

# **Introduction to STARK**

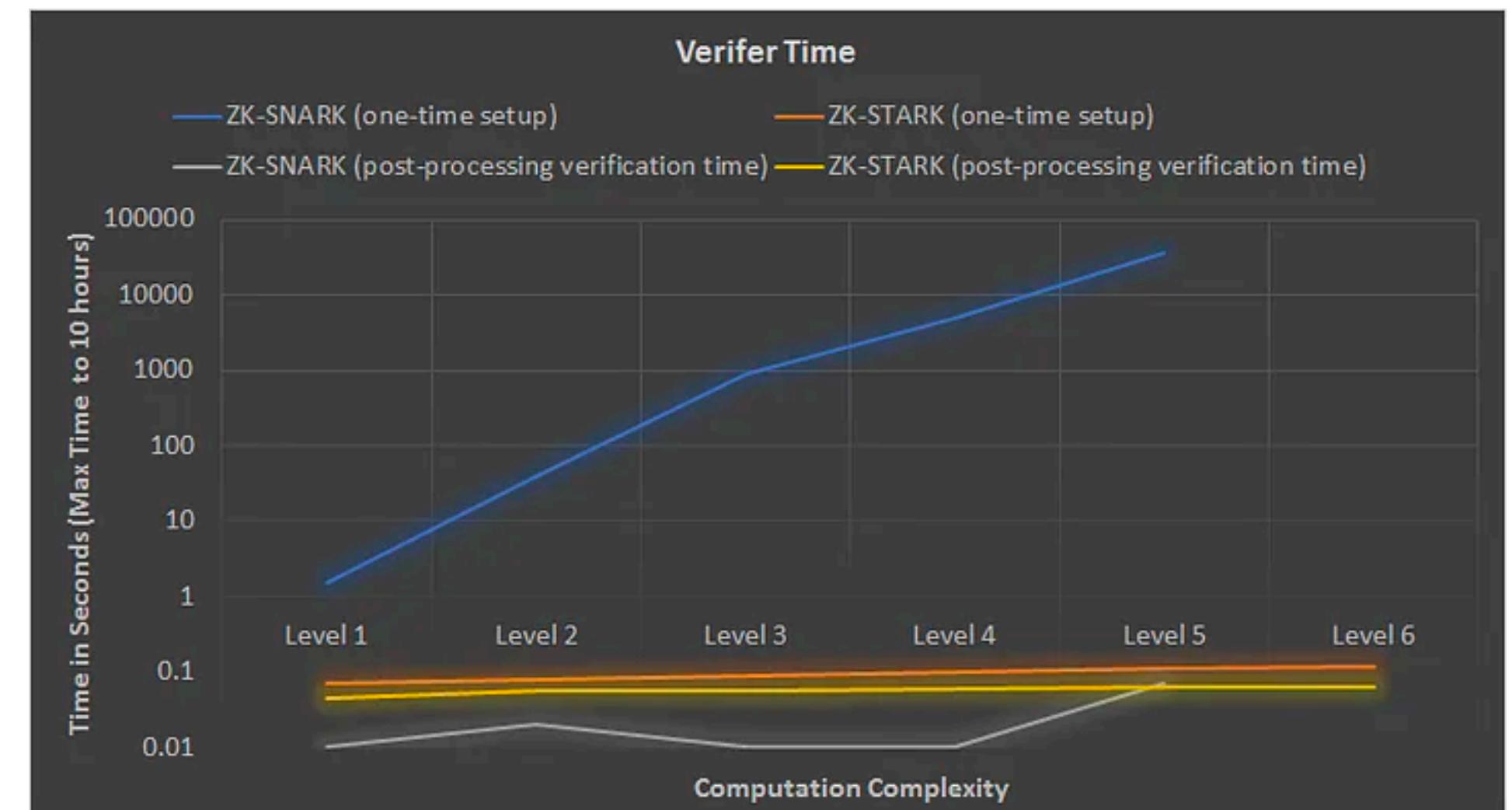
# STARK vs SNARK

## Pros

- **Quantum resistant**
- **Transparent (don't use trusted setup)**
- **Prover complexity:** ZK-STARKs are ~10x faster than ZK-SNARKs as computation size increases
- **Verifier complexity:** As computation size grows, ZK-STARKs grow only slightly vs. ZK-SNARKS

## Cons

- **Proof size**





# Stark 101: Part 1

**Statement, LDE and Commitment**

# FibonacciSq

## (Fibonacci Square)

# FibonacciSq (Fibonacci Square)

FibonacciSq:

$$a_{n+2} = a_{n+1}^2 + a_n^2$$

- Represented as:  $a_0, a_1, a_2, a_3, \dots$
- Determined by first two elements
- Example:
  - 1, 3, 10, 109, 11981, 143556242,...

# Tiny Problem



$a_{10} = 10585384481491331545443435980195330168085$   
227108560824098919278258215839789697544114437  
130080556524289168854586579782387518129922282  
832261605608145523797747714827465842570005148  
785265883367108772402086618503369319342561663  
36593387070293738452872952783090264176685

# FibonacciSq Mod Prime

FibonacciSq mod prime:  $a_{n+2} = a_{n+1}^2 + a_n^2 \pmod{\text{prime}}$

Example:

- 1, 3, 10, 109, 11981, 143556242,...

mod 7:

- 1, 3, 3, 4, 4, 4, ...

# FibonacciSq Mod Prime

FibonacciSq mod prime:  $a_{n+2} = a_{n+1}^2 + a_n^2 \pmod{\text{prime}}$

- Example - mod 7:
  - 1, 3, 3, 4, 4, 4, ...

We use  $\text{prime} = 3 \cdot 2^{30} + 1 = 322122547$

Finite field  $F$

# Statement

# Statement to Prove

There is a number  $x$  such that:

For the FibonacciSq mod 3221225473 with

- $a_0 = 1$
- $a_1 = x$

we have  $a_{1022} = 2338775057$

X = 3141592



# STARK Protocol - Part I

- LDE - Low Degree Extension
- Commitment

# Low Degree Extension (LDE)

# LDE in 3 Steps

1. Generate input
2. Interpolate
3. Extend

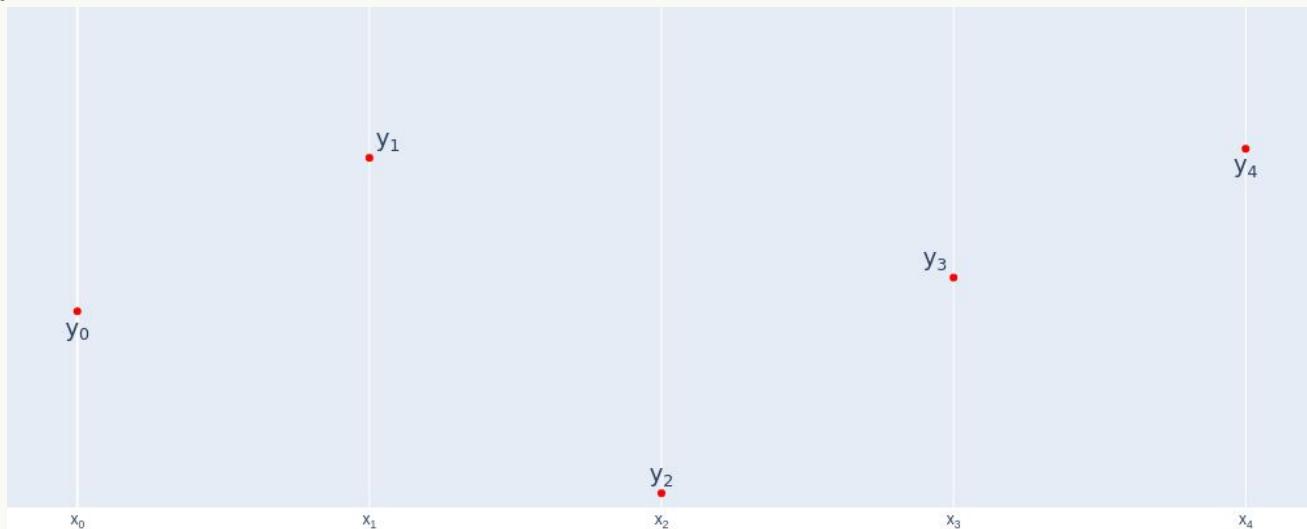
# LDE - General

# LDE Step 1 - Generate Input

**Input:**  $y_0, y_1, y_2, y_3, y_4, \dots$

**Choose:**  $x_0, x_1, x_2, x_3, x_4, \dots$

| $x$   | $y$   |
|-------|-------|
| $x_0$ | $y_0$ |
| $x_1$ | $y_1$ |
| $x_2$ | $y_2$ |
| $x_3$ | $y_3$ |
| $x_4$ | $y_4$ |

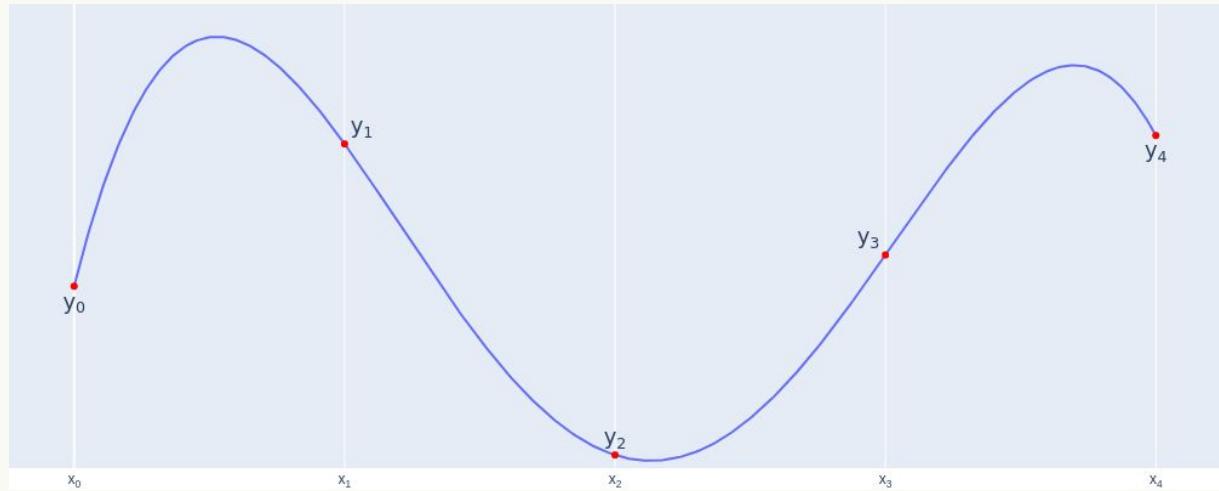


# LDE Step 2 - Interpolate Polynomial

Interpolate a polynomial  $f$ :

