University of Science, VNU-HCM

Faculty of Information Technology

80\$€03



REPORT

COURSE: ARTIFICIAL INTELLIGENCE

PROJECT 02: COLORING PUZZLE

Theory Lecurer: Nguyen Ngoc Thao

Assistant Lecturer: Ho Thi Thanh Tuyen

List Student: Hoang Huu Minh An 20127102

Nguyen Vo Minh Tri 20127364

Phan Minh Xuan 20127395

80 803



Contents

| I. GROUP INFORMATION | 3 |
|--------------------------|----|
| II. ABOUT REQUIREMENTS | 3 |
| III. COMPLETE | 3 |
| IV. DESCRIBE CNFS DETAIL | 4 |
| V. PYSAT | 5 |
| VI. A* ALGORITHM | 6 |
| VII. BACKTRACKING | 6 |
| VIII. BRUTEFORCE | 6 |
| IX. DEMO | 7 |
| X. EXPERIMENT | 20 |
| XI. CONCLUSION | 21 |
| XII. REFERENCE | 21 |



I. GROUP INFORMATION

| ID | Name | Mail |
|----------|--------------------|-------------------------------|
| 20127102 | Hoang Huu Minh An | 20127102@student.hcmus.edu.vn |
| 20127364 | Nguyen Vo Minh Tri | 20127364@student.hcmus.edu.vn |
| 20127395 | Phan Minh Xuan | 20127395@student.hcmus.edu.vn |

II. ABOUT REQUIREMENTS

1. Problem

We are asked to build a coloring puzzle solver by using the first order logic to CNF as described below:

- ✓ Given a matrix of size $m \times n$, where each cell will be a non-negative integer or zero (empty cell). Each cell is considered to be adjacent to itself and 8 surrounding cells.
- ✓ Your puzzle needs to color all the cells of the matrix with either blue or red, so that the number inside each cell corresponds to the number of blue squares adjacent to that cell.

2. Approaching

- + Analyze how to define CNFs.
- + Design UIs.
- + Use the pysat library and CNFs to solve the problem.
- + Design brute-force and backtracking algorithms.
- + Implement brute-force and backtracking.
- + Connect the logic and how the UI works to emit the application.

III. COMPLETE

| No | Content | Complete | Notes |
|----|--|----------|-----------------------------|
| 1 | Describe the logical principles for generating CNFs correctly. | | Done |
| 2 | Generate CNFs automatically. | 100% | Done |
| 3 | Use the pysat library to solve CNFs correctly. | 100% | Done |
| 4 | Use A* to solve CNFs without a library. | 80% | We still some shortcomings. |
| 5 | Program brute-force algorithm. | 100% | Done |
| 6 | Program backtracking algorithm. | 100% | Done |



| 7 | Graphic interface. | 100% | Done |
|----|---|------|------|
| 8 | Give 5 test cases with different sizes. | 100% | Done |
| 9 | Comparing result and performance. | 100% | Done |
| 10 | Comply with the regulations of submission requirements. | 100% | Done |

IV. DESCRIBE CNFs DETAIL

- *Example*:
- Value of each cell start from 1 to 9.

| 2 | | |
|---|---|--|
| | 3 | |
| | | |

- Clauses that color red for the cells adjacent to a cell are a set of boolean expressions.
- At the position of No.1 has a value = 2, which means must be 2 adjacent cells it is colored green.
- KB(2):

Case 1:

- Normal clause: $(1 \land 2) \Rightarrow (\neg 4 \land \neg 5)$
- CNF: $(\neg 1 \lor \neg 2 \lor \neg 4) \land (\neg 1 \lor \neg 2 \lor \neg 5)$
- Mean: (1 \Lambda 2) is green, No.4 and No.5 are red

Case 2:

- Normal clause: $(1 \land 4) \Rightarrow (\neg 2 \land \neg 5)$
- CNF: $(\neg 1 \lor \neg 4 \lor \neg 2) \land (\neg 1 \lor \neg 2 \lor \neg 5)$
- Mean: (1 \wedge 4) is green, No.2 and No.5 are red

Case 3:

- Normal clause: $(1 \land 5) \Rightarrow (\neg 2 \land \neg 4)$
- CNF: $(\neg 1 \lor \neg 5 \lor \neg 2) \land (\neg 1 \lor \neg 5 \lor \neg 4)$



- Mean: (1 \Lambda 5) is green, No.2 and No.4 are red
- At the postion of No.5 has a value= 3, which means must be 3 adjacent cells it is colored green.
- KB(3):

Case 1:

