

# universite Paris-saclay



# **Constraint & Data Mining**

Cours2-3

#### Master 2 - DS

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#### **Definition and steps**

 Knowledge Discovery in Databases (KDD) revolves around the investigation and creation of knowledge, processes, algorithms, and the mechanisms for retrieving potential knowledge from data collections

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Preparation Steps

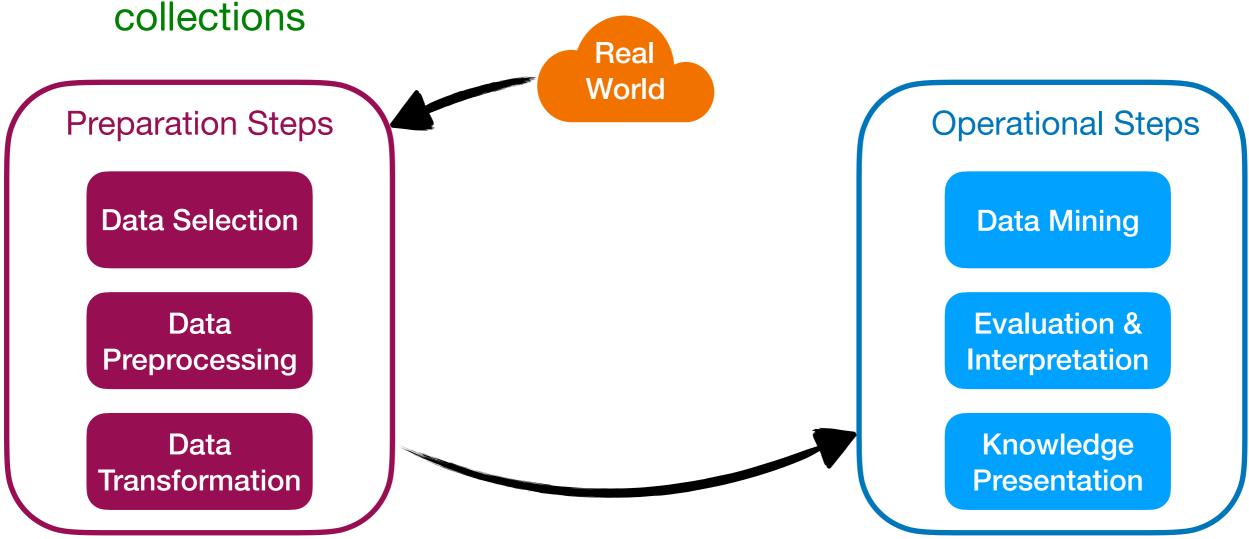
Data Selection

Data Preprocessing

Data Transformation

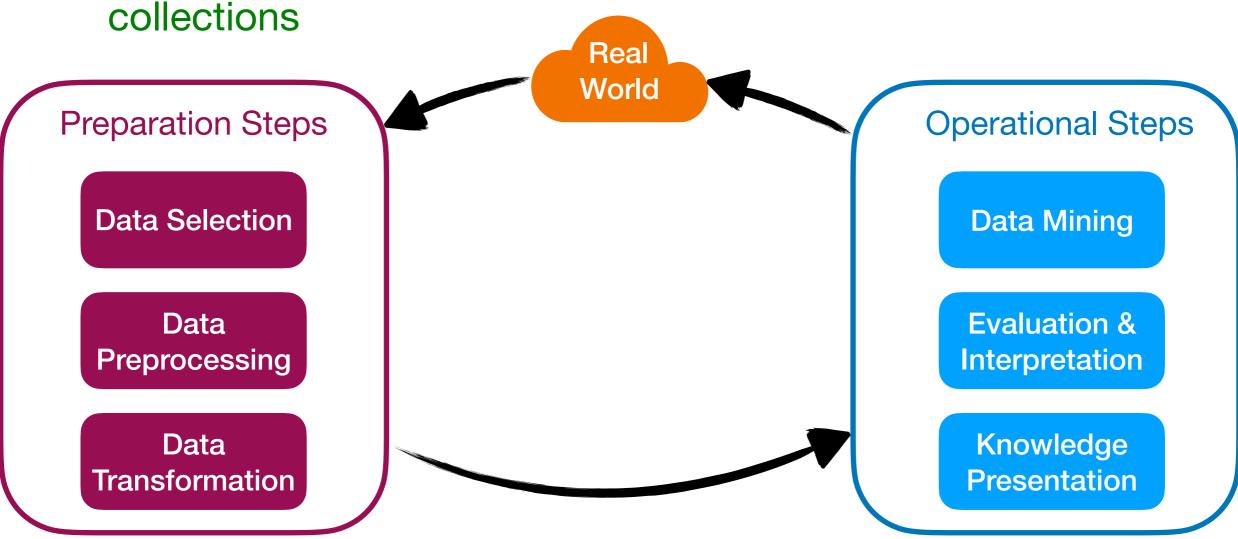
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#### **Motivations**

- Purpose: A key method in market basket analysis, aimed at identifying patterns in customer purchasing behavior
- Goal: To discover frequently co-occurring products, which are items that customers often buy together
- Applications of Frequent Itemsets:
  - Product Arrangement: Optimize product placement on store shelves or within catalogs
  - Cross-Selling: Suggest related products to customers (e.g., recommending additional products during online shopping)
  - Product Bundling: Offer product bundles based on co-purchase patterns
  - Other Applications: Fraud detection, dependency analysis, fault localization, and more

#### **Basic Notions**

• Items: 
$$I = \{p_1, ..., p_n\}$$

• Itemset, transaction: 
$$P, T \subseteq I$$

• Transactional dataset: 
$$D = \{T_1, ..., T_m\}$$

• Language of itemsets: 
$$\mathcal{L}_I = 2^I$$

• Cover of an itemset: 
$$cover(P) = \{T_i \in D : P \subseteq T_i\}$$

• Absolute Frequency: 
$$freq(P) = |cover(P)|$$

. Relative Frequency: 
$$freq(P) = \frac{|\mathit{cover}(P)|}{|D|}$$

#### **Problem Definition**

#### Given:

- A set of items  $I = \{p_1, ..., p_n\}$
- A transactional dataset  $D = \{T_1, ..., T_m\}$
- A minimum support  $\, lpha \,$

#### The need:

• The set of itemset P s.t.:  $freq(P) \ge \alpha$ 

#### Example(1)

- $I = \{a, b, c, d, e\}$
- $D = \{T_1, ..., T_{10}\}$

#### Example(1)

• 
$$I = \{a, b, c, d, e\}$$

• 
$$D = \{T_1, ..., T_{10}\}$$

#### $H_D$

а			d	е
	b	С	d	
а		С		е
a		С	d	е
a				е
а		С	d	
	b	С		
a		С	d	е
	b	С		е
а			d	E
	a a a a	b a a a b a b	b c a c a c a c a c a c b c a c	b c d a c a c d a a c d b c b c a c d b c a c d

#### Example(1)

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$$I = \{a, b, c, d, e\}$$

• 
$$D = \{T_1, ..., T_{10}\}$$

#### $H_D$

1:	a			d	е
2:		b	С	d	
3:	a		С		е
4:	a		С	d	е
5:	a				е
6:	a		С	d	
7:		b	С		
8:	a		С	d	е
9:		b	С		е
10:					

#### $V_D$

а	b	С	d	е
1	2	2	1	1
1 3 4 5 6 8	7	3	2	3
4	9	4	2 4 6	4
5		6	6	4 5 8
6		7	8	
8		8	10	9
10		9		10

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$$I = \{a, b, c, d, e\}$$

• 
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#### $H_D$

1:	a			d	е
2:		b	С	d	
3:	a		С		е
4:	a		С	d	е
5:	a				е
6:	a		С	d	
7:		b	С		
8:	a		С	d	е
9:		b	С		е
10:	а			d	E

#### $V_D$

а	b	С	d	е
1	2	2	1	1
3	7	3	2	3
4	9	4	4	4
3 4 5 6 8		6	1 2 4 6 8	5 8
6		7	8	8
8		8	10	9
10		9		10

#### $M_D$

	а	b	С	d	е
1:	1	0	0	1	1
2:	0	1	1	1	0
3:	1	0	1	0	1
4:	1	0	1	1	1
5:	1	0	0	0	1
6:	1	0	1	1	0
7:	0	1	1	0	0
8:	1	0	1	1	1
9:	0	1	1	0	1
10:	1	0	0	1	1

## Example(1)

• 
$$I = \{a, b, c, d, e\}$$

• 
$$D = \{T_1, ..., T_{10}\}$$

#### $H_D$

1:	a			d	е
2:		b	С	d	
3:	a		С		е
4:	a		С	d	е
5:	a				е
6:	a		С	d	
7:		b	С		
8:	a		С	d	е
9:		b	С		е
10:	а			d	E

cover(bc) = ?

freq(bc) = ?

#### $V_D$

а	b	С	d	е
1	2	2	1	1
3	7	2 3	2	3
4	9	4 6	4	4 5
4 5 6 8			6	5
6		7	8	8
8		8	10	9
10		9		10

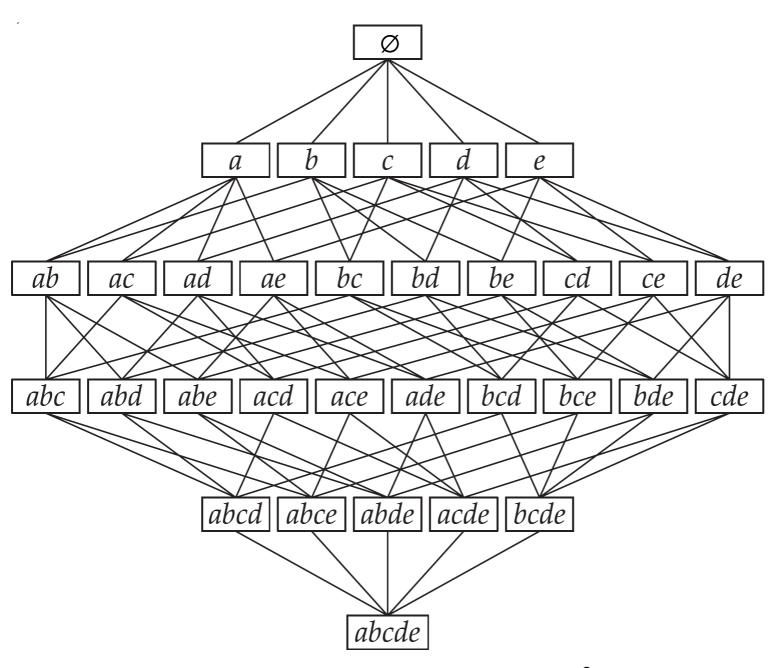
#### $M_D$

	а	b	С	d	е
1:	1	0	0	1	1
2:	0	1	1	1	0
3:	1	0	1	0	1
4:	1	0	1	1	1
5:	1	0	0	0	1
6:	1	0	1	1	0
7:	0	1	1	0	0
8:	1	0	1	1	1
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10:	1	0	0	1	1

Example(1)

$H_D$							
1:	а			d	е		
2:		b	С	d			
3:	a		С		е		
4:	a		С	d	е		
5:	a				е		
6:	a		С	d			
7:		b	С				
8:	а		С	d	е		
9:		b	С		е		
10:	а			d	E		

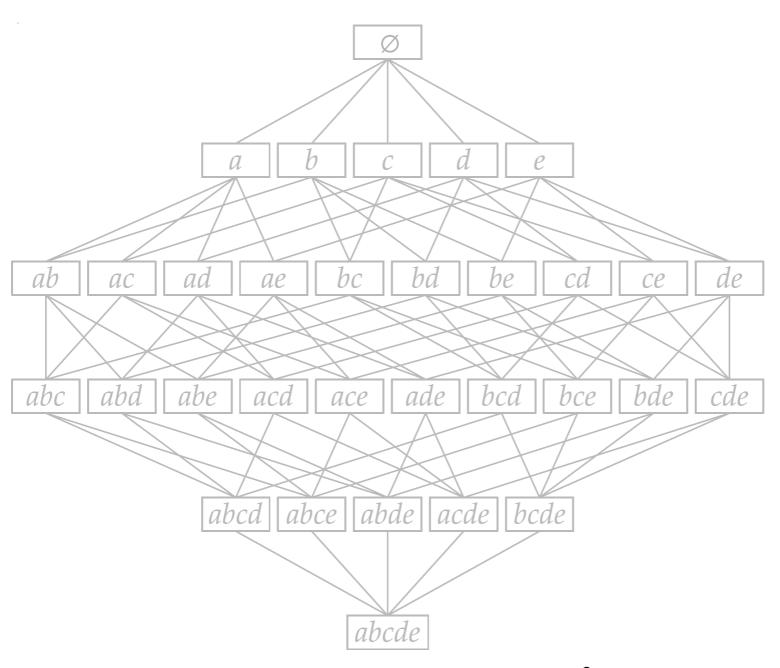
## Example(1)



$H_D$							
a			d	е			
	b	С	d				
a		С		е			
a		С	d	е			
a				е			
a		С	d				
	b	С					
a		С	d	е			
	b	С		е			
а			d	Е			
	a a a	a b a b a b	a b c a c a c a b c a c a b c a	a d b c d a c a c d a c d a c d b c b c c d b c a c d			

II

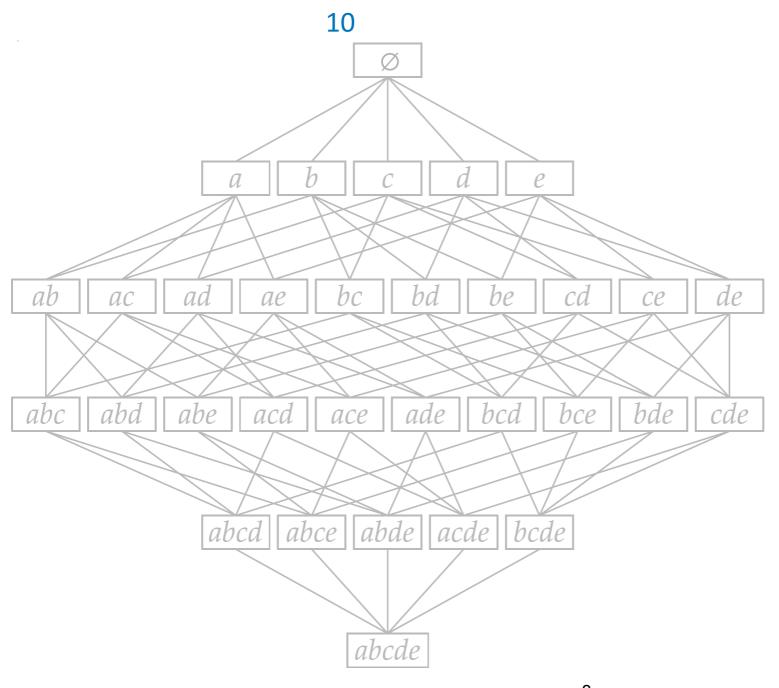
## Example(1)



$n_D$							
1:	a			d	е		
2:		b	С	d			
3:	a		С		е		
4:	a		С	d	е		
5:	a				е		
6:	a		С	d			
7:		b	С				
8:	a		С	d	е		
9:		b	С		е		
10:	a			d	E		

