

1. De las siguientes afirmaciones, indicar cuáles son ciertas y cuáles no:

- a) $n^2 \in O(n^3)$ g) $\log n \in O(n^{1/2})$
 b) $n^3 \in O(n^2)$ h) $n^{1/2} \in O(\log n)$
 c) $2^{(n+1)} \in O(2^n)$ i) $n^2 \in \Omega(n^3)$
 d) $(n+1)! \in O(n!)$ j) $n^3 \in \Omega(n^2)$
 e) $f(n) \in O(n) \Rightarrow 2^{f(n)} \in O(2^n)$ k) $2^{(n+1)} \in \Omega(2^n)$
 f) $3^n \in O(2^n)$

e) $f(n) \in O(n) \Rightarrow 2^{f(n)} \in O(2^n) \rightarrow V$

f) $3^n \in O(2^n) \rightarrow F$

g) $\log_2 n \in O(\sqrt{n}) \rightarrow V$

$\log_{20} n \in O(\sqrt{n}) \rightarrow V$

h) $n^{1/2} \in O(\log n) \rightarrow F$

i) $n^2 \in \Omega(n^3) \rightarrow F$

j) $n^3 \in \Omega(2^n) \rightarrow F$

k) $2^{(n+1)} \in \Omega(2^n) \rightarrow V$

l) $n^{1/2} \in O(\log n) \rightarrow \sqrt{n} \in O(\log n) \rightarrow F$

m) $n^2 \in \Omega(n^3) \rightarrow F$

n) $n^3 \in \Omega(n^2) \rightarrow V$

o) $2^{n+1} \in \Omega(2^n) \rightarrow V$

2. En los siguientes segmentos de código, determinar su tiempo de ejecución en función de n:

a)

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c=1; while (c <= n) { algo.de.O(1); c = 2*c; }
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 $\rightarrow T(n) = \left(\sum_{i=1}^k 4\right) + 1 + 1 = 4k + 2 = 4 \log_2 n + 2$

$2^k \geq n ; \log_2 2^k \geq \log_2 n ; k \geq \log_2 n$

b)

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c=n; while (c > 1) { algo.de.O(1); c = c/2; }
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 $\rightarrow T(n) = \left(\sum_{i=1}^k 4\right) + 1 + 1 = 4k + 2 = 4 \log_2 n + 2$

$n/2^k \leq 1 ; n \leq 2^k ; \log_2 n \leq \log_2 2^k ; \log_2 n \leq k$

c)

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for (int i=0; i <= n; i++) { c=n; while (c > 1) { algo.de.O(1); c = c/2; } }
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 $\rightarrow T(n) = 2 + \sum_{i=0}^{n-1} (3 + 1 + T_b(n)) = 2 + n(4 + T_b(n)) = 2 + n(4 + (4 \log_2 n + 2)) = 2 + 6n + 4n \log_2 n$

$n/2^k \leq 1 \rightarrow \log_2 n \leq k$

Pregunta por

d)

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for (int i=1; i <= n; i++) { c=i; while (c > 1) { algo.de.O(1); c = c/2; } }
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 Pista: para n suficientemente grande, $\log(n) \approx \log n - n$

3

$T(n) = 2T(n-1) + 3^n(n+5)$ $n > 0$ $\left\{ \begin{array}{l} T(n) - 2T(n-1) = 3^n(n+5) \\ a_2 = -2 ; k = 1 ; h(n) = 3^n(n+5) \end{array} \right.$

$T(0) = 0$ Caso base

$3^n(n+5) = \sum_{i=1}^n S_i \rightarrow \frac{1}{2} (n+5) \cdot 3^n \rightarrow m_1 = 1 ; S_1 = 3$

$c_1(x) = x^k + a_{k-1}x^{k-1} + \dots + a_0$
 $c_2(x) = (x - S_1)^{m_1+1} + \dots + (x - S_k)^{m_k+1}$

$c_1(x) = x - 2 \cdot x^{1-1} = x - 2$ $\rightarrow c_1 \cdot c_2 = (x-2) \cdot (x-3)^2 \rightarrow R = \{2, 3, 3\}$
 $c_2(x) = (x-3)^2$ $m_1 = 1 ; m_2 = 2$