- Normal clause: $(1 \land 2 \land 5) \Rightarrow (\neg 3 \land \neg 4 \neg 6 \land \neg 7 \neg 8 \land \neg 9)$
- CNF: $(\neg 1 \lor \neg 2 \lor \neg 5 \lor \neg 3) \land (\neg 1 \lor \neg 2 \lor \neg 5 \lor \neg 4) \land (\neg 1 \lor \neg 2 \lor \neg 5 \lor \neg 6) \land (\neg 1 \lor \neg 2 \lor \neg 5 \lor \neg 7) \land (\neg 1 \lor \neg 2 \lor \neg 5 \lor \neg 8) \land (\neg 1 \lor \neg 2 \lor \neg 5 \lor \neg 9)$
- Mean: $(1 \land 2 \land 5)$ is green, No.3, No.4, No.6, No.7, No.8, and No.9 are red Case 2:
 - Normal clause: $(1 \land 4 \land 5) \Rightarrow (\neg 2 \land \neg 3 \neg 6 \land \neg 7 \neg 8 \land \neg 9)$
 - CNF: $(\neg 1 \lor \neg 4 \lor \neg 5 \lor \neg 2) \land (\neg 1 \lor \neg 4 \lor \neg 5 \lor \neg 3) \land (\neg 1 \lor \neg 4 \lor \neg 5 \lor \neg 6) \land (\neg 1 \lor \neg 4 \lor \neg 5 \lor \neg 7) \land (\neg 1 \lor \neg 4 \lor \neg 5 \lor \neg 8) \land (\neg 1 \lor \neg 4 \lor \neg 5 \lor \neg 9)$
- Mean: $(1 \land 4 \land 5)$ is green, No.2, No.3, No.6, No.7, No.8, and No.9 are red In another case in KB (3), we do similar to case 1 and case 2.

Conclution:

- Clauses that color green for the cells adjacent to a cell is all the combinations (k-combination of set S) is a set of conjunction of a set of disjunction of a combination:
 - S is a set of all cells adjacent to that cell.
 - N is the number of cells adjacent to that cell (includes ifself).
 - M is the number of cells that need to color green.
 - K is the difference between n and m.

We have KB (2)

 $KB (C) = KB (2) \land KB (3)$

 $KB(D) = KB(C) \land KB(4)$ (with KB(4) in the case knowledge base of the next cell which has value).

•

Finally, we will find the CNFs clause's correct requirements.

V. PYSAT

- Using Glucose 3 in pysat.solver lib.
- Using add_clause() to add each clause in a set of CNF clause.
- Using solve() function to solve the problem. This function will return True or return False.
- Finally, use get_model() to get a value of all variables in CNF clause. Accordingly, it is possible to derive a matrix of numbers 0 and 1, representing the colors, 1 being green and 1 being red.

5

Class 20CLC – Term III/2021-2022 Course: CSC14003 – Artificial Intelligence



VI. A* ALGORITHM

- Algorithm A* is an optimization algorithm that uses heuristic evaluation to find the next move. So, the A* algorithm is heavily dependent on the heuristic functions. If the difference in heuristic estimates is large enough, it will lead to false results.
- Heristic function is built based on a total number of cell was colored by blue with the number of red cell was valid colored and subtract the number of cells is visual in each clause.

VII. BACKTRACKING

- Loop through all cells that contain a number from left to right and top to bottom.
- At each cell, get all the combinations (k-combination of a set S):
 - + S is a set of all the red cells adjacent to that cell (includes itself).
 - + k is the number of remaining cells that need to color green.
- Coloring all the cells in a combination (set of numbers) to green and jump to the next cell that contains a number.
- If the current cell has the number of the green cells adjacent to itself that is greater than the number inside or looped through all combinations of that cell, then the backtracking function will go back to the previous cell and try another combination.
- The backtracking function will run until the completion of coloring the last cell that contains a number.

VIII. BRUTEFORCE

- Coloring all cells in a combination to green and go to the next cell that contains a number.
- If the current cell has the number of the green cells adjoining to itself that is higher than the wide variety interior or looped via all combinations of that cell, then the brute pressure function will go back to the preceding cell and try some other combination.
- The brute force feature will run till the completion of coloring the last cell that includes a variety.

