

Constraint & Data Mining

Cours1

Master 2 - DS

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Plan

Constraint Programming

Data Mining

DM \leftrightarrow CP



Constraint Programming

« The user states the problem,
the computer solve it! »

Constraint Programming

Constraint Programming

Definition

- A powerful declarative programming paradigm combining AI, OR and Logic Programming
- Formalism to model (**constraint network**) and to solve (**constraint solver**) combinatorial problems (**scheduling, planning,...**).
- Examples of constraints :

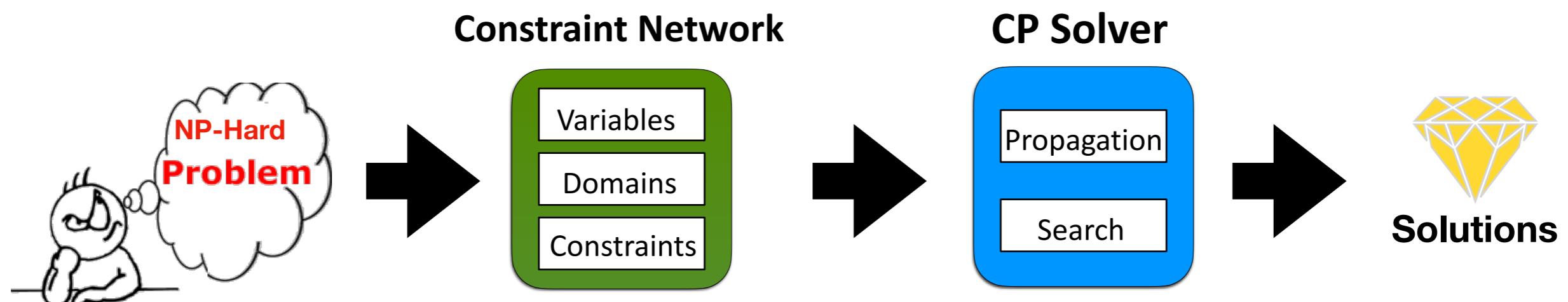
$$x_1 + x_2 + x_3 = 5; \quad allDifferent(X); \quad c(x_i, x_j) = \{(1,1), (2,1), (2,3), (4,4)\}$$

Constraint Programming

Definition

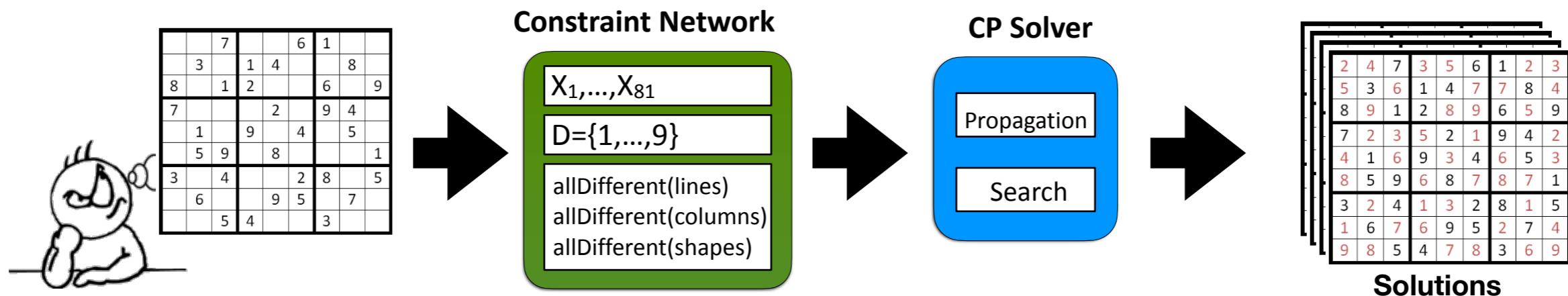
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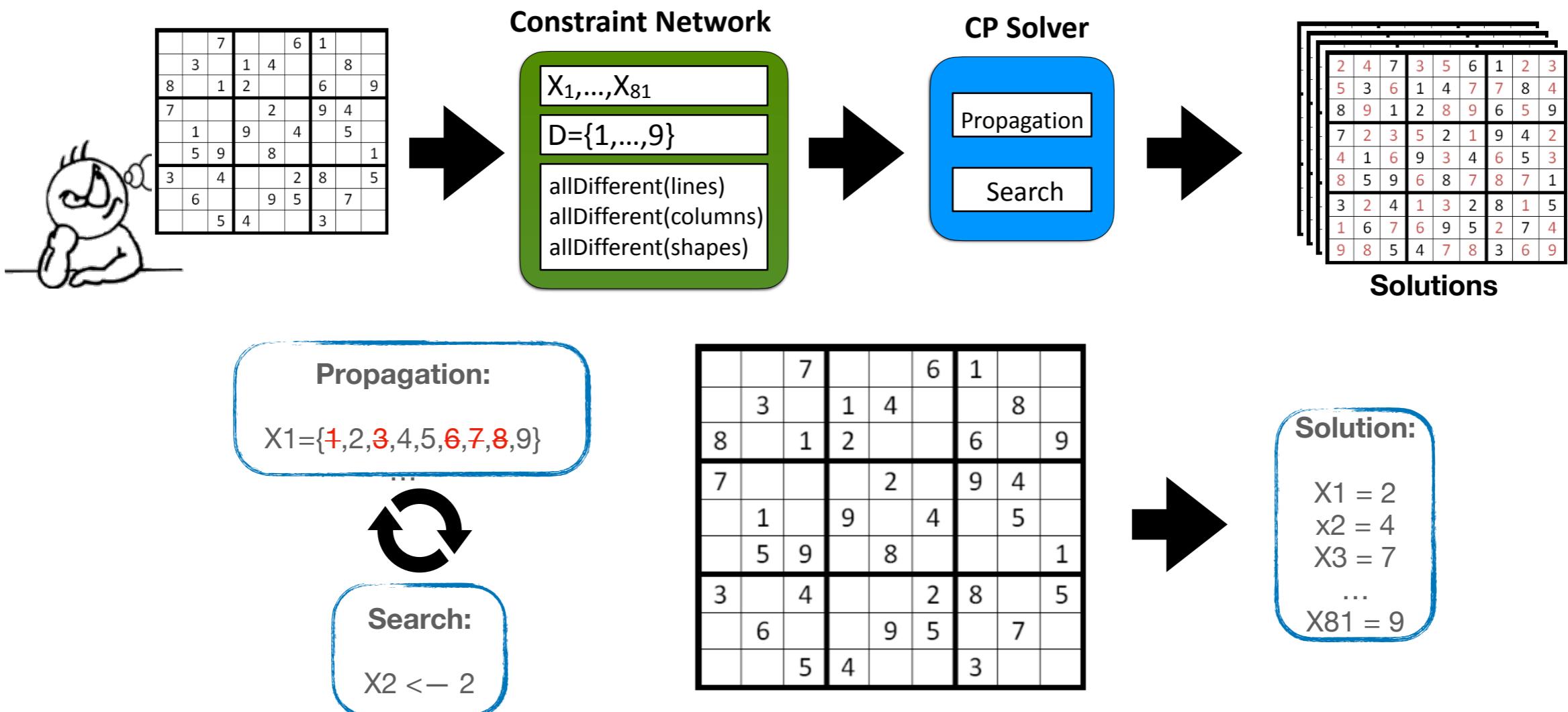
Constraint Programming

Example



Constraint Programming

Example



Constraint Programming

Why CP?

- ▶ Combinatorial problems can be solved by Integer Linear Programming (ILP) or by propositional satisfiability (SAT)
- ▶ Advantages of CP :
 - Compactness
 - Expressiveness
 - And (often) efficiency

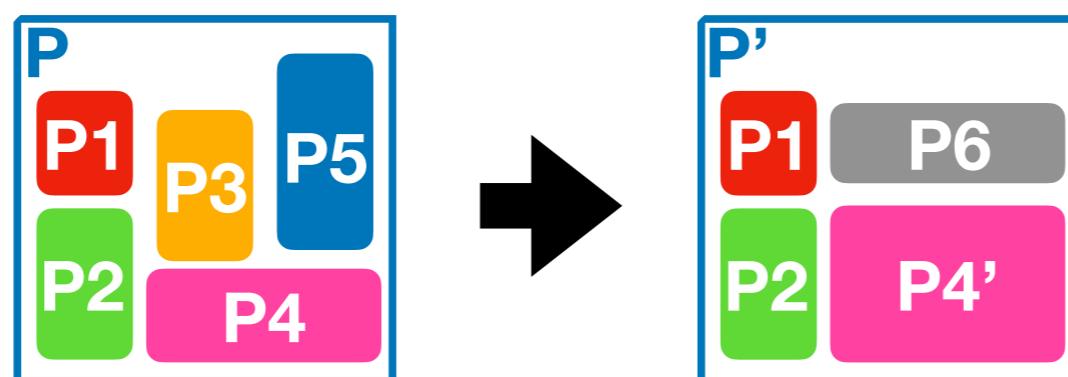
Constraint Programming

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Constraint Programming

Why CP?

- ▶ Sudoku 9x9: SAT p cnf 729 3240
- ▶ CP: [sudoku.mod](#) with n=9

Constraint Programming

Why CP?

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Variables $X = \{X_{1,1}, \dots, X_{n,n}\}$

Domaine des variables : $\{1, \dots, n\}$

Contraintes :

`allDifferent(Xi,1, ..., Xi,n), $\forall i$ (rows)`

`allDifferent(Xi,1, ..., Xi,n), $\forall i$ (columns)`

`allDifferent(Xi,j), $\forall i, j$ (squares)`

[sudoku.mod](#)

Sudoku 36x36

Constraint Programming

Why CP?

► Sudoku 36x36: SAT p cnf 46656 821664

$$\#vars = n^3$$

$$\#clauses = n^3(\text{unaries}) + n^2 \frac{n(n - 1)}{2}(\text{binaries}) + 4n^2(\text{n-aries})$$

► CP: [sudoku.mod](#) with n=9

Variables X= { $x_{1,1}, \dots, x_{n,n}$ }

Domaine des variables : {1, ..., n}

Contraintes :

allDifferent($x_{i,1}, \dots, x_{i,n}$), $\forall i$ (rows)
allDifferent($x_{1,j}, \dots, x_{n,j}$), $\forall j$ (columns)
allDifferent($x_{i,j}$), $\forall i, j$ (squares)

[sudoku.mod](#)

Constraint Programming

Why CP?

► Nurse Rostering Problem

► Best Linear Program: more than 10 000 lines

Employee shift rostering

Populate each work shift with a nurse.

	Maternity nurses			Emergency nurses			Basic nurses		
	A Ann	B Beth	C Cory	D Dan	E Elin	G Greg	H Hue	I Ilse	
Largest staff first									
Maternity nurses	1 C A B	1 C A	2 A C B	1 D G E	2 D G E	1 D E	1 H I	2 C A B	1 C A
Emergency nurses	2 D G E	2 D G E	1 D E	2 D G E	2 D G E	1 D G	1 H I	2 C A B	1 C A
Any nurses	1 H I	1 H I	1 H I	1 H I	1 H I	1 H I	1 H I	1 H I	1 H I
Drools Planner									
Maternity nurses	1 C A B	1 C A	2 A C B	1 D G E	2 D G E	1 D E	1 H I	1 C A B	1 C A
Emergency nurses	2 D G E	2 D G E	1 D E	2 D G E	2 D G E	1 D G	1 H I	2 C A B	1 C A
Any nurses	1 H I	1 H I	1 H I	1 H I	1 H I	1 H I	1 H I	1 H I	1 H I

Constraint Programming

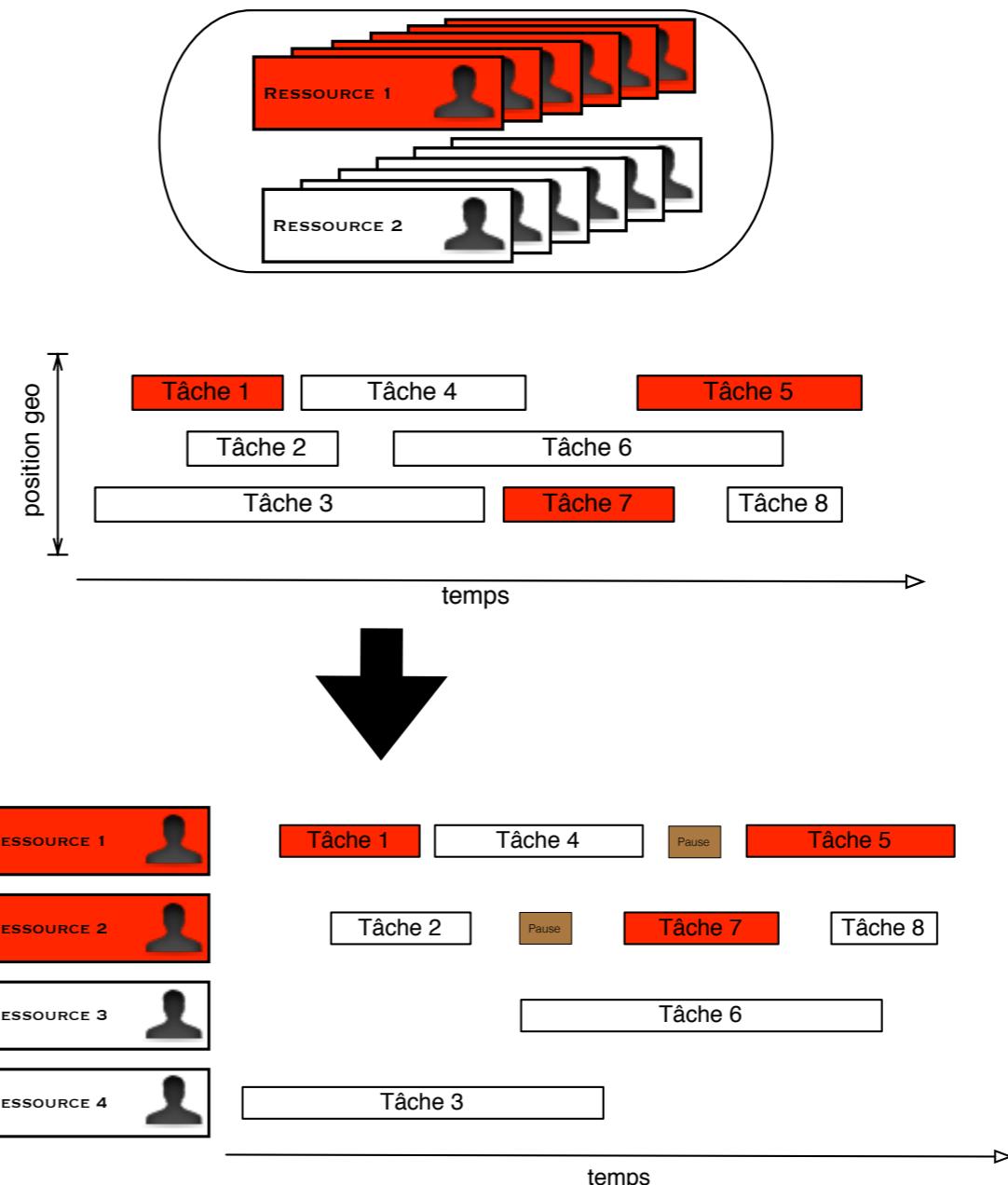
Solvers and CP platforms

IBM ILOG CP Optimizer (Java, C++, Python, .NET)	
Google OR-Tools (C++, Java, C#, Python)	 Google OR-Tools
Artelys Kalis (Java, C++, Python)	
SICStus Prolog (CLPFD bib in prolog)	
Gecode (C++)	
Choco (Java)	
Minizinc (high-level, solver-independent)	

Constraint Programming

SNCF train driver planning using CP

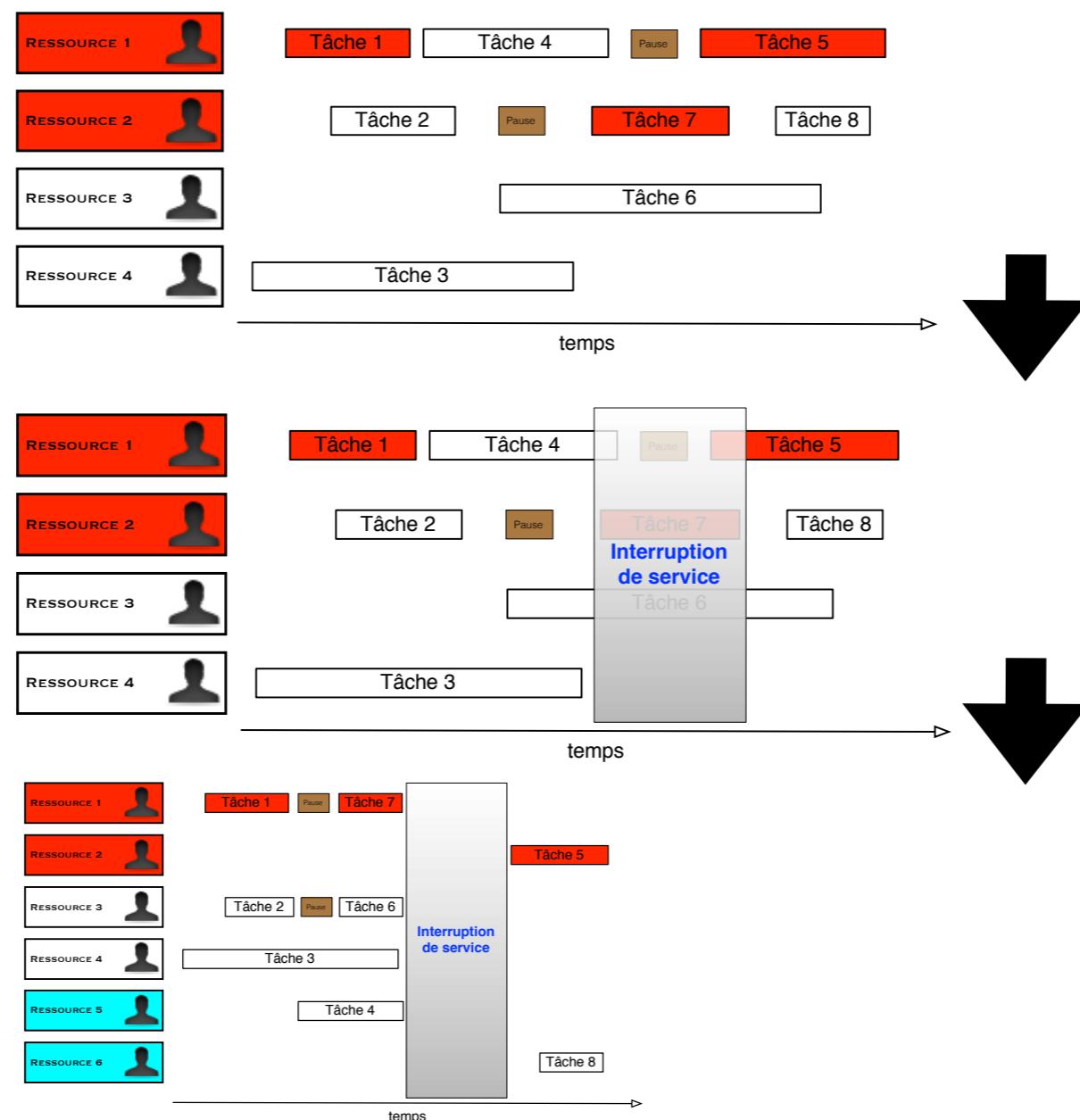
- Legal rules constraints:
 - daily working time
 - daily rest periods
 - ...
- Preferences:
 - Type of lines
 - Day of rest
 - Place of rest
 - ...
- Solution:
 - Best in terms of ressources
 - Robust one
 - Cheapest one
 - ...



Constraint Programming

SNCF train driver planning using CP

- Maintaining an existing planning



Constraint Programming

ABB Robotics partner projects

► SWMOD: Test Case Execution Scheduling with CP

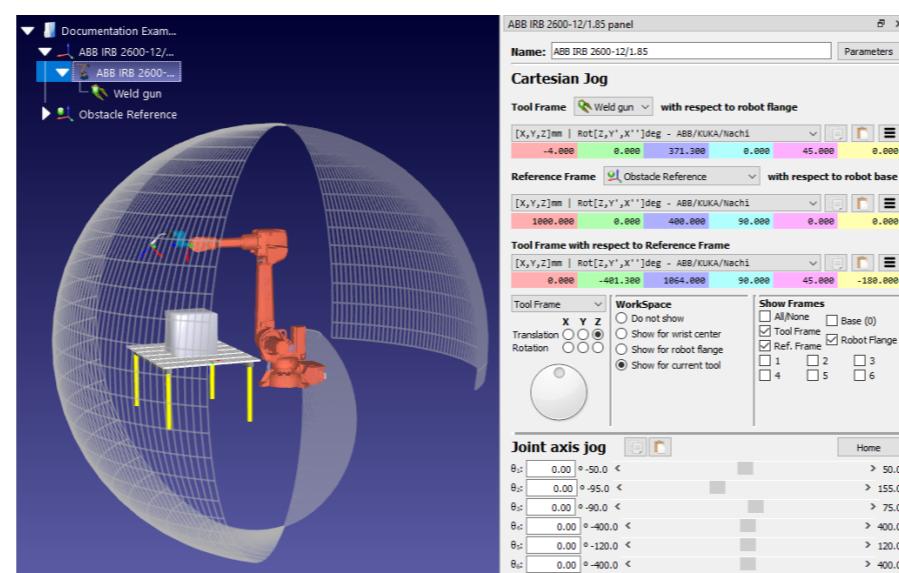


"SWMOD deployed at ABB Robotics and used every day to schedule tests throughout several ABB centers in the world (Norway, Sweden, India, China)"



<https://github.com/Makouno44/Robtest>

► Robtest: Optimal Stress Test Trajectories for Robots with CP

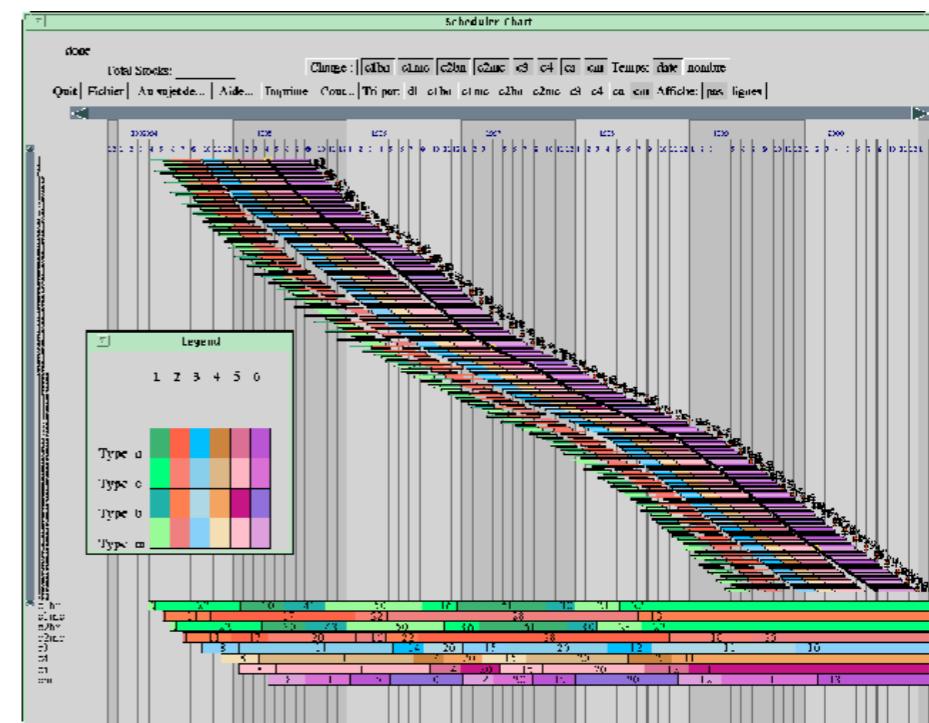


Constraint Programming

Aircraft Industry - Dassault Aviation



Assembly of Mirage aircraft

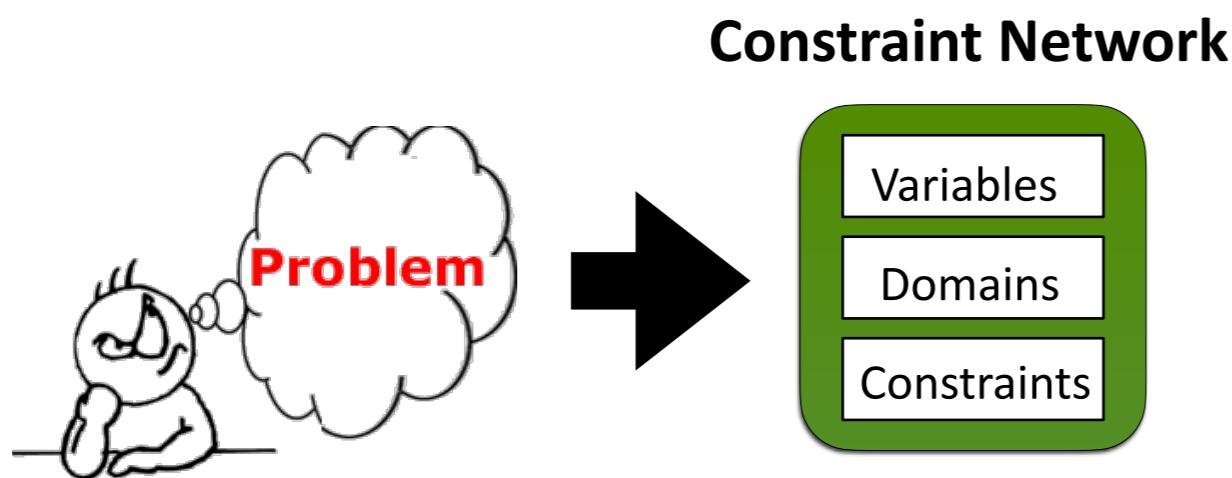


Constraint Programming

Modeling

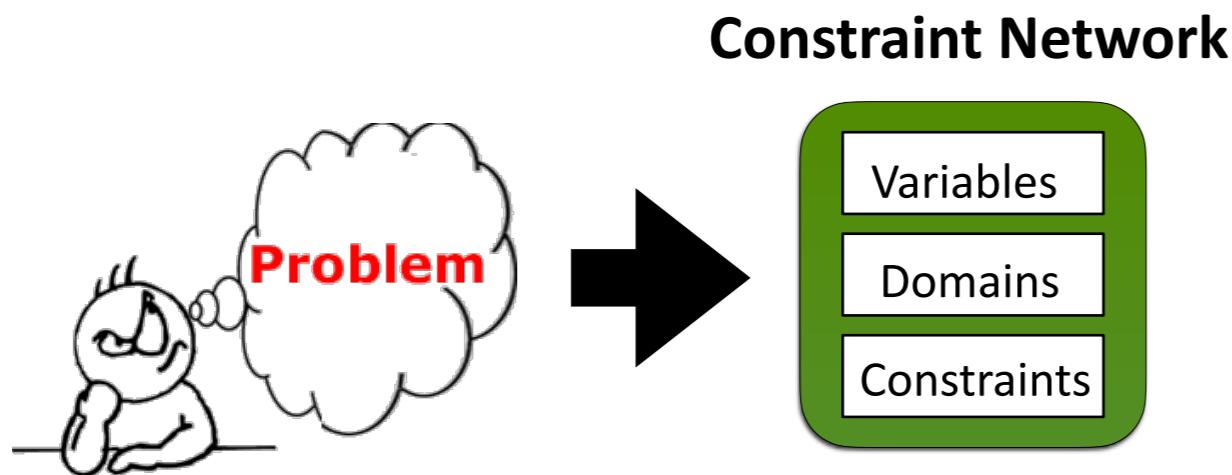
Constraint Programming

Modeling



Constraint Programming

Modeling



Constraint Network $N=(X, D, C)$

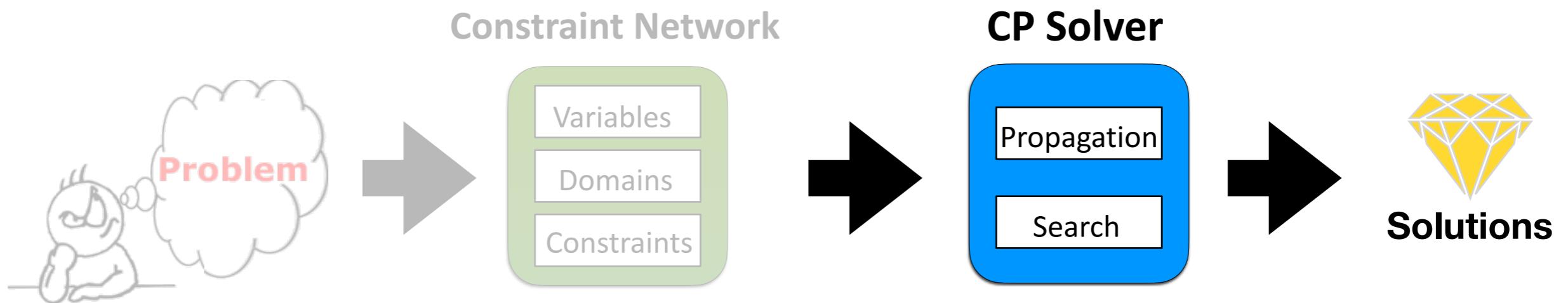
X: finite set of variables

D: domain on X, where $D(X_i)$ is a finite set of values for x_i

C: a set of constraints

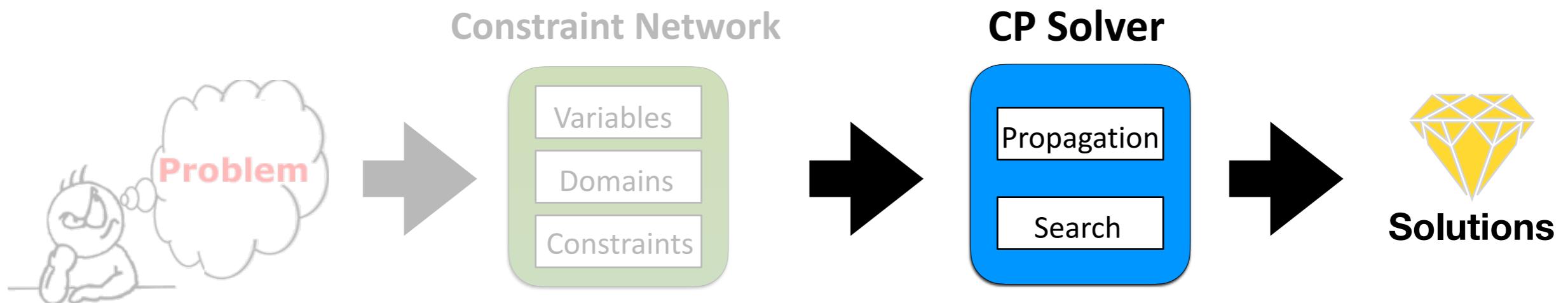
Constraint Programming

Solving



Constraint Programming

Solving



Instantiation I

- Complete: assignment on X
- Partial: assignment on $Y \subset X$
- Valid: values in D
- Violating c : assignment on $\text{var}(c)$ is not in c
- Locally consistent: assignment not violating any constraint on Y
- Solution: locally consistent on X

Constraint Programming

Solving - Backtracking (BT)

- A general-purpose search algorithm
- Systematically explores the search space
- Utilizes a depth-first search strategy
- May encounter inefficiency on large search spaces
- Suitable for a wide range of constraint problems
- May require additional pruning techniques

Constraint Programming

Solving - Backtracking (BT)

Constraint Programming

Solving - Backtracking (BT)

BT(<X,D,C>, I):

If I is complete then return true

Select a variable X_i not in I

ForEach v in $D(X_i)$ do

If I union $\langle X_i, v \rangle$ is locally consistent then

If **BT(<X,D,C>, I union <Xi, v>)** then
return true

return false

Constraint Programming

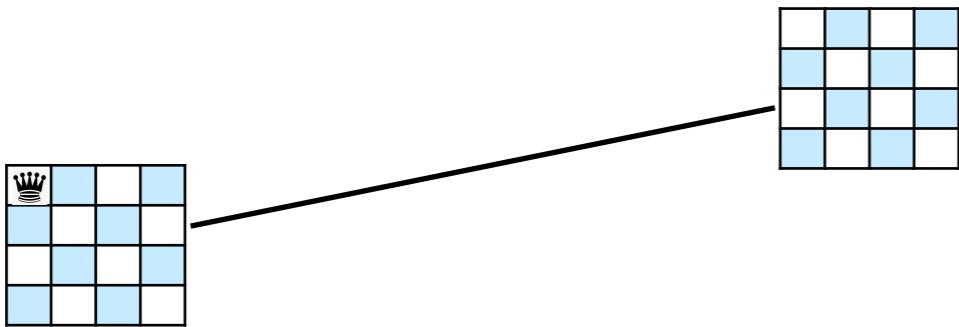
Solving - Backtracking (BT)

