

4.6

2.

```
EDGE(I, J)
RESULT ← F
Y ← VERT[I]
WHILE(Y ≠ 0)
    IF(HEAD(X) == J) THEN
        RESULT ← T
    ELSE
        X ← NEXT(X)
```

3. On average, EDGE must look at the average number of edges from any vertex. If R has P edges and N vertices, then EDGE examines $\frac{\sum P_{ij}}{N} = \frac{P}{N}$ edges on average.

4.

```
LOOK(NUM, NEXT, START, N, K)
X ← START
WHILE(X ≠ 0)
    IF(K == NUM(X)) THEN
        RETURN X
    ELSE
        X ← NEXT(X)
PRINT("NOT FOUND")
```

6.

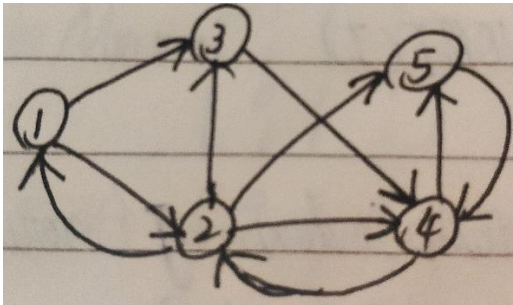
一种解法：

VERT	TAIL	HEAD	NEXT
2	4	1	9

4	1	1	3
7	1	2	5
1	2	1	6
	1	3	0
	2	4	8
	3	4	10
	2	3	0
	4	3	0
	3	3	0

8.

$$M_R = \begin{pmatrix} 0 & 1 & 1 & 0 & 0 \\ 1 & 0 & 1 & 1 & 1 \\ 0 & 0 & 0 & 1 & 0 \\ 0 & 1 & 0 & 0 & 1 \\ 0 & 0 & 0 & 1 & 0 \end{pmatrix}$$



12.

一种解法：

VERT	TAIL	HEAD	NEXT
1	F	F	2
4	F	M	3

7	F	W	0
11	M	M	5
	M	F	6
	M	W	0
	R	F	8
	R	M	9
	R	R	10
	R	W	0
	W	F	12
	W	M	13
	W	W	0

4.7

7. (a) $\{(2, 1), (3, 1), (3, 2), (3, 3), (4, 2), (4, 3), (4, 4), (1, 4)\}$.

(b) $\{(1, 1), (1, 2), (2, 2), (2, 3), (2, 4), (4, 1), (3, 4)\}$.

(c) $\{(1, 1), (1, 2), (1, 3), (1, 4), (2, 2), (2, 3), (2, 4), (3, 1), (3, 2), (3, 4), (4, 1), (4, 4)\}$.

(d) $\{(1, 1), (2, 1), (2, 2), (1, 4), (4, 1), (2, 3), (3, 2), (1, 3), (4, 2), (3, 4), (4, 4)\}$.

8.

(a) $\{(a,c),(a,e) , (b,a),(b,c) , (b,d),(c,a) , (c,c),(c,d) , (c,e),(d,a) , (d,b),(d,e) , (e,c),(e,d) , (e,e)\}$

(b) $\{(a,a), (a,d),(c,b) , (e,a),(e,b)\}$

(c) $\{(a,a),(a,b),(a,d) ,(a,e),(b,a),(b,b),(b,e),(c,b),(d,c),(d,d),(e,a),(e,b),(e,d),(e,e)\}$

(d) $\{(a,a), (e,a), (d,a), (a,b), (b,c), (d,c), (e,e), (a,e), (b,e), (d,e)\}$

12.

