

A

40	35	30	15	10	25	5
1	2	3	4	5	6	7

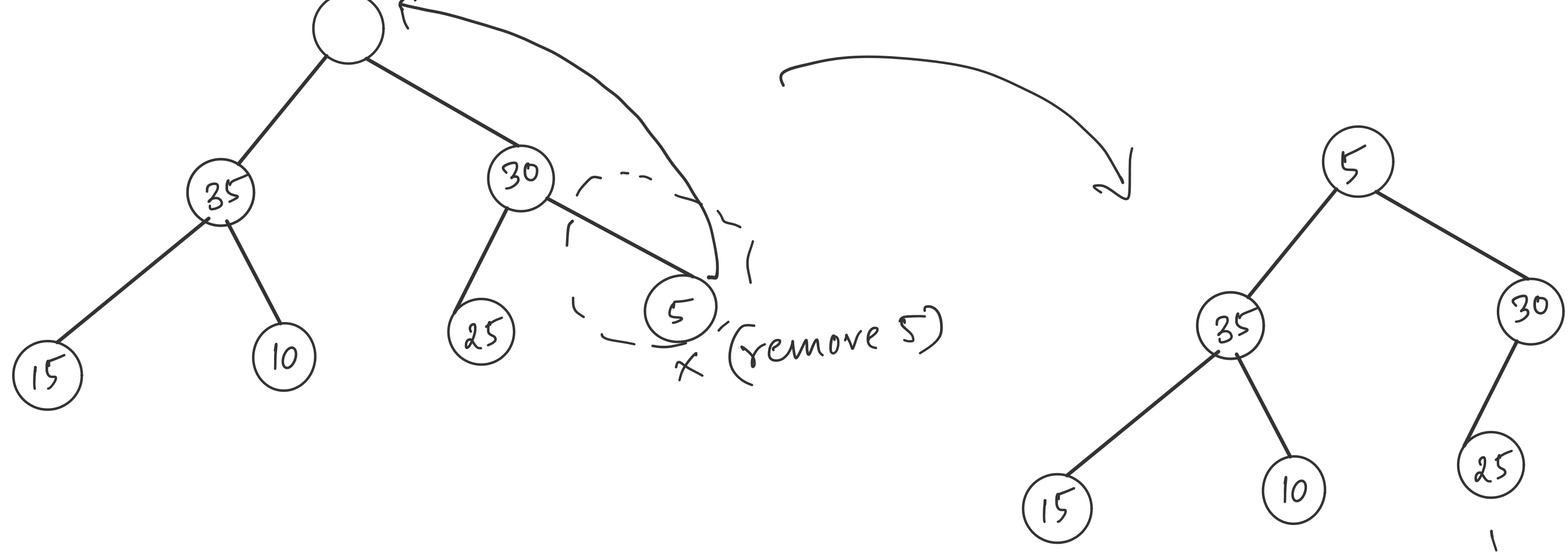
Note  $\Rightarrow$  from heap we can delete only root (i.e. the largest element in case of max heap).

It means we can delete highest priority element.

delete 40  $\Rightarrow$

Step 1  $\Rightarrow$  store 40 in a temporary variable and delete 40 from root.

Step 2  $\Rightarrow$  35 should take its place but this will make the binary tree incomplete.



