma \rangle]$
FUNCTION TYPE STORE			
$\mathbb{F}$	$::=$		
		$\emptyset$	
		$\mathbb{F}[f \mapsto fdec(\langle \tau_1, \ell_1, \sigma_1 \rangle, \dots, \langle \tau_n, \ell_n, \sigma_n \rangle) : \langle b, \ell \rangle]$	

**Type Lattice**

$$\begin{array}{c}
\frac{n_1 < n_2}{\text{UINT}_{n_1} <_{\tau} \text{UINT}_{n_2}} \quad \frac{n_1 < n_2}{\text{INT}_{n_1} <_{\tau} \text{INT}_{n_2}} \quad \frac{}{\text{UINT}_n <_{\tau} \text{INT}_{2n}} \quad \frac{}{\text{PUBLIC} <_{\ell} \text{SECRET}} \\
\\
\frac{}{\text{MUT} <_{\sigma} \text{CONST}} \quad \frac{\Gamma \mid \mu \vdash e : \langle \tau_1, \ell_1, \sigma_1 \rangle \quad \tau_1 \leq_{\tau} \tau_2 \quad \ell_1 \leq_{\ell} \ell_2 \quad \sigma_1 \leq_{\sigma} \sigma_2}{\Gamma \mid \mu \vdash e : \langle \tau_2, \ell_2, \sigma_2 \rangle} \quad \frac{}{\ell \cup \ell = \ell} \\
\\
\frac{}{\ell \cup \text{SECRET} = \text{SECRET}} \quad \frac{}{\text{SECRET} \cup \ell = \text{SECRET}}
\end{array}$$

**Expressions**

$$\boxed{\Gamma \mid \mu \vdash e : \langle \tau, \ell, \sigma \rangle}$$

$$\begin{array}{c}
\text{VAR} \quad \frac{\mu(x) = \langle \tau, \ell, \sigma \rangle}{\Gamma \mid \mu \vdash x : \langle \tau, \ell, \text{CONST} \rangle} \quad \text{UNOP} \quad \frac{\Gamma \mid \mu \vdash e : \langle \tau_1, \ell_1, \sigma_1 \rangle \quad \ominus : \langle \tau_1, \ell_1, \sigma_1 \rangle \rightarrow \langle \tau_2, \ell_2, \sigma_2 \rangle}{\Gamma \mid \mu \vdash \ominus e : \langle \tau_2, \ell_2, \sigma_2 \rangle} \\
\\
\text{BINOP} \quad \frac{\Gamma \mid \mu \vdash e_1 : \langle \tau_1, \ell_1, \sigma_1 \rangle \quad \Gamma \mid \mu \vdash e_2 : \langle \tau_2, \ell_2, \sigma_2 \rangle \quad \oplus : \langle \tau_1, \ell_1, \sigma_1 \rangle \rightarrow \langle \tau_2, \ell_2, \sigma_2 \rangle \rightarrow \langle \tau_3, \ell_3, \sigma_3 \rangle}{\Gamma \mid \mu \vdash e_1 \oplus e_2 : \langle \tau_3, \ell_3, \sigma_3 \rangle} \\
\\
\text{TERNOP} \quad \frac{\Gamma \mid \mu \vdash e_1 : \langle \tau_1, \ell_1, \sigma_1 \rangle \quad \Gamma \mid \mu \vdash e_2 : \langle \tau_2, \ell_2, \sigma_2 \rangle \quad \Gamma \mid \mu \vdash e_3 : \langle \tau_3, \ell_3, \sigma_3 \rangle \quad ? : \langle \tau_1, \ell_1, \sigma_1 \rangle \rightarrow \langle \tau_2, \ell_2, \sigma_2 \rangle \rightarrow \langle \tau_3, \ell_3, \sigma_3 \rangle \rightarrow \langle \tau_4, \ell_4, \sigma_4 \rangle}{\Gamma \mid \mu \vdash e_1 ? e_2 : e_3 : \langle \tau_4, \ell_4, \sigma_4 \rangle} \\
\\
\text{ARRGET} \quad \frac{\mu(x) = \langle \text{ARR}\langle b, n \rangle, \ell, \sigma \rangle \quad \Gamma \mid \mu \vdash e : \langle \text{UINT}, \text{PUBLIC}, \text{CONST} \rangle \quad \text{SMT}(e < n)}{\Gamma \mid \mu \vdash x[e] : \langle b, \ell, \text{CONST} \rangle} \\
\\
\text{ARRCOMP} \quad \frac{\Gamma \mid \mu[x \mapsto \langle b_x, \ell, \text{CONST} \rangle] \vdash e : \langle b, \ell, \text{CONST} \rangle}{\Gamma \mid \mu \vdash \langle b, n, b_x \rangle x \Rightarrow e : \langle \text{ARR}\langle b, n \rangle, \ell, \text{MUT} \rangle} \\
\\
\text{ARRVIEW} \quad \frac{\mu(x) = \langle \text{ARR}\langle b, n \rangle, \ell, \sigma \rangle \quad \Gamma \mid \mu \vdash e : \langle \text{UINT}, \text{PUBLIC}, \text{CONST} \rangle \quad \text{SMT}(e + n' < n)}{\Gamma \mid \mu \vdash \text{VIEW}(x, e, n') : \langle \text{ARR}\langle b, n' \rangle, \ell, \sigma \rangle} \\
\\
\text{MUTREF} \quad \frac{\mu(x) = \langle \tau, \ell, \text{MUT} \rangle}{\Gamma \mid \mu \vdash \text{REF } x : \langle \tau, \ell, \text{MUT} \rangle} \quad \text{FNCALL} \quad \frac{\mathbb{F}(f) = fdec(\langle \tau_1, \ell_1, \sigma_1 \rangle, \dots, \langle \tau_n, \ell_n, \sigma_n \rangle) : \langle b, \ell \rangle \quad \Gamma \mid \mu \vdash e_1 : \langle \tau_1, \ell_1, \sigma_1 \rangle \quad \dots \quad \Gamma \mid \mu \vdash e_n : \langle \tau_n, \ell_n, \sigma_n \rangle}{\Gamma \mid \mu \vdash f(e_1, \dots, e_n) : \langle b, \ell, \text{CONST} \rangle} \\
\\
\text{TRUE} \quad \frac{}{\Gamma \mid \mu \vdash \text{TRUE} : \langle \text{BOOL}, \text{PUBLIC}, \text{CONST} \rangle} \\
\\
\text{FALSE} \quad \frac{}{\Gamma \mid \mu \vdash \text{FALSE} : \langle \text{BOOL}, \text{PUBLIC}, \text{CONST} \rangle} \quad \text{POSNUMBER} \quad \frac{c \geq 0 \quad n = \lceil \log_2 c \rceil}{\Gamma \mid \mu \vdash c : \langle \text{UINT}_n, \text{PUBLIC}, \text{CONST} \rangle} \\
\\
\text{NEGNUMBER} \quad \frac{c < 0 \quad n = \lceil \log_2 |c| \rceil + 1}{\Gamma \mid \mu \vdash c : \langle \text{INT}_n, \text{PUBLIC}, \text{CONST} \rangle}
\end{array}$$