For each  $i : f(x_i) = y_i$

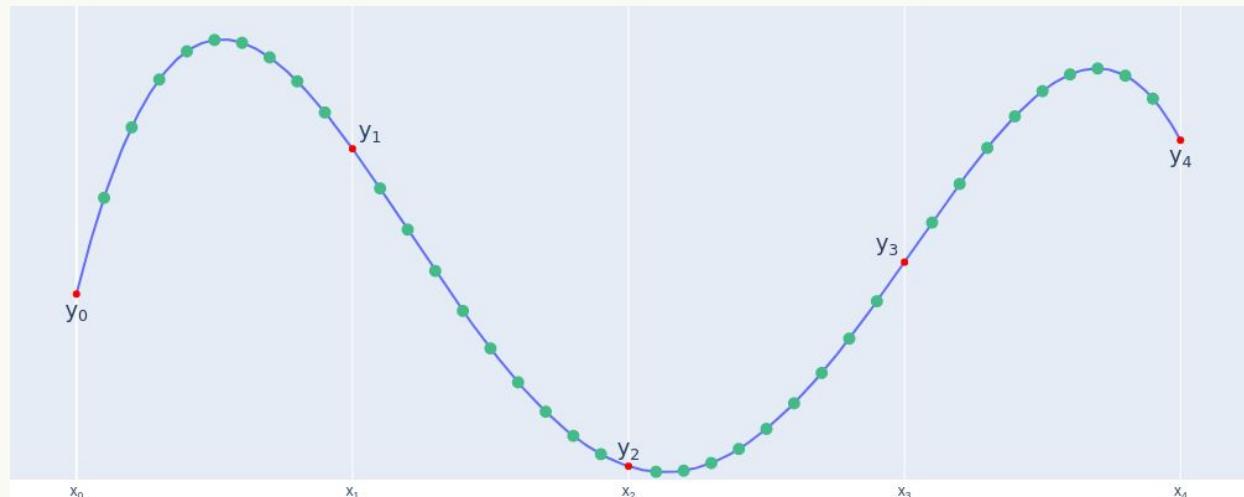
| $x$   | $f(x)$ |
|-------|--------|
| $x_0$ | $y_0$  |
| $x_1$ | $y_1$  |
| $x_2$ | $y_2$  |
| $x_3$ | $y_3$  |
| $x_4$ | $y_4$  |



# LDE Step 3 - Extend

- Pick a larger evaluation domain  $\{x_j\}$
- Output:  $\{f(x_j)\}$

| $x'$   | $f(x')$   |
|--------|-----------|
| $x'_0$ | $f(x'_0)$ |
| $x'_1$ | $f(x'_1)$ |
| $x'_2$ | $f(x'_2)$ |
| $x'_3$ | $f(x'_3)$ |
| ...    | ...       |



# LDE in STARK

# LDE for STARK Step 1 - Generate Input

**Input:**  $a_0, a_1, a_2, \dots, a_{1022}$

The Trace

We choose:  $1, g, g^2, g^3, \dots, g^{1022}$

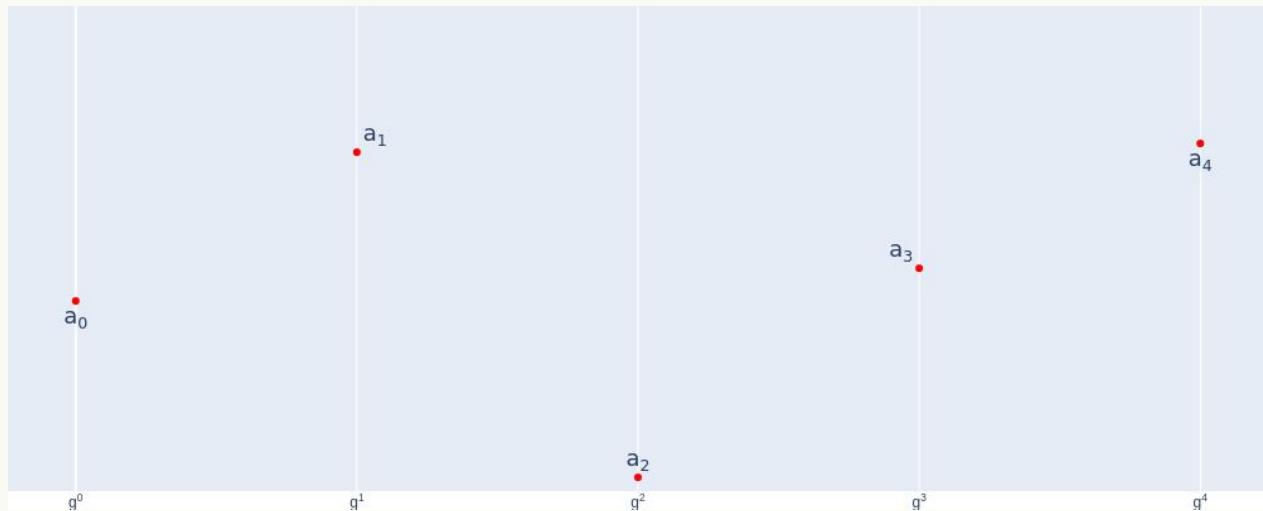
$g$  - element from  $F$

# LDE for STARK Step 1 - Generate Input

**Input:**  $a_0, a_1, a_2, \dots, a_{1022}$

**We choose:**  $1, g, g^2, g^3, \dots, g^{1022}$

| $x$        | $f(x)$     |
|------------|------------|
| $g^0$      | $a_0$      |
| $g^1$      | $a_1$      |
| $g^2$      | $a_2$      |
| ...        | ...        |
| $g^{1022}$ | $a_{1022}$ |

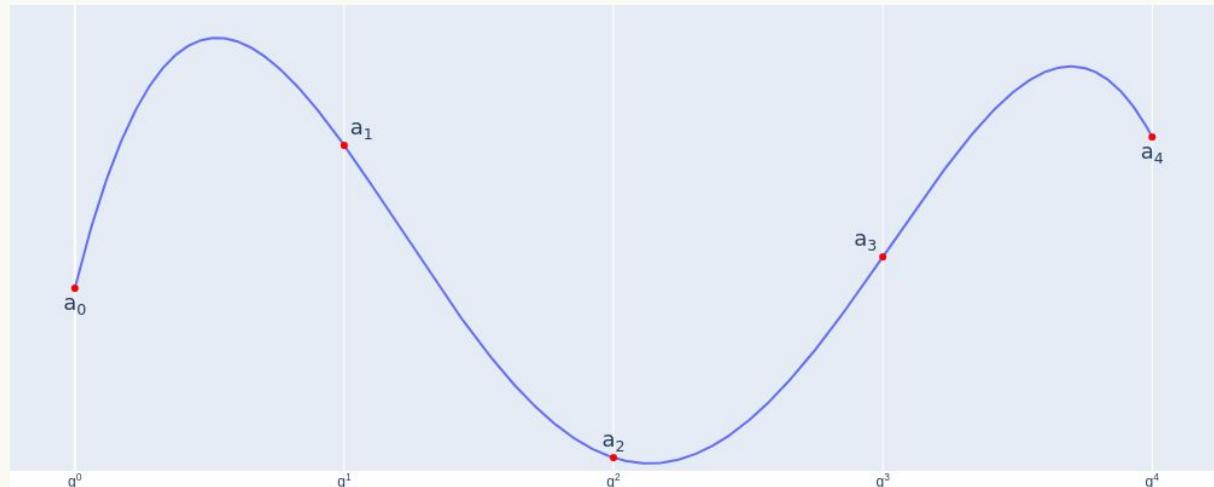


# LDE for STARK Step 2 - Interpolate Poly

Interpolate a polynomial  $f$ :

for each  $i : f(g^i) = a_i$

| $x$        | $f(x)$     |
|------------|------------|
| $g^0$      | $a_0$      |
| $g^1$      | $a_1$      |
| $g^2$      | $a_2$      |
| ...        | ...        |
| $g^{1022}$ | $a_{1022}$ |



# LDE for STARK Step 3 - Extend

- Pick a larger evaluation domain (8k)
- $\{x_i\} = w, w \cdot h, w \cdot h^2, \dots, w \cdot h^{8191}$