 $H_{-}$ 

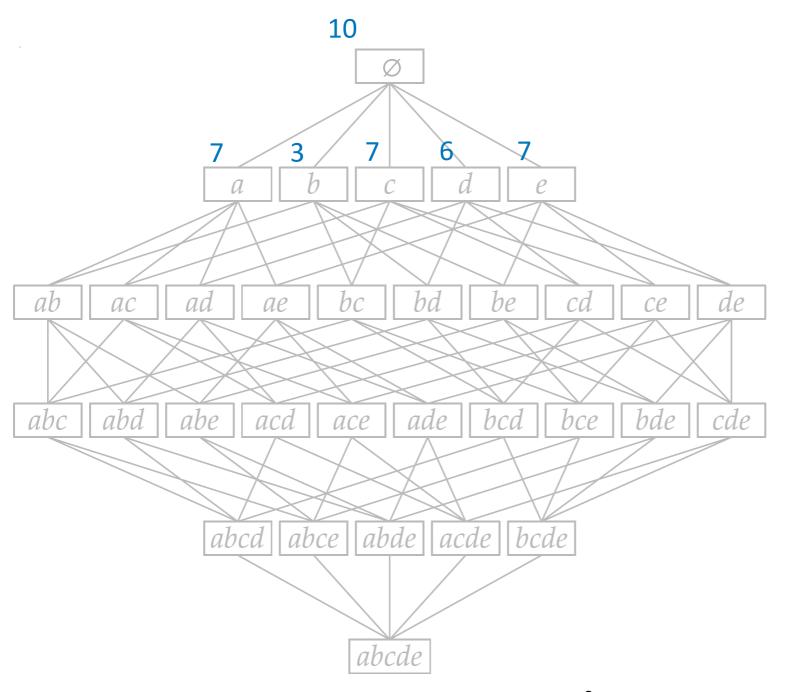
## Example(1)



$H_D$						
1:	a			d	е	
2:		b	С	d		
3:	a		С		е	
4:	а		С	d	е	
5:	а				е	
6:	a		С	d		
7:		b	С			
8:	a		С	d	е	
9:		b	С		е	
10:	а			d	E	

II

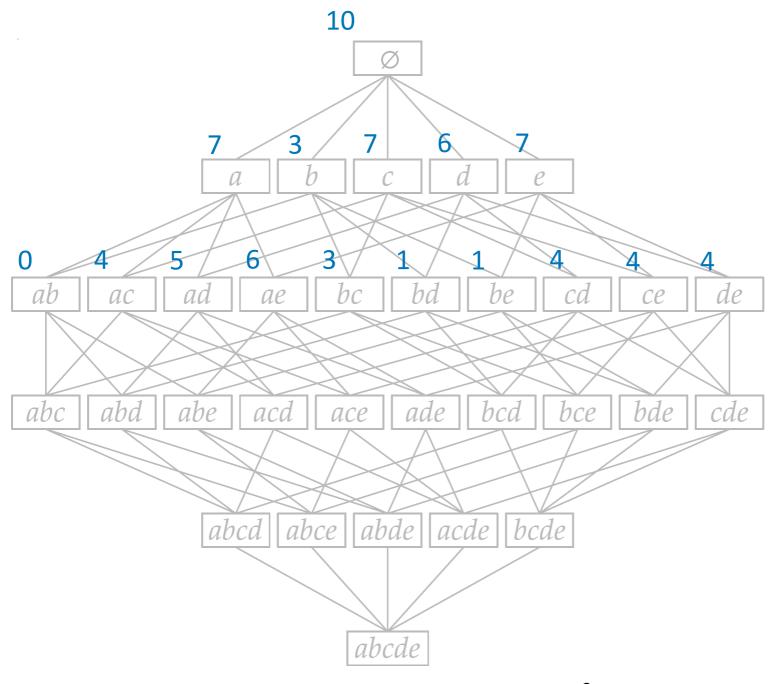
## Example(1)



$H_D$						
1:	a			d	e	
2:		b	С	d		
2: 3: 4: 5: 6:	a		С		е	
4:	a		С	d	е	
5:	a				е	
6:	a		С	d		
7:		b	С			
8: 9:	a		С	d	е	
9:		b	С		е	
10:	a			d	E	

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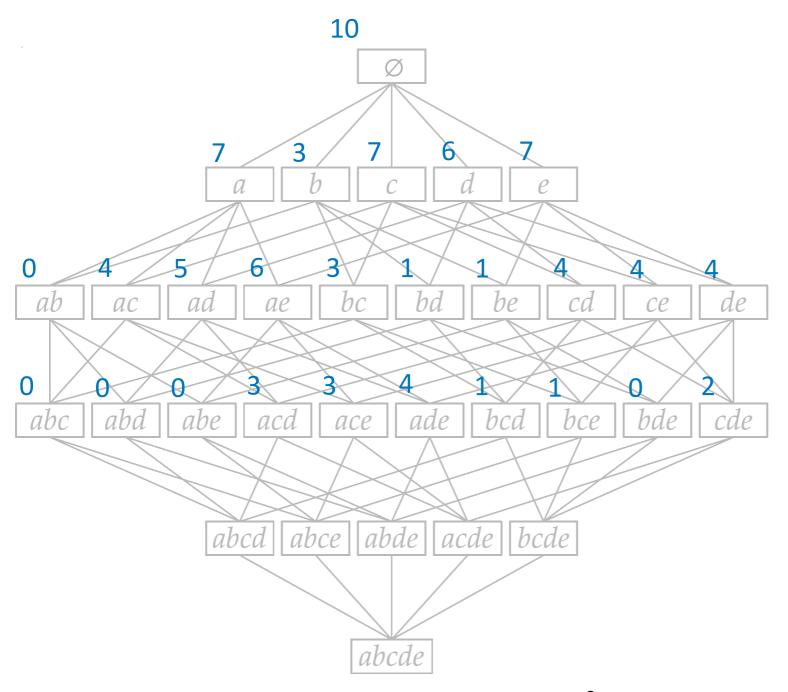
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$H_D$							
1:	a			d	е		
2:		b	С	d			
3:	a		С		е		
4:	a		С	d	е		
5:	a				е		
6:	a		С	d			
7:		b	С				
8:	a		С	d	е		
9:		b	С		е		
10:	а			d	E		

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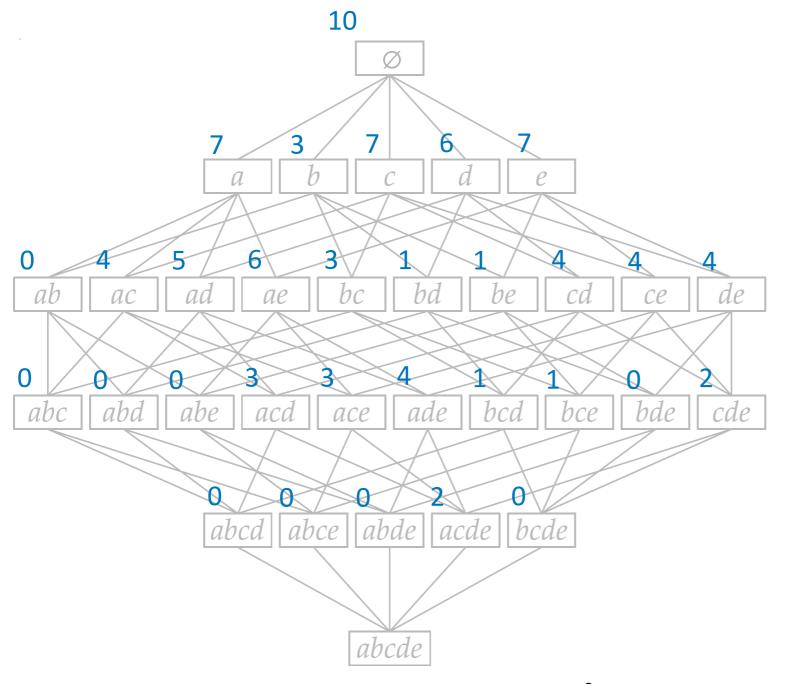
### Example(1)



)				
a			d	е
	b	С	d	
a		С		е
a		С	d	е
a				е
a		С	d	
	b	С		
a		С	d	е
	b	С		е
a			d	E
	a a a	a b a b a b	a b c a c a c a c a c a c b c a c b c a c	a d b c d a c a c d a c d b c a c d b c a c d b c

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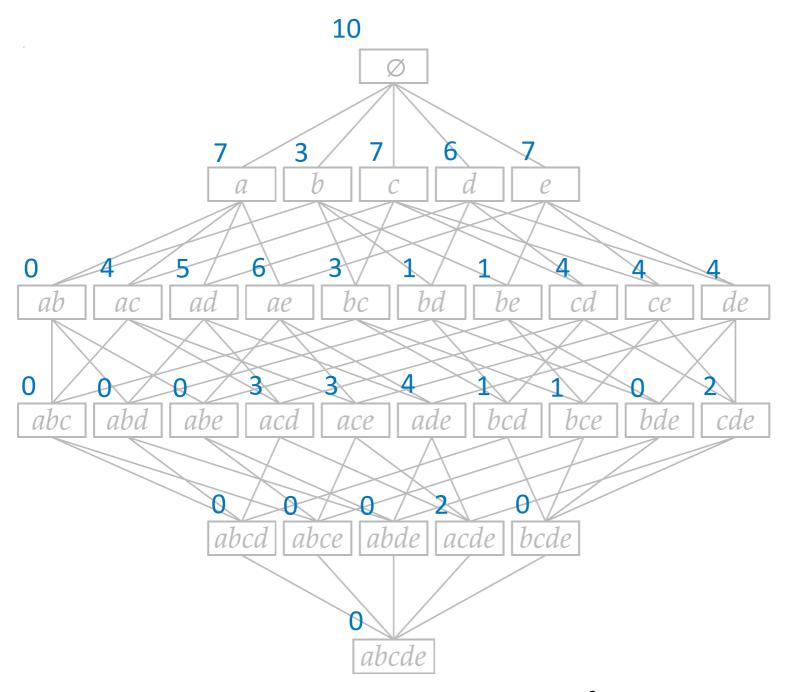
## Example(1)



$H_D$						
1:	a			d	e	
2:		b	С	d		
2: 3: 4: 5: 6:	a		С		е	
4:	a		С	d	е	
5:	a				е	
6:	a		С	d		
7:		b	С			
8: 9:	a		С	d	е	
9:		b	С		е	
10:	a			d	E	

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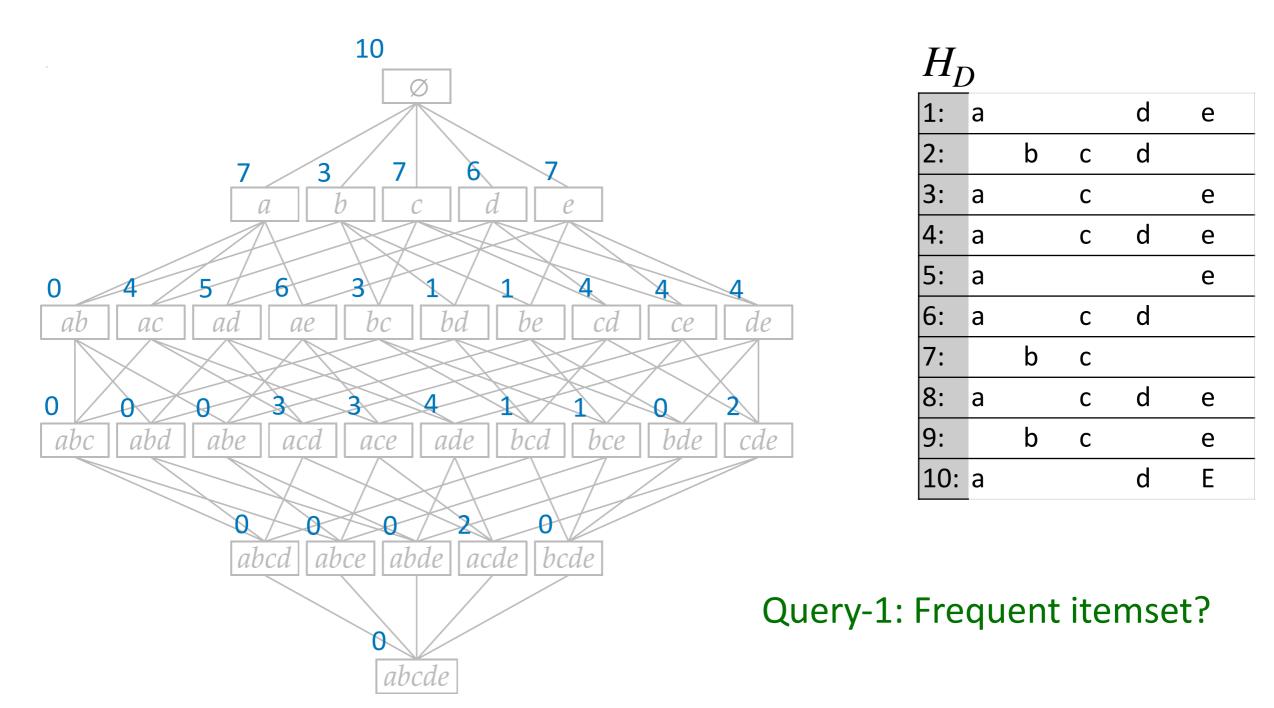
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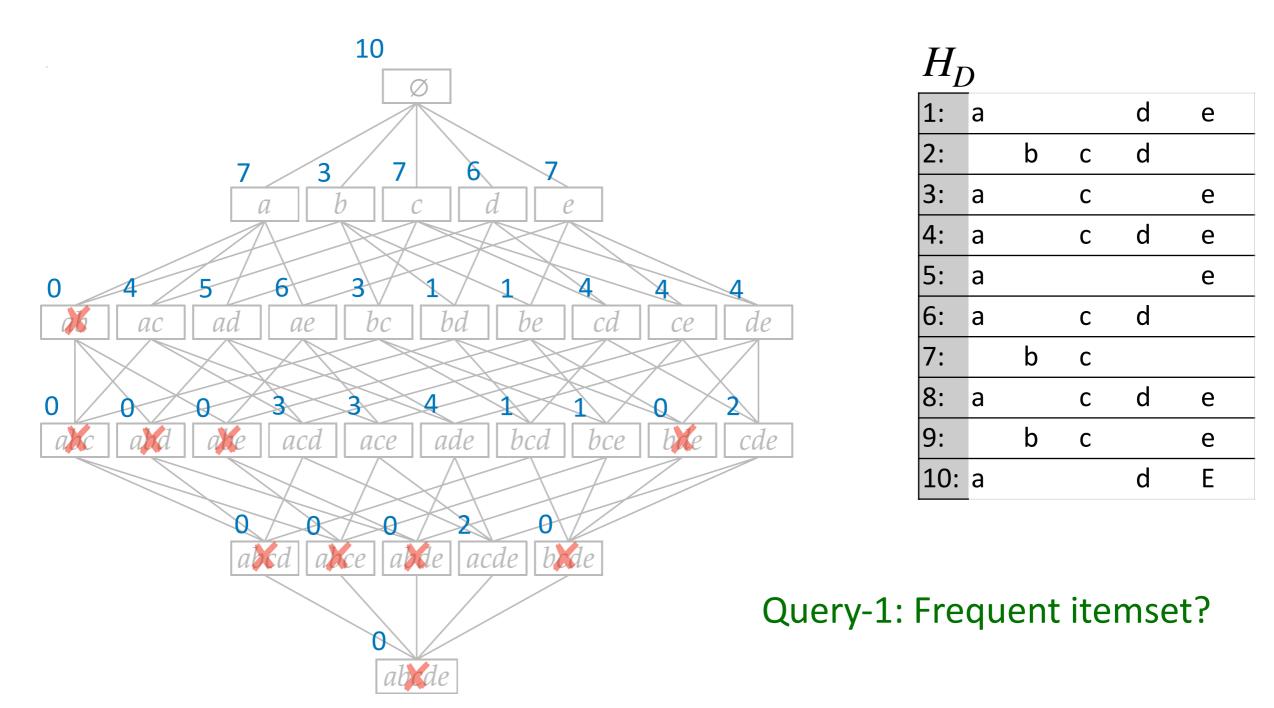
1:       a       d       e         2:       b       c       d         3:       a       c       e         4:       a       c       d       e         5:       a       c       d       e         6:       a       c       d       e         7:       b       c       e         9:       b       c       e         10:       a       d       E	$n_L$	)				
3: a c e 4: a c d e 5: a e 6: a c d 7: b c 8: a c d e 9: b c e	1:	a			d	е
4: a c d e 5: a e 6: a c d 7: b c 8: a c d e 9: b c e	2:		b	С	d	
5: a e 6: a c d 7: b c 8: a c d e 9: b c e	3:	a		С		е
6: a c d 7: b c 8: a c d e 9: b c e	4:	a		С	d	е
7: b c 8: a c d e 9: b c e	5:	a				е
8: a c d e 9: b c e	6:	a		С	d	
9: b c e	7:		b	С		
	8:	a		С	d	е
10: a d E	9:		b	С		е
	10:	a			d	E

 $\boldsymbol{H}$ 

#### Example(1)

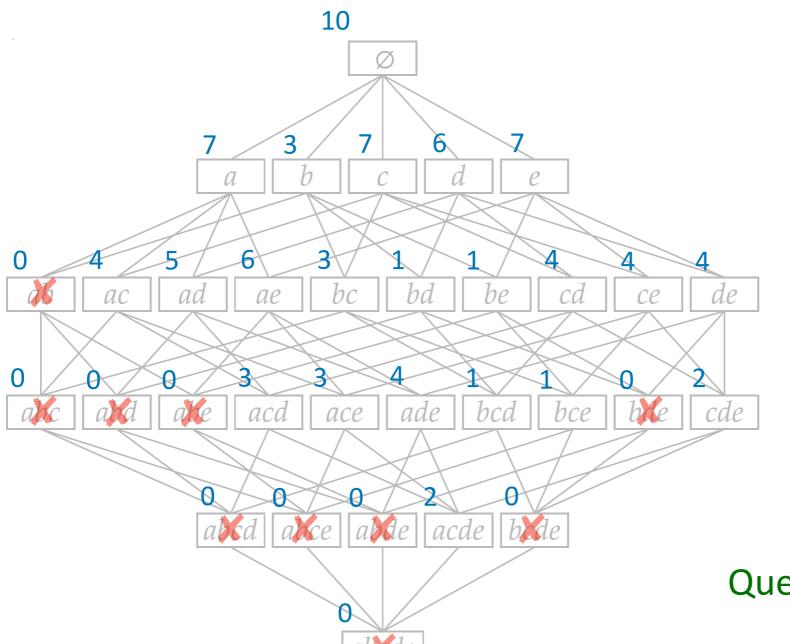


#### Example(1)



#### Example(1)

Query-2: Frequent itemset with minimum support  $\alpha = 3$ ?