$$\mathbf{M}_{R \cap S} = \begin{pmatrix} 0 & 0 & 0 \\ 0 & 0 & 1 \\ 0 & 0 & 0 \\ 0 & 1 & 0 \end{pmatrix} \quad \mathbf{M}_{R \cup S} = \begin{pmatrix} 1 & 1 & 1 \\ 1 & 1 & 1 \\ 0 & 1 & 1 \\ 1 & 1 & 1 \end{pmatrix} \quad \mathbf{M}_{R^{-1}} = \begin{pmatrix} 0 & 0 & 0 & 1 \\ 1 & 1 & 0 & 1 \\ 0 & 1 & 1 & 1 \end{pmatrix} \quad \mathbf{M}_S^{-1} = \begin{pmatrix} 0 & 1 & 0 \\ 0 & 1 & 0 \\ 1 & 0 & 1 \\ 1 & 0 & 1 \end{pmatrix}$$

14.

$$R \cap S = \{(1,1), (1,2), (2,1), (2,2), (3,3), (4,4), (5,5), (6,6)\}$$

$$\{\{1,2\}, \{3\}, \{4\}, \{5\}, \{6\}\}$$

19. The definitions of irreflexive, asymmetric, and antisymmetric each require that a certain pair does not belong to R . We cannot “fix” this by including more pairs in R .

20.(a)成立, (b)不成立

23. (a) Reflexive. $a R a \wedge a S a \Rightarrow a S \circ R a$.
 Irreflexive. No. $1 R 2 \wedge 2 S 1 \Rightarrow 1 S \circ R 1$.
 Symmetric. No. $1 R 3, 3 R 1, 3 S 2, 2 S 3 \Rightarrow 1 S \circ R 2$, but $2 \not S \circ R 1$.
 Asymmetric. No. $R = \{(1, 2), (3, 4)\}$ and $S = \{(2, 3), (4, 1)\}$ provide a counterexample.
 Antisymmetric. No. $R = \{(a, b), (c, d)\}$ and $S = \{(b, c), (d, a)\}$ provide a counterexample.
 Transitive. No. $R = \{(a, d), (b, e)\}$ and $S = \{(d, b), (e, c)\}$ provide a counterexample.
- (b) No, symmetric and transitive properties are not preserved.

24.

$$\begin{array}{ll}
 \text{(a)} \begin{pmatrix} 1 & 1 & 1 & 1 & 1 \\ 1 & 1 & 1 & 1 & 0 \\ 1 & 0 & 1 & 1 & 1 \\ 1 & 0 & 1 & 1 & 1 \\ 1 & 1 & 1 & 1 & 1 \end{pmatrix} & \text{(c)} \begin{pmatrix} 1 & 1 & 1 & 1 & 1 \\ 1 & 0 & 1 & 0 & 0 \\ 1 & 1 & 1 & 1 & 1 \\ 1 & 0 & 1 & 1 & 0 \\ 1 & 1 & 1 & 1 & 1 \end{pmatrix} \\
 \text{(b)} \begin{pmatrix} 1 & 0 & 1 & 1 & 1 \\ 1 & 0 & 1 & 1 & 1 \\ 1 & 0 & 1 & 1 & 1 \\ 1 & 1 & 1 & 1 & 1 \\ 1 & 1 & 1 & 1 & 1 \end{pmatrix} & \text{(d)} \begin{pmatrix} 1 & 1 & 1 & 1 & 1 \\ 1 & 0 & 1 & 1 & 0 \\ 1 & 0 & 1 & 1 & 0 \\ 1 & 1 & 1 & 1 & 1 \\ 1 & 0 & 0 & 1 & 1 \end{pmatrix}
 \end{array}$$

26.

If $a(R \cap S)b$, then $a R b$, since R is asymmetric, there is no $b R a$. Then there is no $b(R \cap S)a$, so $R \cap S$ is asymmetric.

Let $R = \{(1, 2)\}$, $S = \{(2, 1)\}$, so $R \cup S = \{(1, 2), (2, 1)\}$, as we can see, $R \cup S$ is not asymmetric.

27. $R \cap S$ is antisymmetric. If $a(R \cap S)b$ and $b(R \cap S)a$, then $a R b$ and $b R a$. Hence $a = b$ because R is antisymmetric. $R \cup S$ may not be antisymmetric. Let $R = \{(1, 2)\}$, $S = \{(2, 1)\}$.

