

9. Lazy Evaluation

When passing **parameters** to a function, programming language design offers **two options** which are not mutually exclusive (both strategies can be implemented in the language):

- **applicative** (also called **strict**) evaluation strategy:
 - parameters are always evaluated **first**
 - can be further refined into: **call-by-value** (e.g. as it happens in *C*) and **call-by-reference** (e.g. as it happens for objects in *Java*).
- **normal** evaluation strategy (also called **non-strict**, and when implemented as part of the *PL* - **call-by-name**)
 - the function is always evaluated **first**
 - can be further refined into: **lazy**, which ensures that each expression is evaluated **at most once**

For more details, see the lecture on lazy evaluation. In *Haskell*, the default evaluation strategy is **lazy**.

There are ways in which we can force evaluation to be **strict**, however, in this lab, we will only explore several programming constructs which benefit from lazy evaluation.

9.0. Evaluation

9.0.0. Describe the **evaluation strategy** of the following expressions:

a.)

```
foldr ((||).(==1)) False [2,1,3,4,5,6,7]
```

b.)

```
foldl (flip ((||).(==1))) False [2,1,3,4,5,6,7]
```

9.1. Streams

There exists an explicit **Stream** type in Haskell, but for this lab we can think of **streams** as synonymous with **infinite** lists.

If you want to see their content, use the **take** function (**:t take**) to check the first n values.

9.1.1. Define the stream of **natural** numbers.

```
nats :: [Integer]
```

9.1.2. Using the stream of **natural** numbers, define the stream of **odd** numbers and **perfect squares**. Hint: use *higher order functions* for a quick solution.

```
-- 1, 3, 5, 7, ...
odds  :: [Integer]

-- 0, 1, 4, 9, ...
squares :: [Integer]
```

9.1.3. Define the stream of **Fibonacci** numbers. Hint: `:t zipWith`.

```
fibs :: [Integer]
```

9.2. Infinite Binary Trees

You can use the following snippet to implement a `Show` instance for `BTree` that pretty prints the tree, you don't have to understand it, it should just format the Tree a bit nicer whenever it will get shown.
[Click to hide](#)

```
data BTree = Node Int BTree BTree | Nil

data PrintInfo = PrintInfo {
  len  :: Int,
  center :: Int,
  text :: [String]
}

pp :: BTree -> PrintInfo
pp Nil = PrintInfo 3 2 ["Nil"]
pp (Node x l r) = seq check PrintInfo nlen ncenter ntext
  where
    check = if (length (show x)) > nlen then error "Nice try" else ()
    pp_l = pp l
    pp_r = pp r
    nlen = len pp_l + len pp_r + 1
    ncenter = len pp_l + 1
```

```

ntext = aligned_x : center_line : dotted_line : down_lines : combined_lines
where
    aligned_x = replicate (ncenter - (div (length (show x)) 2) - 1) ' ' ++ show x
    center_line = replicate (ncenter - 1) ' ' ++ "|"
    dotted_line = replicate (center pp_l - 1) ' ' ++
        replicate (nlen - center pp_l - center pp_r + 2) '-'
    down_lines = replicate (center pp_l - 1) ' ' ++ "|" ++
        replicate (nlen - center pp_l - center pp_r) ' ' ++ "|"
    combined_lines = zipPad "" (combine) (text pp_l) (text pp_r)
where
    zipPad :: a -> (a -> a -> a) -> [a] -> [a] -> [a]
    zipPad pad f [] [] = []
    zipPad pad f [] (y:ys) = y : zipPad pad f [pad] ys
    zipPad pad f (x:xs) [] = x : zipPad pad f xs [pad]
    zipPad pad f (x:xs) (y:ys) = f x y : zipPad pad f xs ys
    combine l r = l ++ replicate (len pp_l - length l + 1) ' ' ++ r

instance Show BTree where
    show = unlines . text . pp

tree :: BTree
tree = Node 1 (Node 2 Nil Nil) (Node 3 Nil Nil)

```

Defining a simple **binary tree** structure in *Haskell* is easy.
Take this for example:

```

data BTree = Node Int BTree BTree | Nil deriving Show

```

```

{-
    1
   /
  -----
 /      \
2         3
-----  -----
/  /  /  /

```

```

Nil Nil Nil Nil
-}

tree :: BTree
tree = Node 1 (Node 2 Nil Nil) (Node 3 Nil Nil)

```

But what if we want to enforce an **infinite binary tree**?
The solution is eliminating the need for the empty node `Nil`:

```

-- This is an infinite data type, no way to stop generating the tree
data StreamBTree = StreamNode Int StreamBTree StreamBTree

sbtree = StreamNode 1 sbtree sbtree