	R1	R2	R3	R4
p1				
p2				
p3				
p4				

N-queens problem

Constraint Programming

Solving - Backtracking (BT)

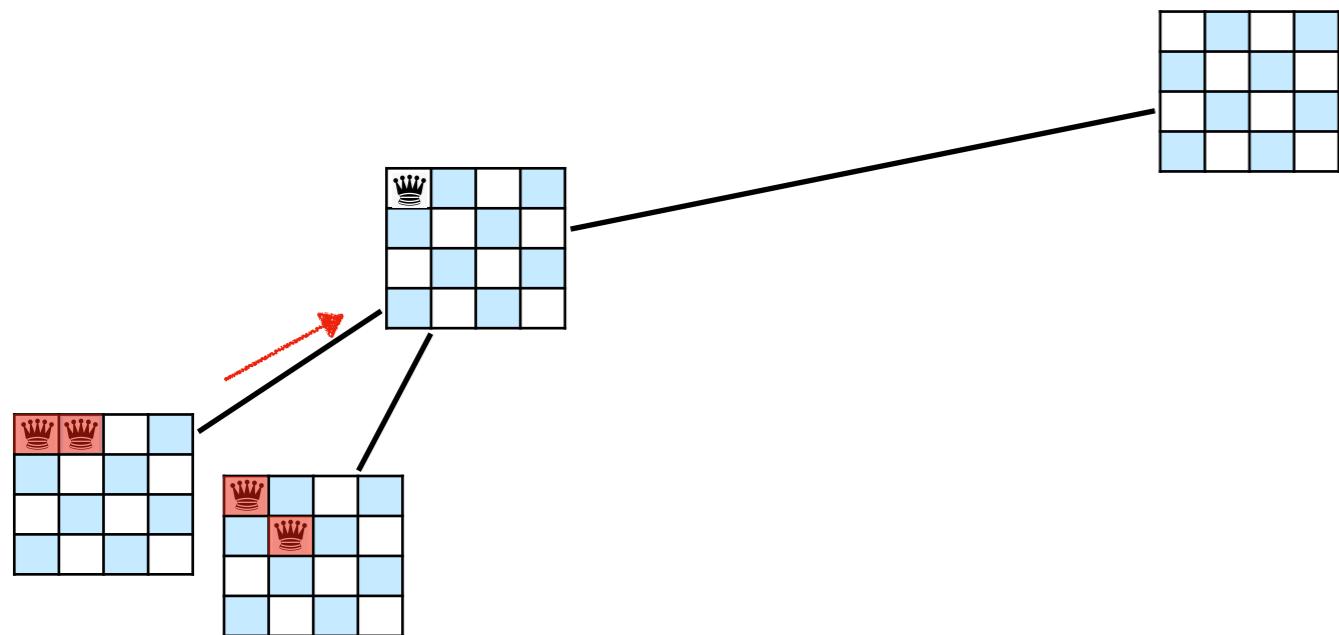


	R1	R2	R3	R4
p1				
p2		X		
p3			X	
p4	X			

N-queens problem

Constraint Programming

Solving - Backtracking (BT)

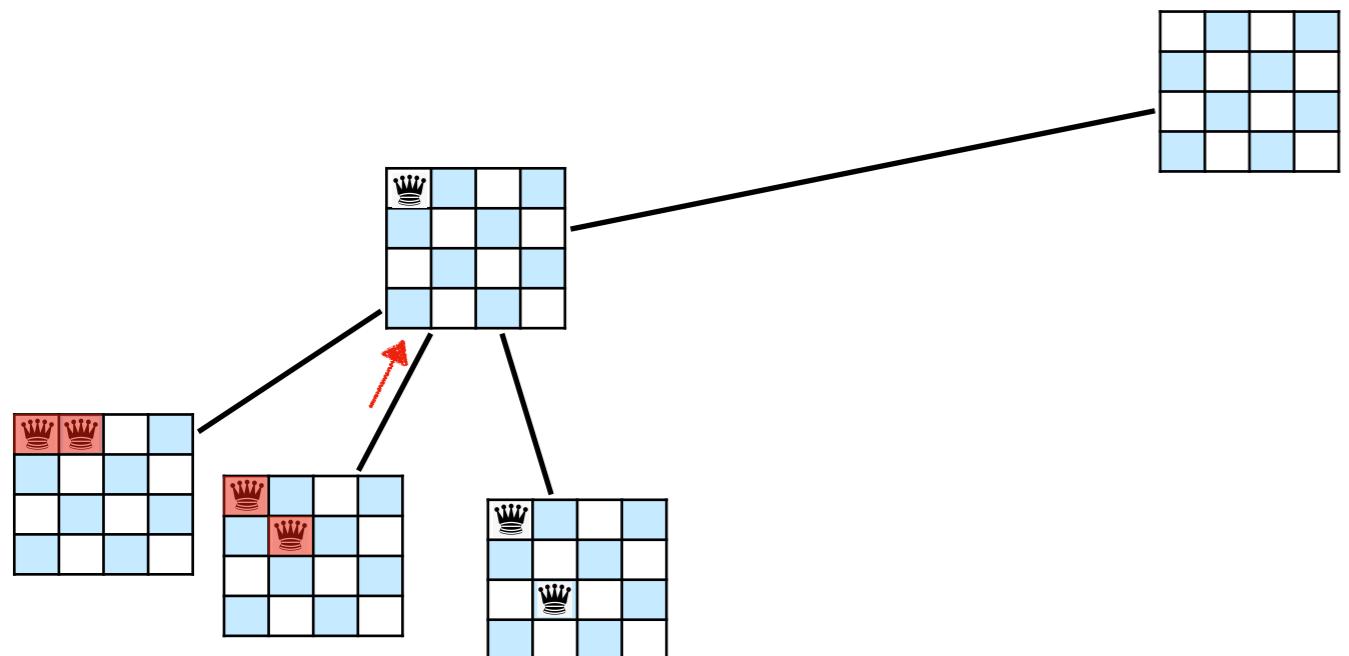


	R1	R2	R3	R4
p1				
p2				
p3				
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Constraint Programming

Solving - Backtracking (BT)

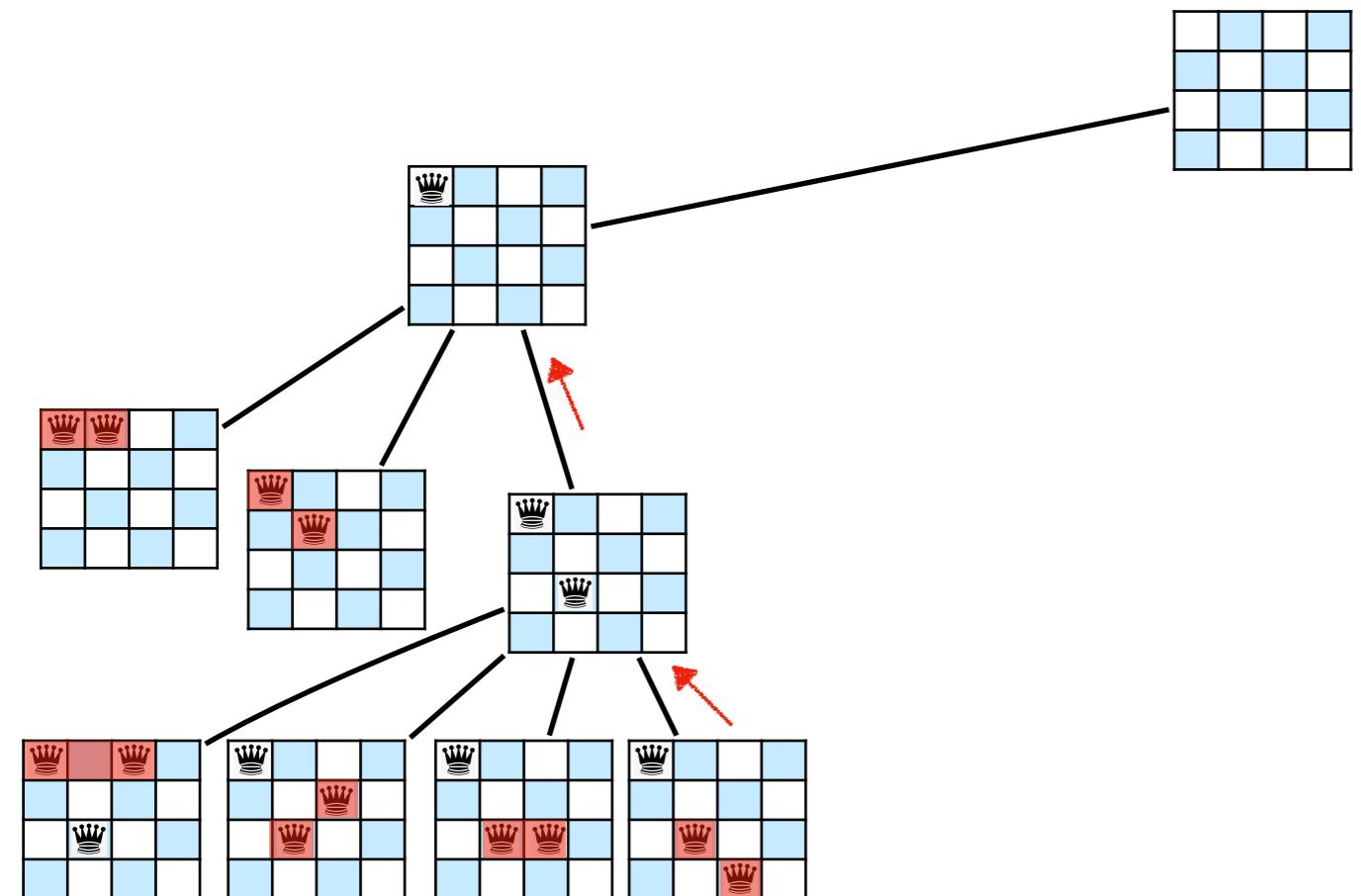


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Constraint Programming

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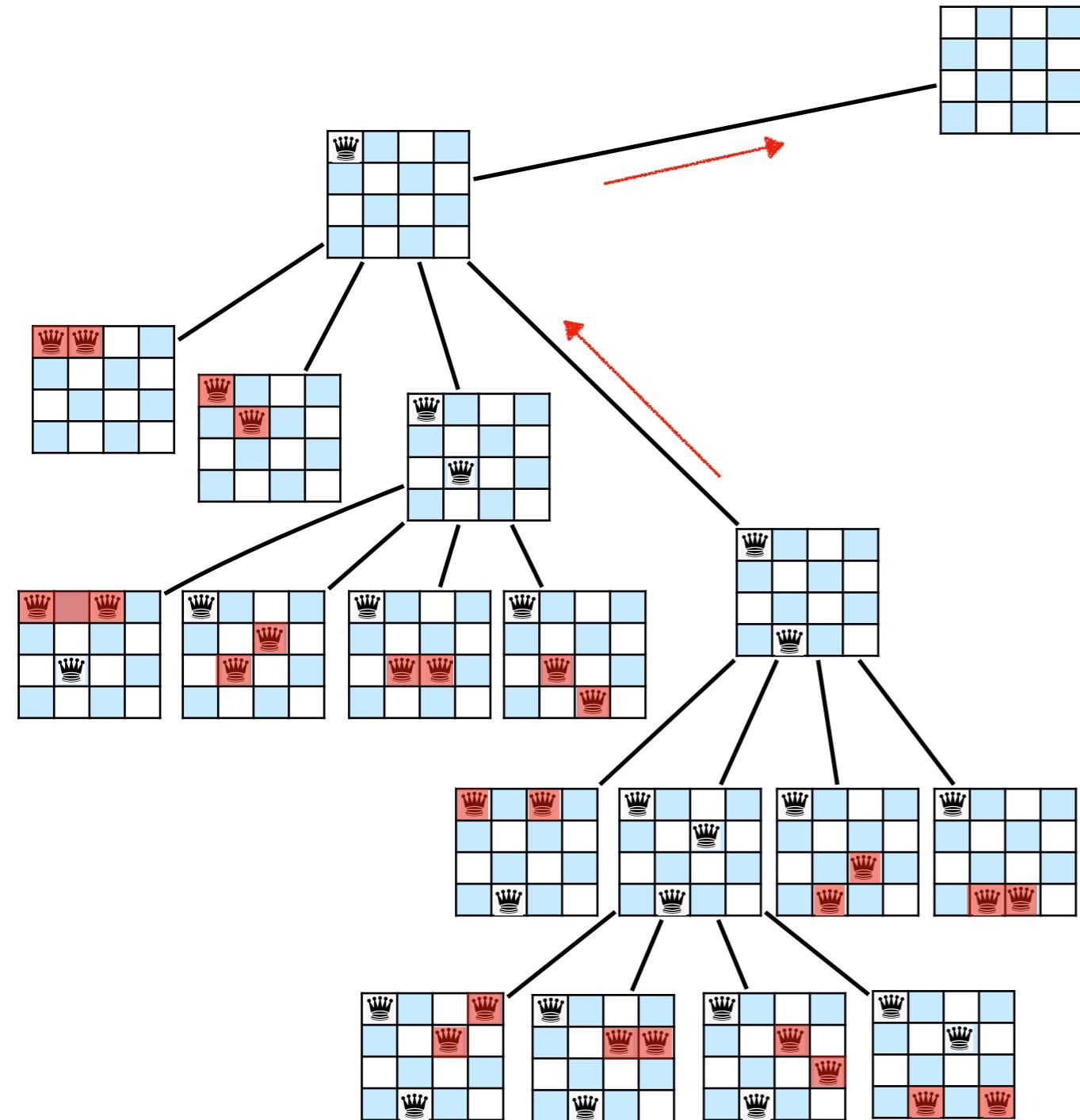


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Constraint Programming

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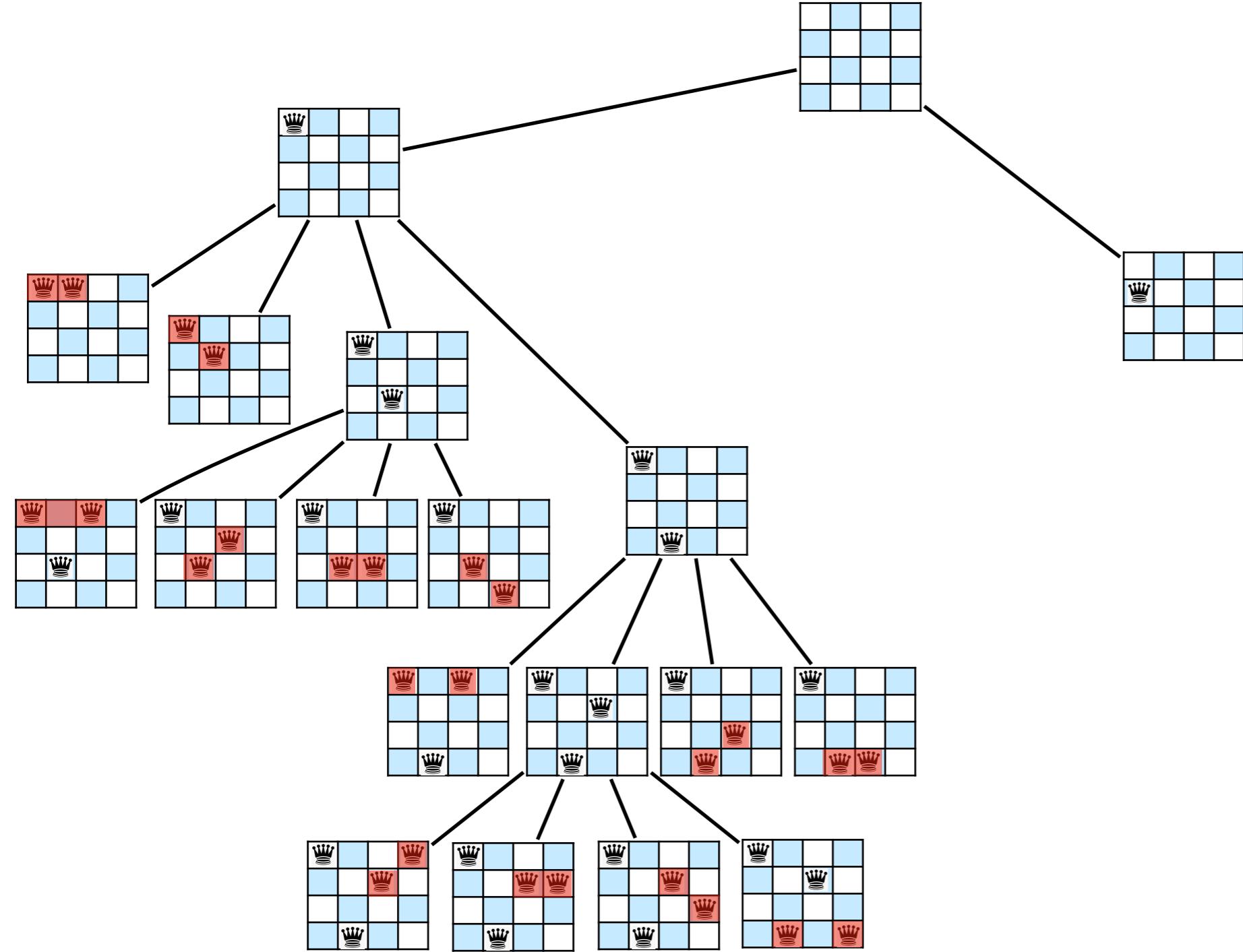


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p1				
p2				
p3				
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Constraint Programming

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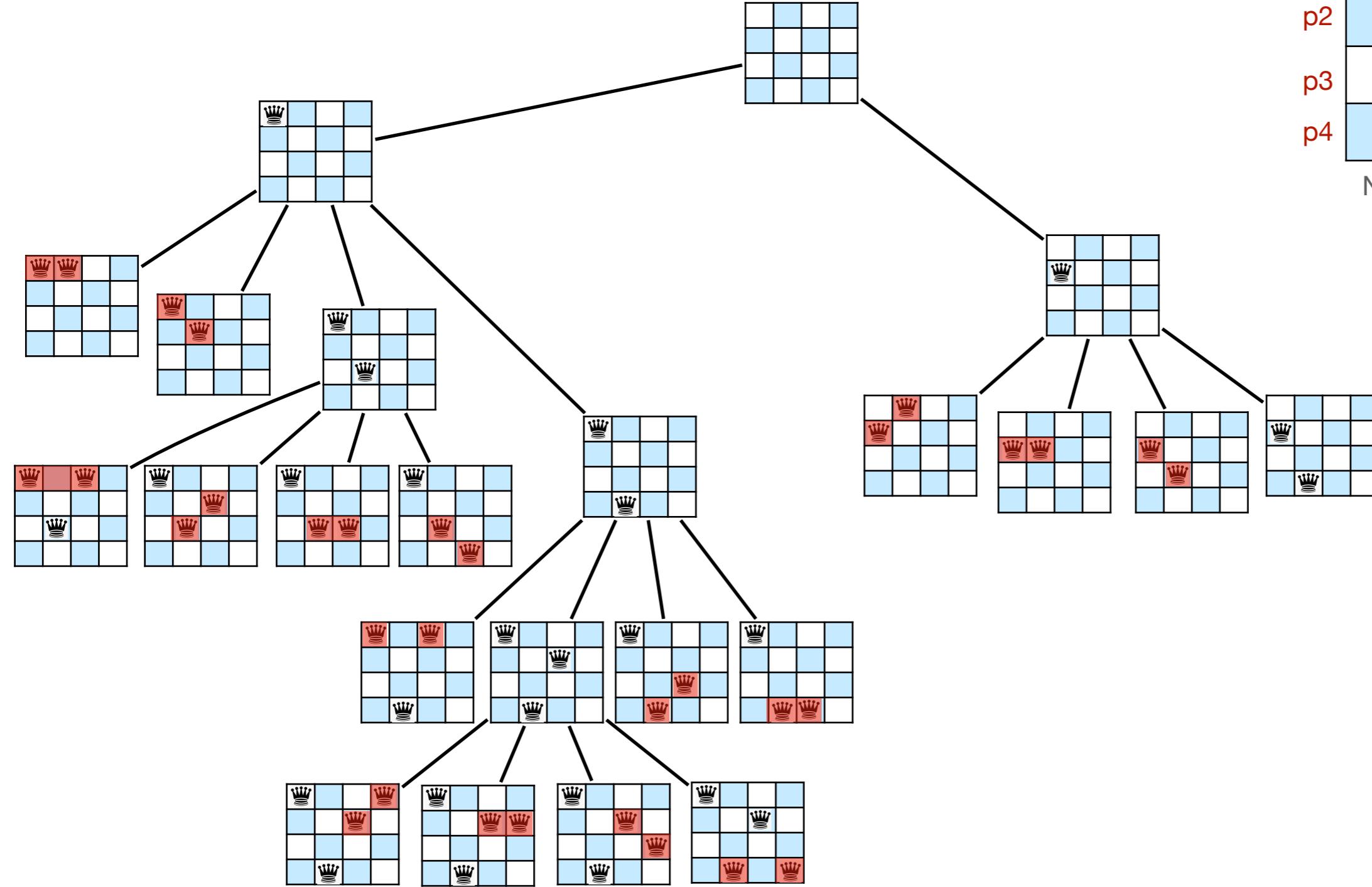


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p1				
p2				
p3				
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Constraint Programming

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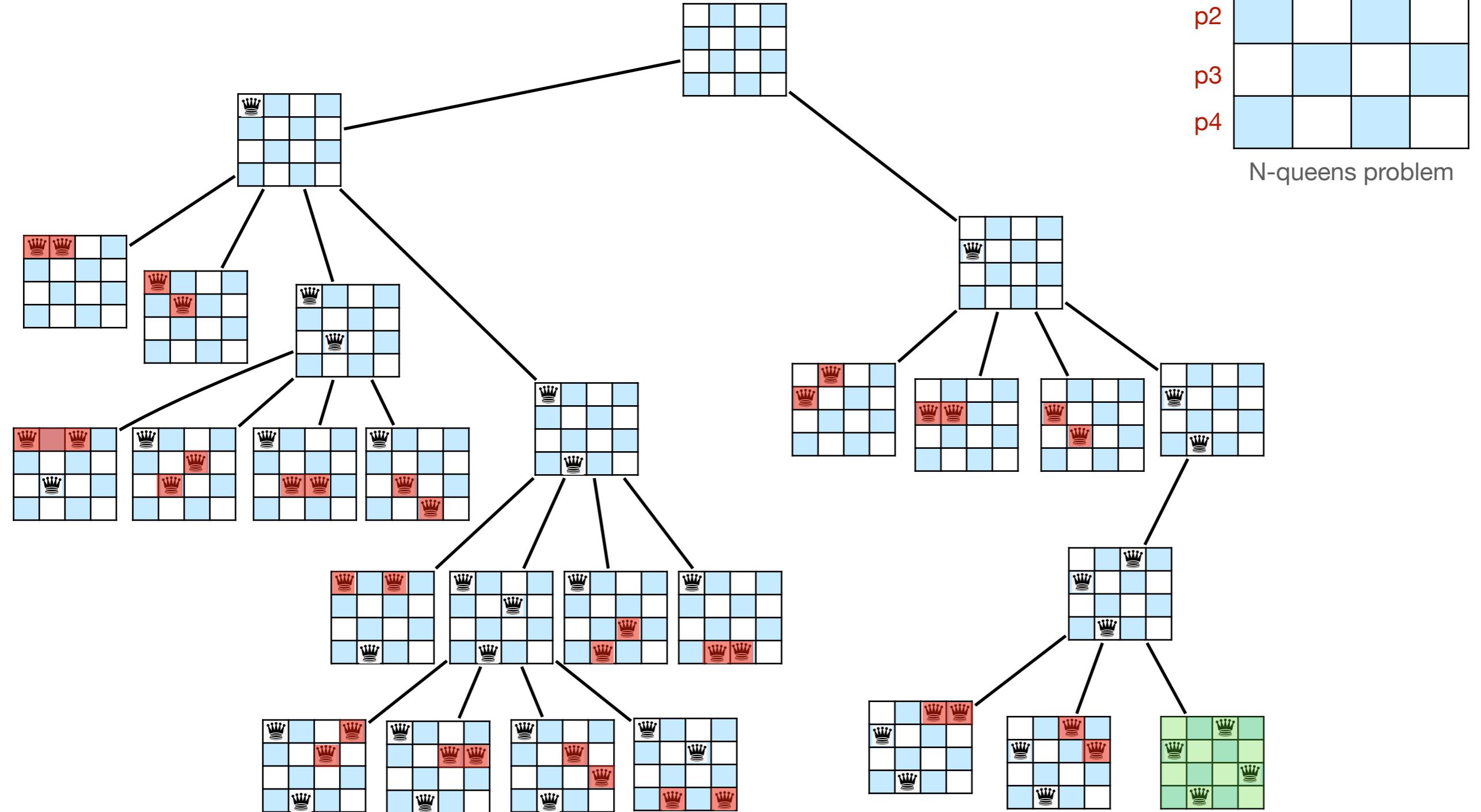


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N-queens problem

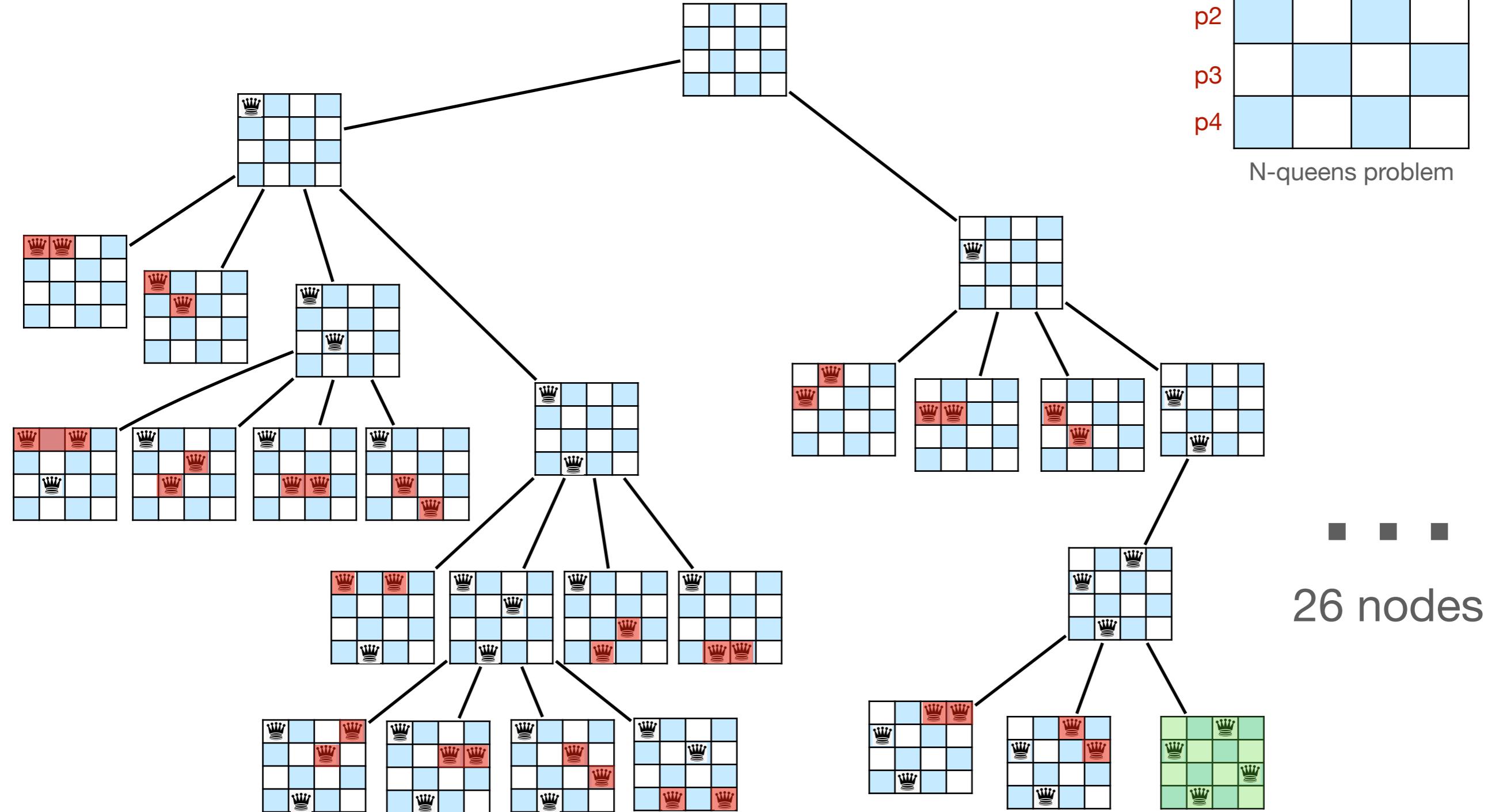
Constraint Programming

Solving - Backtracking (BT)



Constraint Programming

Solving - Backtracking (BT)



Constraint Programming

Solving - propagation (local consistency)

- AC closure implemented by several AC algorithms

Constraint Programming

Solving - propagation (local consistency)

- AC closure implemented by several AC algorithms

AC(<X, D, C>):

[Mackworth77]

Q gets $\{(X_i, c) \mid c \text{ in } C, X_i \text{ in } \text{var}(c)\}$

While Q != emptyset **do**

 pick (X_i, c) in Q

If revise(X_i, c) **then**

If $D(X_i) = \text{emptyset}$ **then** return false

Else Q gets Q union $\{(X_j, c') \mid c' \text{ in } C \text{ and } X_i, X_j \text{ in } \text{var}(c)\}$

 return true

Constraint Programming

Solving - propagation (local consistency)

Constraint Programming

Solving - propagation (local consistency)

```
revise(Xi, c):
```

[Mackworth77]

CHANGE gets false

```
ForEach v in D(Xi) do
```

```
If no allowed pairs (v, vi) for c then
```

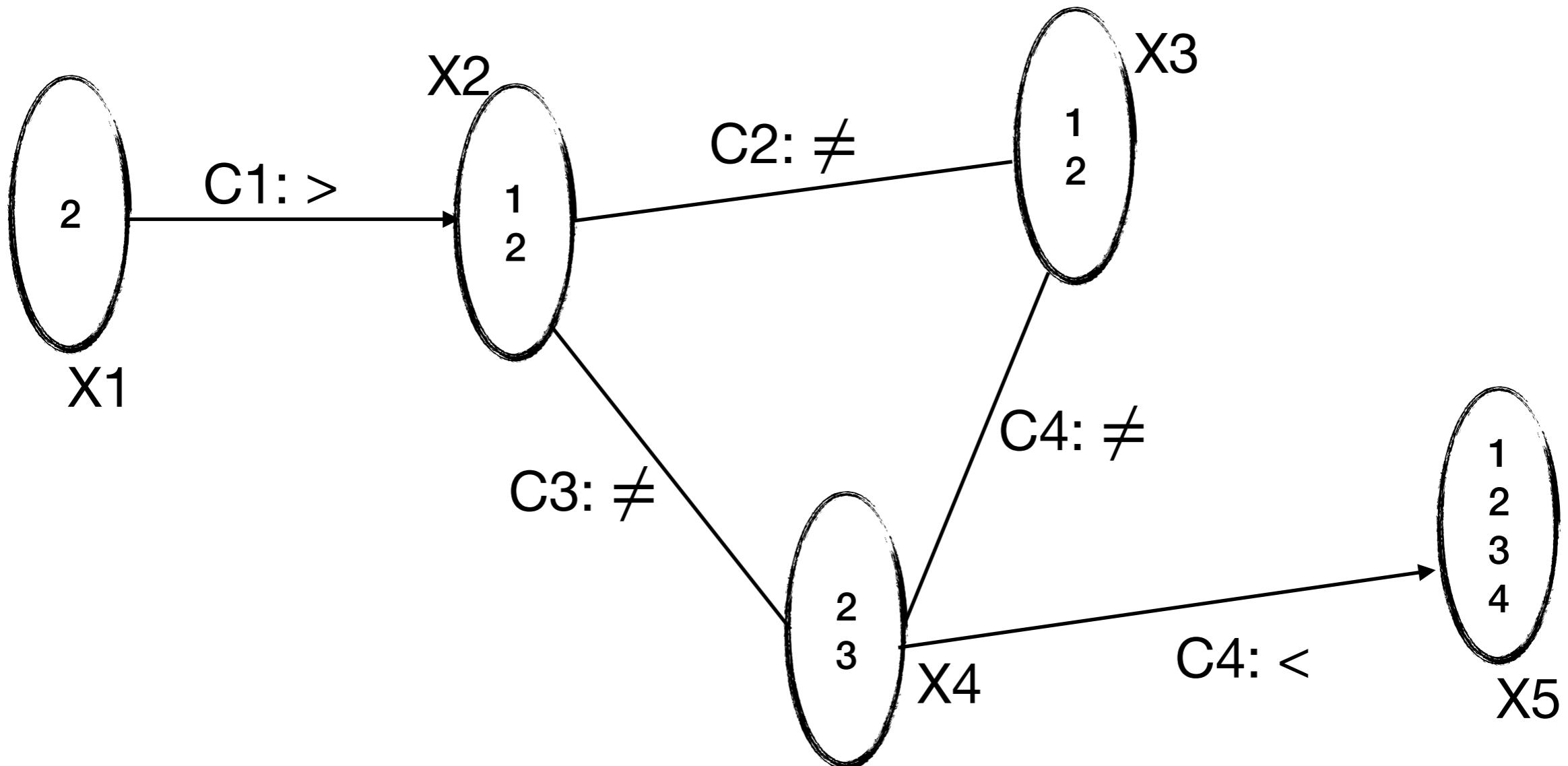
Remove v from D(Xi)

