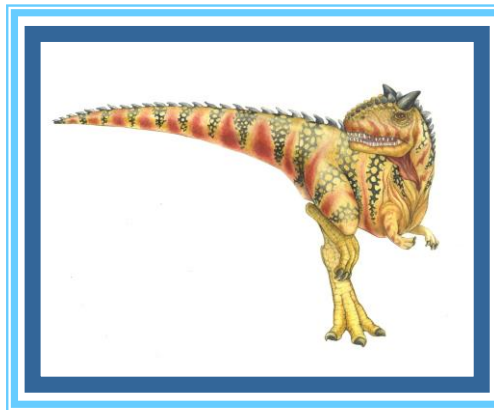


# Chapter 8: Main Memory

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# Chapter 8: Memory Management

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- Background
- Swapping
- Contiguous Memory Allocation
- Segmentation
- Paging
- Structure of the Page Table
- Example: The Intel 32 and 64-bit Architectures
- Example: ARM Architecture





# Objectives

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- To provide a detailed description of various ways of organizing memory hardware
- To discuss various memory-management techniques, including paging and segmentation
- To provide a detailed description of the Intel Pentium, which supports both pure segmentation and segmentation with paging





# Background

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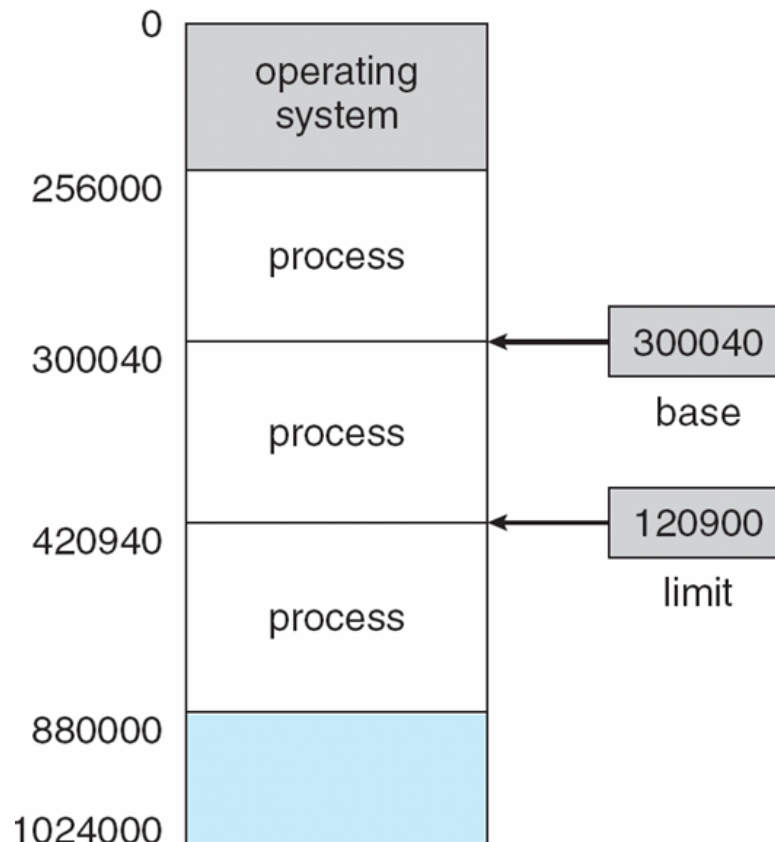
- Program must be brought (from disk) into memory and placed within a process for it to be run
- Main memory and registers are only storage CPU can access directly
- Memory unit only sees a stream of addresses + read requests, or address + data and write requests
- Register access in one CPU clock (or less)
- Main memory can take many cycles, causing a **stall**
- **Cache** sits between main memory and CPU registers
- Protection of memory required to ensure correct operation





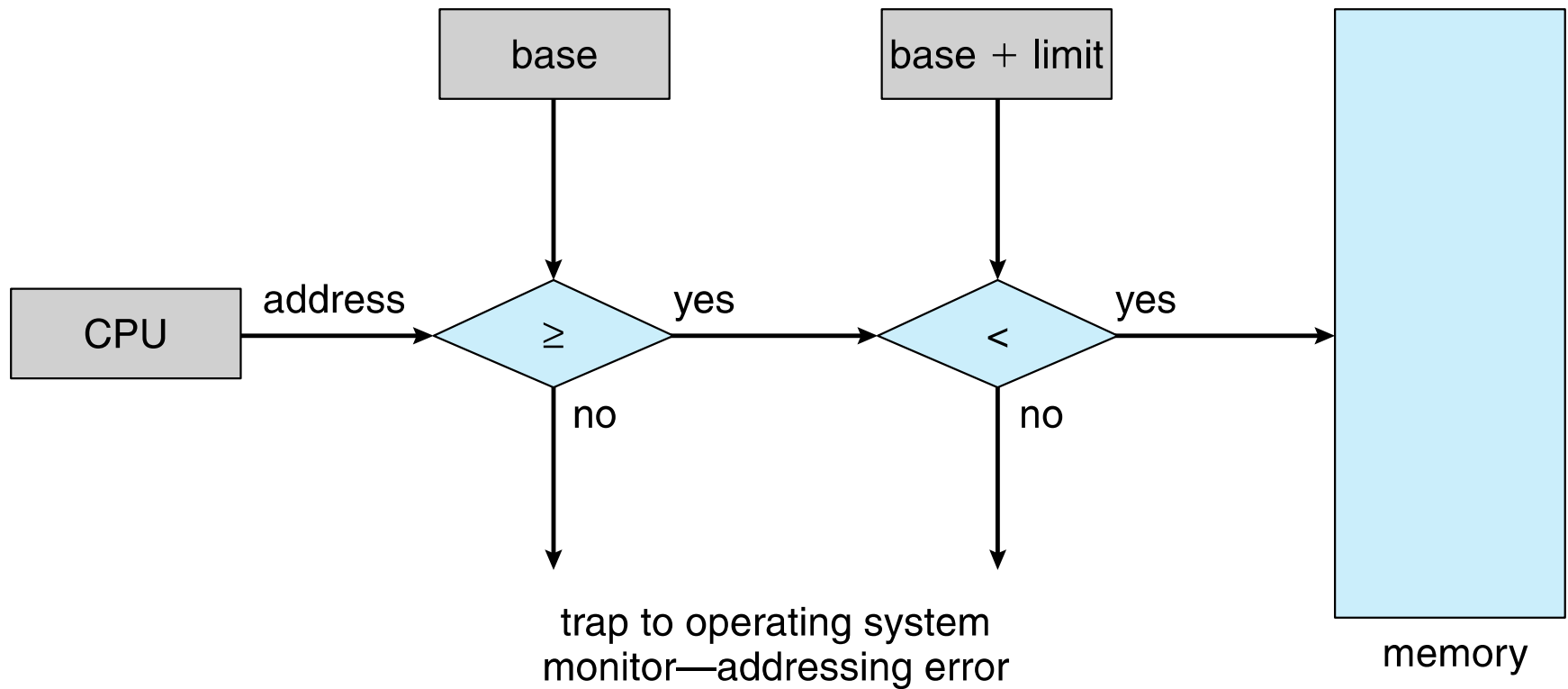
# Base and Limit Registers

- A pair of **base** and **limit registers** define the physical address space corresponds to logical address space of a process
- CPU must check every memory access generated in user mode to be sure it is between base and limit for that user





# Hardware Address Protection





# Address Binding

- Programs on disk, ready to be brought into memory to execute form an **input queue**
  - Without support, must be loaded into address 0000
  
- Further, addresses represented in different ways at different stages of a program's life
  - Source code addresses usually symbolic
  - Compiled code addresses **bind** to relocatable addresses
    - ▶ i.e. “14 bytes from beginning of this module”
  - Linker or loader will bind relocatable addresses to absolute addresses
    - ▶ i.e. 74014
  - Each binding maps one address space to another





# Binding of Instructions and Data to Memory

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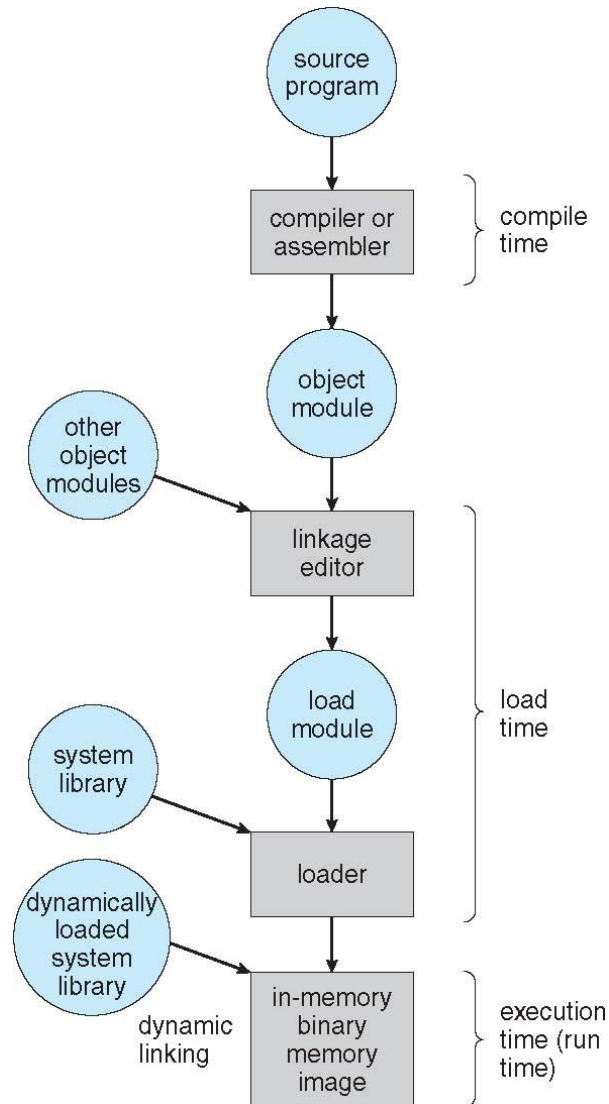
- Address binding of instructions and data to memory addresses can happen at three different stages
  - **Compile time:** If memory location known a priori, **absolute code** can be generated; must recompile code if starting location changes
  - **Load time:** Compiler generate **relocatable code** if memory location is not known at compile time. Loader generate absolute code from relocatable code.
  - **Execution time:** Binding delayed until run time if the process can be moved during its execution from one memory segment to another
    - ▶ Need hardware support for address maps (e.g., base and limit registers)







# Multistep Processing of a User Program





# Logical vs. Physical Address Space

---

- The concept of a logical address space that is bound to a separate **physical address space** is central to proper memory management
  - **Logical address** – generated by the CPU; also referred to as **virtual address**
  - **Physical address** – address seen by the memory unit
- Logical and physical addresses are the same in compile-time and load-time address-binding schemes; logical (virtual) and physical addresses differ in execution-time address-binding scheme
- **Logical address space** is the set of all logical addresses generated by a program
- **Physical address space** is the set of all physical addresses generated by a program