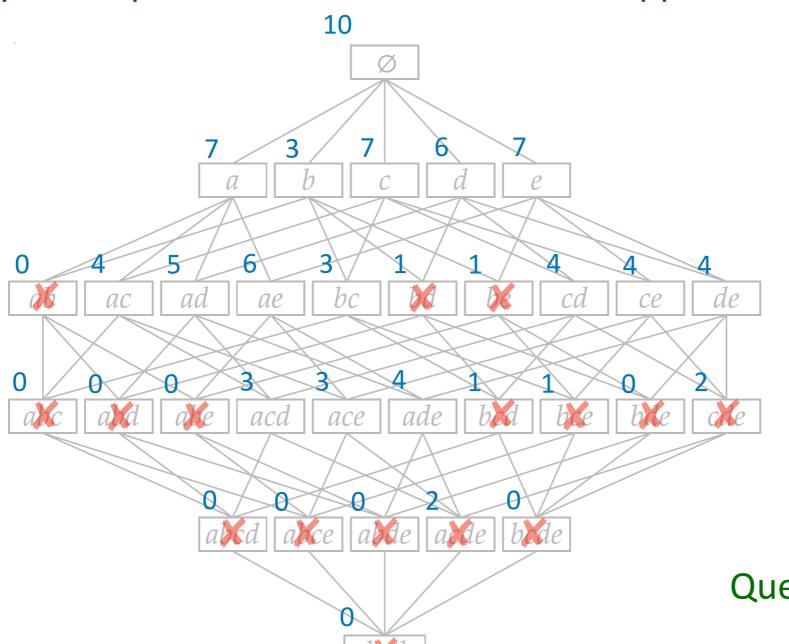


$H_D$						
1:	а			d	е	
2:		b	С	d		
3:	a		С		е	
4:	a		С	d	е	
5:	a				е	
6:	а		С	d		
7:		b	С			
8:	а		С	d	е	
9:		b	С		е	
10:	а			d	E	

Query-1: Frequent itemset?

### Example(1)

Query-2: Frequent itemset with minimum support  $\alpha = 3$ ?



$H_D$						
1:	а			d	е	
2:		b	С	d		
3:	a		С		е	
4:	a		С	d	е	
5:	a				е	
6:	a		С	d		
7:		b	С			
8:	a		С	d	е	
9:		b	С		е	
10:	а			d	E	

Query-1: Frequent itemset?

TT

#### Naïve Search

 A naïve search that consists of enumerating and testing the frequency of itemset candidates in a given dataset is usually infeasible

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 A naïve search that consists of enumerating and testing the frequency of itemset candidates in a given dataset is usually infeasible

Number of items (n)	Search space (2 <sup>n</sup> )
10	≈ 10 <sup>3</sup>
20	≈ <b>10</b> <sup>6</sup>
30	≈ 10 <sup>9</sup>
100	≈ <b>10</b> <sup>30</sup>
128	≈ $10^{68}$ (atoms in the universe)
1000	≈ <b>10</b> <sup>301</sup>

#### **Anti-Monotonicity Property**

Given a transaction database D over items I and two itemsets P and Q:

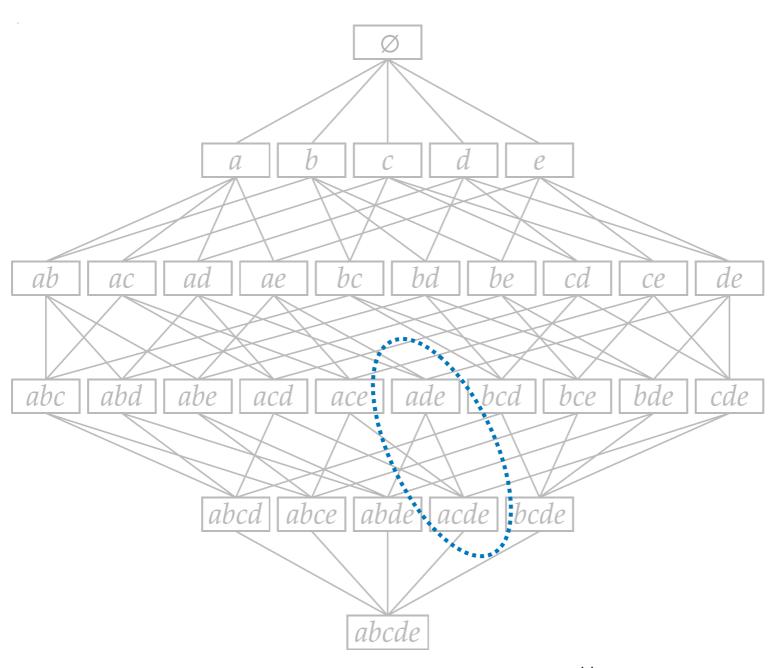
$$Q \subseteq P \Rightarrow cover(P) \subseteq cover(Q)$$

That is,

$$Q \subseteq P \Rightarrow freq(P) \leq freq(Q)$$

## **Anti-Monotonicity Property**

### Example(2)

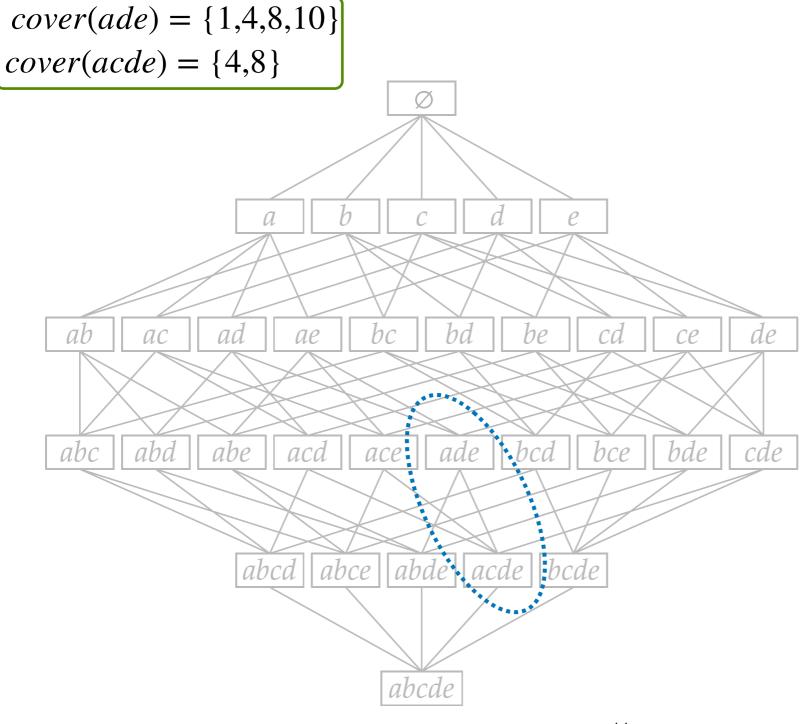


$II_L$	)				
1:	a			d	е
2:		b	С	d	
3:	a		С		е
4:	a		С	d	е
5:	a				е
6:	a		С	d	
7:		b	С		
8:	a		С	d	е
9:		b	С		е
10:	a			d	E

 $\boldsymbol{H}$ 

## **Anti-Monotonicity Property**

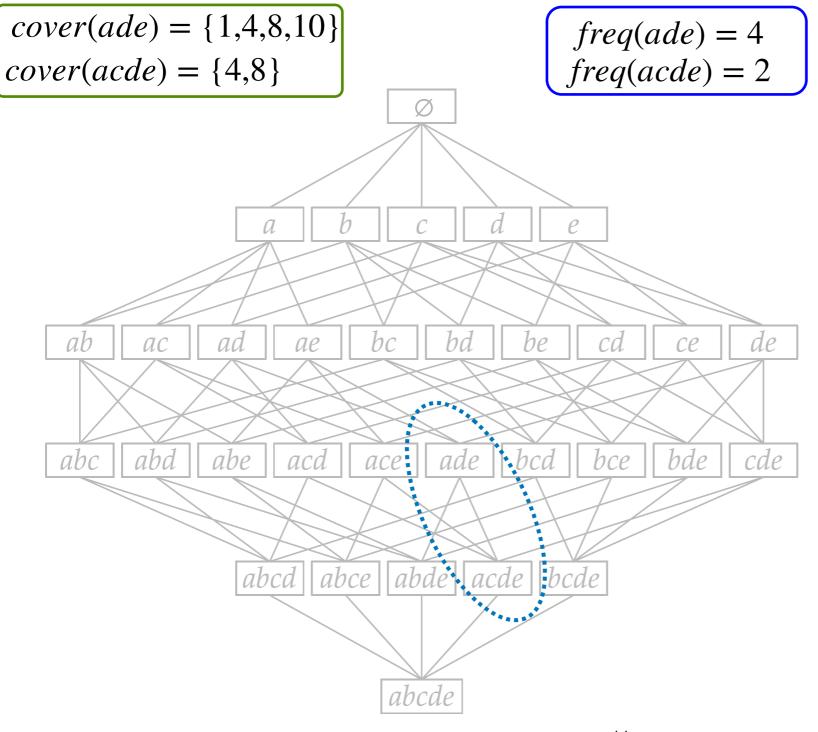
#### Example(2)



$H_{L}$	)				
1:	а			d	е
2:		b	С	d	
3:	a		С		е
4:	a		С	d	е
5:	a				е
6:	a		С	d	
7:		b	С		
8:	a		С	d	е
9:		b	С		е
10:	a			d	E

# **Anti-Monotonicity Property**

#### Example(2)



#### d е b d С 3: С a e d a С e a e d a C b C d a С e b 9: e d 10: a Ε

#### **Apriori Property**

• Given a transaction database D over items I, a minsup  $\alpha$  and two itemsets P and Q:

$$Q \subseteq P \Rightarrow freq(P) \leq freq(Q)$$

• It follows:  $Q \subseteq P \land freq(P) \ge \alpha \Rightarrow freq(Q) \ge \alpha$ 

• Contraposition:  $Q \subseteq P \land freq(Q) < \alpha \Rightarrow freq(P) < \alpha$ 

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All subsets of a frequent itemset are frequent!

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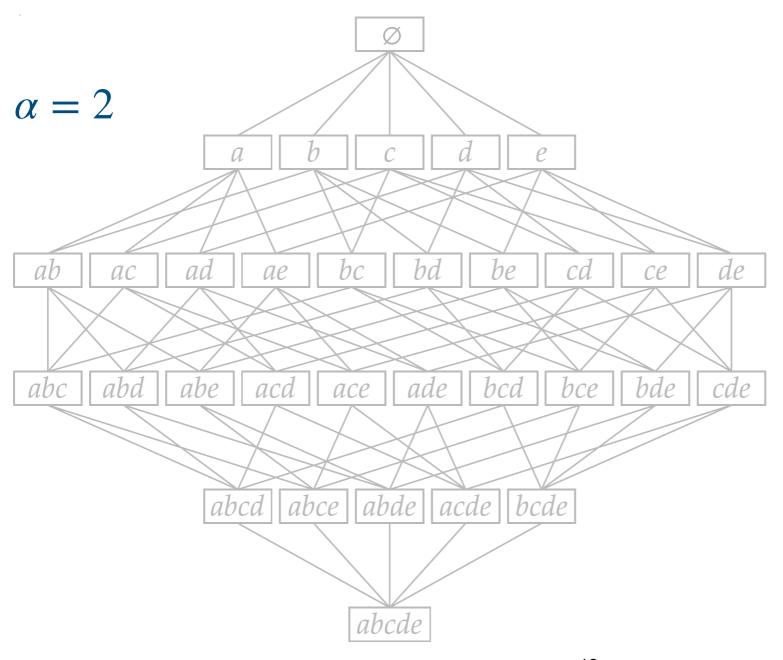
All subsets of a frequent itemset are frequent!

• Contraposition:  $Q \subseteq P \land freq(Q) < \alpha \Rightarrow freq(P) < \alpha$ 

All supersets of an infrequent itemset are infrequent!

## Example(3)

All subsets of a frequent itemset are frequent!

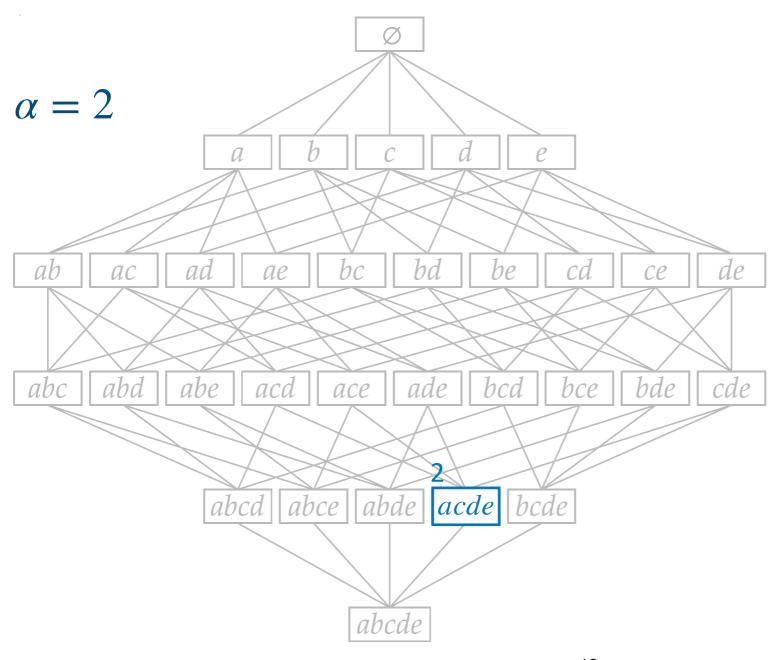


$n_D$							
1:	а			d	е		
2: 3:		b	С	d			
3:	a		С		е		
4: 5: 6:	a		С	d	е		
5:	a				е		
6:	а		С	d			
7:		b	С				
8: 9:	а		С	d	е		
9:		b	С		е		
10:	а			d	E		

 $\boldsymbol{H}$ 

### Example(3)

All subsets of a frequent itemset are frequent!