CHANGE gets true

```
return CHANGE
```

Constraint Programming

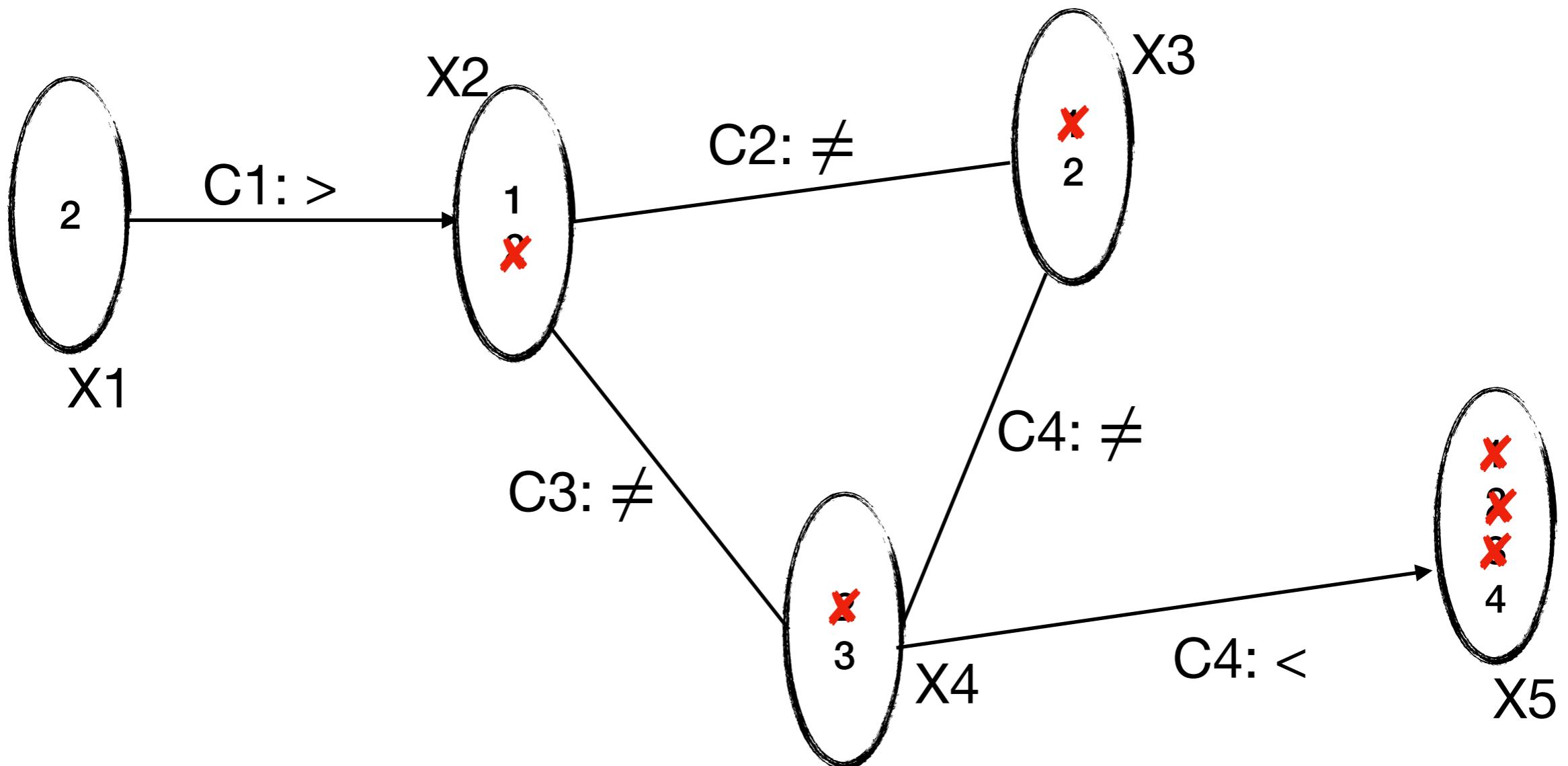
Solving - propagation (local consistency)



Situation 1

Constraint Programming

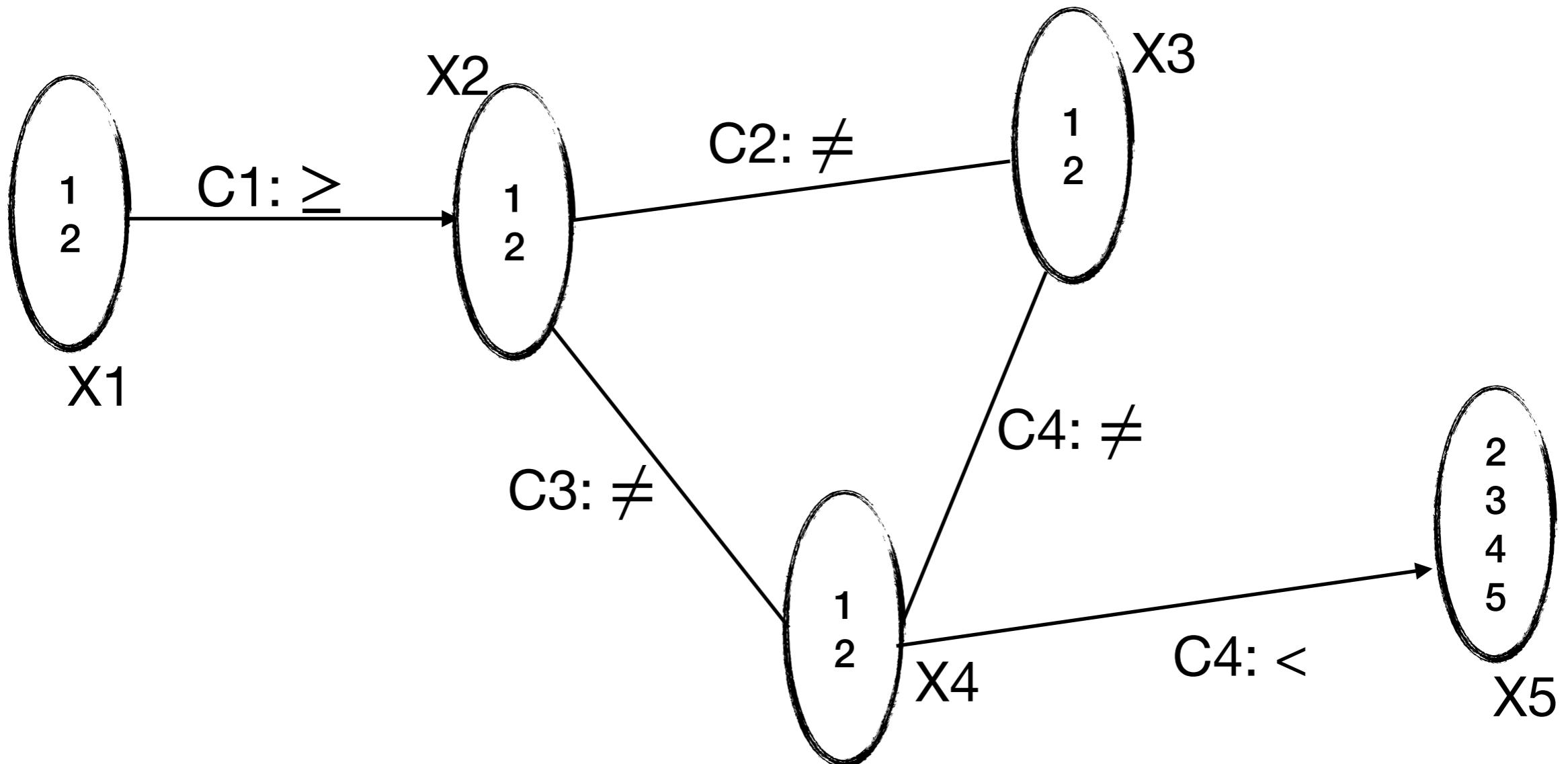
Solving - propagation (local consistency)



Situation 1 (Solution)

Constraint Programming

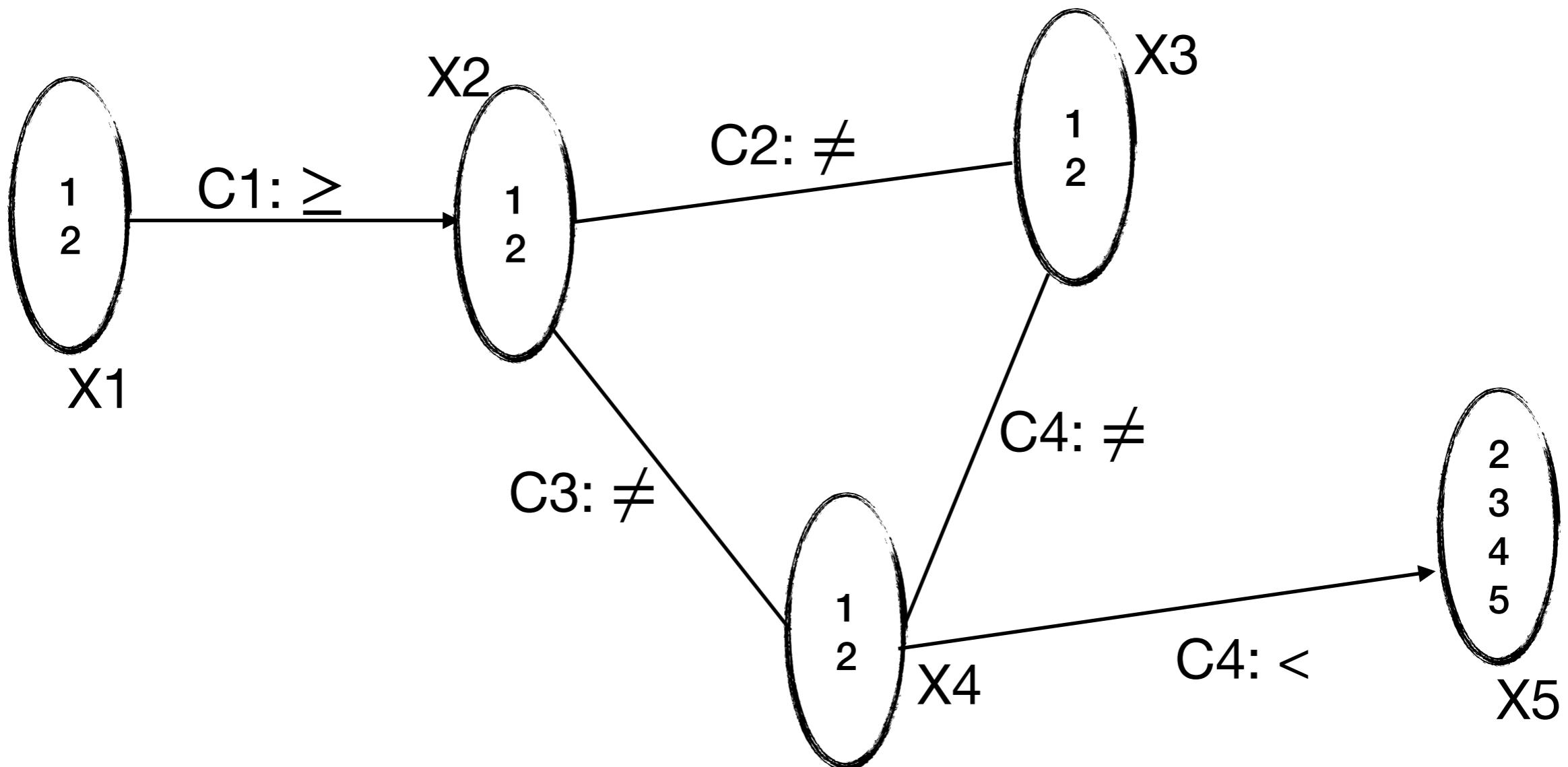
Solving - propagation (local consistency)



Situation 2

Constraint Programming

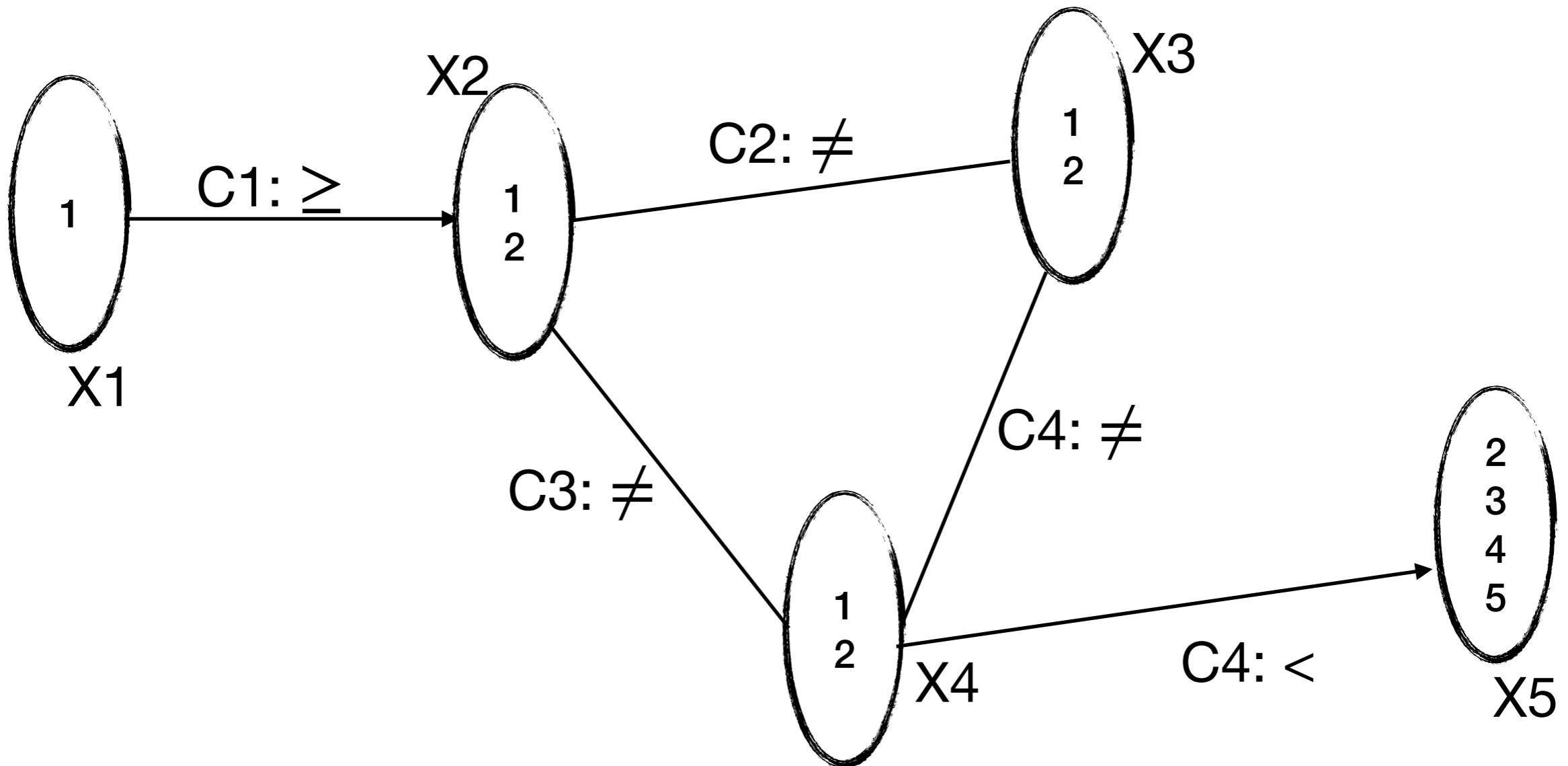
Solving - propagation (local consistency)



Situation 2 (No change)

Constraint Programming

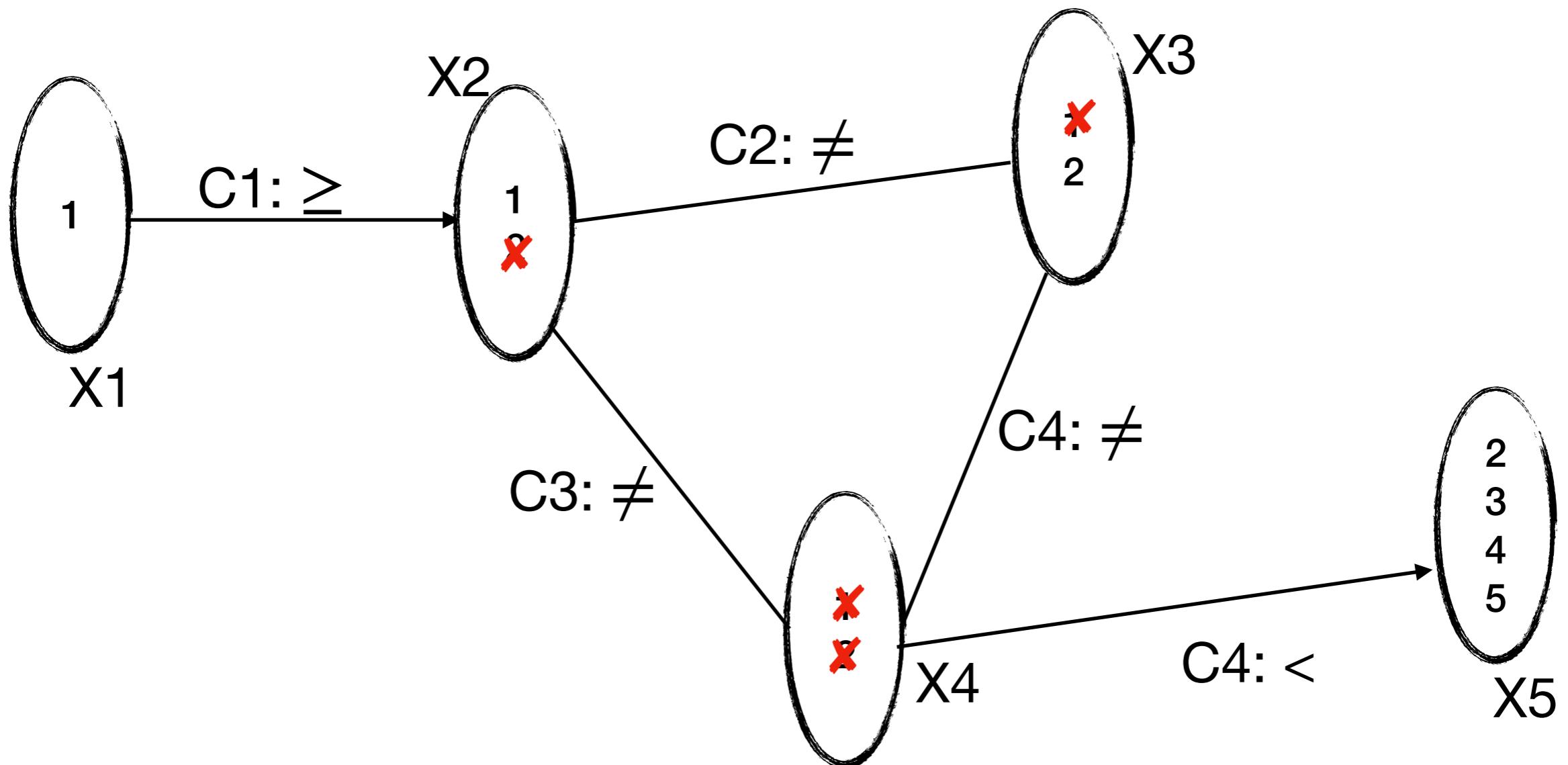
Solving - propagation (local consistency)



Situation 3

Constraint Programming

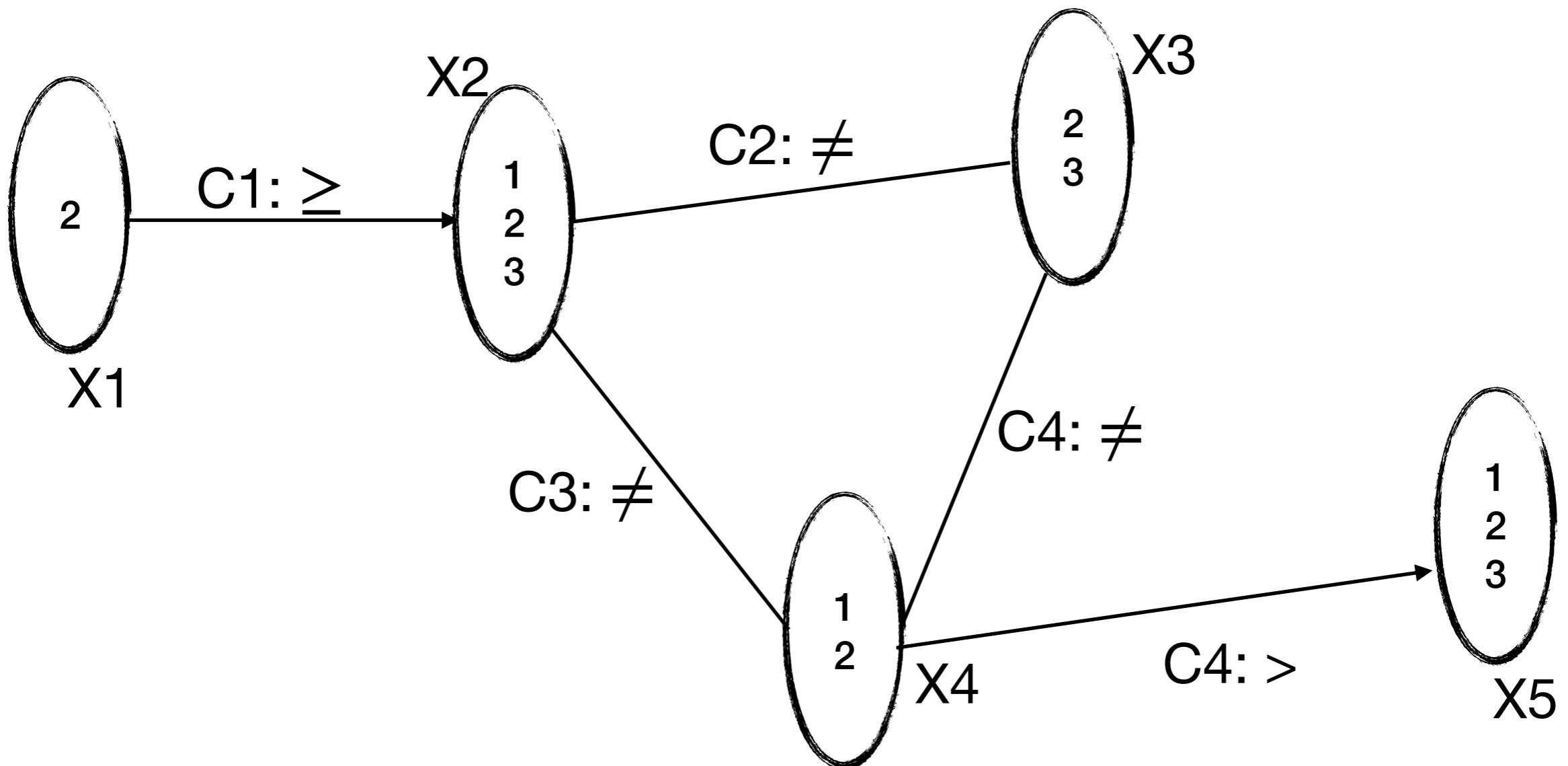
Solving - propagation (local consistency)



Situation 3 (domain wipe-out)

Constraint Programming

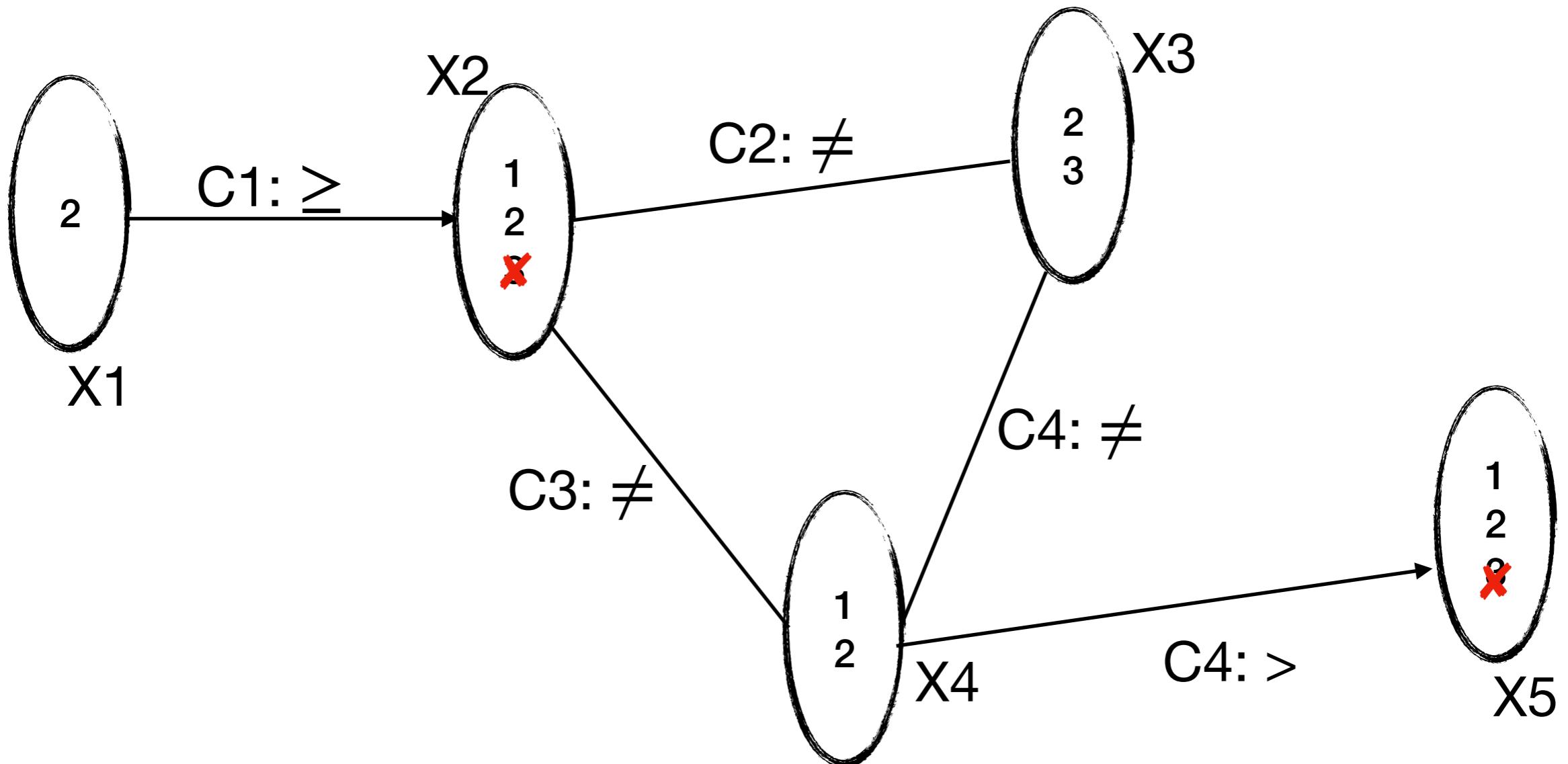
Solving - propagation (local consistency)



Situation 4

Constraint Programming

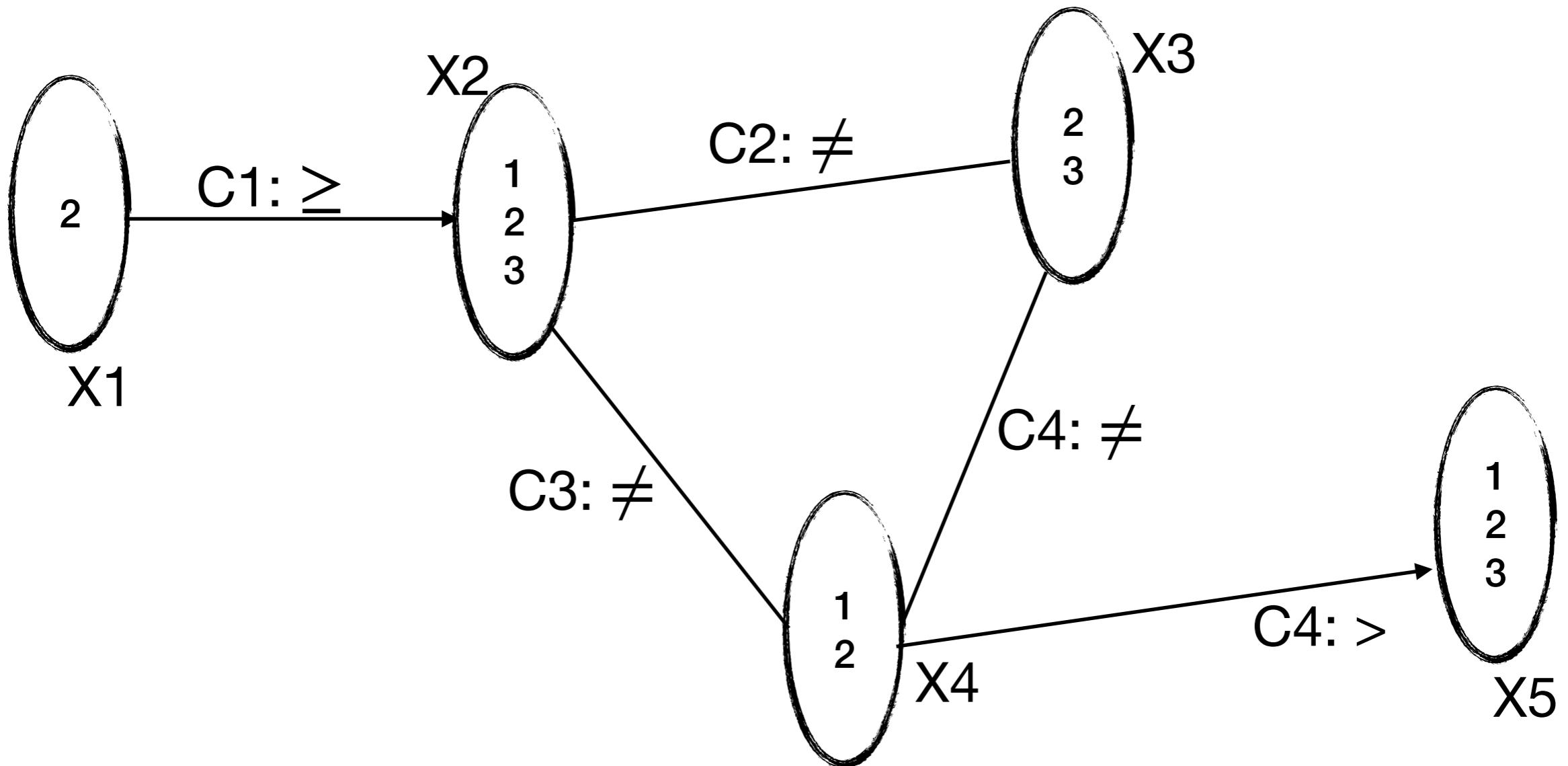
Solving - propagation (local consistency)



Situation 4 (domain reduction)