28.

Let $a(S \cup T) \circ R c$, then there is $a(S \cup T)b$ and $b R c$. Then $a S b$ or $a T b$, so we have $a S \circ R c$ or $a T \circ R c$. Then $a (S \circ R) \cup (T \circ R) c$. Hence, $(S \cup T) \circ R = (S \circ R) \cup (T \circ R)$

30.

Let $a T \circ R c$, so $a R b$ and $b T c$. As $R \subseteq S$, we have $a S b$, then $a T \circ S c$. So if $R \subseteq S$, then $T \circ R \subseteq T \circ S$.

31. (a) Let $\mathbf{M}_{R \cap S} = [m_{ij}]$, $\mathbf{M}_R = [r_{ij}]$, $\mathbf{M}_S = [s_{ij}]$. $m_{ij} = 1$ if and only if $(i, j) \in R \cap S$. $(i, j) \in R$ if and only if $r_{ij} = 1$ and $(i, j) \in S$ if and only if $s_{ij} = 1$. But this happens if and only if the i, j th entry of $\mathbf{M}_R \wedge \mathbf{M}_S$ is 1.

(b) Let $\mathbf{M}_{R \cup S} = [m_{ij}]$, $\mathbf{M}_R = [r_{ij}]$, $\mathbf{M}_S = [s_{ij}]$. $m_{ij} = 1$ if and only if $(i, j) \in R \cup S$. $(i, j) \in R$ if and only if $r_{ij} = 1$ or $(i, j) \in S$ if and only if $s_{ij} = 1$. But this happens if and only if the i, j th entry of $\mathbf{M}_R \vee \mathbf{M}_S$ is 1.

(c) The i, j th entry of $\mathbf{M}_{R^{-1}}$ is 1 if and only if $(i, j) \in R^{-1}$ if and only if $(j, i) \in R$ if and only if the j, i th entry of \mathbf{M}_R is 1 if and only if the i, j th entry of \mathbf{M}_R^T is 1.

(d) The i, j th entry of $\mathbf{M}_{\overline{R}}$ is 1 if and only if $(i, j) \in \overline{R}$ if and only if $(i, j) \notin R$ if and only if the i, j th entry of \mathbf{M}_R is 0 if and only if the i, j th entry of $\overline{\mathbf{M}_R}$ is 1.

36.

If $a R - S b$, then $a R b$ and there is no $a S b$. As R and S are symmetric relations, so $b R a$ and there is no $b S a$. Then $b R - S a$. Hence $R - S$ is also a symmetric relation on A .

37. (a) R is symmetric if and only if $x R y \Rightarrow y R x$ if and only if $R \subseteq R^{-1} \subseteq R$.

(b) Suppose R is antisymmetric. Let $(x, y) \in R \cap R^{-1}$, then $x = y$ and $(x, y) \in \Delta$. Suppose $R \cap R^{-1} \subseteq \Delta$. If $x R y$ and $y R x$, then $x R y$ and $x R^{-1} y$. Thus $(x, y) \in \Delta$, so $x = y$.

(c) Suppose R is asymmetric. Let $(x, y) \in R \cap R^{-1}$. This contradicts the fact that R is asymmetric. Hence $R \cap R^{-1} = \emptyset$. Let $x R y$ and $y R x$. Then $(x, y) \in R \cap R^{-1} = \emptyset$. Hence R is asymmetric.

4.8

8.

Since $S^\infty = S \cup S^2 \cup \dots \cup S^n$, $S = R^2$, hence, $S^\infty = R^2 \cup R^4 \cup \dots \cup R^{2n}$

Therefore, $a S^\infty b \Leftrightarrow a(R^2 \cup R^4 \cup \dots \cup R^{2n})b \Leftrightarrow a R^2 b \vee a R^4 b \vee \dots \vee a R^{2n} b \Leftrightarrow$ there is a path in R from a to b having an even number of edges.

10.

