

$c$	numeric value
$a$	array
$x$	term variable
$f$	function name
$n$	index variable
$k$	index variable
$i$	index variable
$b_h$	boolean value
$c_h$	numeric value
$f_h$	function name
$rval$	“return value” variable
$rnset$	“rval-not-set” variable

$v$	$::=$ $ $ TRUE $ $ FALSE $ $ $c$ $ $ $a$	values bitmask true (0b1111...) bitmask false (0b0000...) numeric value bytearray
$\ominus$	$::=$ $ $ ~	unary operations bitwise not
$\oplus$	$::=$ $ $ + $ $ - $ $ * $ $ << $ $ >> $ $ & $ $   $ $ == <sub>s</sub> $ $ != <sub>s</sub> $ $ > <sub>s</sub> $ $ < <sub>s</sub> $ $ >= <sub>s</sub> $ $ <= <sub>s</sub>	binary operations      bitwise and bitwise or equals (sign extended)
$\{\sigma\}$	$::=$ $ $ $\{x_1/v_1, \dots, x_k/v_k\}$	variable substitution
$fval$	$::=$ $ $ $(x_1, \dots, x_n) : s @e$	function spec
$fndef$	$::=$ $ $ <b>fdef</b> $f fval$	function definition
$program$	$::=$ $ $ $fndef_1; \dots; fndef_n; \textbf{expose } fndef$	program list of fdefs
$\Lambda$	$::=$ $ $ $\emptyset_\Lambda$ $ $ $\Lambda[f \mapsto fval]$	function store empty function store define function
$\Gamma$	$::=$ $ $ $\emptyset_\Gamma$ $ $ $\Gamma[a \mapsto []]$ $ $ $\Gamma[a \mapsto \Gamma(a)[v_1 \mapsto v_2]]$	global memory   new array array update
$\mu$	$::=$ $ $ $\emptyset_\mu$	local memory empty memory

		$\mu[x \mapsto v]$	add/update variable
		$\mu_1 \triangleright \mu_2$	push stack frame
$\kappa$	::=		program transcript
		$\emptyset_\kappa$	empty transcript
		$\kappa \triangleright \ominus$	add $\ominus$ to transcript
		$\kappa \triangleright \oplus$	add $\oplus$ to transcript
		$\kappa \triangleright \mathbf{load}$	add memory load to transcript
		$\kappa \triangleright \mathbf{store}$	add memory store to transcript
		$\kappa \triangleright f$	add call to $f$ to transcript
		$\kappa \triangleright \mathbf{ret}$	add function return to transcript
$e$	::=		expressions
		TRUE	bitmask true (0b1111...)
		FALSE	bitmask false (0b0000...)
		$c$	numeric value
		$a$	bytearray
		$x$	variable
		$a[e]$	array access
		$\ominus e$	unary operation
		$e_1 \oplus e_2$	binary operation
		$f(e_1, \dots, e_n)$	function application
$s$	::=		statements
		<b>skip</b>	skip
		$s_1; s_2$	sequence
		<b>def</b> $x := e$	variable declaration
		<b>ade</b> $x := a$	array declaration
		$x := e$	variable assignment
		$a[e_1] := e_2$	array assignment
		<b>for</b> $x$ <b>from</b> $v_1$ <b>to</b> $v_2$ : $s$	for loop
$v_h$	::=		values
		$b_h$	boolean value
		$c_h$	numeric value
		$a$	bytearray
$\ominus_h$	::=		unary operations
		!	logical not
		~	bitwise not
$\oplus_h$	::=		binary operations
		+	
		-	
		*	
		<<	
		>>	