Constraint Programming

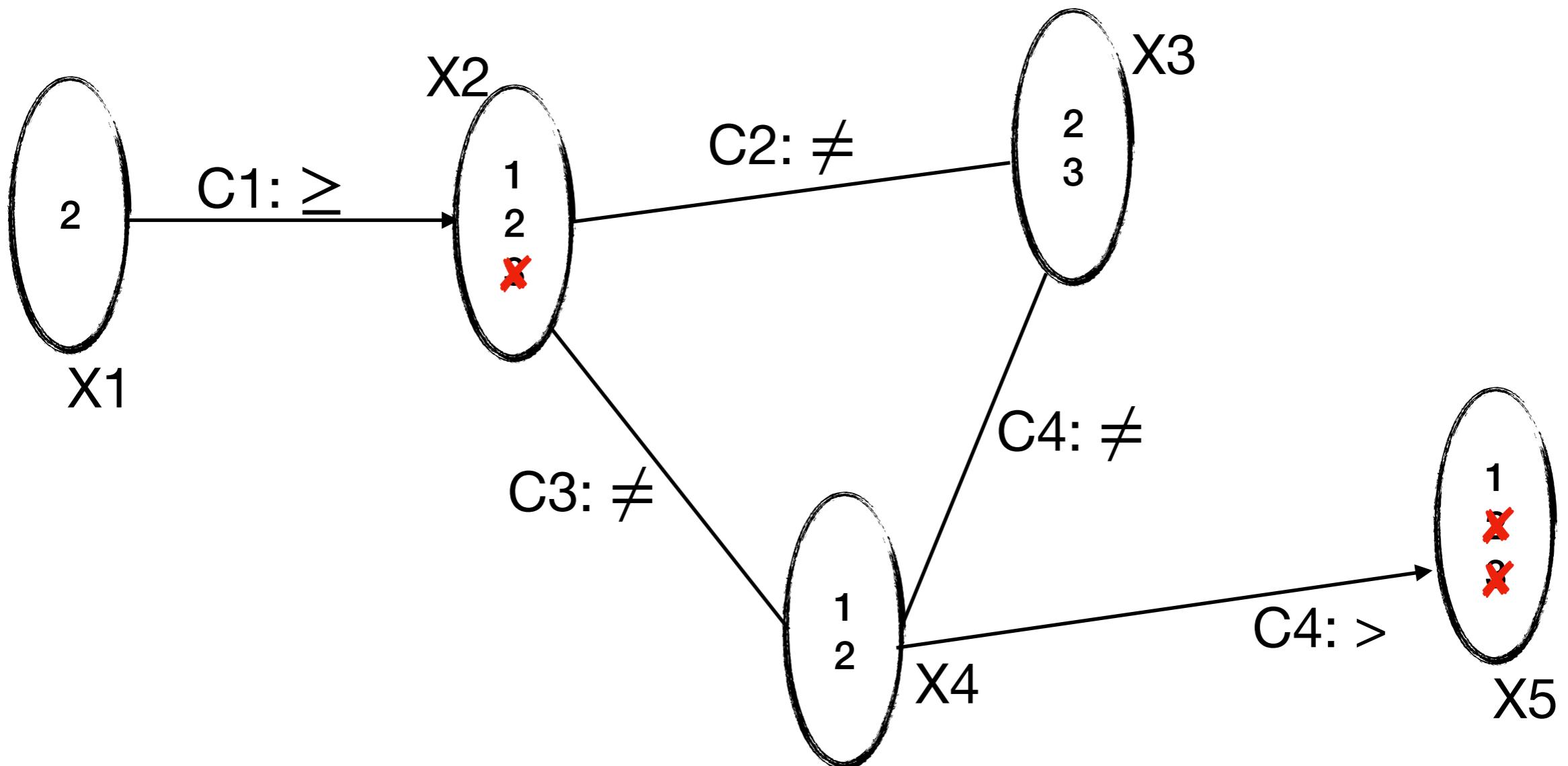
Solving - propagation (local consistency)



Situation 5

Constraint Programming

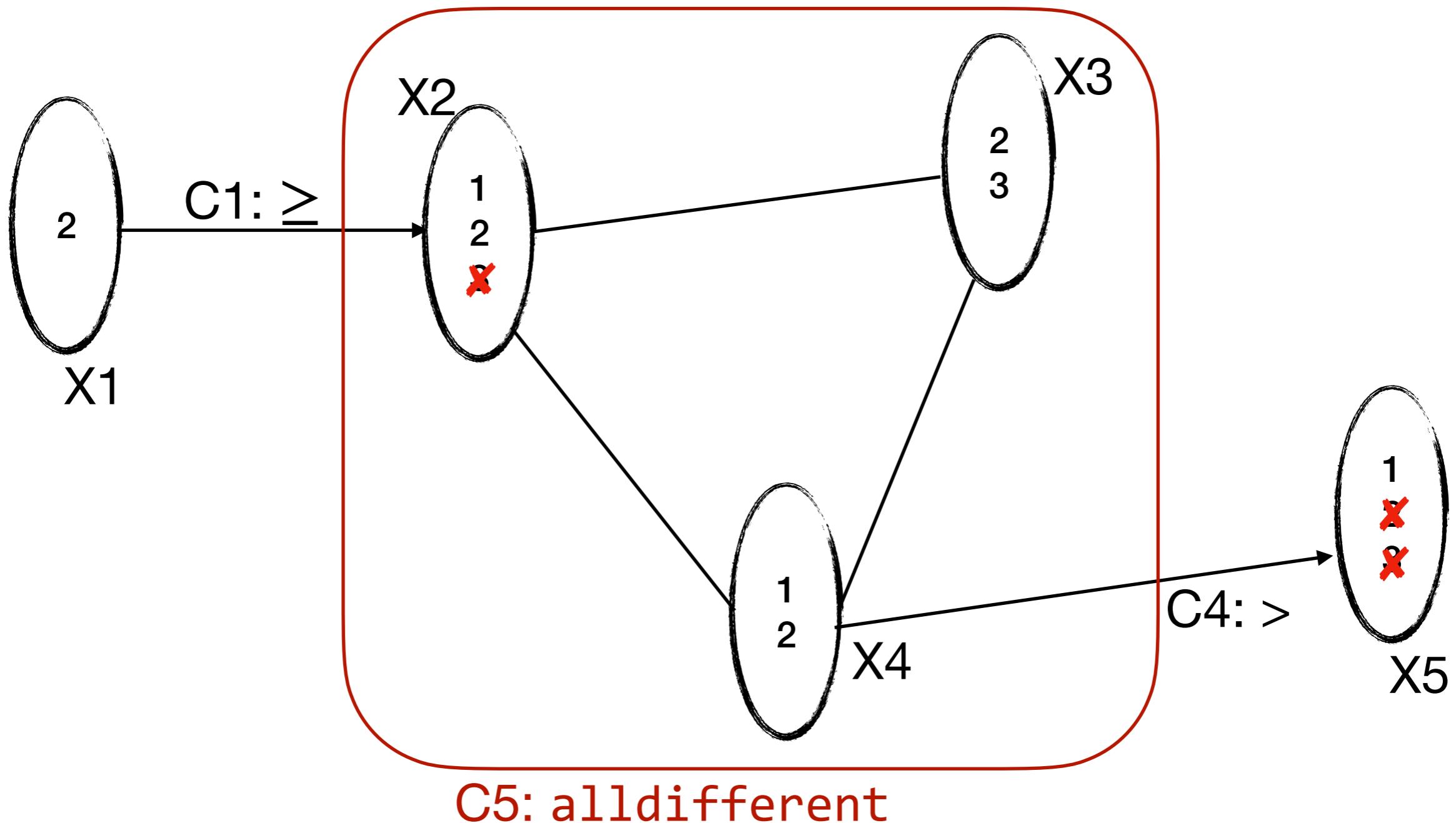
Solving - propagation (local consistency)



Situation 5

Constraint Programming

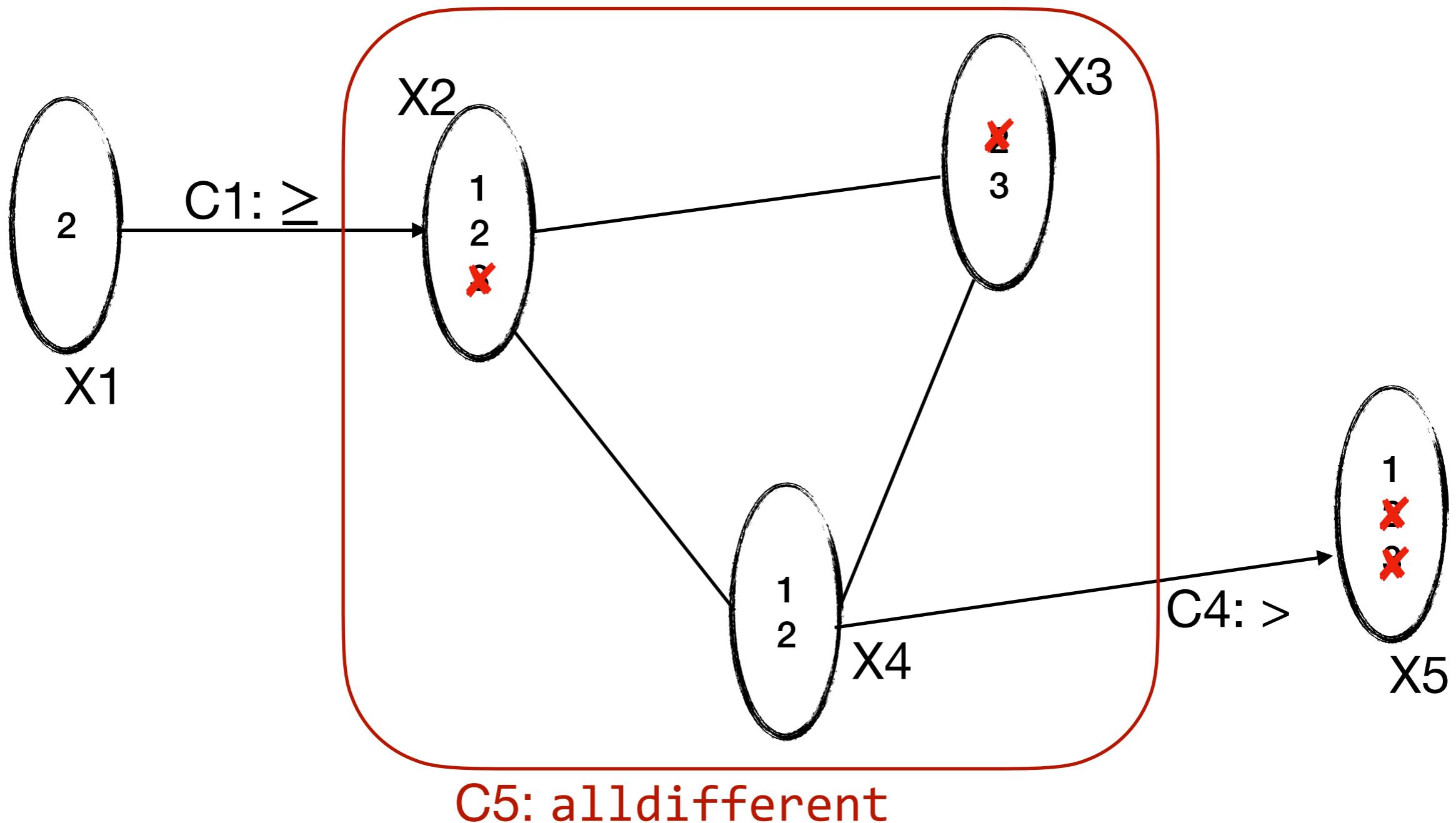
Solving - propagation (local consistency)



Situation 5 (Global constraints)

Constraint Programming

Solving - propagation (local consistency)



Situation 5 (Global constraints)

Constraint Programming

Solving - Forward Checking (FC)

- Checks if a variable's assignment is consistent
- Prunes inconsistent values
- Doesn't necessarily launch assignment propagation
- Simpler and more lightweight compared to MAC
- Suitable for medium-sized constraint problems
- Prunes values after an assignment but not necessarily after each assignment

Constraint Programming

Solving - Forward Checking (FC)

Constraint Programming

Solving - Forward Checking (FC)

FC(<X,D,C>, I): [Haralick&Elliott80, Van Hentenryck89]

If I is complete then return true

Select a variable X_i not in I

ForEach v in $D(X_i)$ do

Remove all inconsistent (v', X_j) with (v, X_i)

If no wape-out then

If **FC(<X,D,C>, I union < $X_i, v>)$** then
return true

Restore all values pruned because of (v, X_i)

return false

Constraint Programming

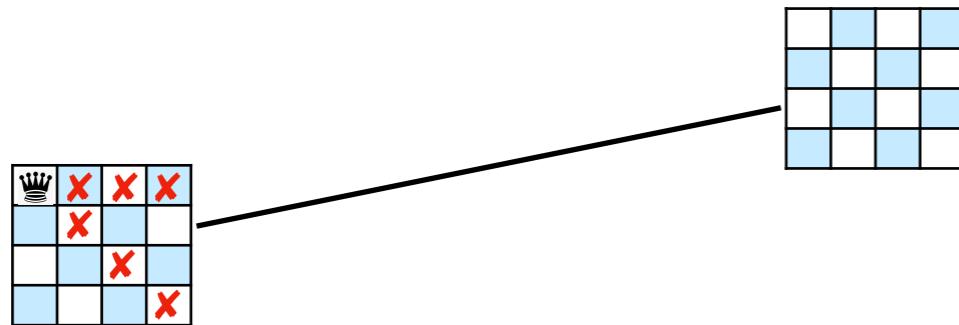
Solving - Forward Checking (FC)

	R1	R2	R3	R4
p1				
p2				
p3				
p4				

N-queens problem

Constraint Programming

Solving - Forward Checking (FC)

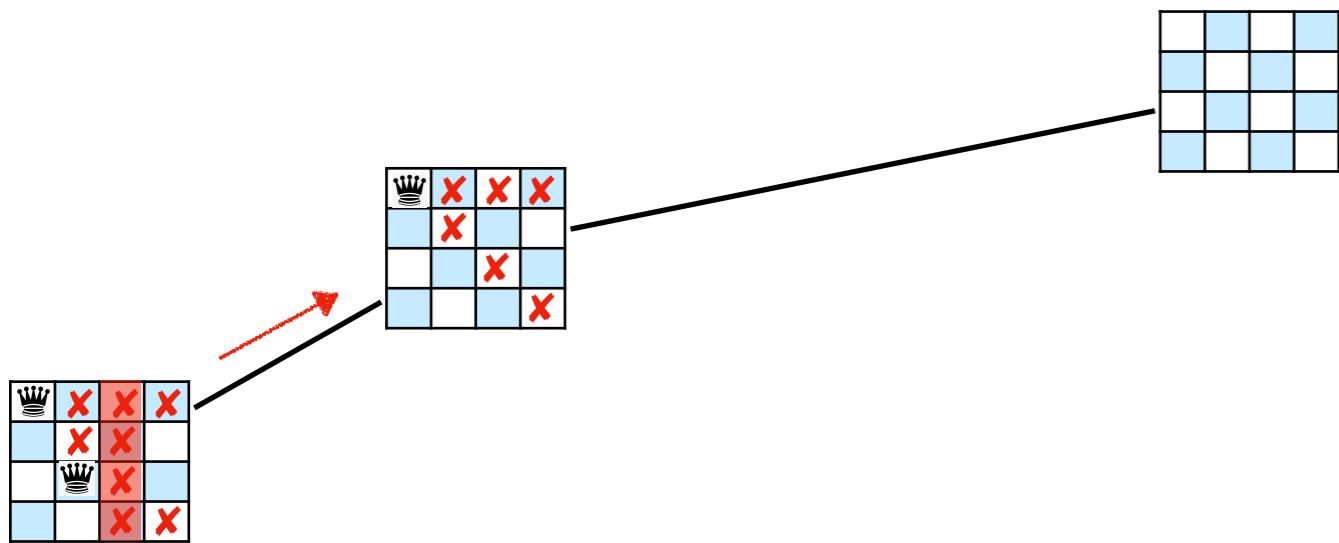


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p1				
p2				
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Constraint Programming

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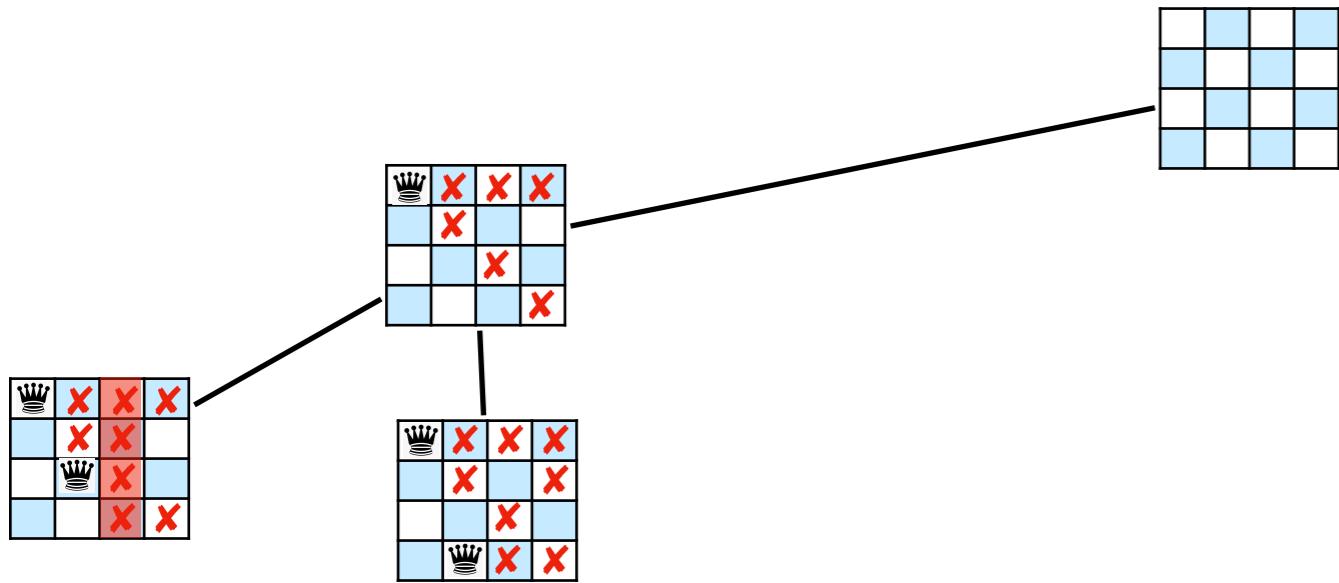


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p1				
p2				
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N-queens problem

Constraint Programming

Solving - Forward Checking (FC)

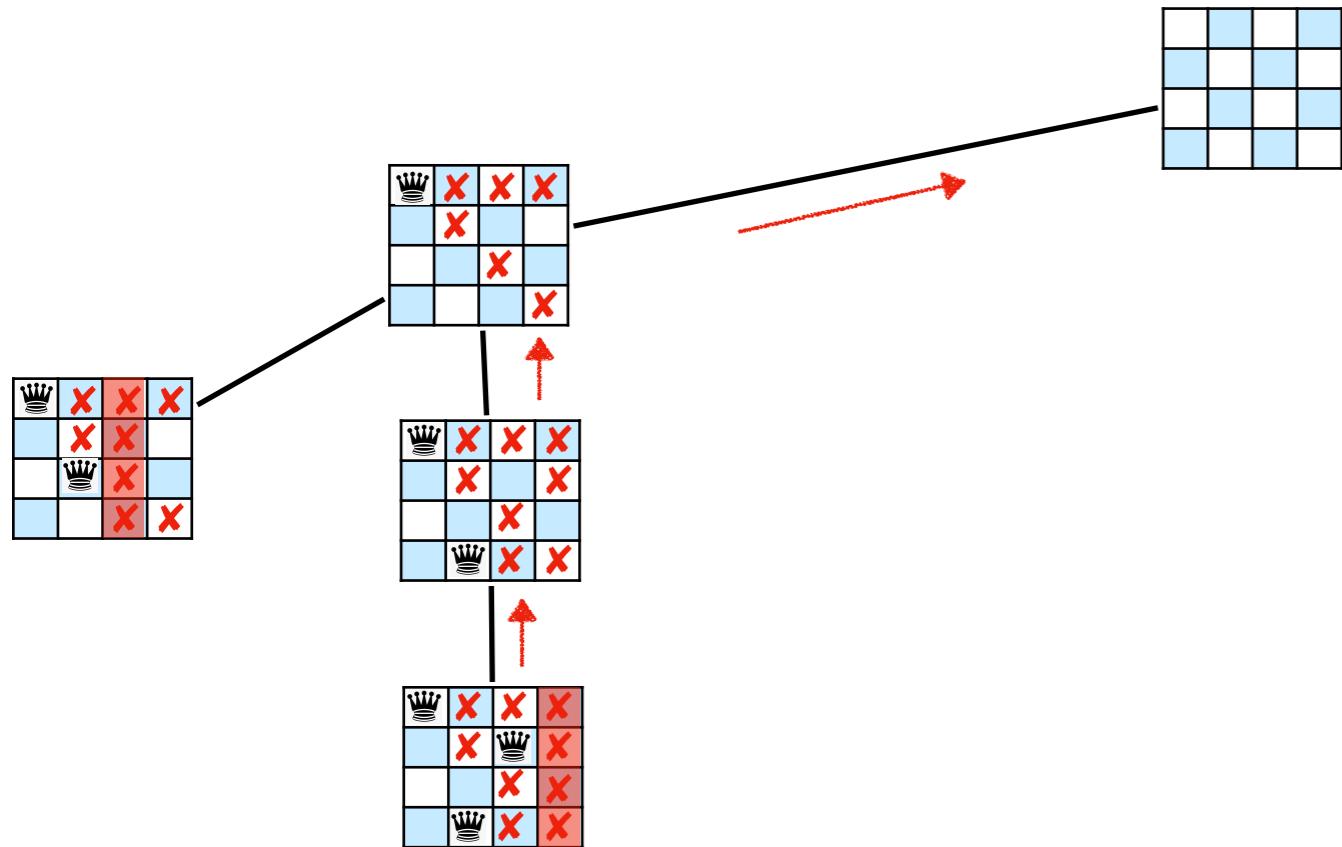


	R1	R2	R3	R4
p1				
p2				
p3				
p4				

N-queens problem

Constraint Programming

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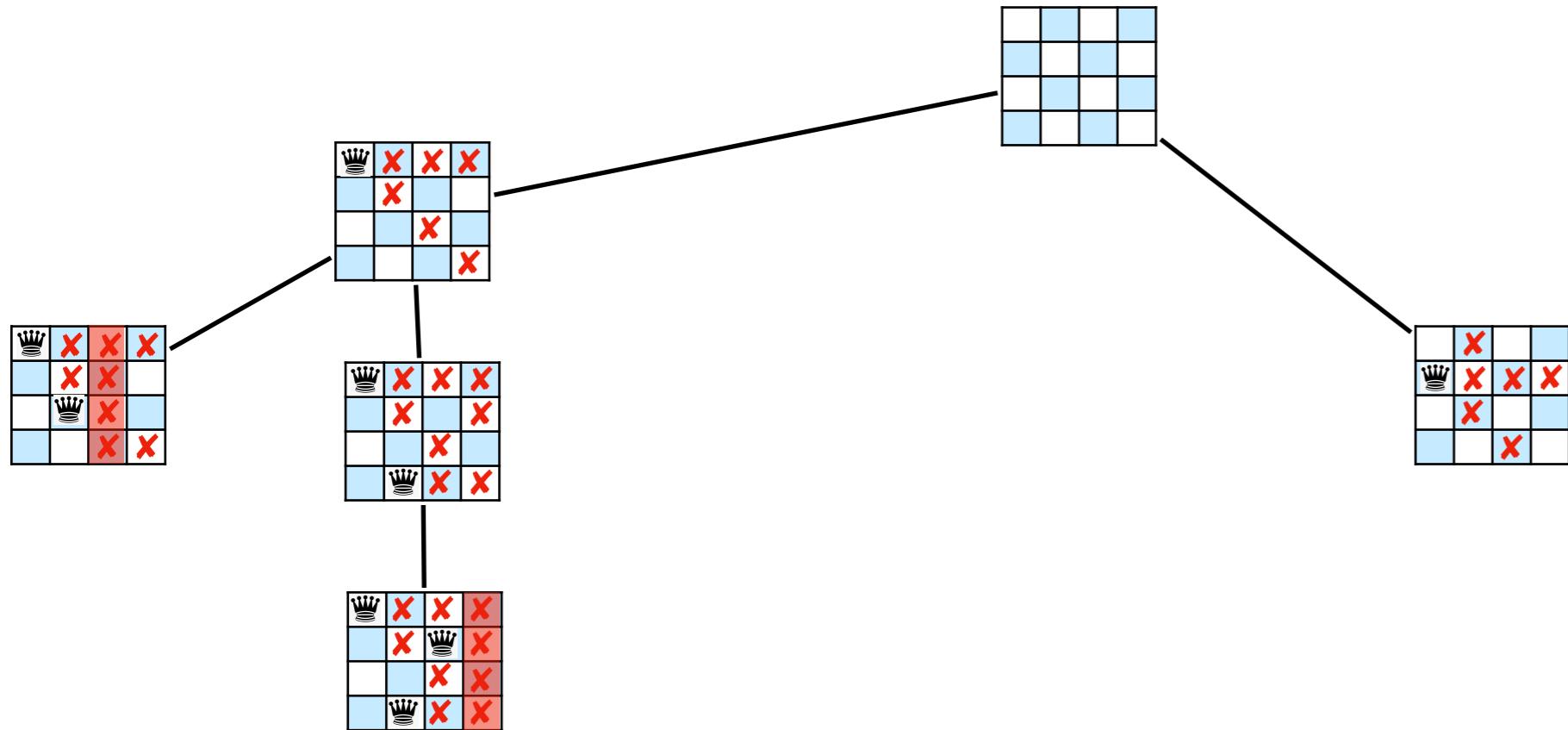


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N-queens problem

Constraint Programming

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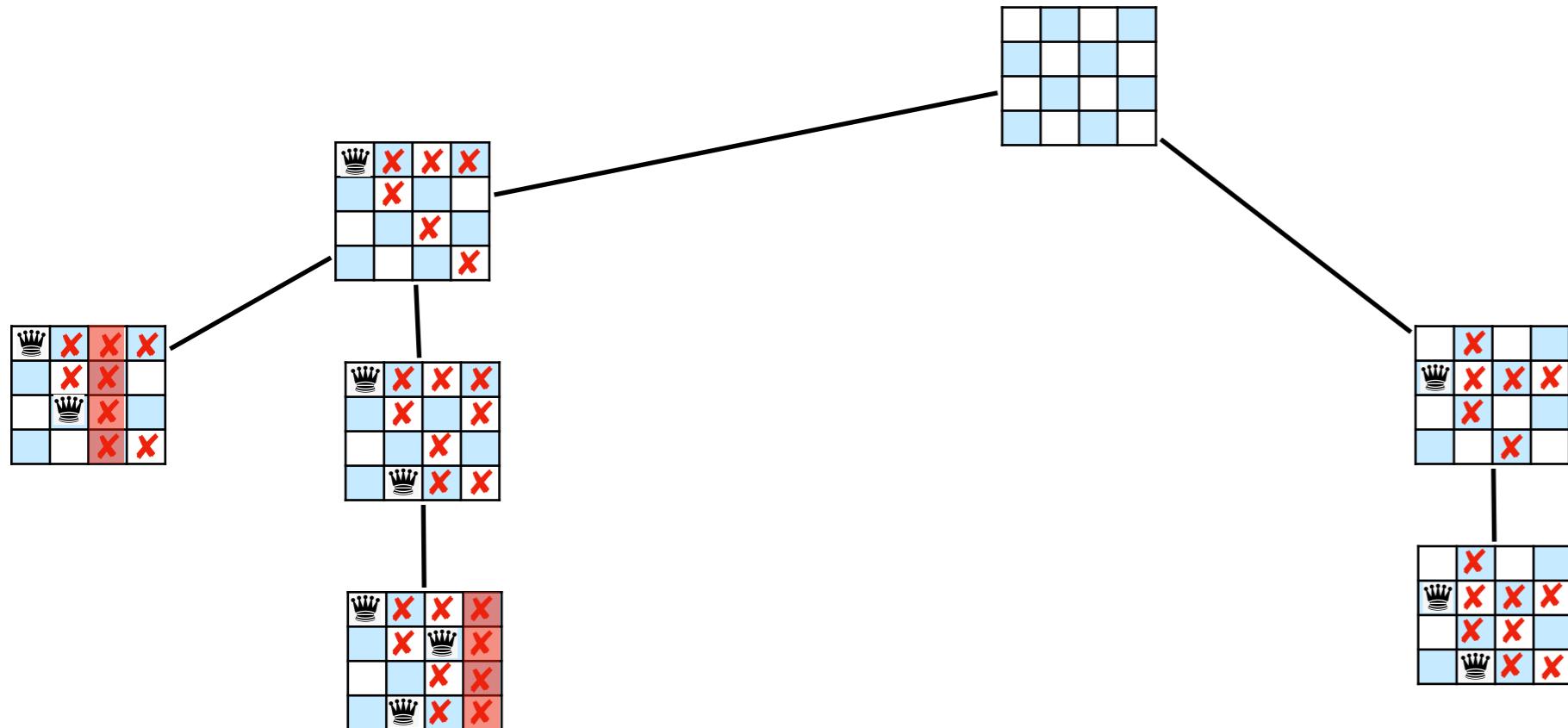


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N-queens problem

Constraint Programming

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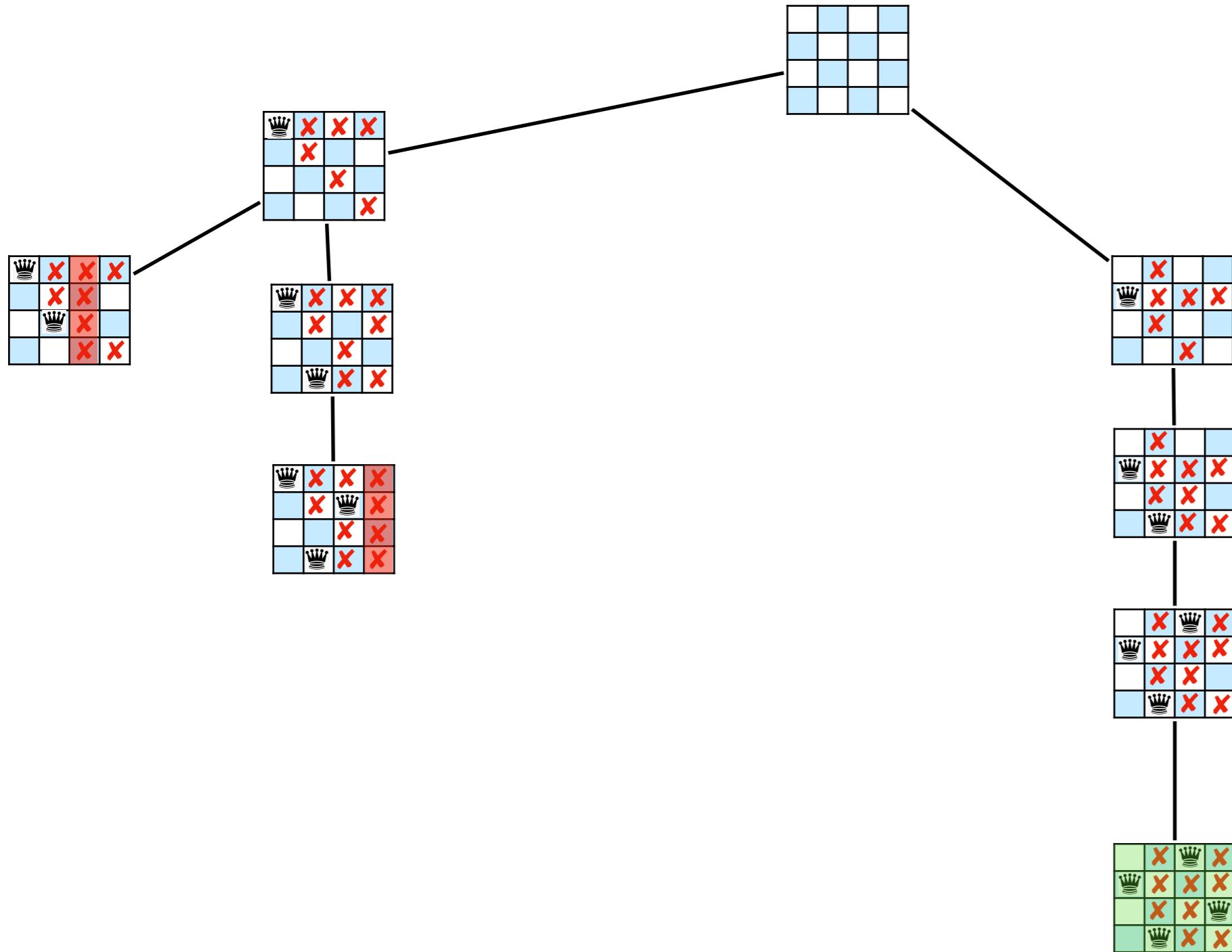