$n = 4$

$$\text{Let } W_0 = M_R = \begin{bmatrix} 1 & 1 & 0 & 0 \\ 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 \end{bmatrix}, \text{ then } W_1 = \begin{bmatrix} 1 & 1 & 0 & 0 \\ 1 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 \end{bmatrix}, W_2 = \begin{bmatrix} 1 & 1 & 0 & 0 \\ 1 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 \end{bmatrix},$$

$$W_3 = W_4 = \begin{bmatrix} 1 & 1 & 0 & 0 \\ 1 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 \end{bmatrix}, \text{ therefore, } M_{R^\infty} = W_4 = \begin{bmatrix} 1 & 1 & 0 & 0 \\ 1 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 \end{bmatrix}$$

12.

$$n = 4$$

$$\text{Let } W_0 = M_R = \begin{bmatrix} 0 & 0 & 0 & 1 \\ 1 & 0 & 0 & 1 \\ 0 & 1 & 0 & 1 \\ 0 & 0 & 1 & 0 \end{bmatrix}, \text{ then } W_1 = \begin{bmatrix} 0 & 0 & 0 & 1 \\ 1 & 0 & 0 & 1 \\ 0 & 1 & 0 & 1 \\ 0 & 0 & 1 & 0 \end{bmatrix}, W_2 = \begin{bmatrix} 0 & 0 & 0 & 1 \\ 1 & 0 & 0 & 1 \\ 1 & 1 & 0 & 1 \\ 0 & 0 & 1 & 0 \end{bmatrix},$$

$$W_3 = \begin{bmatrix} 0 & 0 & 0 & 1 \\ 1 & 0 & 0 & 1 \\ 1 & 1 & 0 & 1 \\ 1 & 1 & 1 & 1 \end{bmatrix}, W_4 = \begin{bmatrix} 1 & 1 & 1 & 1 \\ 1 & 1 & 1 & 1 \\ 1 & 1 & 1 & 1 \\ 1 & 1 & 1 & 1 \end{bmatrix}, \text{ therefore, } M_{R^\infty} = W_4 = \begin{bmatrix} 1 & 1 & 1 & 1 \\ 1 & 1 & 1 & 1 \\ 1 & 1 & 1 & 1 \\ 1 & 1 & 1 & 1 \end{bmatrix}$$

14.

Disproof.

$$M_R = \begin{bmatrix} 0 & 0 & 0 & 0 \\ 1 & 1 & 1 & 0 \\ 0 & 1 & 1 & 0 \\ 0 & 1 & 0 & 0 \end{bmatrix}$$

$$M_{R_t} = \begin{bmatrix} 0 & 0 & 0 & 0 \\ 1 & 1 & 1 & 0 \\ 1 & 1 & 1 & 0 \\ 1 & 1 & 1 & 0 \end{bmatrix}, M_{s(R_t)} = \begin{bmatrix} 0 & 1 & 1 & 1 \\ 1 & 1 & 1 & 1 \\ 1 & 1 & 1 & 1 \\ 1 & 1 & 1 & 0 \end{bmatrix}$$

$$M_{R_s} = \begin{bmatrix} 0 & 1 & 0 & 0 \\ 1 & 1 & 1 & 1 \\ 0 & 1 & 1 & 0 \\ 0 & 1 & 0 & 0 \end{bmatrix}, M_{t(R_s)} = \begin{bmatrix} 1 & 1 & 1 & 1 \\ 1 & 1 & 1 & 1 \\ 1 & 1 & 1 & 1 \\ 1 & 1 & 1 & 1 \end{bmatrix}$$

Therefore, $M_{s(R_t)} \neq M_{t(R_s)}$, the symmetric closure of R_t is not the same relation as the transitive closure of R_s .

