Grammar

Base type

b ::=
- | Bool
- | UInt_n
- | Int_n

Type

τ ::=
- | b
- | Arr$\langle b, n\rangle$

Label

ℓ ::=
- | Public
- | Secret

Mutability

σ ::=
- | Const
- | Mut

Function argument

arg ::=
- | e — by value
- | Mut x — by reference

Expression

e ::=
- | true
- | false
- | c — integer literal
- | x — variable
- | a — array
- | $a[e]$ — array get
- | $\ominus e$ — unary op
- | $e_1 \oplus e_2$ — binary op
- | $f(arg_1, \ldots, arg_n)$ — function call

Array

a ::=
- | $\langle \tau, n\rangle x \Rightarrow e$ — array comprehension
- | view(a, e, n) — array view

Statement

s ::=
- | $s_1; s_2$ — sequence
- | $\langle \tau, \ell, \sigma\rangle x := e$ — variable declaration
- | $x := e$ — variable assignment
- | $a[e_1] = e_2$ — array assignment
- | if e $\{s_1\}$ else $\{s_2\}$ — conditional
- | for $\langle b\rangle x$ from e_1 to e_2 $\{s\}$ — loop
- | return e — return

Function Definition

$fdec$ ::=
- | $\langle b, \ell\rangle f(\langle \tau_1, \ell_1, \sigma_1\rangle x_1, \ldots, \langle \tau_n, \ell_n, \sigma_n\rangle x_n)$ $\{s\}$

Type Lattice

$$\frac{n_1 < n_2}{\text{UInt\_}n_1 <_\tau \text{UInt\_}n_2} \qquad \frac{n_1 < n_2}{\text{Int\_}n_1 <_\tau \text{Int\_}n_2} \qquad \frac{}{\text{UInt\_}n <_\tau \text{Int\_}2n} \qquad \frac{}{\text{Public} <_\ell \text{Secret}}$$

$$\frac{}{\text{Const} <_\sigma \text{Mut}} \qquad \frac{\tau_1 <_\tau \tau_2 \qquad \Gamma \vdash e : \langle \tau_1, \ell\rangle}{\Gamma \vdash e : \langle \tau_2, \ell\rangle} \qquad \frac{}{\ell \cup \ell = \ell} \qquad \frac{}{\ell \cup \text{Secret} = \text{Secret}}$$

Parameter Passing

$$\frac{\Gamma \vdash e : \langle \tau, \ell_1\rangle \qquad \ell_1 \leq_\ell \ell_2}{\langle \tau, \ell_2, \text{Const}\rangle \leftarrow e} \qquad \frac{\mu(x) = \langle \tau, \ell, \text{Mut}\rangle}{\langle \tau, \ell, \text{Mut}\rangle \leftarrow \text{Mut } x}$$

Expressions

$$\textsc{Var}\ \frac{\mu(x) = \langle \tau, \ell, \sigma \rangle}{\Gamma \vdash x : \langle \tau, \ell \rangle} \qquad \textsc{Unop}\ \frac{\Gamma \vdash e : \langle \tau_1, \ell_1 \rangle \qquad \ominus : \langle \tau_1, \ell_1 \rangle \to \langle \tau_2, \ell_2 \rangle}{\Gamma \vdash \ominus e : \langle \tau_2, \ell_2 \rangle}$$

$$\textsc{Binop}\ \frac{\Gamma \vdash e_1 : \langle \tau_1, \ell_1 \rangle \qquad \Gamma \vdash e_2 : \langle \tau_2, \ell_2 \rangle \qquad \oplus : \langle \tau_1, \ell_1 \rangle \to \langle \tau_2, \ell_2 \rangle \to \langle \tau_3, \ell_3 \rangle}{\Gamma \vdash e_1 \oplus e_2 : \langle \tau_3, \ell_3 \rangle}$$

$$\textsc{ArrGet}\ \frac{\mu(a) = \langle \textsc{Arr}\langle b, n \rangle, \ell, \sigma \rangle \qquad \Gamma \vdash e : \langle \textsc{UInt\_}max, \textsc{Public} \rangle}{\Gamma \vdash a[e] : \langle b, \ell \rangle}$$

$$\textsc{FnCall}\ \frac{\begin{array}{c}\mathbb{F}(f) = fdec(\langle \tau_1, \ell_1, \sigma_1 \rangle x_1, \ldots, \langle \tau_n, \ell_n, \sigma_n \rangle x_n) : \langle b, \ell \rangle \\ \langle \tau_1, \ell_1, \sigma_1 \rangle x_1 \leftarrow arg_1 \qquad \cdots \qquad \langle \tau_n, \ell_n, \sigma_n \rangle x_n \leftarrow arg_n\end{array}}{\Gamma \vdash f(arg_1, \ldots, arg_n) : \langle b, \ell \rangle} \qquad \textsc{True}\ \frac{}{\Gamma \vdash \textsc{true} : \langle \textsc{Bool}, \textsc{Public} \rangle}$$