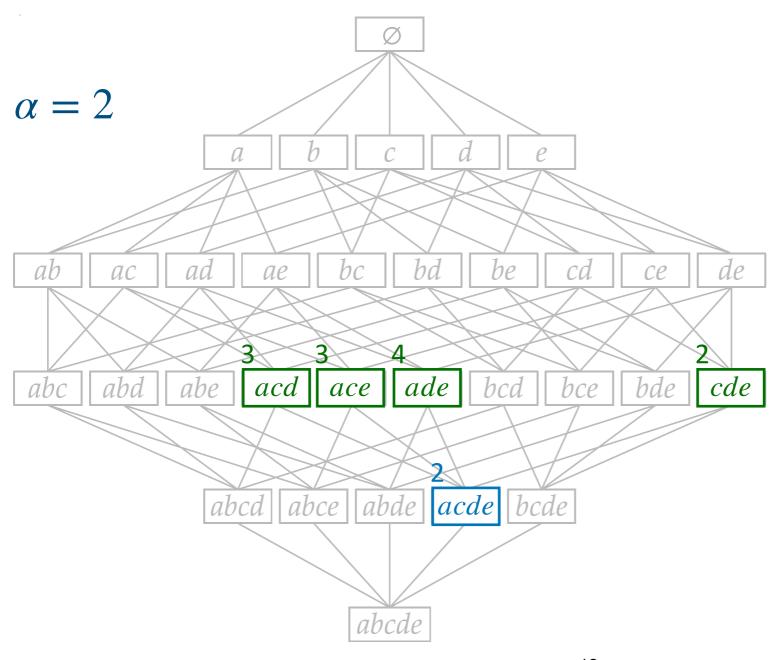


)				
а			d	е
	b	С	d	
a		С		е
a		С	d	е
a				е
a		С	d	
	b	С		
a		С	d	е
	b	С		е
а			d	E
	a a a	a b a b	a b c a c a c a b c a c b c a	a d b c d a c a c d a c d b c b c c d b c c d b c b c b c

 $\boldsymbol{H}$ 

### Example(3)

All subsets of a frequent itemset are frequent!

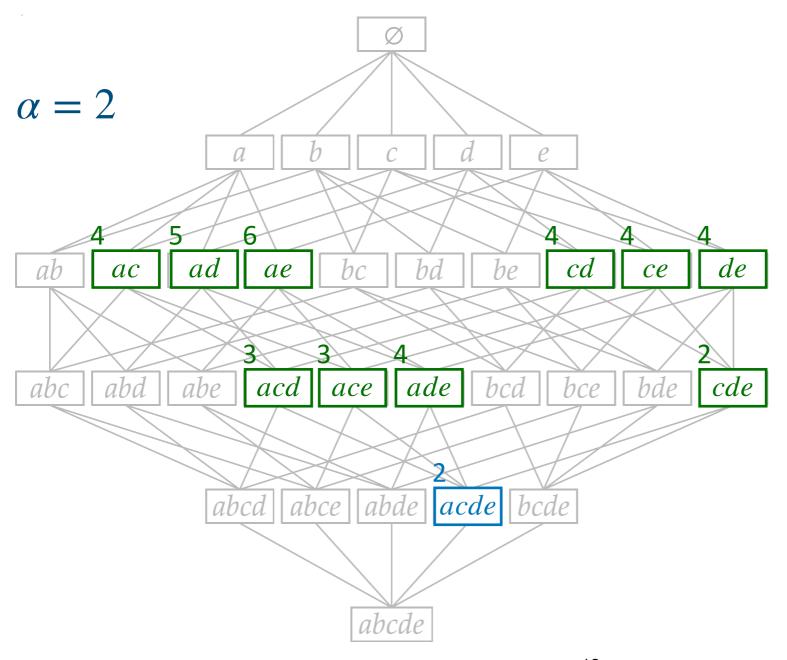


$n_D$							
1:	а			d	е		
2: 3: 4: 5: 6:		b	С	d			
3:	a		С		е		
4:	a		С	d	е		
5:	a				е		
6:	a		С	d			
7:		b	С				
8: 9:	a		С	d	е		
9:		b	С		е		
10:	a			d	E		

 $\boldsymbol{H}$ 

### Example(3)

All subsets of a frequent itemset are frequent!

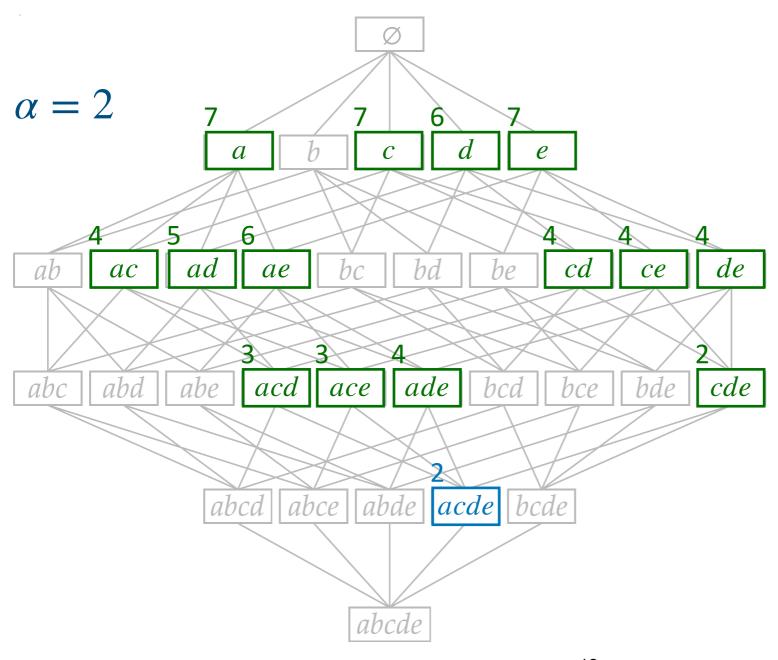


$H_D$							
1:	a			d	е		
2:		b	С	d			
3:	a		С		е		
4:	a		С	d	е		
5:	a				е		
6:	a		С	d			
7:		b	С				
8:	a		С	d	е		
9:		b	С		е		
10:	а			d	E		

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### Example(3)

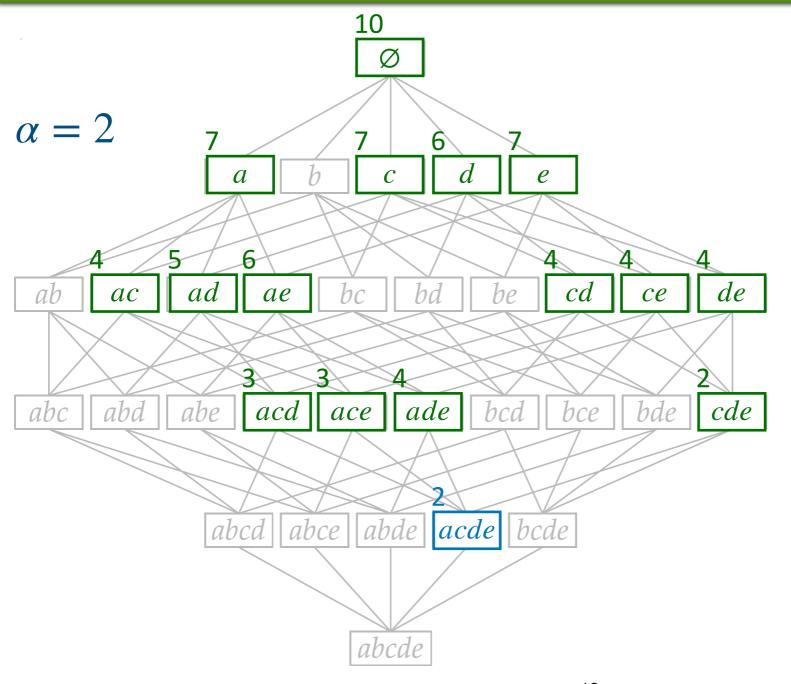
All subsets of a frequent itemset are frequent!



$H_{\underline{D}}$							
1:	a			d	е		
2: 3: 4: 5: 6: 7:		b	С	d			
3:	a		С		е		
4:	a		С	d	е		
5:	a				е		
6:	a		С	d			
7:		b	С				
8: 9:	a		С	d	е		
9:		b	С		е		
10:	a			d	E		

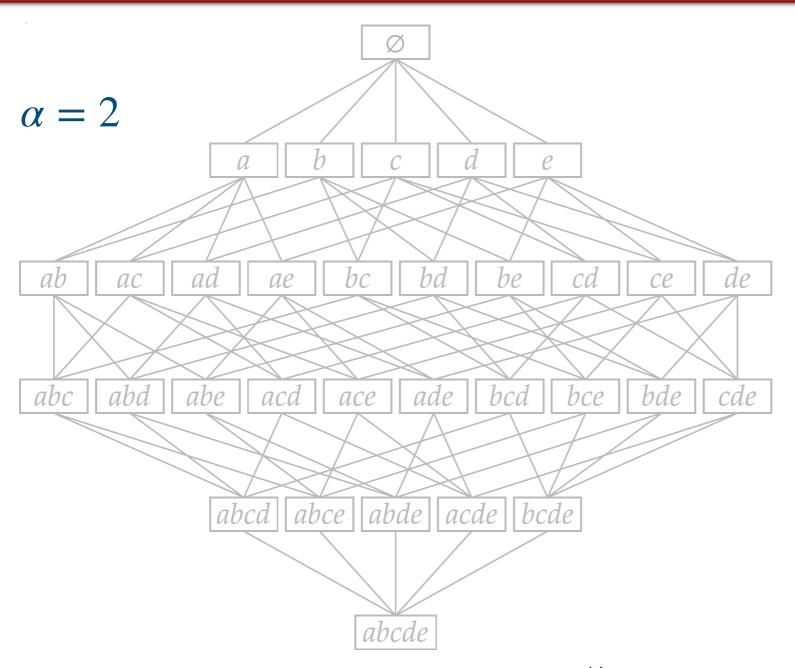
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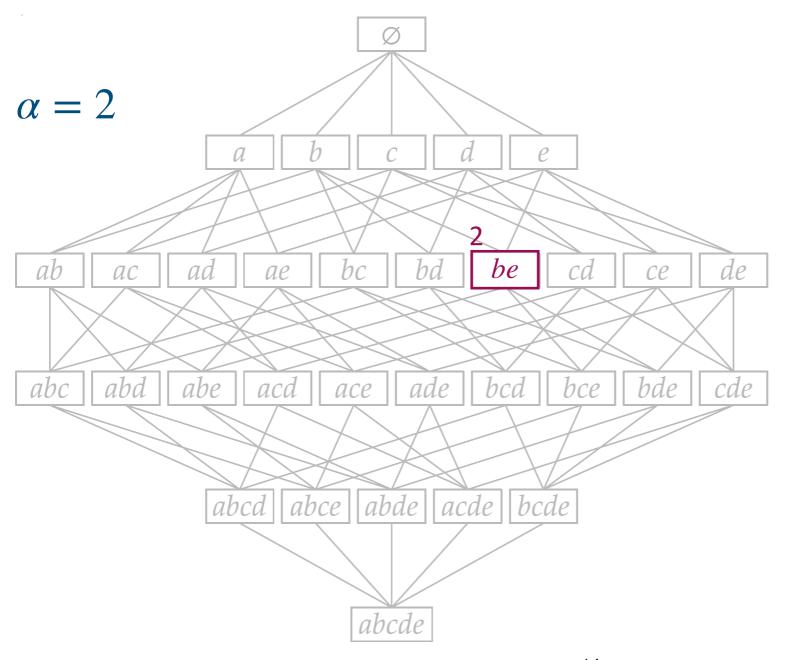
$H_D$								
1:	a			d	е			
2:		b	С	d				
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4:	а		С	d	е			
5:	a				е			
6:	a		С	d				
7:		b	С					
8:	a		С	d	е			
9:		b	С		е			
10:	a			d	Е			

### Example(3)



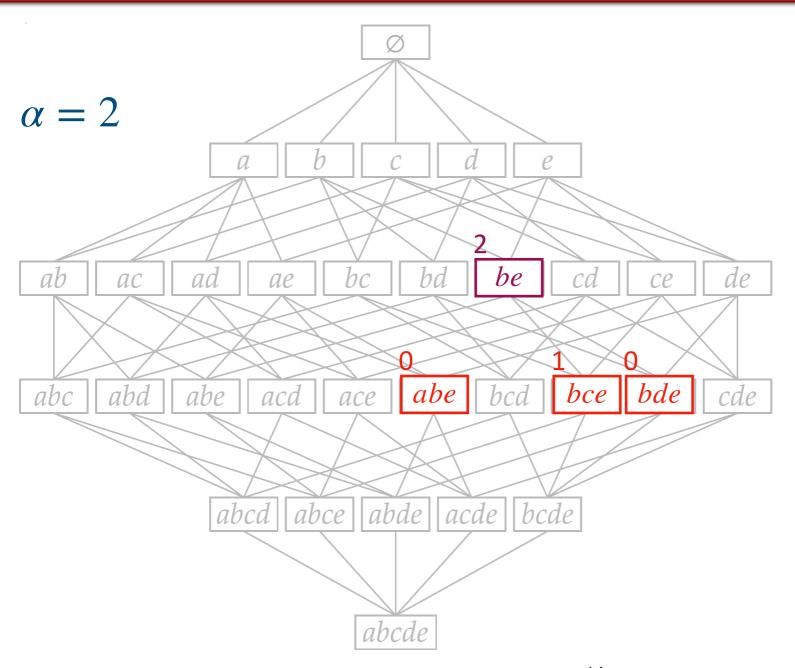
$H_{L}$	)				
1: 2: 3: 4: 5: 6: 7: 8: 9:	а			d	е
2:		b	С	d	
3:	а		С		е
4:	а		С	d	е
5:	a				е
6:	а		С	d	
7:		b	С		
8:	a		С	d	е
9:		b	С		е
10:	a			d	E

### Example(3)



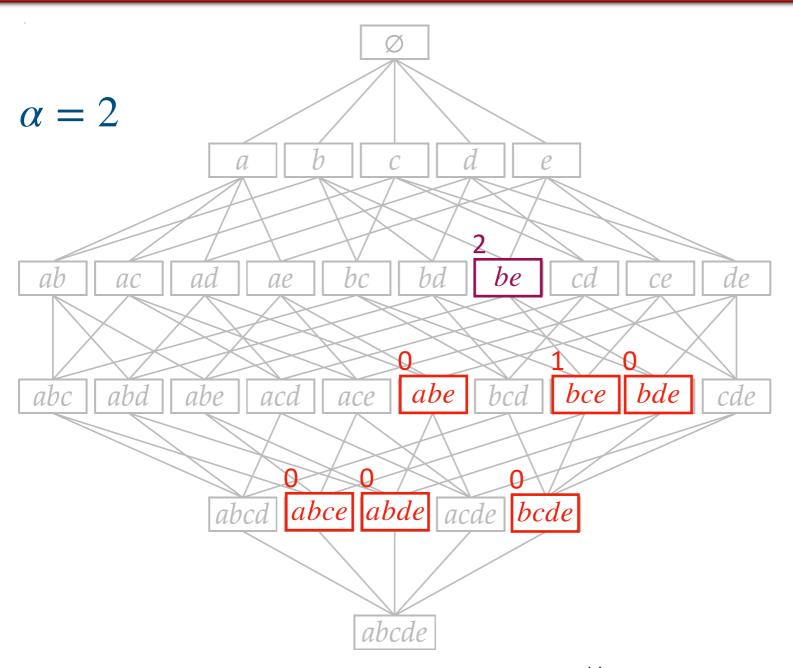
$H_{L}$	)				
1:	a			d	е
2:		b	С	d	
3:	a		С		е
4:	a		С	d	е
5: 6:	a				е
6:	a		С	d	
7:		b	С		
8:	a		С	d	е
9:		b	С		е
10:	а			d	E