18.

$$M_{(R \cup S)^\infty} = \begin{bmatrix} 1 & 1 & 1 & 0 & 0 \\ 1 & 1 & 1 & 0 & 0 \\ 1 & 1 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 & 1 \\ 0 & 0 & 0 & 1 & 1 \end{bmatrix}$$

$$A/R = \{\{1\}, \{2,3\}, \{4,5\}\}, A/S = \{\{1,2\}, \{3\}, \{4\}, \{5\}\}, A/(R \cup S)^\infty = \{\{1,2,3\}, \{4,5\}\}.$$

20.

19 题中的过程仅适用于 R 和 S 均为等价关系，且计算 $(R \cup S)^\infty$ 过程中需要用到 Warshall 算法，所以无法替代 Warshall 算法。

23.

We first show R^∞ is transitive. Then we show it is the smallest relation that contains R . It is a direct proof.

24.

首先证明若从 a 到 b 之间存在环，则存在更短的路径；再证明 a 到 b 之间，若 a, b 不同，则最长路径长度为 $n-1$ ，若 a, b 相同，则最长路径长度为 n ；因此， $\forall k \in \mathbb{Z}, 1 \leq k \leq n, aR^k b$ ，因此 $R^\infty = R \cup R^2 \cup R^n$ 。间接证明。

25.

$$\mathbb{R} \times \mathbb{R}.$$

5.1

5. Each integer has a unique square that is also an integer.

6. For each $a \in \mathbb{R}$, $e^a \in \mathbb{R}$, and for each $a \in \mathbb{R}$, there is a unique e^a .

7. Each $r \in \mathbb{R}$ is either an integer or it is not.

8. For each $a \in \mathbb{R}$, there is a unique integer which is the greatest integer less than or equal to a .

11. (a) Both. (b) Neither.

12. (a) One to one. (b) Onto.

13. (a) Both. (b) Onto.

14. (a) Onto. (b) Neither.

15. (a) Both. (b) Onto.

29. n^m .

30.

If $f : A \rightarrow B$ and $g : B \rightarrow C$ are one-to-one functions, then $g \circ f$ is one to one.

Proof: Let $a_1, a_2 \in A$. Suppose $(g \circ f)(a_1) = (g \circ f)(a_2)$. Then $g(f(a_1)) = g(f(a_2))$ and $f(a_1) = f(a_2)$, because $(g \circ f)(a_1) = g(f(a_1)), (g \circ f)(a_2) = g(f(a_2))$ and g is one-to-one functions. Thus $a_1 = a_2$, because f is one-to-one functions. Hence $g \circ f$ is one to one.

31.

If $f : A \rightarrow B$ and $g : B \rightarrow C$ are onto functions, then $g \circ f$ is onto.

Proof: Choose $x \in C$. Then there exists $y \in B$ such that $g(y) = x$. (Why? because g is onto) Then there exists $z \in A$ such that $f(z) = y$ (why? because f is onto) and $(g \circ f)(z) = x$. Hence, $g \circ f$ is onto.

33.

Suppose $g \circ f$ is onto. Let $c \in C$. Then $\exists a \in A$ such that $(g \circ f)(a) = c$. But $(g \circ f)(a) = g(f(a))$, $f(a) \in B$, so g is onto.

34.

Use mathematical induction to prove, $\forall k \in \mathbb{Z}$, f^k is bijection.

Therefore, if $a_1, a_2 \in A$,

$$O(a_1, f) = \{f^0(a_1), f^1(a_1), f^{-1}(a_1), f^2(a_1), f^{-2}(a_1), \dots\},$$

$$O(a_2, f) = \{f^0(a_2), f^1(a_2), f^{-1}(a_2), f^2(a_2), f^{-2}(a_2), \dots\}$$

If $a_1 \neq a_2$, $\forall k \in \mathbb{Z}$, $f^k(a_1) \neq f^k(a_2)$, then $O(a_1, f) \cap O(a_2, f) = \emptyset$.

Therefore, if $O(a_1, f) \cap O(a_2, f) \neq \emptyset$, then $a_1 = a_2$, then $O(a_1, f) = O(a_2, f)$.

40.

