# Scope



### Recall: substitution model

- Way of giving semantics to the λ-calculus
  - ► E.g.,  $(\lambda x.f \times x)$   $(\lambda y.z)$  →<sub>β</sub>  $f(\lambda y.z)$   $(\lambda y.z)$
- Translate this knowledge to JavaScript functions
  - $> (x => f(x)(x)) (y => z) \rightarrow_{\beta} f(y => z)(y => z)$

### Let's think about this more..

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- Why would you not actually want to do function application in this way for a language like JavaScript?
  - It's super slow! Why?
  - It's actually nonsensical sometimes! When?

# Substitution gone wrong

Consider variable mutation in JavaScript:

```
let y = 1;

let z = 0;

z++;

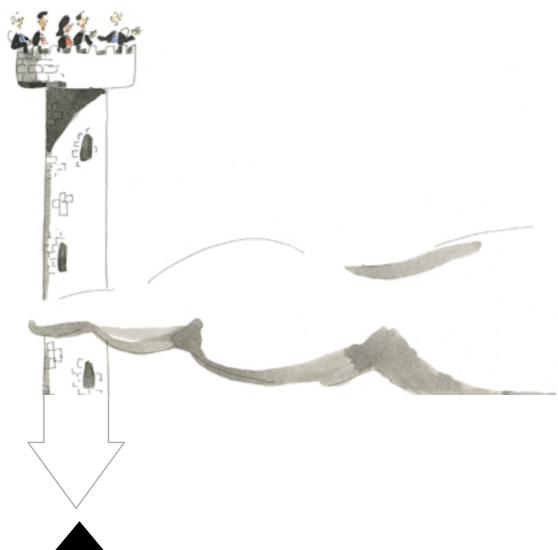
console.log(z);

...
```

- There is nothing wrong with substitution per say
  - It's symbolic evaluation/computation
  - Problem is JavaScript has mutation and not amendable to symbolic evaluation

### What can we do?

λ-calculus



environment model



machine model

### The environment model (by example)

- Anatomy of a scope
- First-order functions
- Free variables
- High-order functions (bonus)

What's the point of a scope (e.g., block scope)?

Recall our previous example:

```
let y = 1;
let z = 0;
z++;
console.log(z);
environment ptr
```

- In this model, we associate an environment (activation record) with the code we're executing
  - Environment contains entries of all variables in scope
  - Environment/stack ptr: points to cur activation record

Recall our previous example:

```
let y = 1;

let z = 0;

z++;

console.log(z);

y 1

z 0

environment ptr
```

- In this model, we associate an environment (activation record) with the code we're executing
  - Environment contains entries of all variables in scope
  - Environment/stack ptr: points to cur activation record

- In the environment model, we can distinguish between values and locations
  - <u>r-values</u>: plain old values; we can reason about them using substitution semantics
  - <u>I-values</u>: refer to locations where r-values are stored; they persist beyond single expressions.
- Why is this important?
  - ➤ It tells us the kind of values operators like ++ must take. Which does ++ take? A: r-values. B: l-values

- In the environment model, we can distinguish between values and locations
  - <u>r-values</u>: plain old values; we can reason about them using substitution semantics
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- Why is this important?
  - It tells us the kind of values operators like ++ must take. Which does ++ take? A: r-values. B: I-values

What's the process for executing z++:

```
let y = 1;

let z = 0;

z + + z = 0;

console.log(z);

y 1

z 0

z environment ptr
```

- Algorithm:
  - Find the current environment
  - Check to see if variable being reference is in env: if so, mutate!

What's the process for executing console.log(z)

```
let y = 1;

let z = 0;

z++;

console.log(z);

y 1

z 1

environment ptr
```

- Algorithm:
  - Find the current environment
  - Check to see if variable being reference is in env: if so, read it!

- This sounds slow!
  - It is!
  - But remember: this is not the machine model, this is still an abstract model!
- Not too far off from machine model
  - In x86, you dereference %esp to figure out where stack is and use offset to that location
  - In JavaScript, you often do table lookup to find location of variables

# The environment model (by example)

- Anatomy of a scope √
- First-order functions
- Free variables
- High-order functions (bonus)

### When do we create an environment?

- A: every time we enter a new block scope
- B: every time we enter a new function scope
- C: A and B
- D: we don't create new environments

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### First-order functions

Consider activation record when calling function:

```
function fact(n) {
    if (n <= 1) {
        return 1;
    } else {
        return n * fact(n-1);
    }
}
fact(3);</pre>
global env
fact(3);
environment ptr
```

What else do we need to keep track of?

### First-order functions

Consider activation record when calling function:

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function fact(n) {
    if (n <= 1) {
        return 1;
    } else {
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    }
}
fact(3);</pre>
global env
control
ret
n
3
environment ptr
```

What else do we need to keep track of?

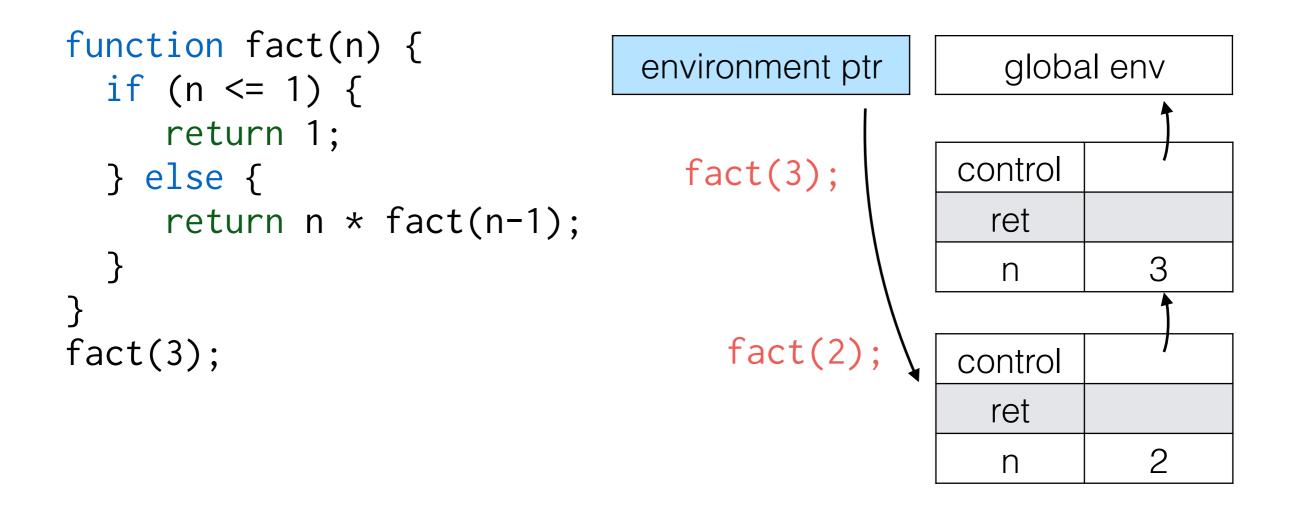
- The parts of an activation record when calling function
  - <u>control link:</u> records where to switch the environment pointer to when we finish evaluating in this scope.
    - Do we need this for block scopes too? A: yes, B:no