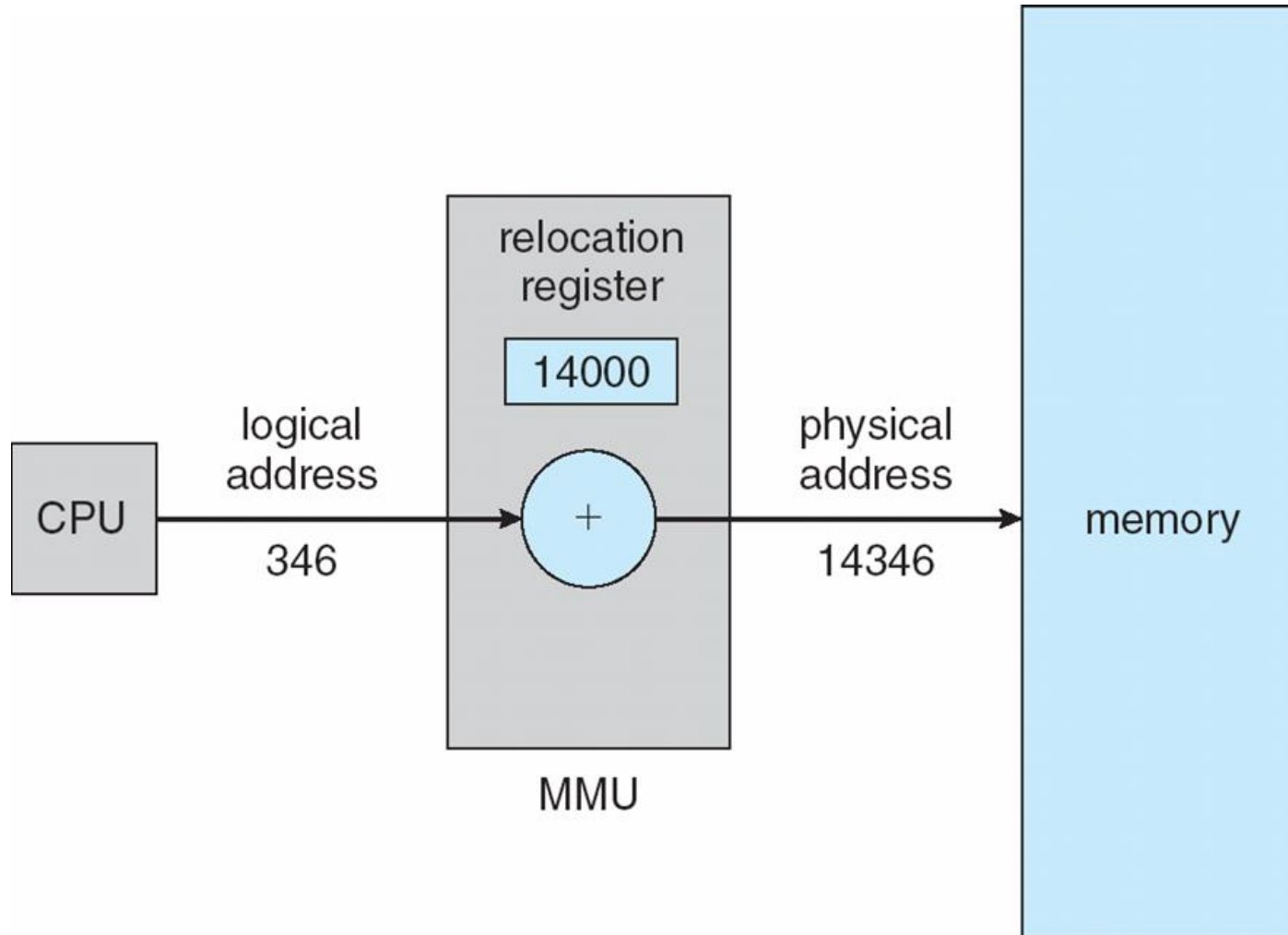
# Memory-Management Unit (MMU)

- Hardware device that at run time maps virtual to physical address
- To start, consider simple scheme where the value in the relocation register is added to every address generated by a user process at the time it is sent to memory
  - Base register now called **relocation register**
  - MS-DOS on Intel 80x86 used 4 relocation registers
- The user program deals with *logical* addresses; it never sees the *real* physical addresses
  - Execution-time binding occurs when reference is made to location in memory
  - Logical address bound to physical addresses





# Memory-Management Unit (MMU)





# Dynamic Loading

---

- Routine is not loaded until it is called
- Better memory-space utilization; unused routine is never loaded
- All routines kept on disk in relocatable load format
- Useful when large amounts of code are needed to handle infrequently occurring cases
- No special support from the operating system is required
  - Implemented through program design
  - OS can help by providing libraries to implement dynamic loading





# Dynamic Linking & Shared Libraries

---

- **Static linking** – system libraries and program code combined by the loader into the binary program image
- Dynamic linking –linking postponed until execution time
- Small piece of code, **stub**, used to locate the appropriate memory-resident library routine
- Stub replaces itself with the address of the routine, and executes the routine
- Operating system checks if routine is in processes' memory address
  - If not in address space, add to address space
- Dynamic linking is particularly useful for libraries
- System also known as **shared libraries**
- Consider applicability to patching system libraries
  - Versioning may be needed





# Swapping

- A process can be **swapped** temporarily out of memory to a backing store, and then brought back into memory for continued execution
  - Total physical memory space of processes can exceed physical memory
- **Backing store** – fast disk large enough to accommodate copies of all memory images for all users; must provide direct access to these memory images
- **Roll out, roll in** – swapping variant used for priority-based scheduling algorithms; lower-priority process is swapped out so higher-priority process can be loaded and executed
- Major part of swap time is transfer time; total transfer time is directly proportional to the amount of memory swapped
- System maintains a **ready queue** of ready-to-run processes which have memory images on disk





# Swapping (Cont.)

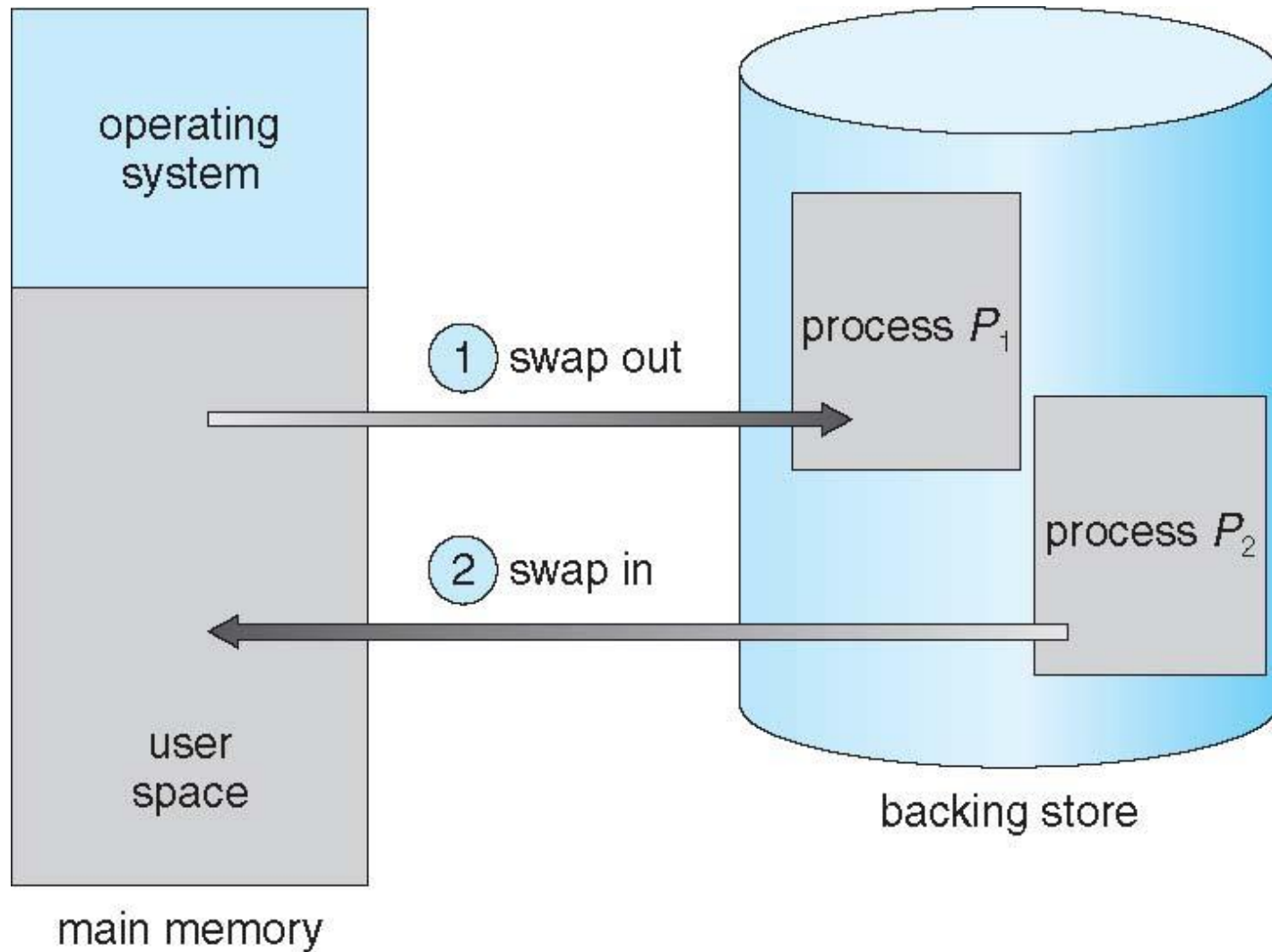
- Does the swapped out process need to swap back in to same physical addresses?
- Depends on address binding method
  - Plus consider pending I/O to / from process memory space
- Modified versions of swapping are found on many systems (i.e., UNIX, Linux, and Windows)
  - Swapping normally disabled
  - Started if more than threshold amount of memory allocated
  - Disabled again once memory demand reduced below threshold







# Schematic View of Swapping





# Context Switch Time including Swapping

---

- If next processes to be put on CPU is not in memory, need to swap out a process and swap in target process
- Context switch time can then be very high
- 100MB process swapping to hard disk with transfer rate of 50MB/sec
  - Swap out time of 2000 ms
  - Plus swap in of same sized process
  - Total context switch swapping component time of 4000ms (4 seconds)
- Can reduce if reduce size of memory swapped – by knowing how much memory really being used
  - System calls to inform OS of memory use via `request_memory()` and `release_memory()`





# Context Switch Time and Swapping (Cont.)

---

- Other constraints as well on swapping
  - Pending I/O – can't swap out as I/O would occur to wrong process
  - Or always transfer I/O to kernel space, then to I/O device
    - ▶ Known as **double buffering**, adds overhead
- Standard swapping not used in modern operating systems
  - But modified version common
    - ▶ Swap only when free memory extremely low





# Swapping on Mobile Systems

---

- Not typically supported
  - Flash memory based
    - ▶ Small amount of space
    - ▶ Limited number of write cycles
    - ▶ Poor throughput between flash memory and CPU on mobile platform
- Instead use other methods to free memory if low
  - iOS **asks** apps to voluntarily relinquish allocated memory
    - ▶ Read-only data thrown out and reloaded from flash if needed
    - ▶ Failure to free can result in termination
  - Android terminates apps if low free memory, but first writes **application state** to flash for fast restart
  - Both OSes support paging as discussed below





# Contiguous Allocation

---

- Main memory must support both OS and user processes
- Limited resource, must allocate efficiently
- Contiguous allocation is one of the early methods
- Main memory usually divided into two **partitions**:
  - Resident operating system, usually held in low memory
  - User processes then held in high memory
  - Each process contained in single contiguous section of memory