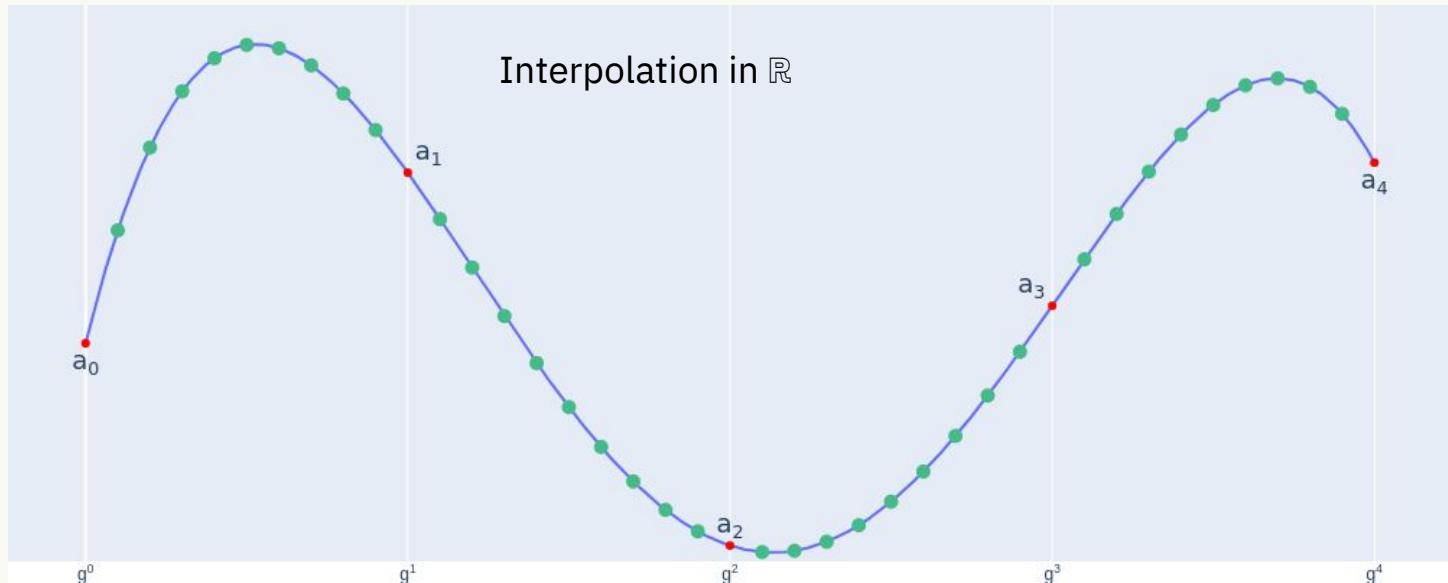
$w, h$  - elements from  $F$

- Result:  $f(w), f(w \cdot h), f(w \cdot h^2), \dots$

Reed-Solomon  
codeword

# LDE for STARK Step 3 - Extend

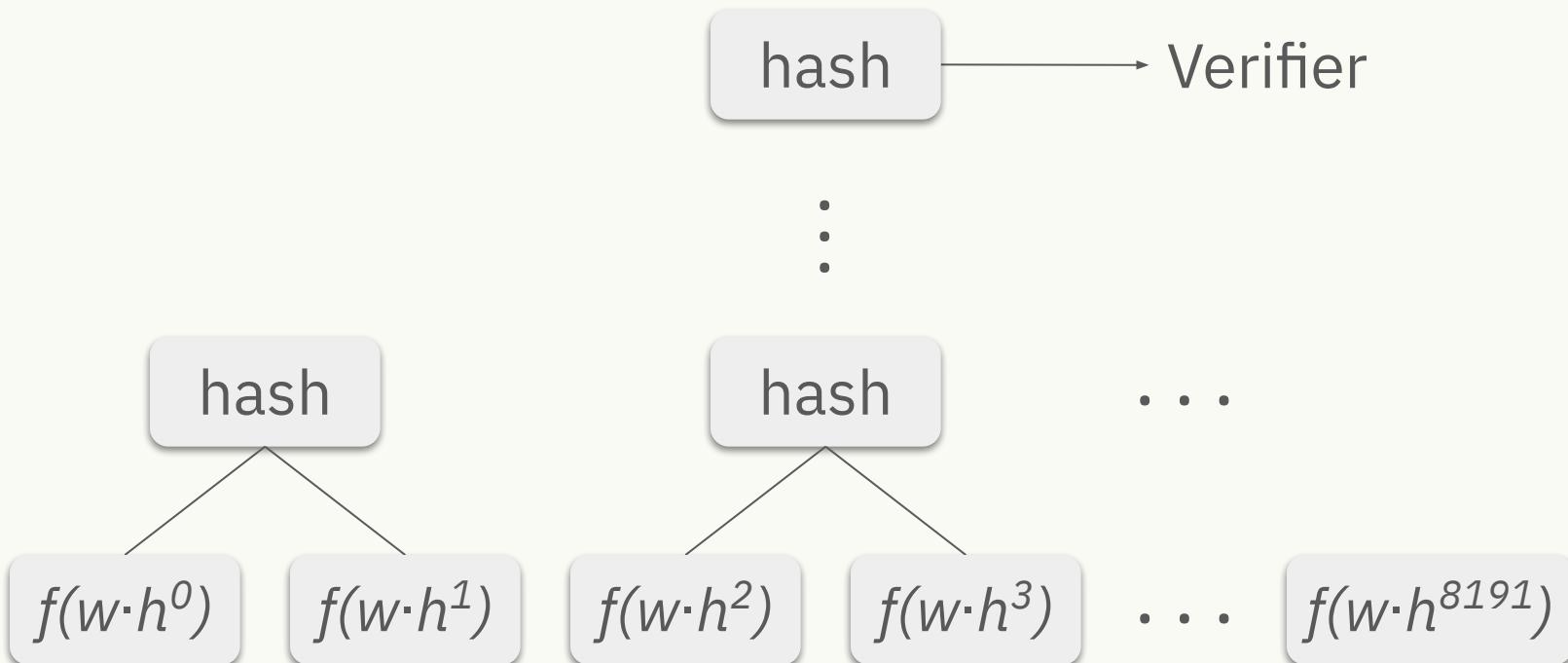
| $x$                | $f(x)$                |
|--------------------|-----------------------|
| $w \cdot h^0$      | $f(w \cdot h^0)$      |
| $w \cdot h^1$      | $f(w \cdot h^1)$      |
| $w \cdot h^2$      | $f(w \cdot h^2)$      |
| ...                | ...                   |
| $w \cdot h^{8191}$ | $f(w \cdot h^{8191})$ |



# Commitment

# Commit on LDE

Merkle Tree



# Summary

- Statement
  - There is  $x$  s.t.  $a_{1022} = 2338775057$  in FibonacciSq mod prime
- STARK protocol - part I:
  - LDE - Low Degree Extension
  - Commitment - Merkle Tree



STARKWARE **STARK** 101

# Stark 101: Part 2

**Polynomial Constraints**

# What Do We Want to Prove?

There is a number  $x$  such that:

$$a_0 = 1$$

$$a_1 = x$$

$$a_{1022} = 2338775057$$

We will use part I:

Trace -  $a$

Generator of  $G$  -  $g$

Trace Polynomial -  $f(x)$

For  $\{a_n\}$  FibonacciSq:  $a_{n+2} = a_{n+1}^2 + a_n^2 \pmod{\text{prime}}$ , for any  $n$

# Constraints on $\{a_n\}$

We need:

$$a_0 = 1$$

$$a_{1022} = 2338775057$$

$$a_{n+2} = a_{n+1}^2 + a_n^2$$

If  $\{a_n\}$  satisfies constraints  $\longrightarrow$  Original statement is true!

# Where are We Heading?

Constraints on  $\{a_n\}$ :

$$a_0 = 1$$

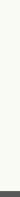
$$a_{1022} = 2338775057$$

$$a_{n+2} = a_{n+1}^2 + a_n^2$$

Reductions

Another statement

Implies



# Where are We Heading?

Constraints on  $\{a_n\}$ :

$$a_0 = 1$$

$$a_{1022} = 2338775057$$

$$a_{n+2} = a_{n+1}^2 + a_n^2$$

Reductions

Exists a polynomial  $f(x)$

such that:

3 rational functions

$p_0(x), p_1(x), p_2(x)$  are **polynomials**

Trace  
polynomial

$$\frac{a(x)}{b(x)}$$

# Step I - From $\{a_n\}$ to $f(x)$

Trace  
polynomial

3 constraints on  $\{a_n\}$    -----> 3 constraints on  $f(x)$

$a_0 = 1$    ----->  $f(x) = 1$ , for  $x = g^0$

$a_{1022} = 2338775057$    ----->  $f(x) = 2338775057$ , for  $x = g^{1022}$

$$a_{n+2} = a_{n+1}^2 + a_n^2$$

| $x$        | $f(x)$     |
|------------|------------|
| $g^0$      | $a_0$      |
| $g^1$      | $a_1$      |
| $g^2$      | $a_2$      |
| ...        | ...        |
| $g^{1022}$ | $a_{1022}$ |

# Step I - From $\{a_n\}$ to $f(x)$

$$a_{n+2} = a_{n+1}^2 + a_n^2 \quad \dashrightarrow \quad f(g^2x) = f(gx)^2 + f(x)^2,$$

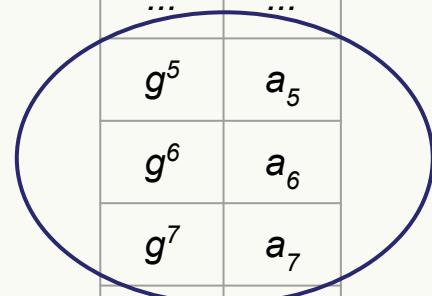
for  $x = g^i$ ,  $0 \leq i \leq 1020$

Example: for  $x = g^5$ :

$$f(g^2 \cdot g^5) = f(g \cdot g^5)^2 + f(g^5)^2$$

$\underbrace{g^7}_{g^7}$        $\underbrace{g^6}_{g^6}$        $\underbrace{g^5}_{g^5}$

| $x$        | $f(x)$     |
|------------|------------|
| $g^0$      | $a_0$      |
| $g^1$      | $a_1$      |
| $g^2$      | $a_2$      |
| ...        | ...        |
| $g^5$      | $a_5$      |
| $g^6$      | $a_6$      |
| $g^7$      | $a_7$      |
| ...        | ...        |
| $g^{1022}$ | $a_{1022}$ |



# Step I - From $\{a_n\}$ to $f(x)$

3 constraints on  $\{a_n\}$    ----->   3 constraints on  $f(x)$

$$a_0 = 1 \quad \text{----->} \quad f(x) = 1, \text{ for } x = g^0$$

$$a_{1022} = 2338775057 \quad \text{----->} \quad f(x) = 2338775057, \text{ for } x = g^{1022}$$

$$a_{n+2} = a_{n+1}^2 + a_n^2 \quad \text{----->} \quad f(g^2x) = f(gx)^2 + f(x)^2, \\ \text{for } x = g^i, 0 \leq i \leq 1020$$

# Step I - From $\{a_n\}$ to $f(x)$

3 constraints on  $\{a_n\}$     →    3 constraints on  $f(x)$

$a_0 = 1$     →     $f(x) = 1$ , for  $x = g^0$

$a_{1022} = 2338775057$     →     $f(x) = 2338775057$ , for  $x = g^{1022}$

$a_{n+2} = a_{n+1}^2 + a_n^2$     →     $f(g^2x) = f(gx)^2 + f(x)^2$ ,  
for  $x = g^i$ ,  $0 \leq i \leq 1020$

If  $f(x)$  satisfies constraints → Original statement is true

## Step II - From Constraints to Roots

$f(x) - 1 = 0$ , for  $x = g^0$     - - - -  $\Rightarrow$  root:  $g^0$

(  $f(x) = 1$ , for  $x = g^0$  )

$z$  is a root of  
 $p(x)$  if  $p(z)=0$

## Step II - From Constraints to Roots

$f(x) - 1 = 0$ , for  $x = g^0$     - - - -  $\Rightarrow$  root:  $g^0$

$f(x) - 2338775057 = 0$ , for  $x = g^{1022}$     - - - -  $\Rightarrow$  root:  $g^{1022}$

## Step II - From Constraints to Roots

$f(x) - 1 = 0$ , for  $x = g^0$   $\dashrightarrow$  root:  $g^0$

$f(x) - 2338775057 = 0$ , for  $x = g^{1022}$   $\dashrightarrow$  root:  $g^{1022}$

$f(g^2x) - f(gx)^2 - f(x)^2 = 0$ , for  $x = g^i$ ,  $0 \leq i \leq 1020$   $\dashrightarrow$  roots:  $\{g^i | 0 \leq i \leq 1020\}$

## Step II - From Constraints to Roots

$f(x) - 1 = 0$ , for  $x = g^0$   $\dashrightarrow$  root:  $g^0$

$f(x) - 2338775057 = 0$ , for  $x = g^{1022}$   $\dashrightarrow$  root:  $g^{1022}$

$f(g^2x) - f(gx)^2 - f(x)^2 = 0$ , for  $x = g^i$ ,  $0 \leq i \leq 1020$   $\dashrightarrow$  roots:  $\{g^i | 0 \leq i \leq 1020\}$

$g^0$  is a root of  $f(x) - 1$

$g^{1022}$  is a root of  $f(x) - 2338775057$

$\{g^i | 0 \leq i \leq 1020\}$  are roots of  $f(g^2x) - f(gx)^2 - f(x)^2$

Original statement is true

# Step III - From Roots to Rational Functions

Thm:  $z$  is a root of  $p(x) \Leftrightarrow (x - z)$  divides  $p(x)$

Def:  $(x - z)$  divides  $p(x)$  if  $p(x) / (x - z)$  is a polynomial

## Polynomial

$$\frac{x^2 - 3x + 2}{x - 2} = \frac{(x - 2)(x - 1)}{x - 2} = x - 1$$

**2 is a root**

## Not polynomial

$$\frac{x^2 - 7x + 6}{x - 2} = \frac{(x - 1)(x - 6)}{x - 2}$$

**2 is NOT a root**

## Step III - From Roots to Rational Functions

Thm:  $z$  is a root of  $p(x) \Leftrightarrow (x - z)$  divides  $p(x)$

Def:  $(x - z)$  divides  $p(x)$  if  $p(x) / (x - z)$  is a polynomial

$g^0$  is a root of  $f(x) - 1 \dashrightarrow \frac{f(x) - 1}{x - g^0}$  is a polynomial

$g^{1022}$  is a root of  $f(x) - 2338775057 \dashrightarrow \frac{f(x) - 2338775057}{x - g^{1022}}$   
is a polynomial

## Step III - From Roots to Rational Functions

$\{g^i \mid 0 \leq i \leq 1020\}$  are roots of  $f(g^2x) - f(gx)^2 - f(x)^2$  →

$$\frac{f(g^2x) - f(gx)^2 - f(x)^2}{\prod_{i=0}^{1020} (x - g^i)}$$

is a polynomial

1023

$$\prod_{i=0}^{1023} (x - g^i) = x^{1024} - 1 \quad \text{fix:}$$

$$\frac{f(g^2x) - f(gx)^2 - f(x)^2}{(x^{1024} - 1)/[(x - g^{1021})(x - g^{1022})(x - g^{1023})]}$$

# 3 Rational Functions

$$p_0(x) = \frac{f(x) - 1}{x - g^0}$$

$$p_1(x) = \frac{f(x) - 2338775057}{x - g^{1022}}$$

$$p_2(x) = \frac{f(g^2 x) - f(gx)^2 - f(x)^2}{(x^{1024} - 1)/[(x - g^{1021})(x - g^{1022})(x - g^{1023})]}$$

If  $p_0(x)$ ,  $p_1(x)$ ,  $p_2(x)$  are polynomials → Original statement is true!