  - return value: I-value where the return value of function should be stored
  - parameters: I-value for each formal parameter
  - local variables: I-values for each let+const declaration

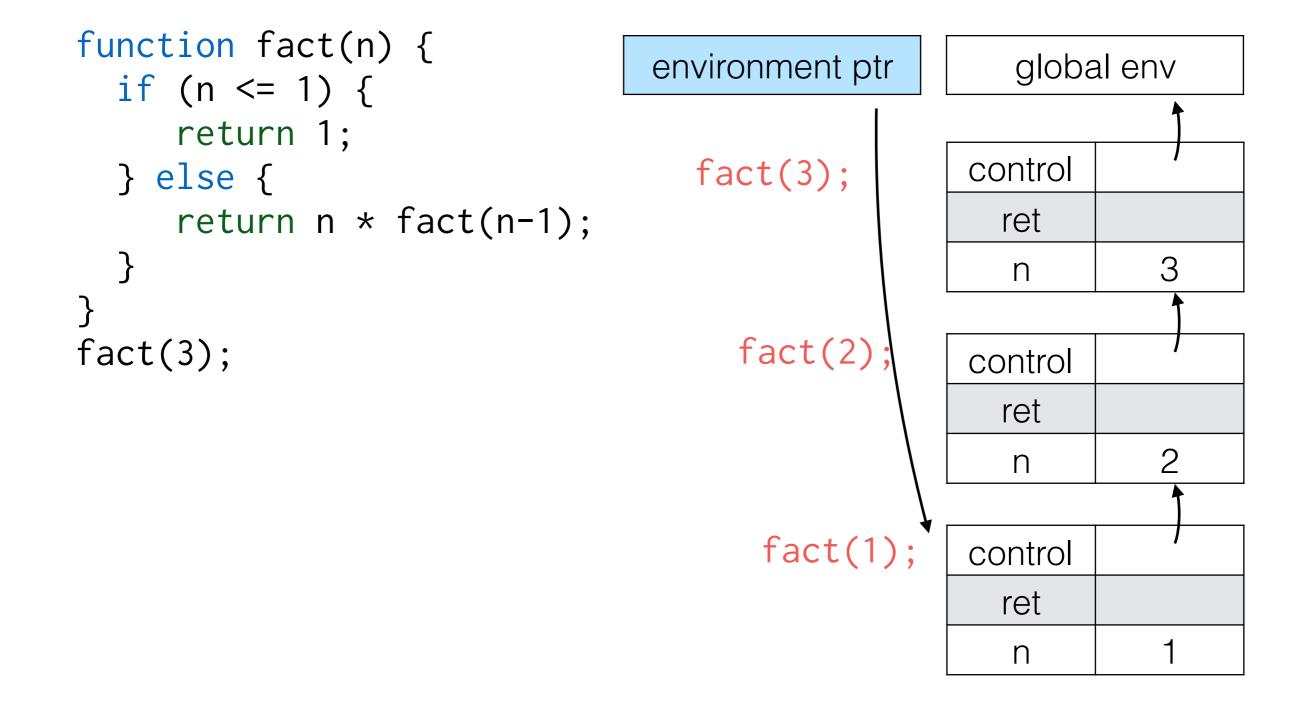
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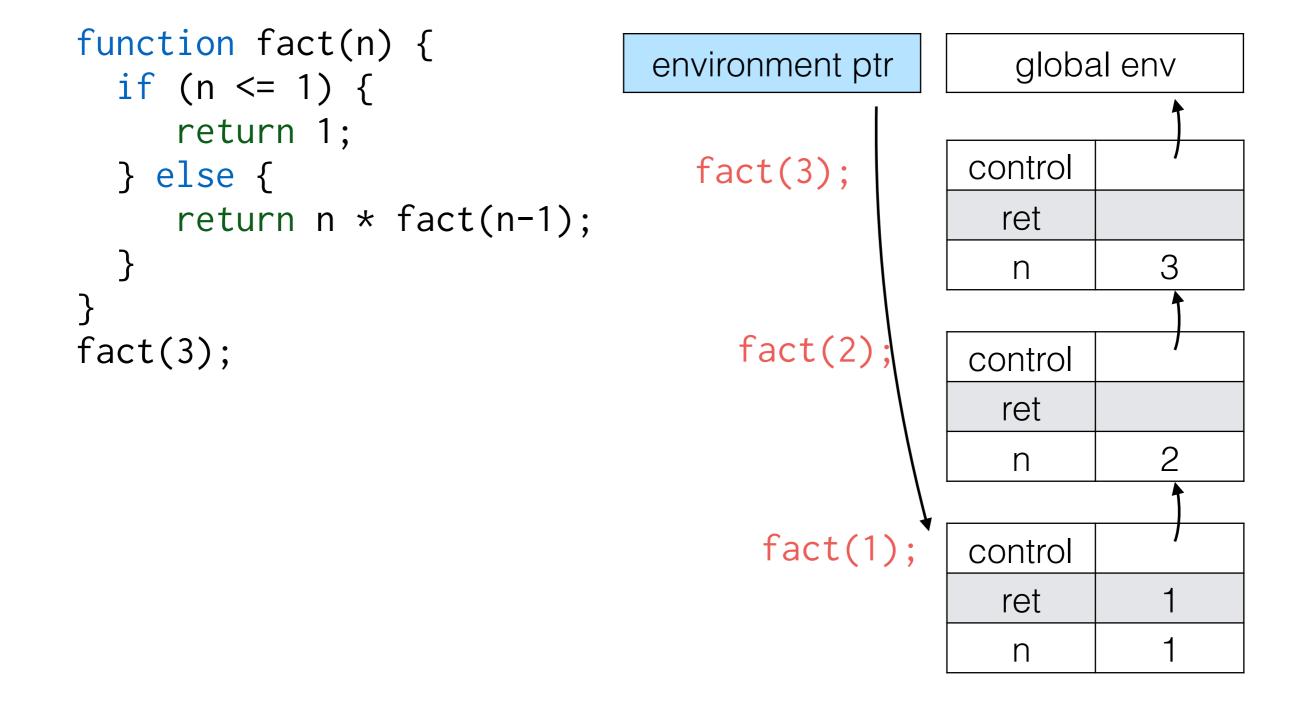
Do we need anything else besides the control link?

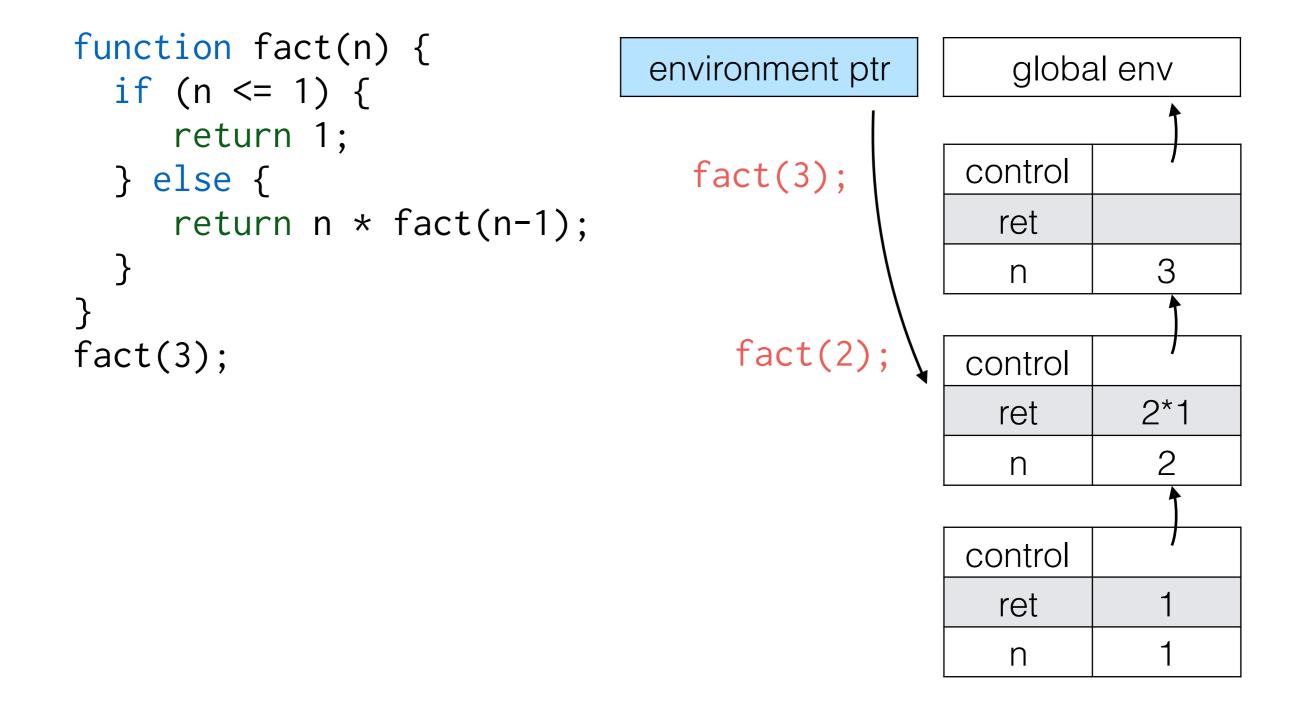
- Do we need anything else besides the control link?
  - Yes! Typically activation records will store the return address where to resume code execution — we'll talk about this in the control flow lecture