fit@hcmus

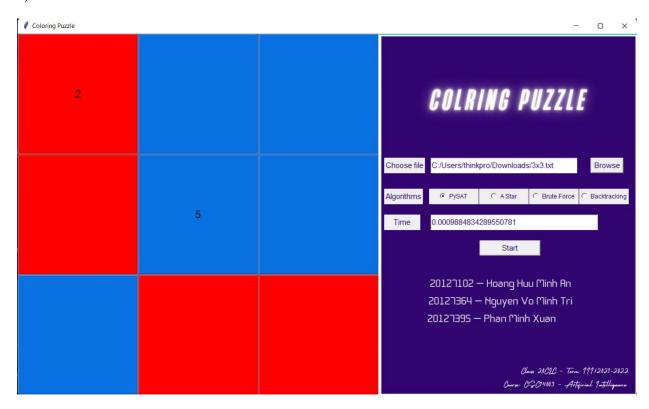
IX. DEMO

User Interface



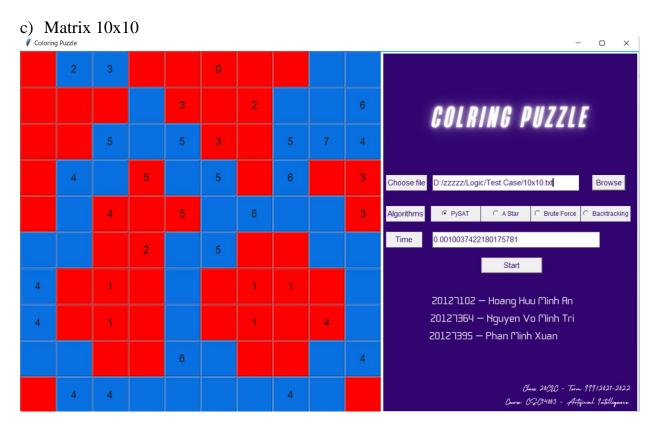
1. Using PySat

a) Matrix 3x3









Start

20127102 — Hoang Huu Minh An

20127364 — Nguyen Vo Minh Tri

Class 20CIC - Term 111/2021-2022 Curse: CGC14003 - Artificial Intelligence

20127395 — Phan Minh Xuan



d) Matrix 15x15 5 4 4 3 5 COLRING PUZZLE 5 4 5 4 Choose file D:/zzzzz/Logic/Test Case/15x15.txt Browse 4 Algorithms 6 5 3 0.0012469291687011719 Time

3 5

4

4

6

5

4

4

5

6 4



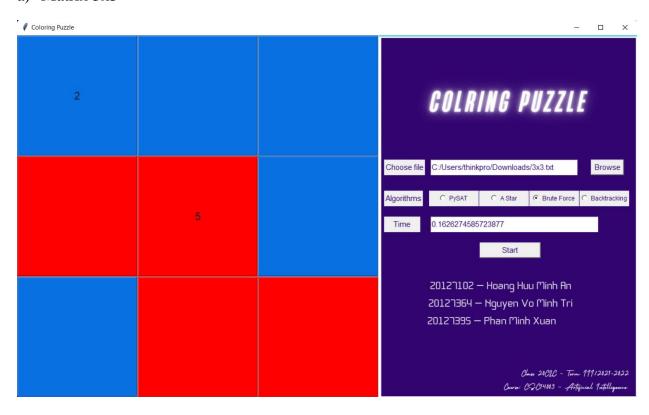


| No. | Input | Output |
|-----|--|---|
| 1 | 3 3 | 011 |
| | 2 -1 -1 | 0 1 1 |
| | -1 5 -1 | 100 |
| | -1 -1 -1 | |
| 2 | 5 5 | 00110 |
| | 0 -1 4 4 -1 | 00110 |
| | -1 4 -1 6 -1 | 11011 |
| | 3 -1 7 6 -1 | 01110 |
| | -1 6 -1 6 5 | 11011 |
| | -1 -1 -1 3 | |
| 3 | 10 10 | 0110000011 |
| | -1 2 3 -1 -1 0 -1 -1 -1 | 0001000111 |
| | -1 -1 -1 -1 3 -1 2 -1 -1 6 | 0011100111 |
| | -1 -1 5 -1 5 3 -1 5 7 4 | 0110110100 |
| | -1 4 -1 5 -1 5 -1 6 -1 3 | 0100011110 |
| | -1 -1 4 -1 5 -1 6 -1 -1 3 | 1100110011 |
| | -1 -1 -1 2 -1 5 -1 -1 -1 | 1000100001 |
| | 4-11-1-1111-1-1 | 1000100001 |
| | 4-11-1-11-14-1 | 1100110011 |
| | -1 -1 -1 -1 6 -1 -1 -1 4 | 011111110 |
| 4 | -1 4 4 -1 -1 -1 -1 4 -1 -1 | 0.00110000011110 |
| 4 | 15 15 | 000110000011110 |
| | 0-1-14321-1-1-1-13-1-1 | 001010011100001 |
| | -1 -1 5 -1 -1 4 -1 -1 4 4 -1 -1 -1 -1 3 -1 5 4 5 4 5 5 -1 5 3 -1 1 2 -1 3 | $egin{array}{c} 1\ 1\ 1\ 1\ 0\ 1\ 1\ 0\ 0\ 0\ 0\ 0\ 1 \\ 1\ 0\ 0\ 0\ 1\ 0\ 1\ 0\ 0\ 0\ 0\ 1\ 1\ 0 \\ \end{array}$ |
| | -1 3 4 3 4 3 3 -1 3 3 -1 1 2 -1 3 4 -1 -1 -1 4 -1 -1 4 2 -1 1 -1 -1 -1 -1 | 011100000110 |
| | -1 -1 5 4 -1 2 2 -1 1 0 -1 -1 7 5 -1 | 010100000001100 |
| | -1 -1 -1 5 4 -1 2 2 -1 1 0 -1 -1 7 5 -1 | 1110100000001110 |
| | 4-1-154200-1-1-156-1-1 | 100110000000001 |
| | 5 -1 -1 6 5 -1 -1 -1 -1 3 3 3 -1 3 | 11110100110000 |
| | -1 -1 5 -1 5 3 -1 -1 -1 -1 -1 3 3 3 1 -1 | 10010011100000 |
| | 5-1-165-135-16-1-10-10 | 111110001010000 |
| | -1 -1 5 -1 4 3 2 4 5 -1 4 -1 -1 1 -1 | 1100100010100 |
| | -17-1-15-1-11-1555-1-1 | 011101000110110 |
| | -1 -1 6 4 4 4 3 1 2 4 -1 -1 6 4 -1 | 011001100011010 |
| | -1 5 -1 6 -1 -1 -1 -1 4 6 -1 -1 -1 | 001110100011001 |
| | -1 -1 -1 -1 -1 -1 3 2 0 -1 4 4 3 -1 2 | |
| 5 | 20 20 | 00000111111110000100 |
| | -1 0 -1 -1 -1 -1 4 -1 -1 4 -1 -1 5 -1 2 -1 -1 2 1 0 | 00001100100111101000 |
| | -1 -1 -1 -1 5 -1 -1 -1 -1 -1 -1 -1 -1 5 -1 -1 4 -1 -1 | 00011001111011111010 |
| | -1 0 2 -1 6 -1 -1 4 -1 -1 6 -1 -1 8 6 6 5 -1 2 -1 | 00010110011111001100 |
| | -1 1 -1 5 -1 6 6 -1 -1 7 7 -1 8 -1 -1 -1 -1 -1 -1 | 00101111101111010110 |
| | 0 -1 -1 -1 -1 6 6 -1 6 -1 7 -1 -1 6 -1 -1 6 6 -1 -1 | 00111001110110111010 |