```

9.2.1. In order to view an **infinite tree**, we need to convert it to a **finite binary tree**. Define the function `sliceTree`, which takes a level `k`, an **infinite tree** and returns the first `k` levels of our tree, in the form of a **finite** one.

```

{-
> sliceTree 2 sbtree

      1
      |
  -----
    /   \
   1     1
  -----
 /  /  /  \
Nil Nil Nil Nil
-}

sliceTree :: Int -> StreamBTree -> BTree

```

9.2.2. Define the `repeatTree` function which takes an `Int k` and generates an **infinite tree**, where each node has the value `k`.

```

{-
> repeatTree 3

```

```

      3
      /
     -----
    /      \
   3        3
  -----  -----
 /  \  /  \ /  \
3   3 3   3 3   3
.   . .   . .   .
.   . .   . .   .
-}

```

```
repeatTree :: Int -> StreamBTree
```

9.2.3. Define the `generateTree` function, which takes a `root k`, a `left generator` function `leftF` and a `right generator` function `rightF`.

For example, let's say we have `k=2`, `leftF=(+1)`, `rightF=(*2)`. This should generate a tree where the *root* is 22, the *left child* is `parent+1` and the *right child* is `parent*2`.

```

{-
> generateTree 2 (+1) (*2)

      2
      /
     -----
    /      \
   3        4
  -----  -----
 /  \  /  \ /  \
4   6 5   8 6   8
  -----  -----
 /  \ /  \ /  \ /  \
5   8 7   12 6   10 9   16
  -----  -----
 /  \ /  \ /  \ /  \ /  \

```

```

6   10  9   16  8   14  13  24           7   12  11  20   10  18  17  32
.   .   .   .   .   .   .   .           .   .   .   .   .   .   .   .
.   .   .   .   .   .   .   .           .   .   .   .   .   .   .   .
-}

generateTree :: Int -> (Int -> Int) -> (Int -> Int) -> StreamBTree

```

9.3. Numerical Approximations

9.3.1. Define the `build` function which takes a generator `g` and an initial value `a0` and generates the **stream**: `[a0, g a0, g (g a0), g (g (g a0)), ...]`.

```

{-
    with this function, you should
    be able to easily define the natural numbers
    with something like: build (+1) 0
-}

build :: (Double -> Double) -> Double -> [Double]

```

9.3.2. Using the `build` function, define the following **streams**:

```

-- 0, 1, 0, 1, 0, 1, 0, ...
alternatingBinary :: [Double]

-- 0, -1, 2, -3, 4, -5, ...
alternatingCons :: [Double]

-- 1, -2, 4, -8, 16, ...
alternatingPowers :: [Double]

```

9.3.3. Define the `select` function which takes a **tolerance** `e` and a **stream** `s` and returns the `nth` element of the **stream** which satisfies the following condition: $|s_n - s_{n+1}| < e$.

```

select :: Double -> [Double] -> Double

```

Mathematical Constants

9.3.4. Knowing that $\lim_{n \rightarrow \infty} \frac{F_{n+1}}{F_n} = \varphi$, where F_n is the n th element of the **Fibonacci** sequence, write an approximation with **tolerance** `e=0.00001` of the **Golden Ration** (φ). Use the previously defined **fibs stream** and the **select** function.

```
phiApprox :: Double
```

9.3.5. Consider the sequence:

$a_{n+1} = a_n + \sin(a_n)$; where a_0 is an *initial approximation*, randomly chosen (but **not** 0 because $a_{n+1} \neq a_n$).

Knowing that $\lim_{n \rightarrow \infty} a_n = \pi$, write an approximation with **tolerance** `e=0.00001` of π . Make sure to use **build** and **select**.

```
piApprox :: Double
```

Square Root

9.3.6. Given a number `k`, we want to create a function which calculates the **square root** of `k`. This is another place where **laziness** and **streams** come into play.

Consider the following sequence:

$a_{n+1} = \frac{1}{2}(a_n + \frac{k}{a_n})$; where a_0 is an *initial approximation*, randomly chosen.

Knowing that $\lim_{n \rightarrow \infty} a_n = \sqrt{k}$, write a function that approximates \sqrt{k} with **tolerance** `e=0.00001`. Use **build** and **select**.

```
sqrApprox :: Double -> Double
```

Derivatives

9.3.7. We can approximate the derivative of a function in a certain point using the definition of the derivative:

$$f'(a) = \lim_{h \rightarrow 0} \frac{f(a+h) - f(a)}{h} \quad f'(a) = \lim_{h \rightarrow 0} \frac{f(a+h) - f(a)}{h}$$

We can obtain better successive approximations of the derivative in a point `a`, using a smaller `h`.

a) generate the sequence: `h0, h0/2, h0/4, h0/8, ...` (where `h0` is a randomly chosen *initial approximation*)

b) generate the list of approximations for $f'(a)$, using the formula above

c) write the function that takes a function `f` and a point `a` and approximates $f'(a)$ with **tolerance** `e=0.00001`, using the previous steps.

```
derivativeApprox :: (Double -> Double) -> Double -> Double
```