# Contiguous Allocation (Cont.)

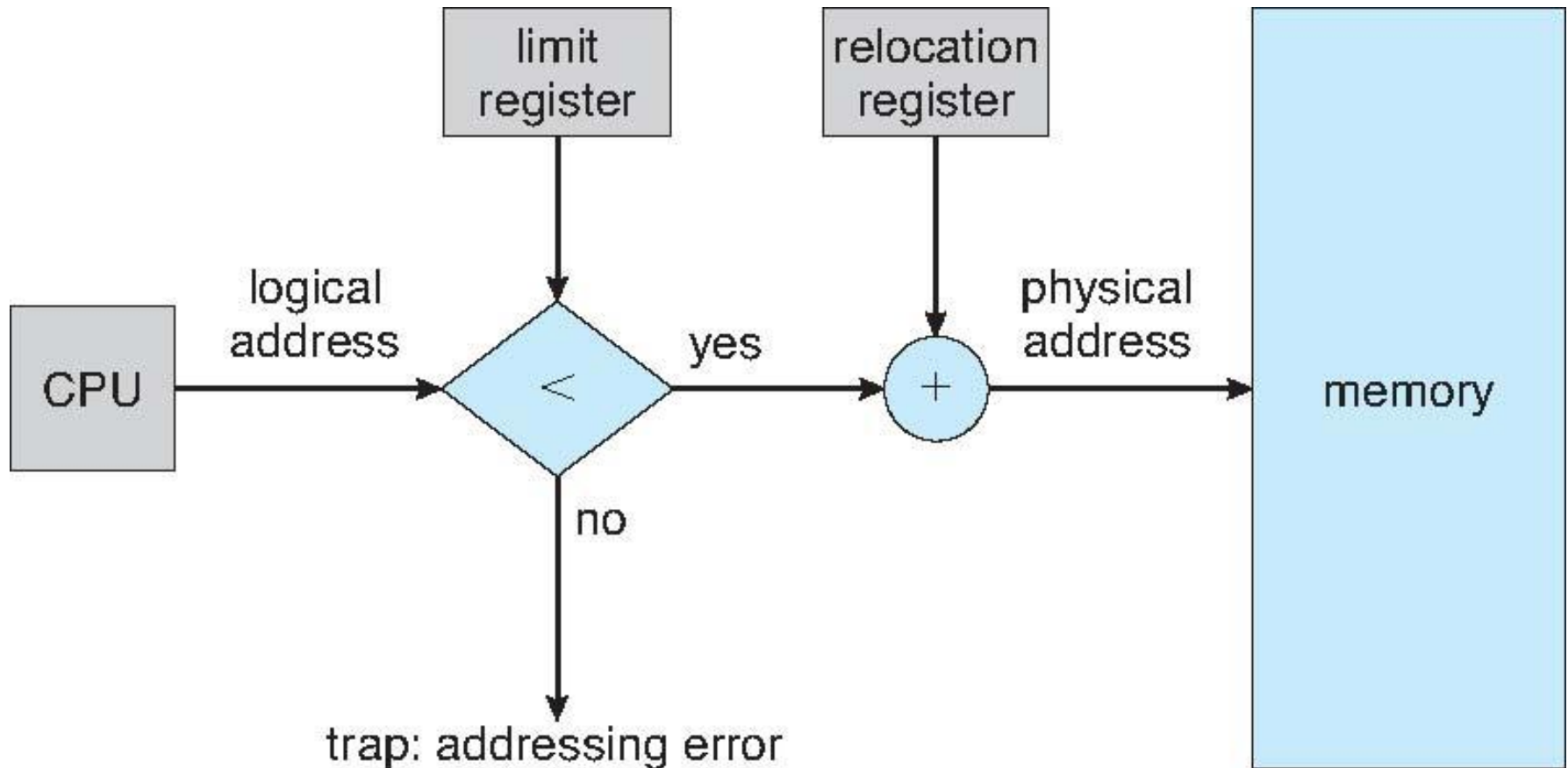
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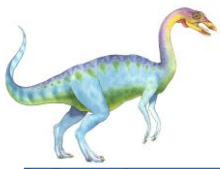
- Relocation registers used to protect user processes from each other, and from changing operating-system code and data
  - Base register contains value of smallest physical address
  - Limit register contains range of logical addresses – each logical address must be less than the limit register
  - MMU maps logical address *dynamically*
  - Can then allow actions such as kernel code being **transient** and kernel changing size





# Hardware Support for Relocation and Limit Registers

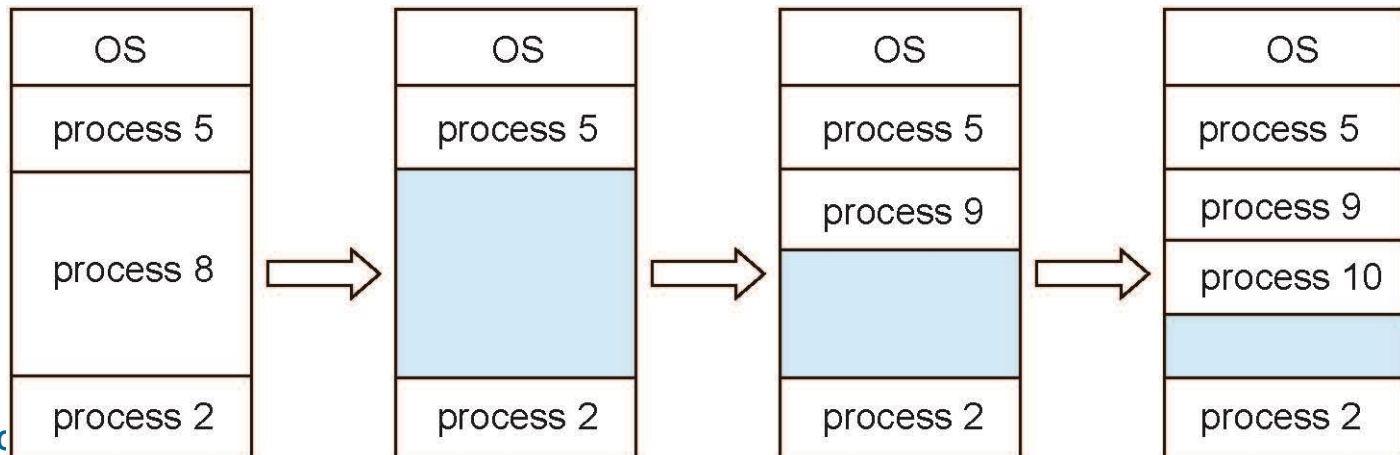




# Multiple-partition allocation

## Multiple-partition allocation

- Degree of multiprogramming limited by number of partitions
- **Variable-partition** sizes for efficiency (sized to a given process' needs)
- **Hole** – block of available memory; holes of various size are scattered throughout memory
- When a process arrives, it is allocated memory from a hole large enough to accommodate it
- Process exiting frees its partition, adjacent free partitions combined
- Operating system maintains information about:  
a) allocated partitions    b) free partitions (hole)







# Dynamic Storage-Allocation Problem

How to satisfy a request of size  $n$  from a list of free holes?

- **First-fit:** Allocate the ***first*** hole that is big enough
- **Best-fit:** Allocate the ***smallest*** hole that is big enough; must search entire list, unless ordered by size
  - Produces the smallest leftover hole
- **Worst-fit:** Allocate the ***largest*** hole; must also search entire list
  - Produces the largest leftover hole

First-fit and best-fit are better than worst-fit in terms of speed and storage utilization





# Fragmentation

---

- **External Fragmentation** – total memory space exists to satisfy a request, but it is not contiguous
- **Internal Fragmentation** – allocated memory may be slightly larger than requested memory; this size difference is memory internal to a partition, but not being used
- First fit analysis reveals that given  $N$  blocks allocated,  $0.5 N$  blocks lost to fragmentation
  - $1/3$  may be unusable -> **50-percent rule**





# Fragmentation (Cont.)

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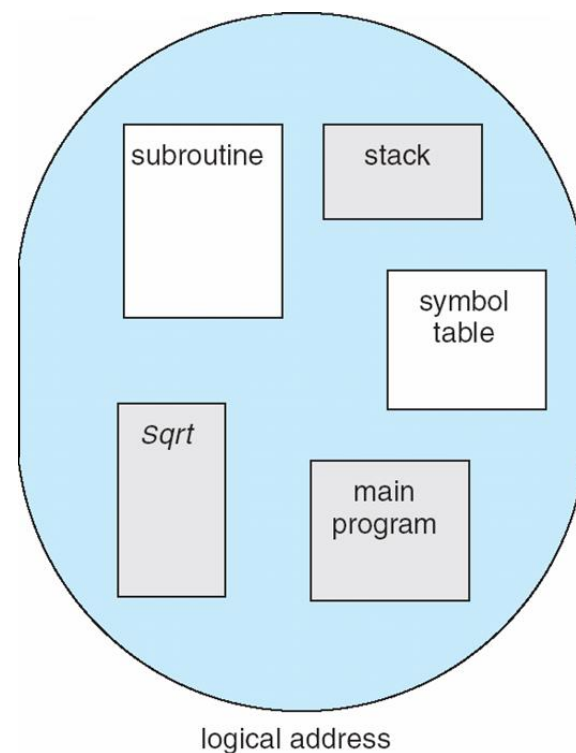
- Reduce external fragmentation by **compaction**
  - Shuffle memory contents to place all free memory together in one large block
  - Compaction is possible *only* if relocation is dynamic, and is done at execution time
  - I/O problem
    - ▶ Latch job in memory while it is involved in I/O
    - ▶ Do I/O only into OS buffers
- Now consider that backing store has same fragmentation problems





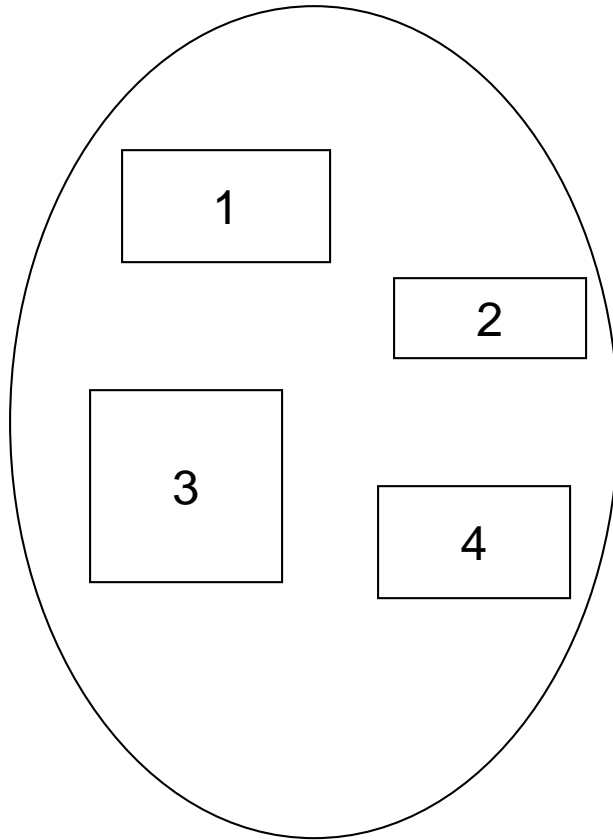
# Segmentation

- Memory-management scheme that supports user view of memory
- A program is a collection of segments
  - A segment is a logical unit such as:
    - main program
    - procedure
    - function
    - method
    - object
    - local variables, global variables
    - common block
    - stack
    - symbol table
    - arrays

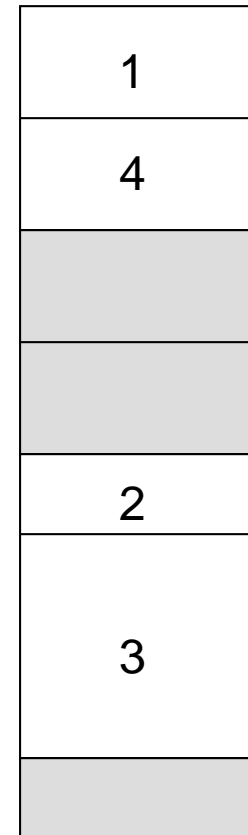




# Logical View of Segmentation



user space



physical memory space





# Segmentation Architecture

---

- Logical address consists of a two tuple:  
    <segment-number, offset>,
- **Segment table** – maps two-dimensional logical addresses; each table entry has:
  - **base** – contains the starting physical address where the segments reside in memory
  - **limit** – specifies the length of the segment
- **Segment-table base register (STBR)** points to the segment table's location in memory
- **Segment-table length register (STLR)** indicates number of segments used by a program;  
    segment number **s** is legal if **s** < **STLR**