### Example(3)



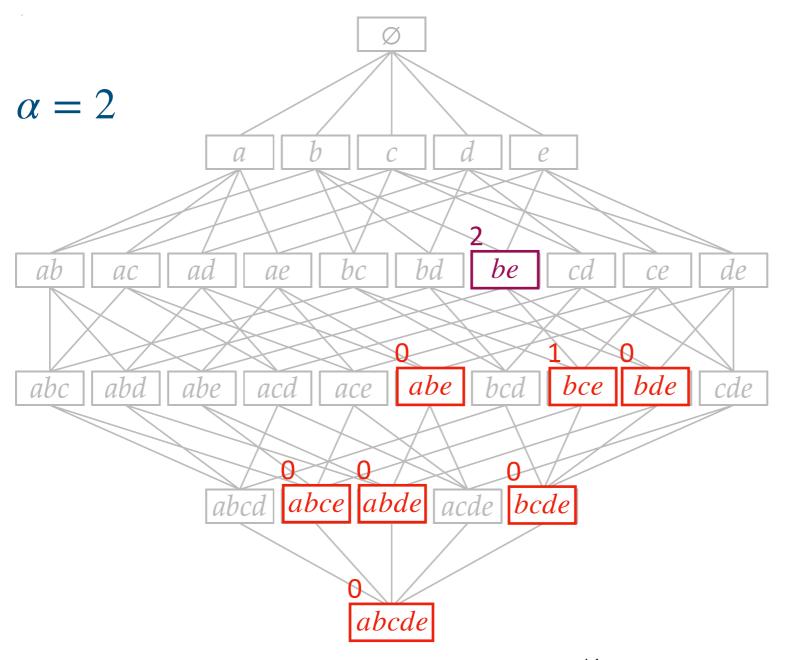
$H_{L}$	)				
1:	a			d	е
2:		b	С	d	
3:	a		С		е
4:	a		С	d	е
5:	a				е
6:	a		С	d	
7:		b	С		
8:	a		С	d	е
9:		b	С		е
10:	a			d	E

## Example(3)



$H_{L}$	)				
1:	a			d	е
2:		b	С	d	
3:	a		С		е
4:	а		С	d	е
5:	a				е
6:	a		С	d	
7: 8: 9:		b	С		
8:	a		С	d	е
9:		b	С		е
10:	a			d	E

### Example(3)



$H_{L}$	)				
1:	a			d	е
2:		b	С	d	
3:	a		С		е
4:	a		С	d	е
5: 6:	a				е
6:	a		С	d	
7: 8:		b	С		
8:	a		С	d	е
9:		b	С		е
10:	a			d	E

Poset  $(2^I, \subseteq)$ 

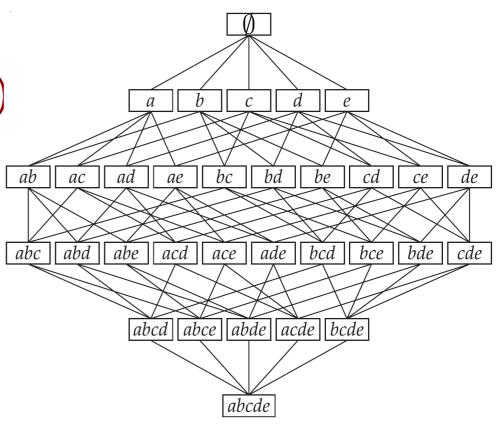
- Test A partial order is a binary relation R over a set S:
- $\forall x, y, z \in S$
- x R x (reflexivity)
- $x R y \land y R x \Rightarrow x = y$  (anti-symmetry)
- $x R y \land y R z \Rightarrow x R z$  (transitivity)
- What are S and R in Itemset Mining?

Poset  $(2^I, \subseteq)$ 

- Test A partial order is a binary relation R over a set S:
- $\forall x, y, z \in S$
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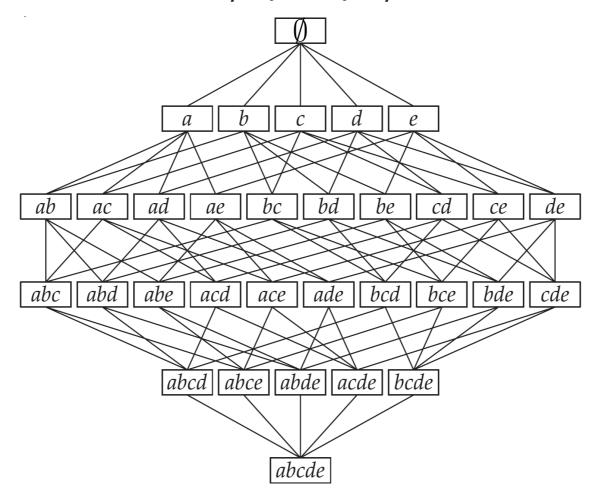
(reflexivity)

- $x R y \land y R x \Rightarrow x = y$  (anti-symmetry)
- $x R y \land y R z \Rightarrow x R z$  (transitivity)
- What are S and R in Itemset Mining?



### Partially ordered sets (review)

- Comparable itemsets:  $x \subseteq y \lor y \subseteq x$
- Incomparable itemsets:  $x \not\subseteq y \land y \not\subseteq x$



#### **Apriori Algorithm** [Agrawal and Srikant 1994]

- Determine the support of the one-element item sets (i.e. singletons) and discard the infrequent items
- Form candidate itemsets with two items (both items must be frequent), determine their support, and discard the infrequent itemsets
- Form candidate item sets with three items (all contained pairs must be frequent), determine their support, and discard the infrequent itemsets
- And so on!

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- And so on!

Based on candidate generation and pruning

#### **Apriori Algorithm** [Agrawal and Srikant 1994]

```
Algorithm 1: Apriori Algorithm
  Input: Transaction database \mathcal{D}, minimum support threshold \alpha
  Output: Frequent itemsets
  k \leftarrow 1;
  L_k \leftarrow \{p_i \mid p_i \in \mathcal{I} \land \mathtt{freq}(p_i) \geq \alpha\} ;
  while L_k \neq \emptyset do
       C \leftarrow \operatorname{aprioriGen}(L_k);
      k \leftarrow k+1;

L_k \leftarrow \{c \mid c \in C \land freq(c) \ge \alpha\};
  return \bigcup_i L_i;
```

#### **Apriori Algorithm** [Agrawal and Srikant 1994]

```
Function aprioriGen(L_k):
      E \leftarrow \emptyset;
      for each pair of itemsets P', P'' \in L_k such that
        P' = \{p_{i_1}, \dots, p_{i_{k-1}}, p_{i_k}\} and P'' = \{p_{i_1}, \dots, p_{i_{k-1}}, p_{i'_k}\} do
            if p_{i_k} \neq p_{i'_k} then

\begin{vmatrix}
P \leftarrow P' \cup P'' ; \\
\mathbf{if} \quad \forall p_i \in P, P \setminus \{p_i\} \in L_k \mathbf{then} \\
E \leftarrow E \cup \{P\};
\end{vmatrix}

      return E;
```

# **Apriori Algorithm** [Agrawal and Srikant 1994] Improving candidates generation

• Using aprioriGen function, an item of k+1 size can be generated in a  $\delta$  possible ways:

$$\delta = \frac{k(k+1)}{2}$$

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• Using aprioriGen function, an item of k+1 size can be generated in a  $\delta$  possible ways:

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6 possibilities to generate (abcd)

	abc	abd	acd	bcd
abc	_	abcd	abcd	abcd
abd	abcd	_	abcd	abcd
acd	abcd	abcd	_	abcd
bcd	abcd	abcd	abcd	_

# **Apriori Algorithm** [Agrawal and Srikant 1994] Improving candidates generation

• Using aprioriGen function, an item of k+1 size can be generated in a  $\delta$  possible ways:

$$\delta = \frac{k(k+1)}{2}$$

Need: Generate itemset candidate only once

How: Assign a unique parent itemset to each candidate, ensuring that it is generated only from its parent itemset

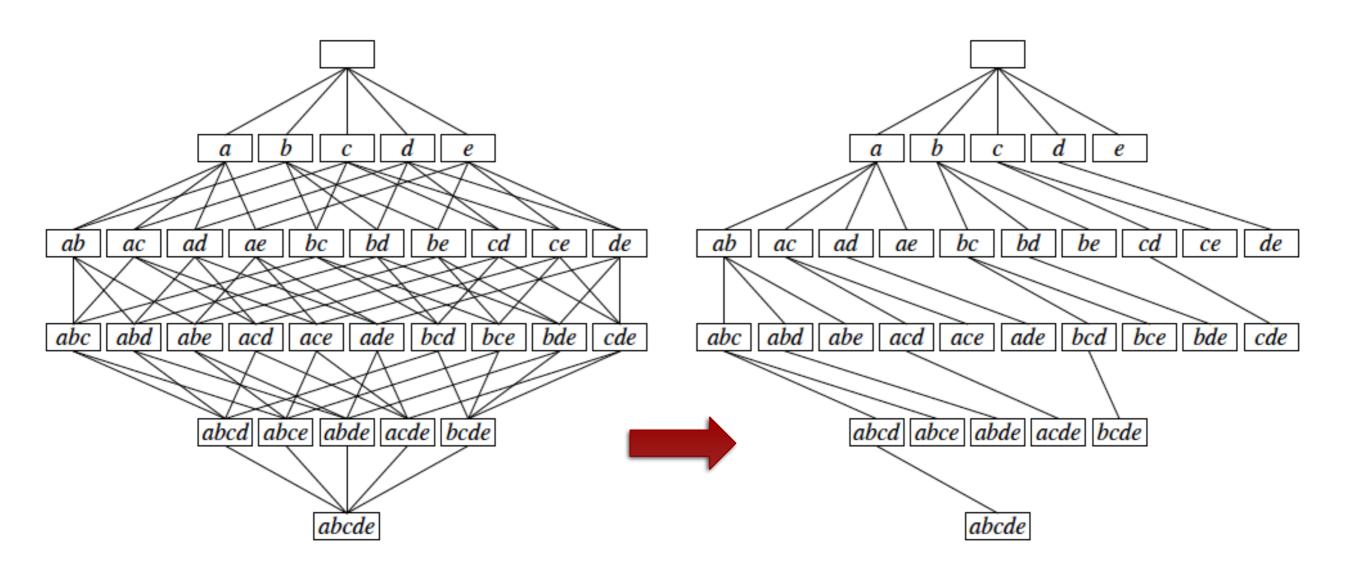
# **Apriori Algorithm** [Agrawal and Srikant 1994] Improving candidates generation

Assigning unique parents turns the poset lattice into a tree:

# Apriori Algorithm [Agrawal and Srikant 1994]

### Improving candidates generation

Assigning unique parents turns the poset lattice into a tree:



## Apriori Algorithm [Agrawal and Srikant 1994]

#### **Canonical form for itemsets**

- ullet An itemset can be represented as a word over an alphabet I
  - Question: how many words of k items can we have?
  - Answer: k-permutations of k items: k!
- By imposing an arbitrary order (e.g., lexicography order) on the items, we can define a canonical form—a unique representation of itemsets that eliminates symmetries
  - Lex on items : abc < acb < bac < bca...
  - $\kappa(abc) = \kappa(acb) = \kappa(bac) = \kappa(bca) = abc$
  - $\kappa(abc,1) = a; \ \kappa(abc,2) = b; \ \kappa(abc,3) = c$

# **Apriori Algorithm** [Agrawal and Srikant 1994] Recursive processing with Canonical forms

 For each itemset of a given size, generate all possible extensions of the itemset by adding one item, subject to the condition that:

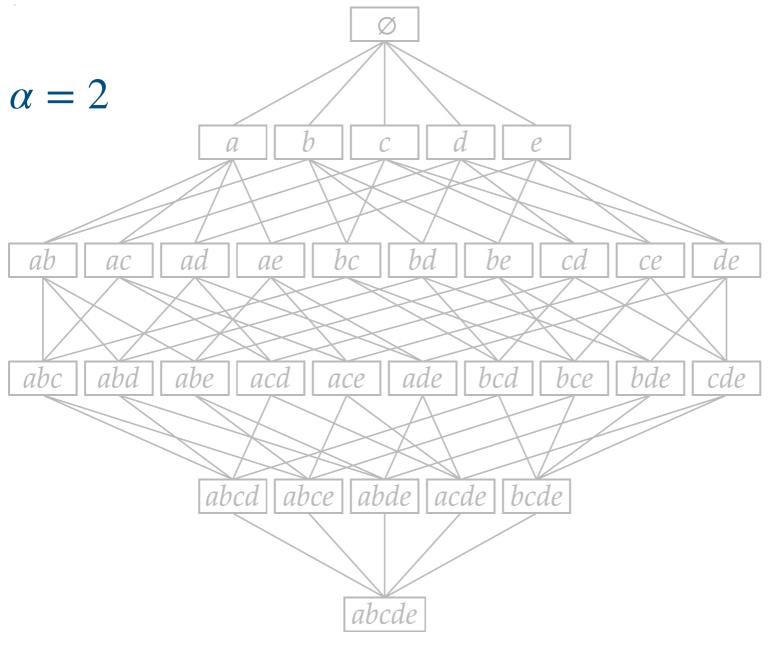
# **Apriori Algorithm** [Agrawal and Srikant 1994] Recursive processing with Canonical forms

• For each itemset of a given size, generate all possible extensions of the itemset by adding one item, subject to the condition that:

$$child(P, \alpha) = \{P' : (P' = P \cup \{p_i\}) \land (\kappa(P, |P|) < p_i) \land (freq(P') \ge \alpha)\}$$

### Example(4)

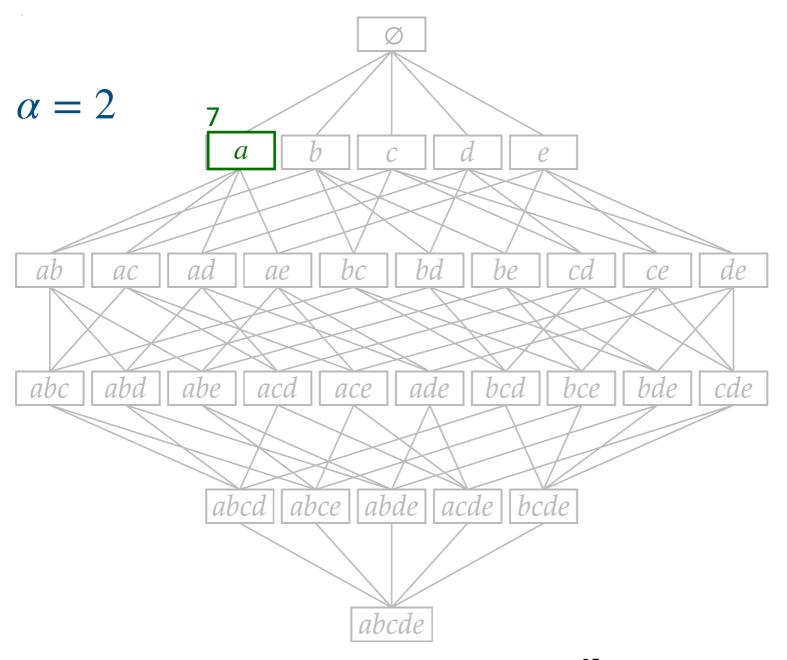
 $child(P,\alpha) = \{P': (P' = P \cup \{p_i\}) \land (\kappa(P, \lfloor P \rfloor) < p_i) \land (freq(P') \ge \alpha)\}$ 



$H_D$							
1:	a			d	е		
2: 3: 4: 5: 6: 7:		b	С	d			
3:	a		С		е		
4:	a		С	d	е		
5:	a				е		
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7:		b	С				
8: 9:	a		С	d	е		
9:		b	С		е		
10:	a			d	E		