```
function fact(n) {
  if (n <= 1) {
    return 1;
  } else {
    return n * fact(n-1);
  }
  fact(3);</pre>
fact(3);
global env
control
  ret
  n 3
```







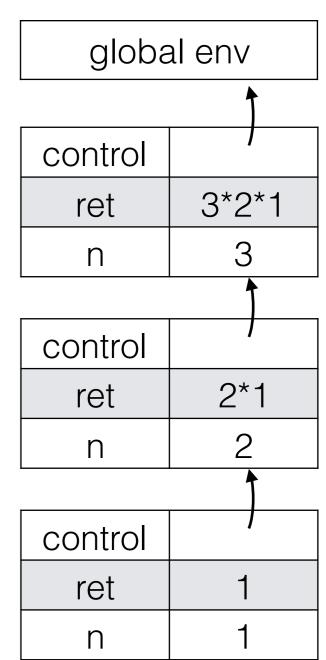


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```

 Do we keep the activation records on the stack after evaluation?

A: yes, B: no

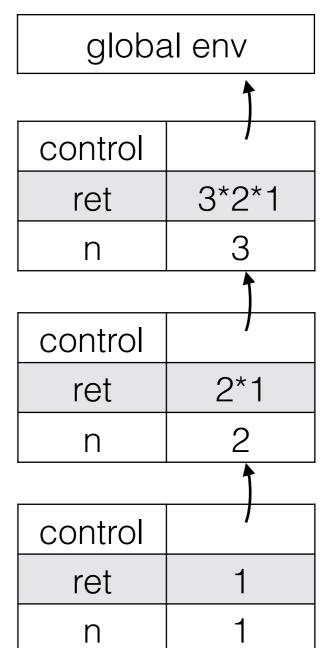


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# The environment model (by example)

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- A: yes
- ➤ B: no

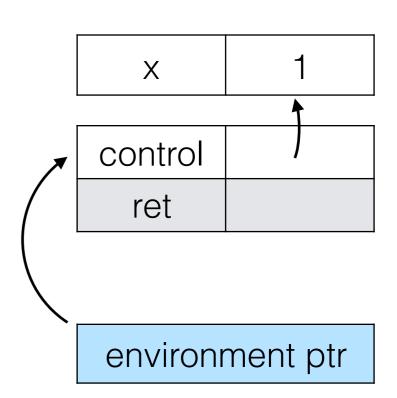
Should we lookup x via the control link?

A: yes

▶ B: no

### Free variables