2. Using Brute-Force

a) Matrix 3x3





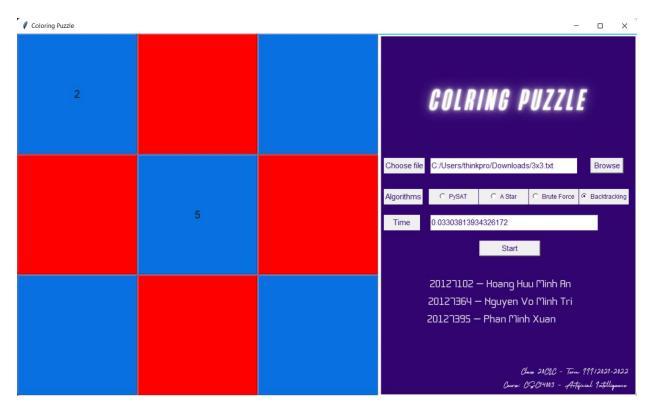
| No. | Input | Output |
|-----|--|-------------------------------|
| 1 | 3 3 | 011 |
| | 2 -1 -1 | 0 1 1 |
| | -1 5 -1 | 100 |
| | -1 -1 -1 | |
| 2 | 5 5 | |
| | 0 -1 4 4 -1 | |
| | -1 4 -1 6 -1 | Running Time: >3600s (1 hour) |
| | 3 -1 7 6 -1 | |
| | -1 6 -1 6 5 | |
| | -1 -1 -1 3 | |
| 3 | 10 10 | |
| | -1 2 3 -1 -1 0 -1 -1 -1 -1 | |
| | -1 -1 -1 -1 3 -1 2 -1 -1 6 | |
| | -1 -1 5 -1 5 3 -1 5 7 4 | |
| | -14-15-16-13 | D |
| | -1 -1 4 -1 5 -1 6 -1 -1 3 | Runnung Time: >3600s (1 hour) |
| | -1 -1 -1 2 -1 5 -1 -1 -1 -1 4 -1 1 -1 -1 1 1 1 -1 -1 | |
| | 4-11-1-111-1-1 | |
| | -1 -1 -1 -1 -1 -1 4 -1 | |
| | -1 -1 -1 -1 -1 -1 -1 -1 -1 -1 4 -1 -1 -1 | |
| 4 | 15 15 | |
| _ | 0 -1 -1 4 3 2 1 -1 -1 -1 -1 3 -1 -1 | |
| | -1 -1 5 -1 -1 4 -1 -1 4 4 -1 -1 -1 3 | |
| | -1545455-153-112-13 | |
| | 4 -1 -1 -1 4 -1 -1 4 2 -1 1 -1 -1 -1 | |
| | -1 -1 5 4 -1 2 2 -1 1 0 -1 -1 7 5 -1 | |
| | -1 -1 -1 5 -1 -1 0 -1 -1 -1 -1 4 5 -1 2 | |
| | 4 -1 -1 5 4 2 0 0 -1 -1 -1 5 6 -1 -1 | |
| | 5 -1 -1 6 5 -1 -1 -1 -1 -1 3 3 3 -1 3 | Runnung Time: >3600s (1 hour) |
| | -1 -1 5 -1 5 3 -1 -1 -1 -1 -1 3 -1 -1 | |
| | 5 -1 -1 6 5 -1 3 5 -1 6 -1 -1 0 -1 0 | |
| | -1 -1 5 -1 4 3 2 4 5 -1 4 -1 -1 1 -1 | |
| | -17-1-15-1-11-1555-1-1-1 | |
| | -1 -1 6 4 4 4 3 1 2 4 -1 -1 6 4 -1 | |
| | -15-16-1-1-1-146-1-1-1 | |
| | -1 -1 -1 -1 -1 -1 3 2 0 -1 4 4 3 -1 2 | |
| 5 | 20 20 | |
| | -1 0 -1 -1 -1 -1 4 -1 -1 4 -1 -1 5 -1 2 -1 -1 2 1 0 -1 -1 -1 5 -1 -1 -1 -1 -1 -1 -1 5 -1 -1 4 -1 -1 | |
| | -1 -1 -1 -1 3 -1 -1 -1 -1 -1 -1 -1 -1 3 -1 -1 4 -1 -1 -1 0 2 -1 6 -1 -1 4 -1 -1 6 -1 -1 8 6 6 5 -1 2 -1 | |
| | -1 1 -1 5 -1 6 6 -1 -1 7 7 -1 8 -1 -1 -1 -1 -1 -1 | |
| | 0-1-1-1-1-1-1 | |
| | -1 -1 5 7 -1 -1 7 7 -1 -1 7 7 -1 5 5 -1 -1 3 | |
| | 1 1 3 1 1 1 -1 -1 -1 3 3 -1 -1 -1 3 | |