(a) Disproof.

$$f(a_1 + a_2) = (a_1 + a_2)^2, f(a_1) + f(a_2) = a_1^2 + a_2^2$$

If $a_1, a_2 \neq 0$, then $(a_1 + a_2)^2 \neq a_1^2 + a_2^2$, then $f(a_1 + a_2) \neq f(a_1) + f(a_2)$.

(b) Proof.

The length of $s_1 \cdot s_2$ is equal to the length of s_1 plus the length of s_2 .

Therefore, $f(s_1 \cdot s_2) = f(s_1) + f(s_2)$.

41. Disproof.

(a) and (b). Consider the table for \diamond .

\diamond	0	1
0	0	1
1	1	0

Since $f(0) = \text{true}$ and $f(1) = \text{false}$, we see this is not the table for either \vee or \wedge .

5.2

7.

$n = ak + r, 0 \leq r < k$. Since $k < 2k < 3k < \dots < ak \leq n$ the number of multiples of k between 1 and n is a . But $\frac{n}{k} = a + \frac{r}{k}$ with $0 \leq \frac{r}{k} < 1$ so $\lfloor \frac{n}{k} \rfloor = a$.

8.

设 $n=2k+1, (k \in \mathbb{Z})$ $\frac{n^2}{4} = \frac{4k^2+4k+1}{4} = k^2+k+\frac{1}{4}$, 向上取整为 k^2+k+1 。

$$\frac{n^2+3}{4} = \frac{4k^2+4k+4}{4} = k^2+k+1$$

$$\text{所以 } \lceil \frac{n^2}{4} \rceil = \frac{n^2+3}{4}$$

18. (例 7.(a)中 $l: A^* \rightarrow \mathbb{Z}$ 为 $l(w)$, 它是字符串 w 的长度)

(a) 由于 A^* 中元素均为字符串, 字符串必有唯一长度, 所以 l 是处处有定义的。

(b) ' ab ' $\in A^*$, ' ba ' $\in A^*$, $l('ab')=2, l('ba')=2$, 所以 l 不是单射。

(c) $l(w)$ 为字符串 w 的长度, $l(w) \geq 0$, 但 $l(w) \in \mathbb{Z}$, 所以 l 不是满射。

20.

存在 2 个不同的关于 p 的布尔函数, $f(p) = p, g(p) = \neg p$

存在 4 个有两个布尔变量的布尔函数, 设这两个布尔变量为 p, q

$$f_1(p, q) = p \& q, f_2(p, q) = p \& \neg q, f_3(p, q) = \neg p \& q, f_4(p, q) = \neg p \& \neg q$$

28.

对任意 S 为 A 的子集, 令 $f(x)=1, x \in S, 0, x \in A \setminus S$; 这样对每个元素 $x \in A, f(x)$ 有两个取值, 0 或 1; 因此, 根据乘法原理, 这样的 f 有 2^n 个, 每个 f 唯一的对应 A 的一个子集, 因此 $|\text{pow}(A)| = 2^n$

29.

$f^{-1}(1)$ is the set of elements of A .

5.3

11.

$\{f_5\}, \{f_6, f_{10}, f_{11}\}, \{f_7\}, \{f_4\}, \{f_8\}, \{f_1\}, \{f_2\}, \{f_3\}, \{f_9\}, \{f_{12}\}.$

12.

$\{f_5\}, \{f_7\}, \{f_4\}, \{f_8\}, \{f_6, f_{10}, f_{11}\}, \{f_{12}\}, \{f_1\}, \{f_2\}, \{f_3\}, \{f_9\}$

13.

$f_1, \Theta(n \lg n); f_2, \Theta(n^2); f_4, \Theta(\lg n); f_5, \Theta(1); f_6, \Theta(n); f_{10}, \Theta(n); f_{11}, \Theta(n).$

20.

$\Theta(n)$

21.

$f(n, m, q) = 1 + nq + 3nmq + 1.$ Let $N = \max(n, m, q)$, then f is $\Theta(N^3)$.

22.