```
let x = 1;
function f() {
  console.log(x)
}
f();
```

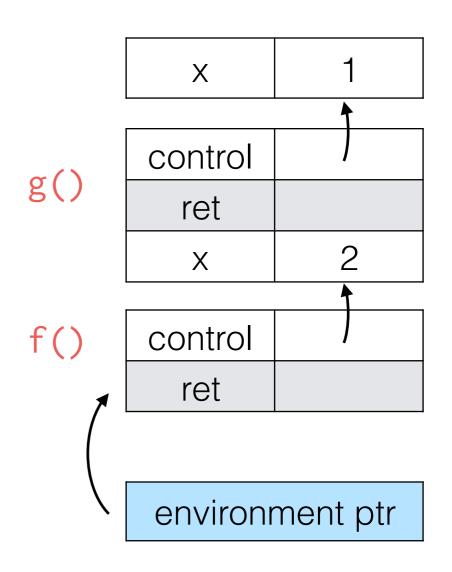


#### Free variables

Consider activation records when calling g:

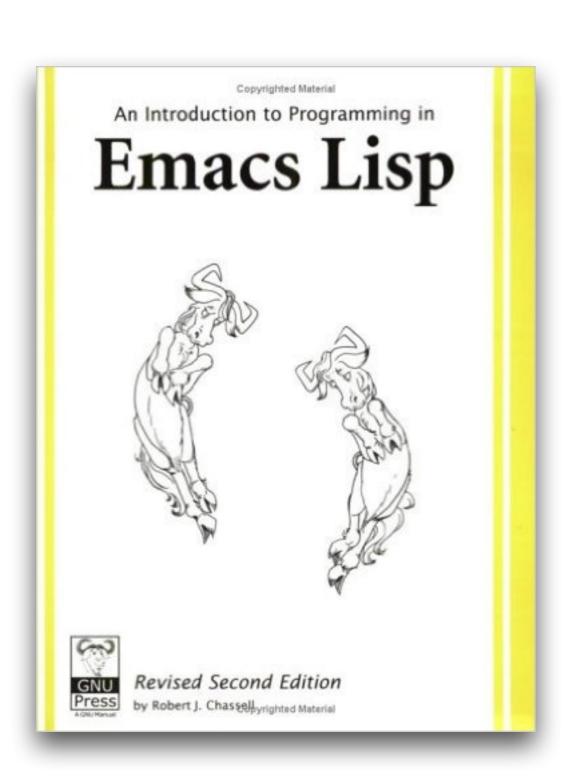
```
let x = 1;
function f() {
  console.log(x)
}

function g() {
  let x = 2;
  f();
}
```



What happens when we follow the control link?

# Congrats, you did it!



You invented dynamic scoping!

#### How do we "fix" this?

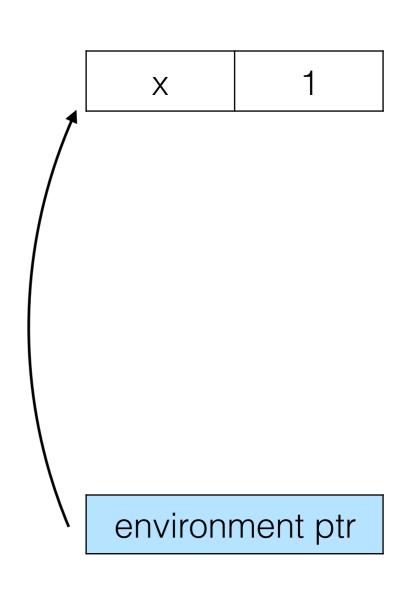
- We need more bookkeeping!
  - <u>access link:</u> reference to activation record of closest enclosing lexical scope
- Modify our lookup algorithm:
  - Find the current environment
  - Check to see if variable being reference is in env
  - If not, follow the access link and repeat

## Retry with access links

Consider activation records when calling g:

```
let x = 1;
function f() {
  console.log(x)
}

function g() {
  let x = 2;
  f();
}
```

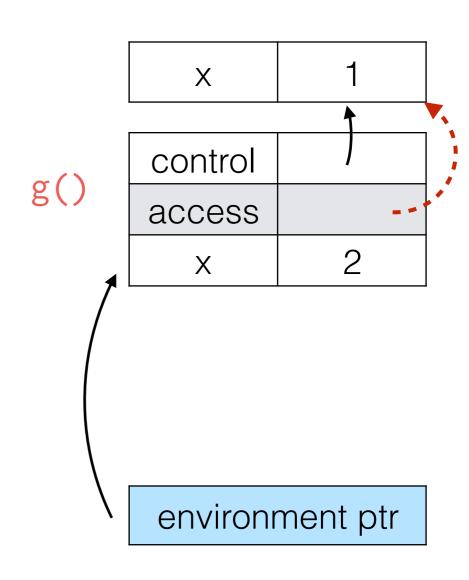


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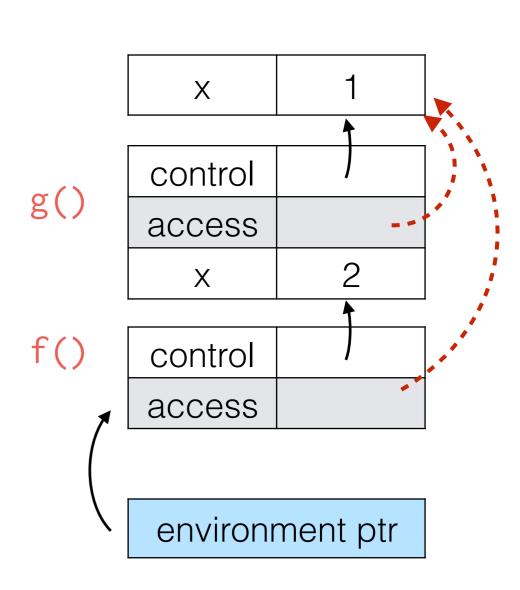


### Retry with access links

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  f();
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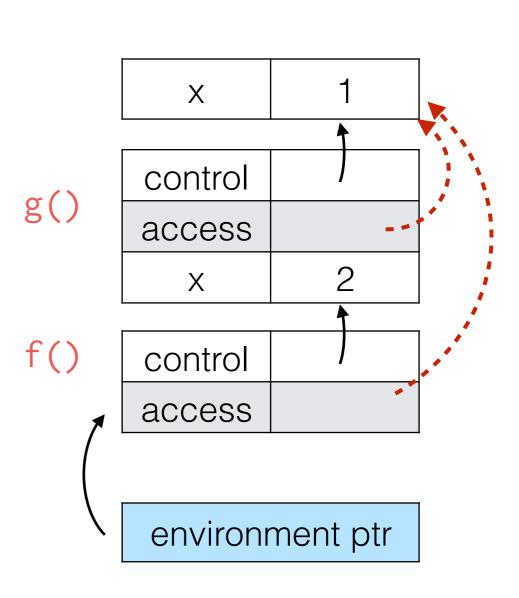


# Wait, there is some magic here

How do we know how to wire up the access links?

```
let x = 1;
function f() {
  console.log(x)
}

function g() {
  let x = 2;
  f();
}
```



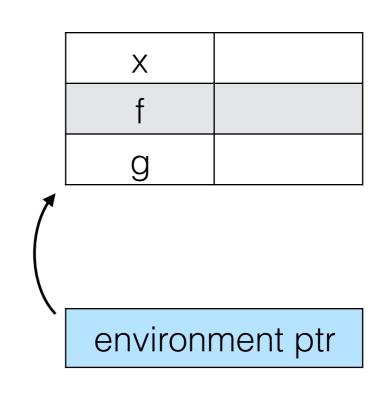
#### Functions are data!

The act of defining a function should include the act of recording the access link associated with the function

• Let's look at the example again, with minor rewrite

```
let x = 1;
let f = () => {
  console.log(x)
}

let g = () => {
  let x = 2;
  f();
}
```



• Let's look at the example again, with minor rewrite