		<b>&amp;</b>	
		<b> </b>	
		<b>&amp;&amp;</b>	
		<b>  </b>	
		<b>==</b>	
		<b>!=</b>	
		<b>&gt;</b>	
		<b>&lt;</b>	
		<b>&gt;=</b>	
		<b>&lt;=</b>	
$e_h$	<b>::=</b>		expressions
		$b_h$	boolean value
		$c_h$	numeric value
		$a$	bytearray
		$x$	variable
		$a[e_h]$	array access
		$\ominus_h e_h$	unary operation
		$e_{h1} \oplus e_{h2}$	binary operation
		$f_h(e_{h1}, \dots, e_{hn})$	function application
$s_h$	<b>::=</b>		statements
		<b>skip</b> <sub><math>h</math></sub>	skip
		$s_{h1}; s_{h2}$	sequence
		<b>def</b> <sub><math>h</math></sub> $x := e_h$	variable declaration
		<b>adef</b> <sub><math>h</math></sub> $x := a$	array declaration
		$x := e_h$	variable assignment
		$a[e_{h1}] := e_{h2}$	array assignment
		<b>for</b> <sub><math>h</math></sub> $x$ <b>from</b> $v_{h1}$ <b>to</b> $v_{h2}$	for loop
		<b>if</b> <sub><math>h</math></sub> $e_h$ <b>then</b> $s_{h1}$ <b>else</b> $s_{h2}$	conditional branch
		<b>return</b> <sub><math>h</math></sub> $e_h$	return
$hfval$	<b>::=</b>		function spec
		$(x_1, \dots, x_n) : s_h$	
$hfndef$	<b>::=</b>		function definition
		<b>fdef</b> <sub><math>h</math></sub> $f_h$ $hfval$	
$hprogram$	<b>::=</b>		program
		$hfndef_1; \dots; hfndef_n; \mathbf{expose} \ hfndef$	list of fdefs
$ctx$	<b>::=</b>		branch context
		$x$	variable
		$\ominus ctx$	unary operation
		$ctx_1 \oplus ctx_2$	binary operation

$\{\Lambda, \Gamma, \mu, \kappa\} e \longrightarrow \{\Lambda', \Gamma', \mu', \kappa'\} e'$   $e$  reduces to  $e'$

$$\frac{\begin{array}{l} \mu = \mu'[x \mapsto v] \\ \kappa' = \kappa \triangleright \mathbf{load} \end{array}}{\{\Lambda, \Gamma, \mu, \kappa\} x \longrightarrow \{\Lambda, \Gamma, \mu, \kappa'\} v} \text{ EXR\_VAR}$$