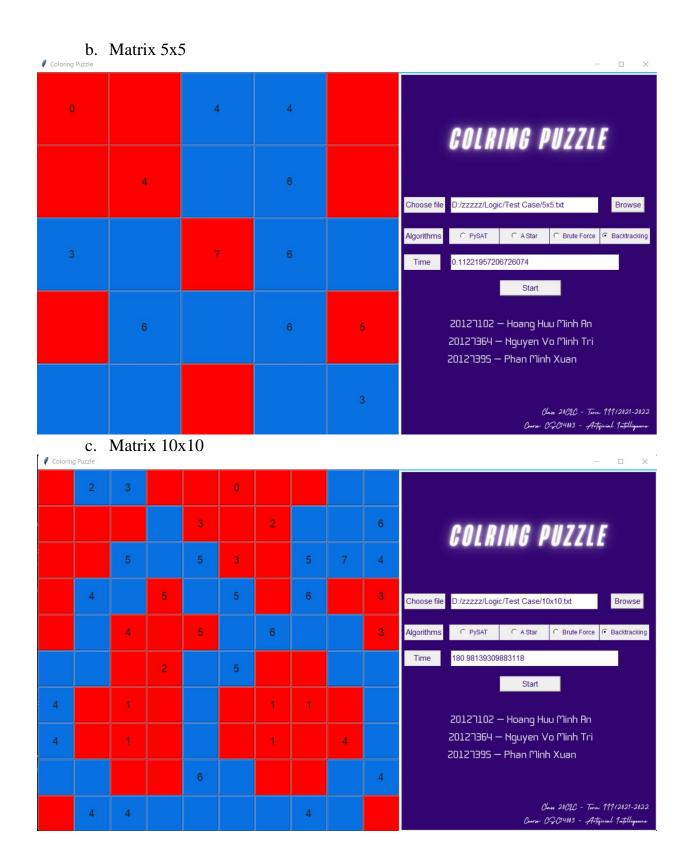
Runnung Time: >3600s (1 hour)

3. Using Backtracking

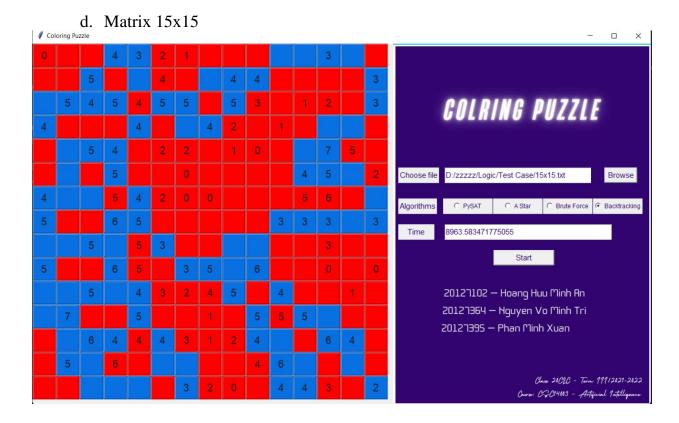
a. Matrix 3x3











| No. | Input | Output |
|-----|----------------------------|---------------------|
| 1 | 3 3 | 011 |
| | 2 -1 -1 | 011 |
| | -1 5 -1 | 100 |
| | -1 -1 -1 | |
| 2 | 5 5 | 00110 |
| | 0 -1 4 4 -1 | 00110 |
| | -1 4 -1 6 -1 | 1 1 0 1 1 |
| | 3 -1 7 6 -1 | 01110 |
| | -1 6 -1 6 5 | 1 1 0 1 1 |
| | -1 -1 -1 -1 3 | |
| 3 | 10 10 | 0110000011 |
| | -1 2 3 -1 -1 0 -1 -1 -1 | 0001000111 |
| | -1 -1 -1 -1 3 -1 2 -1 -1 6 | 0011100111 |
| | -1 -1 5 -1 5 3 -1 5 7 4 | 0110110100 |
| | -1 4 -1 5 -1 5 -1 6 -1 3 | 0100011110 |
| | -1 -1 4 -1 5 -1 6 -1 -1 3 | 1 1 0 0 1 1 0 0 1 1 |
| | -1 -1 -1 2 -1 5 -1 -1 -1 | 1000100001 |
| | 4 -1 1 -1 -1 -1 1 1 -1 -1 | 1 0 0 0 1 0 0 0 0 1 |
| | 4 -1 1 -1 -1 -1 1 -1 4 -1 | 1 1 0 0 1 1 0 0 1 1 |
| | -1 -1 -1 -1 6 -1 -1 -1 4 | 011111110 |



