Experiment No. 5
Data Stream Algorithms: Implement Bloom filter algorithm using any programming language
Date of Performance: 14/08/2023
Date of Submission: 21/08/2023

Aim:

Data Stream Algorithms: Implement bloom filter algorithm using any programming language

Theory:

Bloom filter algorithm approximates the number of unique objects in a stream or a database in one pass. If the stream contains n elements with m of them unique, this algorithm runs in $O(n)$ time and needs $O(\log(m))$ memory.

Algorithm:

1.   Create a bit vector (bit array) of sufficient length L, such that $2L > n$, the number of elements in the stream. Usually a 64-bit vector is sufficient since 264 is quite large for most purposes.

2.   The i-th bit in this vector/array represents whether we have seen a hash function value whose binary representation ends in $0i$. So initialize each bit to

3.   The i-th bit in this vector/array represents whether we have seen a hash function value whose binary representation ends in $0i$. So initialize each bit to

4.   The i-th bit in this vector/array represents whether we have seen a hash function value whose binary representation ends in $0i$. So initialize each bit to

5.   Once input is exhausted, get the index of the first 0 in the bit array (call this R). By the way, this is just the number of consecutive 1s (i.e. we have seen $0,00,...,0R{-}1$ as the output of the hash function) plus one.

6.   Calculate the number of unique words as $2R/\phi$, where ϕ is 0.77351. A proof for this can be found in the original paper listed in the reference section.

7.   The standard deviation of R is a constant: $\sigma(R)=1.12$. (In other words, R can be off by about 1 for 1 - 0.68 = 32% of the observations, off by 2 for about 1 - 0.95 = 5% of the observations, off by 3 for 1 - 0.997 = 0.3% of the observations using the Empirical rule of statistics). This implies that our count can be off by a factor of 2 for 32% of the observations, off by a factory of 4 for 5% of the observations, off by a factor of 8 for 0.3% of the observations and so on.

Code:

```
n = 20 #no of items to add

p = 0.05 #false positive probability

bloomf = BloomFilter(n,p)

print("Size of bit array:{}".format(bloomf.size))

print("False positive Probability:{}".format(bloomf.fp_prob))
print("Number of hash functions:{}".format(bloomf.hash_count))
# words to be added
            word_present = ['abound','abounds','abundance','abundant','accessible',
                        'bloom','blossom','bolster','bonny','bonus','bonuses',
                        'coherent','cohesive','colorful','comely','comfort',
                        'gems','generosity','generous','generously','genial']
# word not added
            word_absent = ['bluff','cheater','hate','war','humanity',
                        'racism','hurt','nuke','gloomy','facebook',
                        'geeksforgeeks','twitter']
    for item in word_present:
        bloomf.add(item)


shuffle(word_present) shuffle(word_absent)
test_words = word_present[:10] + word_absent shuffle(test_words)
for word in test_words:
    if bloomf.check(word):
        if word in word_absent:

            print("'{}' is a false positive!".format(word)) else:
            print("'{}' is probably present!".format(word))

    else:

        print("'{}' is definitely not present!".format(word))
```

Output:

```
ubuntu@ubuntu-HP-Elite-Tower-600-G9-Desktop-PC:~/bloomfilter$ python3 bloom_test.py
Size of bit array:124
False positive Probability:0.05
Number of hash functions:4
'gloomy' is definitely not present!
'cohesive' is probably present!
'geeksforgeeks' is definitely not present!
'hate' is definitely not present!
'bluff' is definitely not present!
'abundant' is probably present!
'nuke' is definitely not present!
'twitter' is a false positive!
'cheater' is definitely not present!
'accessible' is probably present!
'bonus' is probably present!
'generosity' is probably present!
'comely' is probably present!
'genial' is probably present!
'humanity' is a false positive!
'comfort' is probably present!
'war' is definitely not present!
'generous' is probably present!
'bolster' is probably present!
'facebook' is definitely not present!
'hurt' is definitely not present!
'racism' is definitely not present!
ubuntu@ubuntu-HP-Elite-Tower-600-G9-Desktop-PC:~/bloomfilter$ 
```

Conclusion:

Through the use of a programming language, this experiment demonstrates how the Bloom filter approach may be used to efficiently estimate the number of unique objects in a data stream. The Bloom filter uses a bit array and hash algorithms to evaluate whether an element is obviously absent or most likely present in the stream. The code showed how to design a Bloom filter and tested it using a list of words, some of which were present and some of which weren't, with a set false positive probability. The experiment notes that, due to their probabilistic nature, Bloom filters can sometimes yield false positive findings, but it also demonstrates how useful they can be when approximate cardinality estimation and fast membership checks are required.