# Segmentation Architecture (Cont.)

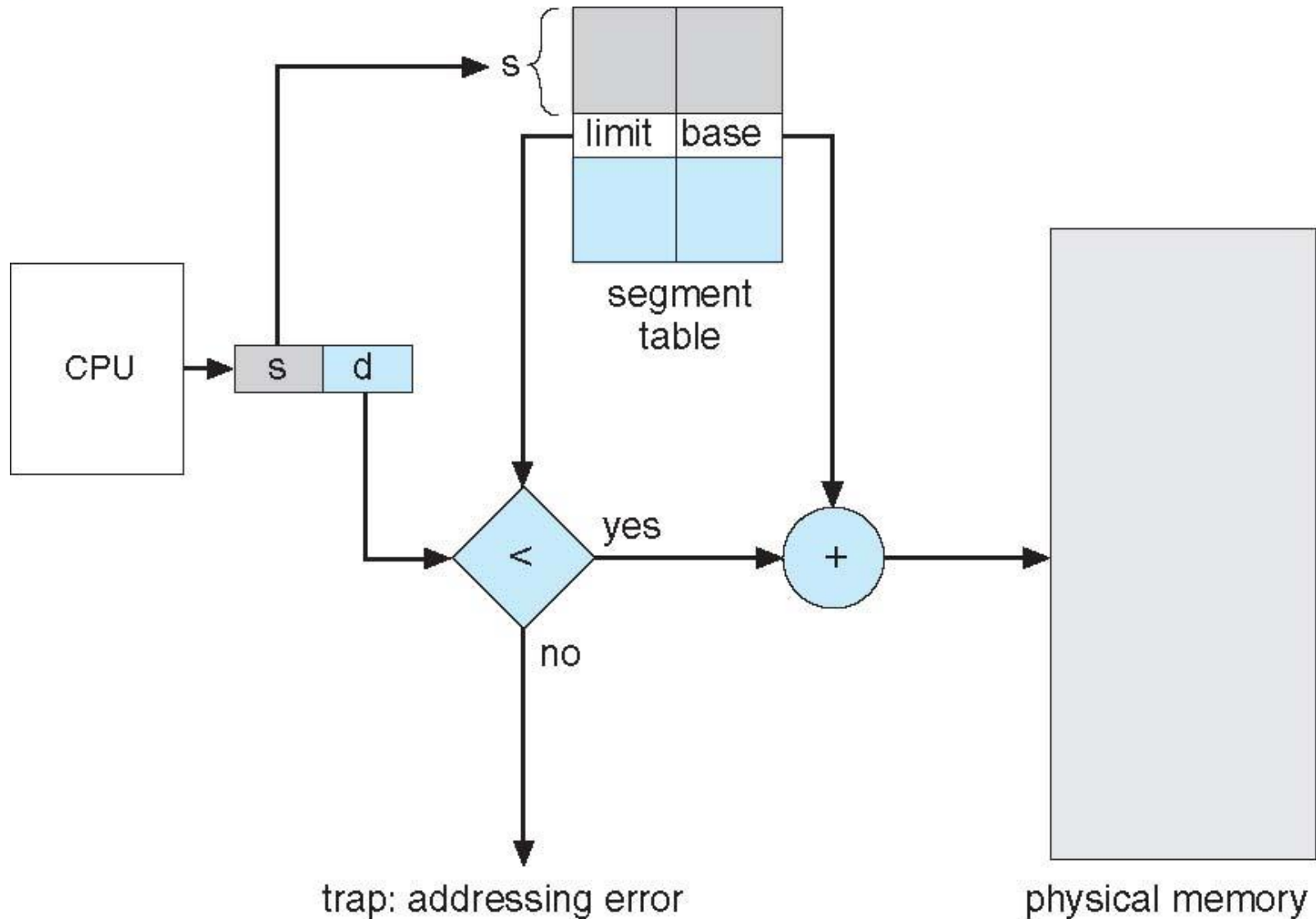
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- Protection
  - With each entry in segment table associate:
    - ▶ validation bit = 0  $\Rightarrow$  illegal segment
    - ▶ read/write/execute privileges
- Protection bits associated with segments; code sharing occurs at segment level
- Since segments vary in length, memory allocation is a dynamic storage-allocation problem





# Segmentation Hardware







# Paging

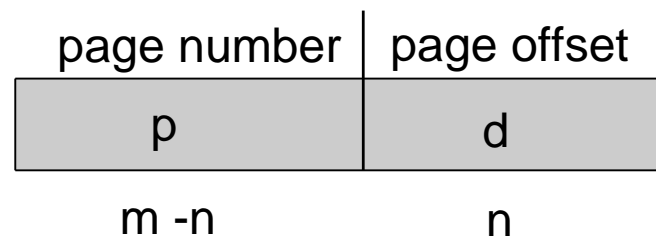
- Physical address space of a process can be noncontiguous; process is allocated physical memory whenever the latter is available
  - Avoids external fragmentation
  - Avoids problem of varying sized memory chunks
- Divide physical memory into fixed-sized blocks called **frames**
  - Size is power of 2, between 512 bytes and 16 Mbytes
- Divide logical memory into blocks of same size called **pages**
- Keep track of all free frames
- To run a program of size  **$N$**  pages, need to find  **$N$**  free frames and load program
- Set up a **page table** to translate logical to physical addresses
- Backing store likewise split into pages
- Still have Internal fragmentation





# Address Translation Scheme

- Address generated by CPU is divided into:
  - **Page number** ( $p$ ) – used as an index into a **page table** which contains base address of each page in physical memory
  - **Page offset** ( $d$ ) – combined with base address to define the physical memory address that is sent to the memory unit

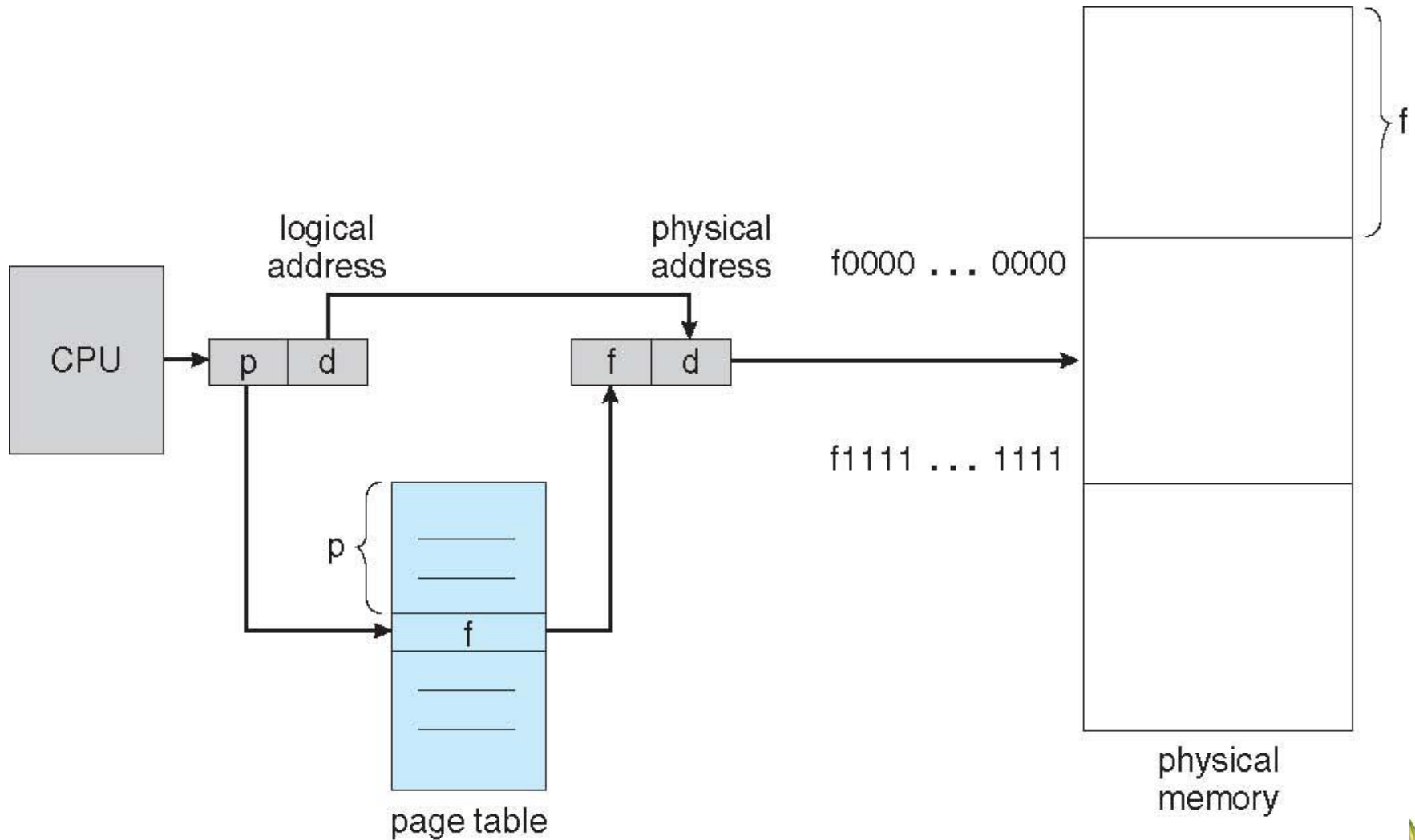


- For given logical address space  $2^m$  and page size  $2^n$



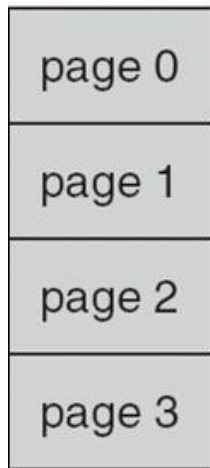


# Paging Hardware





# Paging Model of Logical and Physical Memory

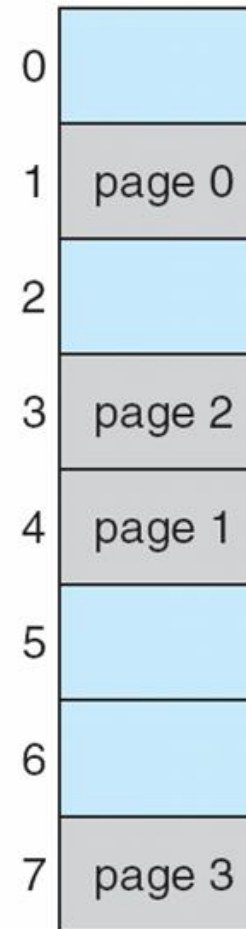


logical  
memory

0	1
1	4
2	3
3	7

page table

frame  
number



physical  
memory





# Paging Example

$n=2$  and  $m=4$

32-byte memory  
and 4-byte pages

0	a
1	b
2	c
3	d
4	e
5	f
6	g
7	h
8	i
9	j
10	k
11	l
12	m
13	n
14	o
15	p

logical memory

0	5
1	6
2	1
3	2

page table

0	
4	i j k l
8	m n o p
12	
16	
20	a b c d
24	e f g h
28	

physical memory



# Paging (Cont.)

- Calculating internal fragmentation
  - Page size = 2,048 bytes
  - Process size = 72,766 bytes
  - 35 pages + 1,086 bytes
  - Internal fragmentation of  $2,048 - 1,086 = 962$  bytes
  - Worst case fragmentation = 1 frame – 1 byte
  - On average fragmentation =  $1 / 2$  frame size
  - So small frame sizes desirable?
  - But each page table entry takes memory to track
  - Page sizes growing over time
    - ▶ Solaris supports two page sizes – 8 KB and 4 MB
- Process view and physical memory now very different
- By implementation process can only access its own memory

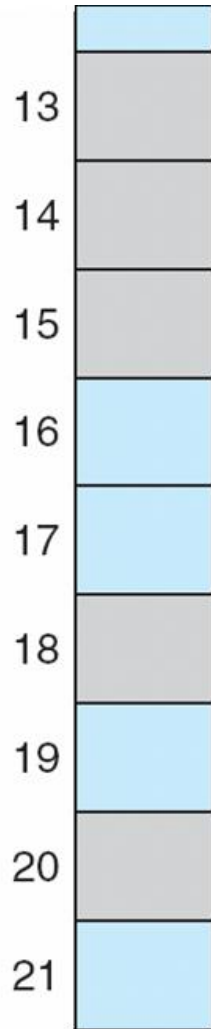
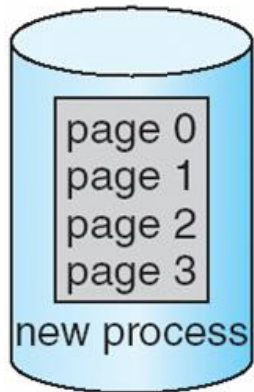




# Free Frames

free-frame list

14  
13  
18  
20  
15

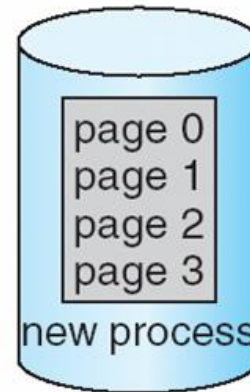


(a)

Before allocation

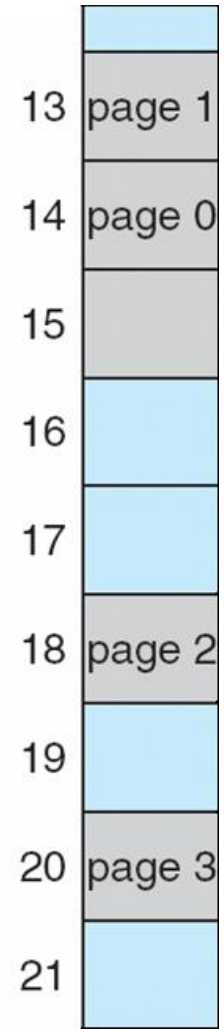
free-frame list

15



0	14
1	13
2	18
3	20

new-process page table



(b)

After allocation





# Implementation of Page Table

---

- Page table is kept in main memory
- **Page-table base register (PTBR)** points to the page table
- **Page-table length register (PTLR)** indicates size of the page table
- In this scheme every data/instruction access requires two memory accesses
  - One for the page table and one for the data / instruction
- The two memory access problem can be solved by the use of a special fast-lookup hardware cache called **associative memory** or **translation look-aside buffers (TLBs)**







