

Functions

First Class

Prefix notation:

```
(predicate argument argument ...)
```

Pass a function as an argument:

```
(predicate arguments
```

```
(predicate (predicate arguments)))
```

First Class

- Homoiconicity
 - Arguments are data
 - Operators are data
- \therefore functions are treated like any other object
 - Passed as an argument
 - Returned as a result
 - Created at runtime

Pure

- Pure function: compiled → evaluated → result
- Side effects
 - Possible in Clojure
 - Do not change the program's state, unless *explicitly* made to do so
 - Example: `println`

Creating Functions

`fn`

`(fn [x] x) - (predicate argument argument)`

Creating Functions

Anonymous Functions

- `((fn [x] x) 2) ⇒ 2`
 - `((predicate argument argument) argument)`
- Parameters can be functions
 - `((fn [x] (zero? x)) 0) ⇒ true`
 - `((fn [x] (zero? x)) 4) ⇒ false`

Creating Functions

Anonymous Functions

- `#()` - for short functions passed as arguments
 - It takes arguments named `%`, `%2`, `%3` ... `%n`.
 - `(#(* 2 %) 3) ⇒ 6`
 - `((fn [x] (even? x)) (#(* 2 %) 3)) ⇒ true`
- An anonymous function has no name, so you don't know what to "call" it!

Creating Functions

Forms Revisited

“ A form can be converted into a function if we can determine the correspondence between the *variables* occurring in the *form* and the *ordered list of arguments* of the desired function. ”

John McCarthy *Recursive Functions of Symbolic Expressions and Their Computation by Machine, Part I* (1960)

Creating Functions

Forms Revisited → Functions

- `((fn [x y] (* (* 2 y) x)) -1 6) ⇒ -12`
 - *Form* - `(* (* 2 y) x)`
 - *Variables* - `[x y]`
 - *Ordered list of arguments* - `-1 6`

Creating Functions

Symbolic Expressions vs. Functions

Symbolic expressions are:

1. Atomic symbols

- `> 3` \Rightarrow `3`
- `> 3 y` (two expressions, not one)
 - \Rightarrow `3`
 - \Rightarrow `Unable to resolve symbol: y`

Creating Functions

Symbolic Expressions vs. Functions

Symbolic expressions are:

2. Formed by using the special characters `(`, `)`, and an infinite set of distinguishable atomic symbols separated by blanks.

- `' (3 4)` \Rightarrow `(3 4)`

- A single form

Creating Functions

Symbols Revisited

`def`

- **Defines** a symbol
- `(def hello-world "Hello World!")`
- `hello-world` \Rightarrow `"Hello World!"`

Creating Functions

def

- Creates or locates a **global var** with the name of **symbol**
- Can name a scalar: `(def x 1): x` \Rightarrow `1`
- Can name a collection: `(def x '(+ 2 3)): x` \Rightarrow `(+ 2 3)`
- Can name a function:
`(def double-num (fn [x] (* x 2))): (double-num 2)` \Rightarrow `4`

Creating Functions

Symbols Revisited

- Symbols are forms. They evaluate to what they *name*.
- Example: `inc` and `+` are symbols that *name* a function

- `(inc 2)` \Rightarrow `3`

- `(+ 2 1)` \Rightarrow `3`

- `inc` \Rightarrow `#object[clojure.core$inc...]`

Creating Functions

The Fast Way

defn

```
(defn double-num [x] (* x 2))
```

≡

```
(def double-num (fn [x] (* x 2)))
```

```
(double-num 3)
```

```
> 6
```

Creating Functions

Multi-arity

```
(defn do-something
  ([] "nothing")
  ([one] "one parameter")
  ([one two] "two parameters")
  ([one two & more] "more than two parameters!"))
```

```
(do-something)
> "nothing"
(do-something 1)
> "one parameter"
(do-something 1 2)
> "two parameters"
(do-something 1 2 3 4)
> "more than two parameters"
```


Creating Functions

Multi-arity

Faster

```
(defn do-something [a & [b c]]  
  (str "Required argument a is " a  
       ". Optional argument b is " b  
       ". C, optional, is " c "."))  
  
(do-something 1 2)  
  
> "Required argument a is 1.  
   Optional argument b is 2. C, optional, is ."
```

Local Bindings

Special Form

let

- Bindings are sequential
- Pairs: **symbols** and **init-exprs**

```
(let [x (+ 1 0)
      y x]
  y)

> 1
```

Local Bindings

Special Form

let

```
(let [double (fn [x] (* 2 x))] (double 21)) ⇒ 42
```

letfn

```
(letfn [(double [x] (* 2 x))] (double 21)) ⇒ 42
```

Controlling Flow

do

- `let` is an implicit `do`
- Evaluates expressions in order
- Fundamentally *imperative*
- Often used to create side effects (ex: print or i/o)

Controlling Flow

do

```
(if true (println "This is true: ") (+ 1 1))  
> This is true:  
> nil  
> ;; nil is the return value
```

VS.

```
(if true (do (println "This is true: ") (+ 1 1)))  
> This is true:  
> 2  
> ;; 2 is the return value
```

Controlling Flow and Scope

```
(defn y []  
  (let [x 1]  
    (println x)  
    (def x 2)  
    x))
```

⇒ 1

```
(defn y []  
  (let [x 1]  
    (println x)  
    (def x 2)  
    x)  
  x)
```

⇒ 2



Recursion



Recursion

- Each recursive call is added to the stack
- Each function is sequentially popped after the final call

```
(defn factorial  
  [n]  
  (if (== 1 n)  
      n  
      (* n (factorial (- n 1)))))
```

```
(factorial 10)  
> 3628800
```

```
(factorial 5000)  
> ERROR: Stack Overflow
```




Recursion

Mutual recursion

```
(letfn [(is-even? [n]
          (if (zero? n)
              true
              (is-odd? (dec n))))
        (is-odd? [n]
          (if (zero? n)
              false
              (is-even? (dec n))))]
  (is-even? 42))
```

```
> true
```

Recursion

`loop` and `recur`

- `recur` must be the last expression evaluated aka the "tail position"
- Form: `loop` \approx `let`
- Arity: the number of bindings.

```
(loop [x 10]  
  (when (> x 1)  
    (println x)  
    (recur (- x 2)))))
```



Recursion

loop and recur

```
(def factorial
  (fn [n]
    (loop [cnt n
          acc 1N]
      (if (>= 0 cnt)
          acc
          (recur (dec cnt) (* acc cnt))))))

(factorial 5000)
> 42285779266055435222...
```



Recursion

`loop` and `recur`

`recur`

- The only non-stack-consuming looping construct in Clojure
- Use in tail-position is verified by the compiler
- Since Clojure uses the Java calling conventions, tail call optimization must be made explicit by `recur`

Recursion vs. Looping

Style: Declarative vs. Imperative

Imperative - uses statements that change a program's state by describing the program's flow

```
var numbers = [1,2,3]
var total = 0

for(var i = 0; i < numbers.length; i++) {
    total += numbers[i]
}
```

Note: `n` and `total` are modified in the loop

Recursion vs. Looping

Style: Declarative vs. Imperative

Declarative - the function expresses the logic of a computation without describing its control flow

```
(fn [numbers]
  (loop [n numbers
        total 0]
    (if (empty? n)
        total
        (recur (rest n) (+ total (first n)))))))
```

Note: `n` and `total` are *not variables*, they are new local bindings in every recursive call

Recursion vs. Looping

Style: Declarative vs. Imperative

Equivalent Functional Solution

```
(reduce + '(1 2 3)) ⇒ 6
```

Note: this is idiomatic Clojure, more on this later!