```
let x = 1;
let f = () => {
  console.log(x)
}

let g = () => {
  let x = 2;
  f();
}
```

X	1
f	
g	
environment ptr	

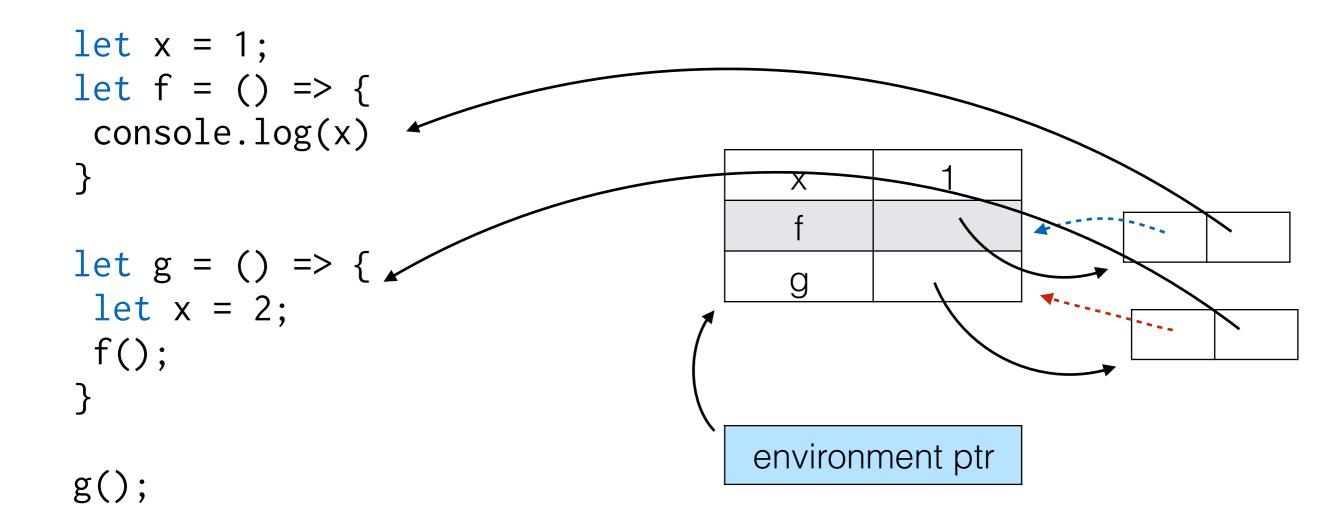
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```
let x = 1;
let f = () => {
  console.log(x)
}

let g = () => {
  let x = 2;
  f();
}

g();
environment ptr
```

Let's look at the example again, with minor rewrite



 When we evaluate function, the access link is set to the pointer in the closure

```
let x = 1;
let f = () => {
  console.log(x)
}

let g = () => {
  let x = 2;
  f();
}
```

environment ptr

```
x 1
f
```

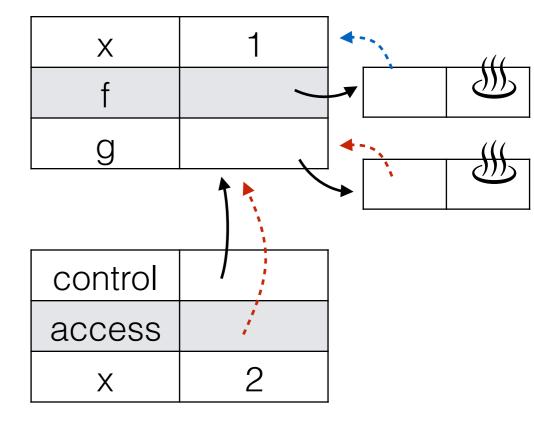
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```
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environment ptr

g()



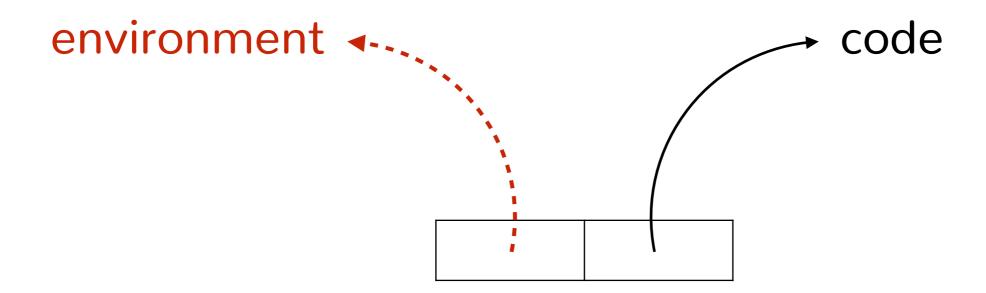
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```
environment ptr
                                               X
let x = 1;
let f = () => {
 console.log(x)
                                    g()
                                            control
let g = () => {
                                            access
 let x = 2;
                                               X
 f();
                                    f()
                                            control
g();
                                            access
```

 When we evaluate function, the access link is set to the pointer in the closure

```
environment ptr
                                              X
let x = 1;
let f = () => {
 console.log(x) // 1
                                    g()
                                            control
let g = () => {
                                            access
 let x = 2;
                                              X
 f();
                                    f()
                                            control
g();
                                            access
```

#### Closures

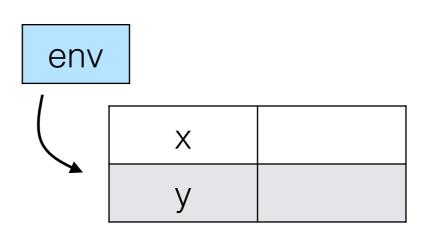


## The environment model (by example)

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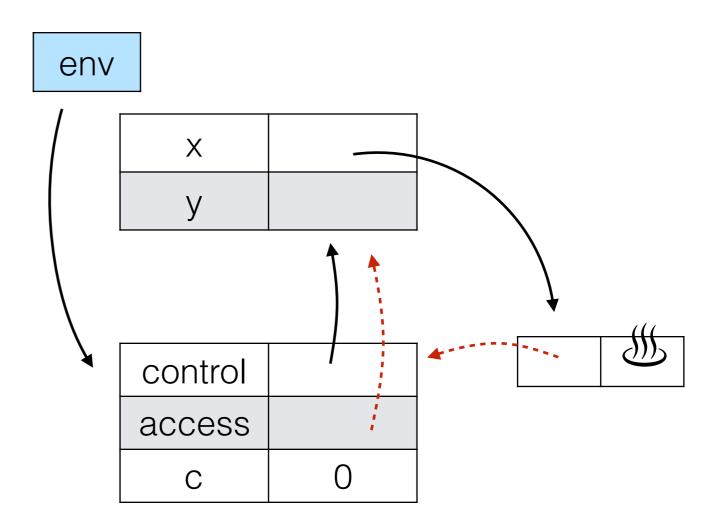
```
function mkCounter(c) {
   return () => {
      return c++;
   };
}