4. Raíces múltiples $\rightarrow B = \{2^n, 3^n, n \cdot 3^n\} \rightarrow T(n) = \lambda_1 2^n + \lambda_2 3^n + \lambda_3 n \cdot 3^n \Rightarrow T(n) = -9 \cdot 2^n + 9 \cdot 3^n + 3n \cdot 3^n = -9 \cdot 2^n + 3^{n+2} + n \cdot 3^{n+1}$

Sea $T(0) = 0 \rightarrow T(0) = \lambda_1 + \lambda_2 = 0 \Rightarrow \lambda_1 = -\lambda_2$

Sea $T(1) = 2T(0) + 3^1(1+5) = 18$

Sea $T(2) = 2T(1) + 3^2(2+5) = 2 \cdot 18 + 9 \cdot 7 = 36 + 63 = 99$

Sea $T(3) = 2T(2) + 3^3(3+5) = 2 \cdot 99 + 27 \cdot 8 = 198 + 216 = 414$

Sistema: $\left\{ \begin{array}{l} \lambda_1 + \lambda_2 = 0 \\ \lambda_1 \cdot 2 + \lambda_2 \cdot 3 + \lambda_3 \cdot 1 \cdot 3 = 18 \\ \lambda_1 \cdot 4 + \lambda_2 \cdot 9 + \lambda_3 \cdot 2 \cdot 9 = 99 \end{array} \right.$

$\left(\begin{array}{ccc|c} 1 & 1 & 0 & 0 \\ 2 & 3 & 3 & 18 \\ 4 & 9 & 18 & 99 \end{array} \right) \xrightarrow{F_3 \rightarrow F_3 - 2F_2} \left(\begin{array}{ccc|c} 1 & 1 & 0 & 0 \\ 2 & 3 & 3 & 18 \\ 0 & 3 & 12 & 63 \end{array} \right) \xrightarrow{F_2 \rightarrow F_2 - 2F_3} \left(\begin{array}{ccc|c} 1 & 1 & 0 & 0 \\ 0 & 1 & 3 & 18 \\ 0 & 3 & 12 & 63 \end{array} \right) \rightarrow$

$$\vec{r}_3 \rightarrow F_3 - 3F_2 \quad \left(\begin{array}{ccc|c} 1 & 1 & 0 & 0 \\ 0 & 1 & 3 & 48 \\ 0 & 0 & 3 & 9 \end{array} \right) \rightarrow \begin{cases} \lambda_1 + \lambda_2 = 0 \rightarrow \lambda_1 = -9 \\ \lambda_2 + 3\lambda_3 = 48 \rightarrow \lambda_2 + 3 \cdot 3 = 48; \lambda_2 = 9 \\ 3\lambda_3 = 9; \lambda_3 = 3 \end{cases}$$

4)

a) $T(n) = 3T(n-1) + 4T(n-2)$ si $n > 1$; $T(0) = 0$; $T(1) = 1$

$$T(n) - 3T(n-1) - 4T(n-2) = 0 \rightarrow a_1 = -3; a_2 = 4; K = 2$$

$$x^2 - 3x - 4 = 0; \quad x = \frac{3 \pm \sqrt{(-3)^2 - 4 \cdot 1 \cdot (-4)}}{2 \cdot 1} = \frac{3 \pm \sqrt{9+16}}{2} = \frac{3 \pm 5}{2} = \begin{cases} x_1 = 4 \\ x_2 = -1 \end{cases}$$

$$R = \{4, -1\} \rightarrow c(x) = (x-4) \cdot (x+1) \rightarrow m_1 = 1; m_2 = 1 \text{ (Raíces simples)}$$

$$B = \{4^n, (-1)^n\} \rightarrow T(n) = \lambda_1 \cdot 4^n + \lambda_2 \cdot (-1)^n \rightarrow T(n) = \frac{4^n}{5} - \frac{(-1)^n}{5} = \frac{1}{5} \cdot (4^n - (-1)^n) \in \Theta(4^n)$$

$$\begin{cases} T(0) = \lambda_1 \cdot 4^0 + \lambda_2 \cdot (-1)^0 \rightarrow \lambda_1 + \lambda_2 = 0 \rightarrow \lambda_1 = -\lambda_2; -4\lambda_2 - \lambda_2 = 1; -5\lambda_2 = 1; \lambda_2 = -\frac{1}{5} \Rightarrow \lambda_1 = \frac{1}{5} \\ T(1) = \lambda_1 \cdot 4^1 + \lambda_2 \cdot (-1)^1 \rightarrow 4\lambda_1 - \lambda_2 = 1 \end{cases}$$