$$\begin{array}{c}
\frac{\{\Lambda, \Gamma, \mu, \kappa\} e \longrightarrow \{\Lambda, \Gamma, \mu, \kappa'\} e'}{\{\Lambda, \Gamma, \mu, \kappa\} a[e] \longrightarrow \{\Lambda, \Gamma, \mu, \kappa'\} a[e']} \quad \text{EXR\_ARR\_GET\_EXPR} \\
\\
\frac{\begin{array}{c} v' = \Gamma(a)[v] \\ \kappa' = \kappa \triangleright \mathbf{load} \end{array}}{\{\Lambda, \Gamma, \mu, \kappa\} a[v] \longrightarrow \{\Lambda, \Gamma, \mu, \kappa'\} v'} \quad \text{EXR\_ARR\_GET\_VAL} \\
\\
\frac{\{\Lambda, \Gamma, \mu, \kappa\} e \longrightarrow \{\Lambda, \Gamma, \mu, \kappa'\} e'}{\{\Lambda, \Gamma, \mu, \kappa\} \ominus e \longrightarrow \{\Lambda, \Gamma, \mu, \kappa'\} \ominus e'} \quad \text{EXR\_UNOP\_EXPR} \\
\\
\frac{\begin{array}{c} v' \equiv \llbracket \ominus v \rrbracket \\ \kappa' = \kappa \triangleright \ominus \end{array}}{\{\Lambda, \Gamma, \mu, \kappa\} \ominus v \longrightarrow \{\Lambda, \Gamma, \mu, \kappa'\} v'} \quad \text{EXR\_UNOP\_VAL} \\
\\
\frac{\{\Lambda, \Gamma, \mu, \kappa\} e_1 \longrightarrow \{\Lambda, \Gamma, \mu, \kappa'\} e'_1}{\{\Lambda, \Gamma, \mu, \kappa\} e_1 \oplus e_2 \longrightarrow \{\Lambda, \Gamma, \mu, \kappa'\} e'_1 \oplus e_2} \quad \text{EXR\_BINOP\_L} \\
\\
\frac{\{\Lambda, \Gamma, \mu, \kappa\} e_2 \longrightarrow \{\Lambda, \Gamma, \mu, \kappa'\} e'_2}{\{\Lambda, \Gamma, \mu, \kappa\} v \oplus e_2 \longrightarrow \{\Lambda, \Gamma, \mu, \kappa'\} v \oplus e'_2} \quad \text{EXR\_BINOP\_R} \\
\\
\frac{\begin{array}{c} v_3 \equiv \llbracket v_1 \oplus v_2 \rrbracket \\ \kappa' = \kappa \triangleright \oplus \end{array}}{\{\Lambda, \Gamma, \mu, \kappa\} v_1 \oplus v_2 \longrightarrow \{\Lambda, \Gamma, \mu, \kappa'\} v_3} \quad \text{EXR\_BINOP\_VAL} \\
\\
\frac{}{\{\Lambda, \Gamma, \mu, \kappa\} \{\} e \longrightarrow \{\Lambda, \Gamma, \mu, \kappa\} e} \quad \text{EXR\_SUBST\_EMPTY} \\
\\
\frac{\{\Lambda, \Gamma, \mu, \kappa\} e \longrightarrow \{\Lambda, \Gamma, \mu, \kappa'\} e'}{\{\Lambda, \Gamma, \mu, \kappa\} \{\sigma\} e \longrightarrow \{\Lambda, \Gamma, \mu, \kappa'\} \{\sigma\} e'} \quad \text{EXR\_SUBST\_EXPR} \\
\\
\frac{\kappa' = \kappa \triangleright \mathbf{load}}{\{\Lambda, \Gamma, \mu, \kappa\} \{x_1/v_1, \dots, x_k/v_k\} x_i \longrightarrow \{\Lambda, \Gamma, \mu, \kappa'\} v_i} \quad \text{EXR\_SUBST\_VAR} \\
\\
\frac{}{\{\Lambda, \Gamma, \mu, \kappa\} \{\sigma\} x \longrightarrow \{\Lambda, \Gamma, \mu, \kappa\} x} \quad \text{EXR\_SUBST\_NO\_VAR} \\
\\
\frac{}{\{\Lambda, \Gamma, \mu, \kappa\} \{\sigma\} v \longrightarrow \{\Lambda, \Gamma, \mu, \kappa\} v} \quad \text{EXR\_SUBST\_VAL} \\
\\
\frac{\{\Lambda, \Gamma, \mu, \kappa\} e_1 \longrightarrow \{\Lambda, \Gamma, \mu, \kappa'\} e'_1}{\{\Lambda, \Gamma, \mu, \kappa\} f(v_1, \dots, v_k, e_1, e_2, \dots, e_n) \longrightarrow \{\Lambda, \Gamma, \mu, \kappa'\} f(v_1, \dots, v_k, e'_1, e_2, \dots, e_n)} \quad \text{EXR\_FN\_EXPR} \\
\\
\frac{\kappa' = \kappa \triangleright \mathbf{store}}{\{\Lambda, \Gamma, \mu, \kappa\} f(x'_1/v'_1, \dots, x'_k/v'_k, v_1, v_2, \dots, v_n) \longrightarrow \{\Lambda, \Gamma, \mu, \kappa'\} f(x'_1/v'_1, \dots, x'_k/v'_k, x_1/v_1, v_2, \dots, v_n)} \quad \text{EXR\_FN\_SUBST} \\
\\
\frac{\begin{array}{c} \Lambda = \Lambda'[f \mapsto (x_1, \dots, x_k) : s @e] \\ \mu' = \mu \triangleright \emptyset_\mu \\ \kappa' = \kappa \triangleright f \end{array}}{\{\Lambda, \Gamma, \mu, \kappa\} f(x_1/v_1, \dots, x_k/v_k) \longrightarrow \{\Lambda, \Gamma, \mu', \kappa'\} \{x_1/v_1, \dots, x_k/v_k\} s @e} \quad \text{EXR\_FN\_CALL} \\
\\
\frac{\{\Lambda, \Gamma, \mu, \kappa\} \{\sigma\} e \longrightarrow \{\Lambda, \Gamma, \mu, \kappa'\} e'}{\{\Lambda, \Gamma, \mu, \kappa\} \{\sigma\} \mathbf{skip} @e \longrightarrow \{\Lambda, \Gamma, \mu, \kappa\} \{\sigma\} \mathbf{skip} @e'} \quad \text{EXR\_SKIP\_EXPR} \\
\\
\frac{\begin{array}{c} \mu = \mu_1 \triangleright \mu_2 \\ \kappa' = \kappa \triangleright \mathbf{ret} \end{array}}{\{\Lambda, \Gamma, \mu, \kappa\} \{\sigma\} \mathbf{skip} @v \longrightarrow \{\Lambda, \Gamma, \mu_1, \kappa'\} v} \quad \text{EXR\_SKIP\_VAL} \\
\\
\frac{\{\Lambda, \Gamma, \mu, \kappa\} \{\sigma\} s_1 @e_0 \longrightarrow \{\Lambda, \Gamma, \mu', \kappa'\} \{\sigma'\} s'_1 @e_0}{\{\Lambda, \Gamma, \mu, \kappa\} \{\sigma\} s_1; s_2 @e_0 \longrightarrow \{\Lambda, \Gamma, \mu', \kappa'\} \{\sigma'\} s'_1; s_2 @e_0} \quad \text{EXR\_SEQ}
\end{array}$$