## Statements

$$\langle \mu, \ell_s, r? \rangle \vdash s \rightarrow \langle \mu', \ell'_s, r'_? \rangle$$

$$\text{SEQ} \quad \frac{\Delta \vdash s_1 \rightarrow \Delta' \quad \Delta' \vdash s_2 \rightarrow \Delta''}{\Delta \vdash s_1; s_2 \rightarrow \Delta''}$$

$$\text{VARDEC} \quad \frac{x \notin \text{Dom}(\mu) \quad \ell_s \leq_\ell \ell \quad \Gamma \mid \mu \vdash e : \langle \tau, \ell, \sigma \rangle}{\langle \mu, \ell_s, r? \rangle \vdash \langle \tau, \ell, \sigma \rangle x := e \rightarrow \langle \mu \mapsto x : \langle \tau, \ell, \sigma \rangle, \ell_s, r? \rangle}$$

$$\text{VARDEC}^* \quad \frac{x \notin \text{Dom}(\mu) \quad \ell_s \leq_\ell \ell \quad \Gamma \mid \mu \vdash e : \langle b, \ell, \text{CONST} \rangle}{\langle \mu, \ell_s, r? \rangle \vdash \langle b, \ell, \text{MUT} \rangle x := e \rightarrow \langle \mu \mapsto x : \langle b, \ell, \text{MUT} \rangle, \ell_s, r? \rangle}$$

$$\text{VARASSIGN} \quad \frac{\mu(x) = \langle b, \ell, \text{MUT} \rangle \quad \Gamma \mid \mu \vdash e : \langle b, \ell, \text{CONST} \rangle}{\Delta \langle \ell_s \rangle \vdash x := e : \text{PUBLIC}}$$

$$\text{ARRASSIGN} \quad \frac{\mu(a) = \langle \text{ARR} \langle b, n \rangle, \ell_1, \text{MUT} \rangle \quad \Gamma \vdash e_1 : \langle \text{UINT}_{max}, \text{PUBLIC} \rangle \quad \Gamma \vdash e_2 : \langle b, \ell_2 \rangle \quad \ell_2 \leq_\ell \ell_1}{\Delta \langle \ell_s \rangle \vdash a[e_1] := e_2 : \text{PUBLIC}}$$

$$\text{IF} \quad \frac{\Gamma \vdash e : \langle \text{BOOL}, \ell \rangle \quad \Delta \langle \ell \cup \ell_s \rangle \vdash s_1 : \ell'_s \quad \Delta \langle \ell \cup \ell_s \rangle \vdash s_2 : \ell''_s}{\Delta \langle \ell_s \rangle \vdash \text{IF } e \{s_1\} \text{ ELSE } \{s_2\} : \ell'_s \cup \ell''_s}$$

$$\text{FOR} \quad \frac{\Gamma \vdash e_1 : \langle b, \text{PUBLIC} \rangle \quad \Gamma \vdash e_2 : \langle b, \text{PUBLIC} \rangle \quad b = \text{UINT}_s \vee b = \text{INT}_s \quad \Delta \langle \ell_s \rangle \vdash s : \ell'_s}{\Delta \langle \ell_s \rangle \vdash \text{FOR } \langle b \rangle x \text{ FROM } e_1 \text{ TO } e_2 \{s\} : \ell'_s \quad \mu(x) = \langle b, \text{PUBLIC}, \text{CONST} \rangle_{(\text{scoping?})}}$$

$$\text{RET} \quad \frac{\Gamma \vdash e : \langle b, \ell_1 \rangle \quad \mathbb{F}(f) = fdec : \langle b, \ell_2 \rangle \quad \ell_1 \leq_\ell \ell_2}{\Delta \langle \ell_s \rangle \vdash \text{RETURN } e : \ell_s}$$