b) $T(n) = 4T(n/2) + n^2$; $n > 1$; $T(1) = 1$; $T(2) = 8$

$$T(n) - 4T(n/2) = n^2; \quad T(2^m) - 4T(2^{m-1}) = 2^{2m}; \quad T(2^m) - 4T(2^{m-1}) = 2^{2m} \rightarrow$$

cambio \downarrow Si $T(2^m) = t(m) \rightarrow t(m) - 4t(m-1) = 4^m \rightarrow a_1 = -4; k = 1; h(m) = 4^m = \sum_{i=1}^m n^0 \cdot 4^m; m_1 = 0; S_1 = 4$

$$c_1(x) = x - 4 \quad c_2(x) = (x-4)^{0+1} = x-4 \rightarrow c_1 \cdot c_2 = (x-4)^2 = c(x) \quad \left\{ \begin{array}{l} \text{multiplicidad} = 2 \\ \text{Raíces múltiples o simples} \end{array} \right.$$

Raíces múltiples \rightarrow Sacar base $B = \{4^m, m \cdot 4^m\} \rightarrow$ Ec. general $T(n) = \lambda_1 \cdot 4^m + \lambda_2 \cdot m \cdot 4^m$

calculo λ_2 $t(m) = \lambda_1 \cdot 4^m + \lambda_2 \cdot m \cdot 4^m \rightarrow \lambda_2 \cdot m \cdot 4^m - 4 \cdot \lambda_2 \cdot (m-1) \cdot 4^{m-1} = 4^m; \quad \lambda_2 \cdot m \cdot 4^m - \lambda_2 \cdot (m-1) \cdot 4^m = 4^m; \quad \lambda_2 \cdot m - \lambda_2 \cdot (m-1) = 1; \quad \lambda_2 = 1$

$\frac{T(n)}{\text{que debería}}$ $t(m) = \lambda_1 \cdot 4^m + m \cdot 4^m$ deja el cambio $t(m) = T(2^m); n = 2^m; \log_2 n = m \rightarrow T(n) = \lambda_1 \cdot 4^{\log_2 n} + \log_2 n \cdot 4^{\log_2 n}$

$$T(n) = \lambda_1 \cdot n^2 + \log_2 n \cdot n^2$$

$$T(1) = 1 = \lambda_1 \cdot 1^2 + \log_2 1 \cdot 1^2; \quad \lambda_1 = 1 \quad \text{Finalmente queda así} \quad T(n) = n^2 + n^2 \log_2 n \in \Theta(n^2 \log n)$$

c) $T(n) = 2T(n/2) + n \cdot \log n$; $n > 1, n = 2^m$

cambio \downarrow $T(2^m) = 2T(2^{m-1}) + 2^m \cdot \log 2^m; \quad T(2^m) = 2T(2^{m-1}) + 2^m \cdot m$

$$\rightarrow T(2^m) = t(m) \rightarrow t(m) = 2t(m-1) + 2^m \cdot m; \quad t(m) - 2t(m-1) = 2^m \cdot m$$

$$a_1 = -2; k = 1; h(m) = 2^m \cdot m \quad \text{como } 2^m \cdot m = \sum_{i=1}^m m \cdot 2^m$$

$$\Rightarrow m_1 = 1, S_1 = 2$$

saco c_1, c_2 $c_1 = x-2 \quad c_2 = (x-2)^2 \quad \left\{ \begin{array}{l} c_1 \cdot c_2 = (x-2) \cdot (x-2)^2 = (x-2)^3 \rightarrow R = \{2, 2, 2\} \end{array} \right. \quad m_2 = 3$

raíces múltiples $B = \{2^m, m \cdot 2^m, m^2 \cdot 2^m\} \rightarrow t(m) = \lambda_1 \cdot 2^m + \lambda_2 \cdot m \cdot 2^m + \lambda_3 \cdot m^2 \cdot 2^m$

raíces múltiples $B = \{z^m, m \cdot z^m, m^2 \cdot z^m\} \rightarrow t(m) = \lambda_1 \cdot z^m + \lambda_2 \cdot m \cdot z^m + \lambda_3 \cdot m^2 \cdot z^m$

