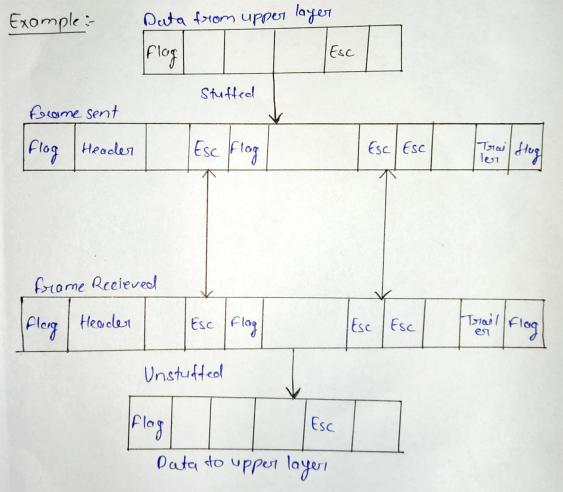
Name: Braner Mishro Registration No: 12/14762 Assignment: 2 Course Code: - CAP275 Teacher Name On, Mithleshs Cowise Title: Data communication & Networking Roll No: RDOC 09 ASS Ques: 1: What is difference between bit stuffing and byte stuffing? Explain with suitable example. The difference between bit stuffing and byte stuffing Ans: can be done under the following-Bit Stuffing: - Mostly flag is a special 8-bit pattern "DIIIIIO" used to detine the beginning and the end of the frame. Problem with the flag is the same as that was in case of byte stuffing. So, In this protocol what we do is, if we encounter 0 and fire consecutive 1 bits, an extra 0 is added after these bits. This extra stuffed bit is removed from the data by the receiver-The extra bit is added after one O followed by five I bits regardless of the values of the next bit. Also, as the sender side always known which sequence is data and which is flag it will only add this extra bit in the data sequence, not in the flag sequence. Data from upper layer Example = 0001111111001111101000 Stuffed Framesent 0001111101100111111001000 Trailer Flog Flag Header Extra 2 Bits Frome Recieved 0001111101100111111001000 Trailer Flog Flog Header Unstuffed 000 111 1111 00 1111101 000 Data to upper layer

Page 1/10

Byte stuffing: A byte, which has a predefined bit puttern is added to the dada section of the frame when there is a character with the same puttern as the flag. Whenever the receiver encounter the ESC (Escape character) character, it removes from the data section and treats the next character as data, not a flag.

But the problem arises when the text contains one on more escape character followed by a flag. To solve this problem, the escape characters that are part of the text are marked by another escape character.

i.e. If the escape character is part of the text, an extra one is added to show that the second one is part of the text.



Ques: 2: What is checksum? Explain it with suitable example.

Ans:- Checksum: Checksum is the everan detection method used by upper layer protocols and is considered to be more reliable the LRC, VRC, and CRC.

This methods makes the use of Checksum Generator on Sender side and Checksum Checker on Receiver

At the sender side, the data is divided into equal subunits of n bit length by the checksum generator. This bit is generally of 16-bit length. These submitted then added together using one's complement units are then added together using one's complement method. This sum is of n bits. The subultant bit is then complemented. This completemented sum which is then complemented. This completemented of original called checksum is appended to the end of original data unit and is then transmitted to Receiver.

The Receiver after oreceiving data t checksum passes it to check sum checker. Checksum checker divides it to check sum checker. Checksum checker divides this data units into various subunits of equal length and adds all these subunits. These subunits also and adds all these sub units. These subunits. The resultant contain checksum as one of the submunits. The resultant bit is then complemented. If the complemented result is zero, it means the data is evolve free, If the result is zero, it means the data contains an evolve and is non zero it means the data contains an evolve and

Ex: If the data unit to be transmitted is 10101001001

11001, the following procedure is used at sender

site and Reciver site.

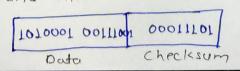
10101001 subunit 1

00111001 subunit 2

11100010 sum (using 1s complement)

00011101 Checksum (complement of sum)

Octa transmitted to Receiver is-



Receiver Site:

10101001 Subunit 1 00111001 Subunit 2 00011101 Checksum

11111111 SAW

00000000 sum's complement

Result is zero, it means no evolon.

Advantage: - The checksum detects all the evirous involving an odd number of bits as well as the eviron involving an even number of bits.

Disadvantage: The main problem is that the evolor goes undetected if one on more bits of a subunits is damaged and corresponding bit on bits of a subunit are damaged and bit on bits of a subunit are also domaged. Value in second subunit are also domaged. This is because the sum of those columns remains unchanged.

Aurs: 3: What is HDLC? Explain the trans format of HDLC with neat and clean diagram.