# Where are We Heading?

Constraints on  $\{a_n\}$ :

$$a_0 = 1$$

$$a_{1022} = 2338775057$$

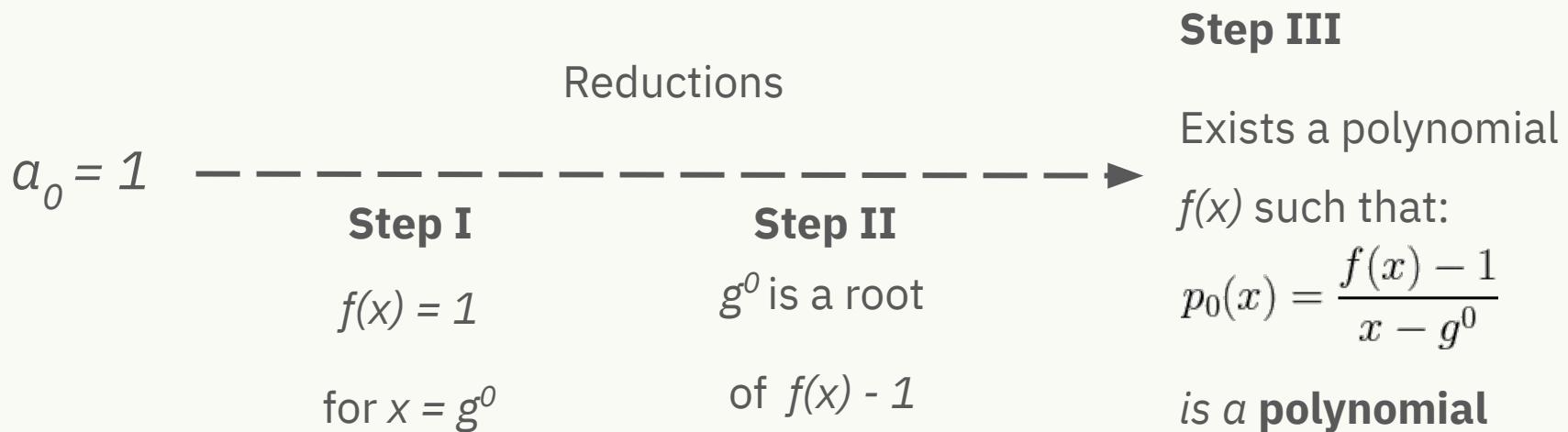
$$a_{n+2} = a_{n+1}^2 + a_n^2$$

Reductions



Exists a polynomial  $f(x)$   
such that:  
3 rational functions  
 $p_0(x), p_1(x), p_2(x)$  are **polynomials**

# Reduction Overview - First Constraint



# 3 Rational Functions

$$p_0(x) = \frac{f(x) - 1}{x - g^0}$$

$$p_1(x) = \frac{f(x) - 2338775057}{x - g^{1022}}$$

$$p_2(x) = \frac{f(g^2 x) - f(gx)^2 - f(x)^2}{(x^{1024} - 1)/[(x - g^{1021})(x - g^{1022})(x - g^{1023})]}$$

If  $p_0(x)$ ,  $p_1(x)$ ,  $p_2(x)$  are polynomials → Original statement is true!

# Combining $p_i(x)$ 's

Random linear combination:

Composition  
Polynomial

$$CP = \alpha_0 \cdot p_0(x) + \alpha_1 \cdot p_1(x) + \alpha_2 \cdot p_2(x)$$

With high probability:

$$CP \text{ is a polynomial} \Leftrightarrow \text{all } p_i \text{'s are polynomials}$$

Committing on  $CP$  with Merkle Tree



# Stark 101: Part 3

**FRI Commitment**

# Recap

**Goal : prove a statement on FibonacciSq**

- Trace in 1023 points
- Create *Trace* polynomial (Lagrange interpolation)
- Evaluate and commit on a larger domain

# Recap

- 3 constraints on  $f(x)$ :

$$f(x) - 1 = 0 \text{ , for } x = 1$$

...

- 3 rational functions from the constraints:

$$p_0(x) = \frac{f(x) - 1}{x - g^0}$$

...

# Recap

- Composition Polynomial:

$$CP(x) = \alpha_0 \cdot p_0(x) + \alpha_1 \cdot p_1(x) + \alpha_2 \cdot p_2(x)$$

- Prover commits on CP
- Goal - show that CP is a **polynomial**
- CP is a **polynomial** → All constraints satisfied

# What Will We Do?

Goal:

Prove that CP is a **polynomial**

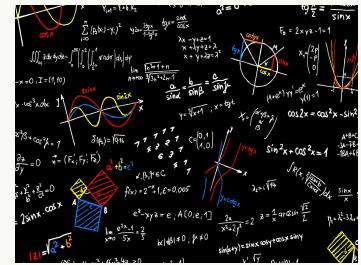


Instead:

Prove that CP is **close** to a **polynomial of low degree**

What is close?

What is low degree?

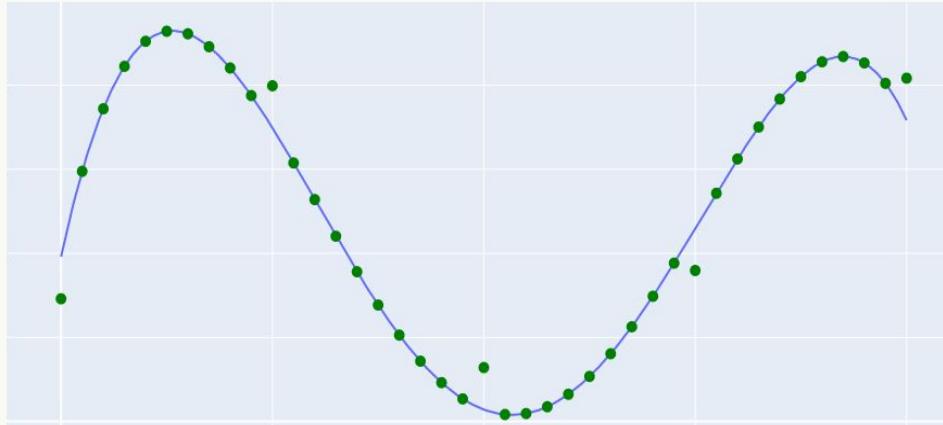


# Proximity to Polynomials

**Distance (def):**

Distance between a function  $f: D \rightarrow F$  to a polynomial  $p$ :

$$D(f, p) := \# \text{ points } x \in D \text{ such that } f(x) \neq p(x)$$



$$D(\mathbf{f}, \mathbf{p}) = 5$$

# Proximity to Polynomials

**Distance (def):**

Distance between a function  $f: D \rightarrow F$  to a polynomial  $p$ :

$$D(f,p) := \# \text{ points } x \in D \text{ such that } f(x) \neq p(x)$$

**Proximity**

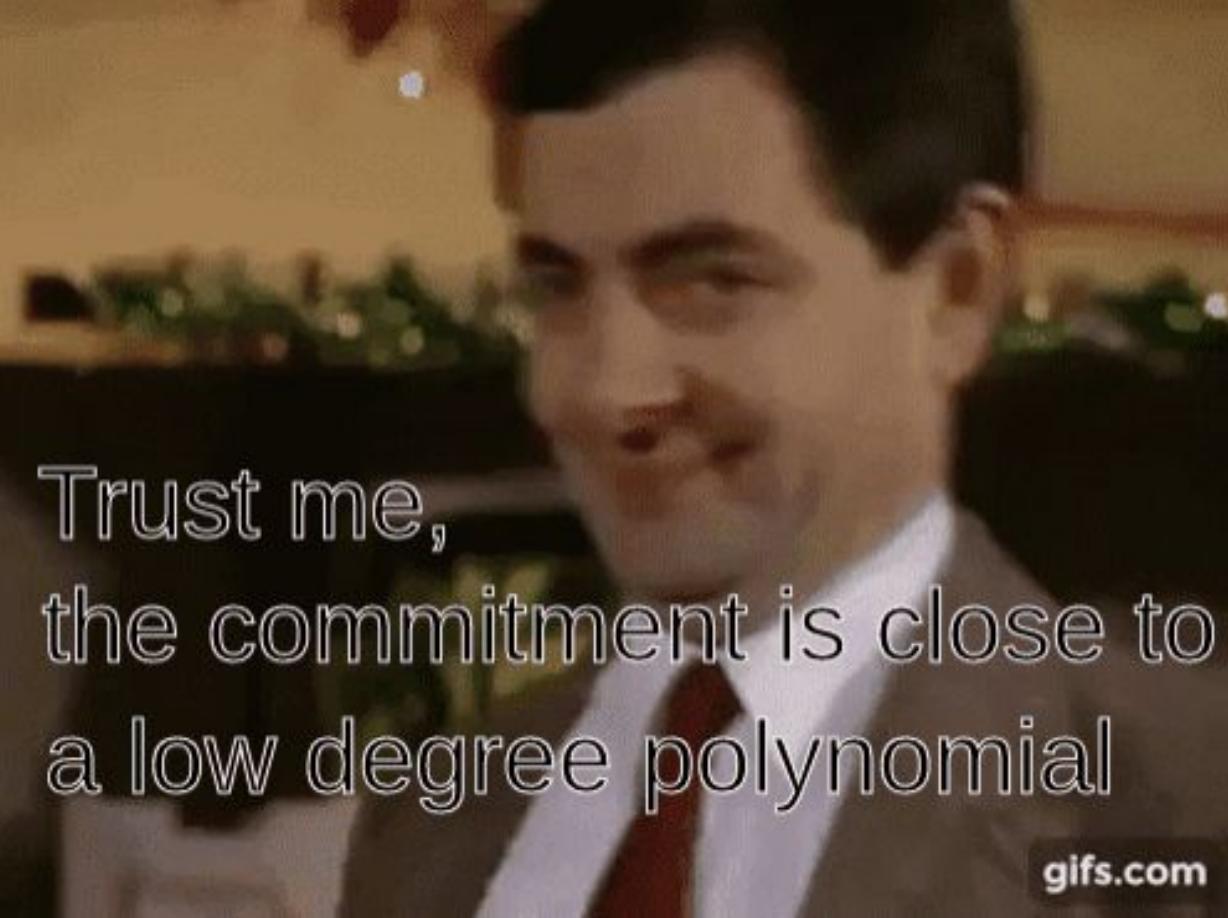
A function  $f: D \rightarrow F$  is **close** to a polynomial  $p$  if:  $D(f,p)$  is **small**

# What Will We Do? - Reminder

Goal:

Prove that CP is **close** to a **polynomial of low degree**

**How?**



Trust me,  
the commitment is close to  
a low degree polynomial

[gifs.com](http://gifs.com)

FRI

# Fast Reed-Solomon Interactive Oracle Proofs of Proximity

By Ben-Sasson, E., Bentov, I., Horesh, Y., & Riabzev, M.