### Example(4)

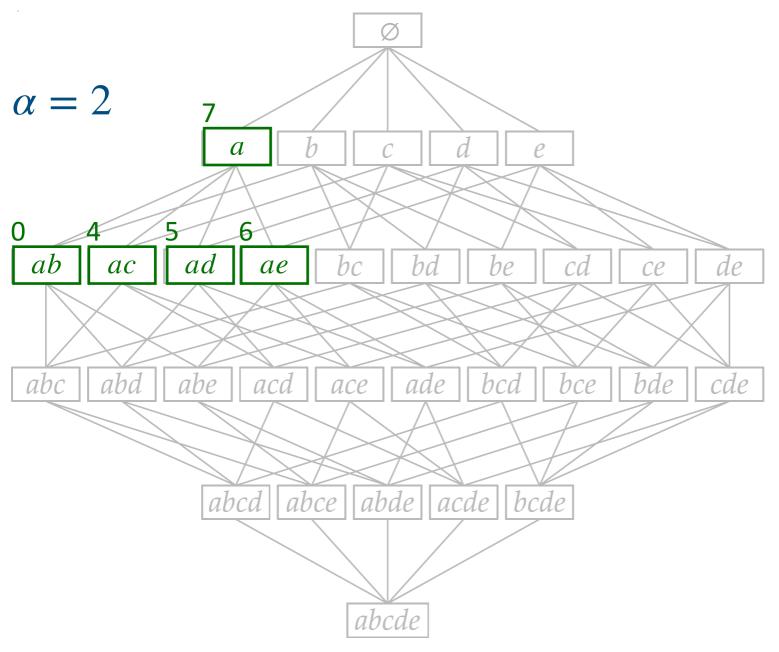
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$H_D$							
a			d	е			
	b	С	d				
a		С		е			
a		С	d	е			
a				е			
a		С	d				
	b	С					
a		С	d	е			
	b	С		е			
а			d	Е			
	a a a	a b a b a b	a b c a c a c a b c a c a b c a	a d b c d a c a c d a c d a c d b c b c a c d b c a c d			

### Example(4)

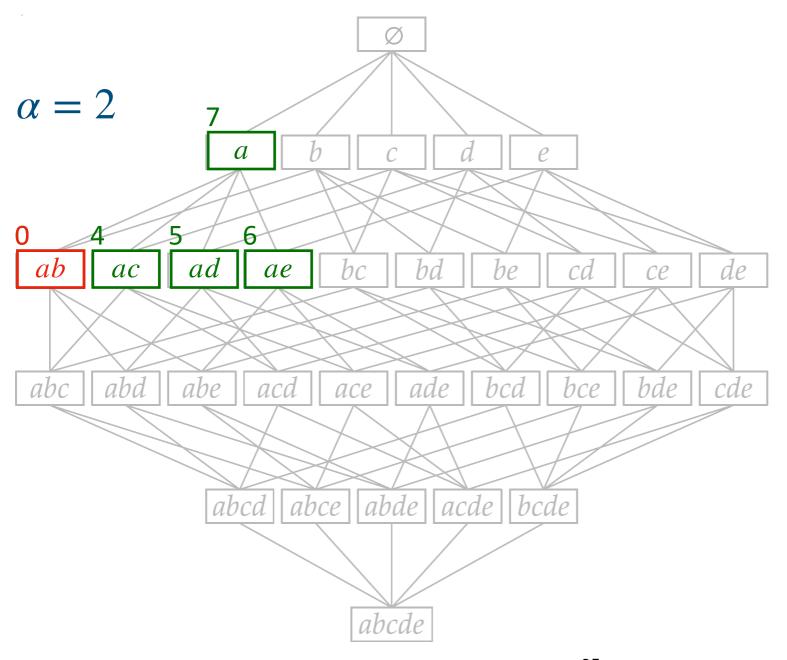
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4:	а		С	d	е			
5:	a				е			
6:	a		С	d				
7:		b	С					
8:	a		С	d	е			
9:		b	С		е			
10:	a			d	Е			

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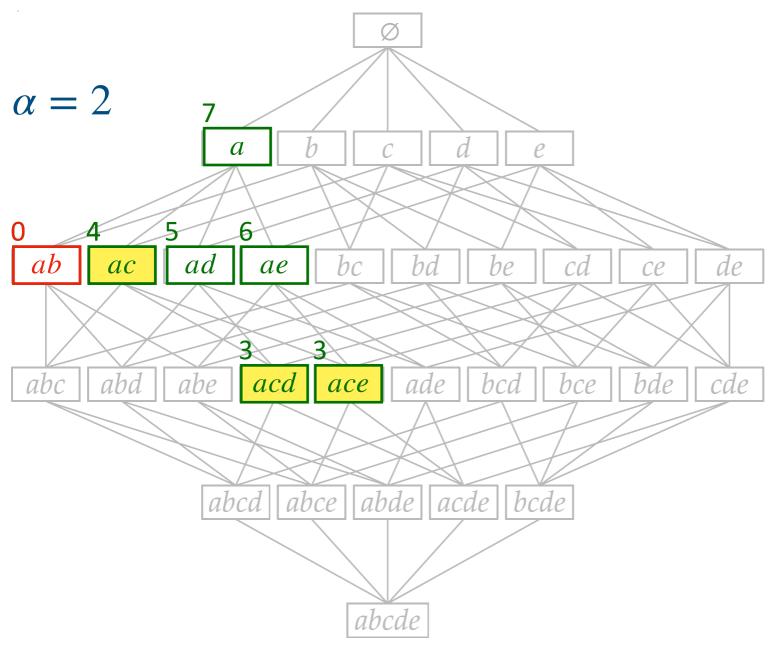
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4:	a		С	d	е		
5:	a				е		
6:	a		С	d			
7:		b	С				
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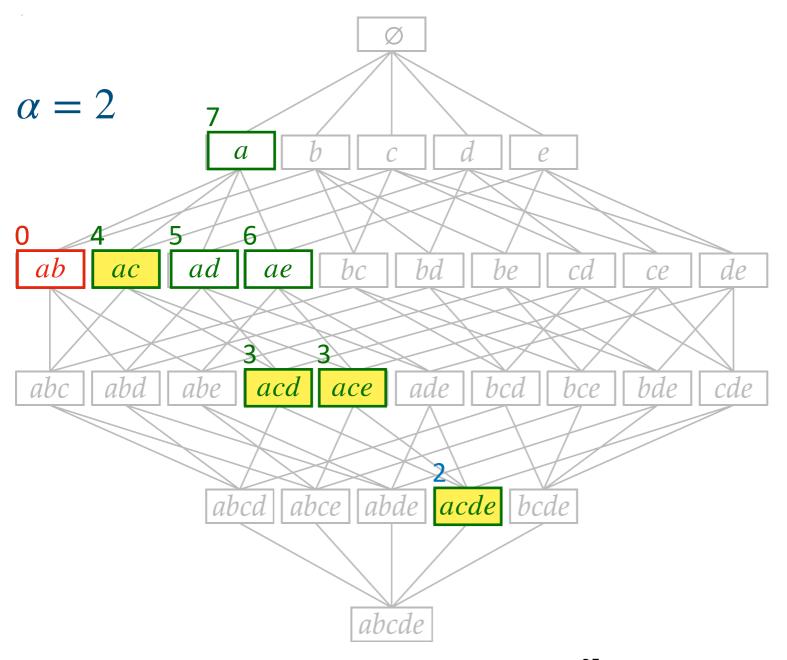
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7:		b	С					
8:	a		С	d	е			
9:		b	С		е			
10:	a			d	Е			

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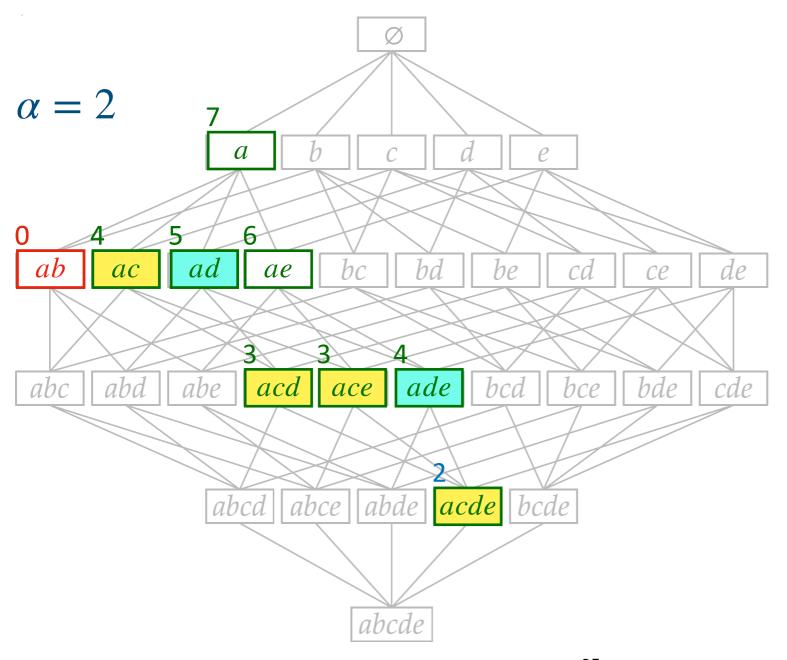
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$H_D$							
a			d	е			
	b	С	d				
a		С		е			
a		С	d	е			
a				е			
a		С	d				
	b	С					
a		С	d	е			
	b	С		е			
а			d	Е			
	a a a	a b a b a b	a b c a c a c a b c a c a b c a	a d b c d a c a c d a c d a c d b c b c a c d b c a c d			

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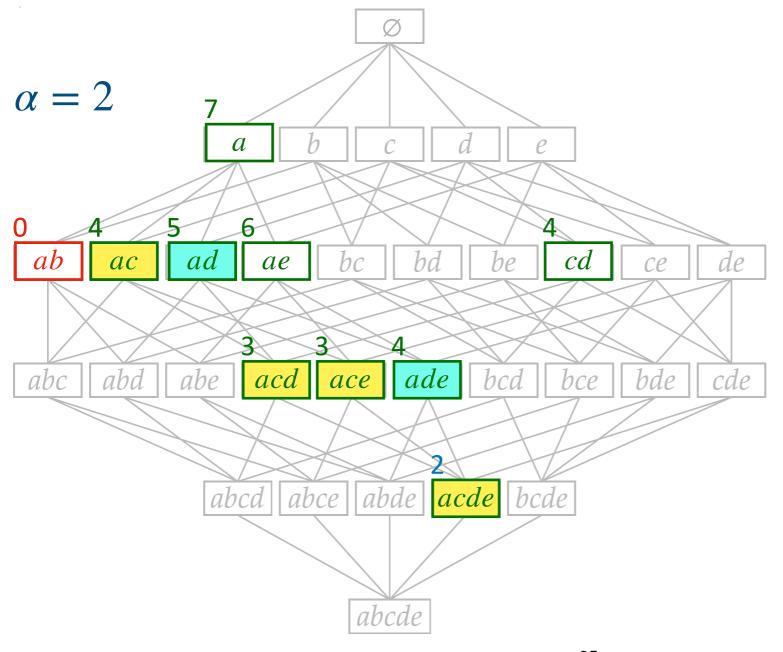
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$H_D$								
1:	а			d	е			
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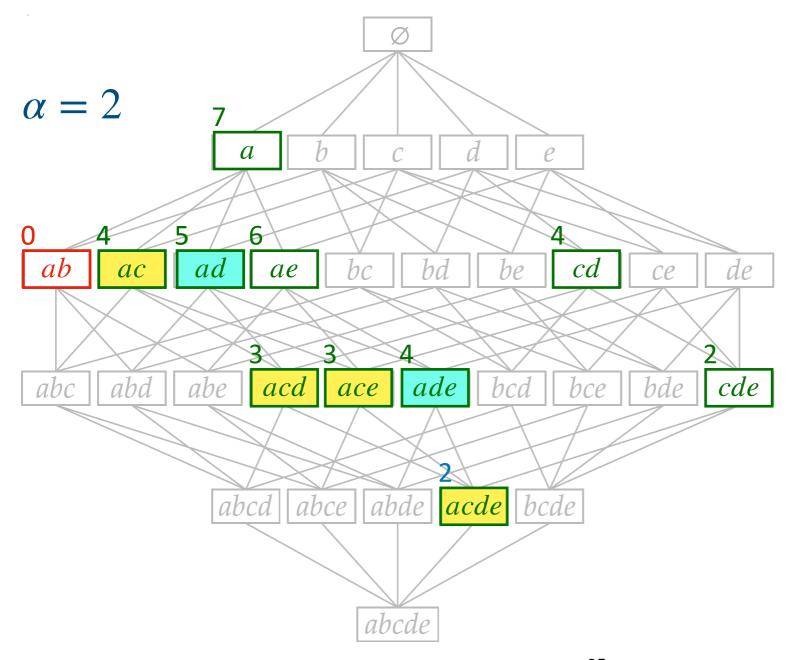
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1: 2: 3:		b	С	d			
3:	a		С		е		
4: 5: 6: 7:	a		С	d	е		
5:	a				е		
6:	a		С	d			
7:		b	С				
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5:	a				е		
6:	a		С	d			
7:		b	С				
8: 9:	a		С	d	е		
9:		b	С		е		
10:	a			d	E		

### **Items Ordering**

- Any order can be used, but:
  - The structure of the search space depends heavely on the chosen order
  - Algorithm efficiency varies significantly based on the item order
- Advanced methods dynamically adjust the order during teh search:
  - Use different but « compatible » orders in different branches

# Searching for Frequent Itemsets

**Items Ordering (heuristics)** 

• Sort the items w.r.t. their frequency (decreasing/increasing) : Frequent itemsets are composed of frequent items

 Sort items based on the size of the area they cover : considering both their frequency and the sizes of the transactions that include them

# Condensed representation of Frequent Itemsets:

Closed and Maximal Itemsets

### **Definition**

The set of Maximal (frequent) Itemsets:

$$M_{\alpha} = \{ P \subset I | freq(P) \geq \alpha \land \forall P' \supset P : freq(P') < \alpha \}$$

That is:

$$\forall \alpha, \forall P \in F_{\alpha} : (P \in M_{\alpha}) \lor (\exists P' \supset P : freq(P') \ge \alpha)$$

(With  $F_{\alpha}$  the set of all frequent itemsets)

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An itemset is considered maximal if it is frequent, and none of its proper supersets is frequent

That is:

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(With  $F_{\alpha}$  the set of all frequent itemsets)

### **Definition**

• Every frequent itemset has a maximal superset:

$$\forall \alpha, \forall P \in F_{\alpha} : (\exists P' \in M_{\alpha} : P \subseteq P')$$