# Implementation of Page Table (Cont.)

- Some TLBs store **address-space identifiers (ASIDs)** in each TLB entry – uniquely identifies each process to provide address-space protection for that process
  - Otherwise need to flush at every context switch
- TLBs typically small (64 to 1,024 entries)
- On a TLB miss, value is loaded into the TLB for faster access next time
  - Replacement policies must be considered
  - Some entries can be **wired down** for permanent fast access





# Associative Memory

## ■ Associative memory – parallel search

Page #	Frame #

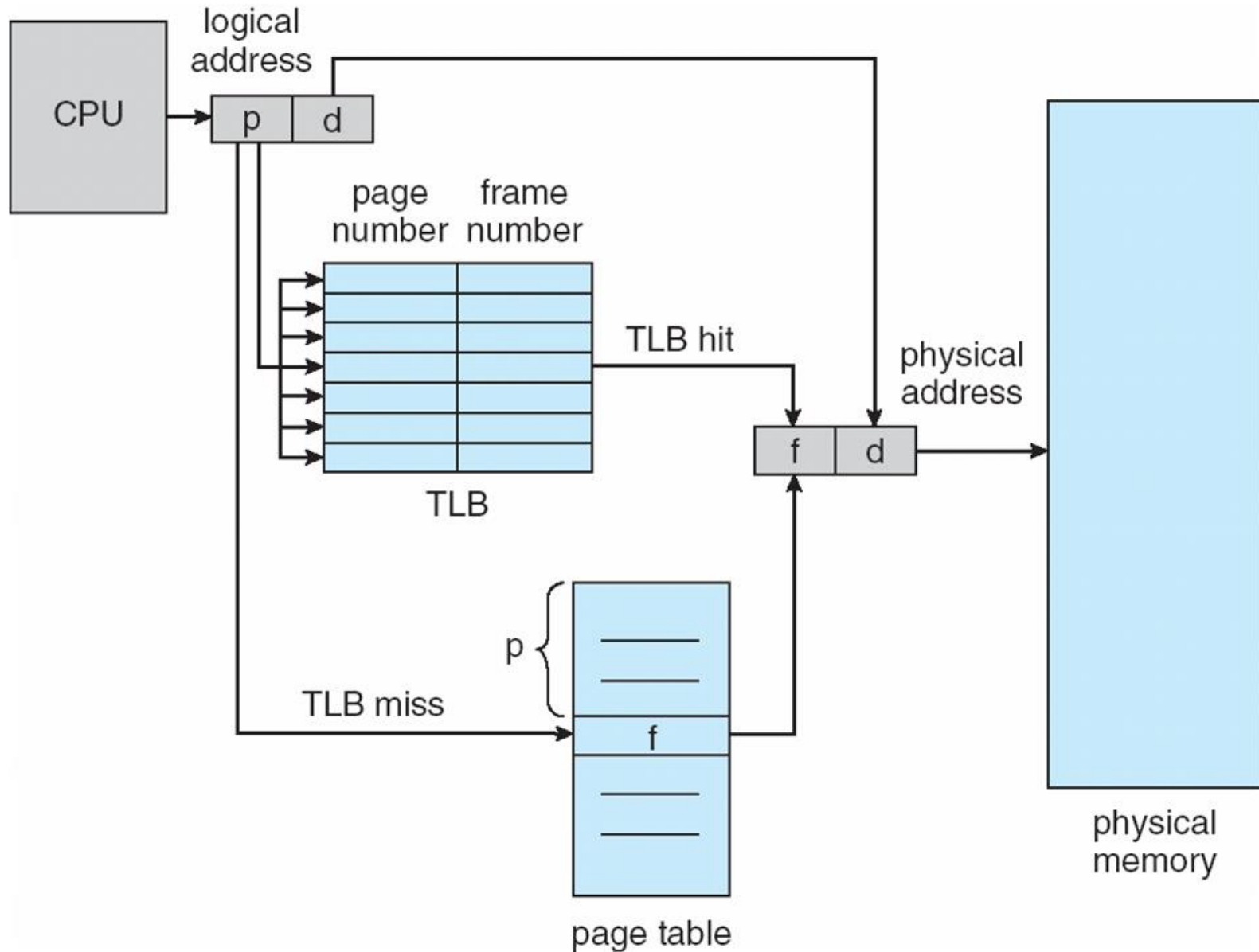
## ■ Address translation (p, d)

- If p is in associative register, get frame # out
- Otherwise get frame # from page table in memory





# Paging Hardware With TLB





# Effective Access Time

- Associative Lookup =  $\varepsilon$  time unit
  - Can be  $< 10\%$  of memory access time
- Hit ratio =  $\alpha$ 
  - Hit ratio – percentage of times that a page number is found in the associative registers; ratio related to number of associative registers
- Consider  $\alpha = 80\%$ ,  $\varepsilon = 20\text{ns}$  for TLB search,  $100\text{ns}$  for memory access
- **Effective Access Time (EAT)**

$$\begin{aligned} \text{EAT} &= (1 + \varepsilon) \alpha + (2 + \varepsilon)(1 - \alpha) \\ &= 2 + \varepsilon - \alpha \end{aligned}$$

- Consider  $\alpha = 80\%$ ,  $\varepsilon = 20\text{ns}$  for TLB search,  $100\text{ns}$  for memory access
  - $\text{EAT} = 0.80 \times 100 + 0.20 \times 200 = 120\text{ns}$
- Consider more realistic hit ratio  $\rightarrow \alpha = 99\%$ ,  $\varepsilon = 20\text{ns}$  for TLB search,  $100\text{ns}$  for memory access
  - $\text{EAT} = 0.99 \times 100 + 0.01 \times 200 = 101\text{ns}$





# Memory Protection

---

- Memory protection implemented by associating protection bit with each frame to indicate if read-only or read-write access is allowed
  - Can also add more bits to indicate page execute-only, and so on
- **Valid-invalid** bit attached to each entry in the page table:
  - “valid” indicates that the associated page is in the process’ logical address space, and is thus a legal page
  - “invalid” indicates that the page is not in the process’ logical address space
  - Or use **page-table length register (PTLR)**
- Any violations result in a trap to the kernel





# Valid (v) or Invalid (i) Bit In A Page Table

00000	page 0
	page 1
	page 2
	page 3
	page 4
10,468	page 5
12,287	

frame number		valid-invalid bit
0	2	v
1	3	v
2	4	v
3	7	v
4	8	v
5	9	v
6	0	i
7	0	i

page table

0	
1	
2	page 0
3	page 1
4	page 2
5	
6	
7	page 3
8	page 4
9	page 5
	⋮
	page <i>n</i>





# Shared Pages

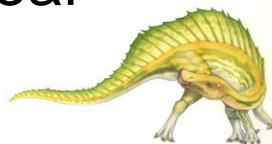
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## ■ Shared code

- One copy of read-only (**reentrant**) code shared among processes (i.e., text editors, compilers, window systems)
- Similar to multiple threads sharing the same process space
- Also useful for interprocess communication if sharing of read-write pages is allowed

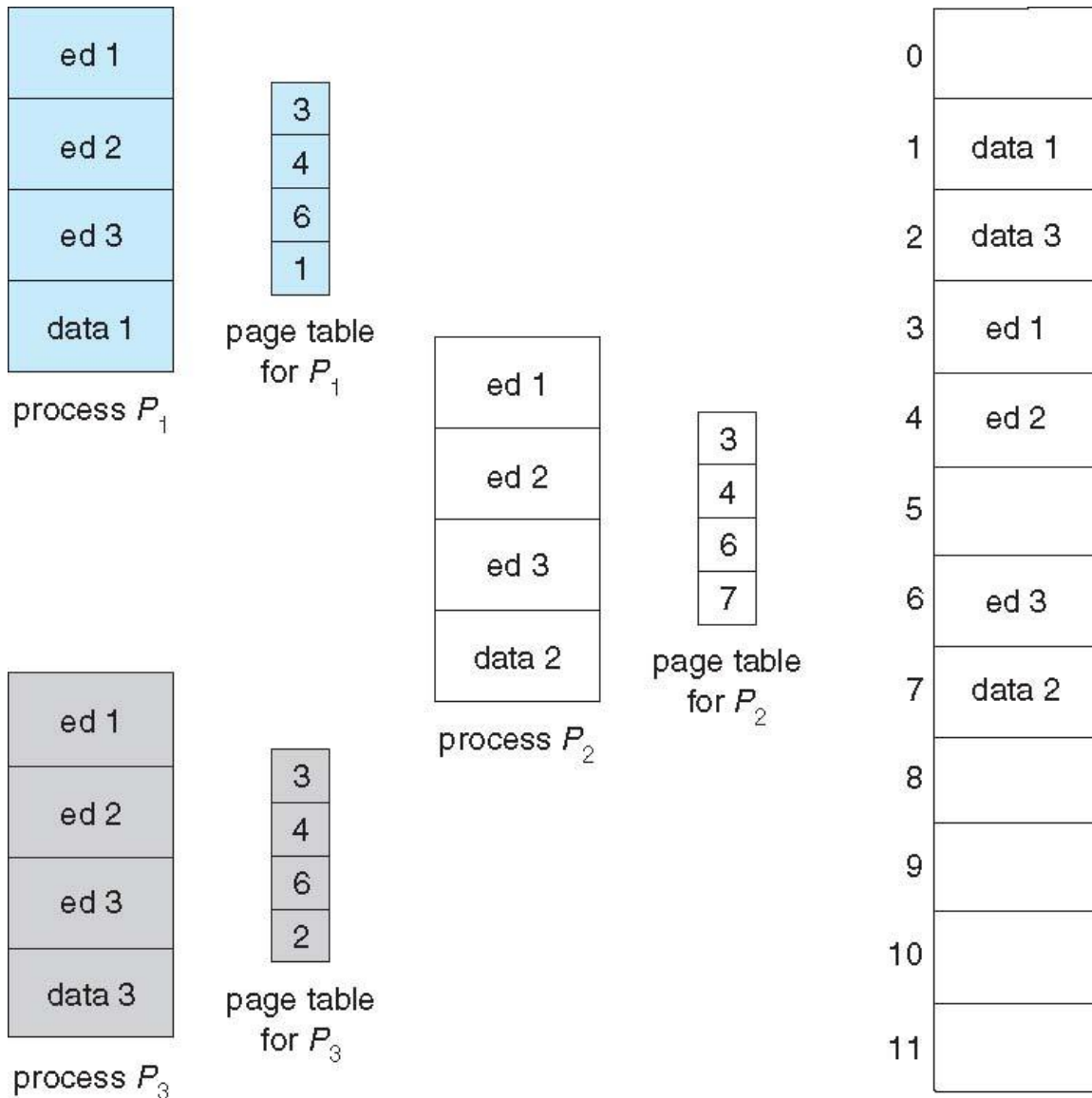
## ■ Private code and data

- Each process keeps a separate copy of the code and data
- The pages for the private code and data can appear anywhere in the logical address space





# Shared Pages Example







# Structure of the Page Table

- Memory structures for paging can get huge using straight-forward methods
  - Consider a 32-bit logical address space as on modern computers
  - Page size of 4 KB ( $2^{12}$ )
  - Page table would have 1 million entries ( $2^{32} / 2^{12}$ )
  - If each entry is 4 bytes -> 4 MB of physical address space / memory for page table alone
    - ▶ That amount of memory used to cost a lot
    - ▶ Don't want to allocate that contiguously in main memory
- Hierarchical Paging
- Hashed Page Tables
- Inverted Page Tables