$$\begin{array}{c}
\frac{}{\{\Lambda, \Gamma, \mu, \kappa\} \{\sigma\} \mathbf{skip}; s @e_0 \longrightarrow \{\Lambda, \Gamma, \mu, \kappa\} \{\sigma\} s @e_0} \text{EXR\_SEQ\_SKIP} \\
\\
\frac{\{\Lambda, \Gamma, \mu, \kappa\} \{\sigma\} e \longrightarrow \{\Lambda, \Gamma, \mu, \kappa'\} e'}{\{\Lambda, \Gamma, \mu, \kappa\} \{\sigma\} \mathbf{def} x := e @e_0 \longrightarrow \{\Lambda, \Gamma, \mu, \kappa'\} \{\sigma\} \mathbf{def} x := e' @e_0} \text{EXR\_DEF\_EXPR} \\
\\
\frac{\mu' = \mu[x \mapsto v] \quad \kappa' = \kappa \triangleright \mathbf{store}}{\{\Lambda, \Gamma, \mu, \kappa\} \{\sigma\} \mathbf{def} x := v @e_0 \longrightarrow \{\Lambda, \Gamma, \mu', \kappa'\} \{\sigma\} \mathbf{skip} @e_0} \text{EXR\_DEF\_VAL} \\
\\
\frac{\Gamma' = \Gamma[a \mapsto []] \quad \mu' = \mu[x \mapsto a] \quad \kappa' = \kappa \triangleright \mathbf{store}}{\{\Lambda, \Gamma, \mu, \kappa\} \{\sigma\} \mathbf{adef} x := a @e_0 \longrightarrow \{\Lambda, \Gamma', \mu', \kappa'\} \{\sigma\} \mathbf{skip} @e_0} \text{EXR\_DEF\_ARR} \\
\\
\frac{\{\Lambda, \Gamma, \mu, \kappa\} \{\sigma\} e \longrightarrow \{\Lambda, \Gamma, \mu, \kappa'\} e'}{\{\Lambda, \Gamma, \mu, \kappa\} \{\sigma\} x := e @e_0 \longrightarrow \{\Lambda, \Gamma, \mu, \kappa'\} \{\sigma\} x := e' @e_0} \text{EXR\_ASSIGN\_EXPR} \\
\\
\frac{\mu' = \mu[x \mapsto v] \quad \kappa' = \kappa \triangleright \mathbf{store}}{\{\Lambda, \Gamma, \mu, \kappa\} \{\sigma\} x := v @e_0 \longrightarrow \{\Lambda, \Gamma, \mu', \kappa'\} \{\sigma\} \mathbf{skip} @e_0} \text{EXR\_ASSIGN\_VAL} \\
\\
\frac{\{\Lambda, \Gamma, \mu, \kappa\} \{\sigma\} e_1 \longrightarrow \{\Lambda, \Gamma, \mu, \kappa'\} e'_1}{\{\Lambda, \Gamma, \mu, \kappa\} \{\sigma\} a[e_1] := e_2 @e_0 \longrightarrow \{\Lambda, \Gamma, \mu, \kappa'\} \{\sigma\} a[e'_1] := e_2 @e_0} \text{EXR\_ARR\_ASSIGN\_EXPR\_L} \\
\\
\frac{\{\Lambda, \Gamma, \mu, \kappa\} \{\sigma\} e_2 \longrightarrow \{\Lambda, \Gamma, \mu, \kappa'\} e'_2}{\{\Lambda, \Gamma, \mu, \kappa\} \{\sigma\} a[v_1] := e_2 @e_0 \longrightarrow \{\Lambda, \Gamma, \mu, \kappa'\} \{\sigma\} a[v_1] := e'_2 @e_0} \text{EXR\_ARR\_ASSIGN\_EXPR\_R} \\
\\
\frac{\Gamma' = \Gamma[a \mapsto \Gamma(a)[v_1 \mapsto v_2]] \quad \kappa' = \kappa \triangleright \mathbf{store}}{\{\Lambda, \Gamma, \mu, \kappa\} \{\sigma\} a[v_1] := v_2 @e_0 \longrightarrow \{\Lambda, \Gamma', \mu, \kappa'\} \{\sigma\} \mathbf{skip} @e_0} \text{EXR\_ARR\_ASSIGN\_VAL} \\
\\
\frac{v_1 < v_2 \quad v'_1 = v_1 + 1 \quad \kappa' = \kappa \triangleright \mathbf{store}}{\{\Lambda, \Gamma, \mu, \kappa\} \{\sigma\} \mathbf{for} x \mathbf{from} v_1 \mathbf{to} v_2 : s @e_0 \longrightarrow \{\Lambda, \Gamma, \mu, \kappa'\} \{\sigma\} (\{x/v_1\} s); \mathbf{for} x \mathbf{from} v'_1 \mathbf{to} v_2 : s @e_0} \text{EXR\_FOR} \\
\\
\frac{\{\sigma_1\} \cap \{\sigma_2\} = \{\} \quad \{\sigma_3\} = \{\sigma_1\} \cup \{\sigma_2\}}{\{\Lambda, \Gamma, \mu, \kappa\} \{\sigma_1\} (\{\sigma_2\} s) @e_0 \longrightarrow \{\Lambda, \Gamma, \mu, \kappa'\} \{\sigma_3\} s @e_0} \text{EXR\_ADD\_SUBST} \\
\\
\frac{v_1 = v_2}{\{\Lambda, \Gamma, \mu, \kappa\} \{\sigma\} \mathbf{for} x \mathbf{from} v_2 \mathbf{to} v_2 : s @e_0 \longrightarrow \{\Lambda, \Gamma, \mu', \kappa\} \{\sigma\} \mathbf{skip} @e_0} \text{EXR\_FOR\_BASE} \\
\\
\boxed{\llbracket e_h \rrbracket_t = e} \quad e_h \text{ is transformed to } e \\
\\
\frac{v \equiv \llbracket v_h \rrbracket_{int}}{\llbracket v_h \rrbracket_t = v} \text{EXT\_VAL} \\
\\
\frac{}{\llbracket x \rrbracket_t = x} \text{EXT\_VAR} \\
\\
\frac{}{\llbracket a \rrbracket_t = a} \text{EXT\_ARR} \\
\\
\frac{\llbracket e_h \rrbracket_t = e}{\llbracket a[e_h] \rrbracket_t = a[e]} \text{EXT\_ARR\_GET}
\end{array}$$