$$\textsc{False}\ \frac{}{\Gamma \vdash \textsc{false} : \langle \textsc{Bool}, \textsc{Public} \rangle} \qquad \textsc{PosNumber}\ \frac{c >= 0 \qquad n = \lceil \log_2 c \rceil}{\Gamma \vdash c : \langle \textsc{UInt\_}n, \textsc{Public} \rangle} \qquad \textsc{NegNumber}\ \frac{c < 0 \qquad n = \lceil \log_2 |c| \rceil + 1}{\Gamma \vdash c : \langle \textsc{Int\_}n, \textsc{Public} \rangle}$$

$$\textsc{ArrComp}\ \frac{\Gamma \vdash e : \langle b, \ell \rangle}{\Gamma \vdash \langle b, n \rangle x \Rightarrow e : \langle \textsc{Arr}\langle b, n \rangle, \ell \rangle}$$

$$\textsc{ArrView}\ \frac{\mu(a) = \langle \textsc{Arr}\langle b, n \rangle, \ell, \sigma \rangle \qquad \Gamma \vdash e : \langle \textsc{UInt\_}max, \textsc{Public} \rangle \qquad n' \leq n}{\Gamma \vdash \textsc{view}(a, e, n') : \langle \tau, \ell \rangle}$$

Statements

$$\textsc{Seq}\ \frac{\Sigma\langle \ell_s \rangle \vdash s_1 : \ell'_s \qquad \Sigma\langle \ell'_s \rangle \vdash s_2 : \ell''_s}{\Sigma\langle \ell_s \rangle \vdash s_1; s_2 : \ell'_s \cup \ell''_s}$$

$$\textsc{VarDec}\ \frac{\Gamma \vdash e : \langle \tau, \ell_1 \rangle \qquad \ell_1 \leq_\ell \ell_2}{\Sigma\langle \ell_s \rangle \vdash \langle \tau, \ell_2, \sigma \rangle x := e : \textsc{Public}} \quad \mu(x) = \langle \tau, \ell_2, \sigma \rangle \text{ (scoping?)}$$

(how to ensure no Mut view of Const array?)

$$\textsc{VarAssign}\ \frac{\mu(x) = \langle b, \ell_1, \textsc{Mut} \rangle \qquad \Gamma \vdash e : \langle b, \ell_2 \rangle \qquad \ell_2 \leq_\ell \ell_1}{\Sigma\langle \ell_s \rangle \vdash x := e : \textsc{Public}}$$

$$\textsc{ArrAssign}\ \frac{\mu(a) = \langle \textsc{Arr}\langle b, n \rangle, \ell_1, \textsc{Mut} \rangle \qquad \Gamma \vdash e_1 : \langle \textsc{UInt\_}max, \textsc{Public} \rangle \qquad \Gamma \vdash e_2 : \langle b, \ell_2 \rangle \qquad \ell_2 \leq_\ell \ell_1}{\Sigma\langle \ell_s \rangle \vdash a[e_1] := e_2 : \textsc{Public}}$$

$$\textsc{If}\ \frac{\Gamma \vdash e : \langle \textsc{Bool}, \ell \rangle \qquad \Sigma\langle \ell \cup \ell_s \rangle \vdash s_1 : \ell'_s \qquad \Sigma\langle \ell \cup \ell_s \rangle \vdash s_2 : \ell''_s}{\Sigma\langle \ell_s \rangle \vdash \textsc{if}\ e\ \{s_1\}\ \textsc{else}\ \{s_2\} : \ell'_s \cup \ell''_s}$$

$$\textsc{For}\ \frac{\Gamma \vdash e_1 : \langle b, \textsc{Public} \rangle \qquad \Gamma \vdash e_2 : \langle b, \textsc{Public} \rangle \qquad b = \textsc{UInt\_}s \vee b = \textsc{Int\_}s \qquad \Sigma\langle \ell_s \rangle \vdash s : \ell'_s}{\Sigma\langle \ell_s \rangle \vdash \textsc{for}\ \langle b \rangle x\ \textsc{from}\ e_1\ \textsc{to}\ e_2\ \{s\} : \ell'_s} \quad \mu(x) = \langle b, \textsc{Public}, \textsc{Const} \rangle \text{ (scoping?)}$$

$$\textsc{Ret}\ \frac{\Gamma \vdash e : \langle b, \ell_1 \rangle \qquad \mathbb{F}(f) = fdec : \langle b, \ell_2 \rangle \qquad \ell_1 \leq_\ell \ell_2}{\Sigma\langle \ell_s \rangle \vdash \textsc{return}\ e : \ell_s}$$