# Hierarchical Page Tables

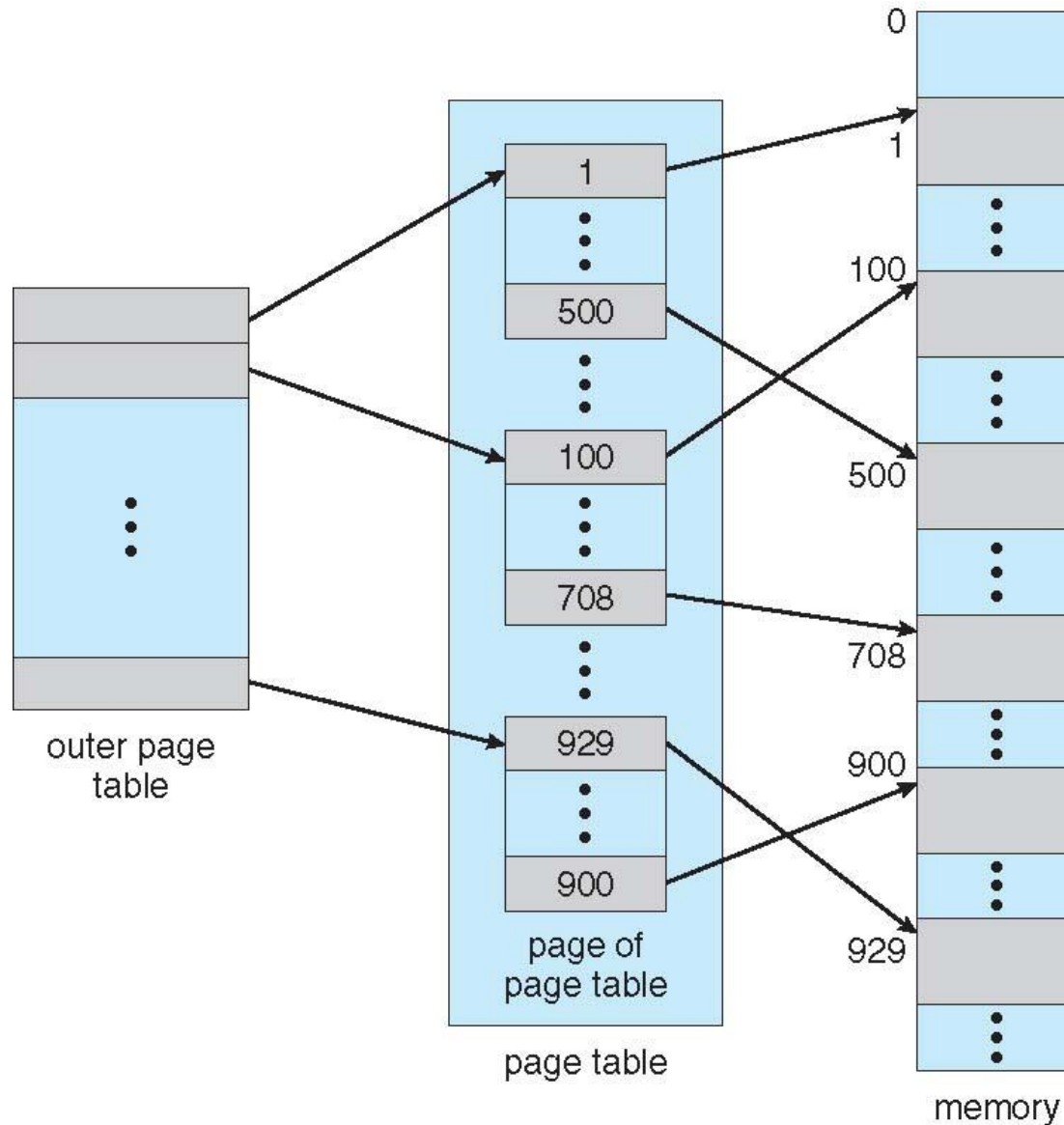
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- Break up the logical address space into multiple page tables
- A simple technique is a two-level page table
- We then page the page table





# Two-Level Page-Table Scheme





# Two-Level Paging Example

- A logical address (on 32-bit machine with 1K page size) is divided into:
  - a page number consisting of 22 bits
  - a page offset consisting of 10 bits
- Since the page table is paged, the page number is further divided into:
  - a 12-bit page number
  - a 10-bit page offset

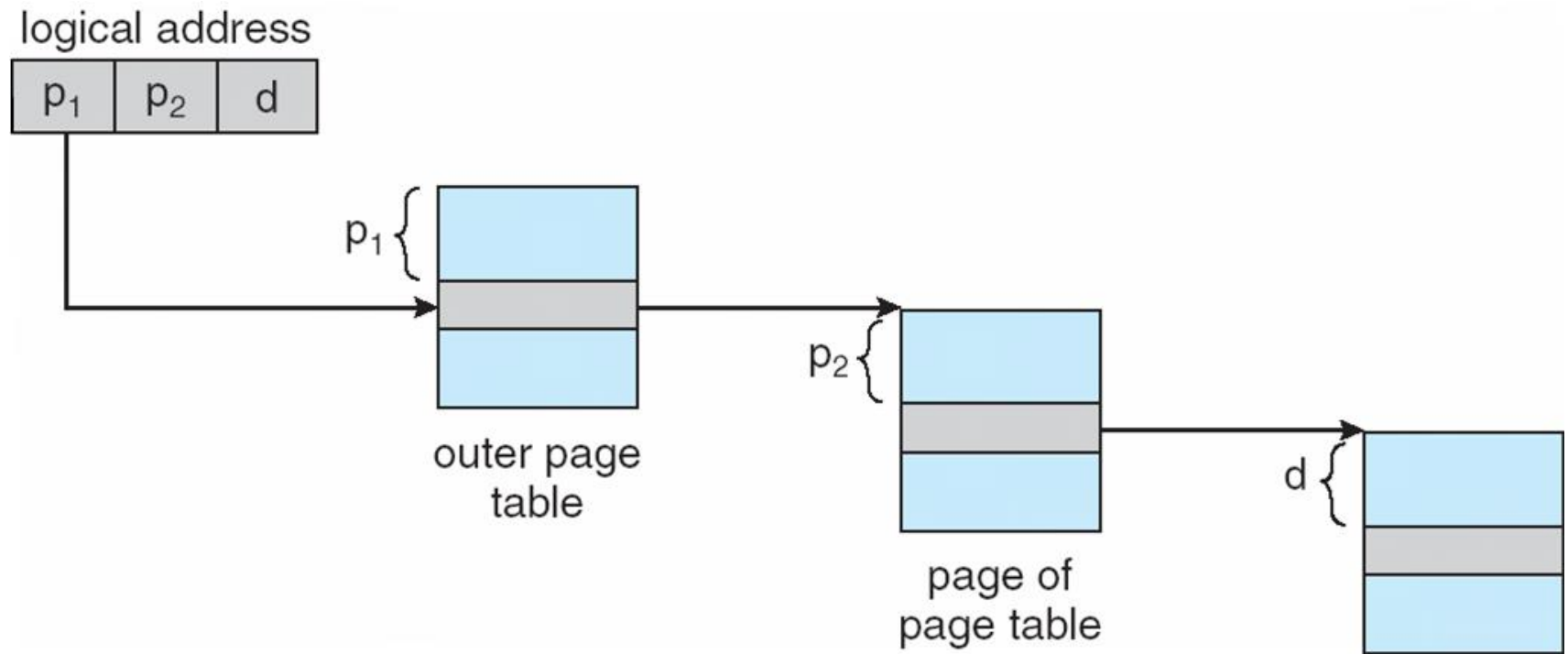
page number		page offset
$p_1$	$p_2$	$d$
12	10	10

- Thus, a logical address is as follows:
- where  $p_1$  is an index into the outer page table, and  $p_2$  is the displacement within the page of the inner page table
- Known as **forward-mapped page table**





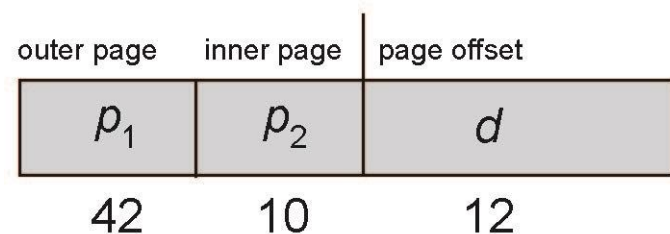
# Address-Translation Scheme





# 64-bit Logical Address Space

- Even two-level paging scheme not sufficient
- If page size is 4 KB ( $2^{12}$ )
  - Then page table has  $2^{52}$  entries
  - If two level scheme, inner page tables could be  $2^{10}$  4-byte entries
  - Address would look like



- Outer page table has  $2^{42}$  entries or  $2^{44}$  bytes
- One solution is to add a 2<sup>nd</sup> outer page table
- But in the following example the 2<sup>nd</sup> outer page table is still  $2^{34}$  bytes in size
  - ▶ And possibly 4 memory access to get to one physical memory location





# Three-level Paging Scheme

outer page	inner page	offset
$p_1$	$p_2$	$d$
42	10	12

2nd outer page	outer page	inner page	offset
$p_1$	$p_2$	$p_3$	$d$
32	10	10	12





# Hashed Page Tables

---

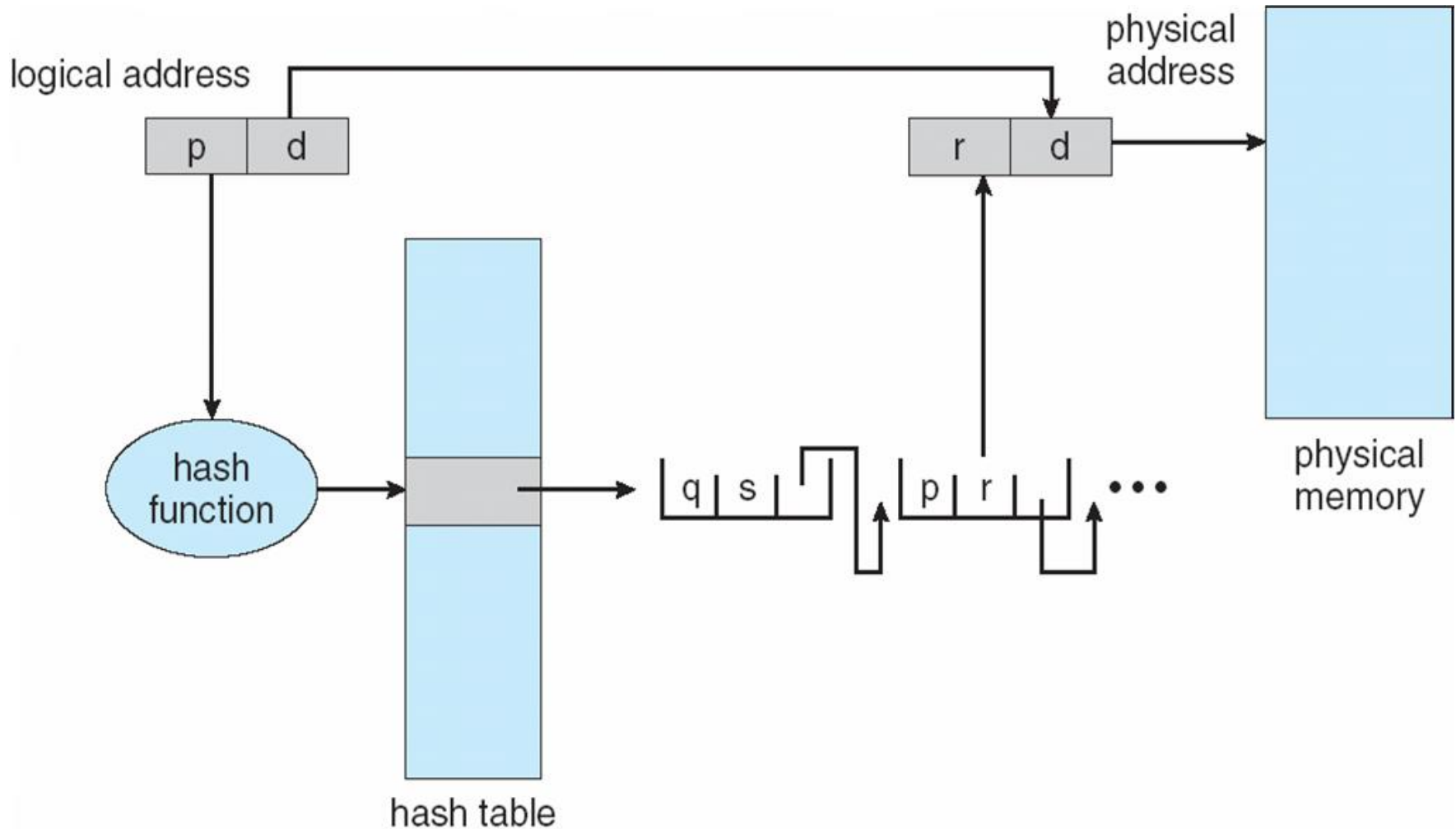
- Common in address spaces  $> 32$  bits
- The virtual page number is hashed into a page table
  - This page table contains a chain of elements hashing to the same location
- Each element contains (1) the virtual page number (2) the value of the mapped page frame (3) a pointer to the next element
- Virtual page numbers are compared in this chain searching for a match
  - If a match is found, the corresponding physical frame is extracted
- Variation for 64-bit addresses is **clustered page tables**
  - Similar to hashed but each entry refers to several pages (such as 16) rather than 1
  - Especially useful for **sparse** address spaces (where memory references are non-contiguous and scattered)