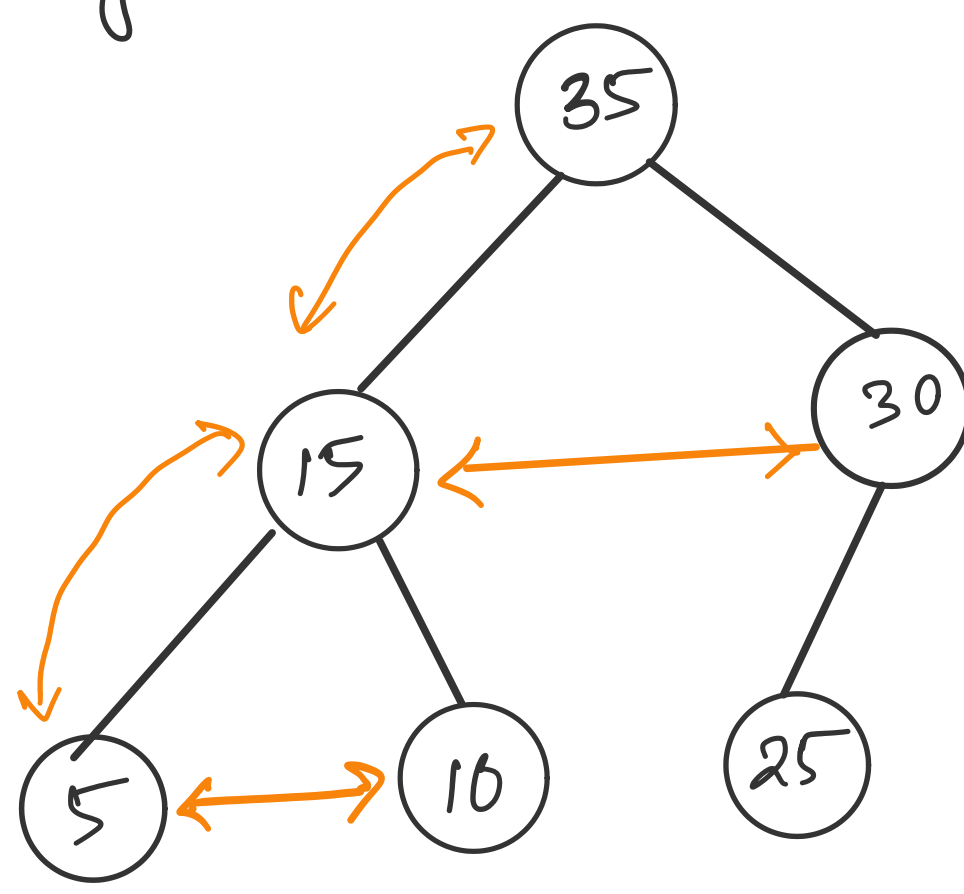
Step 2  $\Rightarrow$  bring 5 to root.