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N-queens problem

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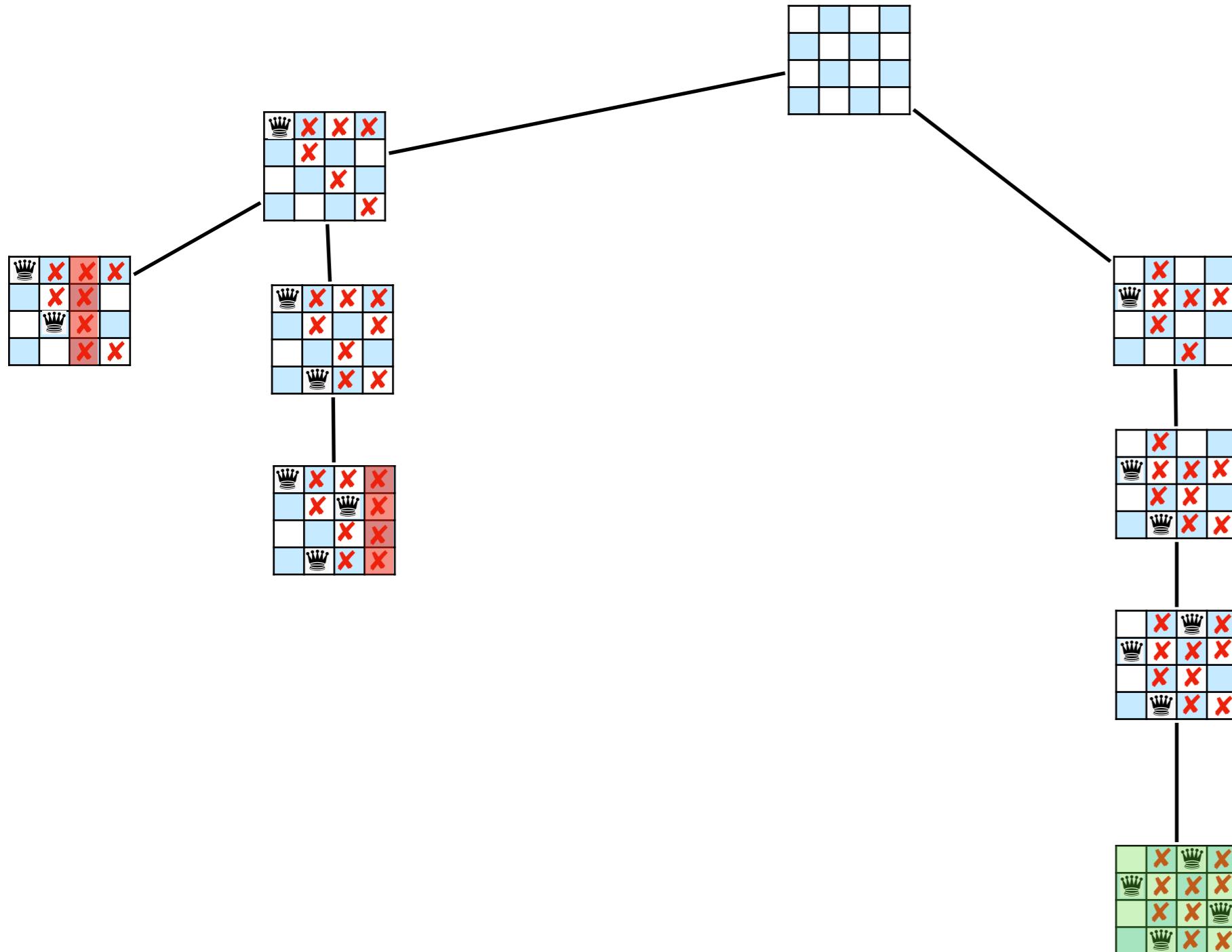


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N-queens problem

Constraint Programming

Solving - Forward Checking (FC)



	R1	R2	R3	R4
p1				
p2				
p3				
p4				

N-queens problem

■ ■ ■

8 nodes

Constraint Programming

Solving - Maintaining Arc Consistency (MAC)

- Enforces arc consistency
- Reduces the search space
- Prunes inconsistent values
- Utilizes an arc consistency restoration algorithm
- Ensures consistency after each assignment
- Suitable for complex constraint problems
- May be more computationally intensive

Constraint Programming

Solving - Maintaining Arc Consistency (MAC)

Constraint Programming

Solving - Maintaining Arc Consistency (MAC)

MAC(<X,D,C>, I):

[Sabin&Freuder94]

If AC(<X,D,C>) then

If I is complete then return true

Select a variable X_i not in I

Select a value v in $D(X_i)$

return MAC(<X,D,C>, I union < X_i, v >)

OR MAC(<X,D,C union $X_i \neq v$ >, I)

return false

Constraint Programming

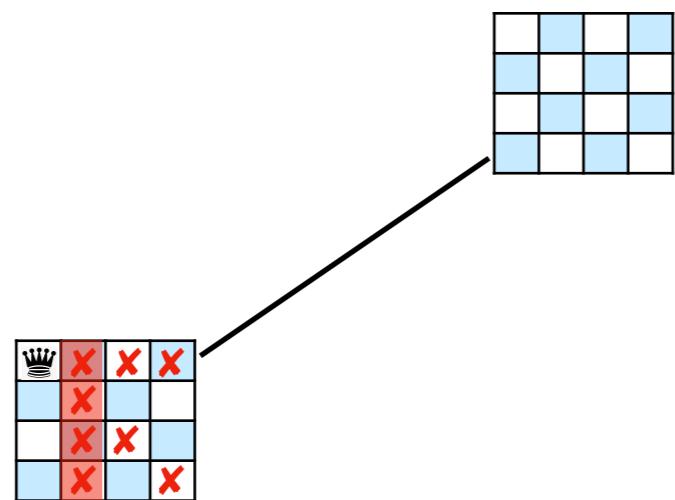
Solving - (MAC)

	R1	R2	R3	R4
p1				
p2				
p3				
p4				

N-queens problem

Constraint Programming

Solving - (MAC)

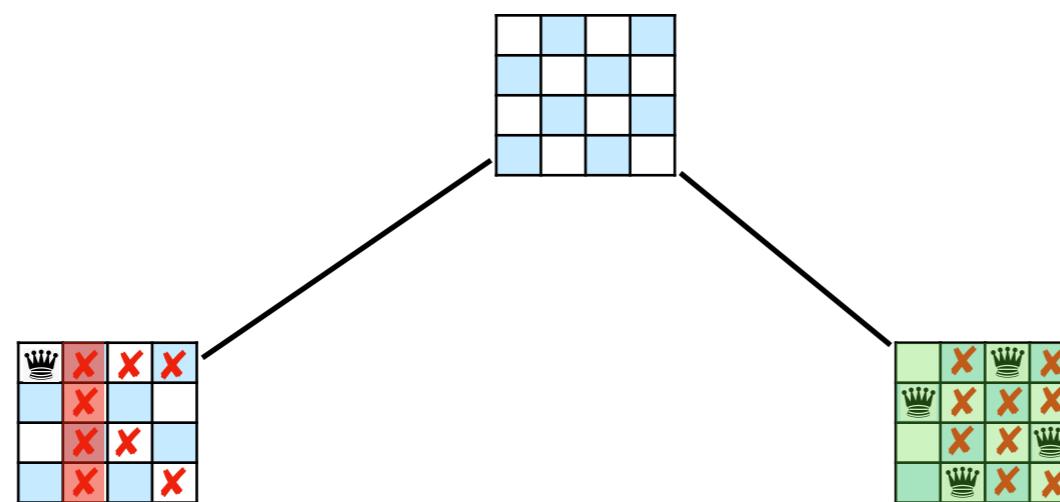


	R1	R2	R3	R4
p1				
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p3				
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N-queens problem

Constraint Programming

Solving - (MAC)

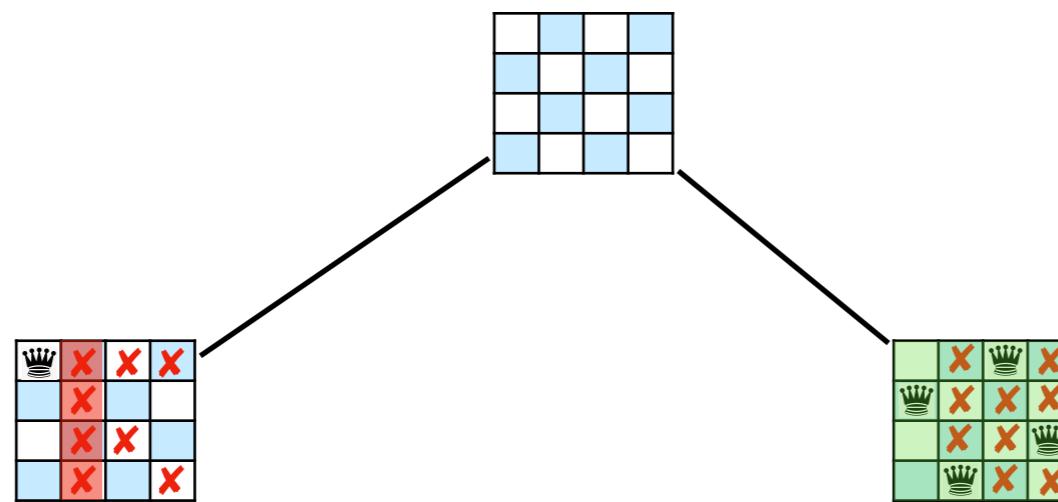


	R1	R2	R3	R4
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N-queens problem

Constraint Programming

Solving - (MAC)



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p1				
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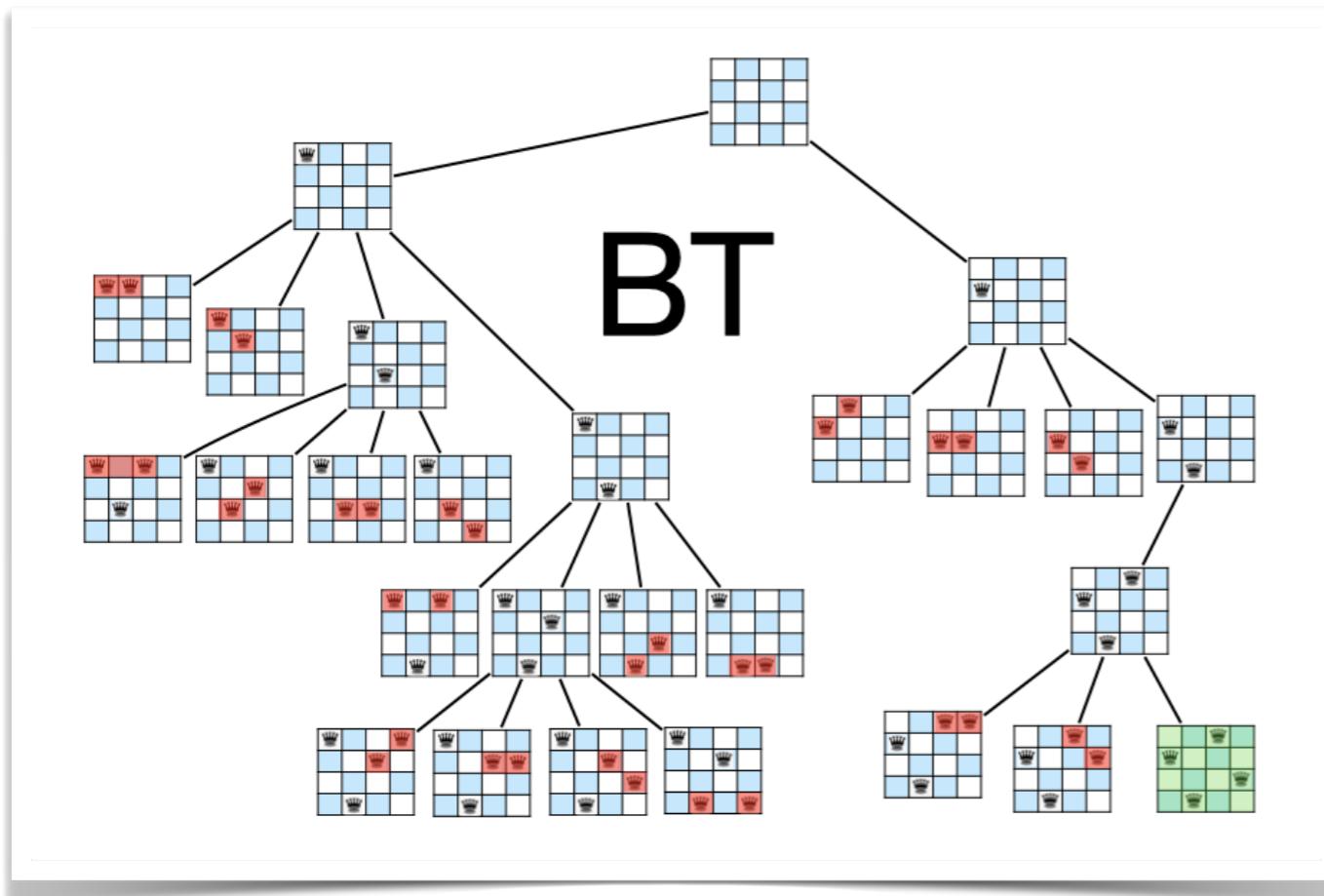
N-queens problem



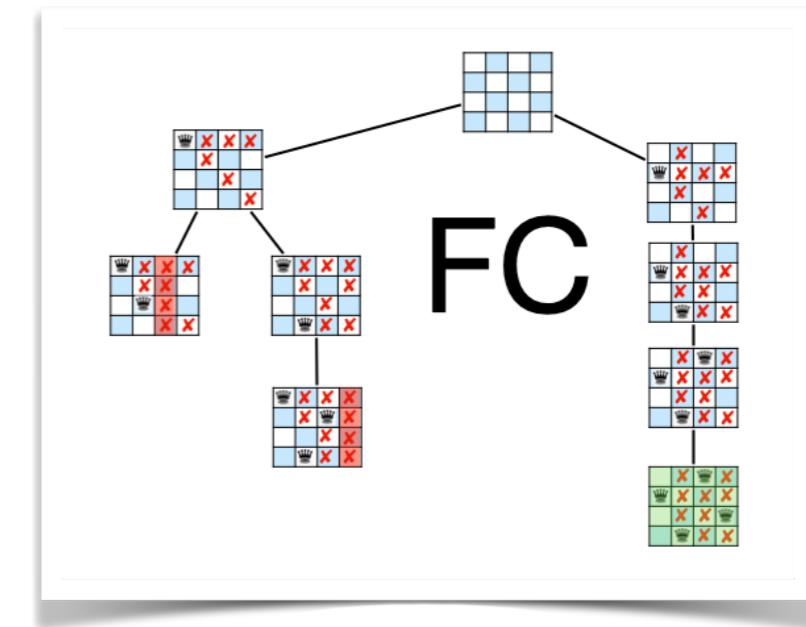
2 nodes

Constraint Programming

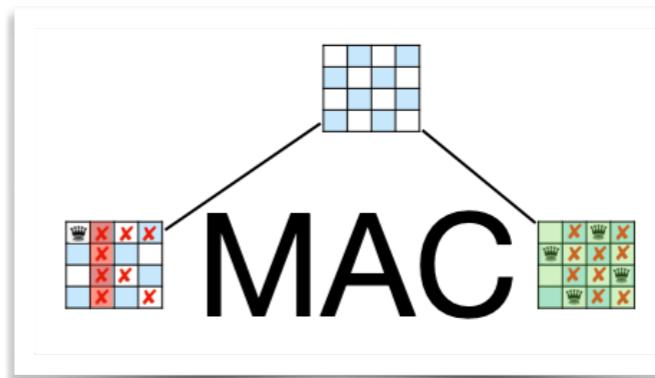
Solving - BT, MAC, FC, When to Use Each



BT



FC



MAC

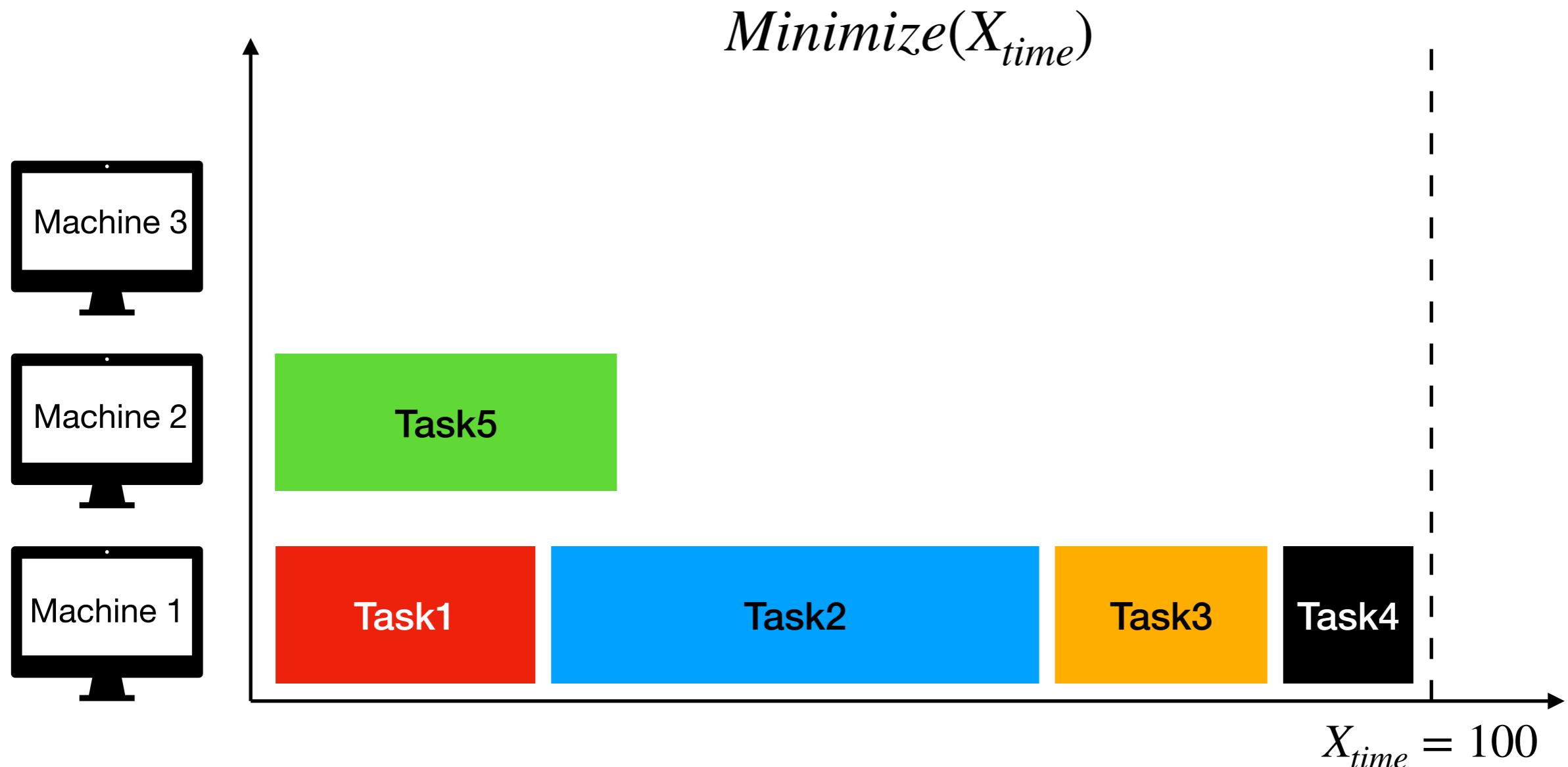
Constraint Programming

Solving - BT, MAC, FC, When to Use Each

- **Use BT** when dealing with various types of problems but be cautious of large search spaces
- **Use FC** for medium-sized problems when you need consistency checking and some pruning
- **Use MAC** for complex problems when you require strong consistency enforcement and are willing to invest more computation
- The choice depends on the problem's size, complexity, and computational resources available.
- It's a trade-off between search efficiency, consistency enforcement, and computational resources.

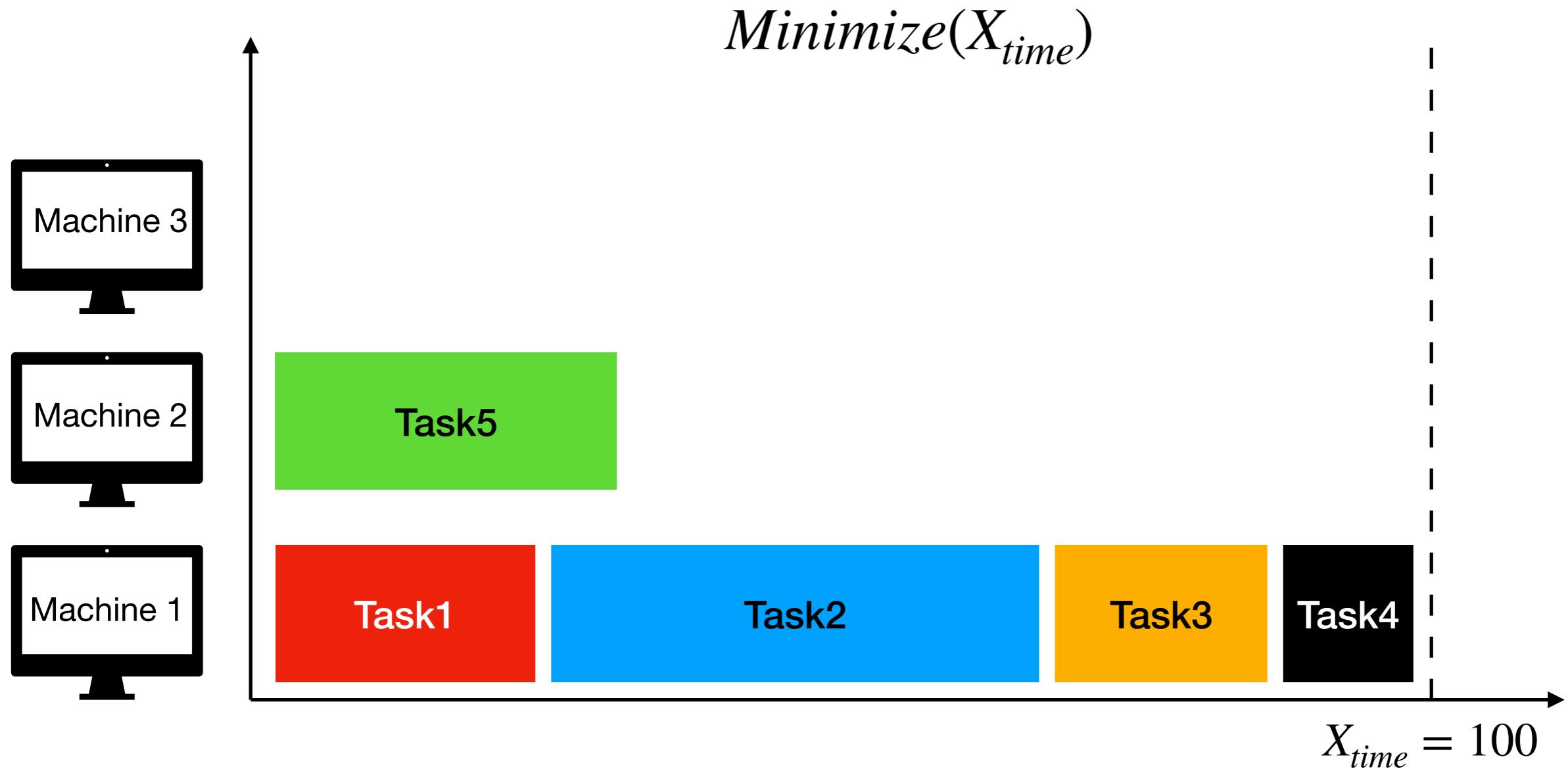
Constraint Programming

Optimization in CP (Banch&Bound)



Constraint Programming

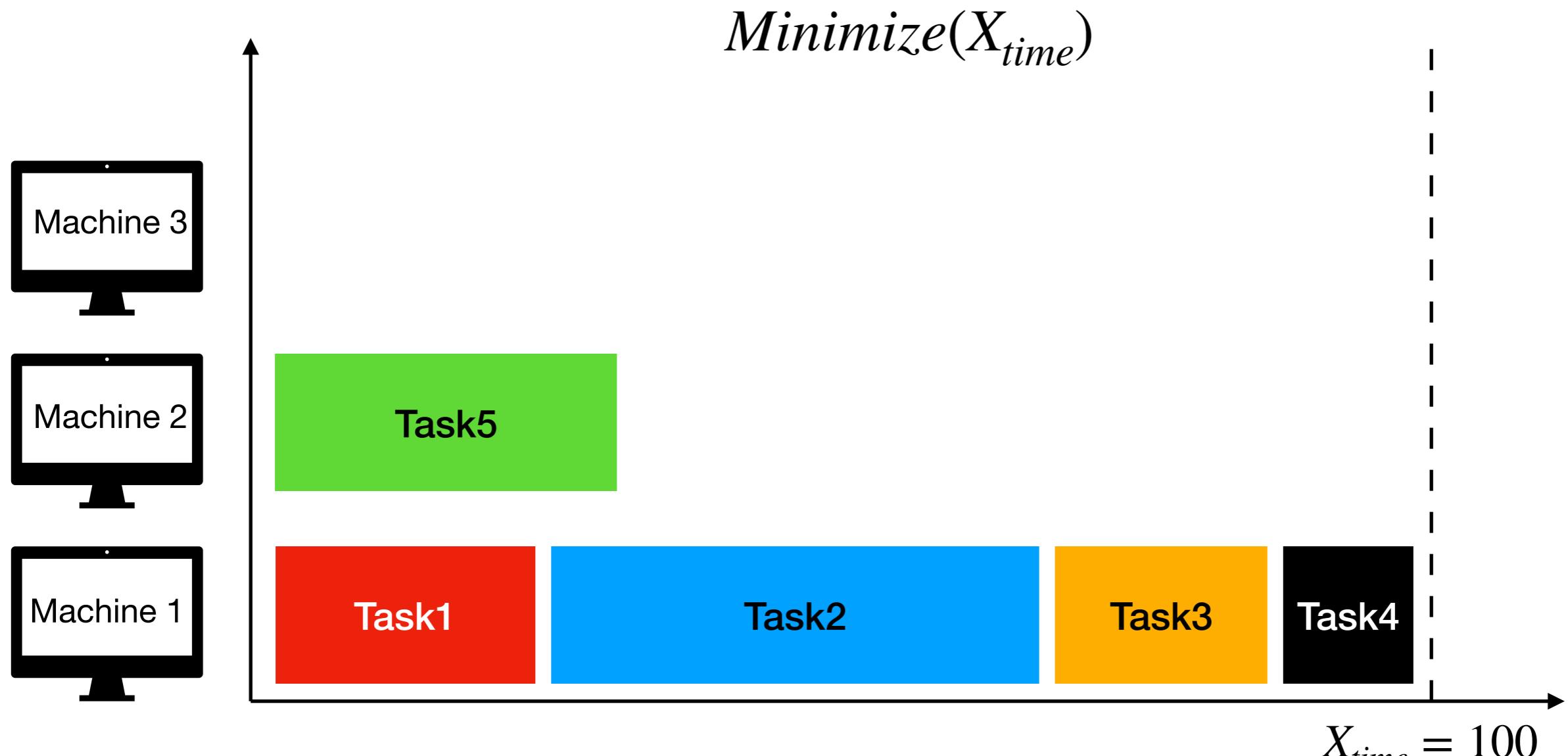
Optimization in CP (Banch&Bound)



$\langle X, D, C \rangle$

Constraint Programming

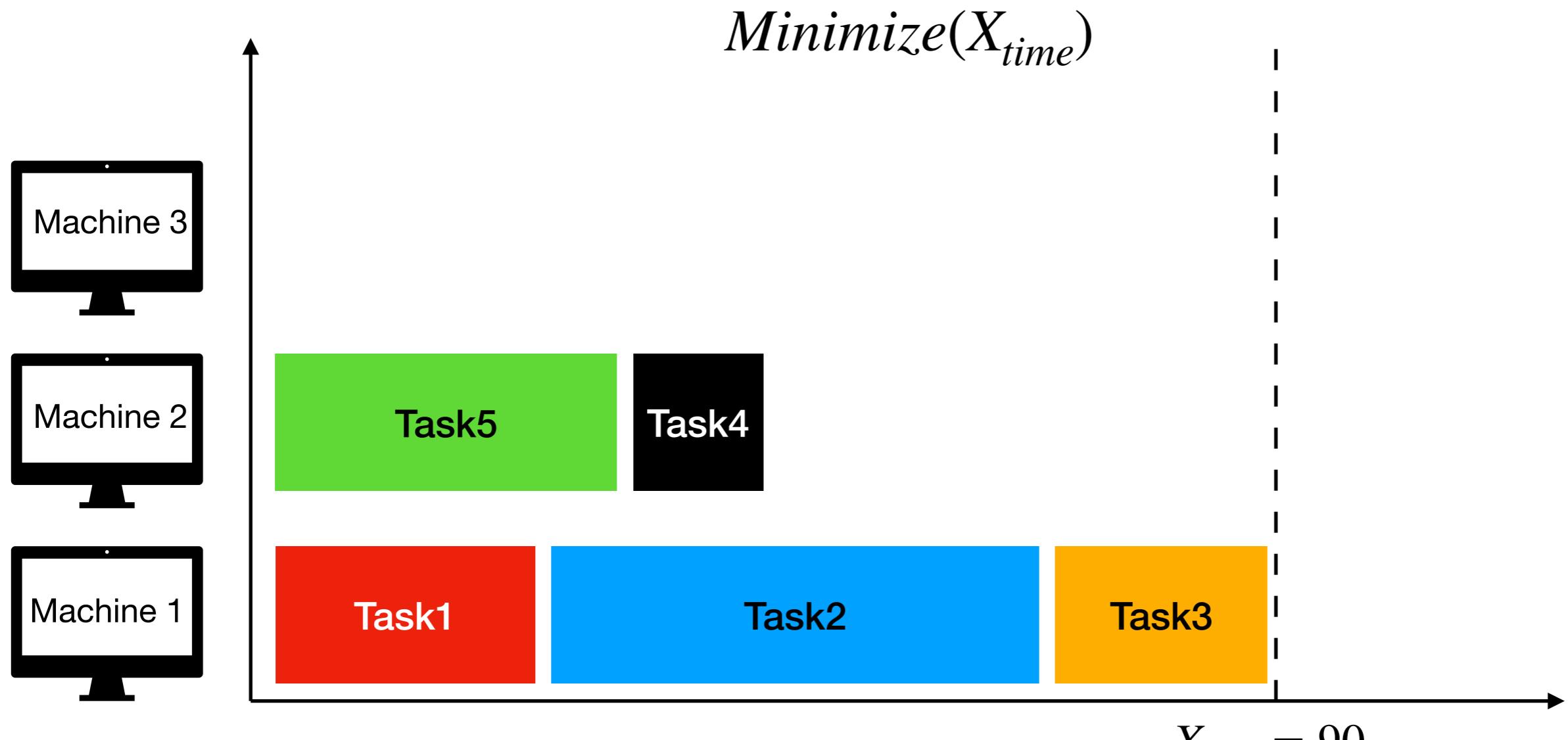
Optimization in CP (Banch&Bound)