# Hashed Page Table





# Inverted Page Table

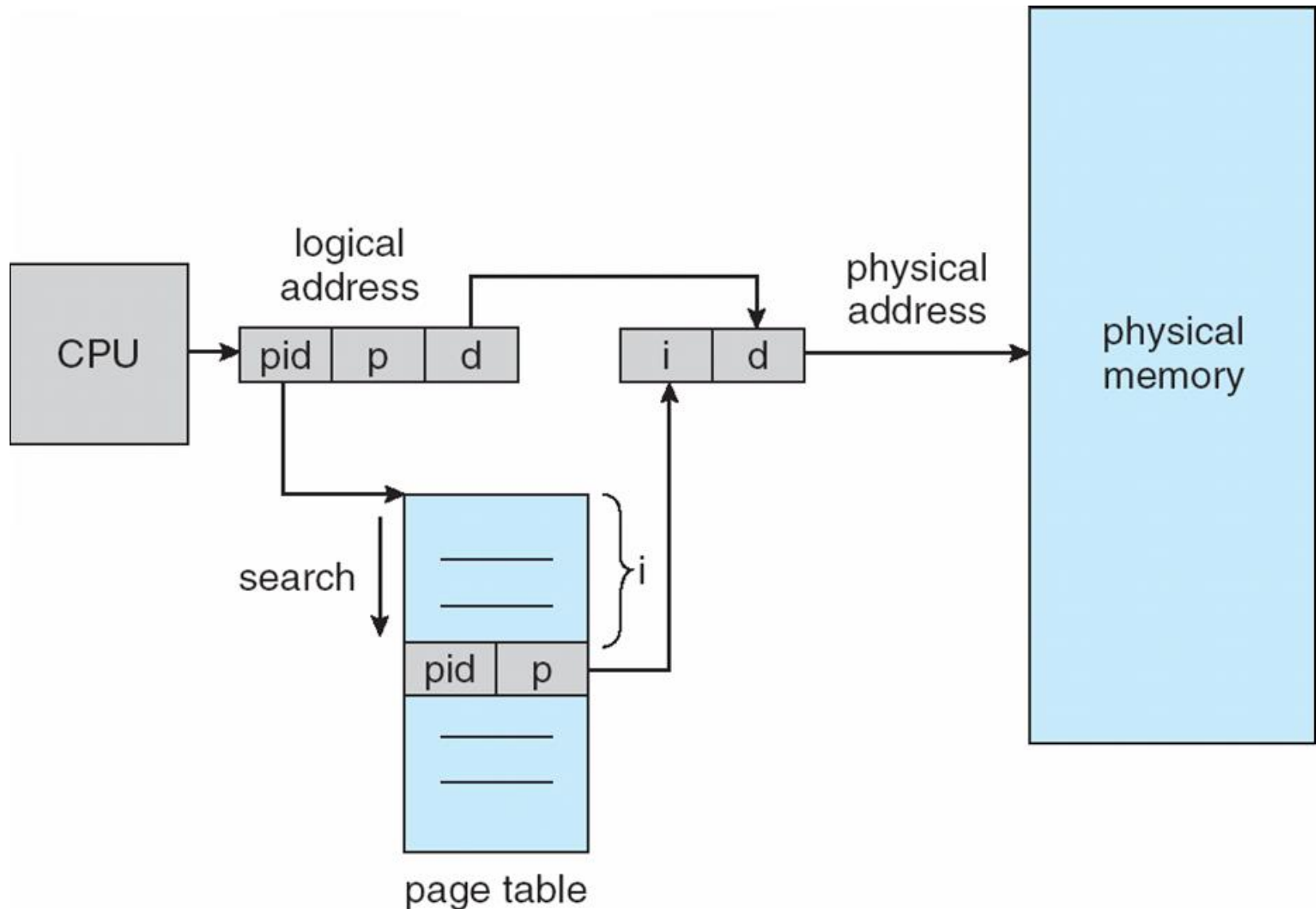
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- Rather than each process having a page table and keeping track of all possible logical pages, track all physical pages
- One entry for each real page of memory
- Entry consists of the virtual address of the page stored in that real memory location, with information about the process that owns that page
- Decreases memory needed to store each page table, but increases time needed to search the table when a page reference occurs
- Use hash table to limit the search to one — or at most a few — page-table entries
  - TLB can accelerate access
- But how to implement shared memory?
  - One mapping of a virtual address to the shared physical address





# Inverted Page Table Architecture





# Oracle SPARC Solaris

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- Consider modern, 64-bit operating system example with tightly integrated HW
  - Goals are efficiency, low overhead
- Based on hashing, but more complex
- Two hash tables
  - One kernel and one for all user processes
  - Each maps memory addresses from virtual to physical memory
  - Each entry represents a contiguous area of mapped virtual memory,
    - ▶ More efficient than having a separate hash-table entry for each page
  - Each entry has base address and span (indicating the number of pages the entry represents)





# Oracle SPARC Solaris (Cont.)

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- TLB holds translation table entries (TTEs) for fast hardware lookups
  - A cache of TTEs reside in a translation storage buffer (TSB)
    - ▶ Includes an entry per recently accessed page
- Virtual address reference causes TLB search
  - If miss, hardware walks the in-memory TSB looking for the TTE corresponding to the address
    - ▶ If match found, the CPU copies the TSB entry into the TLB and translation completes
    - ▶ If no match found, kernel interrupted to search the hash table
      - The kernel then creates a TTE from the appropriate hash table and stores it in the TSB, Interrupt handler returns control to the MMU, which completes the address translation.





## Example: The Intel 32 and 64-bit Architectures

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- Dominant industry chips
- Pentium CPUs are 32-bit and called IA-32 architecture
- Current Intel CPUs are 64-bit and called IA-64 architecture
- ARM architectures





# Example: The Intel IA-32 Architecture

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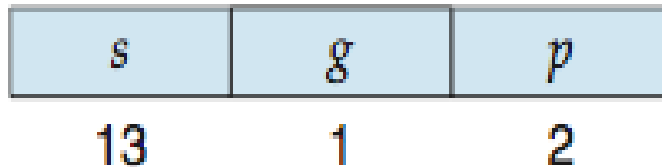
- Supports both segmentation and segmentation with paging
  - Each segment can be 4 GB
  - Up to 16 K segments per process
  - Divided into two partitions
    - ▶ First partition of up to 8 K segments are private to process (kept in **local descriptor table (LDT)**)
    - ▶ Second partition of up to 8K segments shared among all processes (kept in **global descriptor table (GDT)**)





## Example: The Intel IA-32 Architecture (Cont.)

- CPU generates logical address
  - Selector given to segmentation unit
    - ▶ Which produces linear addresses



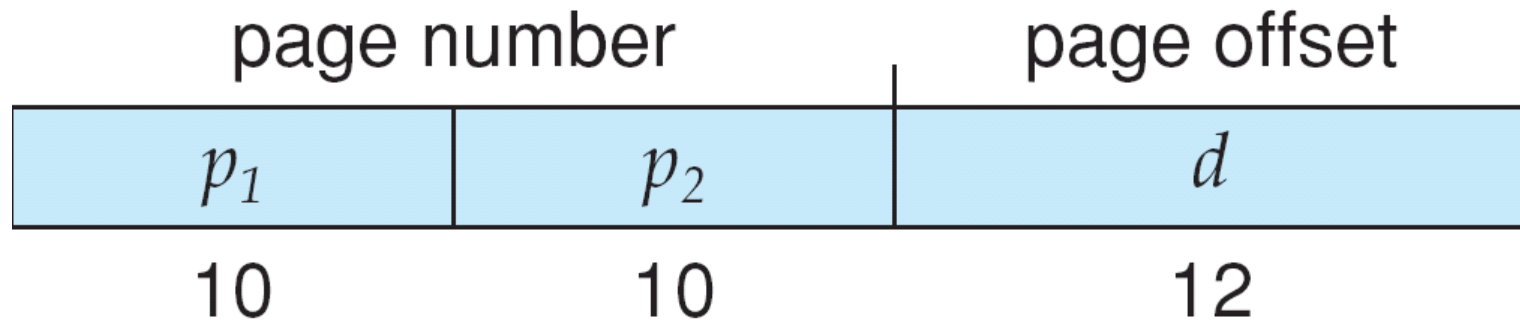
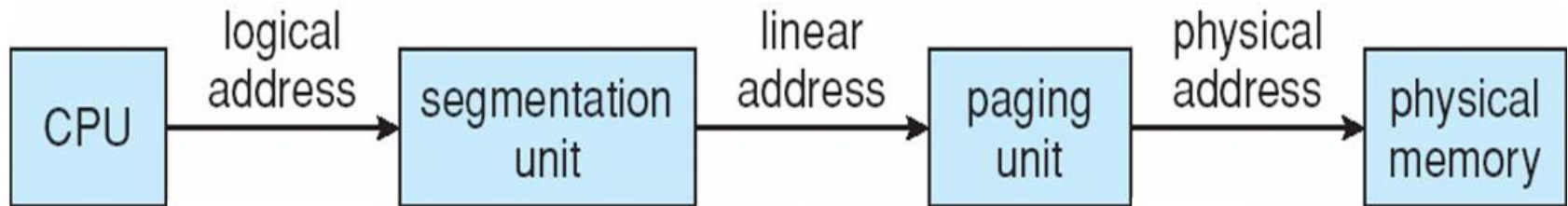
- Linear address given to paging unit
  - ▶ Which generates physical address in main memory
  - ▶ Paging units form equivalent of MMU
  - ▶ Pages sizes can be 4 KB or 4 MB