<https://eccc.weizmann.ac.il/report/2017/134/>

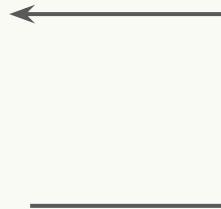
# FRI - Goal

Prover convinces verifier:

**“The commitment is close to a low degree polynomial”**

# FRI - The Protocol

- Receive random  $\beta$
- Apply the FRI operator
- Commit
- Lastly the prover sends the result



Do it repeatedly

# FRI

- FRI operator - motivation
- FRI steps overview
- Deep into the FRI operator

# FRI Operator

# FRI Operator

Goal:

Prove that a function is close to a polynomial of a degree  $< D$



Applying the FRI operator

New Goal:

Prove that a **new** function is close to a **new** polynomial



Half of the domain size



Degree  $< D/2$

# FRI Operator - Example

## Before applying FRI operator

- Prove:

A function is close to a polynomial of a degree < **1024**

where domain size = **8192**

# FRI Operator - Example

~~Before After applying FRI operator~~

- Prove:

A function is close to a polynomial of a degree < ~~1024~~ 512

where domain size = ~~8192~~ 4096



# FRI Steps Overview

# FRI Steps Overview

Showing that  $\deg(CP) < 1024$ ,  $|D| = 8192$

$CP(x)$

**CP Commitment**

$$\beta_0$$

$\deg(CP_1) < 512$ ,  $|D_1| = 4096$

$$\beta_1$$

$\deg(CP_2) < 256$ ,  $|D_2| = 2048$

.

.

.

$$\beta_9$$

**Constant**

$\deg(CP_{10}) < 1$ ,  $|D_{10}| = 8$

# FRI Steps Overview

Showing that  $\deg(CP) < 1024$ ,  $|D| = 8192$

$CP(x)$

**CP Commitment**

$\beta_0$

$\deg(CP_1) < 512$ ,  $|D_1| = 4096$

$\beta_1$

$\deg(CP_2) < 256$ ,  $|D_2| = 2048$

.

.

.

$\beta_9$

$\deg(CP_{10}) < 1$ ,  $|D_{10}| = 8$



# Deep Into the FRI Operator



# FRI Operator - How Does it Work?

- Split to even and odd powers

$$P_0(x) = g(x^2) + xh(x^2)$$

- Get a random  $\beta$

- Consider the new function:

$$P_1(y) = g(y) + \beta h(y)$$

- Example:

$$\begin{array}{r} P_0(x) = 5x^5 + 3x^4 + 7x^3 + 2x^2 + x + 3 \\ \downarrow \qquad \downarrow \qquad \downarrow \qquad \downarrow \qquad \downarrow \\ g(x^2) \qquad 3x^4 \qquad 2x^2 \qquad 3 \\ \downarrow \qquad \downarrow \qquad \downarrow \qquad \downarrow \\ xh(x^2) \qquad 5x^5 \qquad 7x^3 \qquad x \end{array}$$

# FRI Operator - How Does it Work?

- Split to even and odd powers

$$P_0(x) = \mathbf{g}(x^2) + x\mathbf{h}(x^2)$$

- Get a random  $\beta$

- Consider the new function:

$$P_1(y) = \mathbf{g}(y) + \beta\mathbf{h}(y)$$

- Example:

$$\begin{array}{ccccccc} P_0(x) & = & 5x^5 & + & 3x^4 & + & 7x^3 \\ & & \downarrow & & \downarrow & & \downarrow \\ g(x^2) & & 3x^4 & & 2x^2 & & 3 \\ & \downarrow & \downarrow & & \downarrow & & \downarrow \\ g(y) & & 3y^2 & & 2y & & 3 \\ & \downarrow & \downarrow & & \downarrow & & \downarrow \\ xh(x^2) & & 5x^5 & & 7x^3 & & x \end{array}$$

# FRI Operator - How Does it Work?

- Split to even and odd powers

$$P_0(x) = g(x^2) + xh(x^2)$$

- Get a random  $\beta$

- Consider the new function:

$$P_1(y) = g(y) + \beta h(y)$$

- Example:

$$P_0(x) = 5x^5 + 3x^4 + 7x^3 + 2x^2 + x + 3$$

$$\begin{array}{cccccc} & & & & & \\ & g(y) & & 3y^2 & & 2y & \\ & \downarrow & & \downarrow & & \downarrow & \\ xh(x^2) & 5x^5 & & 7x^3 & & x & \\ & & & & & & \end{array}$$

# FRI Operator - How Does it Work?

- Split to even and odd powers

$$P_0(x) = g(x^2) + xh(x^2)$$

- Get a random  $\beta$

- Consider the new function:

$$P_1(y) = g(y) + \beta h(y)$$

- Example:

$$P_0(x) = 5x^5 + 3x^4 + 7x^3 + 2x^2 + x + 3$$

$$\begin{array}{cccccc} & & & & & \\ & g(y) & & 3y^2 & & 2y & \\ \downarrow & & \downarrow & & \downarrow & & \downarrow \\ xh(x^2) & 5x^5 & & 7x^3 & & x & \\ & \downarrow & & \downarrow & & \downarrow & \\ h(y) & 5y^5 & & 7y^3 & & 1 & \end{array}$$

# FRI Operator - How Does it Work?

- Split to even and odd powers

$$P_0(x) = g(x^2) + xh(x^2)$$

- Get a random  $\beta$

- Consider the new function:

$$P_1(y) = g(y) + \beta h(y)$$

- Example:

$$P_0(x) = 5x^5 + 3x^4 + 7x^3 + 2x^2 + x + 3$$

$$\begin{array}{cccccc} & & & & & \\ & g(y) & & 3y^2 & & 2y & & 3 \\ \downarrow & & \downarrow & & \downarrow & & \downarrow & \\ h(y) & & 5y^2 & & 7y & & 1 & \end{array}$$

# FRI Operator - How Does it Work?

- Split to even and odd powers

$$P_0(x) = g(x^2) + xh(x^2)$$

- Get a random  $\beta$

- Consider the new function:

$$P_1(y) = g(y) + \beta h(y)$$

- Example:

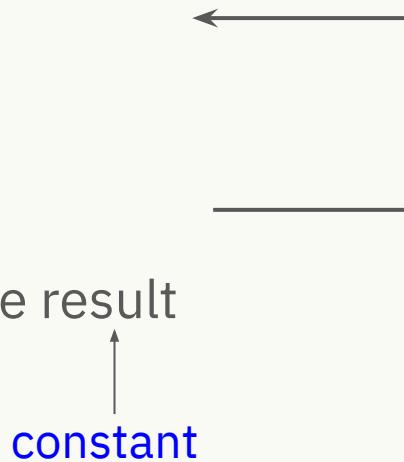
$$P_0(x) = 5x^5 + 3x^4 + 7x^3 + 2x^2 + x + 3$$

$$\begin{array}{cccc} g(y) & 3y^2 & 2y & 3 \\ \downarrow & \downarrow & \downarrow & \downarrow \\ h(y) & 5y^2 & 7y & 1 \end{array}$$

- $P_1(y) = 3y^2 + 2y + 3 + \beta(5y^2 + 7y + 1)$   
 $= (3+5\beta)y^2 + (2+7\beta)y + 3 + \beta$

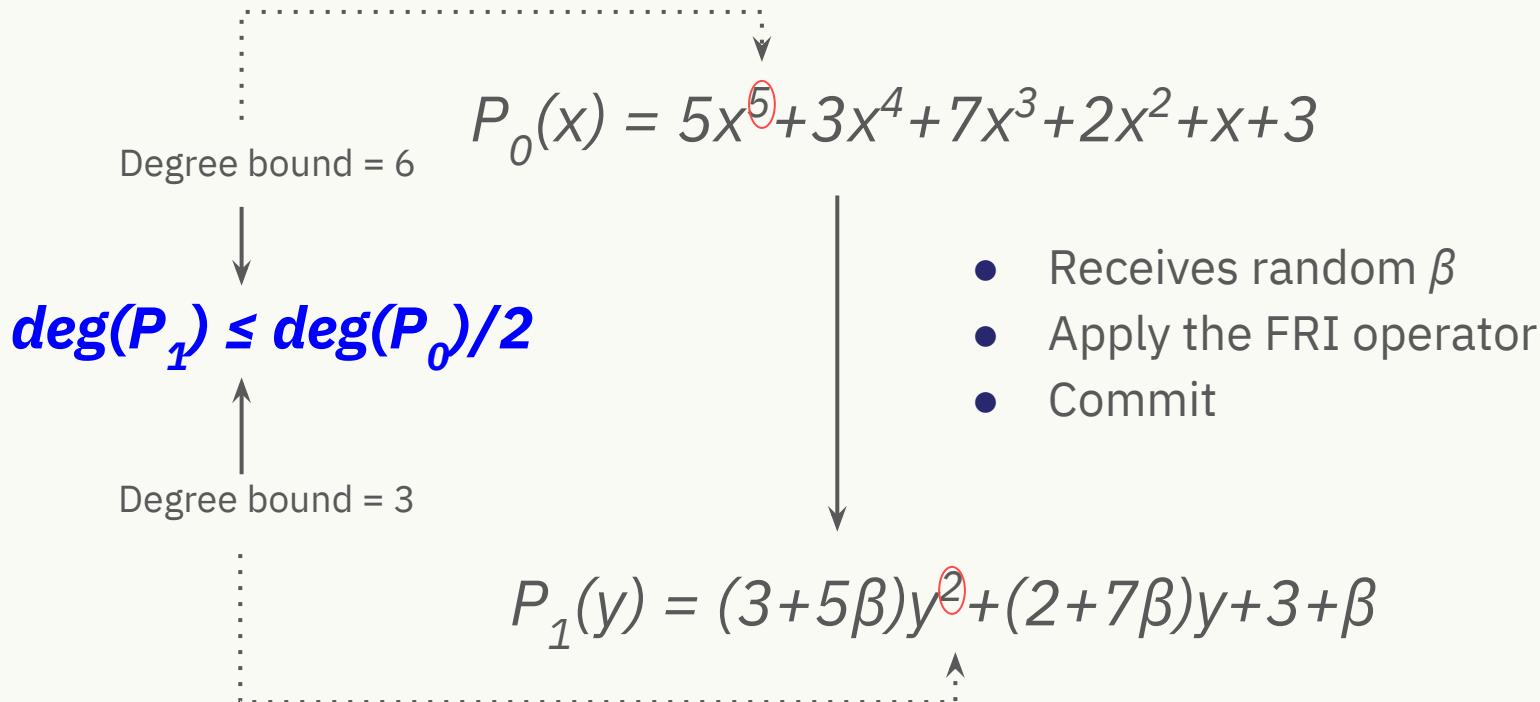
# FRI - The Protocol - Reminder

- Receive random  $\beta$
- Apply the FRI operator
- Commit
- Lastly the prover sends the result