✓
 $\langle X, D, C \rangle + X_{time} < 100$

Constraint Programming

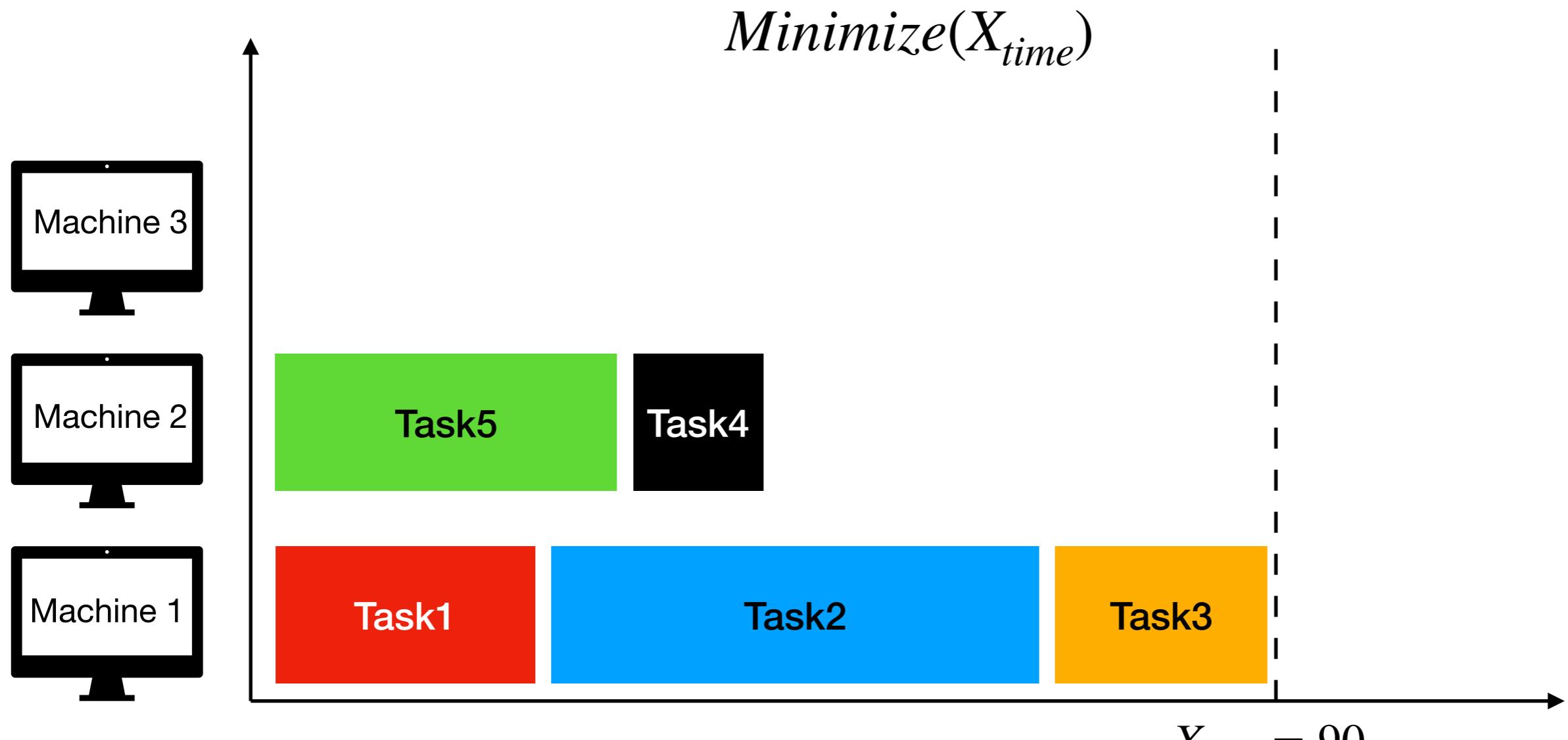
Optimization in CP (Banch&Bound)



$$\checkmark \quad \checkmark$$
$$\langle X, D, C \rangle + X_{time} < 100$$

Constraint Programming

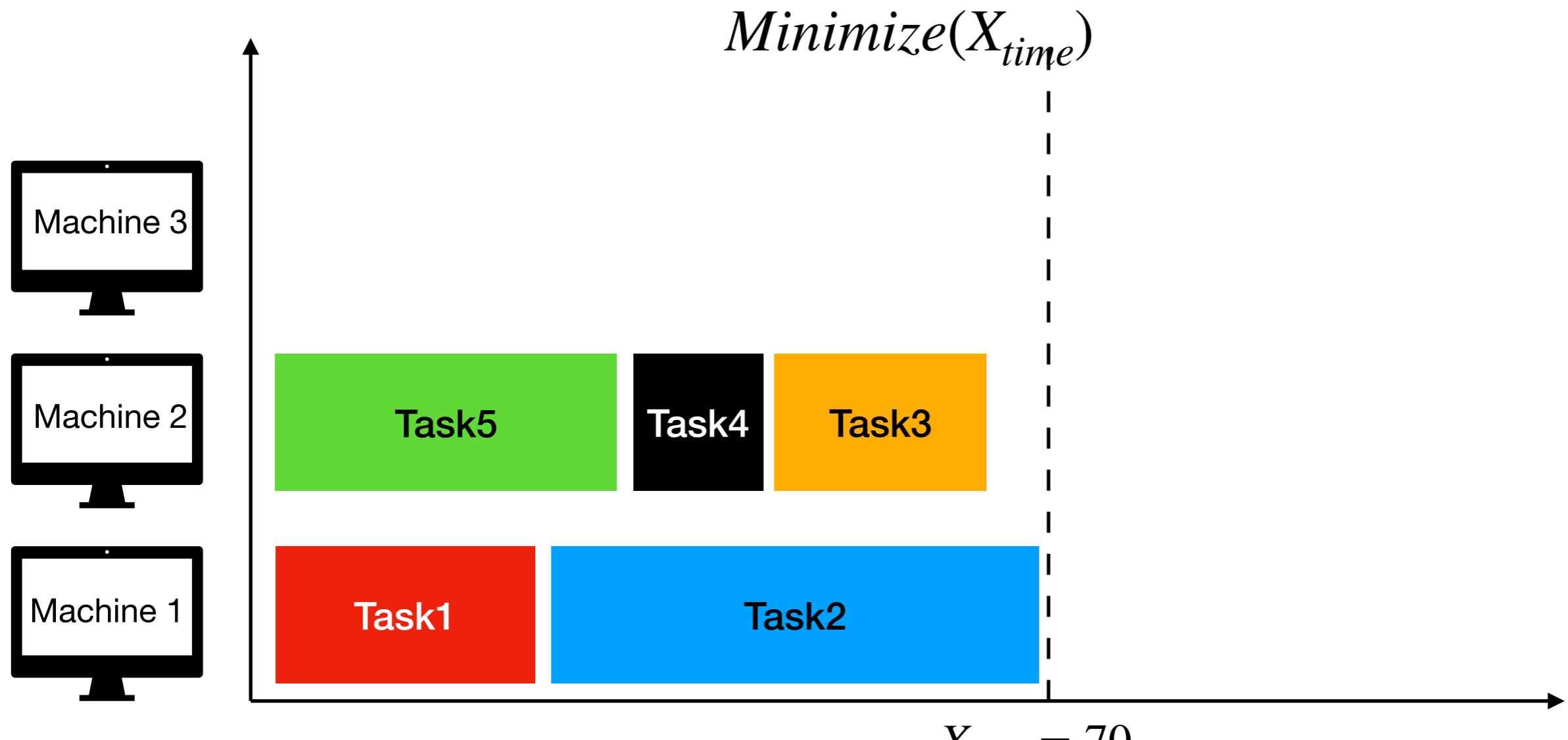
Optimization in CP (Banch&Bound)



$$\checkmark \quad \checkmark$$
$$\langle X, D, C \rangle + X_{time} < 100 + X_{time} < 90$$

Constraint Programming

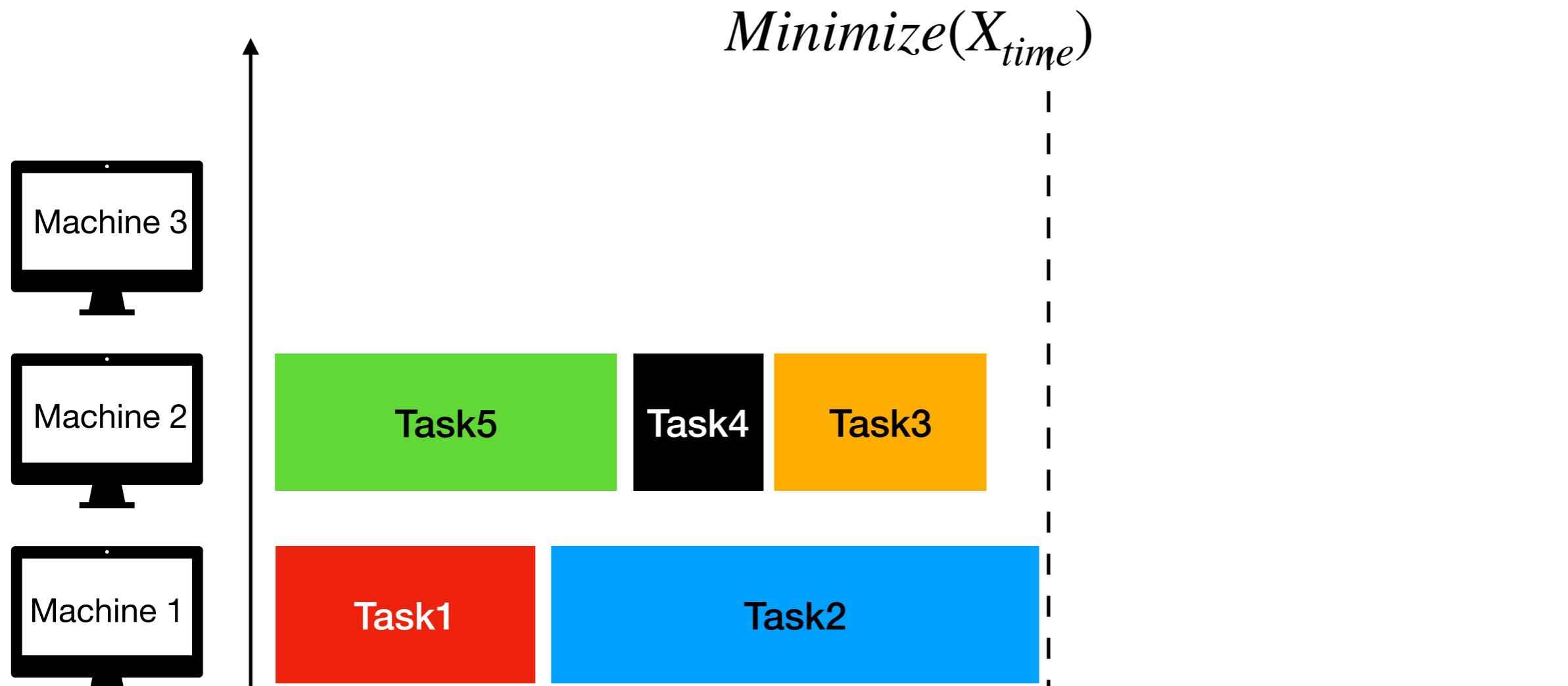
Optimization in CP (Banch&Bound)



$$\langle X, D, C \rangle + X_{time} < 100 + X_{time} < 90$$

Constraint Programming

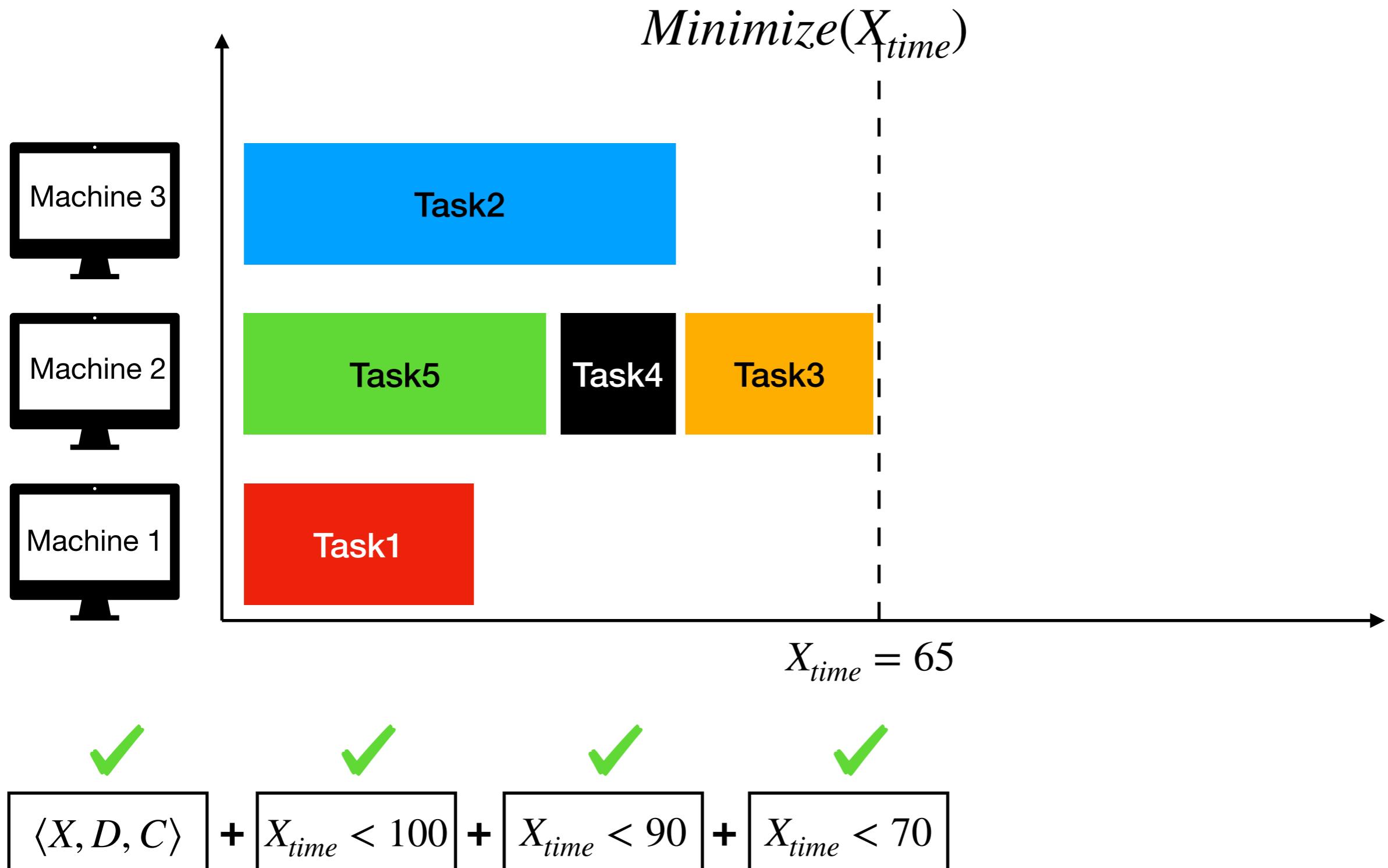
Optimization in CP (Banch&Bound)



$$\checkmark \quad \checkmark \quad \checkmark$$
$$\langle X, D, C \rangle + X_{time} < 100 + X_{time} < 90 + X_{time} < 70$$

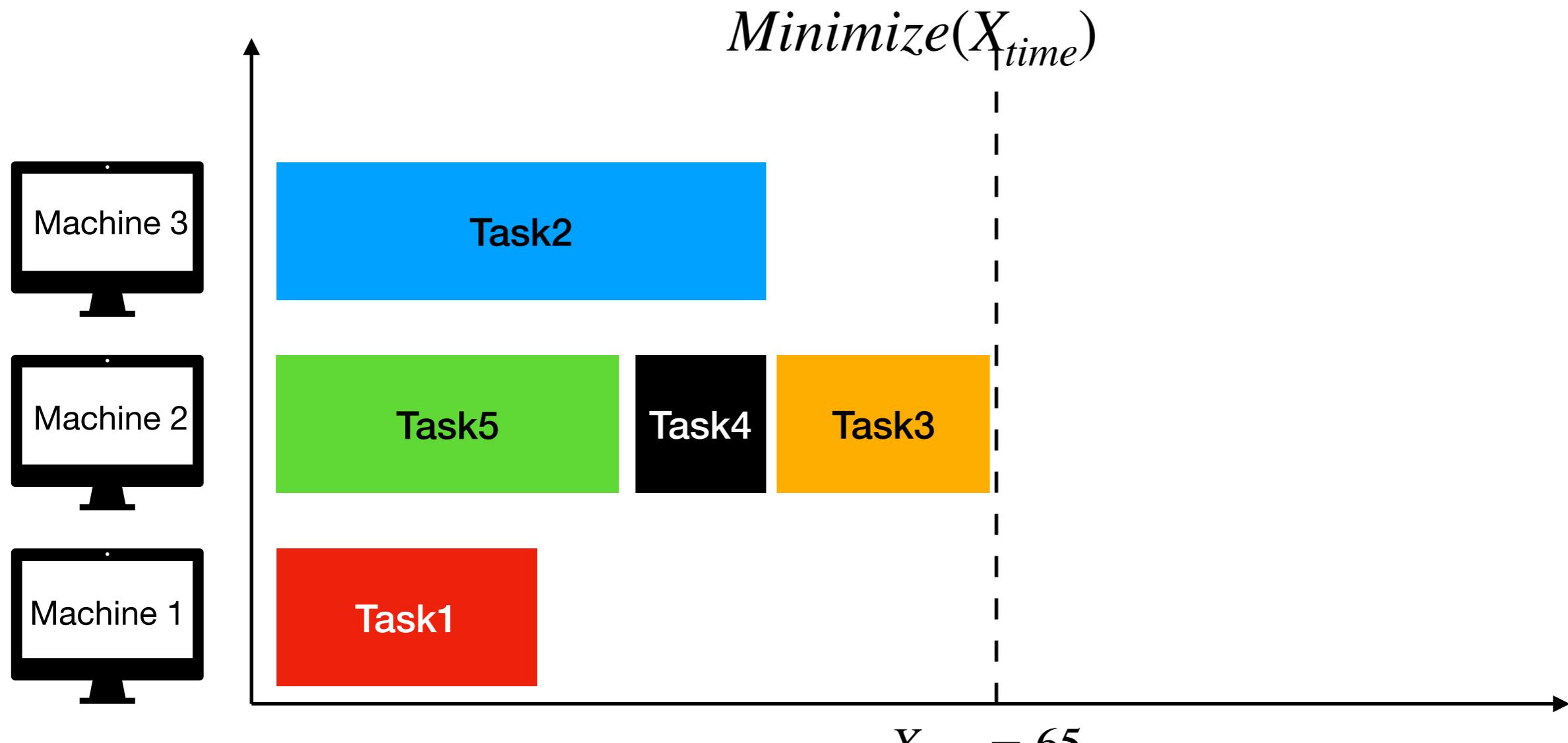
Constraint Programming

Optimization in CP (Banch&Bound)



Constraint Programming

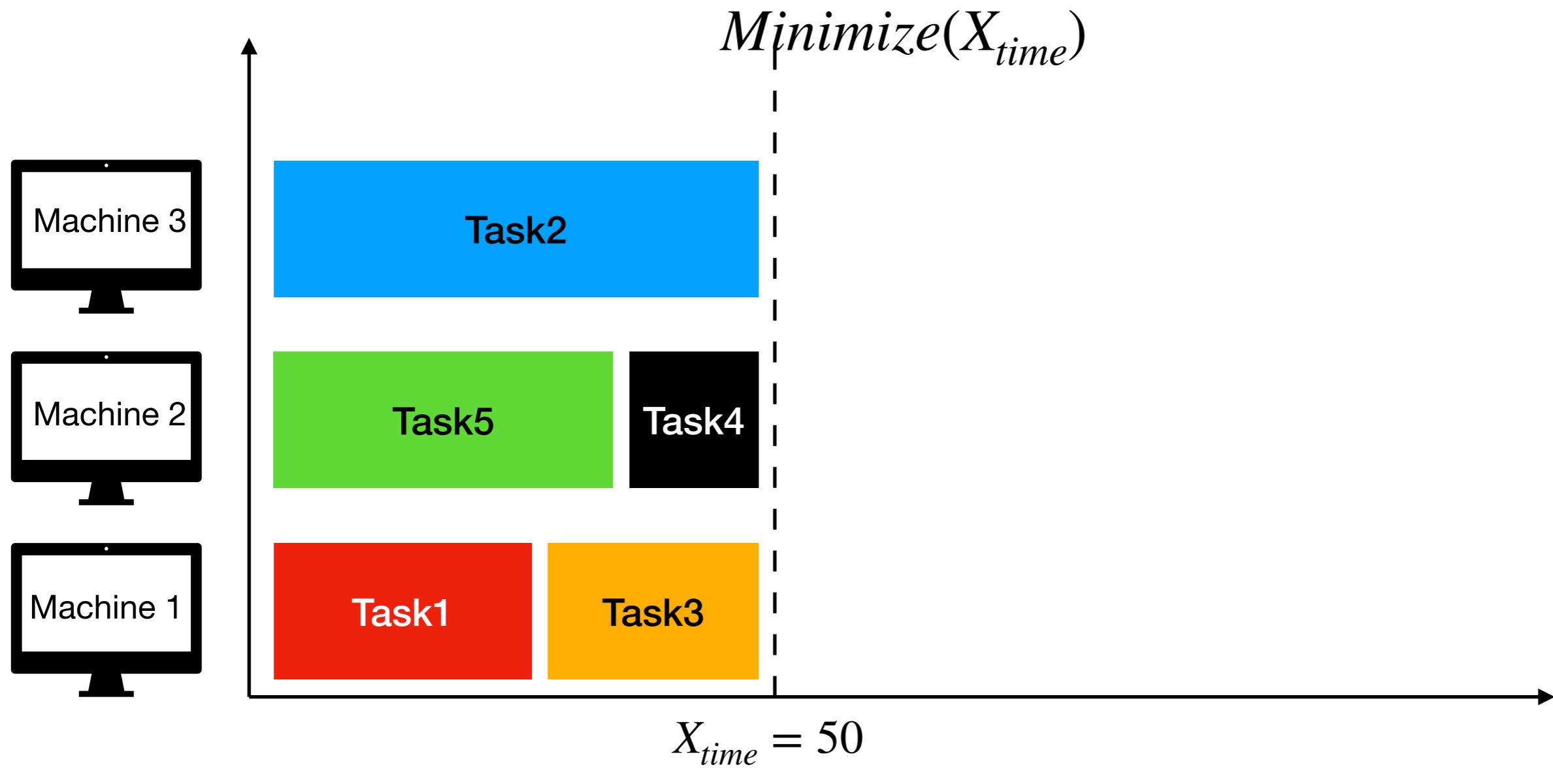
Optimization in CP (Banch&Bound)



$$\checkmark + \checkmark + \checkmark + \checkmark + \checkmark$$
$$\langle X, D, C \rangle + X_{time} < 100 + X_{time} < 90 + X_{time} < 70 + X_{time} < 65$$

Constraint Programming

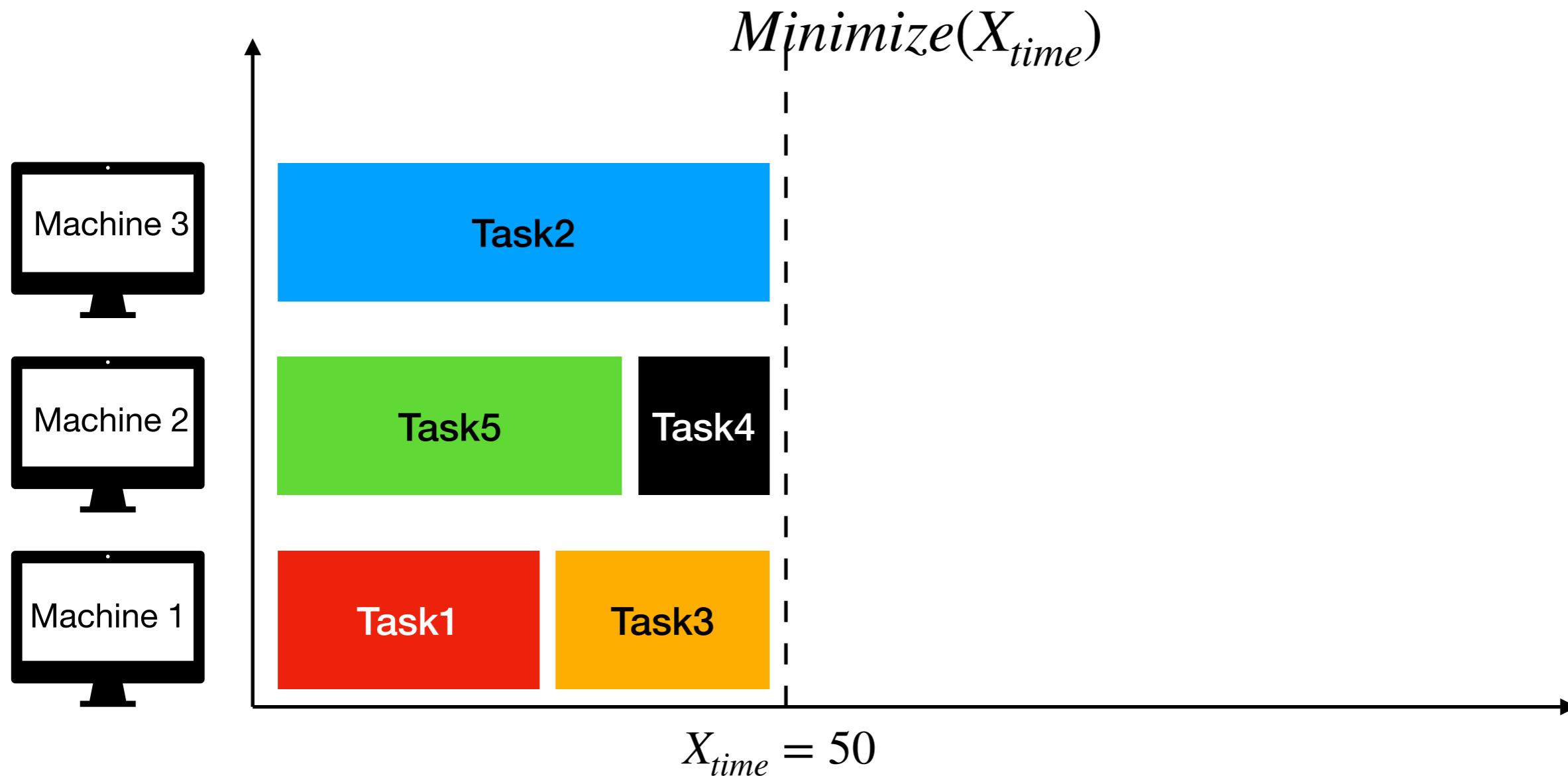
Optimization in CP (Banch&Bound)



$$\checkmark + \checkmark + \checkmark + \checkmark + \checkmark$$
$$\langle X, D, C \rangle + X_{time} < 100 + X_{time} < 90 + X_{time} < 70 + X_{time} < 65$$

Constraint Programming

Optimization in CP (Banch&Bound)

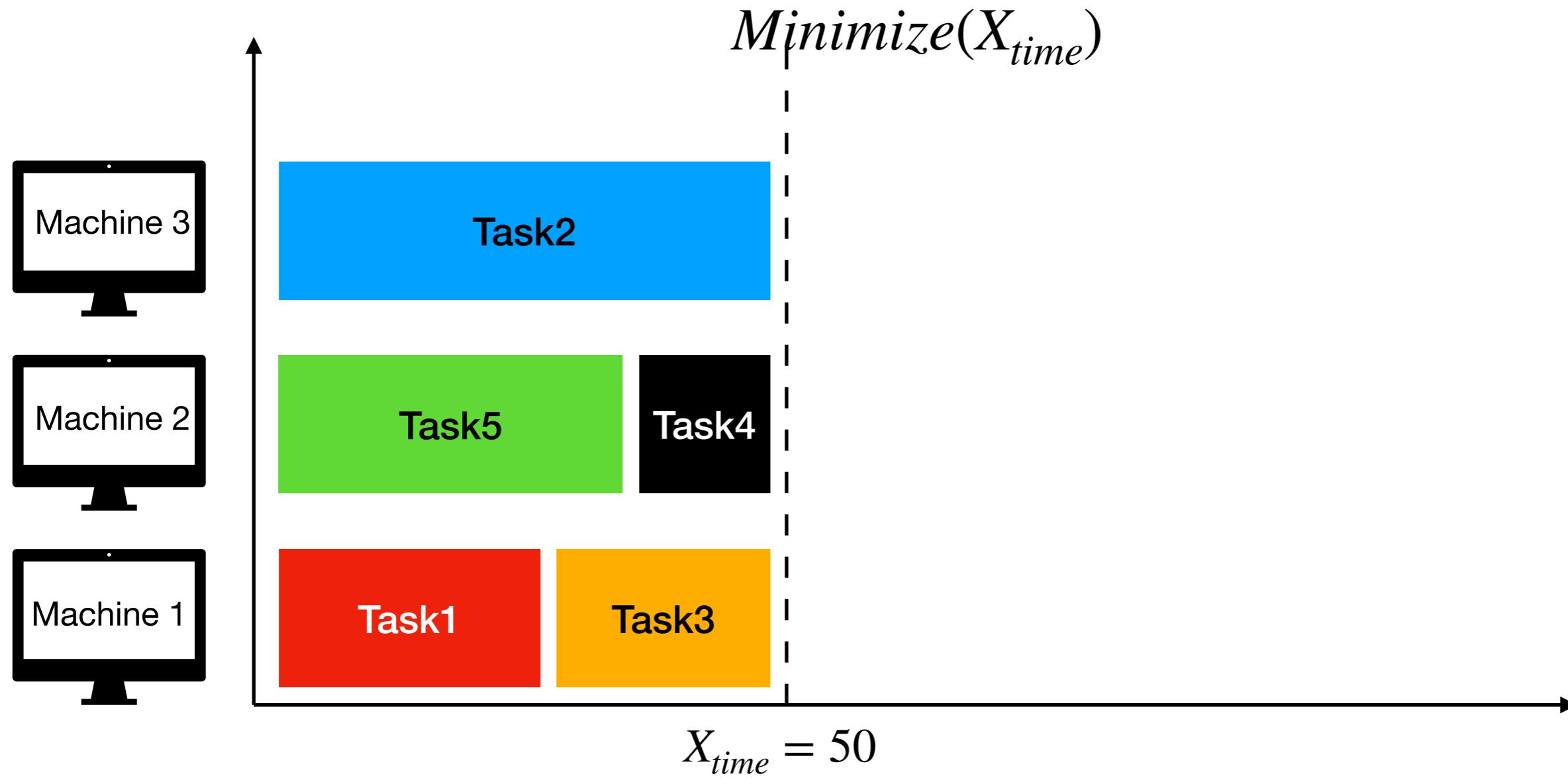


$$\checkmark + \langle X, D, C \rangle + X_{time} < 100 + X_{time} < 90 + X_{time} < 70 + X_{time} < 65 + X_{time} < 50$$

36

Constraint Programming

Optimization in CP (Branch&Bound)



$$\langle X, D, C \rangle + X_{time} < 100 + X_{time} < 90 + X_{time} < 70 + X_{time} < 65 + X_{time} < 50$$

36

Constraint Programming

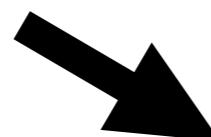
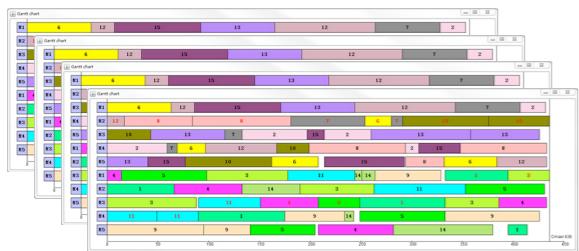
Global constraints



Constraint Programming

Global constraints

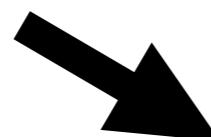
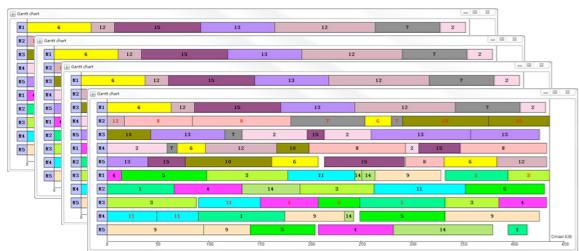
Scheduling instances



Constraint Programming

Global constraints

Scheduling instances

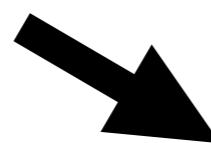
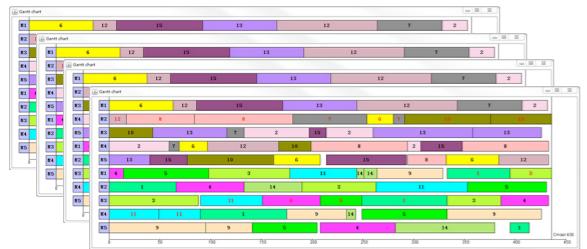


Cumulative

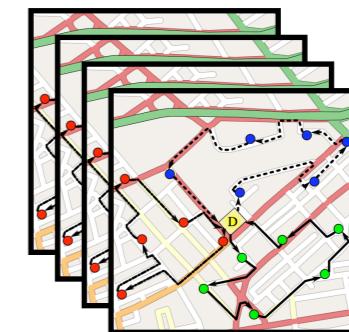
Constraint Programming

Global constraints

Scheduling instances



Vehicule routing instances

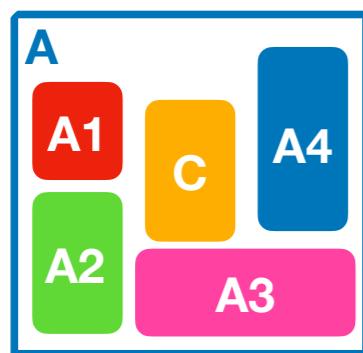
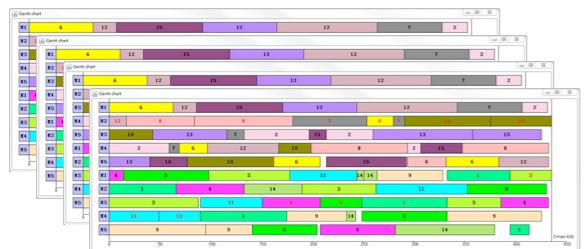


Cumulative

Constraint Programming

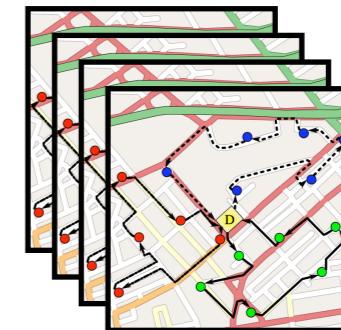
Global constraints

Scheduling instances



Cumulative

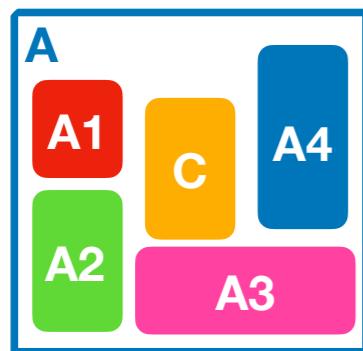
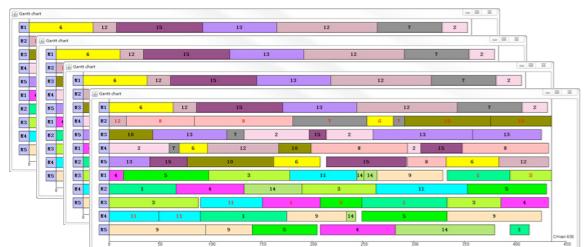
Vehicule routing instances



