

Optimization

Lagrange Multiplier

Def. **Lagrange Multiplier**.

Use to convert a optimization problem with constraints to one without constraints.

$$\begin{aligned} \min_{\omega} \quad & f(\omega) \\ \text{s.t.} \quad & g_i(\omega) \leq 0 \quad i = 1, \dots, k \Rightarrow \\ & h_j(\omega) = 0 \quad j = 1, \dots, l \\ \min_{\omega, \alpha, \beta} \quad & \mathcal{L}(\omega, \alpha, \beta) = f(\omega) + \sum_{i=1}^k \alpha_i g_i(\omega) + \sum_{j=1}^l \beta_j h_j(\omega) \end{aligned}$$

α_i, β_j are so-called Lagrange multiplier, $\alpha_i \geq 0$.

Lagrange Duality

- The solution to the dual problem provides a lower bound to the solution of the primal problem.

Def. Lagrange Dual Function

$$\begin{aligned} \mathcal{G}(\alpha, \beta) &= \inf_{\omega \in \mathcal{D}} \mathcal{L}(\omega, \alpha, \beta) \\ &= \inf_{\omega \in \mathcal{D}} \left(f(\omega) + \sum_{i=1}^k \alpha_i g_i(\omega) + \sum_{j=1}^l \beta_j h_j(\omega) \right) \end{aligned}$$

- The *Lagrange dual problem* with respect to the primal problem. The optimal value is d^* , and $d^* \leq p^*$.

$$\begin{aligned} \max_{\alpha, \beta} \quad & \mathcal{G}(\alpha, \beta) \\ \text{s.t.} \quad & \alpha \geq 0 \quad \forall i = 1, \dots, k \end{aligned}$$

Karush-Kuhn-Tucker (KKT) Conditions

- Let ω^* be a primal optimal point and (α^*, β^*) be a dual optimal solution.

Def. **KKT Conditions**

- Stationarity: $\nabla f(\omega^*) + \sum_{i=1}^k \alpha_i^* \nabla g_i(\omega^*) + \sum_{j=1}^l \beta_j^* \nabla h_j(\omega^*) = 0$
- Primal feasibility: $g_i(\omega^*) \leq 0, \forall i = 1, \dots, k$
 $h_j(\omega^*) = 0, \forall j = 1, \dots, l$
- Dual feasibility: $\alpha_i^* \geq 0, \forall i = 1, \dots, k$
- Complementary slackness: $\alpha_i^* g_i(\omega^*) = 0, \forall i = 1, 2, \dots, k$
- Proof. **Stationarity Condition**
 - ω^* is the minimizer of $\mathcal{L}(\omega, \alpha^*, \beta^*)$ over ω . Thus, $\nabla \mathcal{L} = 0$
- The primal feasibility conditions holds naturally.

- Proof. **Dual Feasibility**
 - If $\alpha \geq 0$ and $\tilde{\omega}$ is feasible, then
 - $f(\tilde{\omega}) \geq \mathcal{L}(\tilde{\omega}, \alpha, \beta) \geq \mathcal{G}(\alpha, \beta) = \inf_{\omega \in \mathcal{D}} \mathcal{L}(\omega, \alpha, \beta)$
- Proof. **Complementary Slackness**
 - If strong duality holds, then
 - $f(\omega^*) = \mathcal{G}(\alpha^*, \beta^*)$
 - $$\leq f(\omega^*) + \sum_{i=1}^k \alpha_i^* g_i(\omega^*) + \sum_{j=1}^l \beta_j^* h_j(\omega^*)$$
 - $$\leq f(\omega^*)$$
 - $\therefore \sum_{i=1}^k \alpha_i^* g_i(\omega^*) = 0$
 - Since each term is nonpositive, $\alpha_i^* g_i(\omega) = 0$.

Convex Optimization

- If objective function $f(\omega)$ and inequality constraints $g_i(\omega)$ are convex, and the equality constraints $h_j(\omega)$ are affine functions. A convex optimization problem can be represented by

$$\begin{array}{ll} \min_{\omega} & f(w) \\ \text{s.t.} & g_i(w) \leq 0, i = 1, \dots, k \\ & Aw - b = 0 \end{array}$$

- where, $A \in \mathbb{R}^{l \times n}$ and $b \in \mathbb{R}^l$.

Theorem. **Slater's Condition** (one of so-called *constraint qualification*, a sufficient condition)

Strong duality holds for a convex problem if it is strictly feasible, i.e.,

$$\exists \omega \in \text{relint } \mathcal{D} : g_i(\omega) < 0, i = 1, \dots, m, Aw = b$$

relint (relative interior)