$$\frac{\begin{array}{c} \llbracket \mathbf{fdef}_h f_h hfval \rrbracket_t = \mathbf{fdef} f fval \\ \llbracket e_{h1} \rrbracket_t = e_1 \quad \dots \quad \llbracket e_{hk} \rrbracket_t = e_k \end{array}}{\llbracket f_h(e_{h1}, \dots, e_{hk}) \rrbracket_t = f(e_1, \dots, e_k)} \quad \text{EXT\_FN\_CALL}$$

$$\boxed{\llbracket s_h \rrbracket_{ctx} = s} \quad s_h \text{ is transformed to } s$$

$$\frac{}{\llbracket \mathbf{skip}_h \rrbracket_{ctx} = \mathbf{skip}} \quad \text{STT\_SKIP}$$

$$\frac{\begin{array}{c} \llbracket s_{h1} \rrbracket_{ctx} = s_1 \\ \llbracket s_{h2} \rrbracket_{ctx} = s_2 \end{array}}{\llbracket s_{h1}; s_{h2} \rrbracket_{ctx} = s_1; s_2} \quad \text{STT\_SEQ}$$

$$\frac{\llbracket e_h \rrbracket_t = e}{\llbracket \mathbf{def}_h x := e_h \rrbracket_{ctx} = \mathbf{def} x := e} \quad \text{STT\_VAR\_DEC}$$

$$\frac{}{\llbracket \mathbf{adef}_h x := a \rrbracket_{ctx} = \mathbf{adef} x := a} \quad \text{STT\_ARR\_DEC}$$