$f(n+2) = 2 * f(n) + f(n+1)$

$f(0) = 0, f(1) = 1$

得 $f(n) = 2^n/6 + (-1)^n/3$

$f(n) \in \Theta(2^n)$

伪代码:

1. $f0 \leftarrow 0$

2. $f1 \leftarrow -1$

3. **FOR** $I=0$ **THRU** N

 a. $f2 = 2*f0 + f1$

 b. $f0 = f1$

 c. $f1 = f2$

4. **RETURN** $f2$

计算 $F(N)$ 所需要的运行时间为 $O(n)$

23.

(a) $P_n = P_{n-1} + (n-2) + (n-3), P_3 = 1, P_4 = 4.$

(b) $\theta(n^2).$

24.

若 $\Theta(n^a)$ 比 $\Theta(n^b)$ 低, 则存在 $n \geq k$ 和常数 $c, |n^a| < c|n^b| = c|n^a| \cdot |n^{b-a}|$

即 $c|n^{b-a}| > 1$, c 为常数, 所以 $b > a > 0$

若 $b > a > 0$, 则存在 $n \geq 2$ 和常数 $c = 1, |n^a| < |n^b| = |n^a| \cdot |n^{b-a}|$

即 $\Theta(n^a)$ 比 $\Theta(n^b)$ 低

所以 $\Theta(n^a)$ 比 $\Theta(n^b)$ 低, 当且仅当 $b > a > 0$ 。

25.

若 $\Theta(a^n)$ 比 $\Theta(b^n)$ 低，则存在 $n \geq k$ 和常数 $c, |a^n| < c|b^n|$ ，即 $c > (a/b)^n$

则 $0 < a < b$ 。

若 $0 < a < b$ ，则存在 $n \geq 2$ 和常数 $c = 1, |a^n| < |b^n|$ ，所以 $\Theta(a^n)$ 比 $\Theta(b^n)$ 低。

综上 $\Theta(a^n)$ 比 $\Theta(b^n)$ 低，当且仅当 $0 < a < b$

26.

设 $g(x) = r f(x)$

若 $r > 0$ 存在 $n \geq 1$ ，和常数 $c = 2r$ ，使得 $|g(x)| = |r f(x)| \leq 2r|f(x)|$ 成立

同时存在 $n \geq 1$ ，和常数 $c = 2/r$ ，使得 $|f(x)| \leq 2/r * |r f(x)| = 2|f(x)|$ 成立

若 $r < 0$ 存在 $n \geq 1$ ，和常数 $c = -2r$ ，使得 $|g(x)| = |r f(x)| \leq -2r|f(x)|$ 成立

同时存在 $n \geq 1$ ，和常数 $c = -2/r$ ，使得 $|f(x)| \leq -2/r * |r f(x)| = 2|f(x)|$ 成立

所以如果 $r \neq 0$ ，那么对于任意函数 f ，有 $\Theta(rf) = \Theta(f)$

27.

Suppose $h(n) > 0, \forall n, \Theta(f)$ lower than (or the same as) $\Theta(g)$. $|f(n)| \leq c \cdot |g(n)|, n \geq k$ (and $|g(n)| \leq d \cdot |f(n)|, n \geq l$). $h(n)|f(n)| \leq c \cdot h(n) \cdot |g(n)|, n \geq k$ (and $h(n)|g(n)| \leq d \cdot h(n) \cdot |f(n)|, n \geq l$). Hence $|f(n) \cdot h(n)| \leq c \cdot |g(n) \cdot h(n)|, n \geq k$ (and $|g(n) \cdot h(n)| \leq d \cdot |f(n) \cdot h(n)|, n \geq l$). Hence $\Theta(fh)$ is lower than (or the same as) $\Theta(gh)$. Note that if $\Theta(f)$ is strictly lower than $\Theta(g)$, then $\Theta(fh)$ must be strictly lower than $\Theta(gh)$.

28.