| | | T |
|---|--|-------------------------------|
| | -1 4 4 -1 -1 -1 4 -1 -1 | |
| 4 | 15 15 | 000110000011110 |
| | 0 -1 -1 4 3 2 1 -1 -1 -1 -1 -1 3 -1 -1 | 001010011100001 |
| | -1 -1 5 -1 -1 4 -1 -1 4 4 -1 -1 -1 3 | 111101101000001 |
| | -1 5 4 5 4 5 5 -1 5 3 -1 1 2 -1 3 | 100010110000110 |
| | 4 -1 -1 -1 4 -1 -1 4 2 -1 1 -1 -1 -1 -1 | 011100000001100 |
| | -1 -1 5 4 -1 2 2 -1 1 0 -1 -1 7 5 -1 | 010100000001110 |
| | -1 -1 -1 5 -1 -1 0 -1 -1 -1 -1 4 5 -1 2 | 1 1 1 0 1 0 0 0 0 0 0 0 0 1 |
| | 4 -1 -1 5 4 2 0 0 -1 -1 -1 5 6 -1 -1 | 100110000011111 |
| | 5 -1 -1 6 5 -1 -1 -1 -1 -1 3 3 3 -1 3 | 11110100110000 |
| | -1 -1 5 -1 5 3 -1 -1 -1 -1 -1 3 -1 -1 | 100100111100000 |
| | 5 -1 -1 6 5 -1 3 5 -1 6 -1 -1 0 -1 0 | 111110001010000 |
| | -1 -1 5 -1 4 3 2 4 5 -1 4 -1 -1 1 -1 | 110010000101100 |
| | -1 7 -1 -1 5 -1 -1 1 -1 5 5 5 -1 -1 -1 | 011101000110110 |
| | -1 -1 6 4 4 4 3 1 2 4 -1 -1 6 4 -1 | 011001100011010 |
| | -1 5 -1 6 -1 -1 -1 -1 4 6 -1 -1 -1 -1 | 001110100011001 |
| | -1 -1 -1 -1 -1 -1 3 2 0 -1 4 4 3 -1 2 | |
| 5 | 20 20 | |
| | -1 0 -1 -1 -1 -1 4 -1 -1 4 -1 -1 5 -1 2 -1 -1 2 1 0 | |
| | -1 -1 -1 -1 5 -1 -1 -1 -1 -1 -1 -1 -1 5 -1 -1 4 -1 -1 | |
| | -1 0 2 -1 6 -1 -1 4 -1 -1 6 -1 -1 8 6 6 5 -1 2 -1 | |
| | -1 1 -1 5 -1 6 6 -1 -1 7 7 -1 8 -1 -1 -1 -1 -1 -1 -1 | |
| | 0 -1 -1 -1 -1 6 6 -1 6 -1 7 -1 -1 6 -1 -1 6 6 -1 -1 | |
| | -1 -1 5 7 -1 -1 -1 7 7 -1 -1 7 7 -1 5 5 -1 -1 -1 3 | |
| | -1 4 6 -1 -1 6 5 -1 5 3 -1 -1 -1 -1 4 5 5 -1 -1 -1 | |
| | -1 -1 -1 -1 8 8 -1 4 -1 -1 -1 -1 7 6 -1 -1 6 6 4 -1 | |
| | -1 -1 -1 6 8 -1 6 -1 -1 -1 4 -1 -1 -1 6 -1 7 -1 -1 -1 | |
| | 5 -1 6 -1 -1 -1 5 -1 2 -1 4 5 -1 7 8 -1 -1 7 -1 -1 | Runnung Time: >3600s (1 hour) |
| | 67-165-1-1-1-1-1-1-175-17-13 | |
| | -1 -1 -1 -1 -1 3 0 -1 2 -1 2 -1 -1 -1 7 -1 -1 6 -1 -1 | |
| | 4 -1 4 6 4 -1 -1 -1 -1 -1 -1 2 4 4 -1 -1 -1 -1 -1 -1 | |
| | -1 -1 -1 -1 -1 -1 -1 3 -1 5 4 -1 -1 -1 -1 5 -1 5 -1 1 | |
| | -1 5 5 -1 7 7 -1 -1 5 5 -1 -1 6 -1 7 -1 -1 4 -1 -1 | |
| | -1 -1 -1 -1 -1 -1 -1 4 -1 -1 -1 6 -1 9 6 6 -1 -1 -1 -1 | |
| | 5 -1 -1 4 7 5 -1 3 3 -1 -1 -1 8 8 5 -1 -1 -1 7 -1 | |
| | 4 -1 -1 4 -1 4 4 2 3 -1 -1 7 7 -1 4 -1 -1 6 -1 6 | |
| | -1 -1 -1 -1 -1 4 -1 -1 0 0 2 -1 -1 -1 -1 -1 -1 -1 5 | |
| | -1 4 -1 5 -1 -1 -1 -1 -1 -1 -1 4 3 -1 -1 3 4 -1 -1 | |