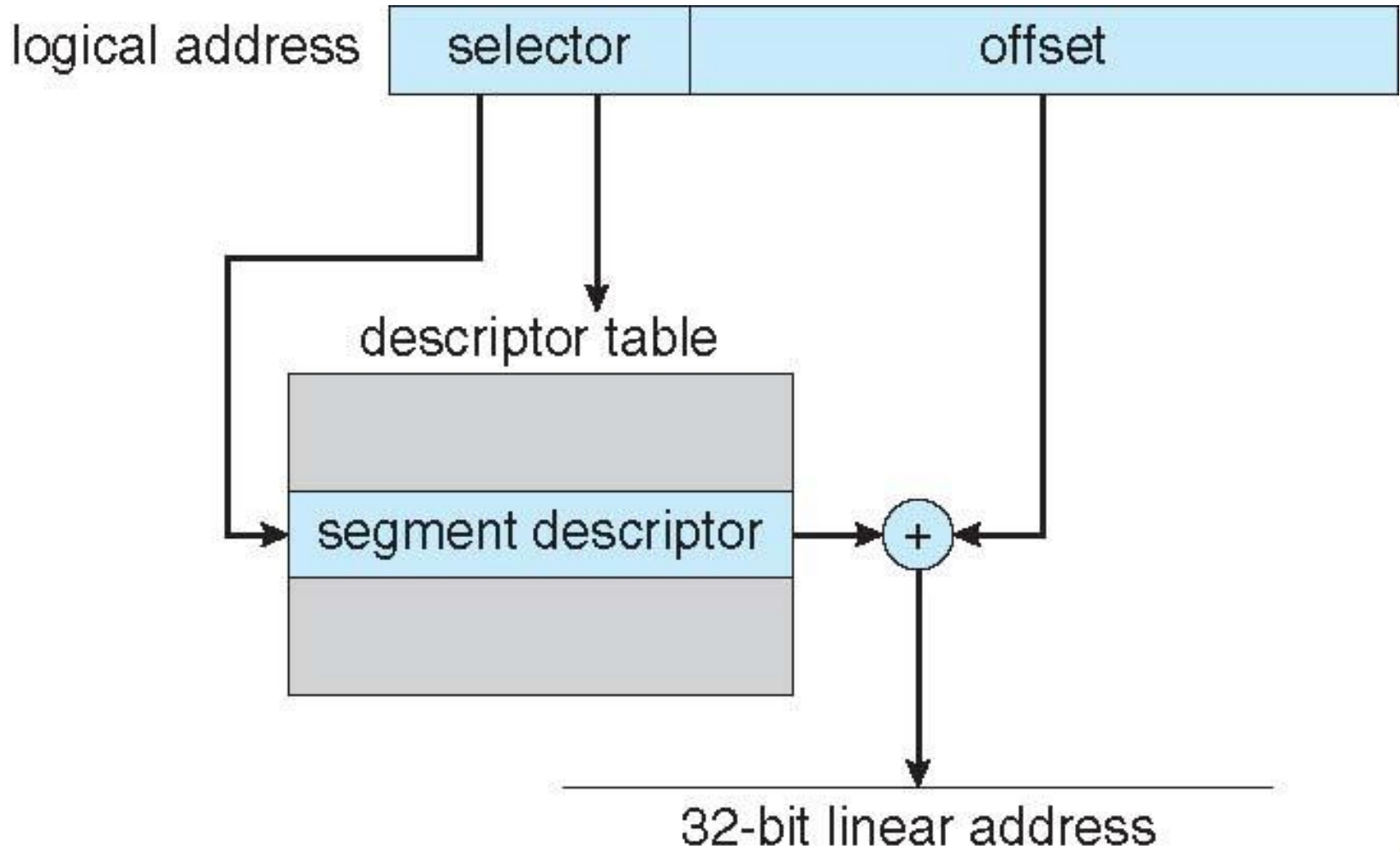


# Logical to Physical Address Translation in IA-32



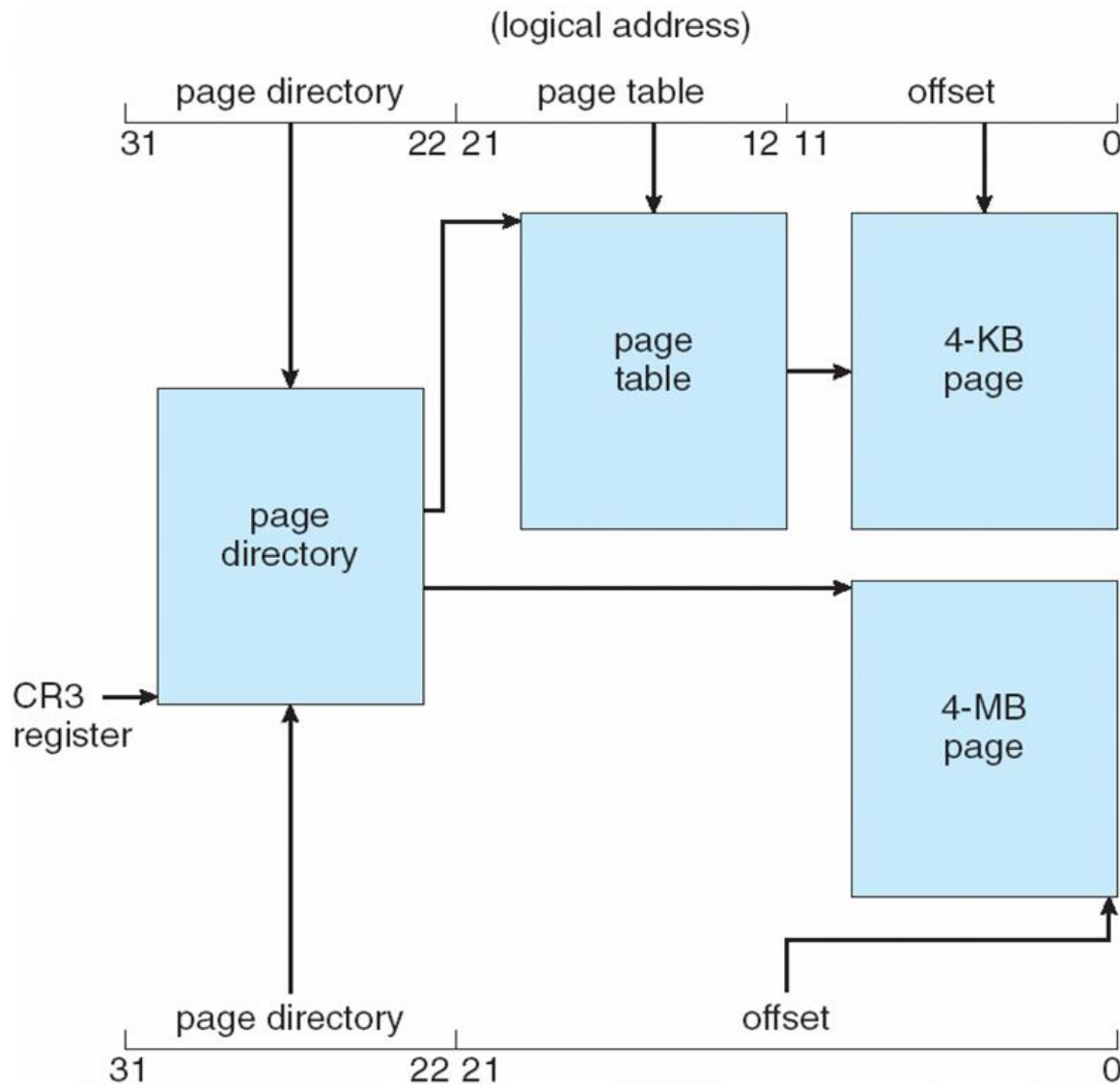


# Intel IA-32 Segmentation





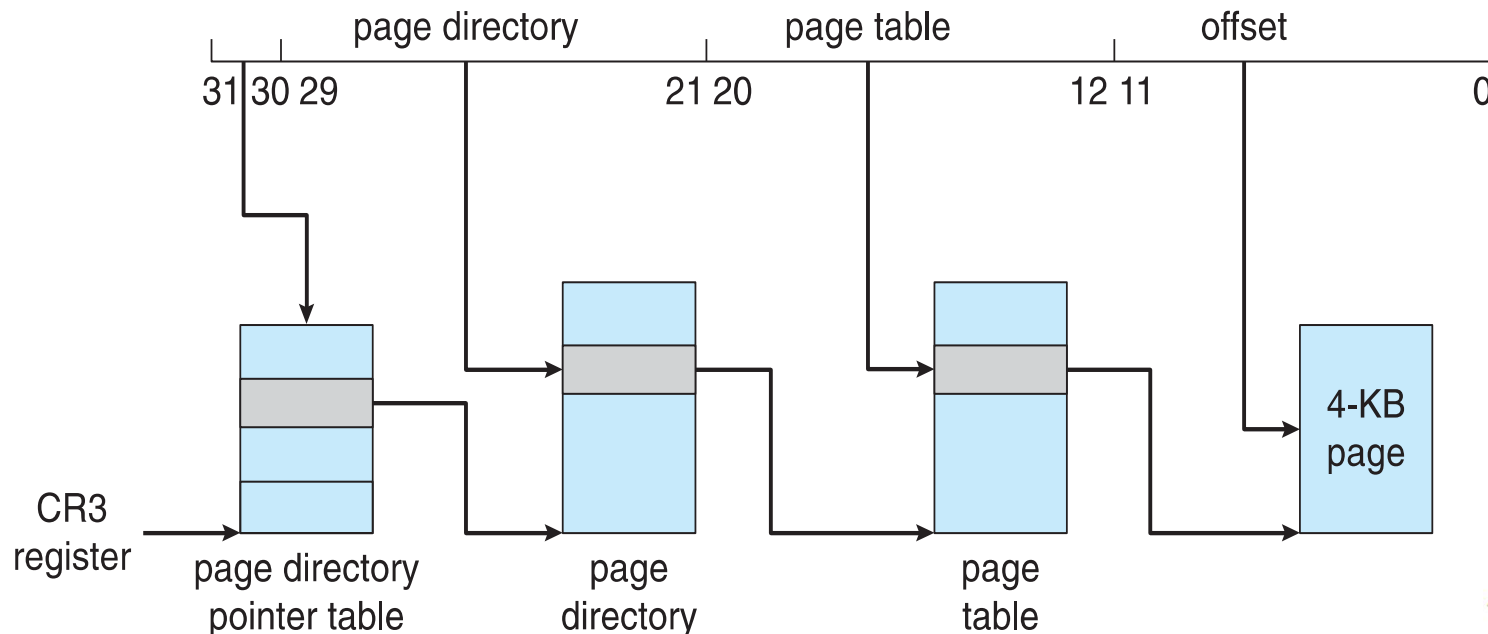
# Intel IA-32 Paging Architecture





# Intel IA-32 Page Address Extensions

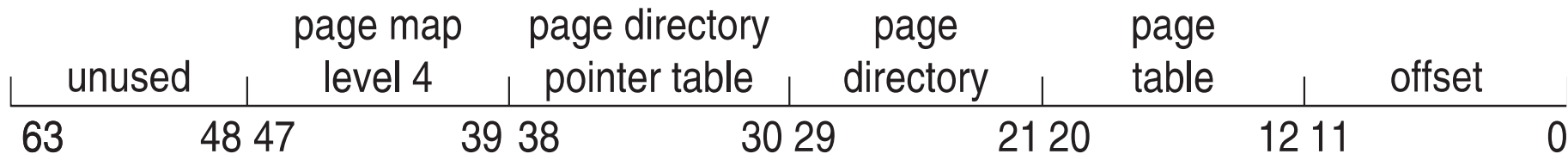
- 32-bit address limits led Intel to create **page address extension (PAE)**, allowing 32-bit apps access to more than 4GB of memory space
  - Paging went to a 3-level scheme
  - Top two bits refer to a **page directory pointer table**
  - Page-directory and page-table entries moved to 64-bits in size
  - Net effect is increasing address space to 36 bits – 64GB of physical memory





# Intel x86-64

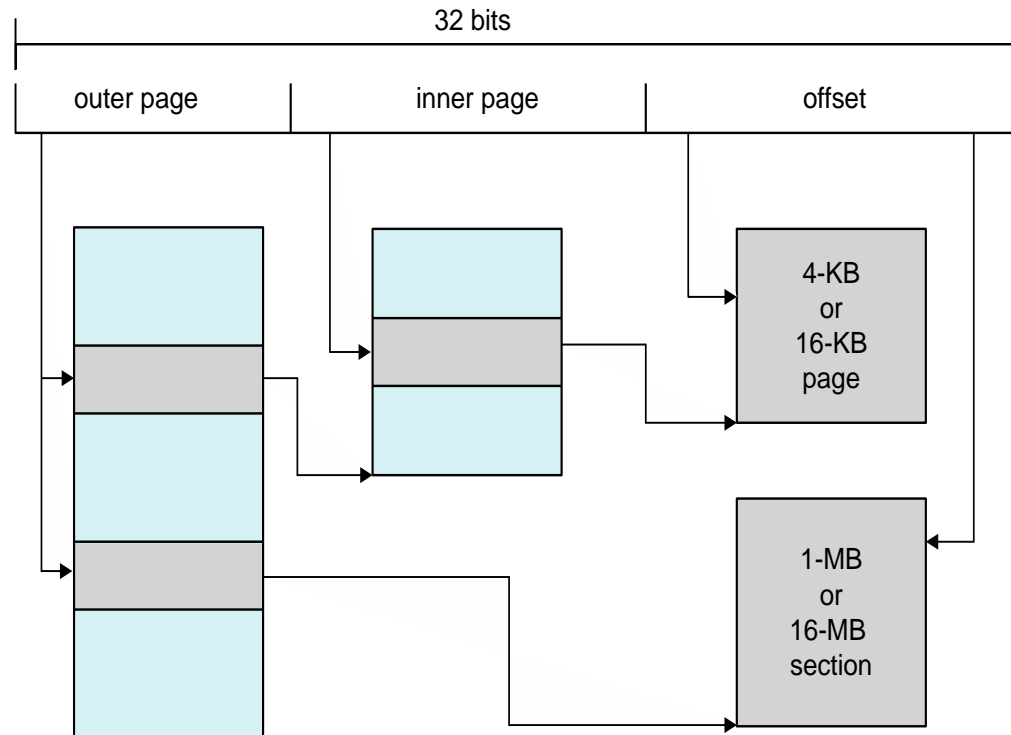
- In practice only implement 48 bit addressing
  - Page sizes of 4 KB, 2 MB, 1 GB
  - Four levels of paging hierarchy
- Can also use PAE so virtual addresses are 48 bits and physical addresses are 52 bits





# Example: ARM Architecture

- Dominant mobile platform chip (Apple iOS and Google Android devices for example)
- Modern, energy efficient, 32-bit CPU
- 4 KB and 16 KB pages
- 1 MB and 16 MB pages (termed **sections**)
- One-level paging for sections, two-level for smaller pages
- Two levels of TLBs
  - Outer level has two micro TLBs (one data, one instruction)
  - Inner is single main TLB
  - First inner is checked, on miss outers are checked, and on miss page table walk performed by CPU



# End of Chapter 8

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