Do it repeatedly  
↑  
 $\text{deg}(\text{poly}) < 1$   
where  
domain size is 8

# FRI - The Protocol - A Single Step





# Stark 101: Part 4

The Proof

# Recap

Trace

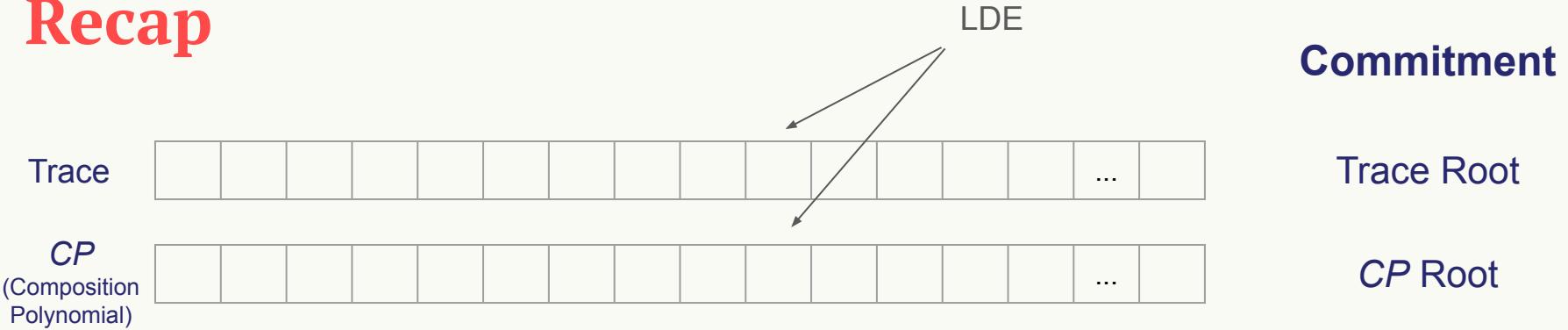


LDE

Commitment

Trace Root

# Recap



# Recap

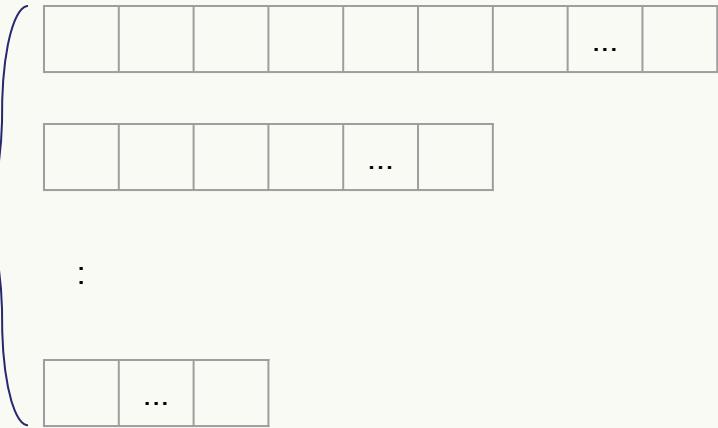
Trace



$CP$   
(Composition  
Polynomial)



FRI



Commitment

Trace Root

$CP$  Root

$CP_1$  Root

$CP_2$  Root

:

$CP_{10}$  Root

# The Entire Proof



Commitment

- Decommitment (Convincing)

# How Can the Prover Convince the Verifier?



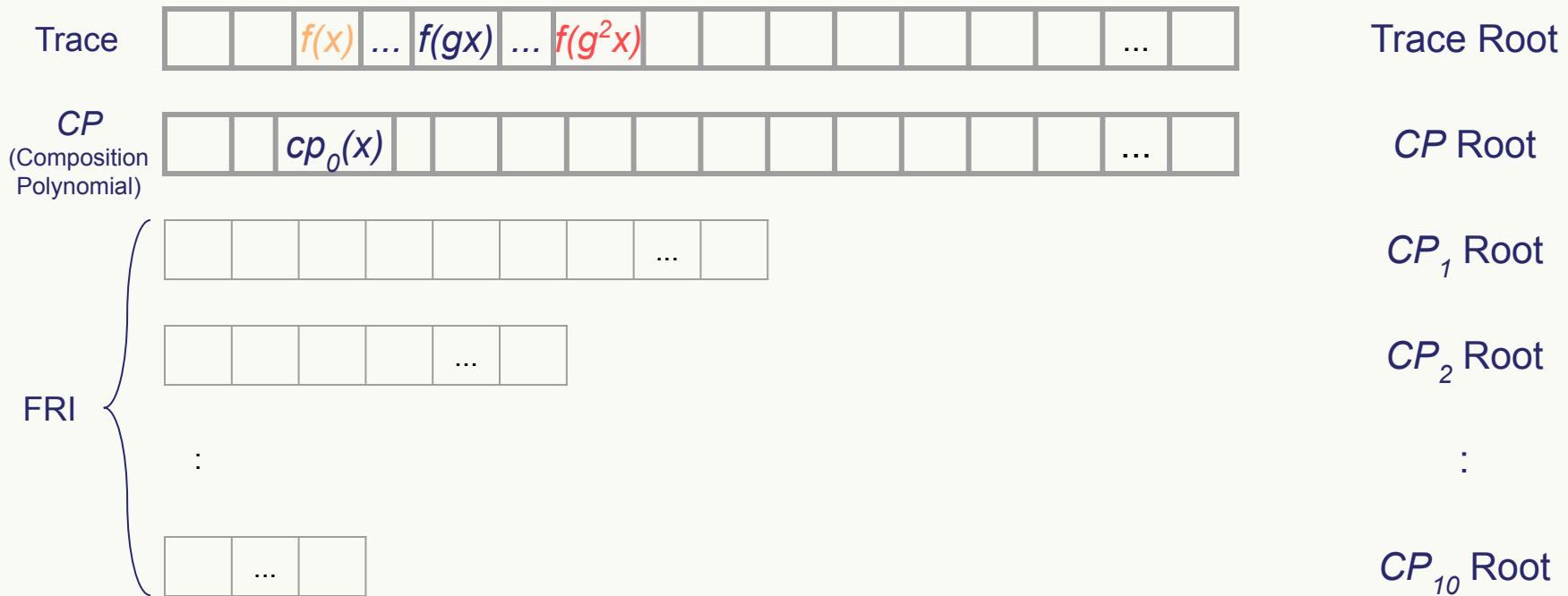
# How Can the Prover Convince the Verifier?

Verifier: Sends  $q$  random elements

Prover: Provides a validation data for each

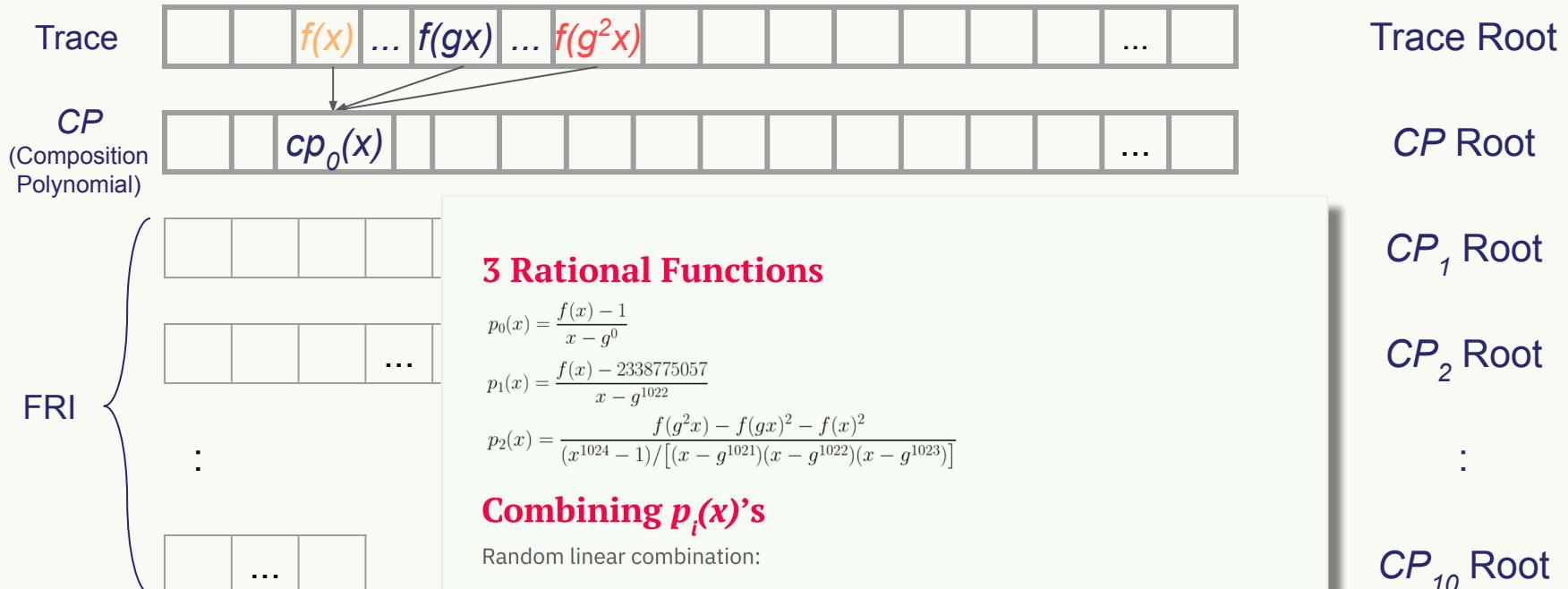
## What is a validation data?