Constraint Programming

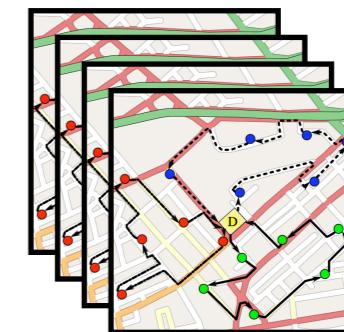
Global constraints

Scheduling instances



Cumulative

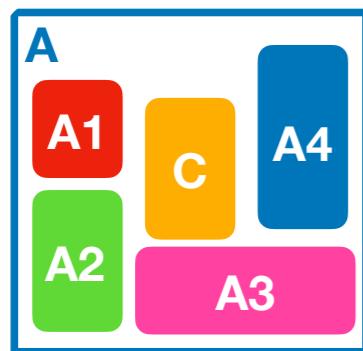
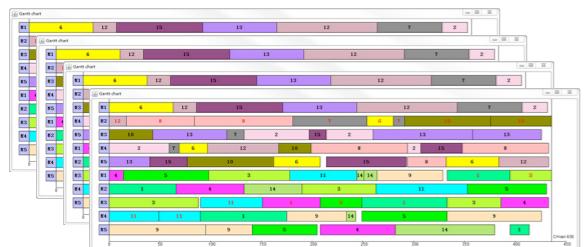
Vehicule routing instances



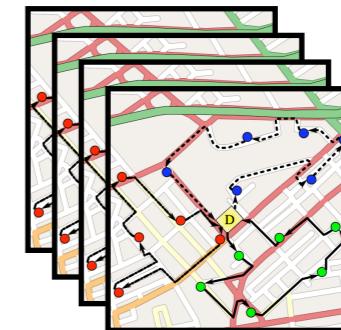
Constraint Programming

Global constraints

Scheduling instances



Vehicule routing instances

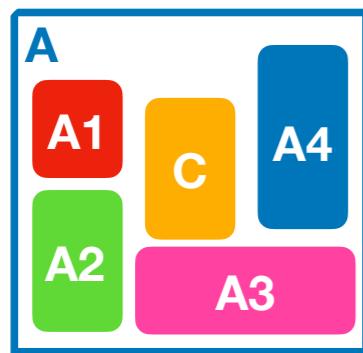
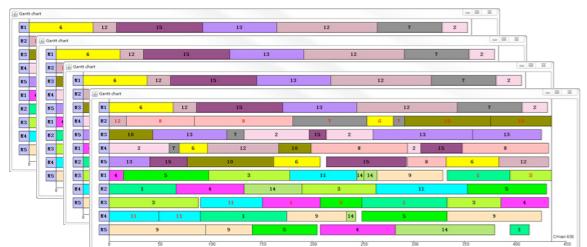


Cumulative
Alldifferent

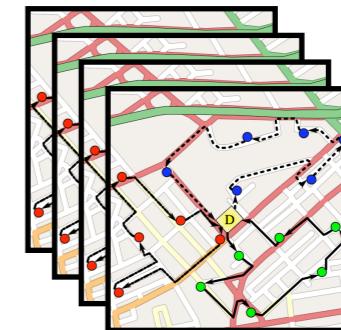
Constraint Programming

Global constraints

Scheduling instances



Vehicule routing instances



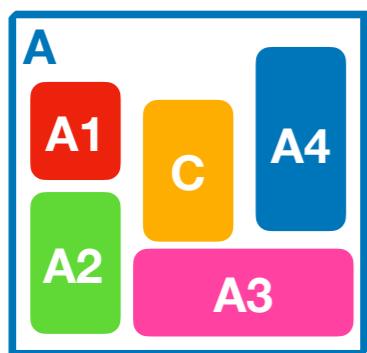
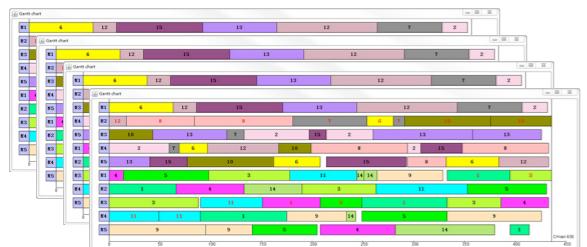
Cumulative
Alldifferent

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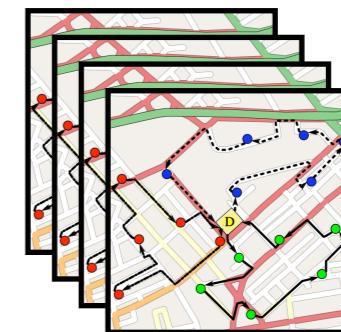
Constraint Programming

Global constraints

Scheduling instances



Vehicule routing instances

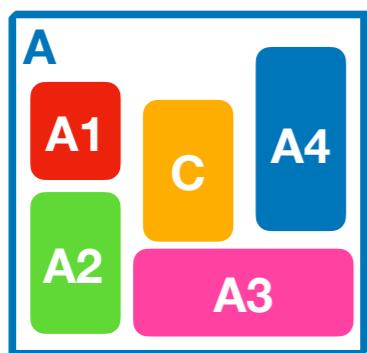
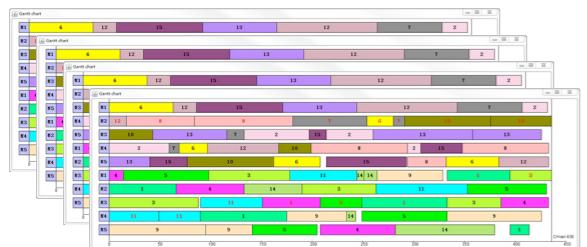


Cumulative
Alldifferent
Sum
knapsack
Element
GlobalCardinality
Regular
...

Constraint Programming

Global constraints

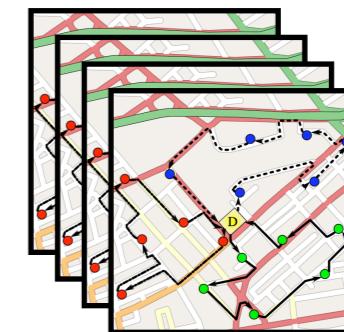
Scheduling instances



toolbox

Cumulative
AllDifferent
Sum
knapsack
Element
GlobalCardinality
Regular
...

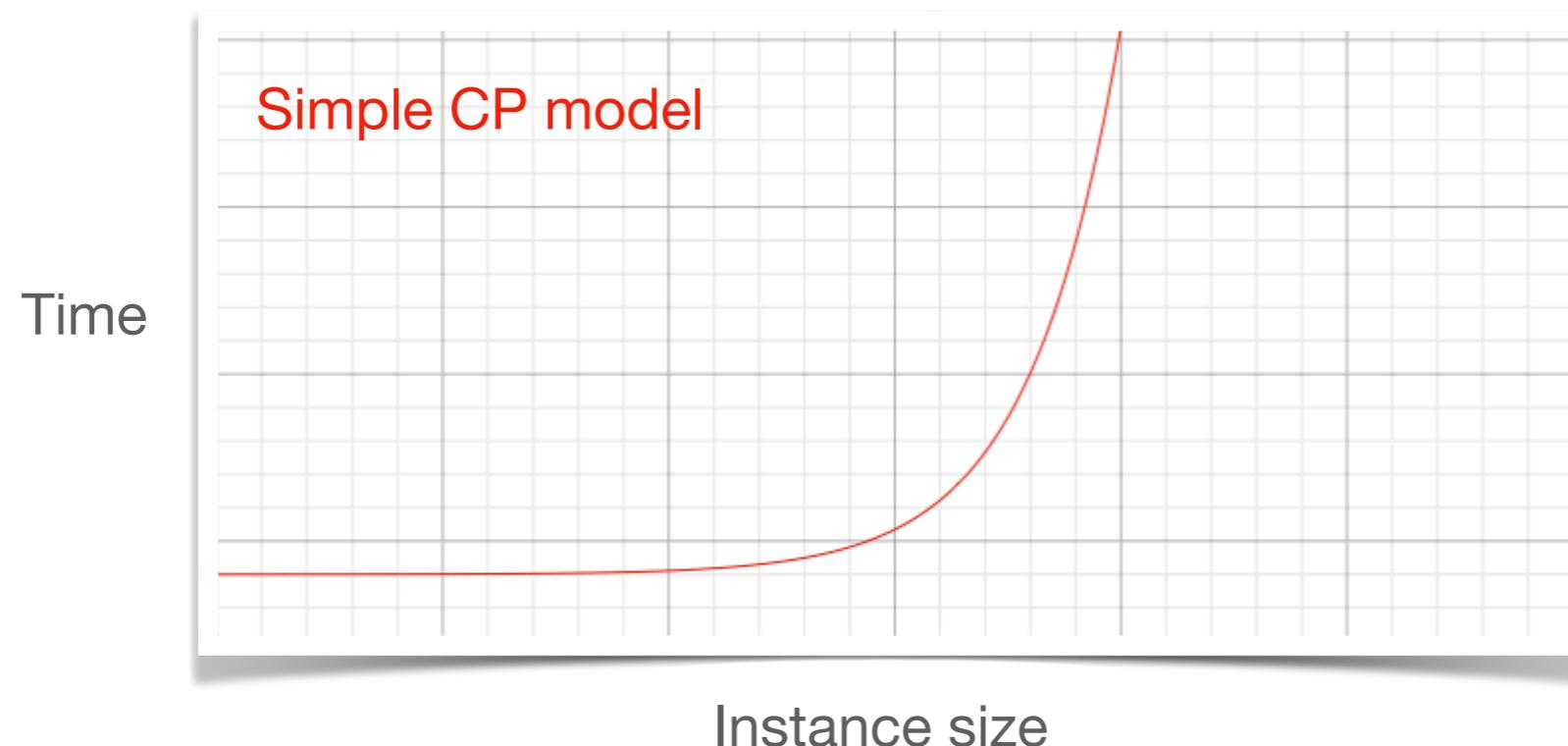
Vehicule routing instances



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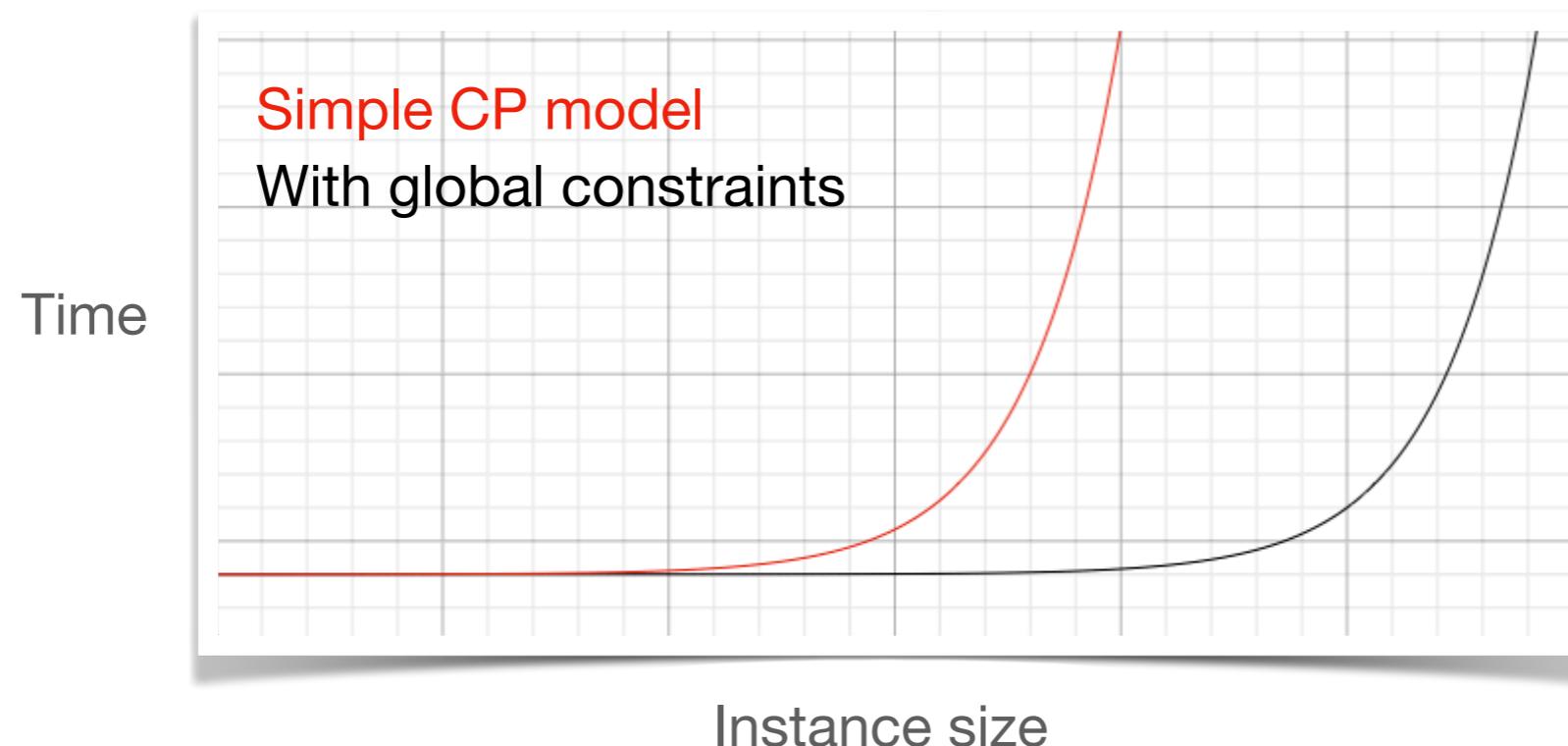
Constraint Programming

Global constraints



Constraint Programming

Global constraints



Global Constraints

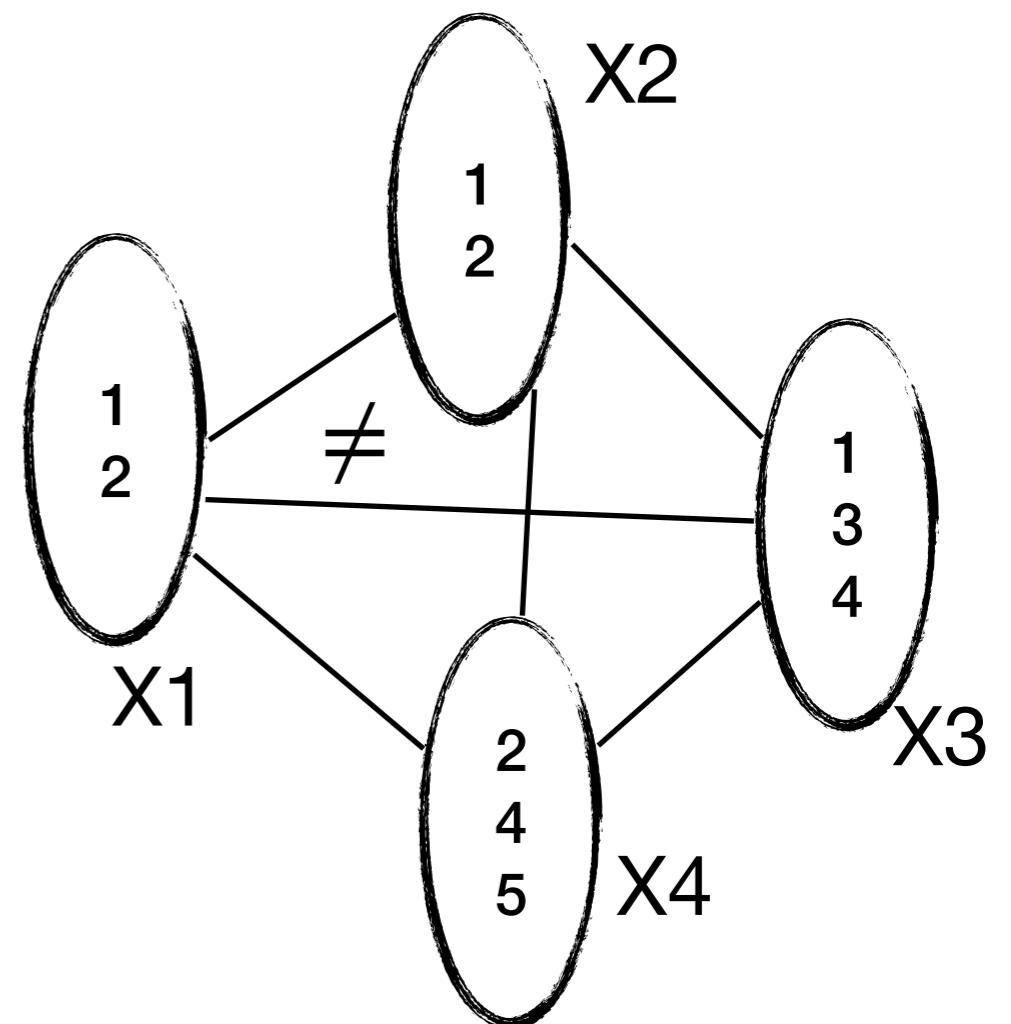
Example: Alldifferent

[Régin 1994]

Global Constraints

Example: Alldifferent

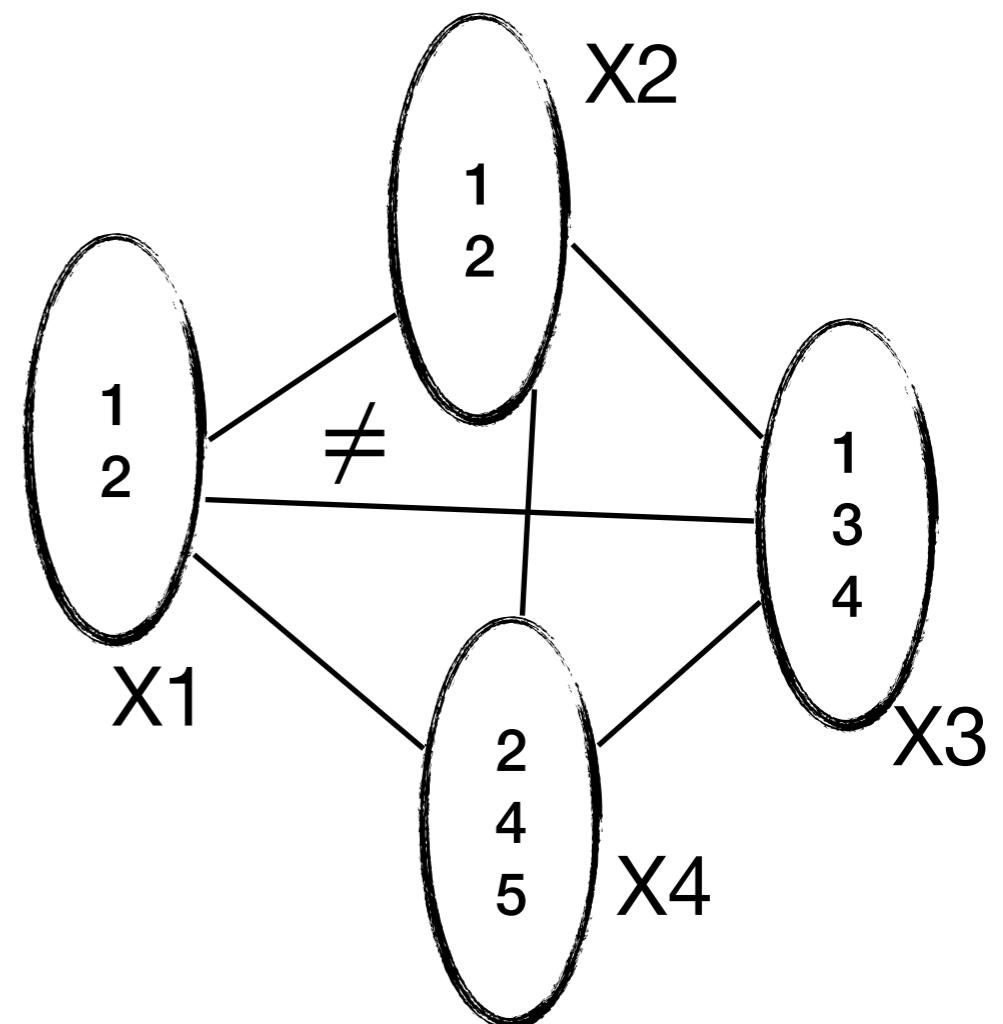
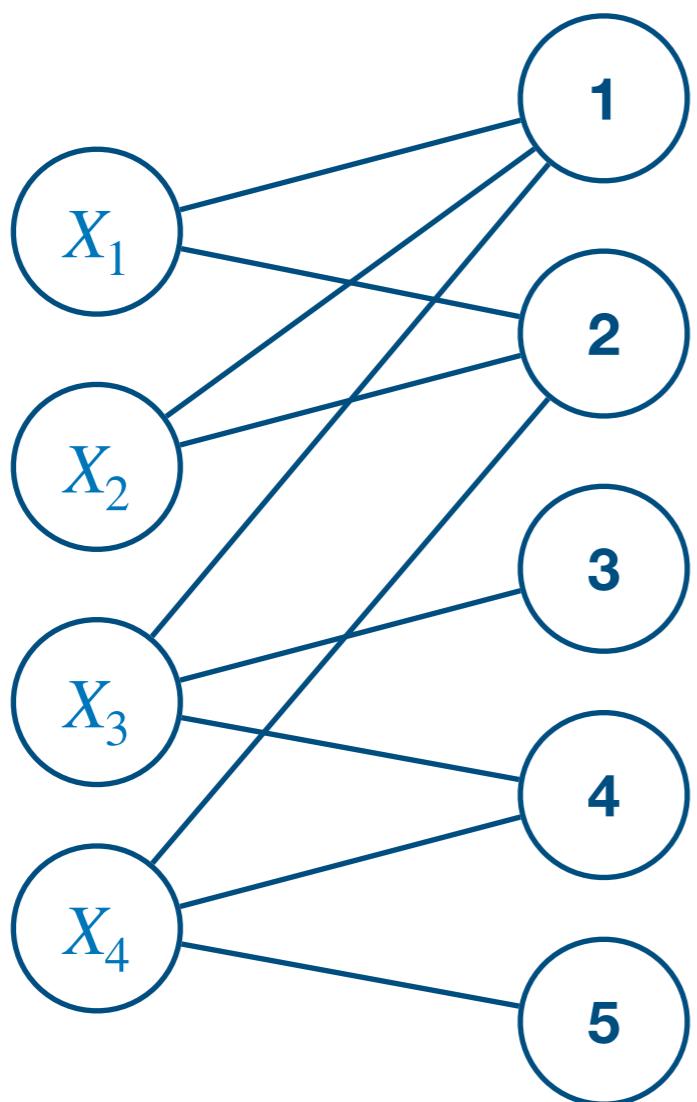
[Régin 1994]



Global Constraints

Example: Alldifferent

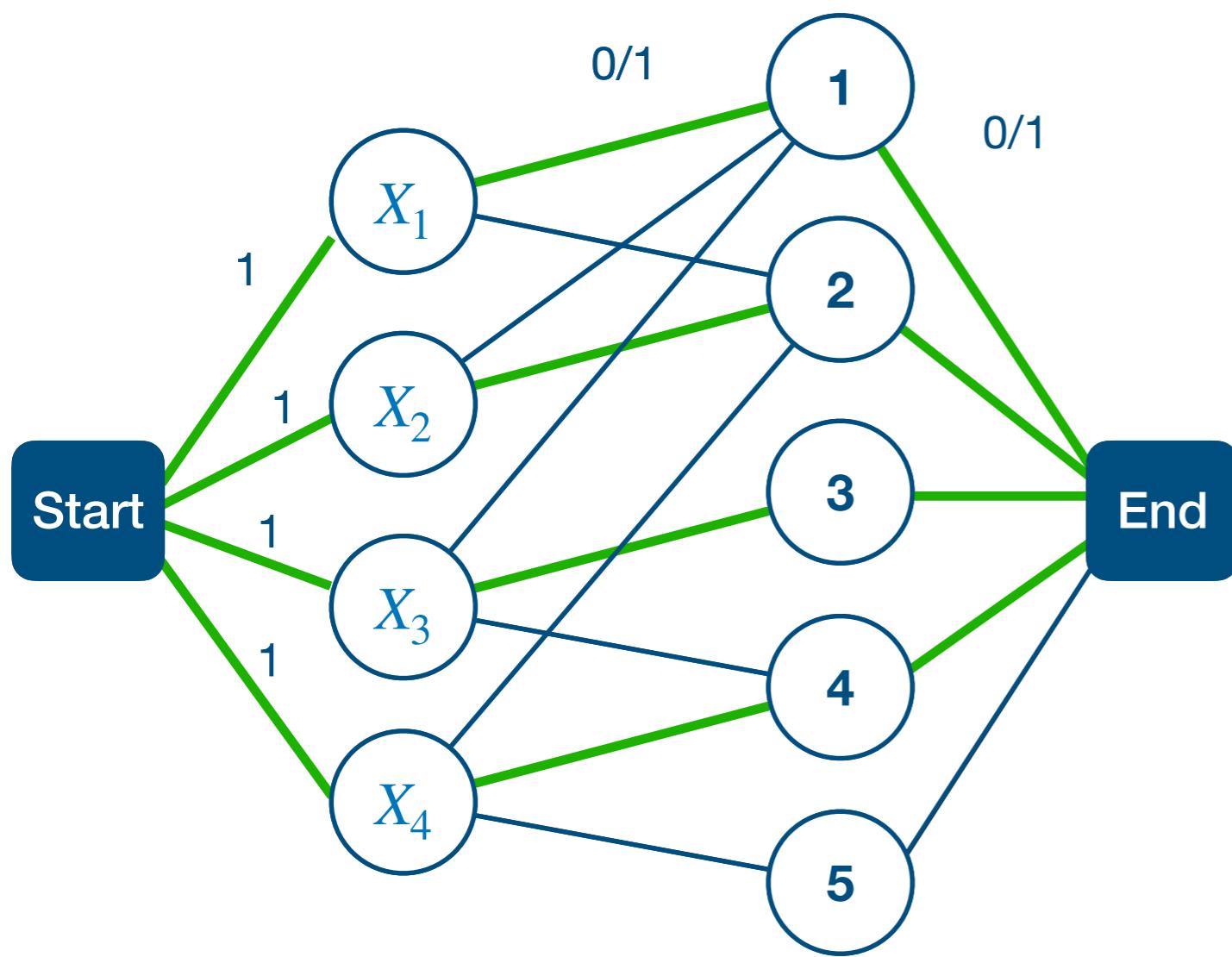
[Régin 1994]



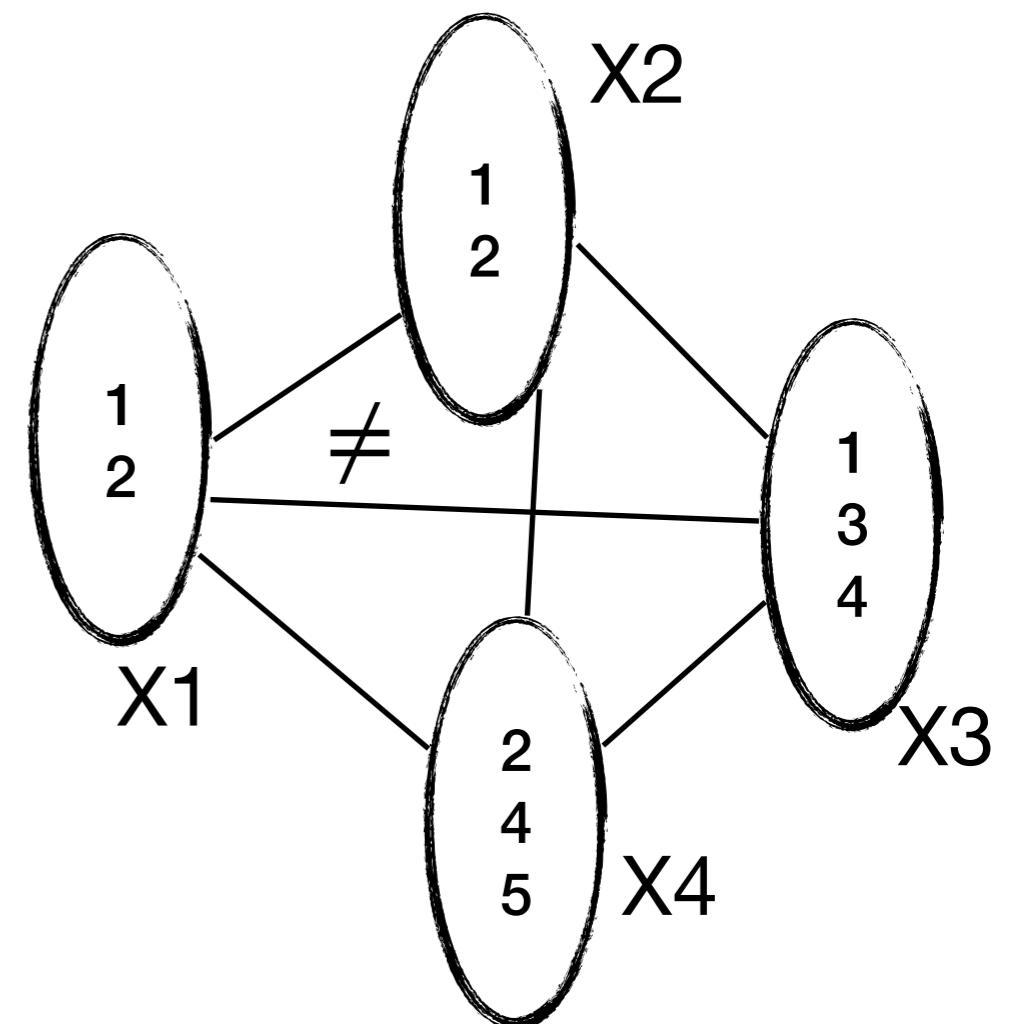
Global Constraints

Example: AllDifferent

[Régin 1994]