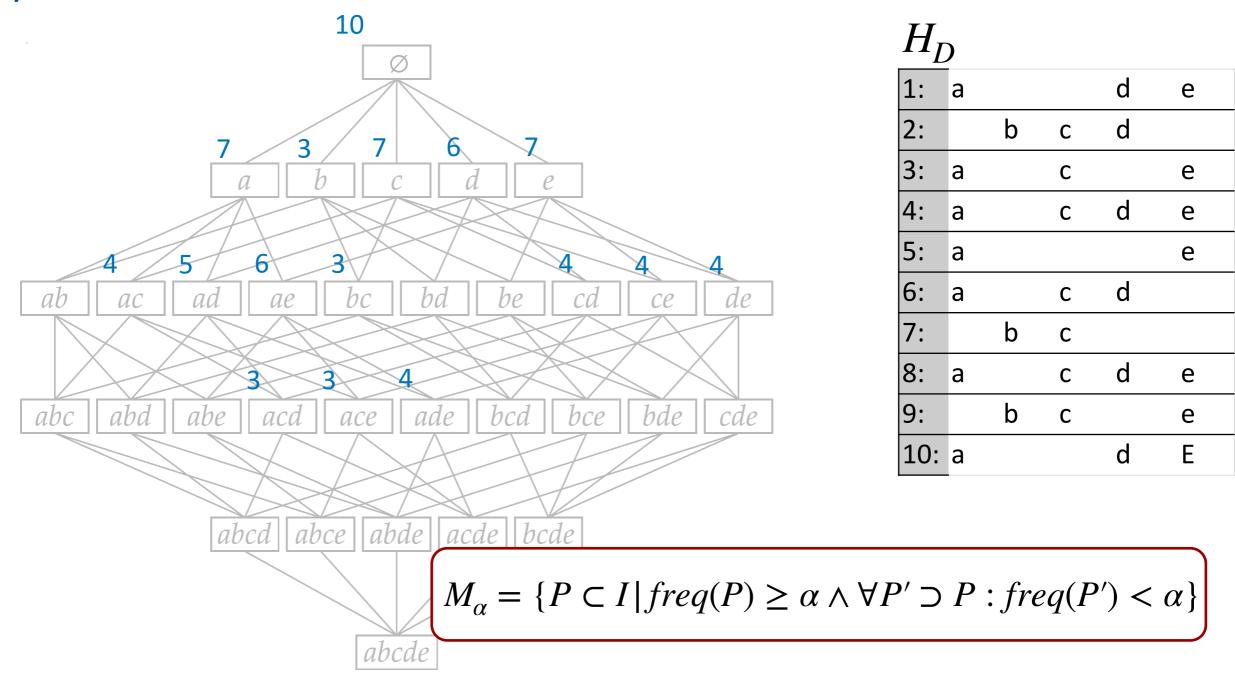
• The maximal itemsets provide a condensed representation of the frequent itemsets, where:

$$\forall \alpha : F_{\alpha} = \bigcup_{P \in M_{\alpha}} 2^{P}$$

# **Maximal Itemset Mining**

### Example(5)

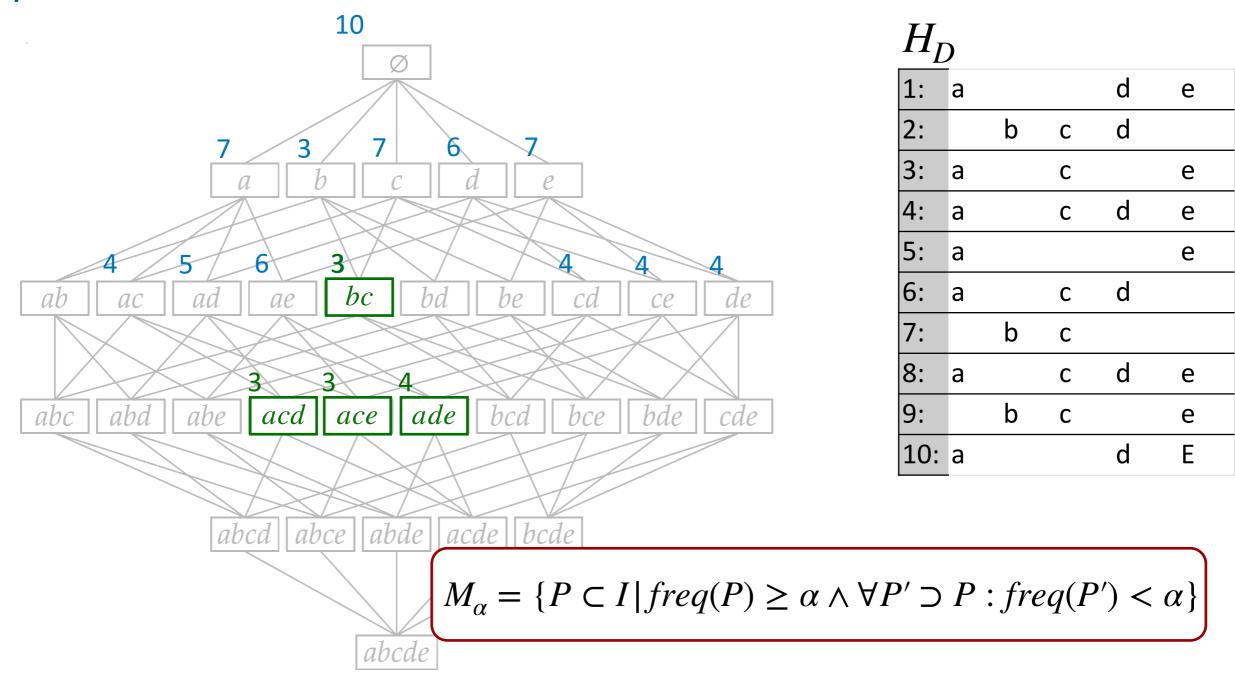
Query-3: Maximal itemsets with  $\alpha = 3$ ?



# **Maximal Itemset Mining**

### Example(5)

Query-3: Maximal itemsets with  $\alpha = 3$ ?



### Limitation

 The set of maximal itemsets captures all frequent itemsets, but it does not retain the support values for all of them

 The Need: Is it possible to have a condensed representation of the frequent itemsets that still preserves the support values for all?"

### Limitation

• The set of Closed (frequent) Itemsets:

$$C_{\alpha} = \{P \subset I | freq(P) \geq \alpha \land \forall P' \supset P : freq(P') < freq(P)\}$$

An itemset is closed if it is frequent, but none of its proper supersets has the same support value

That is:

$$\forall \alpha, \forall P \in F_{\alpha} : (P \in C_{\alpha}) \lor (\exists P' \supset P : freq(P') = freq(P))$$

# **Closed Itemsets**

### **Definition**

• Every frequent itemset has a closed superset:

$$\forall \alpha, \forall P \in F_{\alpha} : (\exists P' \in C_{\alpha} : P \subseteq P')$$

 Closed itemsets provide a condensed representation of frequent itemsets, where:

$$\forall \alpha : F_{\alpha} = \bigcup_{P \in C_{\alpha}} 2^{P}$$

## **Closed Itemsets**

### **Definition**

- Every frequent itemset has a closed superset with the same support
- The set of all closed itemsets preserves the knowledge of support values:

$$\forall \alpha, \forall P \in F_{\alpha} : cover(P) = \max_{P' \in C_{\alpha}, P' \supseteq P} cover(P')$$

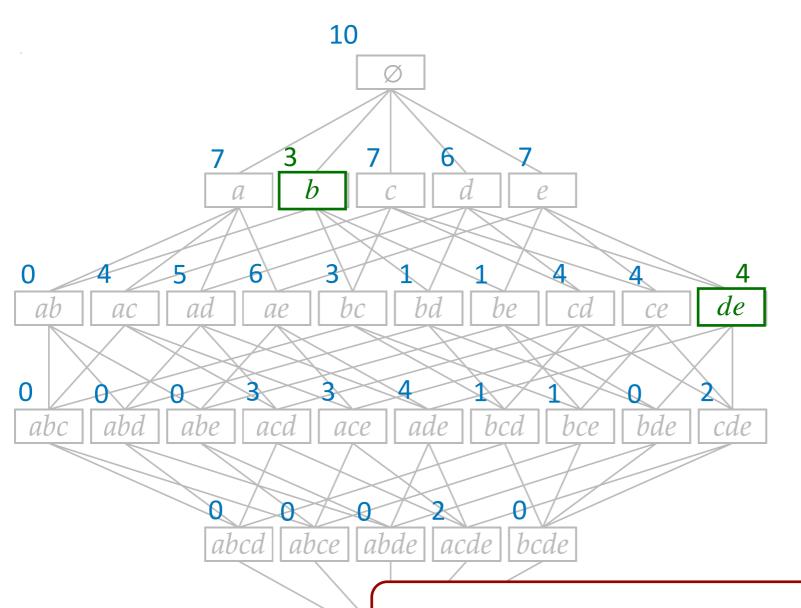
Unlike maximal itemsets:

$$\forall \alpha, \forall P \in F_{\alpha} : cover(P) \supseteq \max_{P' \in M_{\alpha}, P' \supseteq P} cover(P')$$

# **Closed Itemset Mining**

### Example(6)

Q: Are 'b' and 'de' closed itemsets?"



$H_D$								
1:	a			d	е			
1: 2: 3: 4: 5: 6: 7: 8:		b	С	d				
3:	a		С		е			
4:	a		С	d	е			
5:	a				е			
6:	a		С	d				
7:		b	С					
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 $C_{\alpha} = \{P \subset I | freq(P) \geq \alpha \land \forall P' \supset P : freq(P') < freq(P)\}$ 