Desarrollo casillas $t(m) = T(2^m)$; $z^m = n$; $\log z^m = \log n$; $m = \log n \Rightarrow T(n) = \lambda_1 n + \lambda_2 \cdot n \cdot \log n + \lambda_3 n \log^2 n \in \Theta(n \log^2 n)$

d) $T(n) = 3T(n/2) + 5n + 3$ si $n > 1$, $n = 2^m$

$\rightarrow n = 2^m \rightarrow T(2^m) = 3T(2^{m-1}) + 5 \cdot 2^m + 3$; $T(2^m) - 3T(2^{m-1}) = 5 \cdot 2^m + 3$;

$\rightarrow T(2^m) = t(m) \rightarrow t(m) - 3t(m-1) = 5 \cdot 2^m + 3$; $a_1 = -3$; $k = 1$; $h(m) = 5P$

$\rightarrow \sum_{i=0}^1 p(m) S_i^m \rightarrow$

e) $T(n) = 2T(n/2) + \log n$ si $n > 1$, $n = 2^m \rightarrow T(2^m) = 2T(2^{m-1}) + \log 2^m$; $T(2^m) - 2T(2^{m-1}) = m$

$\rightarrow T(2^m) = t(m) \rightarrow t(m) - 2t(m-1) = m$; $a_1 = -2$; $k = 1$; $h(m) = m$

$\rightarrow \sum p(m) \cdot S^m \rightarrow \sum_{i=0}^1 m^i \cdot 1^m \rightarrow m_1 = 1$; $S_1 = 1$; $\rightarrow \begin{cases} c_1(x) = x - 2 \\ c_2(x) = (x - 1)^2 \end{cases} \mid c_1 \cdot c_2 = (x - 2)(x - 1)^2$

$R = \{2, 1, 1\} \rightarrow m_1 = 1$; $m_2 = 2$; $B = \{2^m, 1, m\}$

$t(m) = \lambda_1 \cdot 2^m + \lambda_2 + \lambda_3 \cdot m$ Desarrollo casillas $T(2^m) = \lambda_1 \cdot 2^m + \lambda_2 + \lambda_3 \cdot m$ $\xrightarrow{z^m = n, m = \log n} T(n) = \lambda_1 \cdot n + \lambda_2 + \lambda_3 \cdot \log(n) \in \Theta(\log(n))$ [si $\lambda_1 = 0$]
 $\in \Theta(n)$ [si $\lambda_1 \neq 0$]

f) $T(n) = T(n-1) + 2T(n-2) - 2T(n-3)$; $n > 2$; $T(n) = 9n^2 - 15n + 106$ si $n = 0, 1, 2$

$T(n) - T(n-1) - 2T(n-2) + 2T(n-3) = 0$; $a_1 = -1$, $a_2 = -2$, $a_3 = 2$, $k = 3$, $h(n) = 0$

$x^3 - x^2 - 2x + 2 = 0$;

1	-1	-2	2
1	0	-2	0

$\rightarrow x^2 - 2 = 0$; $x^2 = 2$; $x = \pm \sqrt{2}$

$R = \{1, \sqrt{2}, -\sqrt{2}\} \rightarrow B = \{1, \sqrt{2}^n, (-\sqrt{2})^n\} \rightarrow T(n) = \lambda_1 + \lambda_2 \cdot (\sqrt{2})^n + \lambda_3 \cdot (-\sqrt{2})^n$

$T(0) = 9 \cdot 0^2 - 15 \cdot 0 + 106 = 106$

$T(1) = 9 \cdot 1^2 - 15 \cdot 1 + 106 = 9 - 15 + 106 = 100$

$T(2) = 9 \cdot 2^2 - 15 \cdot 2 + 106 = 36 - 30 + 106 = 112$

$\left\{ \begin{array}{l} 106 = \lambda_1 + \lambda_2 \cdot (\sqrt{2})^0 + \lambda_3 \cdot (-\sqrt{2})^0; \quad 106 = \lambda_1 + \lambda_2 + \lambda_3 \\ 100 = \lambda_1 + \lambda_2 \cdot (\sqrt{2})^1 + \lambda_3 \cdot (-\sqrt{2})^1; \quad 100 = \lambda_1 + \sqrt{2} \cdot \lambda_2 - \sqrt{2} \cdot \lambda_3 \\ 112 = \lambda_1 + \lambda_2 \cdot (\sqrt{2})^2 + \lambda_3 \cdot (-\sqrt{2})^2; \quad 112 = \lambda_1 + 2\lambda_2 + 2\lambda_3 \end{array} \right\} \rightarrow$