4. Using A*
a) Matrix 3x3











| No. | Input | Output |
|-----|----------------------------|---------------------|
| 1 | 3 3 | 0 1 1 |
| | 2 -1 -1 | 0 1 1 |
| | -1 5 -1 | 100 |
| | -1 -1 -1 | |
| 2 | 5 5 | 00110 |
| | 0 -1 4 4 -1 | 00110 |
| | -1 4 -1 6 -1 | 11011 |
| | 3 -1 7 6 -1 | 01110 |
| | -1 6 -1 6 5 | 11011 |
| | -1 -1 -1 3 | |
| 3 | 10 10 | 0110000011 |
| | -1 2 3 -1 -1 0 -1 -1 -1 | 0001000111 |
| | -1 -1 -1 -1 3 -1 2 -1 -1 6 | 0011100111 |
| | -1 -1 5 -1 5 3 -1 5 7 4 | 0110110100 |
| | -1 4 -1 5 -1 5 -1 6 -1 3 | 0100011110 |
| | -1 -1 4 -1 5 -1 6 -1 -1 3 | 1100110011 |
| | -1 -1 -1 2 -1 5 -1 -1 -1 | 1000100001 |
| | 4 -1 1 -1 -1 -1 1 1 -1 -1 | 1000100001 |
| | 4 -1 1 -1 -1 -1 1 -1 4 -1 | 1 1 0 0 1 1 0 0 1 1 |

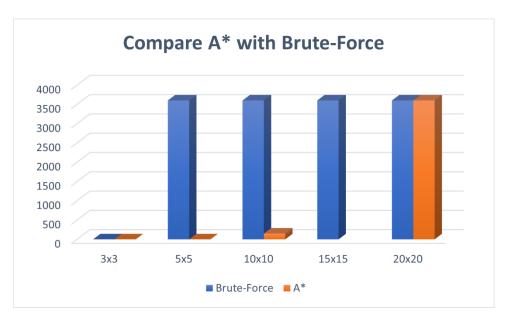


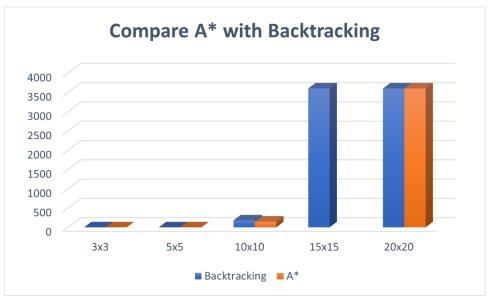
| | -1 -1 -1 -1 6 -1 -1 -1 4 | 011111110 |
|---|--|---------------------------------|
| | -1 4 4 -1 -1 -1 4 -1 -1 | |
| 4 | 15 15 | |
| • | 0 -1 -1 4 3 2 1 -1 -1 -1 -1 3 -1 -1 | |
| | -1 -1 5 -1 -1 4 -1 -1 4 4 -1 -1 -1 3 | |
| | -1545455-153-112-13 | |
| | 4-1-1-14-1-142-11-1-1-1 | |
| | -1 -1 5 4 -1 2 2 -1 1 0 -1 -1 7 5 -1 | |
| | -1 -1 -1 5 -1 -1 0 -1 -1 -1 4 5 -1 2 | |
| | 4-1-154200-1-1-156-1-1 | Running Time: >3600s (1 hour) |
| | 5 -1 -1 65 -1 -1 -1 -1 3 3 3 -1 3 | |
| | -1 -1 5 -1 5 3 -1 -1 -1 -1 -1 3 -1 -1 | |
| | 5 -1 -1 65 -1 35 -1 6 -1 -1 0 -1 0 | |
| | -1 -1 5 -1 4 3 2 4 5 -1 4 -1 -1 1 -1 | |
| | -17-1-15-1-11-1555-1-1-1 | |
| | -1 -1 6 4 4 4 3 1 2 4 -1 -1 6 4 -1 | |
| | -1 5 -1 6 -1 -1 -1 -1 4 6 -1 -1 -1 -1 | |
| | -1 -1 -1 -1 -1 -1 3 2 0 -1 4 4 3 -1 2 | |
| 5 | 20 20 | |
| | -1 0 -1 -1 -1 -1 4 -1 -1 4 -1 -1 5 -1 2 -1 -1 2 1 0 | |
| | -1 -1 -1 -1 5 -1 -1 -1 -1 -1 -1 -1 -1 5 -1 -1 4 -1 -1 | |
| | -1 0 2 -1 6 -1 -1 4 -1 -1 6 -1 -1 8 6 6 5 -1 2 -1 | |
| | -1 1 -1 5 -1 6 6 -1 -1 7 7 -1 8 -1 -1 -1 -1 -1 -1 -1 | |
| | 0 -1 -1 -1 6 6 -1 6 -1 7 -1 -1 6 -1 -1 6 6 -1 -1 | |
| | -1 -1 5 7 -1 -1 -1 7 7 -1 -1 7 7 -1 5 5 -1 -1 -1 3 | |
| | -1 4 6 -1 -1 6 5 -1 5 3 -1 -1 -1 -1 4 5 5 -1 -1 -1 | |
| | -1 -1 -1 -1 88 -1 4 -1 -1 -1 7 6 -1 -1 6 6 4 -1 | |
| | -1 -1 -1 68 -1 6 -1 -1 -1 4 -1 -1 -1 6 -1 7 -1 -1 -1 | D 5000 |
| | 5 - 1 6 - 1 - 1 5 - 1 2 - 1 4 5 - 1 7 8 - 1 - 1 7 - 1 - 1 | Runnung Time: >3600s (1 hour) |
| | 67-165-1-1-1-1-1-1-175-17-13 | |
| | -1 -1 -1 -1 -1 3 0 -1 2 -1 2 -1 -1 7 -1 -1 6 -1 -1 | |
| | 4-1464-1-1-1-1-1244-1-1-1-1-1-1 | |
| | -1 -1 -1 -1 -1 -1 -1 3 -1 5 4 -1 -1 -1 5 -1 5 -1 1 155 177 1 155 1 16 17 1 14 1 1 | |
| | -1 5 5 -1 7 7 -1 -1 5 5 -1 -1 6 -1 7 -1 -1 4 -1 -1 -1 -1 -1 -1 -1 -1 -1 4 -1 -1 -1 6 -1 9 6 6 -1 -1 -1 -1 | |
| | 5-1-1475-133-1-1-1885-1-1-17-1 | |
| | 4-1-14-14423-1-177-14-1-16-16 | |
| | -1 -1 -1 -1 -1 4 -1 -1 0 0 2 -1 -1 -1 -1 -1 -1 -1 5 | |
| | -1 4 -1 5 -1 -1 -1 -1 -1 -1 -1 -1 -1 -1 -1 -1 -1 | |
| | -1 4 -1 3 -1 -1 -1 -1 -1 -1 -1 4 3 -1 -1 3 4 -1 -1 | |