Ans:- HDLC (High-Level Data Link Control):- High-level Data Link
Control is basically bit-

bit stuffing to achieve data transparency over very point to point and multipoint links in Data Link Lagur (DLL).

Transparency is basically separation of data from control signals. HDLC was being derived from Synchronous Data Link Control (SDLC).

HDLC is most important and essential protocol in data link layer. Duta is also engonized and divided into small units also known as duto frames and is transerved across network to destination that vovides and ensure its successful arrival.

There are three types of HDLC frames:

- -> Information frames/User data (I-frames)
- Supervisory frames/Control data (s-frames)
- + Unnumbered fromes (U-fromes)

HDLC Frame :-

HDLC is a bit-ordented protocol where each frame contains up to six fields. The structure varies according to the type of frame.

The fields of a HDLC frome are-

HDLC Frame

Flog	Address	Control	Pay load	FCS	Flag
1 byte	lbyte	1 by te	vou'a ble	2 on 4 bytes	1 byte
(01111110)				(0	177777

Page 5/10

- Flag-It is an 8-bit sequence that masks the beginning and the end of the frome. The bit pattern of the flag is 01111110.
- Address: It contains the address of the succeiver. If the frame is sent by the primary station, it contains the address (es) of the secondary station, it contains it is sent by the secondary station, it contains the address of the primary station. The address field may be from 1 byte to several bytes.
 - Control: It is 1 or 2 bytes containing flow and evolone Control information.
 - Payload: This causies the data from the network logur.

 Its length may vary from one network to another.
 - FCS:- It is a 2 byte on 4 bytes frame check sequence for evolor detection. The standard code used is CRC Ccyclic Redundancy code).

Ques: 4: What is clossful addressing scheme? Find the class of tollowing:

(A) find the class of each address with their subnet mask.

(d) 00000001 00001011 00001011 111011111

Ans: Classful Addressing: IP address is an address having information about how to reach a specific host, especially outside the LAN. An IP address is a 32 bit unique address having an address space

of 232.

Grenerally, there are two notation in which IP address is written, dotted decimal notation and hexadecimal notation

Dotted Decimal Notation:

120-11.3.31

Hexadecimal Notation:

DX751D95EA

Classful

The 32 bit IP address is divided into five subclasses. These we following:

+ Class A

+ Class B

+ Class C

d 22 pl) +

+ Class E

(d) 000000001 0000101T 0000101T 11T01171

1.11.11.239

This is from class A address and the subnet mask of this 255-0.0.0

(b) 110000001 10000011 00011011 11111111 193-131-27-255

This is from class C address and the subnet mask of the class C is 255.255.25.0

(C) 10100111 11011011 1 0001011 0110 11111 167-219-139-109

This is from class B address and the subnet mosk of class B is 255, 255, 0.0

(d) 11110011 10011011 11111011 00001111 243.155.251-15

This is from class D address and the subnet mask of class D is 255.255.255.255.

Ques:-5 An address in a block is given as 73.22.17.25. Find the number of address in the block, the First address and the last address.

Ans: - Given IP: 73.22.17.25

1- This address is belongs to class A as it lies in range 0-127.

2- As we know, In class the subnet mask is 8 n=8
73.22.17.25/8

Number of Adobress: N= 232-n

Here n=8, therefore

 $2^{32-8} = 2^{24} \Rightarrow 16777216$

First address: To get first address, we add IP address to subnet mask of that class IP address.

Subnet mosk for class A is: 255.0.0.0

· Add 73.22.17-25 with 255.0.0.0 perform bitwise and operation

73.22.17.25

• Finst address is 73.0.0.018 or 73.0.0.0

Last address is - 73.255.255.255

2) An address in a block is given as 172.16.5.1. Find the number of address in the block, the first address and the last address.

Ans: - Given IP = 172.16.5.1

7 This address belongs to class B because it lies blw 120-191

Number of Address: N= 232-n

 $= 2^{32-16} = 2^{16} = 65536$

So, Number of address in this b

First address: Subnet mask class B is 255.255.0-0

172-16-5-1
255-255.0.0

First IP adobuss is 172.16.0.0

Last address:

So, Last address 1s + 172.16.255.255

B) An address in a block is given as 192. 168. 5.1. Find the number of address in the block, the first address and the last address.

Ans: - Given IP address is 192.168.5.1 + This belongs to class C, because it has lies between 192-233.

-> Défault subnet value for classe 15 24 n= 24

So, we can easily find number of address in this block

Number of Address:

N= 232-n

Here n=24

 $2^{32-24} = 2^{8} = 256$

Total no. of address At 256

First IP address: Default subnet mask for class Cis 255.255.255.0

> 192.160.5.1 255.255.255.0

First Address 1's = 192,168,50

Last Address: Last address is = 192.160.5.255

As last address is used as broadcast network, so

Mactically last address is 192.168.5.254 Page 10/10