# Trace $\rightarrow CP$



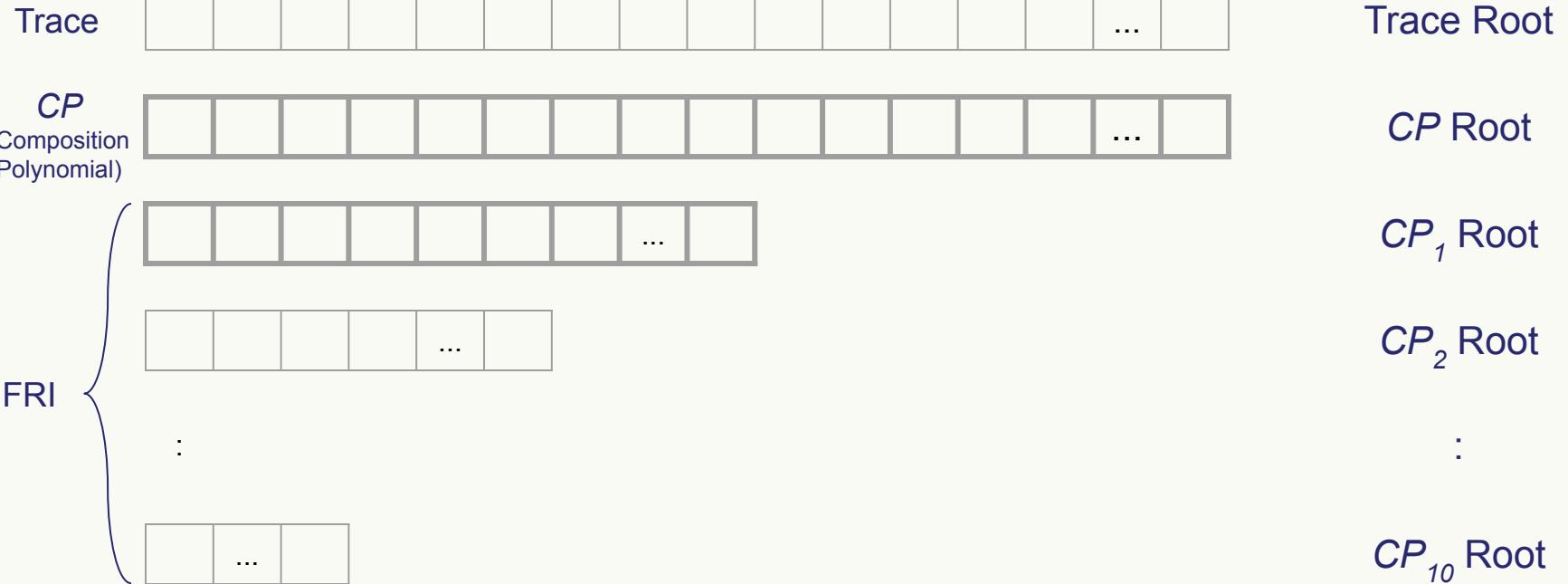
# Trace -> CP

Commitment



# FRI Step

Commitment



# FRI Step

$$\begin{cases} CP_i(x) = g(x^2) + x h(x^2) \\ CP_i(-x) = g(x^2) - x h(x^2) \end{cases} \quad \rightarrow \quad \begin{cases} g(x^2) = \frac{CP_i(x) - CP_i(-x)}{2} \\ h(x^2) = \frac{CP_i(x) - CP_i(-x)}{2x} \end{cases}$$

Reminder:

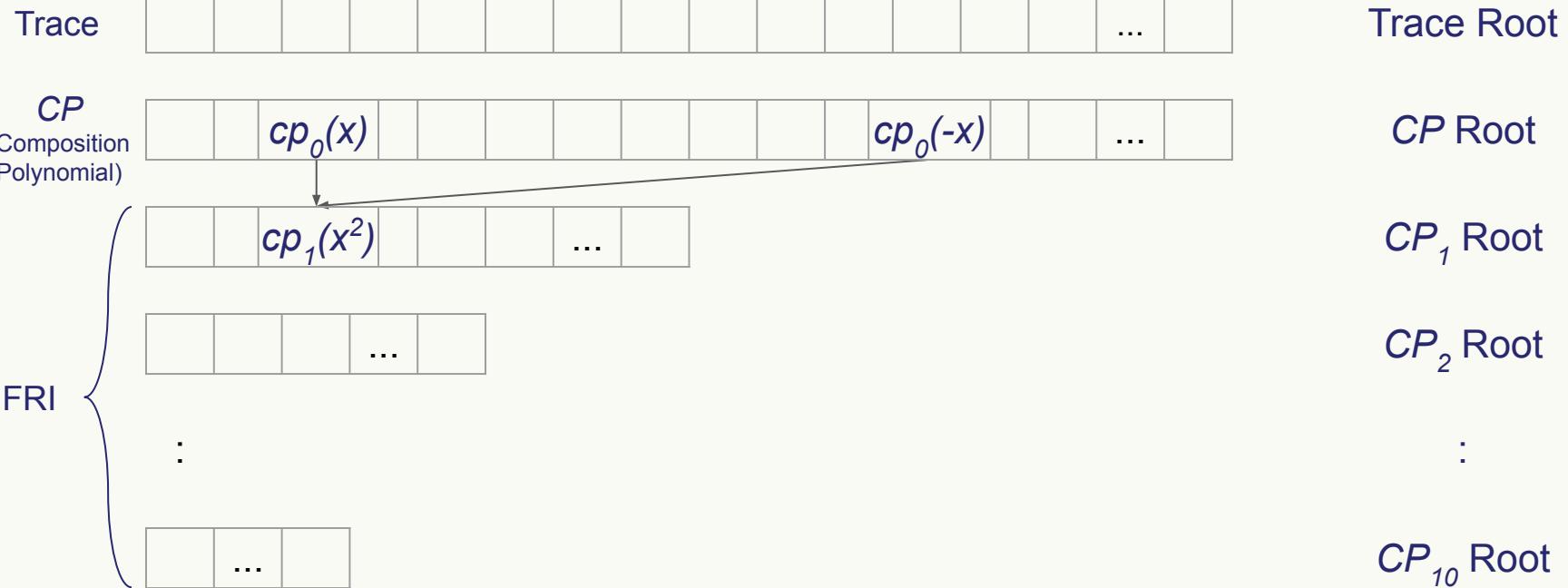
$$CP_{i+1}(x^2) = g(x^2) + \beta h(x^2)$$

Bottom line:

To compute  $CP_{i+1}(x^2)$  we only need  $CP_i(x)$  and  $CP_i(-x)$

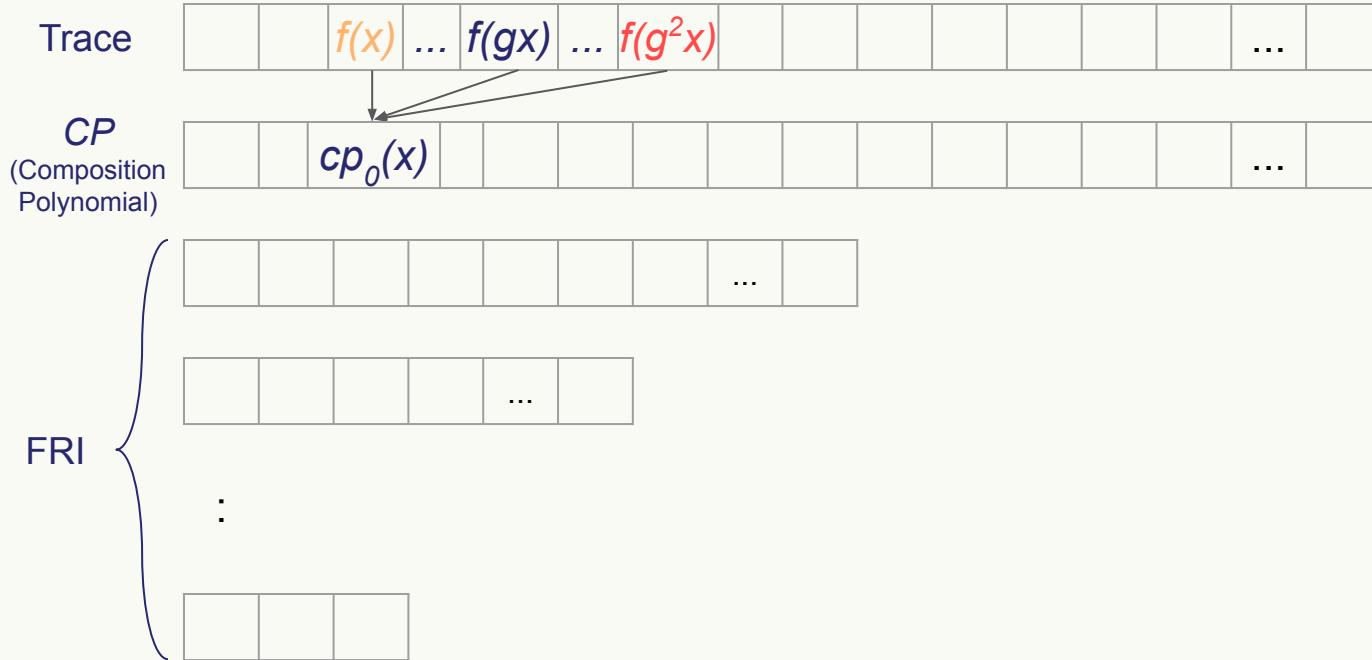
# FRI Step

Commitment



# Decommitment Phase (for query $x$ )

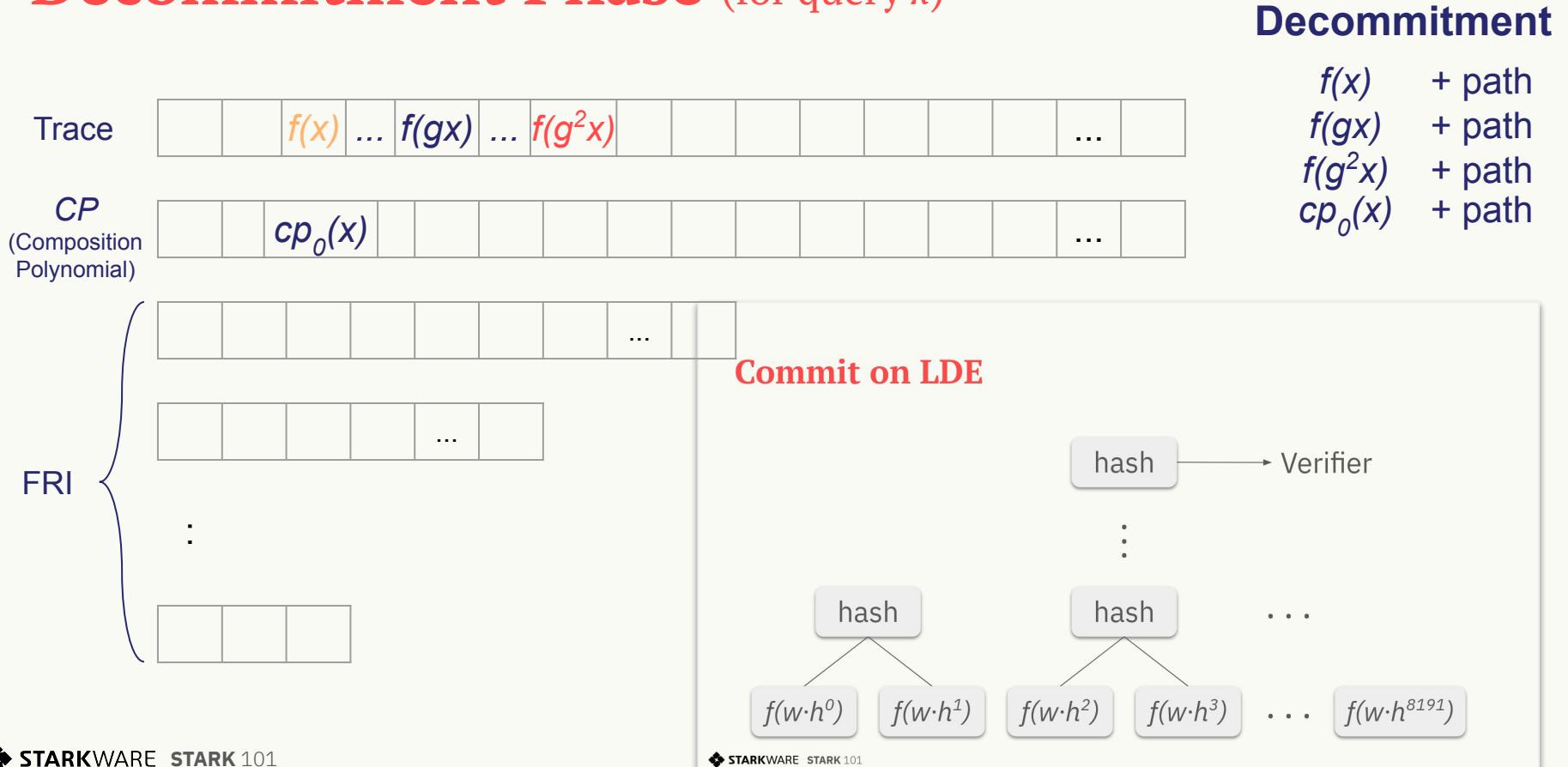
# Decommitment Phase (for query $x$ )