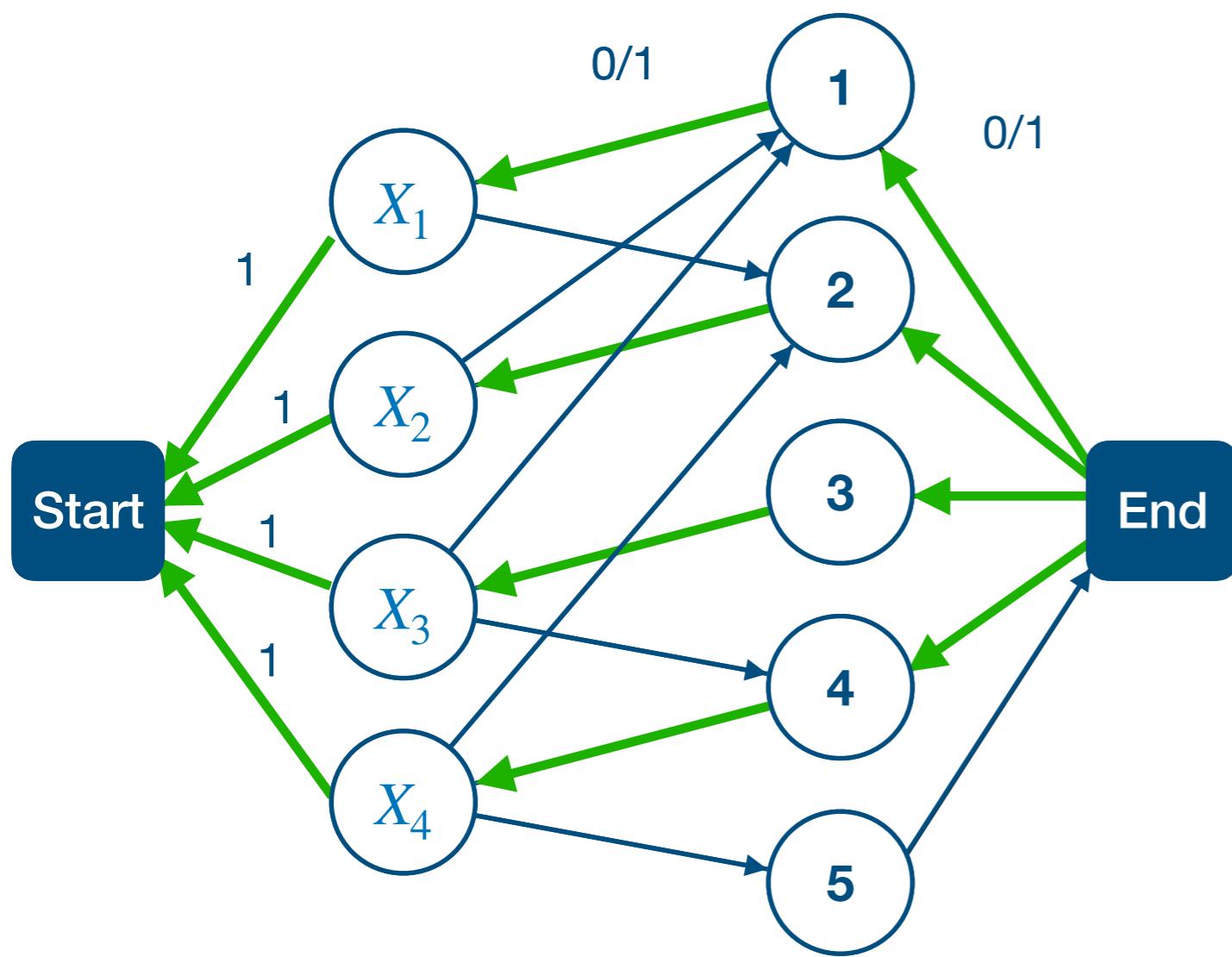
Max Flow



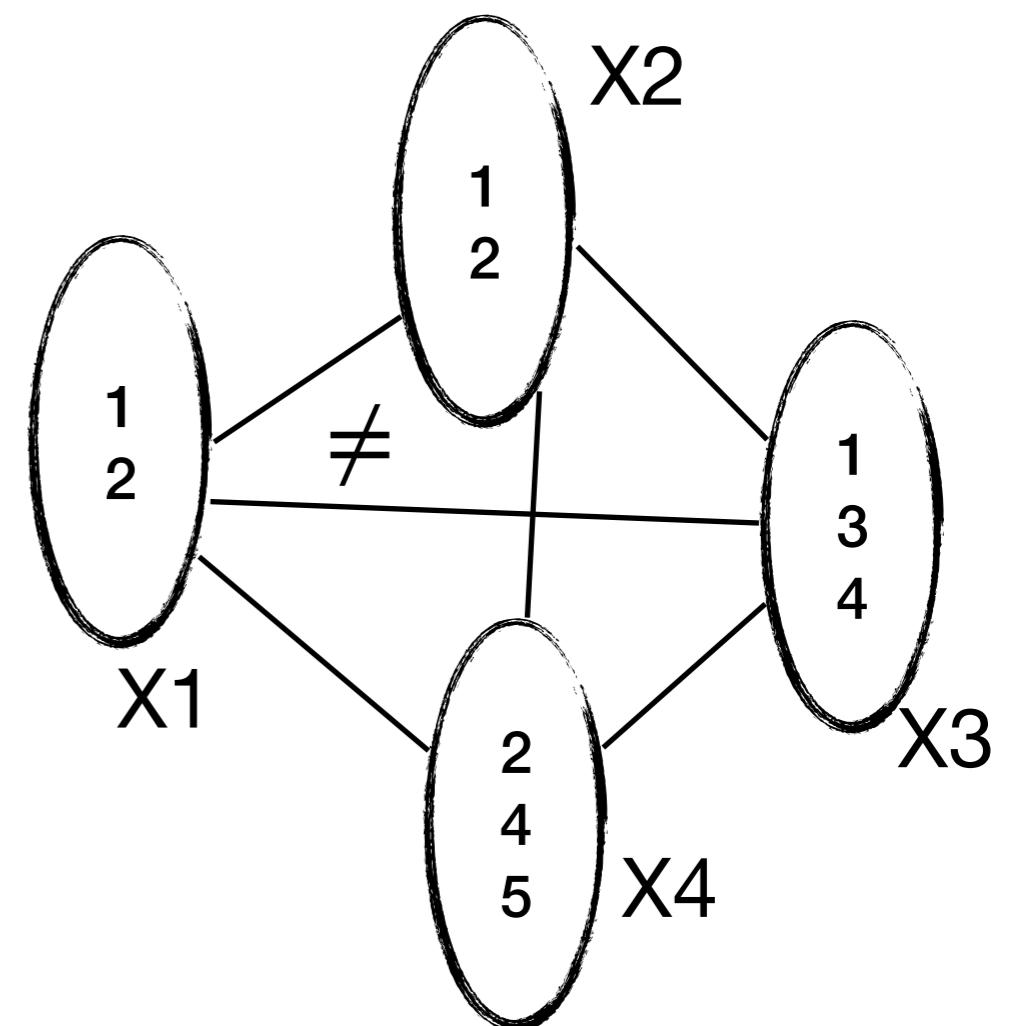
Global Constraints

Example: Alldifferent

[Régin 1994]



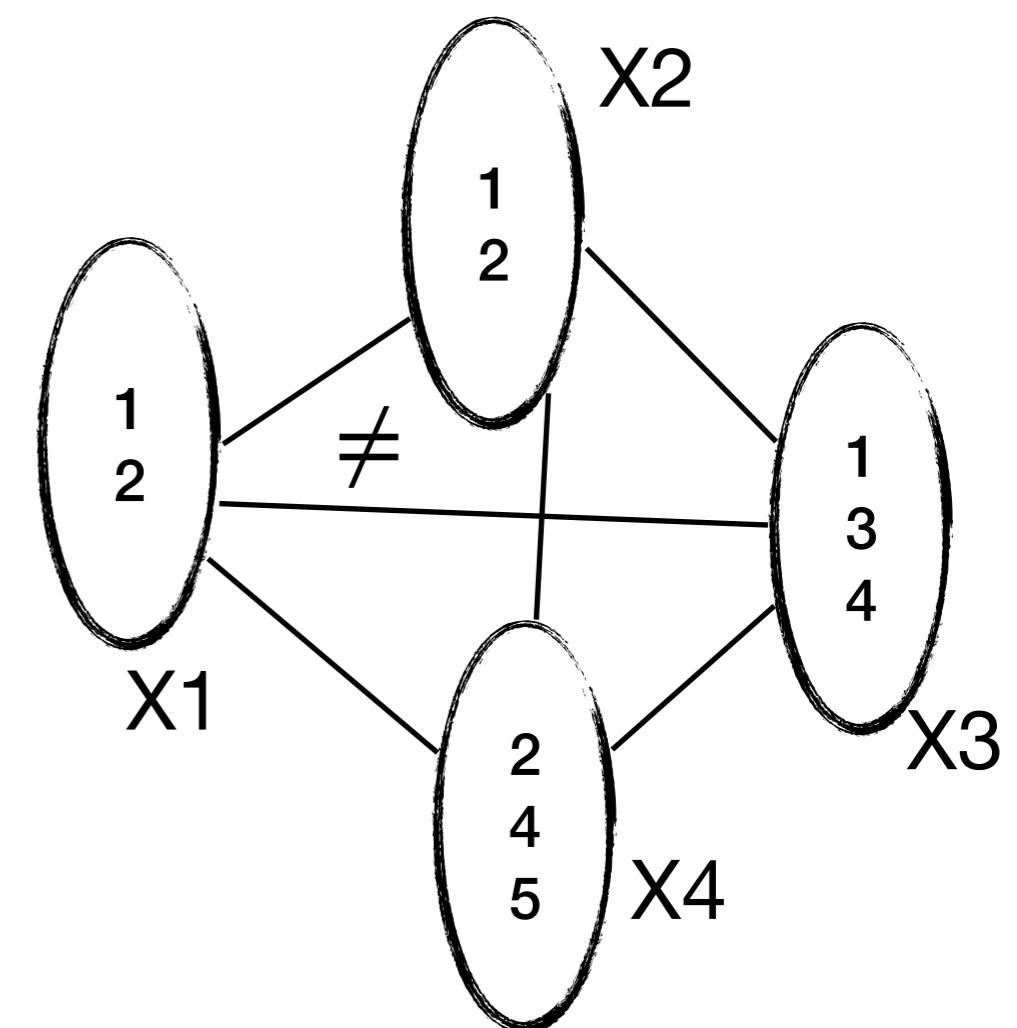
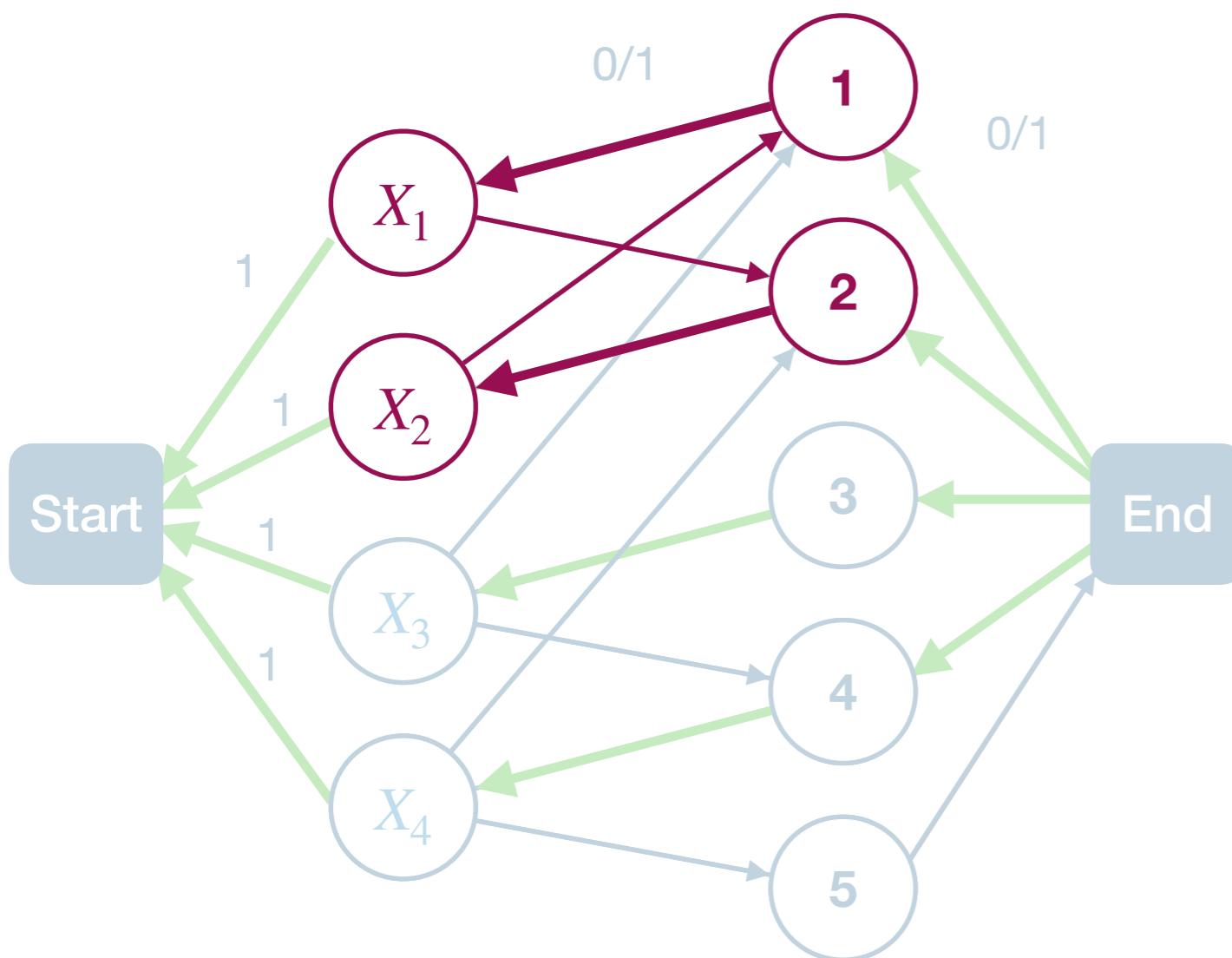
Residual Graph



Global Constraints

Example: Alldifferent

[Régin 1994]

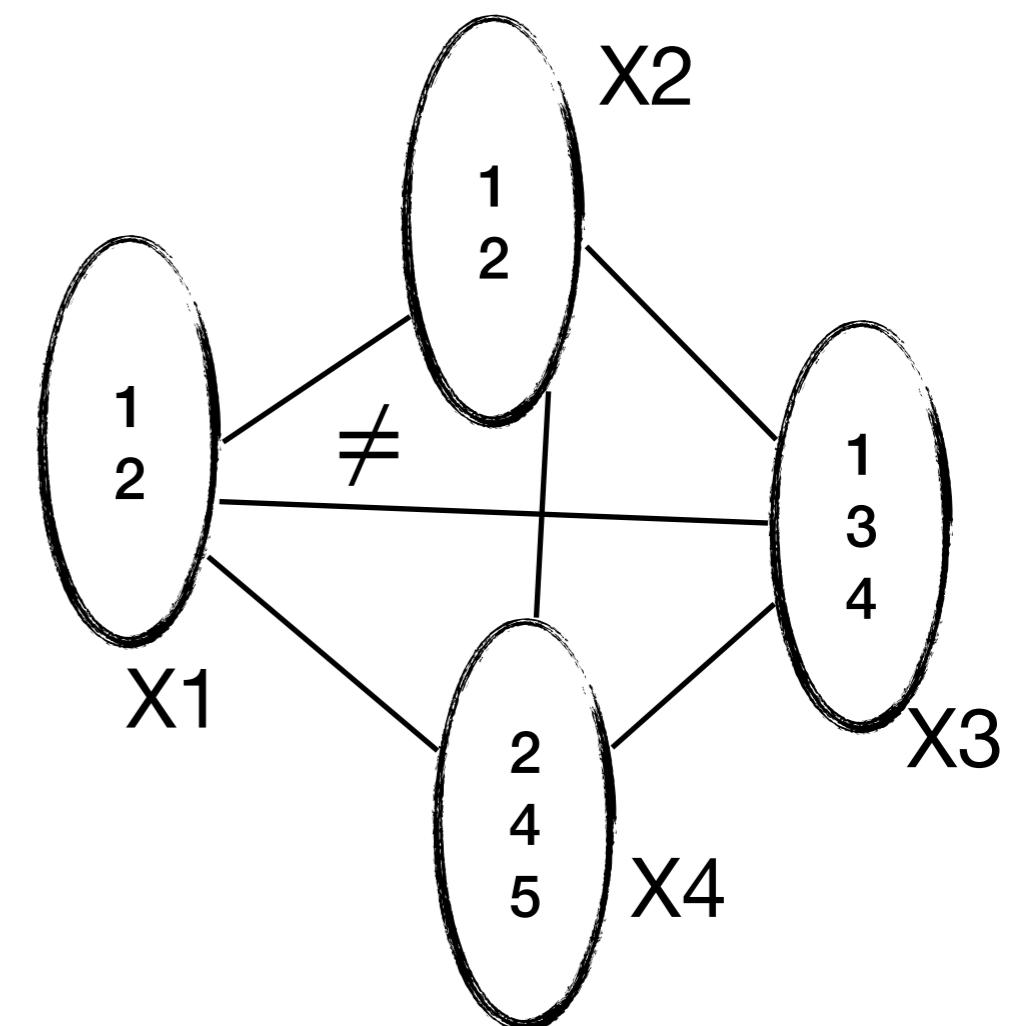
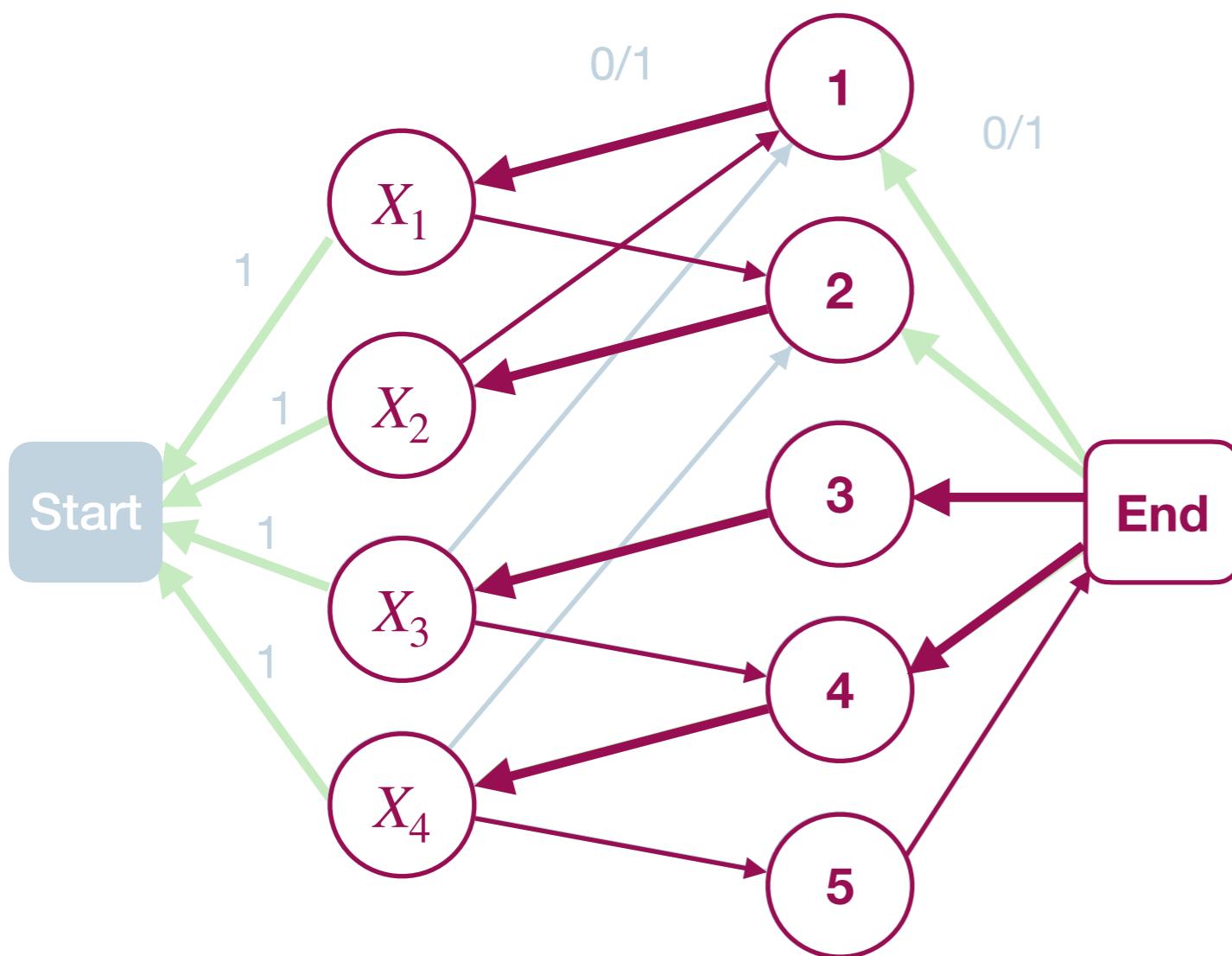


Strongly connected components

Global Constraints

Example: Alldifferent

[Régin 1994]

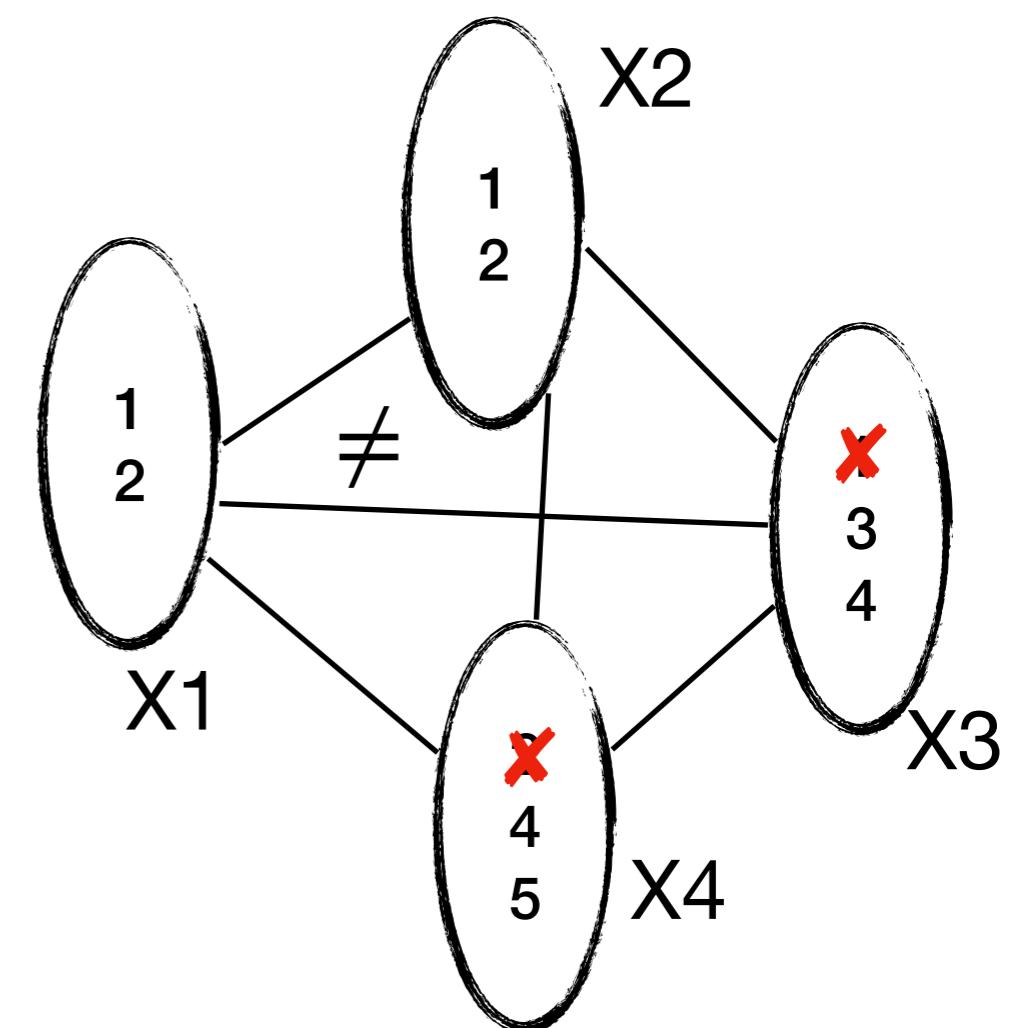
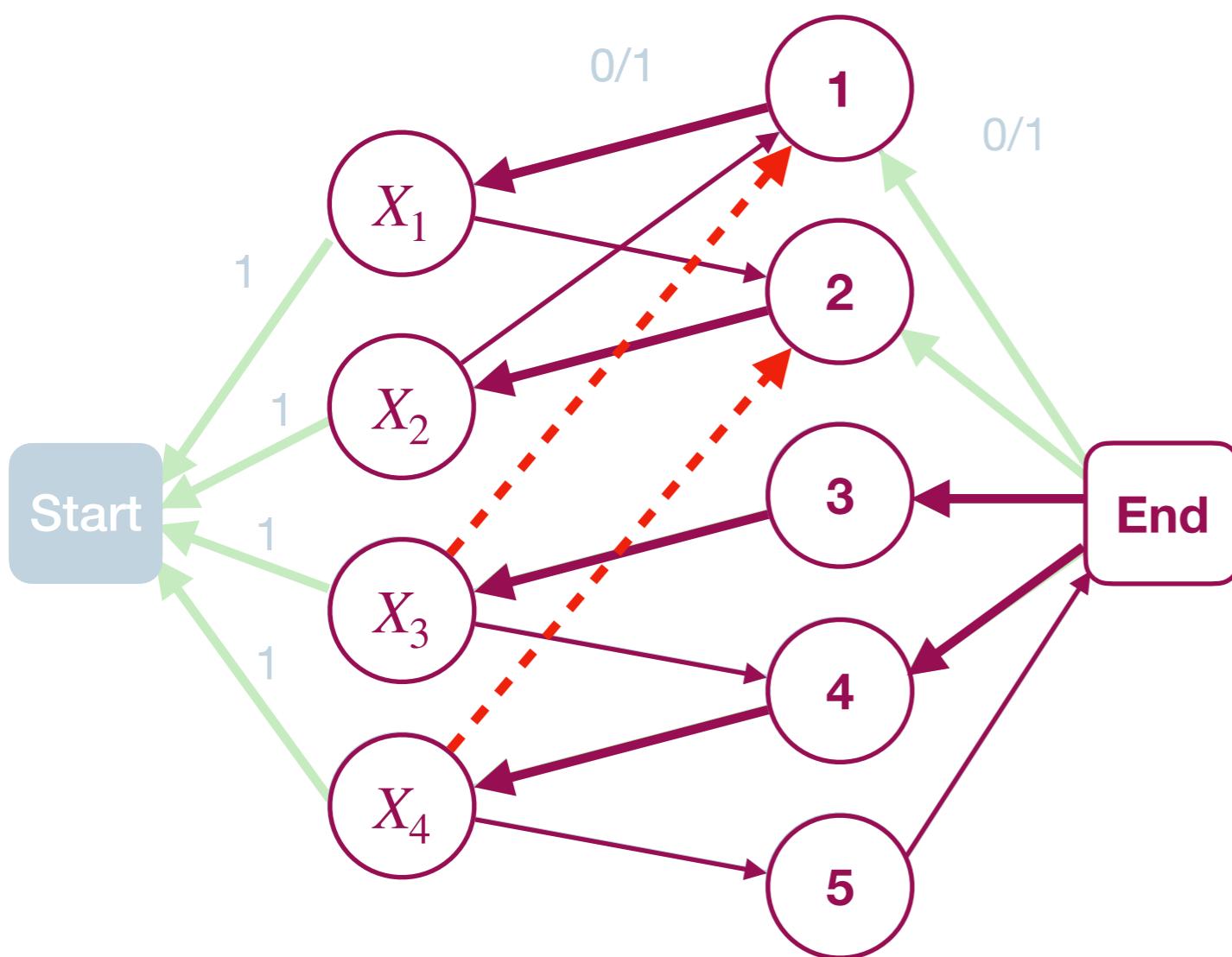


Strongly connected components

Global Constraints

Example: Alldifferent

[Régin 1994]

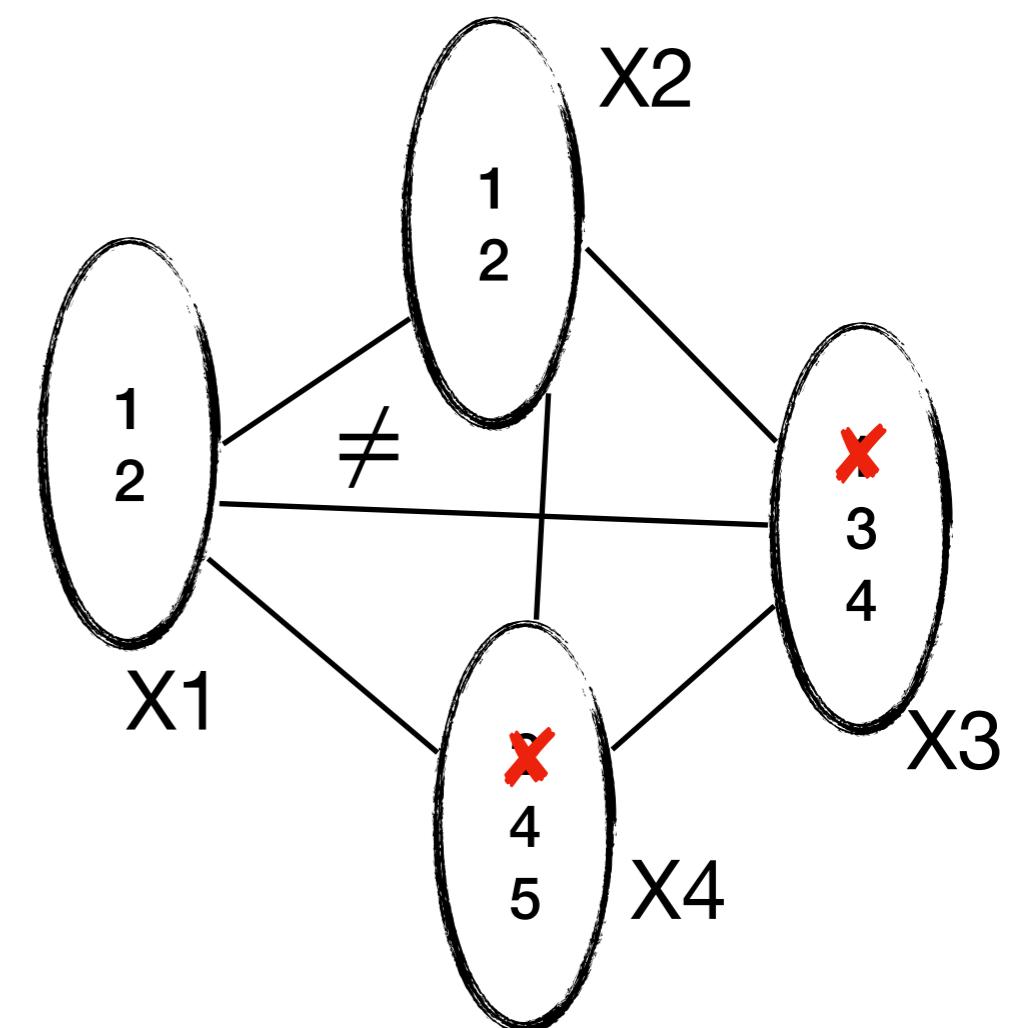
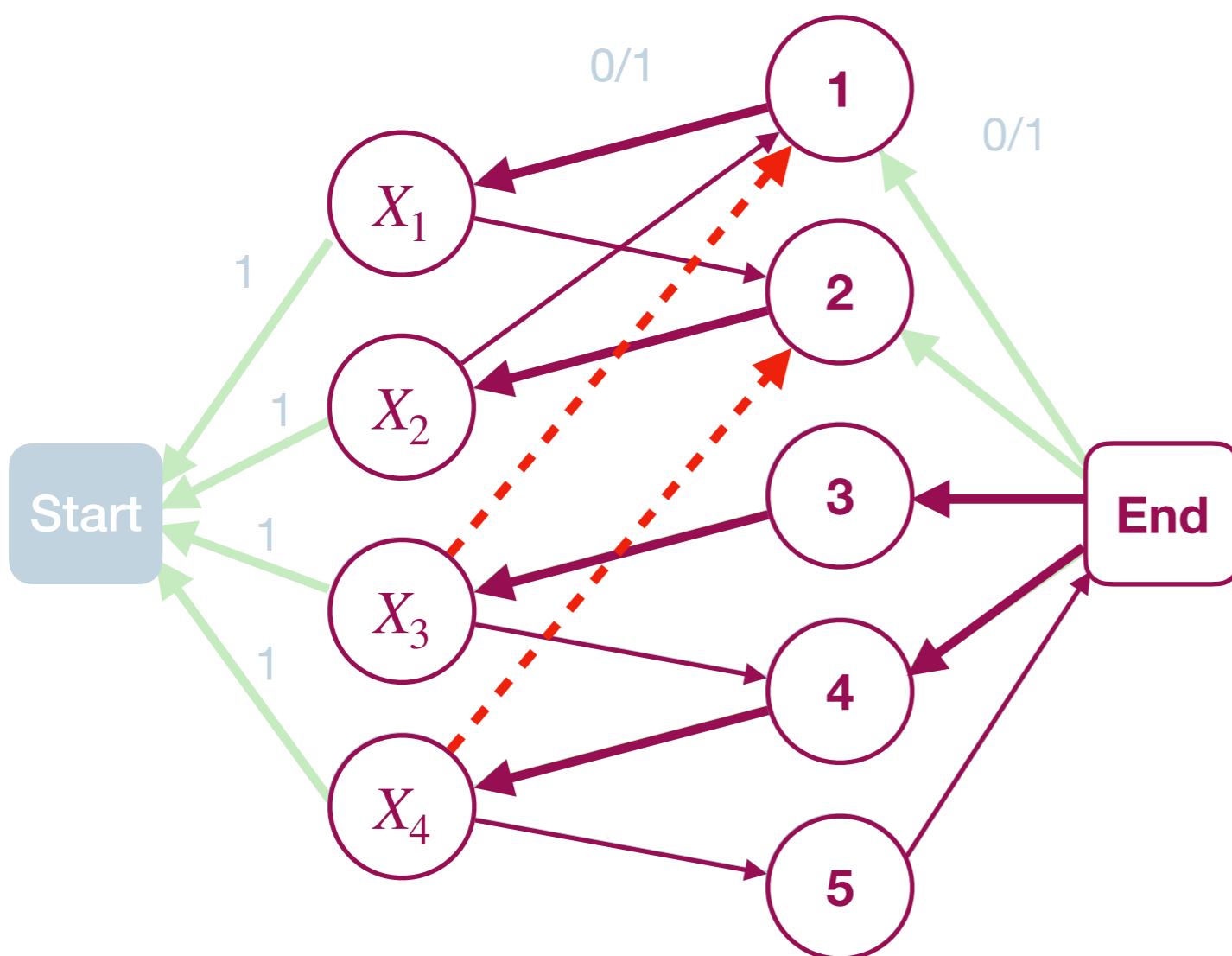


Strongly connected components

Global Constraints

Example: Alldifferent

[Régin 1994]



Strongly connected components

Constraint Programming

References & links

- **Handbook of Constraint Programming.** Francesca Rossi Peter van Beek Toby Walsh. Elsevier Science 2006



- ACP: [Association for Constraint Programming](#) 
- [Guide to Constraint Programming](#)
- [CP conference Series](#)
- [Global Constraint Catalog](#)