## **Closed Itemsets**

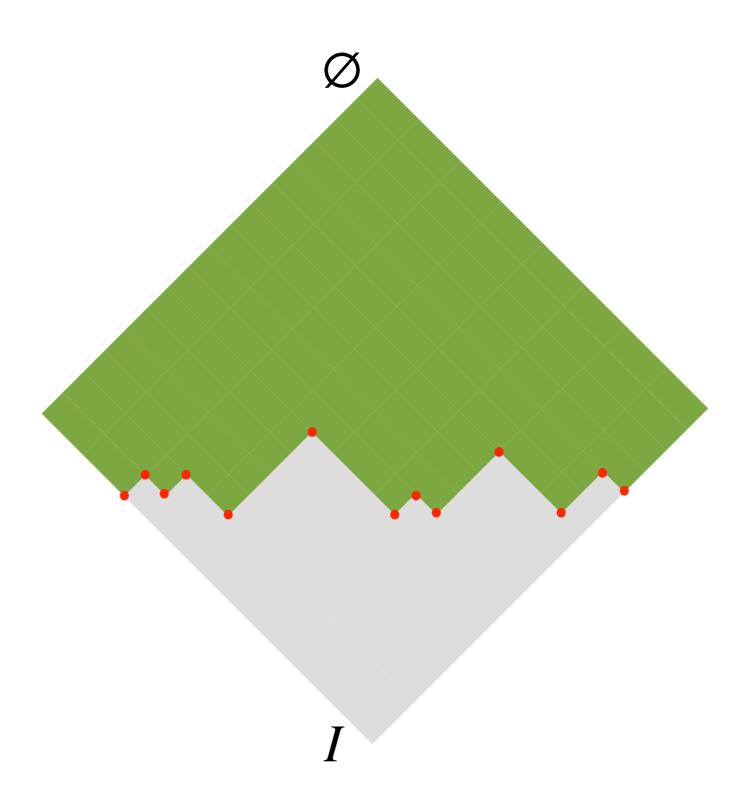
### **Definition**

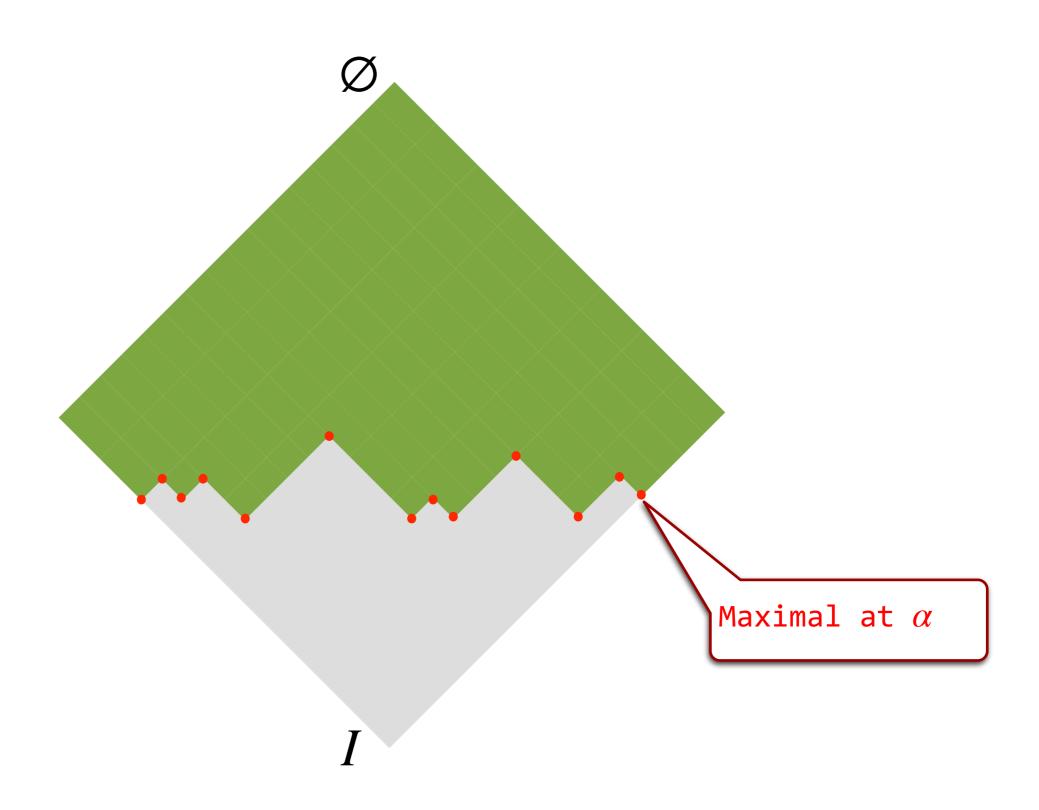
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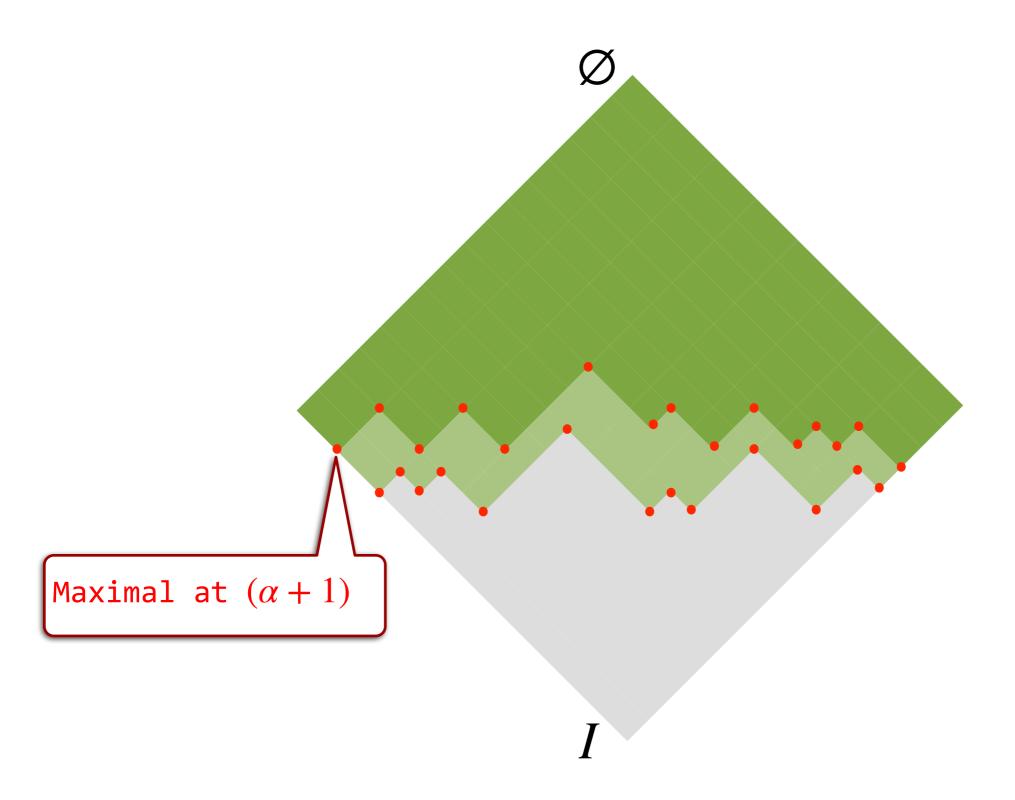
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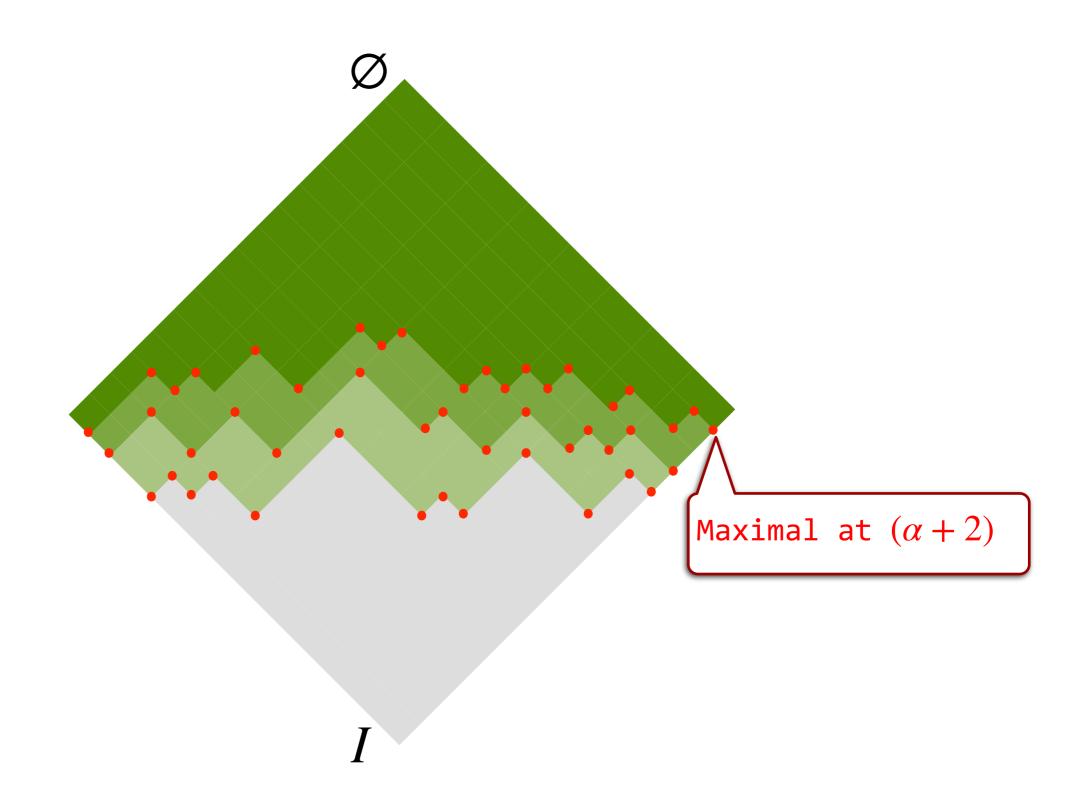
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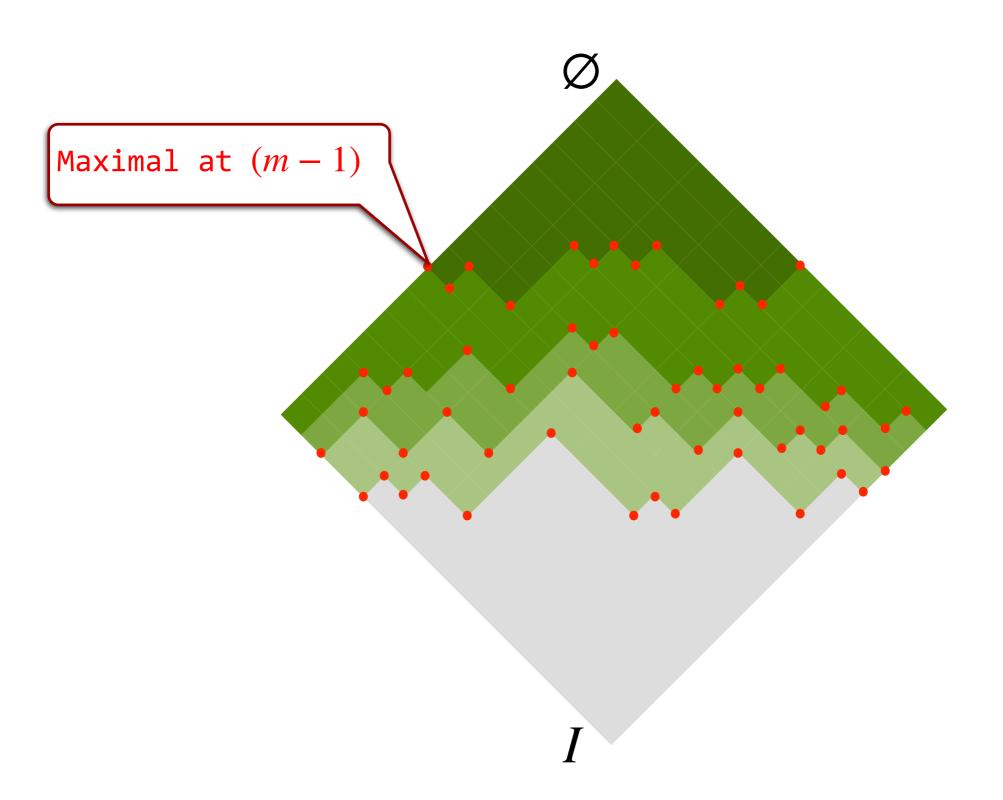
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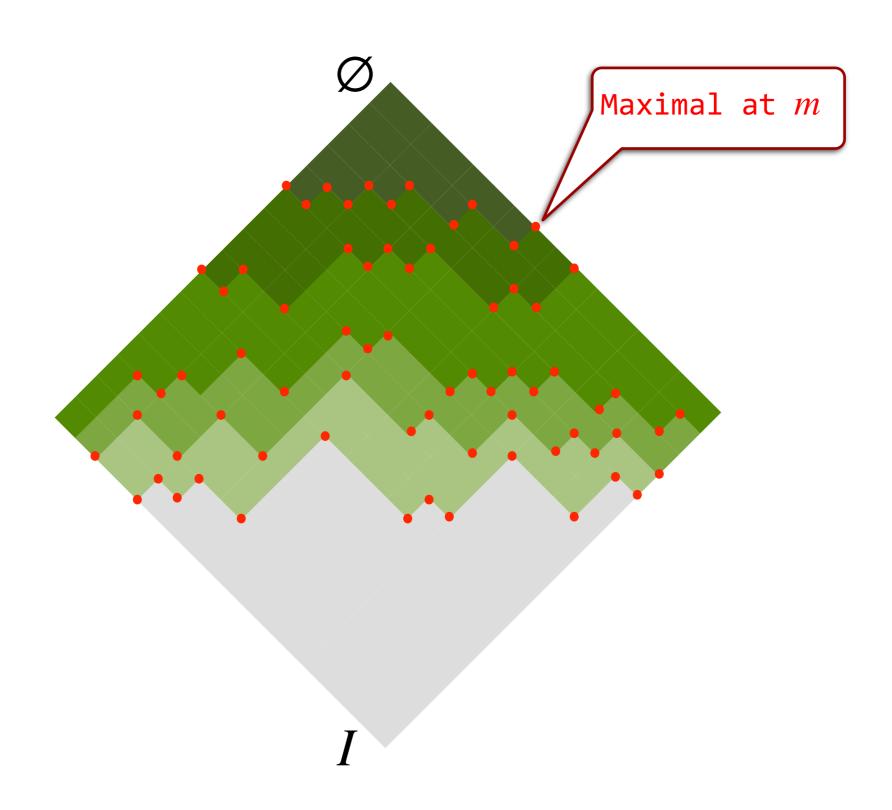


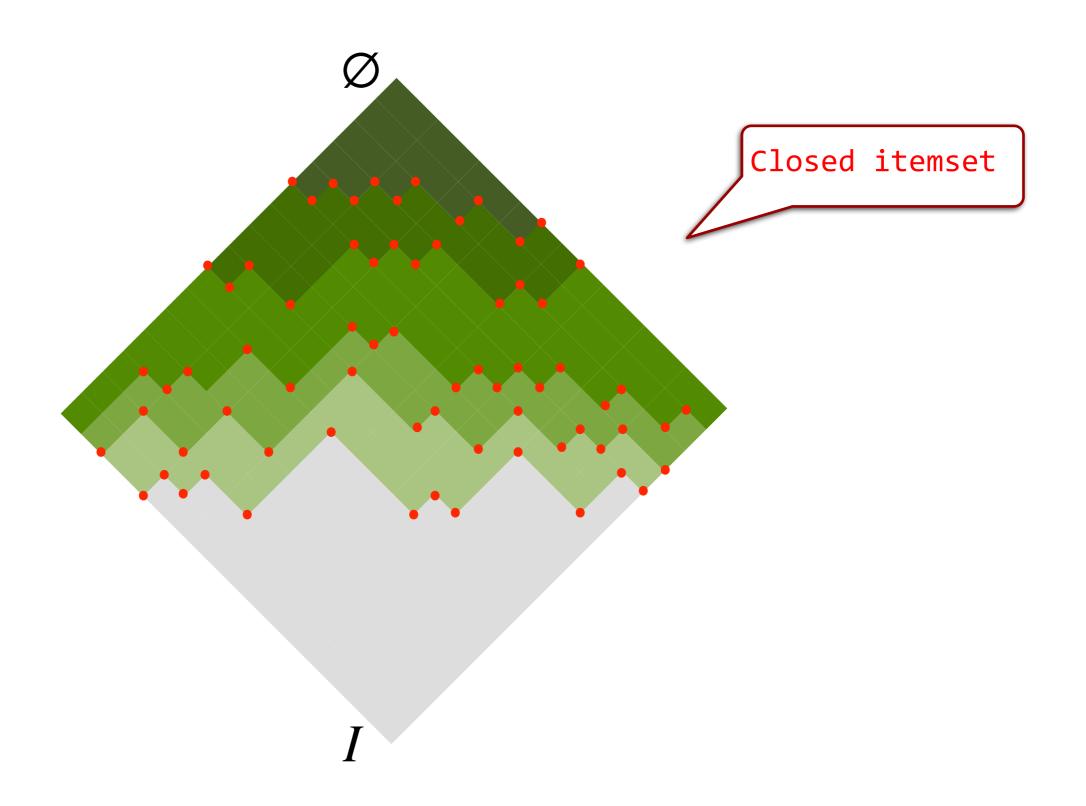


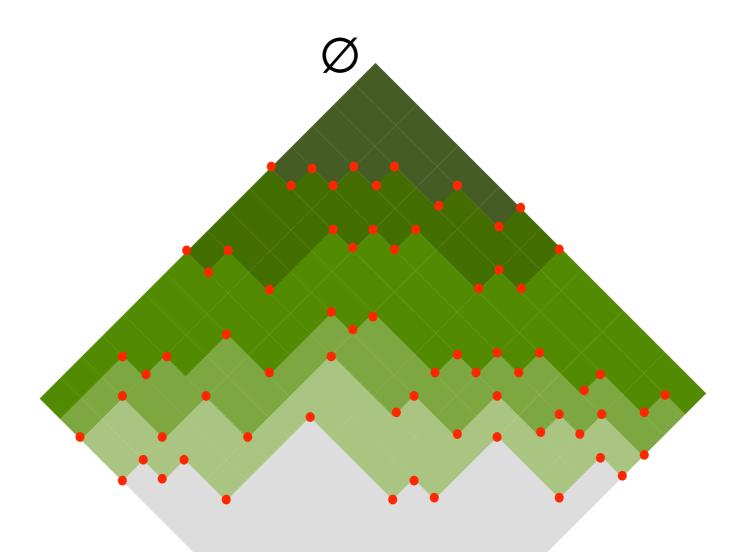












$$C_{\alpha} = \bigcup_{i \in \{\alpha, \dots, m\}} M_i$$

# Frequent/Closed/Maximal

### **Itemsets**

Dataset	#Frequent	#Closed	#Maximal
Zoo-1	151 807	3 292	230
Mushroom	155 734	3 287	453
Lymph	9 967 402	46 802	5 191
Hepatitis	27 . 107+	1 827 264	189 205

http://fimi.ua.ac.be/data/

# LCM Algorithm Linear Closed Item Set Miner

[Uno et al., 03] (version 1) [Uno et al., 04, 05] (versions 2 & 3)

### **Basic Ideas**

- Itemset candidates are evaluated in a given order (depth-first traversal of the prefix tree—lex order for instance—)
- Items are eliminated step-by-step from the transaction database through recursive processing of conditional transaction databases
- Maintains both horizontal and vertical representations of the transaction database in parallel
- Vertical representation: Used to filter transactions based on the chosen split item
- Horizontal representation: Used to populate the vertical representation for the next recursion step (avoiding intersection operations as in the Eclat algorithm).

### **Notations**

• Tail: The tail of an itemset P is the last item in P:

$$P = \{a_1, ... a_k\} : tail(P) = a_k$$

• Sub-itemsets: The sub-itemset of P up to the  $j^{th}$  item is defined as:

$$P = \{a_1, ..., a_k\} : P(a_i) = \{a_1, ..., a_i\}$$

• **Prefix:** The prefix of *P* in the prefix tree is defined as:

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### Closure

A set S has a closure under an operation f if:

$$\forall x \in S, f(x) \in S$$

- In this case, we say that the set S is closed under f
- A closure operation is:
  - Increasing (or extensive): The closure of an object contains the object itself, i.e.,  $S \subseteq \text{closure}(S)$
  - **Idempotent:** The closure of the closure of an object is the closure itself, i.e., closure(closure(S)) = closure(S)
  - Monotone: If  $X \subseteq Y$ , then  $closure(X) \subseteq closure(Y)$

### **Itemset Closure**

$$closure(P) = \bigcap_{t \in cover(P)} t$$

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$$closure(P) = \bigcap_{t \in cover(P)} t$$

t1:	В	C		Е	F	G	Н
t2: A			D			G	
t3: A		C	D				Н
<b>t4:</b> A				Е	F		
t5:	В			E	F		
t6:	В			Е	F	G	

### **Itemset Closure**

$$closure(P) = \bigcap_{t \in cover(P)} t$$

t1:	В	C		E	F	G	Н
t2: A			D			G	
t3: A		C	D				Н
t4: A				Е	F		
t5:	В			E	F		
t6:	В			Е	F	G	

closure(EG) = ?

### **Itemset Closure**

$$closure(P) = \bigcap_{t \in cover(P)} t$$

t1:	В	C	Е	F	G	Н
t5:	В		E	F		
t6:	В		 Е	F	G	

closure(EG) = ?

### **Itemset Closure**

$$closure(P) = \bigcap_{t \in cover(P)}$$

t1:	В	C	Е	F	G	Н
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closure(EG) = ?

closure(EG) = BEFG

### **Itemset Closure**

$$closure(P) = \bigcap_{t \in cover(P)}$$

t1:		В	C		E	F	G	Н
t2:	A			D			G	
t3:	A		С	D				Н
t4:	A				Е	F		
t5:		В			Е	F		
t6:		В			Е	F	G	

$$closure(EG) = ?$$

$$closure(EG) = BEFG$$

$$closure(EF) = ?$$

### **Itemset Closure**

$$closure(P) = \bigcap_{t \in cover(P)} t$$

t1:	В	C		Е	F	G	Н
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<b>t4:</b> A				Е	F		
t5:	В			Е	F		
t6:	В			Е	F	G	

closure(EG) = ?

closure(EG) = BEFG

closure(EF) = ?

- Closure Extension: A rule for constructing a closed itemset from an existing one:
  - Simply add an item and compute its closure!

### Closure Extension [Pasquier et al., 99]

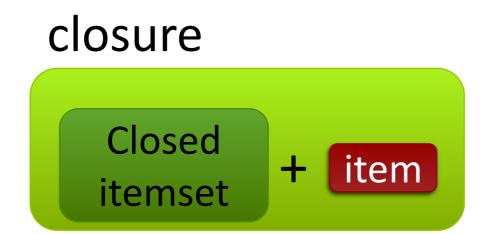
- Closure Extension: A rule for constructing a closed itemset from an existing one:
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Closed itemset

- Closure Extension: A rule for constructing a closed itemset from an existing one:
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### **Algorithm**

```
Algorithm 1: LCM

InOut: P: Closed Frequent Itemset;
In: \alpha: minsup

print(P)
foreach p_i > \text{tail}(P) do

if freq(P \cup \{p_i\}) \ge \alpha then \text{LCM}(\text{closure}(P \cup \{p_i\}, \alpha))
```

[Uno et al., 03]



# universite Paris-saclay



# **Constraint & Data Mining**

Cours2

### Master 2 - DS

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