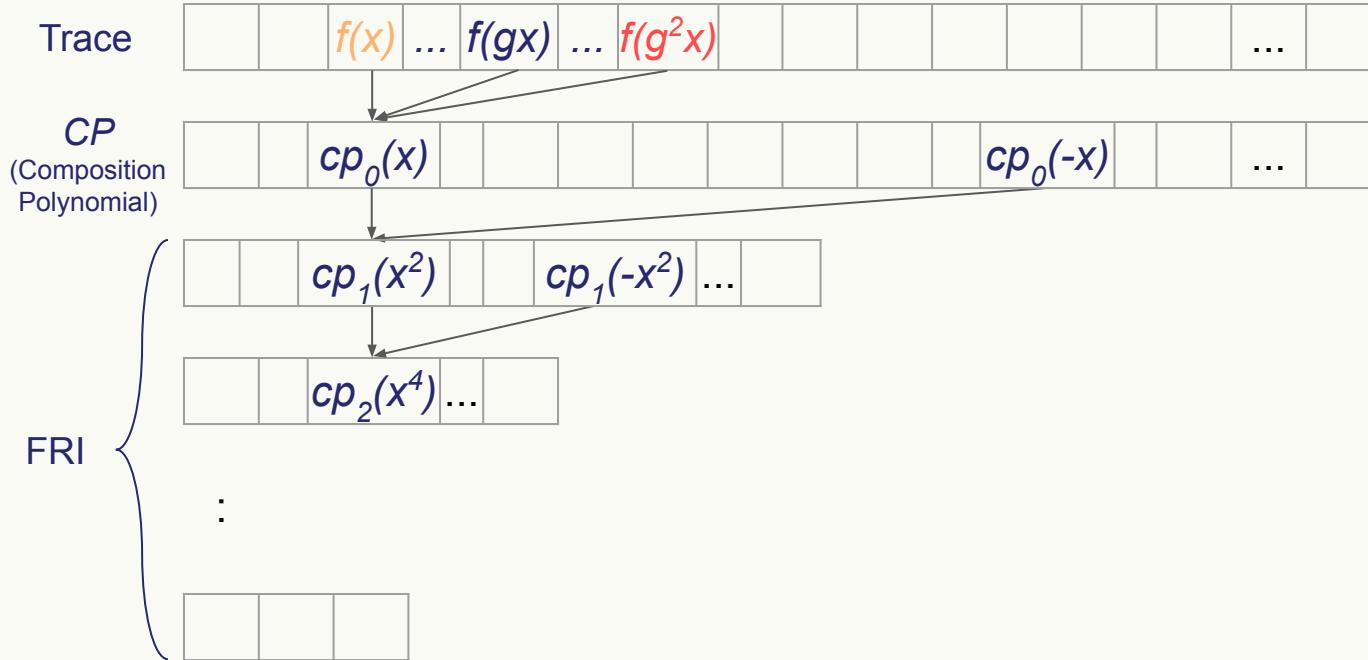
## Decommitment

|           |        |
|-----------|--------|
| $f(x)$    | + path |
| $f(gx)$   | + path |
| $f(g^2x)$ | + path |
| $cp_0(x)$ | + path |

# Decommitment Phase (for query $x$ )



# Decommitment Phase (for query $x$ )



## Decommitment

|              |          |
|--------------|----------|
| $f(x)$       | + path   |
| $f(gx)$      | + path   |
| $f(g^2x)$    | + path   |
| $cp_0(x)$    | + path   |
| $cp_0(-x)$   | + path   |
| $cp_1(x^2)$  | + path   |
| $cp_1(-x^2)$ | + path   |
| $cp_2(x^4)$  | + path   |
| $cp_2(-x^4)$ | + path   |
| $\vdots$     | $\vdots$ |

$cp_{10}(x^{1024})$  + path

# The Entire Proof

 Commitment

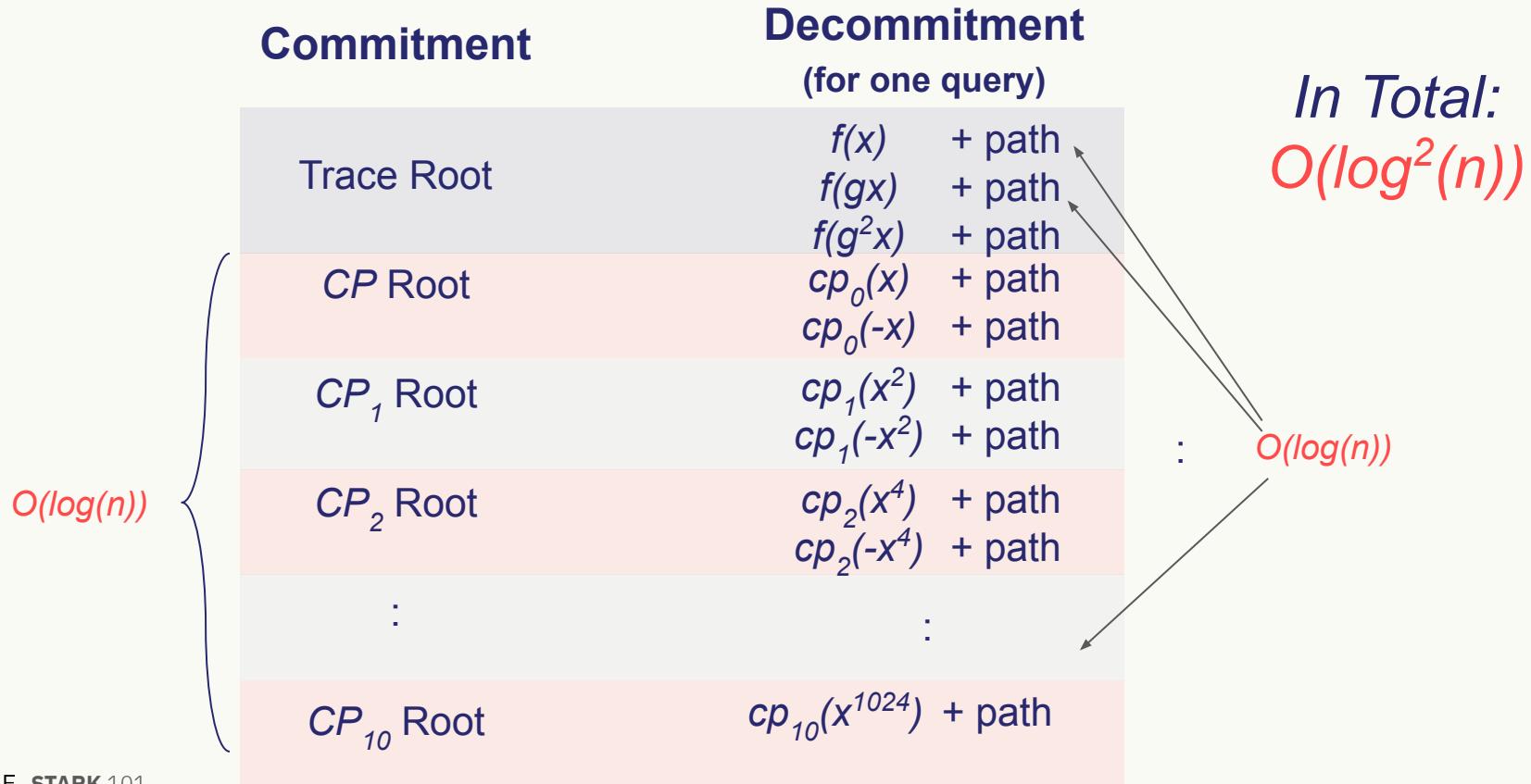
 Decommitment

- Get  $q$  random elements
- Provide a validation data for each

# Proof Length

| Commitment   | Decommitment<br>(for one query) |
|--|---------------------------------|
| O( $\log(n)$ )<br>$n = \text{trace length (1023)}$ | Trace Root $f(x)$ + path        |
|  | $f(gx)$ + path                  |
|  | $f(g^2x)$ + path                |
|  | $cp_0(x)$ + path                |
|  | $cp_0(-x)$ + path               |
|  | $cp_1(x^2)$ + path              |
|  | $cp_1(-x^2)$ + path             |
|  | $cp_2(x^4)$ + path              |
|  | $cp_2(-x^4)$ + path             |
|  | :                               |
|  | $cp_{10}(x^{1024})$ + path      |

# Proof Length



# Proof Length

| Commitment     |                                | Decommitment<br>for q queries |                                |                            |
|----------------|--------------------------------|-------------------------------|--------------------------------|----------------------------|
| Trace Root     | $f(x)$<br>$f(gx)$<br>$f(g^2x)$ | + path<br>+ path<br>+ path    | $f(x)$<br>$f(gx)$<br>$f(g^2x)$ | + path<br>+ path<br>+ path |
| CP Root        | $cp_0(x)$<br>$cp_0(-x)$        | + path<br>+ path              | $cp_0(x)$<br>$cp_0(-x)$        | + path<br>+ path           |
| $CP_1$ Root    | $cp_1(x^2)$<br>$cp_1(-x^2)$    | + path<br>+ path              | $cp_1(x^2)$<br>$cp_1(-x^2)$    | + path<br>+ path           |
| $CP_2$ Root    | $cp_2(x^4)$<br>$cp_2(-x^4)$    | + path<br>+ path              | $cp_2(x^4)$<br>$cp_2(-x^4)$    | + path<br>+ path           |
| :              | :                              | ...                           | :                              |                            |
| $CP_{10}$ Root | $cp_{10}(x^{1024})$            | + path                        | $cp_{10}(x^{1024})$            | + path                     |

$O(\log^2(n))$

# Thank you