A

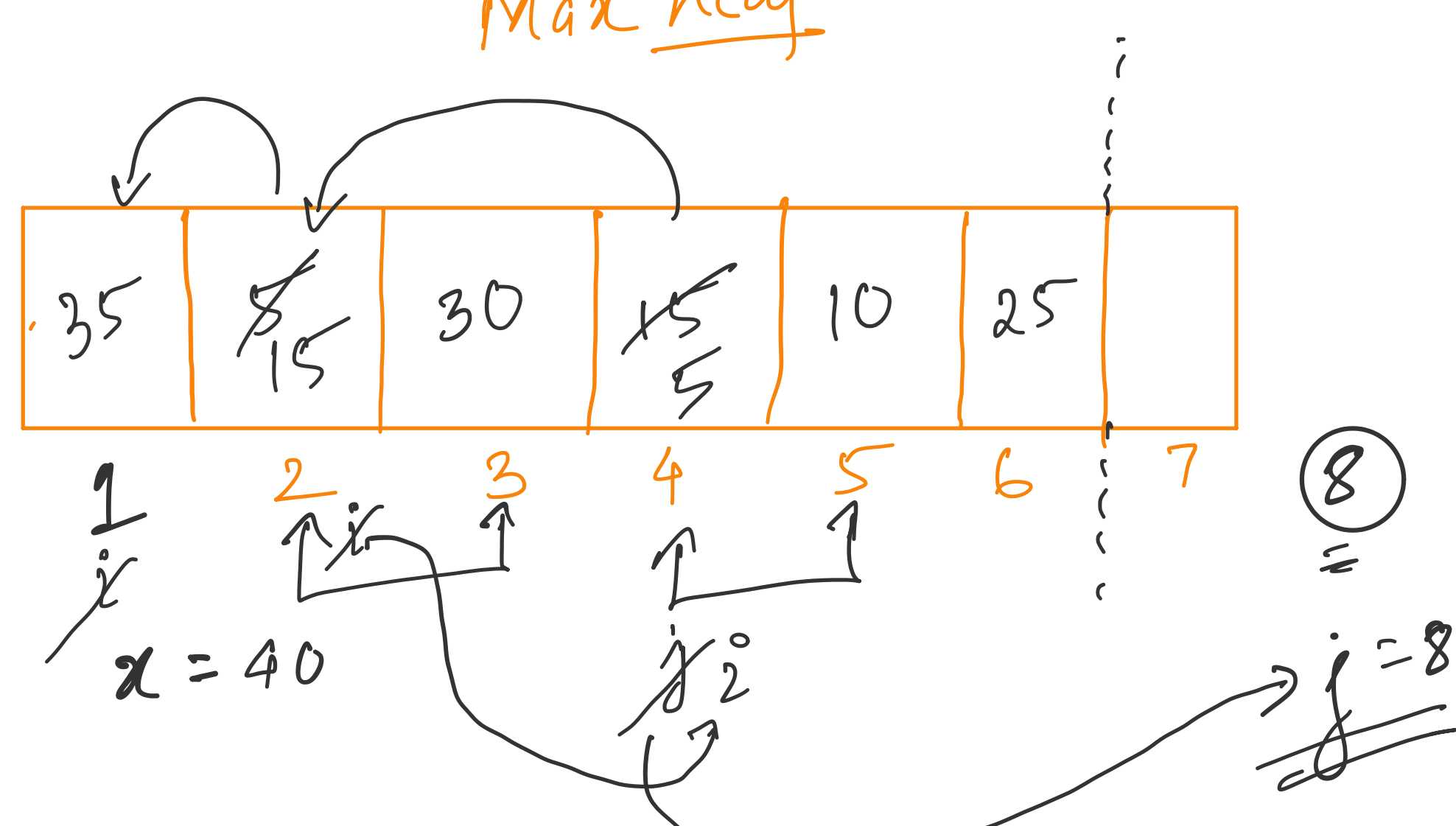
5	35	30	15	10	25	
1	2	3	4	5	6	7

Step 3  $\Rightarrow$  to rearrange and make it a max heap compare the children of new root i.e. 5. (35 and 30).

Now  $35 > 30$ ; compare 35 with 5. Now interchange 5 and 35. Now again compare children of 5 i.e. 15 and 10.  $15 > 10$  compare 15 and 5 and interchange them as  $5 < 15$ .



Max heap.



void Delete (int A[], int n)

```
{
    int x, i, j;
    x = A[n];
    A[1] = A[n];
    i = 1; j = 2 * i;
    while (j < n - 1)
    {
        if (A[j+1] > A[j])
            j = j + 1;
        if (A[i] < A[j])
            swap(A[i], A[j]);
            i = j;
            j = 2 * j;
        else
            break;
    }
```

Heap sort  $\Rightarrow$  whenever we delete an element from the heap we get a free space at the end now if we go on storing the deleted element in this free space then our array becomes sorted and this is called heap sort.

Implementation  $\Rightarrow$

void Delete (int A[], int n)

```
{
    int x, i, j;
    x = A[n];
    A[1] = A[n];
    i = 1; j = 2 * i;
    while (j < n - 1)
    {
        if (A[j+1] > A[j])
            j = j + 1;
        if (A[i] < A[j])
            swap(A[i], A[j]);
            i = j;
            j = 2 * j;
        else
            break;
    }
    A[n] = x; }
```

Heap sort  $\Rightarrow$

- 1) createHeap of n elements
- 2) delete n elements one by one

Time complexity  $\Rightarrow O(n \log n)$   
( $2n \log n$ )