若 $\Theta(f) = \Theta(h)$ ，则 $f = O(h)$ ，有 $n \geq k_1$ ，和常数 c_1 ，使得 $|f(n)| \leq c_1|h(n)|$

同样若 $\Theta(g) = \Theta(h)$ ，则 $g = O(h)$ ，有 $n \geq k_2$ ，和常数 c_2 ，使得 $|g(n)| \leq c_2|h(n)|$

则存在 $n \geq \max\{k_1, k_2\}$ 和常数 c_1, c_2 , 使得 $|f(n)| + |g(n)| \leq (c_1 + c_2)|h(n)|$

$$|f(n) + g(n)| \leq |f(n)| + |g(n)|$$

则结论 如果 $\Theta(f) = \Theta(g) = \Theta(h)$, 那么 $f+g$ 是 $O(h)$ 的 成立。

29.

26 题已证明定理 6, 结合关系 Θ 是传递的, 结论得证。

5.4

12.

(a) $(6, 8) \circ (2, 3) \circ (1, 4, 5)$

(b) $(1, 2, 3, 4) \circ (5, 7, 8, 6)$

13.

(a) $(1, 6, 3, 7, 2, 5, 4, 8)$

(b) $(5, 6, 7, 8) \circ (1, 2, 3)$

14.

(a) (a, g, e, c, b, d)

(b) (a, d, b, e, g, c)

15.

(a) $(2, 6) \circ (2, 8) \circ (2, 5) \circ (2, 4) \circ (2, 1).$

(b) $(3, 6) \circ (3, 1) \circ (4, 5) \circ (4, 2) \circ (4, 8).$

16.

RAEEU YEO HRW

20.

(a)将置换写成不相交循环的积

$(1,4,6,8,3) = (1,3)。(1,8)。(1,6)。(1,4)$ 为偶置换

(b)将置换写成不相交循环的积

$(1,7,6,8,5)。(2,3,4) = (1,5)。(1,8)。(1,6)。(1,7)。(2,4)。(2,3)$ 为偶置换

26.

p 是集合 A 到自身的双射。

$p^2 = p \circ p$, 也是集合 A 到自身的双射。

所以 p^2 是置换。

28.

(a) $(2,3,5)。(1,4)$

(b)

1	2	3	4	5	6
4	5	2	1	3	6

(c)

1	2	3	4	5	6
1	5	2	4	3	6

(d) p 的周期为 $2 \times 3 = 6$

29.

(a) Basis step: $n = 1$. If p is a permutation of a finite set A , then p^1 is a permutation of A is true.

Induction step: The argument in Exercise 26 also shows that if p^{n-1} is a permutation of A , then $p^{n-1} \circ p$ is a permutation of A . Hence p^n is a permutation of A .

(b) If $|A| = n$, then there are $n!$ permutations of A . Hence, the sequence $1_A, p, p^2, p^3, \dots$ is finite and $p^i = p^j$ for some $i \neq j$. Suppose $i < j$. Then $p^{-i} \circ p^i = 1_A = p^{-i} \circ p^j$. So $p^{j-i} = 1_A, j-i \in \mathbb{Z}$.

30.

证自反性。对于 A 中任一元素 a , 有 $p^0(a) = a$, 所以 $a R a$ 。自反性成立。

证对称性。若 $a R b$, 则 $p^n(a) = b$, 则 $p^{-n}(b) = a$, $b R a$ 。对称性成立。

证传递性。若 $a R b, b R c$, 则 $p^{n_1}(a) = b, p^{n_2}(b) = c$, 则 $p^{n_1} \circ p^{n_2}(a) = c$,

即 $p^{n_1+n_2}(a) = c$, 即 $a R c$ 。传递性成立。

所以 R 是一个等价关系。

37.

(a) 3

(b) 6

38.

10

39.

- . For each increasing sequence of length $\lceil \frac{n}{2} \rceil$, there is exactly one associated up-down permutation of A , because there is just one way to arrange the remaining elements of A in decreasing order and insert them to fill the even positions.