let x = mkCounter(0);
let y = mkCounter(2);
console.log(x());
```



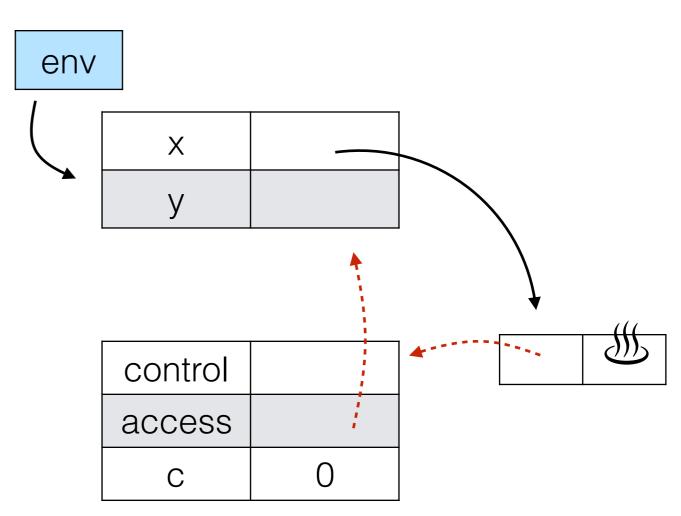
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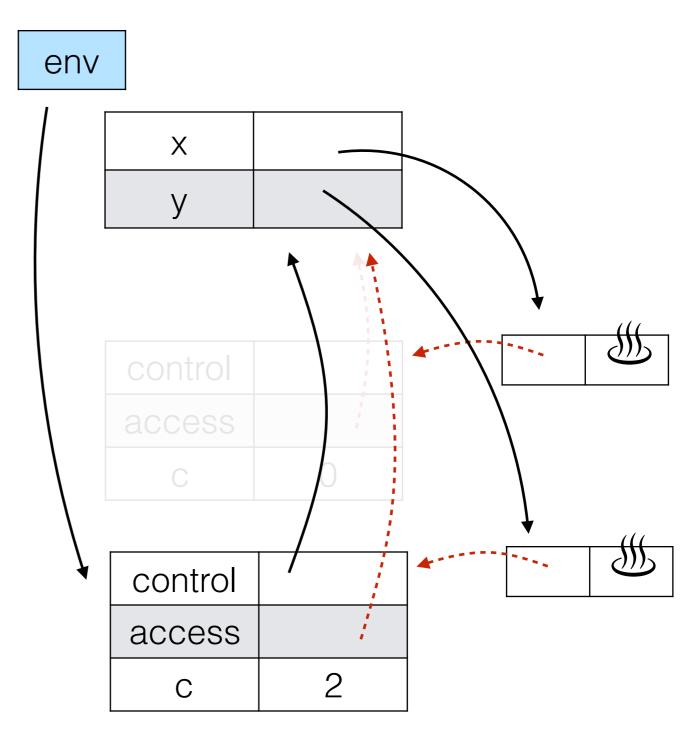
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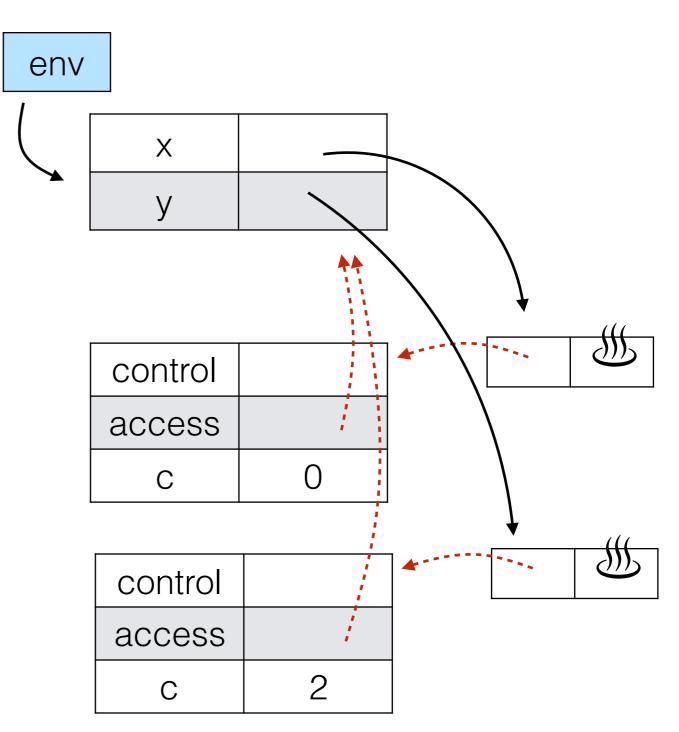
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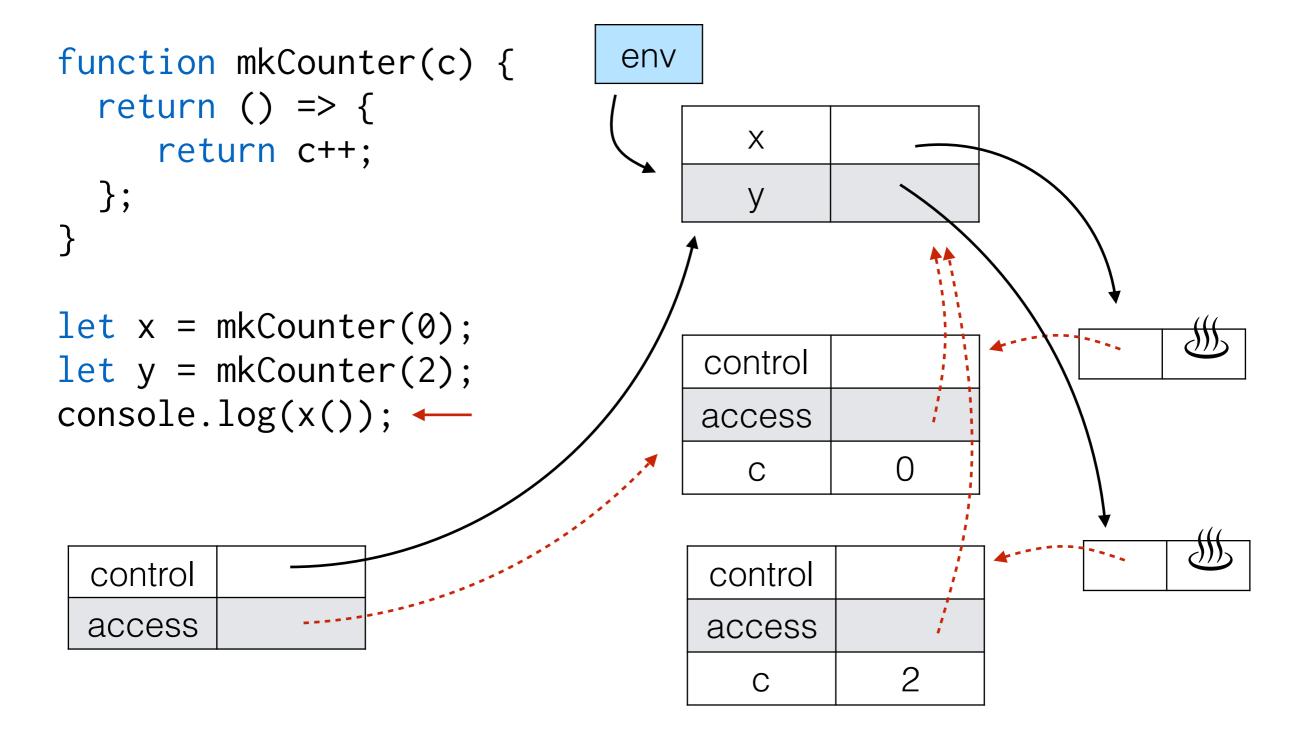
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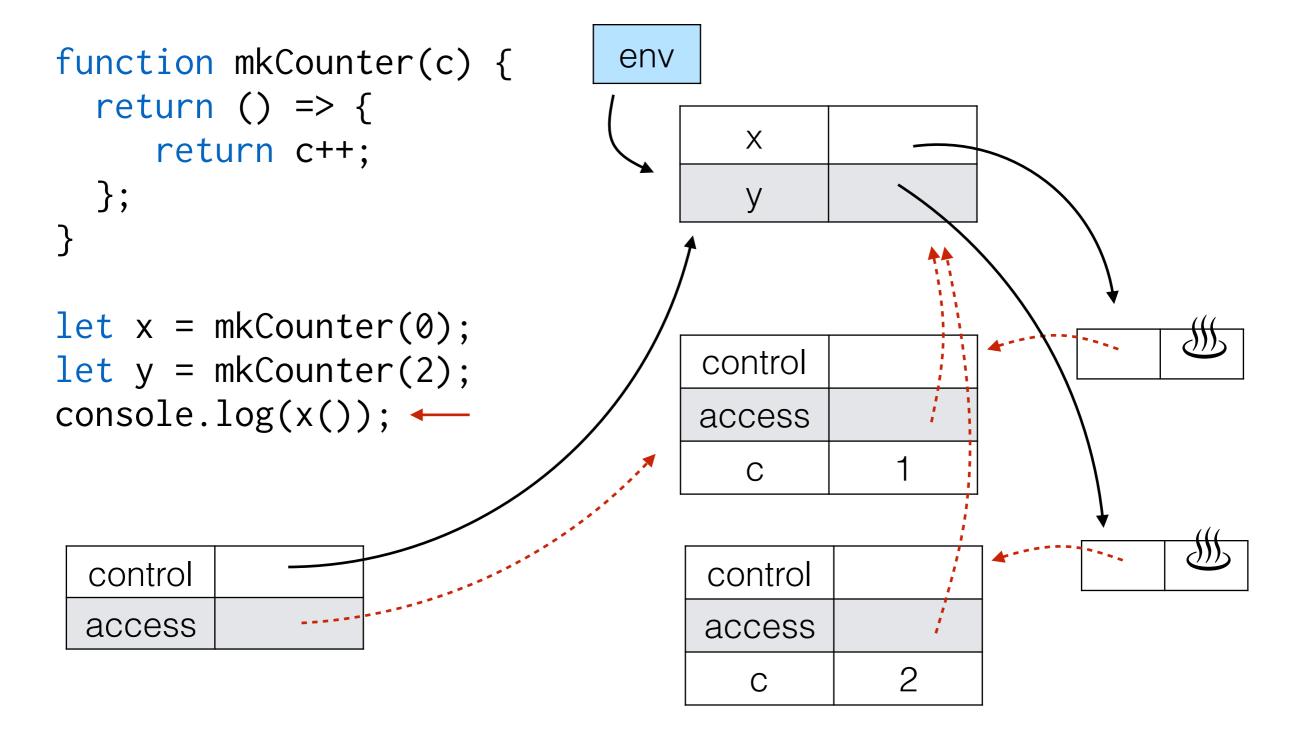


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