$$\frac{\begin{array}{c} \llbracket e_h \rrbracket_t = e \\ e' = ctx \& rnset \\ e'' = e \& e' \\ e''' = x \& (\sim e') \end{array}}{\llbracket x := e_h \rrbracket_{ctx} = x := (e'' \mid e''')} \quad \text{STT\_VAR\_ASSIGN}$$

$$\frac{\begin{array}{c} \llbracket e_{h1} \rrbracket_t = e_1 \\ \llbracket e_{h2} \rrbracket_t = e_2 \\ e' = ctx \& rnset \\ e'' = e_2 \& e' \\ e''' = a[e_1] \& (\sim e') \end{array}}{\llbracket a[e_{h1}] := e_{h2} \rrbracket_{ctx} = a[e_1] := (e'' \mid e''')} \quad \text{STT\_ARR\_ASSIGN}$$

$$\frac{\begin{array}{c} \llbracket v_{h1} \rrbracket_t = v_1 \\ \llbracket v_{h2} \rrbracket_t = v_2 \\ \llbracket s_h \rrbracket_{ctx} = s \end{array}}{\llbracket \mathbf{for}_h x \mathbf{from} v_{h1} \mathbf{to} v_{h2} \rrbracket_{ctx} = \mathbf{for} x \mathbf{from} v_1 \mathbf{to} v_2 : s} \quad \text{STT\_FOR}$$

$$\frac{\begin{array}{c} \llbracket e_h \rrbracket_t = e \\ \llbracket s_{h1} \rrbracket_{(ctx' \& ctx)} = s_1 \\ \llbracket s_{h2} \rrbracket_{(ctx' \& ctx)} = s_2 \end{array}}{\llbracket \mathbf{if}_h e_h \mathbf{then} s_{h1} \mathbf{else} s_{h2} \rrbracket_{ctx} = \mathbf{def} ctx' := e; s_1; ctx' := (\sim ctx'); s_2} \quad \text{STT\_IF}$$

$$\frac{\begin{array}{c} \llbracket e_h \rrbracket_t = e \\ e' = ctx \& rnset \\ e'' = e \& e' \end{array}}{\llbracket \mathbf{return}_h e_h \rrbracket_{ctx} = rval := (e'' \mid rval); rnset := (rnset \& (\sim ctx))} \quad \text{STT\_RET}$$

$$\boxed{\llbracket hfndef \rrbracket_t = fndef} \quad hfndef \text{ is transformed to } fndef$$

$$\frac{\llbracket s_h \rrbracket_{\text{TRUE}} = s}{\llbracket \mathbf{fdef}_h f_h (x_1, \dots, x_k) : s_h \rrbracket_t = \mathbf{fdef} f (x_1, \dots, x_k) : \mathbf{def} rval := \text{FALSE}; \mathbf{def} rnset := \text{TRUE}; s @rval} \quad \text{FDEFT\_FDEF}$$

$$\boxed{\llbracket \ominus_h \rrbracket_t = \ominus} \quad \ominus_h \text{ is transformed to } \ominus$$

$$\frac{}{\llbracket ! \rrbracket_t = \sim} \quad \text{UNOPT\_LNOT}$$

$$\frac{}{\llbracket \ominus_h \rrbracket_t = \ominus} \quad \text{UNOPT\_UNOP}$$

$\boxed{\llbracket \oplus_h \rrbracket_t = \oplus}$      $\oplus_h$  is transformed to  $\oplus$

$\overline{\llbracket \&\& \rrbracket_t = \&}$	BINOPT_LAND
$\overline{\llbracket \mid \mid \rrbracket_t = \mid}$	BINOPT_LOR
$\overline{\llbracket == \rrbracket_t = ==_s}$	BINOPT_EQ
$\overline{\llbracket != \rrbracket_t = !=_s}$	BINOPT_NEQ
$\overline{\llbracket > \rrbracket_t = >_s}$	BINOPT_GT
$\overline{\llbracket < \rrbracket_t = <_s}$	BINOPT_LT
$\overline{\llbracket >= \rrbracket_t = >=_s}$	BINOPT_GTE
$\overline{\llbracket <= \rrbracket_t = <=_s}$	BINOPT_LTE
$\overline{\llbracket \oplus_h \rrbracket_t = \oplus}$	BINOPT_BINOP

Definition rules:            57 good      0 bad  
Definition rule clauses: 125 good      0 bad