X. EXPERIMENT

| Algorithms | Running Time (seconds) | | | | |
|--------------------|-------------------------|----------|-----------|-----------|----------|
| | 3x3 5x5 10x10 | | 15x15 | 20x20 | |
| Pysat | 0.000988 | 0.001041 | 0.001003 | 0.0012469 | 0.003014 |
| Brute-force | 0.126274 | > 3600s | > 3600s | > 3600s | > 3600s |
| Backtracking | 0.033038 | 0.112219 | 180.98139 | 8963.5834 | > 3600s |
| A* | A* 0.017 0.17742 | | 149.39893 | >3600s | >3600s |





Class 20CLC – Term III/2021-2022 Course: CSC14003 – Artificial Intelligence



XI. CONCLUSION

- Algorithm A* is an optimization algorithm that uses heuristic evaluation to find the next move. So, the A* algorithm is heavily dependent on the heuristic functions. If the difference in heuristic estimates is large enough, it will lead to false results. The disadvantage of the A* algorithm is that it consumes a lot of memory.
- Algorithms Brute-Force and Backtracking applied in this problem will always have the right results, but their disadvantages is that the running time is very long, depending on the size of the matrix. We are able to use CNF clause with pysat module of Python to solve satisfied problems.

XII. REFERENCE

Material of lecturer:

https://l.facebook.com/l.php?u=https%3A%2F%2Fdrive.google.com%2Fdrive%2Ffolders%2F1OmW_a959zfyCfWP_agipFzfEZX_eV1S%3Ffbclid%3DIwAR2Zfdd40XKkkqm9w-

BgKgYdVy8HuAT1EXkR4BNO9IGSArloavoF9JtE52g&h=AT2Jk4i7ueaCNw HenSrJE111W4IlTnmh3XvSw3iMnynnDwfhCUiAci6zEMf8Y0XCT1DKBLU3 nKF1fKl8qc6Z4YHGsaun8WSNtTDmzTgsFgUFzNMvjEvRRT4z049bAgsRvs4 y0ALzVrxBBO5sQY6iYg

- A* Search Brilliant Math & Science Wifi:

https://brilliant.org/wiki/a-star-

<u>search/?fbclid=IwAR0DRY9FXAyIxPArXYJv6itfMEwFfoUxe6UHasLyGrtHf5</u>
<u>LNbWveNvEFYpE</u>

- Python – GUI Programming (Tkinter)
https://www.tutorialspoint.com/python/python_gui_programming.htm?fbclid=Iw
AR0WHvK015v37WH6LUcUHKs_rUofJCZrLlGmH-

bXE2P2gyxXH6TZTGee-8Y

- Backtracking – Wifipedia

 $\frac{https://en.wikipedia.org/wiki/Backtracking?fbclid=IwAR2Zfdd40XKkkqm9w-BgKgYdVy8HuAT1EXkR4BNO9IGSArloavoF9JtE52g}{}$

- Brute-force Wikipedia

https://en.wikipedia.org/wiki/Brute-

force_search?fbclid=IwAR0oRD_PviqRbwoKuQcT-

mxby3tCOZDODZ6cSo9BnsagnMt3olUDiMEvHOM

∞END∝